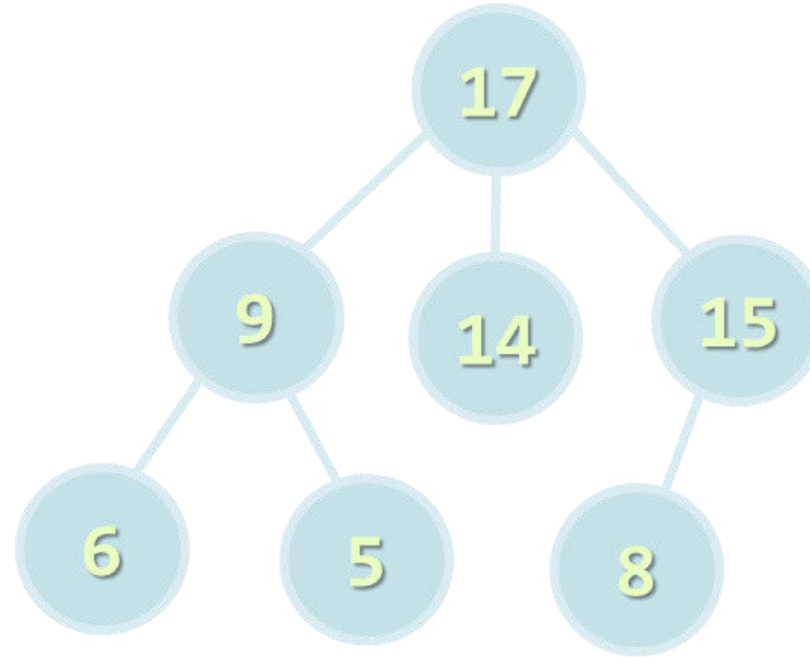


Data Structures and Algorithms

LECTURE 09: BINARY TREES, HEAPS, BINARY SEARCH TREES

- Binary Trees
 - Traversal algorithms
- Heaps
 - Binary heap, Min/Max heaps
- Binary Search Trees

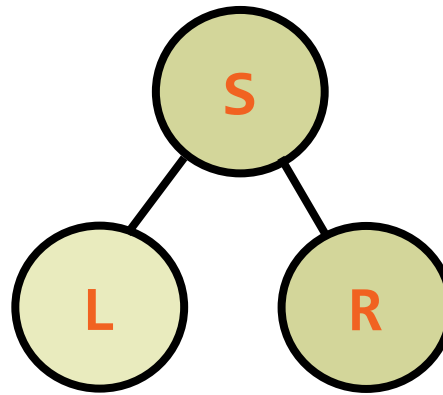
Binary Trees and BT Traversal



Preorder, In-Order, Post-Order

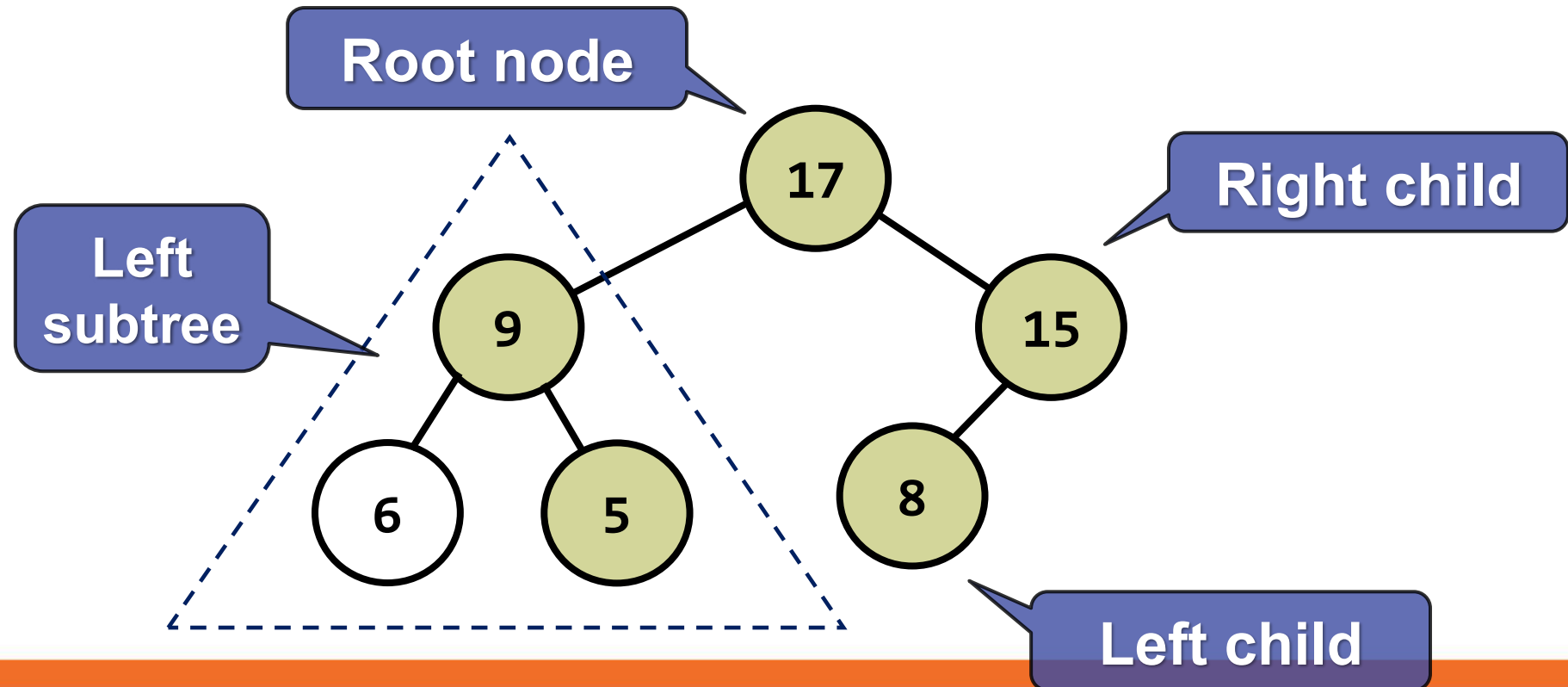
Binary Tree

- ADS representing tree like hierarchy
- Each node has **at most two** children
 - Children are called **left** and **right**
 - The **parent** is also called **source**



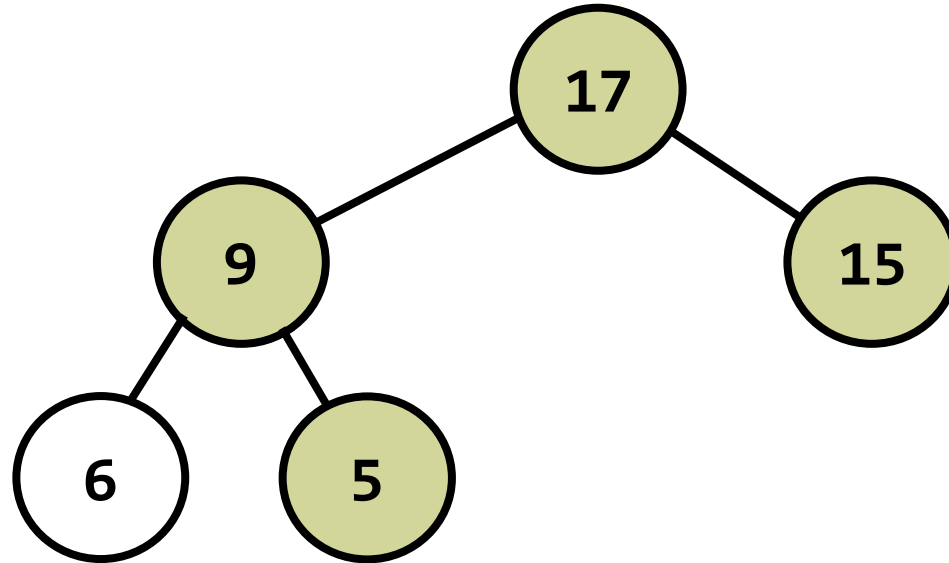
Binary Trees

- **Binary trees:** the most widespread form
 - Each node has at most 2 children (left and right)



