

# *Group 6*



Nguyễn Duy Hưng  
20194436



Trần Quốc Lập  
20194443



Hoàng Thiện Tâm  
20194450

## *MINI-PROJECT 20*

Balanced Article Assignment



## Problem

### Contents

Different approaches

Running time

Problem

Input

Problem analysis



Problem

Different approaches

Running time

## Contents

← Problem

 $N$  articles $M$  scientists

DISCARD

NEXT

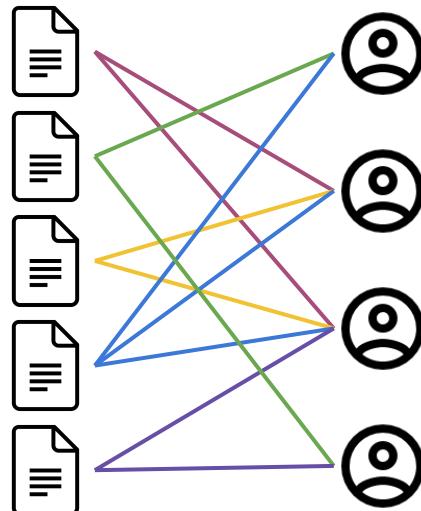


Different approaches

Running time

## Contents

← Problem

 $N$  articles $M$  scientists

1. Each article **can only be given** to a **certain** list of scientists.

DISCARD

NEXT



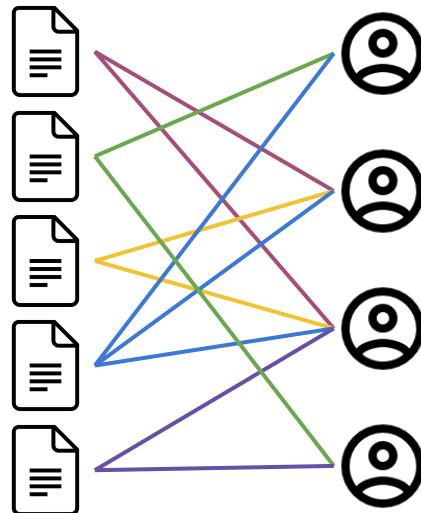
Problem

Different approaches

Running time

## Contents

← Problem

 $N$  articles $M$  scientists

1. Each article **can only be given** to a **certain** list of scientists.
2. Every article must be given to **at least  $K$**  scientists

DISCARD

NEXT



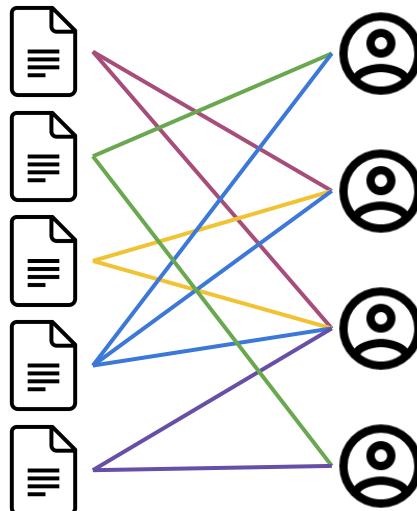
Problem

Different approaches

Running time

## Contents

← Problem

 $N$  articles $M$  scientists

1. Each article **can only be given** to a **certain** list of scientists.
2. Every article must be given to **at least  $K$**  scientists

Find a way of assigning:

the **maximum** of scientists' workload

is **minimized**.

DISCARD

NEXT

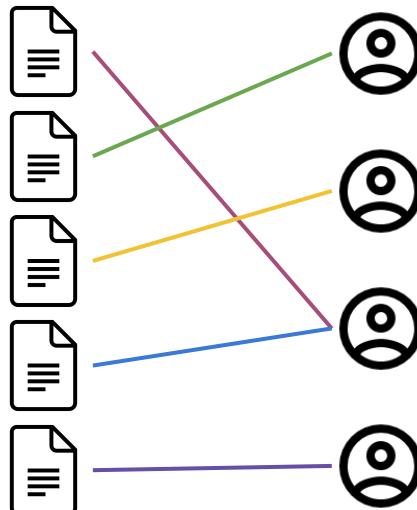


Different approaches

Running time

## Contents

← Problem

 $N$  articles $M$  scientists

1. Each article **can only be given** to a **certain** list of scientists.
2. Every article must be given to **at least  $K$**  scientists

Find a way of assigning:

the **maximum** of scientists' workload

is **minimized**.

