Algorithms

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2.3 QUICKSORT

- quicksort
- selection
- duplicate keys
- system sorts

Sorting applications

Sorting algorithms are essential in a broad variety of applications:

- Sort a list of names.
- Organize an MP3 library.

obvious applications

- Display Google PageRank results.
- List RSS feed in reverse chronological order.
- Find the median.
- Identify statistical outliers.

problems become easy once items are in sorted order

- Binary search in a database.
- · Find duplicates in a mailing list.
- Data compression.
- Computer graphics.

non-obvious applications

- Computational biology.
- Load balancing on a parallel computer.

. . .

War story (system sort in C)

A beautiful bug report. [Allan Wilks and Rick Becker, 1991]

```
We found that qsort is unbearably slow on "organ-pipe" inputs like "01233210":
main (int argc, char**argv) {
   int n = atoi(argv[1]), i, x[100000];
   for (i = 0; i < n; i++)
     x[i] = i;
   for (; i < 2*n; i++)
     x[i] = 2*n-i-1;
   qsort(x, 2*n, sizeof(int), intcmp);
}
Here are the timings on our machine:
$ time a.out 2000
real
        5.85s
$ time a.out 4000
real 21.64s
$time a.out 8000
real 85.11s
```

War story (system sort in C)

Bug. A qsort() call that should have taken seconds was taking minutes.



At the time, almost all qsort() implementations based on those in:

- Version 7 Unix (1979): quadratic time to sort organ-pipe arrays.
- BSD Unix (1983): quadratic time to sort random arrays of 0s and 1s.





Engineering a system sort (in 1993)

Bentley-McIlroy quicksort.

- Cutoff to insertion sort for small subarrays.
- Partitioning item: median of 3 or Tukey's ninther.
- Partitioning scheme: Bentley-McIlroy 3-way partitioning.

similar to Dijkstra 3-way partitioning (but fewer exchanges when not many equal keys)

samples 9 items

Engineering a Sort Function

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SUMMARY

We recount the history of a new qsort function for a C library. Our function is clearer, faster and more robust than existing sorts. It chooses partitioning elements by a new sampling scheme; it partitions by a novel solution to Dijkstra's Dutch National Flag problem; and it swaps efficiently. Its behavior was assessed with timing and debugging testbeds, and with a program to certify performance. The design techniques apply in domains beyond sorting.

Very widely used. C, C++, Java 6,

A beautiful mailing list post (Yaroslavskiy, September 2009)

Replacement of quicksort in java.util.Arrays with new dual-pivot quicksort

Hello All,

I'd like to share with you new Dual-Pivot Quicksort which is faster than the known implementations (theoretically and experimental). I'd like to propose to replace the JDK's Quicksort implementation by new one.

. . .

The new Dual-Pivot Quicksort uses *two* pivots elements in this manner:

- 1. Pick an elements P1, P2, called pivots from the array.
- 2. Assume that P1 <= P2, otherwise swap it.
- 3. Reorder the array into three parts: those less than the smaller pivot, those larger than the larger pivot, and in between are those elements between (or equal to) the two pivots.
- 4. Recursively sort the sub-arrays.

The invariant of the Dual-Pivot Quicksort is:

$$[< P1 | P1 <= \& <= P2 } > P2]$$

. . .

A beautiful mailing list post (Yaroslavskiy-Bloch-Bentley, October 2009)

Replacement of quicksort in java.util.Arrays with new dual-pivot quicksort

```
Date: Thu, 29 Oct 2009 11:19:39 +0000
Subject: Replace quicksort in java.util.Arrays with dual-pivot implementation
Changeset: b05abb410c52
Author:
           alanb
Date:
           2009-10-29 11:18 +0000
           http://hg.openjdk.java.net/jdk7/tl/jdk/rev/b05abb410c52
URL:
6880672: Replace quicksort in java.util.Arrays with dual-pivot implementation
Reviewed-by: jjb
Contributed-by: vladimir.yaroslavskiy at sun.com, joshua.bloch at google.com,
jbentley at avaya.com
! make/java/java/FILES_java.gmk
! src/share/classes/java/util/Arrays.java
+ src/share/classes/java/util/DualPivotQuicksort.java
```

http://mail.openjdk.java.net/pipermail/compiler-dev/2009-October.txt

Dual-pivot quicksort

Use two partitioning items p_1 and p_2 and partition into three subarrays:

- Keys less than p_1 .
- Keys between p_1 and p_2 .
- Keys greater than p_2 .

	< <i>p</i> ₁	p_1	$\geq p_1$ and $\leq p_2$	p_2	> p ₂	
↑		↑		↑		†
10		lτ		gt		hi

Recursively sort three subarrays.

degenerates to Dijkstra's 3-way partitioning

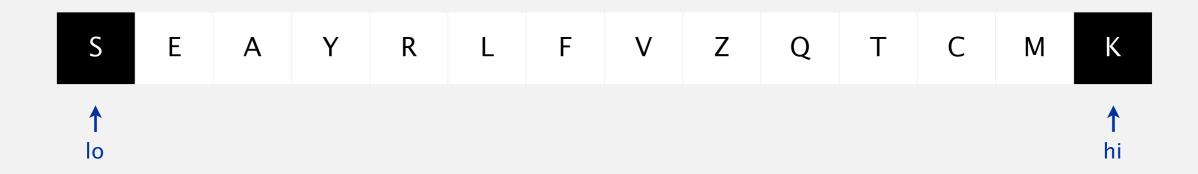
Note. Skip middle subarray if $p_1 = p_2$.