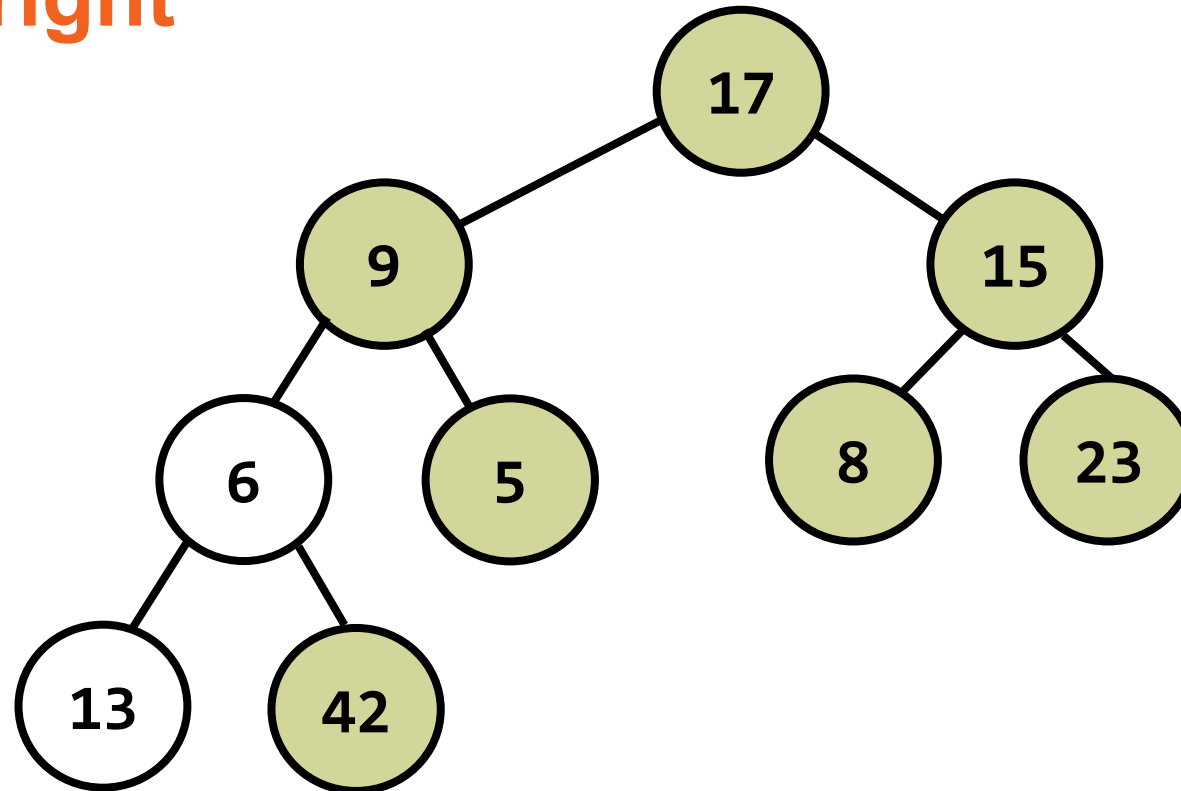
Types of Binary Trees

- **Full** – each node has 0 or 2 children



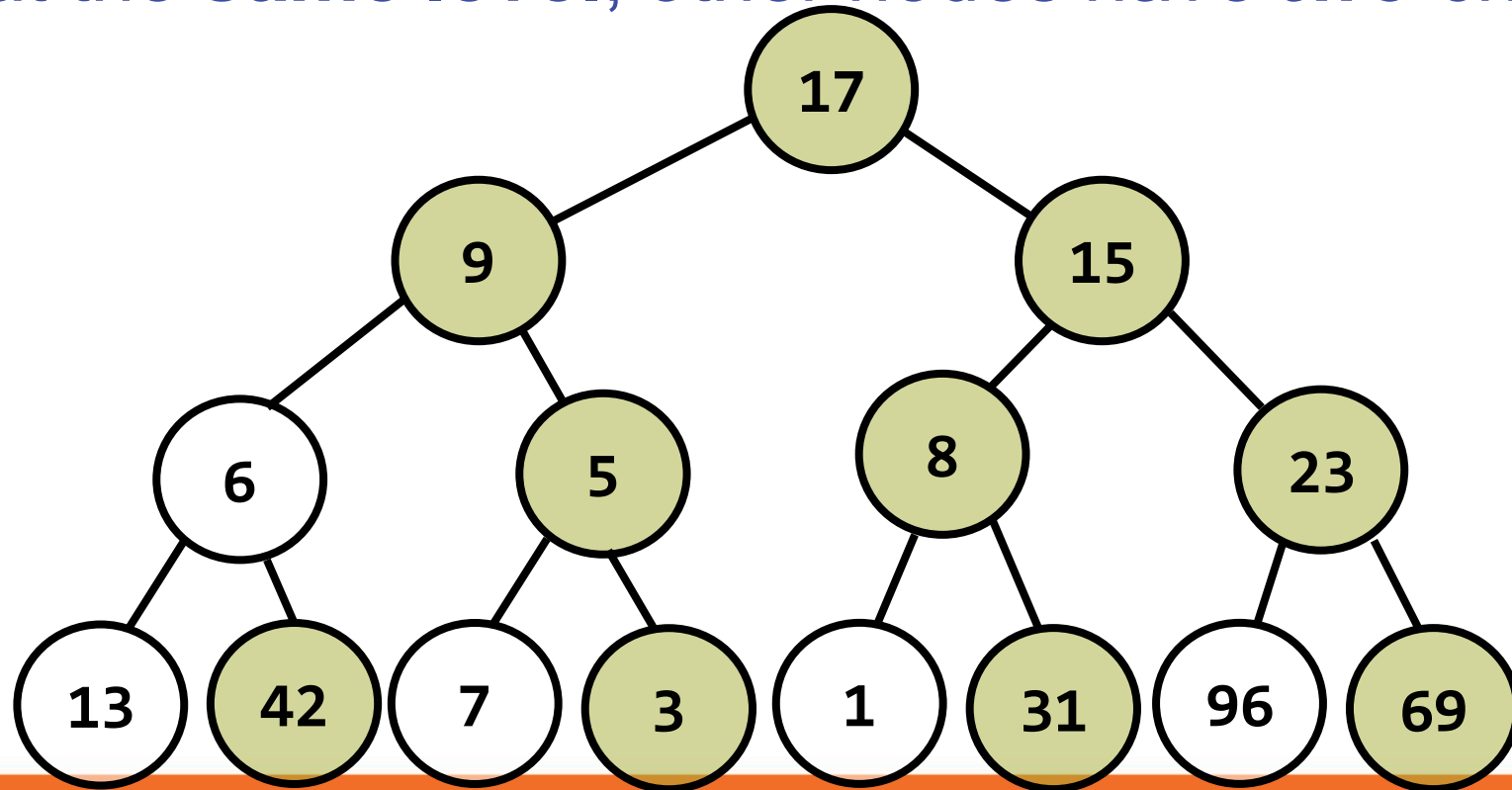
Types of Binary Trees

- **Complete** – nodes are filled **top to bottom** and **left to right**



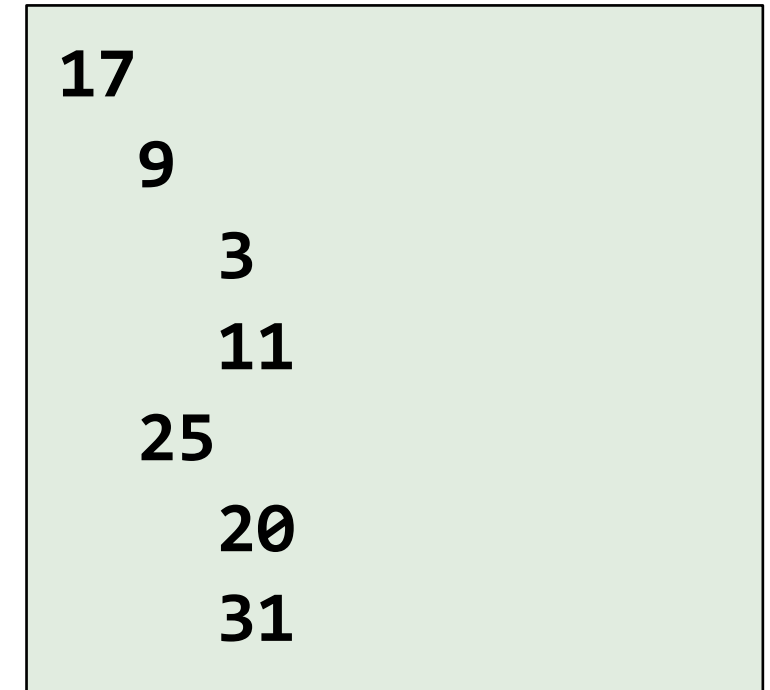
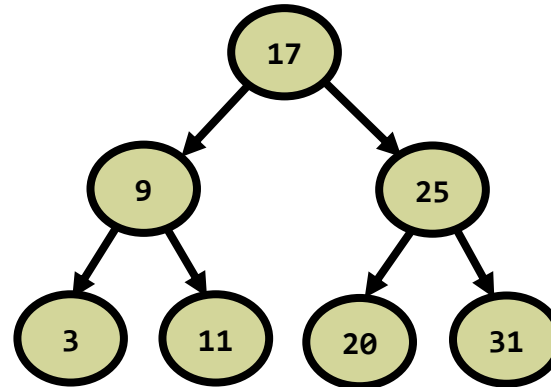
Types of Binary Trees

- **Perfect** – combines **complete** and **full**
 - leafs are at the **same level**, other nodes have **two** children



Problem: Binary Tree Traversals

- Inside the given skeleton
 - Implement **AbstractBinaryTree<E>**
 - Implement **asIndentedPreOrder**, each level indented +2



- **preOrder, inOrder and postOrder**
 - Return the nodes as list **List<AbstractBinaryTree<E>>**

Solution: BT Traversals - Constructor

- Fields and constructor:

```
public class BinaryTree<E> implements AbstractBinaryTree<E> {  
    private E key;  
    private BinaryTree<E> left;  
    private BinaryTree<E> right;  
  
    public BinaryTree(E key, BinaryTree<E> left, BinaryTree<E> right) {  