DISCARD

NEXT

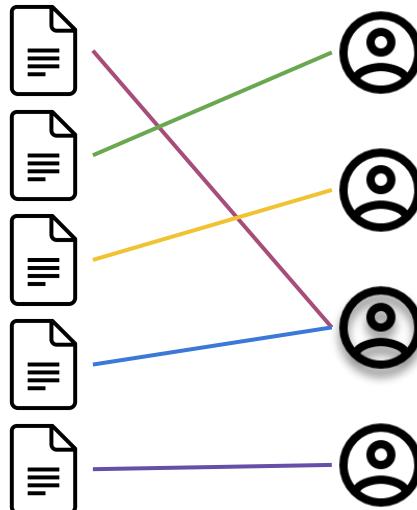


Different approaches

Running time

## Contents

← Problem

 $N$  articles $M$  scientists

1. Each article **can only be given** to a **certain** list of scientists.
2. Every article must be given to **at least  $K$**  scientists

Find a way of assigning:

the **maximum** of scientists' workload

is **minimized**.

DISCARD

NEXT



Different approaches

Running time

## Contents

← Input

10            7            3

5 1 2 4 6 7  
6 1 3 7 2 4 5  
4 2 6 7 1  
5 4 5 2 1 7  
5 3 6 7 1 2  
5 4 1 5 3 6  
3 1 2 6  
4 2 4 7 5  
4 5 3 1 4  
5 5 4 2 1 3

DISCARD

NEXT



Different approaches

Running time

## Contents

← Input

$N = 10 \quad M = 7 \quad K = 3$

A[1] = [1 2 4 6 7]  
A[2] = [1 3 7 2 4 5]  
A[3] = [2 6 7 1]  
A[4] = [4 5 2 1 7]  
A[5] = [3 6 7 1 2]  
A A[6] = [4 1 5 3 6]  
A[7] = [1 2 6]  
A[8] = [2 4 7 5]  
A[9] = [5 3 1 4]  
A[10] = [5 4 2 1 3]

