

GENERIC PROGRAMMING

PROGRAMMAZIONE CONCORRENTE E DISTR.

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SUMMARY

NATION IN THE PARTY OF THE PART

- Introduction
- Generic classes and methods
- Type variables bounds
- Type erasure
- Restrictions and limitations
- Generics and inheritance
- Wildcard types

INTRODUCTION

- o Generic programming means writing code that can be reused for object of many different types
 - Added to the JDK from version 5.0
 - It tooks 5 years to be formalized (from 1999)
 - First added to support strongly typed collections (as in C++)
 - No generic code is obscure and not typesafe

```
// Before generic classes
ArrayList files = new ArrayList();
// You do not know which kind of object are store inside the list
String filename = (String) files.get(0);
```

- Before generics there was no typesafe checking
- Errors will be find only at runtime
- Violation of FAIL FAST principle

INTRODUCTION



- o Generics solution: type parameters
 - The compiler can perform typesafe checks

```
// After generic classes
ArrayList<String> files = new ArrayList<>();
// You do not need the cast operation anymore
String filename = files.get(0);
```

- Generics make the code safer and easier to read
- But how will you use generic programming?
 - Basic level: just use generic classes
 - Intermediate level: when you encounter your first enigmatic error using generic classes
 - Advanced level: implement your own generic classes

INTRODUCTION





GENERIC CLASSES



- A class with one or more type variable
 - Type variable are introduced after class name, enclosed by angle brackets
 - Type variables are used throughout the class definition

```
class Pair<T, U> {
    private T first;
    private U second;
    // ...
}
```

- o It is common to use upper case letters for type variables
- A type variables is instantiated by substituting types
 - Generic classes act as a factory for ordinary classes

```
Pair<Integer, String> pair = new Pair<>();
Integer first = pair.getFirst();
String second = pair.getSecond();
```

GENERIC METHODS



Also single methods can be defined as generics

 Type variable definition stands behind method's return type

```
public static <T> T getMiddle(T... a) { /* ... */ }
String middle = getMiddle("Riccardo", "Cardin", "Professor");
```

- Can be placed inside both generic an ordinary classes
- You can omit the type variable instantiation
 - But sometimes the compiler gets it wrong...

```
double middle = ArrayAlg.getMiddle(3.14, 1729, 0);
```

- Which type is inferred by the compiler? Number!!!
- In C++, type parameters are after method name
 - This can lead to ambiguites

Type variables bounds



 Rescriction of a type variable to a class that is a subtype of an another type

- Use extends keyword both for classes and interfaces
- It is possibile to use multiple bounds

```
T extends Comparable & Serializable
```

At most one bound can be a class



- The JVM does not have objects of generic types
 - Type variables are erased and replaced with bounding variables or Object type if it ain't any
 - For each generic type a raw type is provided automatically

```
// Remember generic class Pair<T, U> ?
class Pair {
   private Object first;
   private Object second;
   // ...
}
```

- Casts are inserted if necessary to preserve type safety
- Bridge methods are generated to preserve polymorphism
- No new class are created for parametrized types



Translating generic expression

 Compiler inserts casts when the return the return type has been erased

```
Pair<Employee, Salary> buddies = /* ... */;
Employee buddy = buddies.getFirst();
// After type erasure will become
Pair buddies = /* ... */;
Employee buddy = (Employee) buddies.getFirst(); // Returns an Object
```

Translating generic methods

A generic method is not a family of methods

```
public static <T extends Comparable> T min(T[] a)
// becomes after type erasure
public static Comparable min(Comparable[] a)
```

This fact has a series of implications...



Translating generic methods

 Type erasure interferes with polymorphism, then the compiler generates synthetic bridge methods

```
public class Node<T> {
   public T data;
    public Node(T data) { this.data = data; }
    public void setData(T data) {
        System.out.println("Node.setData");
        this.data = data;
public class MyNode extends Node<Integer> {
    public MyNode(Integer data) { super(data); }
    @Override
    public void setData(Integer data) {
        System.out.println("MyNode.setData");
        super.setData(data);
```