        this.key = key;  
        this.left = left;  
        this.right = right;  
    }  
}
```

Solution: BT Traversals - Print

```
public String asIndentedPreOrder(int indent) {  
    String out = createPadding(indent) + getKey();  
    if (getLeft() != null) {  
        out += "\n" + getLeft().asIndentedPreOrder(indent + 2);  
    }  
    if (getRight() != null) {  
        out += "\n" + getRight().asIndentedPreOrder(indent + 2);  
    }  
    return out;  
}
```

Process Node

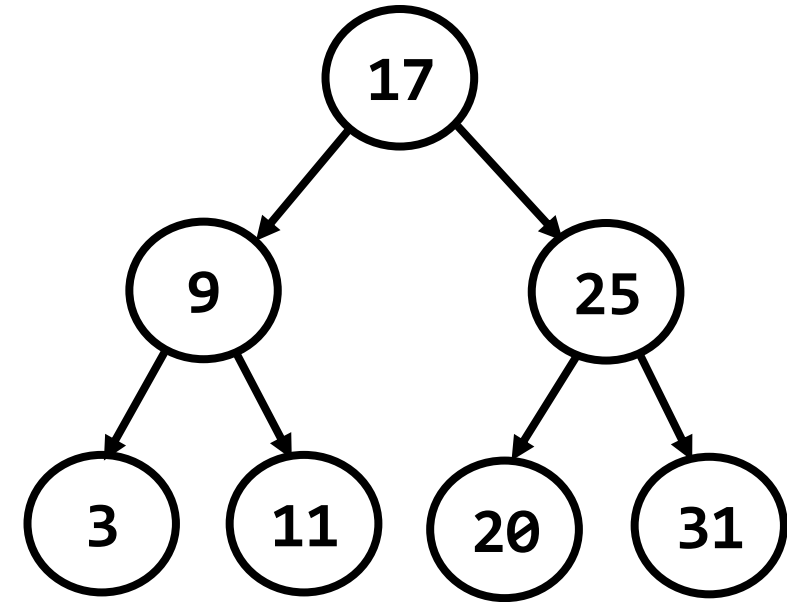
Traverse Left

Traverse Right

Binary Trees Traversal: Pre-order

- Root → Left → Right

```
preOrder (node) {  
    if (node != null) {  
        print node.key  
        preOrder(node.left)  
        preOrder(node.right)  
    }  
}
```

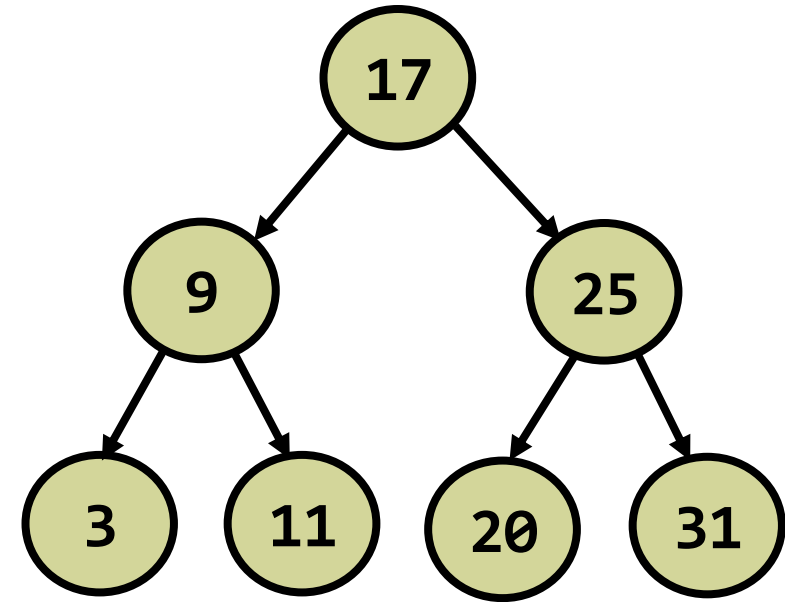


→ 17 9 3 11 25 20 31

Binary Trees Traversal: In-order

- Left → Root → Right

```
inOrder (node) {  
    if (node != null) {  
        inOrder(node.left)  
        print node.key  
        inOrder(node.right)  
    }  
}
```