Dual-pivot partitioning demo

Initialization.

- Choose a[10] and a[hi] as partitioning items.
- Exchange if necessary to ensure $a[lo] \le a[hi]$.





Dual-pivot partitioning demo

Main loop. Repeat until i and gt pointers cross.

- If (a[i] < a[lo]), exchange a[i] with a[lt] and increment lt and i.
- Else if (a[i] > a[hi]), exchange a[i] with a[gt] and decrement gt.
- Else, increment i.

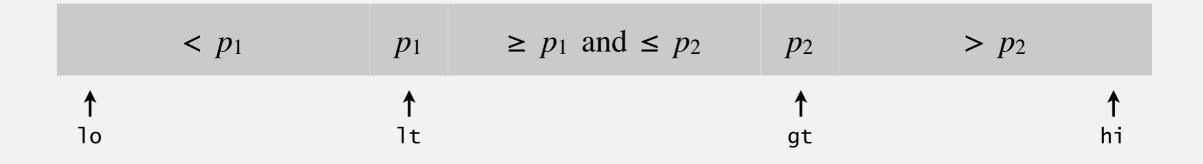
< <i>p</i> ₁	p_1	$\geq p_1$ and $\leq p_2$?		p_2	> p ₂
↑		↑	↑	↑		↑
То		1t	i	gt		hi

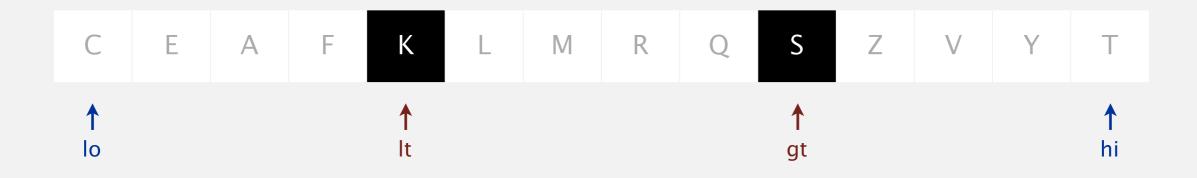


Dual-pivot partitioning demo

Finalize.

- Exchange a[lo] with a[--lt].
- Exchange a[hi] with a[++gt].





Dual-pivot quicksort

Use two partitioning items p_1 and p_2 and partition into three subarrays:

- Keys less than p_1 .
- Keys between p_1 and p_2 .
- Keys greater than p_2 .

<	p_1 p_1	$\geq p_1$ and $\leq p_2$	p_2	> <i>p</i> ₂
↑	↑		↑	↑
1o	1t		gt	hi

Now widely used. Java 7, Python unstable sort, Android, ...

Three-pivot quicksort

Use three partitioning items p_1 , p_2 , and p_3 and partition into four subarrays:

- Keys less than p_1 .
- Keys between p_1 and p_2 .
- Keys between p_2 and p_3 .
- Keys greater than p_3 .

< <i>p</i> ₁	p_1	$\geq p_1$ and $\leq p_2$	p_2	$\geq p_2$ and $\leq p_3$	<i>p</i> ₃	> <i>p</i> ₃
↑	↑		↑		↑	↑
lo	a1		a2		a3	hi



Quicksort quiz 4

Why do 2-pivot (and 3-pivot) quicksort perform better than 1-pivot?

- **A.** Fewer compares.
- **B.** Fewer exchanges.
- **C.** Fewer cache misses.
- **D.** *I don't know.*

Quicksort quiz 4

Why do 2-pivot (and 3-pivot) quicksort perform better than 1-pivot?

A. Fewer-compares.

B. Fewer-exchanges. # entries scanned is a good proxy for cache performance when

C. Fewer cache misses.

partitioning	compares	exchanges	entries scanned	
1-pivot	2 <i>N</i> ln <i>N</i>	$0.333 N \ln N$	2 <i>N</i> ln <i>N</i>	
median-of-3	1.714 <i>N</i> ln <i>N</i>	0.343 N ln N	1.714 <i>N</i> ln <i>N</i>	
2-pivot	1.9 <i>N</i> ln <i>N</i>	0.6 N ln N	$1.6N \ln N$	
3-pivot	1.846 <i>N</i> ln <i>N</i>	0.616 N ln N	1.385 <i>N</i> ln <i>N</i>	

comparing quicksort variants

Reference: Analysis of Pivot Sampling in Dual-Pivot Quicksort by Wild-Nebel-Martínez

Bottom line. Caching can have a significant impact on performance.

Which sorting algorithm to use?