Translating generic methods

 After type erasure setData method has wrong type and cannot override parent method

```
public class Node {
    public Object data;
    public Node(Object data) { this.data = data; }
    public void setData(Object data) {
        System.out.println("Node.setData");
        this.data = data;
public class MyNode extends Node {
    public MyNode(Integer data) { super(data); }
    // Not polymorphic anymore :0
    public void setData(Integer data) {
        System.out.println("MyNode.setData");
        super.setData(data);
```

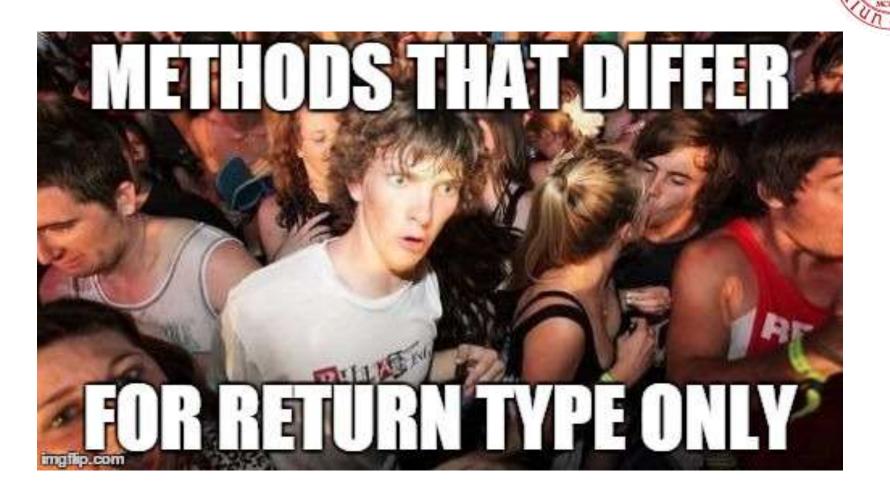


- Translating generic methods
 - Compiler insert a bridge method that overrides the parent method, resolving subtyping
 - Delegation design pattern

```
class MyNode extends Node {
    // Bridge method generated by the compiler
    public void setData(Object data) {
        setData((Integer) data);
    }
    // ...
```

- There are cases that are stanger
 - The type parameters appears only in method return type

```
class DateInterval extends Pair<Date> {
    public Date getSecond() {
        return (Date) super.getSecond().clone(); }
}
Let's try it
yourself
```



RESTRICTIONS AND LIMITATIONS

STATE OF STA

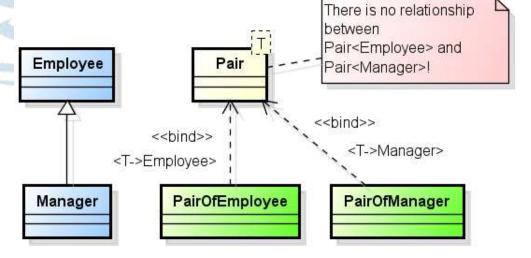
- Type Parameters Cannot Be Instantiated with Primitive Types
- Runtime Type Inquiry Only Works with Raw Types
- You Cannot Create Arrays of Parameterized Types
- Varargs Warnings...
- You Cannot Instantiate Type Variables
- Type Variables Are Not Valid in Static Contexts of Generic Classes
- You Cannot Throw or Catch Instances of a Generic Class
- Beware of Clashes after Erasure

GENERICS AND INHERITANCE



- Generics and inheritance works togheter in an unintuitive way
 - In general, there is no relationship between Pair<S> and Pair<T>, no matter how S and T are related
 Necessary for type safety
 - Generics are said invariant

```
// Type safety restriction
Pair<Manager> managerBuddies =
    new Pair<> (ceo, cfo);
// Illegal, but suppose it
// wasn't
Pair<Employee> employeeBuddies =
    managerBuddies;
employeeBuddies.setFirst(
    lowlyEmployee);
```



GENERICS AND INHERITANCE



 Unfortunately there is a way to bypass this type safety controls

```
Pair<Manager> managerBuddies = new Pair<>(ceo, cfo);
Pair rawBuddies = managerBuddies; // OK
// Only a compile-time warning. A ClassCastException
// is thrown at runtime
rawBuddies.setFirst(new File(". . ."));
```

- But, this is the same behaviour we obtain using older version of Java (≤ 1.5)
 - NEVER use raw type of generics if you can
- Generic classes can extend or implement other generic types
 - For example, ArrayList<T> implements List<T>
 - o So, an ArrayList<Manager> is subtype of List<Manager>