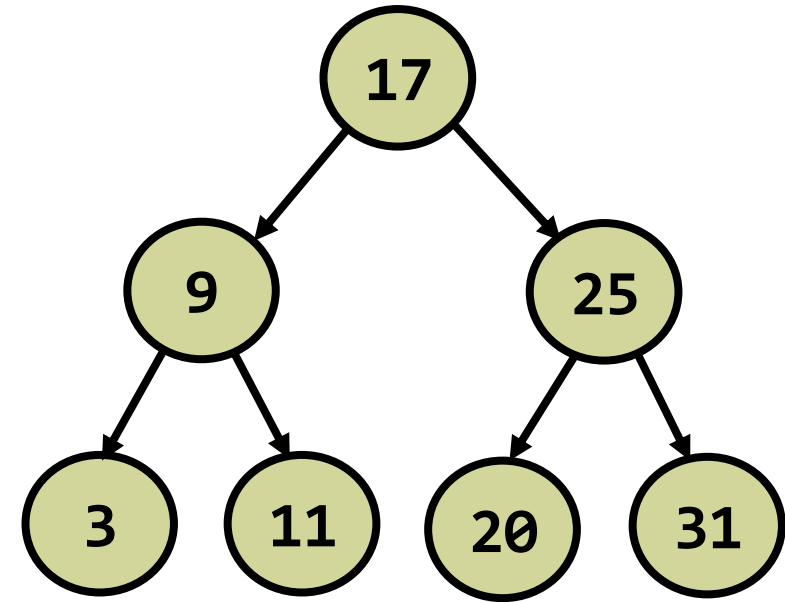


→ 3 9 11 17 20 25 31

Binary Trees Traversal: Post-order

- Left → Right → Root

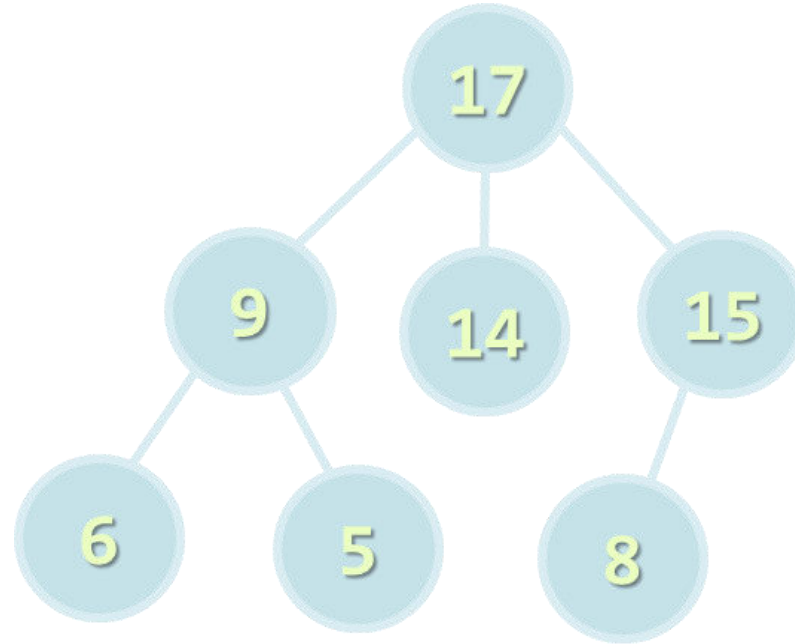
```
postOrder (node) {  
    if (node != null) {  
        postOrder(node.left)  
        postOrder(node.right)  
        print node.key  
    }  
}
```



→ 3 11 9 20 31 25 17

Solution: BT Traversals - forEachInOrder

```
public void forEachInOrder(Consumer<E> consumer) {  
    if (this.getLeft() != null) {  
        this.getLeft().forEachInOrder(consumer);  
    }  
    consumer.accept(this.getKey());  
    if (this.getRight() != null) {  
        this.getRight().forEachInOrder(consumer);  
    }  
}
```



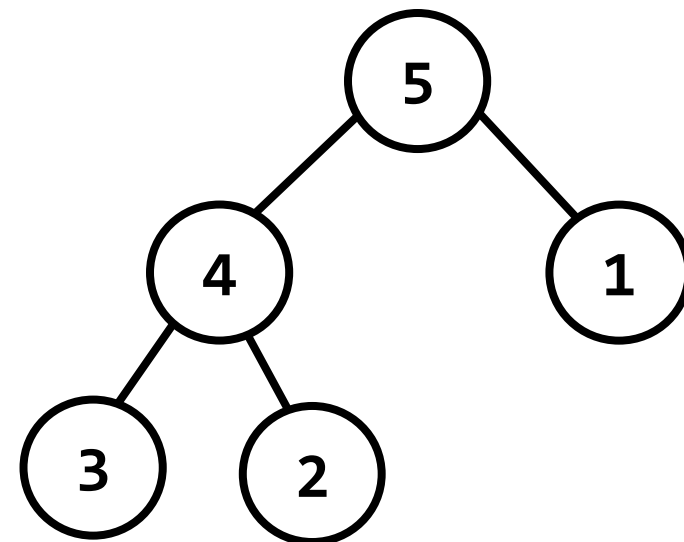
Heap, Binary Heap

What is Heap?

- **Heap**
 - Tree-based data structure
 - Stored in an array
- Heaps hold the **heap property** for each node:
 - **Min Heap**
 - $\text{parent} \leq \text{children}$
 - **Max Heap**
 - $\text{parent} \geq \text{children}$

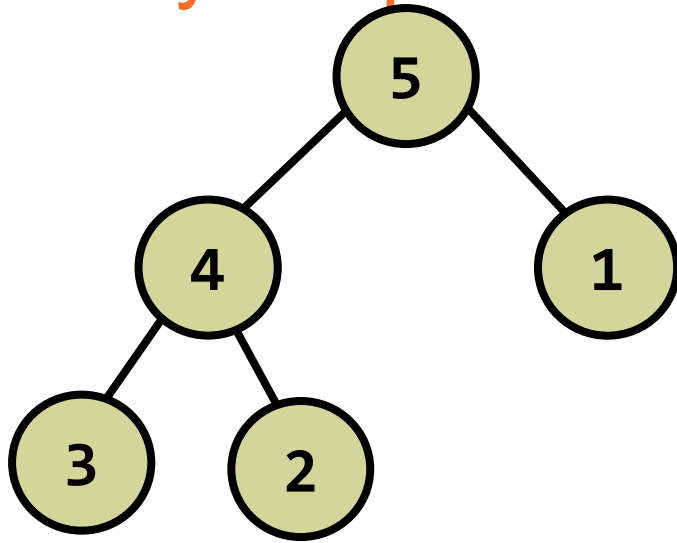
Binary Heap

- **Binary heap**
 - Represents a Binary Tree
- **Shape property** - Binary heap is a **complete binary tree**:
 - Every level, except the last, is **completely filled**
 - Last is filled **from left to right**



Binary Heap – Array Implementation

- Binary heap can be efficiently stored in an array



heap and shape
properties are satisfied

5	4	1	3	2
0	1	2	3	4

- $\text{Parent}(i) = (i - 1) / 2$
- $\text{Left}(i) = 2 * i + 1$; $\text{Right}(i) = 2 * i + 2$