DISCARD

NEXT



Different approaches

Running time

## Contents

← Input

$$N = 10 \quad M = 7 \quad K = 3$$

$$A[1] = [1 2 4 6 7]$$

$$A[2] = [1 3 7 2 4 5]$$

$$A[3] = [2 6 7 1]$$

$$A[4] = [4 5 2 1 7]$$

$$A[5] = [3 6 7 1 2]$$

$$A[6] = [4 1 5 3 6]$$

$$A[7] = [1 2 6]$$

$$A[8] = [2 4 7 5]$$

$$A[9] = [5 3 1 4]$$

$$A[10] = [5 4 2 1 3]$$

$$L = [[1 1 0 1 0 1 1]$$

$$[1 1 1 1 1 0 1]$$

$$[1 1 0 0 0 1 1]$$

$$[1 1 0 1 1 0 1]$$

$$[1 1 1 0 0 1 1]$$

$$[1 0 1 1 1 1 0]$$

$$[1 1 0 0 0 1 0]$$

$$[0 1 0 1 1 0 1]$$

$$[1 0 1 1 1 0 0]$$

$$[1 1 1 1 1 0 0]]$$

DISCARD

NEXT



Different approaches

Running time

## Contents

← Input

$$N = 10 \quad M = 7 \quad K = 3$$

$$A[1] = [1 2 4 6 7]$$

$$A[2] = [1 3 7 2 4 5]$$

$$A[3] = [2 6 7 1]$$

$$A[4] = [4 5 2 1 7]$$

$$A[5] = [3 6 7 1 2]$$

$$A[6] = [4 1 5 3 6]$$

$$A[7] = [1 2 6]$$

$$A[8] = [2 4 7 5]$$

$$A[9] = [5 3 1 4]$$

$$A[10] = [5 4 2 1 3]$$

$$L = [[1 1 0 1 0 1 1]$$

$$[1 1 1 1 1 0 1]$$

$$[1 1 0 0 0 1 1]$$

$$[1 1 0 1 1 0 1]$$

$$[1 1 1 0 0 1 1]$$

$$[1 0 1 1 1 1 0]$$

$$[1 1 0 0 0 1 0]$$

$$[1 1 0 0 0 1 0]$$

$$[0 1 0 1 1 0 1]$$

$$[1 0 1 1 1 0 0]$$

$$[1 1 1 1 1 0 0]]$$

$$X = [[1 1 0 1 0 0 0]$$

$$[1 1 1 0 0 0 0]$$

$$[1 1 0 0 0 1 0]$$

$$[1 1 0 1 0 0 0]$$

$$[0 0 1 0 0 1 1]$$

$$[0 0 1 0 1 1 0]$$

$$[1 1 0 0 0 1 0]$$

$$[1 1 0 0 0 1 0]$$

$$[0 0 1 0 1 0 1]$$

$$[0 0 1 1 1 0 0]$$

$$[0 0 1 1 1 0 0]]$$

DISCARD

NEXT



Different approaches

Running time

## Contents

← Input

$N = 10 \quad M = 7 \quad K = 3$

$A[1] = [1 \ 2 \ 4 \ 6 \ 7]$	$L = [[1 \ 1 \ 0 \ 1 \ 0 \ 1 \ 1]$	$X = [[1 \ 1 \ 0 \ 1 \ 0 \ 0 \ 0]$
$A[2] = [1 \ 3 \ 7 \ 2 \ 4 \ 5]$	$[1 \ 1 \ 1 \ 1 \ 1 \ 0 \ 1]$	$[1 \ 1 \ 1 \ 0 \ 0 \ 0 \ 0]$
$A[3] = [2 \ 6 \ 7 \ 1]$	$[1 \ 1 \ 0 \ 0 \ 0 \ 1 \ 1]$	$[1 \ 1 \ 0 \ 0 \ 0 \ 1 \ 0]$
$A[4] = [4 \ 5 \ 2 \ 1 \ 7]$	$[1 \ 1 \ 0 \ 1 \ 1 \ 0 \ 1]$	$[1 \ 1 \ 0 \ 1 \ 0 \ 0 \ 0]$
$A[5] = [3 \ 6 \ 7 \ 1 \ 2]$	$[1 \ 1 \ 1 \ 0 \ 0 \ 1 \ 1]$	$[0 \ 0 \ 1 \ 0 \ 0 \ 1 \ 1]$
$A[6] = [4 \ 1 \ 5 \ 3 \ 6]$	$[1 \ 0 \ 1 \ 1 \ 1 \ 1 \ 0]$	$[0 \ 0 \ 1 \ 0 \ 1 \ 1 \ 0]$
$A[7] = [1 \ 2 \ 6]$	$[1 \ 1 \ 0 \ 0 \ 0 \ 1 \ 0]$	$[1 \ 1 \ 0 \ 0 \ 0 \ 1 \ 0]$
$A[8] = [2 \ 4 \ 7 \ 5]$	$[0 \ 1 \ 0 \ 1 \ 1 \ 0 \ 1]$	$[0 \ 0 \ 1 \ 0 \ 1 \ 0 \ 1]$
$A[9] = [5 \ 3 \ 1 \ 4]$	$[1 \ 0 \ 1 \ 1 \ 1 \ 0 \ 0]$	$[0 \ 0 \ 1 \ 1 \ 1 \ 0 \ 0]$
$A[10] = [5 \ 4 \ 2 \ 1 \ 3]$	$[1 \ 1 \ 1 \ 1 \ 1 \ 0 \ 0]]$	$[0 \ 0 \ 1 \ 1 \ 1 \ 0 \ 0]$

```
optimal_solution = [[1 2 4] [1 2 3] [1 2 6] [1 2 4] [3 6 7]
                     [3 5 6] [1 2 6] [4 5 7] [3 4 5] [3 4 5]]
```

```
objective_value = 5
```

```
workload = [5, 5, 5, 5, 4, 4, 2]
```

DISCARD

NEXT



Different approaches

Running time

## Contents

[← Problem analysis](#)**Claim**

If  $\exists$  at least 1 solution for the assigning  $\Rightarrow$  also  $\exists$  at least 1 optimal solution.

In that optimal solutions, every article has **exactly  $K$**  scientists to take.

In a certain way of assigning:

- The number of articles that the scientist  $j$  takes on:  $\sum_{i=1}^N x_{ij} l_{ij}$
- The number of scientists that the article  $i$  is given to:  $\sum_{j=1}^M x_{ij} l_{ij}$

[DISCARD](#)[NEXT](#)



Different approaches

Running time

## Contents

← Problem analysis

## Claim

If  $\exists$  at least 1 solution for the assigning  $\Rightarrow$  also  $\exists$  at least 1 optimal solution.

