

Data Structures and Algorithms

LECTURE 04: ALGORITHM ANALYSIS

- Algorithmic Complexity
- Time Complexity
- Asymptotic notations
- Brute Force

Algorithm Analysis

- Why should we analyze algorithms?
 - Predict the **resources** the algorithm will need
 - Computational time (**CPU** consumption)
 - Memory space (**RAM** consumption)
 - Communication **bandwidth** consumption
 - **Hard disk** operations

Problem: Get Number of Steps

- Calculate maximum steps to find the result

```
long getOperationsCount(int n) {  
    long counter = 0;  
    for (int i = 0; i < n; i++)  
        for (int j = 0; j < n; j++)  
            counter++;  
    return counter;  
}
```

Solution:

$$T(n) = 3(n^2) + 3n + 3$$

- The input(n) of the function is the main source of steps growth

Simplifying Step Count

- Some parts of the equation **grow much faster** than others
 - $T(n) = 3(n^2) + 3n + 3$
 - We can **ignore** some part of this equation
 - Higher terms **dominate** lower terms – $n > 2$, $n^2 > n$, $n^3 > n^2$
 - Multiplicative constants can be **omitted** – $12n \rightarrow n$, $2n^2 \rightarrow n^2$
- The previous solution becomes $\approx n^2$

Time Complexity

- **Worst-case**
 - An **upper** bound on the running time
- **Average-case**
 - **Average** running time
- **Best-case**
 - The **lower** bound on the running time
(the optimal case)

Time Complexity

- Therefore, we need to measure **all** the possibilities:



Time Complexity

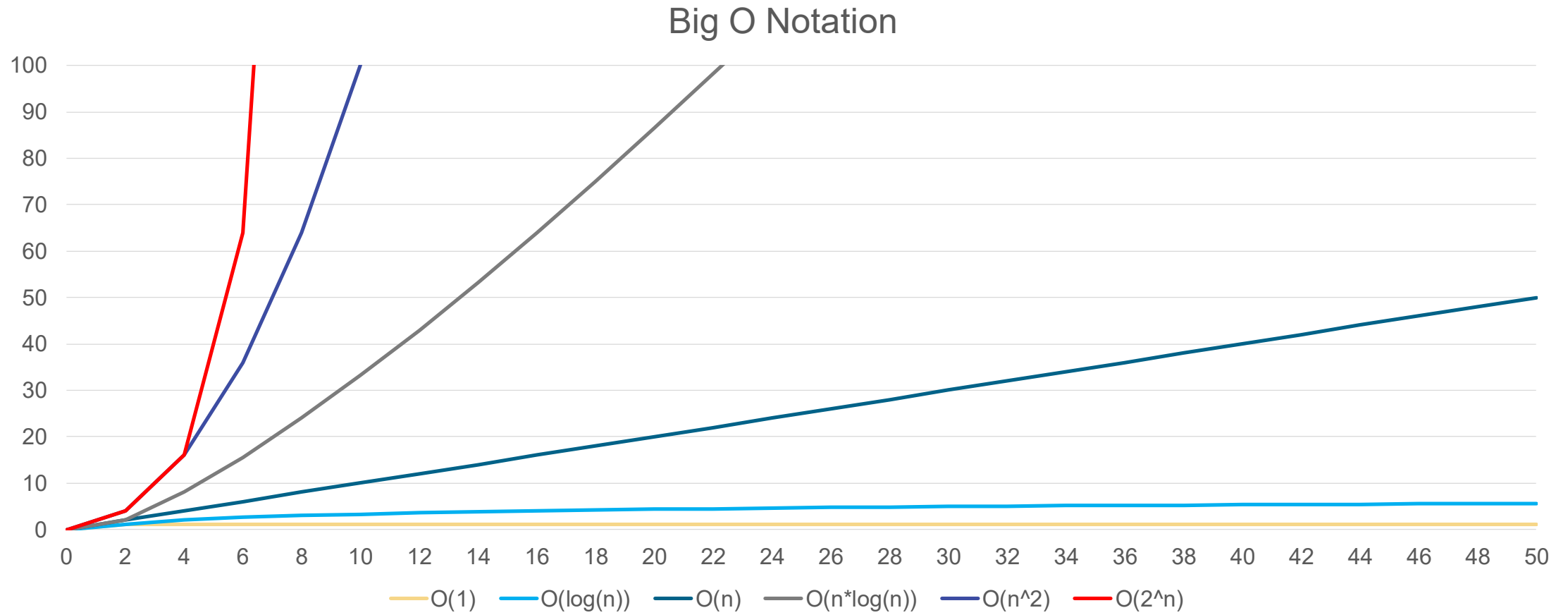
- From the previous chart we can deduce:
 - For smaller size of the input (n) we **don't care much for the runtime**. So we measure the time as n approaches **infinity**
 - If an algorithm **has to scale**, it **should compute** the result within a **finite and practical time**
 - We're concerned about the **order of an algorithm's complexity**, not the actual time in terms of **milliseconds**

Asymptotic notations

- **Asymptotic notations** are descriptions that allow us to examine an algorithm's running time by expressing its **performance** as the input size, **n**, of an algorithm or a function **f increases**. There are **three** common asymptotic notations:
 - Big O – $O(f(n))$
 - Big Theta – $\Theta(f(n))$
 - Big Omega – $\Omega(f(n))$

Asymptotic Functions

- Below are some examples of common algorithmic growth:



Typical Complexities

Complexity	Notation	Description
constant	$O(1)$	$n = 1\ 000 \rightarrow 1\text{-}2$ operations
logarithmic	$O(\log n)$	$n = 1\ 000 \rightarrow 10$ operations
linear	$O(n)$	$n = 1\ 000 \rightarrow 1\ 000$ operations
linearithmic	$O(n \cdot \log n)$	$n = 1\ 000 \rightarrow 10\ 000$ operations
quadratic	$O(n^2)$	$n = 1\ 000 \rightarrow 1\ 000\ 000$ operations
cubic	$O(n^3)$	$n = 1\ 000 \rightarrow 1\ 000\ 000\ 000$ operations
exponential	$O(n^n)$	$n = 10 \rightarrow 10\ 000\ 000\ 000$ operations

Time Complexity and Program Speed

Complexity	10	20	50	100	1 000	10 000	100 000
$O(1)$	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s
$O(\log n)$	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s
$O(n)$	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s
$O(n \cdot \log n)$	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s
$O(n^2)$	< 1 s	< 1 s	< 1 s	< 1 s	< 1 s	2 s	3-4 min
$O(n^3)$	< 1 s	< 1 s	< 1 s	< 1 s	20 s	5 hours	231 days
$O(2^n)$	< 1 s	< 1 s	260 days	hangs	hangs	hangs	hangs
$O(n!)$	< 1 s	hangs	hangs	hangs	hangs	hangs	hangs
$O(n^n)$	3-4 min	hangs	hangs	hangs	hangs	hangs	hangs

Brute-Force Algorithms

- Trying all possible combinations
- Picking the best solution
- Usually slow and inefficient



Brute-Force Algorithms



Brute-Force Algorithms

0 0 0 0 1

Brute-Force Algorithms

0 0 0 0 2

Brute-Force Algorithms



$10 \times 10 \times 10 \times 10 \times 10 = 100,000$ combinations

Summary

- **Algorithmic Complexity**
- Time Complexity