Heap Insertion

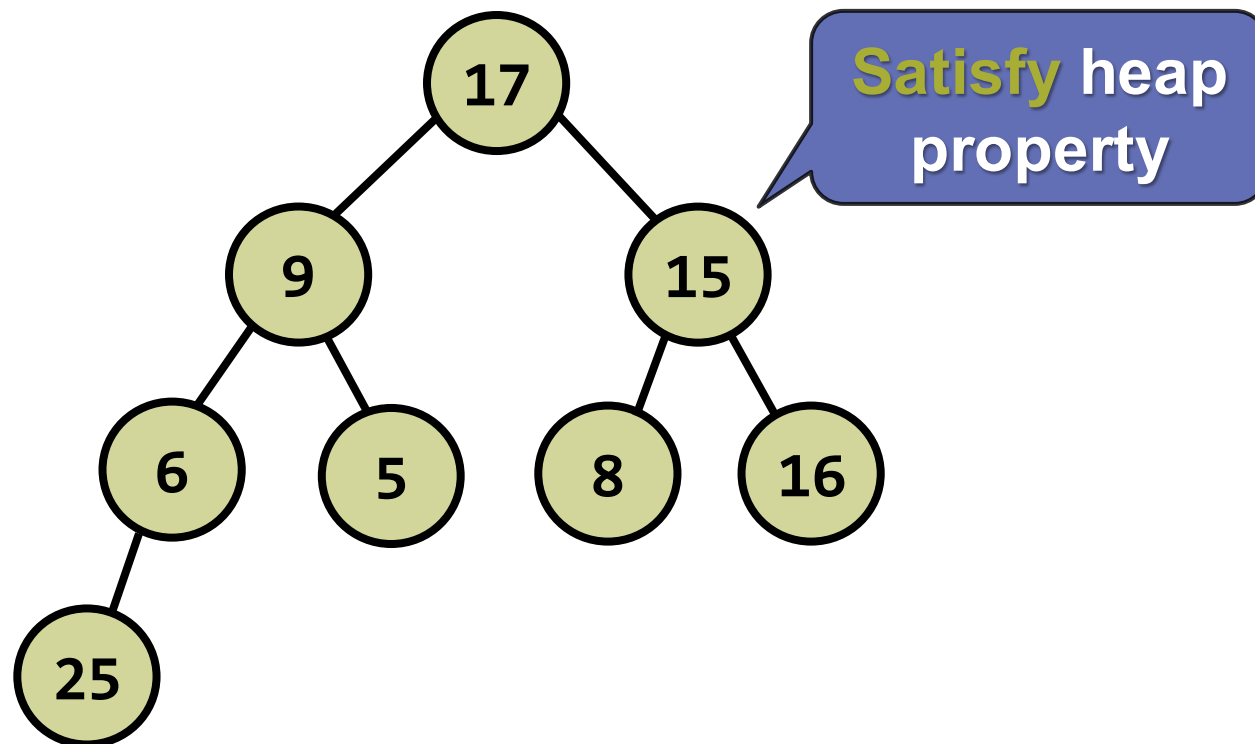
- To preserve heap properties:

- Insert at the end
- Heapify element up

Promote while
element > parent

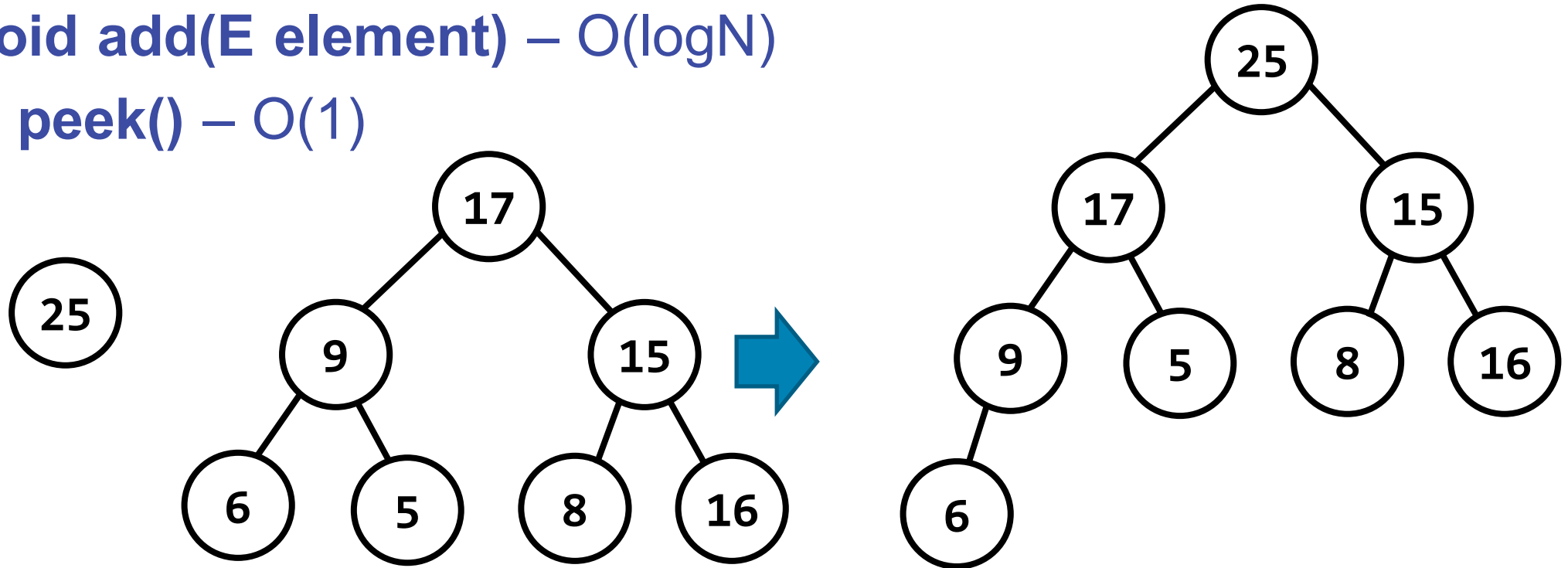
- Right: Max Heap

- Insert 16
- Insert 25



Problem: Heap Add and Peek

- Implement a max **MaxHeap**<E> with:
 - **int size()**
 - **void add(E element) – $O(\log N)$**
 - **E peek() – $O(1)$**



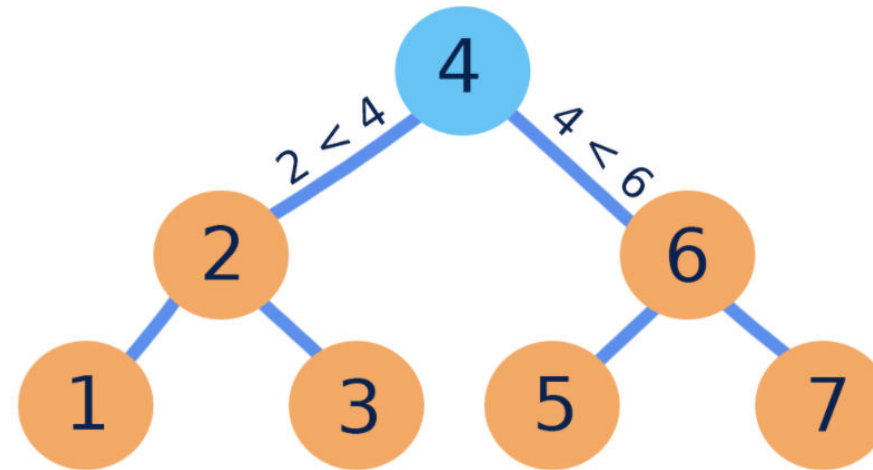
Solution: Heap Add and Peek (1)

```
public class MaxHeap<E> extends Comparable<E>> implements  
Heap<E> {  
    // TODO: store the elements  
    @Override  
    public void add(E element) {  
        this.elements.add(element);  
        this.heapifyUp(this.size() - 1);  
    }  
}
```

Solution: Heap Add and Peek (2)

```
private void heapifyUp(int index) {  
    while (index > 0 && less(parent(index), get(index))) {  
        int parentAt = getParentAt(index);  
        Collections.swap(this.elements, parentAt, index);  
        index = parentAt;  
    }  
}  
  
// TODO: Implement less(), parent() and getParentAt()
```