In that optimal solutions, every article has **exactly**  $K$  scientists to take.

Mathematical model:

- 

Minimize

$$\max_{j=1, M} \sum_{i=1}^N x_{ij} l_{ij}$$

- 

Subject to

$$\sum_{j=1}^M x_{ij} l_{ij} = K \quad \forall i = \overline{1, N}$$

DISCARD

NEXT



Problem

Different approaches

Running time

Contents

Ortools using MIP solver

Heuristics

Backtracking

Special algorithm

Problem

Different approaches

Running time

## Ortools using MIP solver

```
def ortools():
    """OR-tools using Mixed-Integer Programming."""
    global objective_value, optimal_solution

    solver = pywraplp.Solver.CreateSolver('SCIP')
    INF = solver.infinity()

    X = np.array([[None]*M for i in range(N)])
    Y = solver.IntVar(-INF, INF, 'Y') # maintaining the objective function

    for i in range(N):
        for j in range(M):
            X[i, j] = solver.IntVar(0, int(L[i, j]), 'X[{}, {}]'.format(i, j))

    for i in range(N):
        solver.Add(K == sum(L[i, j]*X[i, j] for j in range(M)))

    for j in range(M):
        solver.Add(Y >= sum(L[i, j]*X[i, j] for i in range(N)))

    solver.Minimize(Y)
    status = solver.Solve()

    if status == pywraplp.Solver.OPTIMAL:
        objective_value = solver.Objective().Value()
        optimal_solution = [[j+1 for j in range(M) if X[i,j].solution_value() for i in range(N)]
```



Problem

Different approaches

Running time

Contents

Ortools using MIP solver

Heuristics

Backtracking

Special algorithm



Problem

Different approaches

Running time

## Heuristics

## Algorithm description

- Consider the article  $i$  that has not been taken by any scientists yet:
  - Pick out of  $A[i]$  a list of  $K$  scientists who are taking the lightest workload as possible. Assign the article  $i$  for them.
  - Increase the workload of these scientists by 1.
- Consider the article  $i + 1$ .
- ...

## Pseudo-code

Heuristics:

```
if FEASIBLE:  
    workload[j] <- 0 initialize for each scientist j  
  
    for each article i:  
        list <- sorted( [workload[j], j] for j in A[i] )  
        optimal_solution[i] <- list[:K]  
  
        for each scientist j in optimal_solution[i]:  
            increase workload[j] by 1  
        objective_value = max(workload)
```



Problem

Different approaches

Running time

Contents

Ortools using MIP solver

Heuristics

Backtracking

Special algorithm

## Problem

## Different approaches

## Running time

## Backtracking

## Algorithm description

- With an article  $i$ :
  - Consider a  $K$ -subset of  $A[i]$  corresponding to a way of choosing  $K$  scientists who can take on the article  $i$ . There are totally  $\text{len}(A[i]) C_K$  subsets of that way.
    - Assigning the article  $i$  to them by  $\text{solution}[i] \leftarrow$  that  $K$ -subset
    - Increase  $\text{workload}[j]$  by 1.
      - Move to considering the article  $i + 1$ .
      - ...
        - Ultimately, if  $\max(\text{workload}) <$  the smallest  $\text{objective\_value}$  that has been recorded so far, we update.
  - Roll back to considering another  $K$ -subset of  $A[i]$ .
  - ...

Problem

Different approaches

Running time

## Backtracking

### Pseudo-code

Backtracking:

```
initialize workload[j] <- 0 for each scientist j
if the problem is feasible:
    Try(1, 1)      i.e finding the 1st scientist to take the article 1

Try(i, t):
    for each scientist j in A[i] that hasn't been assigned i yet:
        if workload[j] + 1 < objective_value:
            solution[i][t] <- j      i.e set him as the (t)th scientist to takes the article i
            increase workload[j] by 1
            if i = N and t = K:    then Update
            else if t < K:
                Try(i, t + 1)      i.e find (t + 1)th scientist for the article i
            else if i = N:
                Try(i + 1, 1)      i.e move to the next article
                solution[i][t] <- NULL   i.e omit j in order to try other scientists
                decrease workload[j] by 1
```

Update:

```
if max(workload) < objective_value:
    objective_value <- max(workload)
    optimal_solution <- solution
```