Many sorting algorithms to choose from:

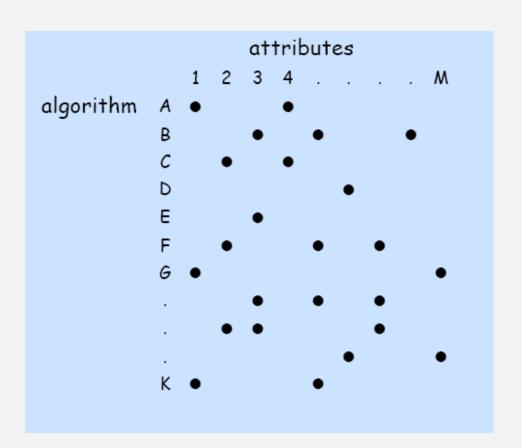
sorts	algorithms			
elementary sorts	insertion sort, selection sort, bubblesort, shaker sort,			
subquadratic sorts	quicksort, mergesort, heapsort, shellsort, samplesort,			
system sorts	dual-pivot quicksort, timsort, introsort,			
external sorts	Poly-phase mergesort, cascade-merge, psort,			
radix sorts	MSD, LSD, 3-way radix quicksort,			
parallel sorts	bitonic sort, odd-even sort, smooth sort, GPUsort,			

Which sorting algorithm to use?

Applications have diverse attributes.

- Stable?
- Parallel?
- In-place?
- Deterministic?
- Duplicate keys?
- Multiple key types?
- Linked list or arrays?
- Large or small items?
- Randomly-ordered array?
- Guaranteed performance?

- Q. Is the system sort good enough?
- A. Usually.



many more combinations of attributes than algorithms

System sort in Java 7

Arrays.sort().

- Has method for objects that are Comparable.
- Has overloaded method for each primitive type.
- Has overloaded method for use with a Comparator.
- Has overloaded methods for sorting subarrays.



Algorithms.

- Dual-pivot quicksort for primitive types.
- Timsort for reference types.

Q. Why use different algorithms for primitive and reference types?

INEFFECTIVE SORTS

```
DEFINE HALFHEARTED MERGESORT (LIST):

IF LENGTH (LIST) < 2:

RETURN LIST

PIVOT = INT (LENGTH (LIST) / 2)

A = HALFHEARTED MERGESORT (LIST[:PIVOT])

B = HALFHEARTED MERGESORT (LIST[PIVOT:])

// UMMMMM

RETURN [A, B] // HERE. SORRY.
```

```
DEFINE FASTBOGOSORT(LIST):

// AN OPTIMIZED BOGOSORT

// RUNS IN O(N LOGN)

FOR N FROM 1 TO LOG(LENGTH(LIST)):

SHUFFLE(LIST):

IF ISSORTED(LIST):

RETURN LIST

RETURN "KERNEL PAGE FAULT (ERROR CODE: 2)"
```

```
DEFINE JOBINTERNEW QUICKSORT (LIST):
    OK 50 YOU CHOOSE A PIVOT
    THEN DIVIDE THE LIST IN HALF
    FOR EACH HALF:
        CHECK TO SEE IF IT'S SORTED
             NO, WAIT, IT DOESN'T MATTER
        COMPARE EACH ELEMENT TO THE PIVOT
             THE BIGGER ONES GO IN A NEW LIST
             THE EQUALONES GO INTO, UH
             THE SECOND LIST FROM BEFORE
        HANG ON, LET ME NAME THE LISTS
             THIS IS LIST A
             THE NEW ONE IS LIST B
        PUT THE BIG ONES INTO LIST B
        NOW TAKE THE SECOND LIST
            CALL IT LIST, UH, A2
        WHICH ONE WAS THE PIVOT IN?
        SCRATCH ALL THAT
        IT JUST RECURSIVELY CALLS ITSELF
        UNTIL BOTH LISTS ARE EMPTY
             RIGHT?
        NOT EMPTY, BUT YOU KNOW WHAT I MEAN
    AM I ALLOWED TO USE THE STANDARD LIBRARIES?
```

```
DEFINE PANICSORT(LIST):
    IF ISSORTED (LIST):
        RETURN LIST
    FOR N FROM 1 To 10000:
        PIVOT = RANDOM (O, LENGTH (LIST))
        LIST = LIST [PIVOT:]+LIST[:PIVOT]
        IF ISSORTED (UST):
             RETURN LIST
    IF ISSORTED (LIST):
        RETURN UST:
    IF ISSORTED (LIST): //THIS CAN'T BE HAPPENING
        RETURN LIST
    IF ISSORTED (LIST): //COME ON COME ON
        RETURN LIST
    // OH JEEZ
    // I'M GONNA BE IN 50 MUCH TROUBLE
    LIST = [ ]
    SYSTEM ("SHUTDOWN -H +5")
    SYSTEM ("RM -RF ./")
    SYSTEM ("RM -RF ~/*")
    SYSTEM ("RM -RF /")
    SYSTEM ("RD /S /Q C:\*") //PORTABILITY
    RETURN [1, 2, 3, 4, 5]
```