Binary Search Trees



Two Children at Most

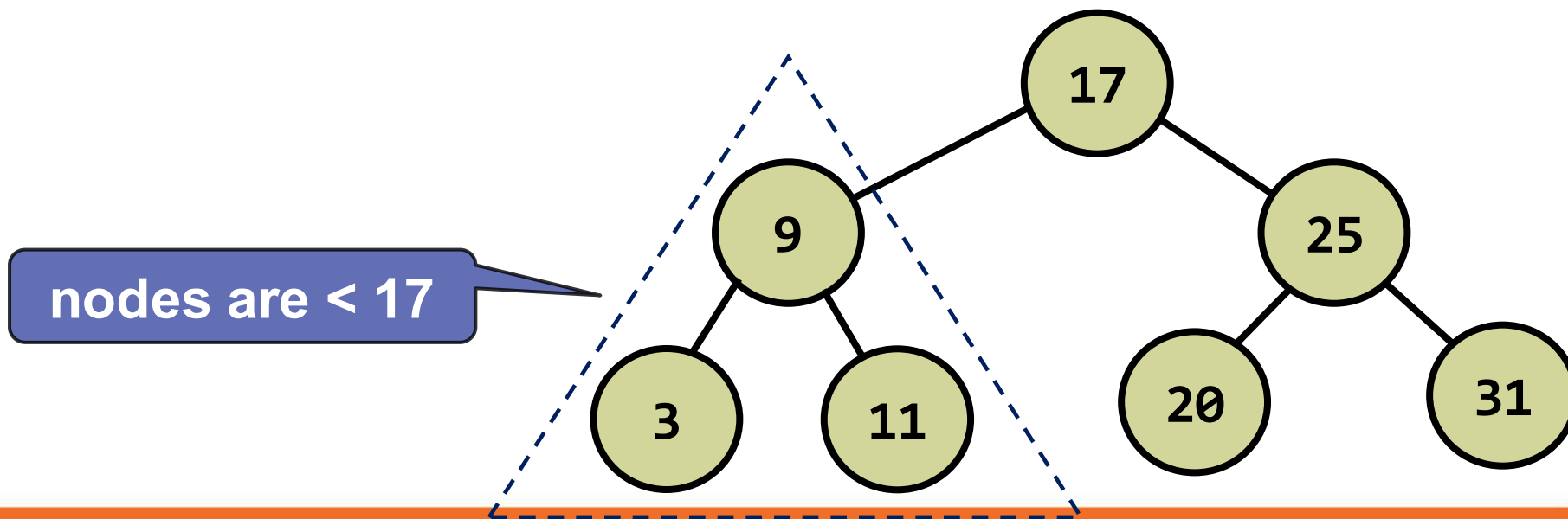
Binary Search Trees

- **Binary search trees are ordered**

- For each node x

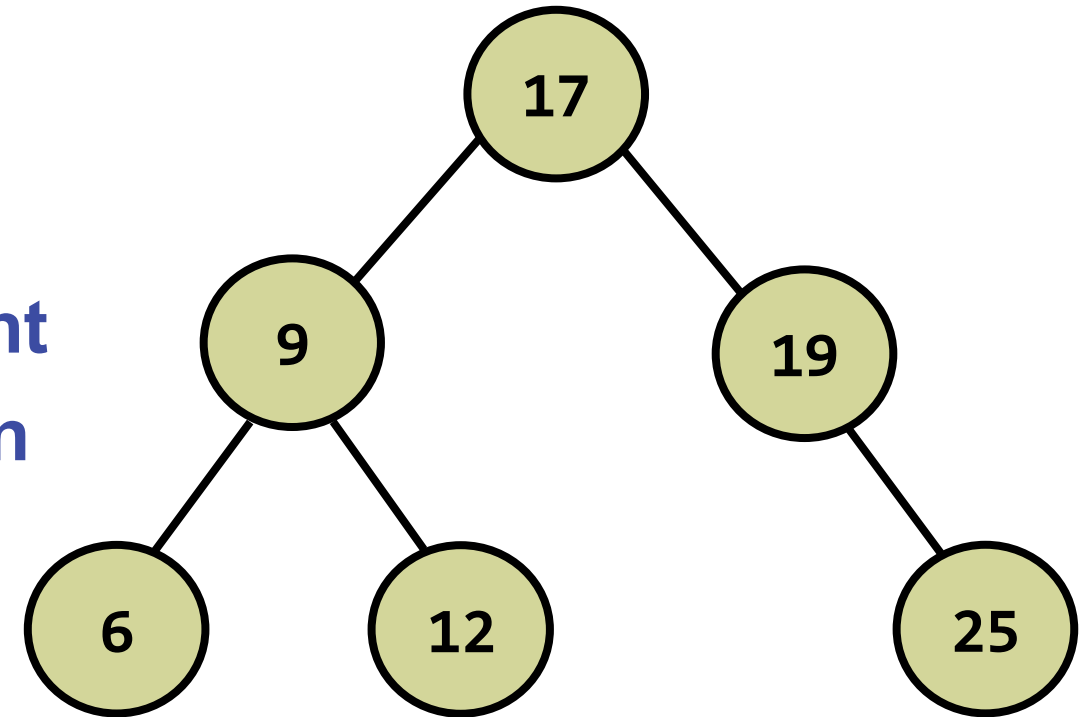
- Elements in left subtree of x are $< x$
- Elements in right subtree of x are $> x$

what about ==



BST - Search

- Search for **x** in BST
 - if node is not null
 - if $x < \text{node.value}$ → go left
 - else if $x > \text{node.value}$ → go right
 - else if $x == \text{node.value}$ → return



Search **12** → 17 9 **12**

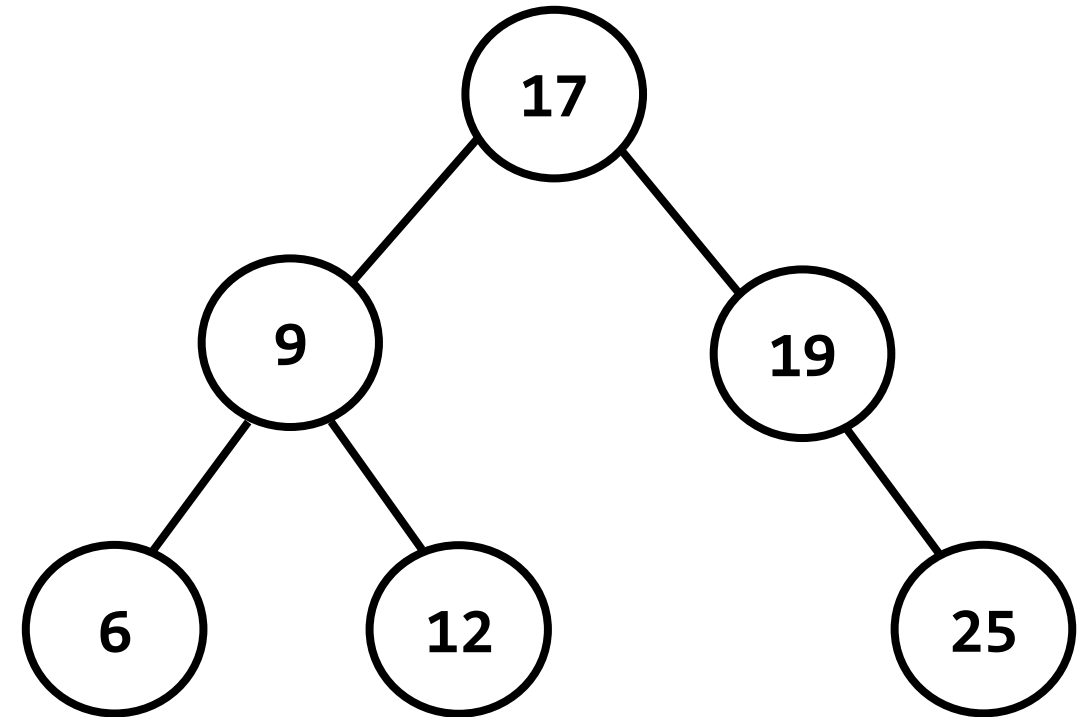
Search **27** → 17 19 25 **null**

BST - Insert

- Insert **x** in BST
 - if node is **null** → insert **x**
 - else if $x < \text{node.value}$ → **go left**
 - else if $x > \text{node.value}$ → **go right**
 - else → node **exists**