Problem

Different approaches

Running time

Contents

Ortools using MIP solver

Heuristics

Backtracking

Special algorithm

Problem

Different approaches

Running time

## Special algorithm

## Algorithm description

$$L = [[0 \ 4 \ 5 \ 4 \ 3] \\ [2 \ 0 \ 1 \ 1 \ 0] \\ [4 \ 1 \ 1 \ 1 \ 1] \\ [4 \ 1 \ 1 \ 1 \ 1] \\ [3 \ 1 \ 1 \ 1 \ 0] \\ [3 \ 1 \ 1 \ 0 \ 1]]$$

Initial state

Ideal:

- We aim to minimize the\_max\_the\_article\_taken\_among\_scientists  
So, I decided to reduce the max of the article till I can't reduce it any more.  
It's not mean that I can reduce the max to any value arbitrarily  
I decided to reduce the max to the predecessor(the number that has the value just only less than the max number. Ex:....)  
If after reduction, we can't reach the predecessor, so we find the solution...  
Else we continue reduce them :)))

Another case: If we don't have predecessor(ex: a sequence numbers have the same value)

So, we reduce 'slowly': each value reduce only 1  
If after reduce, which value doesn't change is our solution

Problem

Different approaches

Running time

## Special algorithm

## Algorithm description

$$L = [[0\ 4\ 5\ 4\ 3]\ [2\ 0\ 1\ 1\ 0]\ [4\ 1\ 1\ 1\ 1]\ [4\ 1\ 1\ 1\ 1]\ [3\ 1\ 1\ 1\ 0]\ [3\ 1\ 1\ 0\ 1]]$$

Initial state

$$L = [[0\ 3\ 4\ 4\ 5]\ [2\ 0\ 0\ 1\ 1]\ [3\ 0\ 1\ 1\ 1]\ [3\ 1\ 1\ 0\ 1]\ [4\ 1\ 1\ 1\ 1]\ [4\ 1\ 1\ 1\ 1]]$$

Problem

Different approaches

Running time

## Special algorithm

## Algorithm description

$$L = [[0\ 4\ 5\ 4\ 3]\newline [2\ 0\ 1\ 1\ 0]\newline [4\ 1\ 1\ 1\ 1]\newline [4\ 1\ 1\ 1\ 1]\newline [3\ 1\ 1\ 1\ 0]\newline [3\ 1\ 1\ 0\ 1]]$$

Initial state

$$L = [[0\ 3\ 4\ 4\ 5]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [4\ 1\ 1\ 1\ 1]\newline [4\ 1\ 1\ 1\ 1]]$$

$$L = [[0\ 3\ 4\ 4\ 4]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [4\ 1\ 1\ 1\ 1]\newline [3\ 1\ 1\ 1\ 0]]$$

Problem

Different approaches

Running time

## Special algorithm

## Algorithm description

$$L = [[0\ 4\ 5\ 4\ 3]\n[2\ 0\ 1\ 1\ 0]\n[4\ 1\ 1\ 1\ 1]\n[4\ 1\ 1\ 1\ 1]\n[3\ 1\ 1\ 1\ 0]\n[3\ 1\ 1\ 0\ 1]]$$

Initial state

$$L = [[0\ 3\ 4\ 4\ 5]\n[2\ 0\ 0\ 1\ 1]\n[3\ 0\ 1\ 1\ 1]\n[3\ 1\ 1\ 0\ 1]\n[4\ 1\ 1\ 1\ 1]\n[4\ 1\ 1\ 1\ 1]]$$
$$L = [[0\ 3\ 4\ 4\ 4]\n[2\ 0\ 0\ 1\ 1]\n[3\ 0\ 1\ 1\ 1]\n[3\ 1\ 1\ 0\ 1]\n[3\ 1\ 1\ 1\ 0]\n[4\ 1\ 1\ 1\ 1]]$$

$$L = [[0\ 3\ 4\ 4\ 4]\n[2\ 0\ 0\ 1\ 1]\n[3\ 0\ 1\ 1\ 1]\n[3\ 1\ 1\ 0\ 1]\n[4\ 1\ 1\ 1\ 1]\n[3\ 1\ 1\ 1\ 0]]$$

Problem

Different approaches

Running time

## Special algorithm

## Algorithm description

$$L = [[0\ 4\ 5\ 4\ 3]\n[2\ 0\ 1\ 1\ 0]\n[4\ 1\ 1\ 1\ 1]\n[4\ 1\ 1\ 1\ 1]\n[3\ 1\ 1\ 1\ 0]\n[3\ 1\ 1\ 0\ 1]]$$