 The type system derived is too rigid: wildcard types help to safetly relax some constraint

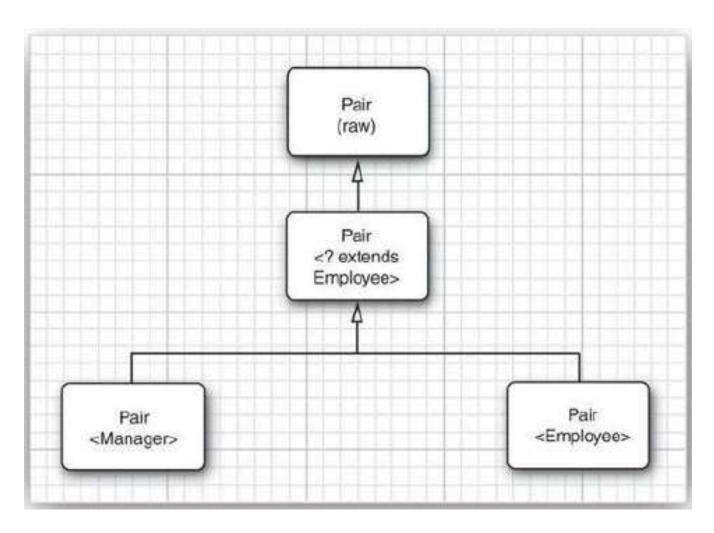
```
// Any generic Pair whose type parameter is a subclass of Employee
Pair<? extends Employee>
```

- Why using wildcard types can we force type safety?
 - The compiler cannot guess the real type the object passed

```
Pair<Manager> managerBuddies = new Pair<> (ceo, cfo);
Pair<? extends Employee> wildcardBuddies = managerBuddies; // OK
wildcardBuddies.setFirst(lowlyEmployee); // compile-time error
```

- Use wildcard type as return types in methods
 - In return types the compiler needs only to know which is the supertype, to allow assignments
 - Type? extends T is said covariant wrt type T







- O How can we use wildcard type for parameters?
 - In Java are available supertype bounds

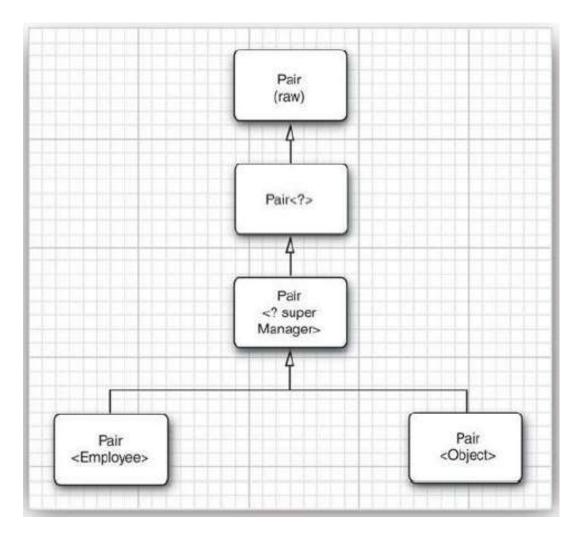
```
// Any generic Pair whose type parameter is restricted to all
// supertypes of Manager
Pair<? super Manager>
```

- A wildcard with a super type bound gives you a behavior that is opposite to that of the wildcards
- You can supply parameters to methods, but you can't use the return values

```
// Any generic Pair whose type parameter is restricted to all
// supertypes of Manager
Pair<? super Manager>
```

- As a return values the only possible assignee type is Object
- o Type ? super T is said contravariant wrt type T







You can even use wildcards with no bounds

```
// It is different from the raw type Pair
Pair<?>
```

- As a return value, can only be assigned to an Object
- As a parameter, no type matches, not even Object
 - A method with a parameter with an unbounded wildcard type can never be called
 - Actually, you can call it passing a null reference...
- Remember, an unbounded wildcard is not a type
 - So you can not use it when a type is requested

```
// This code is not valid, use type variables instead
? t = p.getFirst(); // ERROR
p.setFirst(p.getSecond());
p.setSecond(t);
```





EXAMPLES





https://github.com/rcardin/pcd-snippets

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