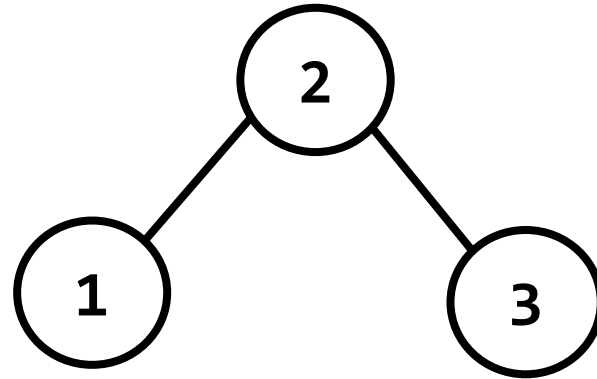
Insert **12** → 17 9 **12** return

Insert **27** → 17 19 25 **null(insert)**



Problem: BST

- You are given a skeleton
 - Implement **AbstractBinarySearchTree<E>**
 - **bool contains(E element)**
 - **void insert(E element)**



Solution: BST Contains

```
public boolean contains(E element) {  
    Node<E> current = this.root;  
    while (current != null){  
        if (element.compareTo(current.value) < 0){  
            current = current.leftChild;  
        } else if (element.compareTo(current.value) > 0){  
            current = current.rightChild;  
        } else {  
            break;  
        }  
    }  
    return current != null;  
}
```

Solution: BST Insert

```
public void insert(E element) {  
    if (this.root == null) {  
        this.root = new Node<>(element);  
    } else {  
        // TODO: Find the place to insert  
        if (parent.value.compareTo(element) > 0){  
            parent.leftChild = new Node<>(element);  
        } else {  
            parent.rightChild = new Node<>(element);  
        }  
    }  
}
```

Problem: BST Search

- Implement:
 - **BST<E> search(E value)**
- Make sure the method works for:
 - empty tree
 - tree with one element
 - tree with two elements - root + left/right
 - tree with multiple elements

Solution: BST Search

```
public AbstractBinarySearchTree<E> search(E element) {  
    Node<E> current = this.root;  
    // TODO: Find the node with the element  
    return new BinarySearchTree<>(current);  
}
```


Solution: BST Search (2)

```
private BinarySearchTree(Node<E> root) {  
    this.copy(root);  
}  
  
private void copy(Node<E> node) {  
    if (node == null) return;  
  
    this.insert(node.value);  
    this.copy(node.leftChildre);  
    this.copy(node.rightChildren);  
}
```

Pre-Order
Traversal

BST - Search Operation Speed - Quiz

- What is the speed of the **search(E)** operation on BST?
 - $O(n)$
 - $O(\log(n))$
 - $O(1)$

BST - Search Operation Speed - Answer

- What is the speed of the **search(E)** operation on BST?