Initial state

$$L = [[0\ 3\ 4\ 4\ 5]\n[2\ 0\ 0\ 1\ 1]\n[3\ 0\ 1\ 1\ 1]\n[3\ 1\ 1\ 0\ 1]\n[4\ 1\ 1\ 1\ 1]\n[4\ 1\ 1\ 1\ 1]]$$

$$L = [[0\ 3\ 4\ 4\ 4]\n[2\ 0\ 0\ 1\ 1]\n[3\ 0\ 1\ 1\ 1]\n[3\ 1\ 1\ 0\ 1]\n[4\ 1\ 1\ 1\ 1]\n[3\ 1\ 1\ 1\ 0]]$$
$$L = [[0\ 3\ 4\ 4\ 4]\n[2\ 0\ 0\ 1\ 1]\n[3\ 0\ 1\ 1\ 1]\n[3\ 1\ 1\ 0\ 1]\n[3\ 1\ 1\ 1\ 0]\n[4\ 1\ 1\ 1\ 1]]$$

$$L = [[0\ 3\ 3\ 3\ 3]\n[2\ 0\ 0\ 1\ 1]\n[3\ 0\ 1\ 1\ 1]\n[3\ 1\ 1\ 0\ 1]\n[2\ 1\ 0\ 1\ 0]\n[2\ 1\ 1\ 0\ 0]]$$

## Problem

## Different approaches

## Running time

## Special algorithm

## Algorithm description

$$L = [[0\ 4\ 5\ 4\ 3]\newline [2\ 0\ 1\ 1\ 0]\newline [4\ 1\ 1\ 1\ 1]\newline [4\ 1\ 1\ 1\ 1]\newline [3\ 1\ 1\ 1\ 0]\newline [3\ 1\ 1\ 0\ 1]]$$

Initial state

$$L = [[0\ 3\ 4\ 4\ 5]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [4\ 1\ 1\ 1\ 1]\newline [4\ 1\ 1\ 1\ 1\ \textcolor{red}{1}]]$$

$$L = [[0\ 3\ 4\ 4\ 4]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [4\ 1\ 1\ 1\ 1]\newline [3\ 1\ 1\ 1\ 1\ \textcolor{brown}{0}]]$$
$$L = [[0\ 3\ 4\ 4\ 4]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [3\ 1\ 1\ 1\ 0]\newline [4\ 1\ 1\ 1\ 1\ \textcolor{red}{1}]]$$

$$L = [[0\ 3\ 3\ 3\ 3]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [2\ 1\ 0\ 1\ 0]\newline [2\ 1\ 1\ 0\ 1\ \textcolor{brown}{0}]]$$
$$L = [[0\ 3\ 3\ 3\ 3]\newline [2\ 0\ 0\ 1\ 1]\newline [2\ 1\ 0\ 1\ 0]\newline [2\ 1\ 1\ 0\ 0]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1\ \textcolor{red}{1}]]$$

Problem

Different approaches

Running time

## Special algorithm

## Algorithm description

$$L = [[0\ 4\ 5\ 4\ 3]\newline [2\ 0\ 1\ 1\ 0]\newline [4\ 1\ 1\ 1\ 1]\newline [4\ 1\ 1\ 1\ 1]\newline [3\ 1\ 1\ 1\ 0]\newline [3\ 1\ 1\ 0\ 1]]$$

Initial state

$$L = [[0\ 3\ 4\ 4\ 5]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [4\ 1\ 1\ 1\ 1]\newline [4\ 1\ 1\ 1\ 1\ \textcolor{red}{1}]]$$


$$L = [[0\ 3\ 4\ 4\ 4]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [4\ 1\ 1\ 1\ 1]\newline [3\ 1\ 1\ 1\ 1\ \textcolor{brown}{0}]]$$

$$L = [[0\ 3\ 4\ 4\ 4]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [3\ 1\ 1\ 1\ 0]\newline [4\ 1\ 1\ 1\ 1\ \textcolor{red}{1}]]$$


$$L = [[0\ 3\ 3\ 3\ 3]\newline [2\ 0\ 0\ 1\ 1]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1]\newline [2\ 1\ 0\ 1\ 0]\newline [2\ 1\ 1\ 0\ 1\ \textcolor{brown}{0}]]$$

$$L = [[0\ 3\ 3\ 3\ 3]\newline [2\ 0\ 0\ 1\ 1]\newline [2\ 1\ 0\ 1\ 0]\newline [2\ 1\ 1\ 0\ 0]\newline [3\ 0\ 1\ 1\ 1]\newline [3\ 1\ 1\ 0\ 1\ \textcolor{red}{1}]]$$


$$L = [[0\ 3\ 3\ 2\ 2]\newline [2\ 0\ 0\ 1\ 1]\newline [2\ 1\ 0\ 1\ 0]\newline [2\ 1\ 1\ 0\ 0]\newline [2\ 0\ 1\ 0\ 1]\newline [2\ 1\ 1\ 0\ 0\ \textcolor{brown}{1}]]$$

