


# **Write-up\_picoCTF\_2021**

# Chall: Rolling My Own

You can find out the challenge's source and solving script here( [source](#), [script](#))

 Firstly, I think two hints of this challenge is very important. Firstly, you should deeply the [paper](#) related directly this challenge. Besides, the second hint *"Here's the start of the password: D1v1 "* is also worthy.

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## The main idea of paper

This paper suggested an anti-disassembly technique by using cryptographic hash. Particularly, a key and salt were chosen to generate the hash digest which contains some bytes of machine code( also called "running code"). In program there are a function which pick some bytes at arbitrary position in those hash digests ( these positions will be decided by programmer before) and create running code. The disassembly of IDA or well-known disassembler will failed in disassembling those bytes because the byte arrays which they received are hash digests, no machine code.

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## The key steps to right password

### The idea of challenge

- Based on above paper, author claims the right keys from us to concatenate them with available salts to create plaintexts, get MD5 digests of these plaintexts.
- In each generated md5 digests, there is a **running** code. Merging all running code at specified positions in MD5 digests, I have a function which can call flag-opening function.

### Step-by-step to solve

## Determine the salts and positions of running code in MD5 digest

- i The pseudo-code generated by IDA was edited( function's name, variable's name...) to simplify analysis process.

```
v17 = 'WEjMaLpG';  
v18 = 'kmnnjOVp';  
v19 = '6pdeliGR';  
v20 = 'slxzecvM';  
v21 = 0;  
v11 = 8;  
v12 = 2;  
v13 = 7;  
v14 = 1;  
memset(pwdInput, 0, 0x40uLL);  
memset(&dest, 0, 0x40uLL);  
printf("Password: ", 0LL);  
fgets(pwdInput, 64, stdin);
```

Storing salts and positions of running code on stack

- While debugging and analyzing with arbitrary password, I recognized that v17 - v20 are salts of plaintext( easily concluded there are four 16-byte MD5 digests) and v11 - v14 are starting positions of running code in those digests in corresponding.

## Determine the running code to open flag

```
for ( i = 0; i <= 3; ++i )
{
    strncat(&dest, &passwdInput[4 * i], 4uLL);
    strncat(&dest, (const char *)&v17 + 8 * i, 8uLL);
}
```

- Above snipped-code will concatenate the part of password( each 4 bytes at one time) with its salt( 8 bytes). Finally, we will have 4 12-byte plaintexts to do cryptographic hash

```
someCryptoHashStuff((__int64)msg_Digest, (__int64)&dest, v3);
for ( ia = 0; ia <= 3; ++ia )
{
    for ( j = 0; j <= 3; ++j )
        *((_BYTE *)&runArray + 4 * j + ia) = msg_Digest[16 * j + ia + *(&v11 + j)];
}
v4 = (void (__fastcall *))(unsigned __int64 (__fastcall *)(__int64), signed __int64))mmap(0LL, 0x10uLL, 7, 34, -1, 0LL);
v5 = v16;
*(_QWORD *)v4 = runArray;
*(_QWORD *)v4 + 1 = v5;
v4(getFlag, 16LL);
```

- From the loops above, I easily discovered that there are 16 bytes running code( get from 4 bytes of each MD5 digests).

Conclude about running code: ***its length is 16 bytes and it can call the flag-opening function. (\*)***

There are two noticeable points which I haven't yet mention above. They are: challenge's second hint and the flag-opening function.

## 1 - Hint:

 ***Here's the start of the password: D1v1***

- I will calculate the MD5 digest of the first pass of password concatenating its corresponding salt "GpLaMjEW".  
MD5 digest in hexa: **23 f1 44 e0 8b 60 3e 72 48 89 fe 48 9f 78 fa 53**
- If I get 4 adjacent bytes starting digest[8]( because the specified of running code in first MD5 digest is 8), I will have **48 89 fe 48**

, try disassembling it at [here](#):

0: 48 89 fe            mov rsi,rdi

3: 48            rex.W( maybe miss some bytes to create a completed instruction)

## 2 - The flag- opening function

```
getFlag proc near
```

```
var_A8= qword ptr -0A8h
```

```
stream= qword ptr -98h
```

```
s= byte ptr -90h
```

```
var_8= qword ptr -8
```

```
; __unwind {
```

```
push     rbp
```

```
mov      rbp, rsp
```

```
sub      rsp, 0B0h
```

```
mov      [rbp+var_A8], rdi
```

```
mov      rax, fs:28h
```

```
mov      [rbp+var_8], rax
```

```
xor      eax, eax
```

```
mov      rax, 7B3DC26F1h
```

```
cmp      [rbp+var_A8], rax
```

```
jz       short loc_5631E3D4106D
```

```
loc_5631E3D4106D:
```

```
lea      rsi, modes            ; "r"
```

```
lea      rdi, filename        ; "flag"
```

```
call     _fopen
```

```
mov      [rbp+stream], rax
```

```
cmp      [rbp+stream], 0
```

```
jnz      short loc_5631E3D410A7
```

- To open **flag.txt** file, %rdi register must be assigned equally 0x7b3dc26f1. This is a task of running code. (\*\*)

From (\*) and (\*\*), I will create running code appropriately to satisfy all above conditions:

```
1 #Before, %rdi stored the address of flag-opening function
2 mov     rsi,rdi            #%rsi = &flag-opening func
```

```

3  movabs rdi,0x7b3dc26f1    #%rdi = 0x7b3dc26f1
4  jmp     rsi               #jump to flag-opening func
5  ret
6
7  #and this is running code:
8  0:  48 89 fe              mov     rsi,rdi
9  3:  48 bf f1 26 dc b3 07    movabs rdi,0x7b3dc26f1
10 a:  00 00 00
11 d:  ff d6                  call    rsi
12 f:  c3                     ret

```

## Brute-force three remaining keys

- Brute-force 3 remaining keys to create MD5 digest containing running code at specified positions.