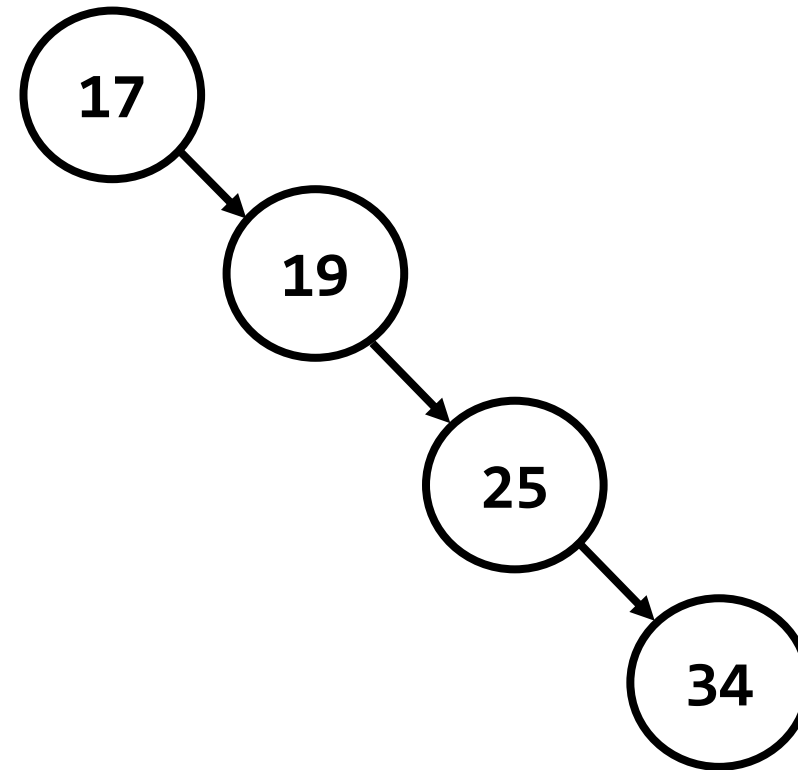
- $O(n)$



- $O(\log(n))$

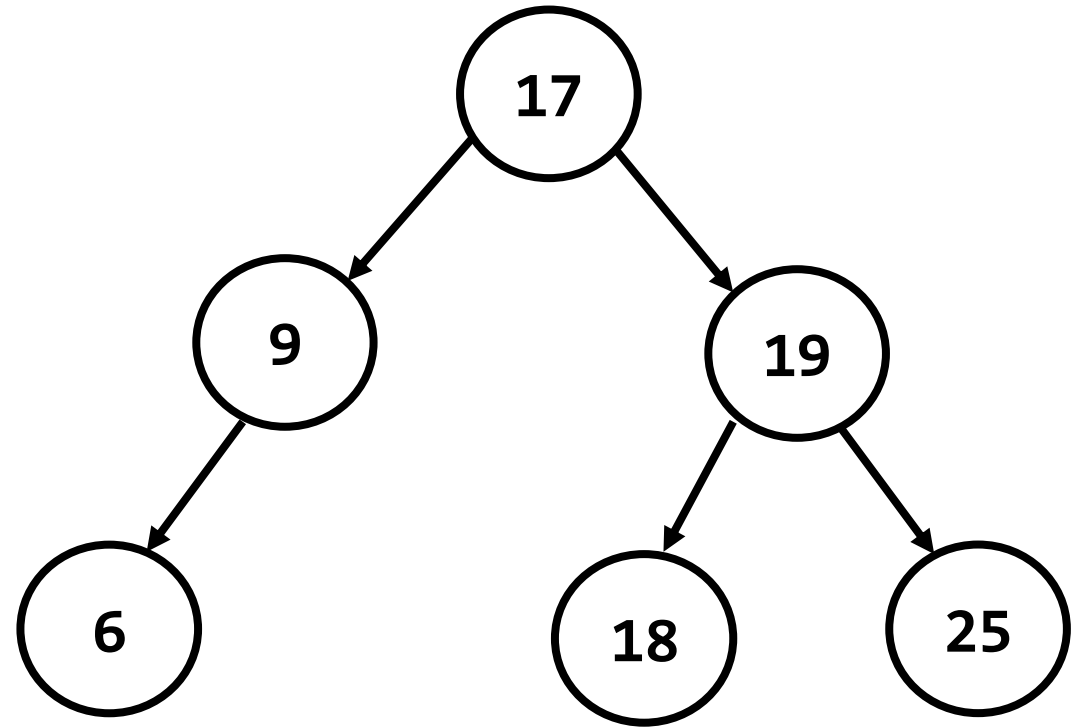


- $O(1)$



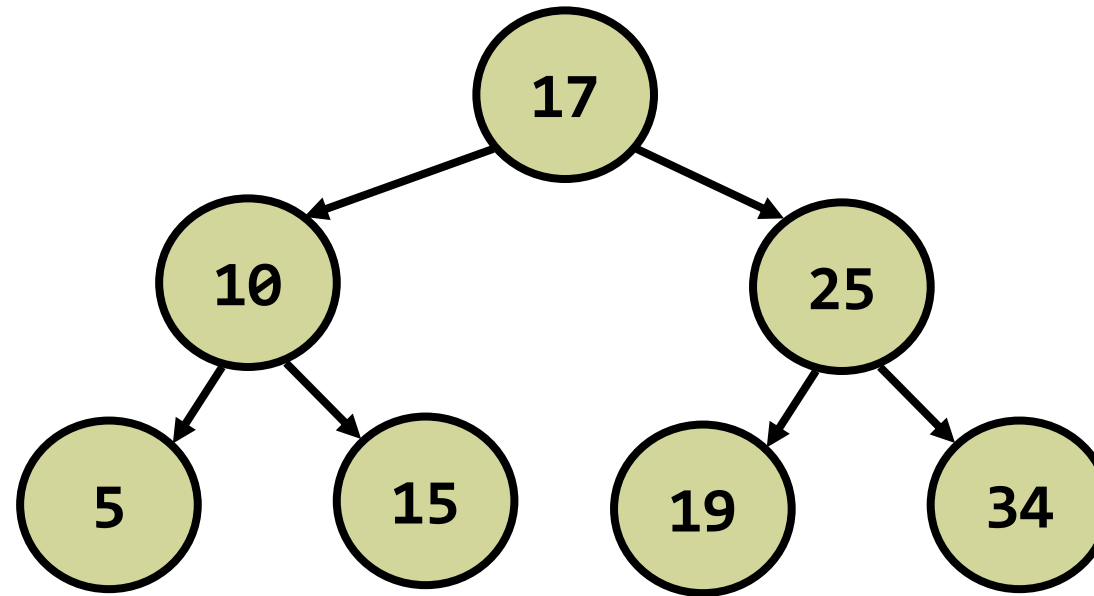
Binary Search Trees – Operation Speed

- Insert – **height** of tree
- Search – **height** of tree



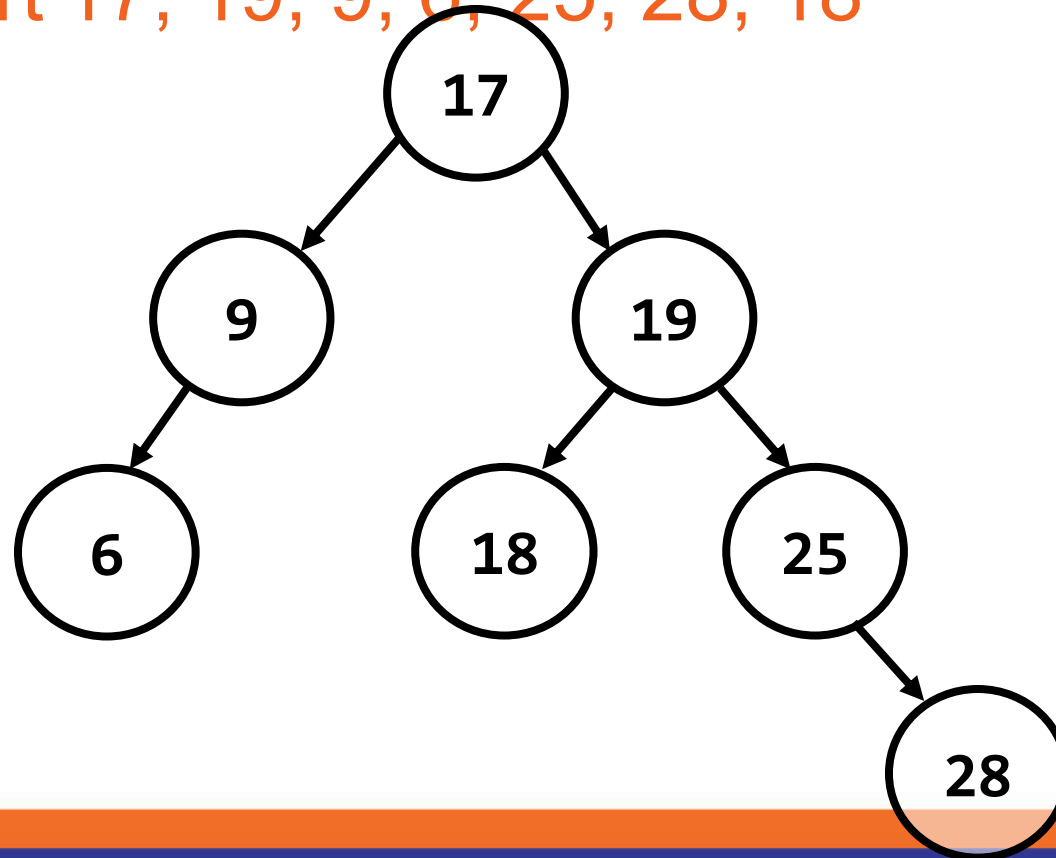
Binary Search Trees – Best Case

- Example: Insert 17, 10, 25, 5, 15, 19, 34



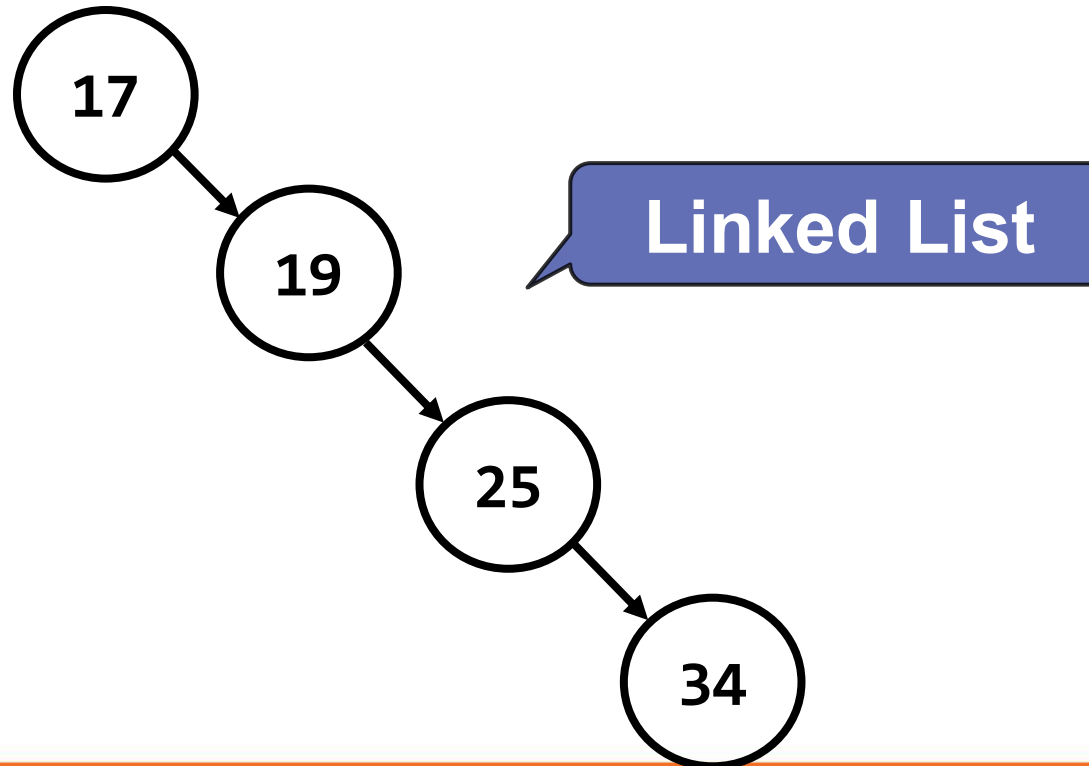
Binary Search Trees – Average Case

- You can insert values in ever **random** order
- Example: Insert 17, 19, 9, 6, 25, 28, 18



Binary Search Trees – Worst Case

- You can insert values in ever **increasing/decreasing** order
- Example: Insert 17, 19, 25, 34



Balanced Binary Search Trees

- Binary search trees can be **balanced**
 - Balanced trees have for each node
 - Nearly equal number of nodes in its subtrees
 - **Balanced trees have height of $\sim \log(n)$**

Summary

- **Binary** trees have 0 or 2 children
- **Heaps** are used to implement priority queues
- Binary Heaps have tree-like structure
- **Efficient** operations
 - Add
 - Find min
 - Remove min