## Problem

## Different approaches

## Running time

## Special algorithm

## Algorithm description

$$L = [[0 \ 4 \ 5 \ 4 \ 3] \\ [2 \ 0 \ 1 \ 1 \ 0] \\ [4 \ 1 \ 1 \ 1 \ 1] \\ [4 \ 1 \ 1 \ 1 \ 1] \\ [3 \ 1 \ 1 \ 1 \ 0] \\ [3 \ 1 \ 1 \ 0 \ 1]]$$

Initial state

$$L = [[0 \ 2 \ 2 \ 3 \ 3] \\ [2 \ 1 \ 1 \ 0 \ 0] \\ [2 \ 1 \ 0 \ 1 \ 0] \\ [2 \ 0 \ 0 \ 1 \ 1] \\ [2 \ 0 \ 1 \ 0 \ 1] \\ [2 \ 0 \ 0 \ 1 \ 1]]$$

final state

$$L = [[0 \ 3 \ 4 \ 4 \ 5] \\ [2 \ 0 \ 0 \ 1 \ 1] \\ [3 \ 0 \ 1 \ 1 \ 1] \\ [3 \ 1 \ 1 \ 0 \ 1] \\ [4 \ 1 \ 1 \ 1 \ 1] \\ [4 \ 1 \ 1 \ 1 \ 1 \ 1]]$$


$$L = [[0 \ 3 \ 4 \ 4 \ 4] \\ [2 \ 0 \ 0 \ 1 \ 1] \\ [3 \ 0 \ 1 \ 1 \ 1] \\ [3 \ 1 \ 1 \ 0 \ 1] \\ [4 \ 1 \ 1 \ 1 \ 1] \\ [3 \ 1 \ 1 \ 1 \ 0 \ 1]]$$


$$L = [[0 \ 3 \ 4 \ 4 \ 4] \\ [2 \ 0 \ 0 \ 1 \ 1] \\ [3 \ 0 \ 1 \ 1 \ 1] \\ [3 \ 1 \ 1 \ 0 \ 1] \\ [3 \ 1 \ 1 \ 1 \ 0] \\ [4 \ 1 \ 1 \ 1 \ 1 \ 1]]$$


$$L = [[0 \ 3 \ 3 \ 3 \ 3] \\ [2 \ 0 \ 0 \ 1 \ 1] \\ [2 \ 1 \ 0 \ 1 \ 0] \\ [2 \ 1 \ 1 \ 0 \ 0] \\ [3 \ 0 \ 1 \ 1 \ 1] \\ [3 \ 1 \ 1 \ 0 \ 1]]$$


$$L = [[0 \ 3 \ 3 \ 2 \ 2] \\ [2 \ 0 \ 0 \ 1 \ 1] \\ [2 \ 1 \ 0 \ 1 \ 0] \\ [2 \ 1 \ 1 \ 0 \ 0] \\ [2 \ 0 \ 1 \ 0 \ 1] \\ [2 \ 1 \ 1 \ 0 \ 0 \ 1]]$$

## Problem

## Experiments

## Different approaches

## Running time

Data	Size	MIP	Backtracking	Special approach	Heuristics		Opt
					Time	Result	
Data 1	5	0.031	0.00016	0.0004	0.00011	3	3
Data 2	10	0.029	0.08074	0.0007	0.00013	5	5
Data 3	20	0.036	–	0.0018	0.00021	9	8
Data 4	50	0.297	–	0.0066	0.00081	22	22
Data 5	100	1.047	–	0.0509	0.00174	44	43
Data 6	500	59.149	–	1.1152	0.01484	129	129

# *Group 6*



Nguyễn Duy Hưng  
20194436



Trần Quốc Lập  
20194443



Hoàng Thiện Tâm  
20194450

*Thanks for watching!!!*

Love you all! <3 <3 <3