

How to derive Aiyagari's NBL

Let $w := w(e)$ and $R := 1+r$

Budget constraint:

$$c_t + a_{t+1} \leq w e_t + R a_t$$

Question to ask: What is the most one could borrow sustainably (i.e., in a P.V. budget feasibility sense) ?

Let this maximal borrowing at date t be denoted by \underline{a}_t .

Intuitively, if I hit the "shortest" asset position \underline{a}_t now, the best I need to do is not to consume any of my current and future endowments:

$$\boxed{c_t = 0 \text{ for all } t \geq 0}$$

The P.V. budget constraint is then:

$$0 + \underline{a}_{t+1} = we_t + R \underline{a}_t$$

$$\begin{aligned}\Rightarrow \underline{a}_t &= \frac{1}{R} \left[\underline{a}_{t+1} - we_t \right] \\ &= \frac{1}{R} \left[\frac{1}{R} \left(\underline{a}_{t+2} - we_{t+1} \right) - we_t \right] \\ &= \frac{1}{R^2} \left[\frac{1}{R} \left(\underline{a}_{t+3} - we_{t+2} \right) - we_{t+1} \right] - \frac{we_t}{R} \\ &\vdots\end{aligned}$$

$$\Rightarrow \boxed{\underline{a}_t = -\frac{1}{R} \sum_{j=0}^{\infty} \frac{we_{t+j}}{R^j}}$$

NBL \equiv Natural borrowing limit

$$\text{If } c_t = 0 \text{ for } t \geq 0 \Leftrightarrow a_t = \underline{a}_t$$

$$\text{Then } c_t > 0 \text{ for } t \geq 0 \Leftrightarrow a_t > \underline{a}_t$$

So, for $c_t \geq 0$:

$$a_t \geq \underline{a}_t \equiv -\frac{1}{1+r} \sum_{j=0}^{\infty} \frac{w(e) e_{t+j}}{(1+r)^j}$$

as in the slides, p.10.

A Karush-Kuhn-Tucker problem:

Given r :

$$v(z, e) = \max_{\hat{a}'} \left\{ \begin{array}{l} u(z - \hat{a}') + \beta \mathbb{E}\{v(z', e') | e\} \\ \text{s.t.} \\ (1) \ z' = we' + (1+r)\hat{a}' - r\phi \\ (2) \ \hat{a}' \leq z \\ (3) \ 0 \leq \hat{a}' \end{array} \right\}$$

$\cdot \mathbb{I}(z, s)$
 def: (on p. 15)

We can write this as a KKT problem:

$$v(z, e) = \max_{\hat{a}'} \left\{ \begin{array}{l} u(z - \hat{a}') \\ + \beta \mathbb{E}\{v[we' + (1+r)\hat{a}' - r\phi, e'] | e\} \\ - \bar{\lambda} [\hat{a}' - z] \\ - \underline{\lambda} [0 - \hat{a}'] \end{array} \right\}$$

Envelope BS:

$$v_z(z, e) = u_c(z - \hat{a}') \equiv u_c(c)$$

$$\Rightarrow v_z(z', e') = u_c(z' - \hat{a}'') \equiv u_c(c') \quad (4)$$

FoCs

$$\begin{aligned} -u_c(c) + \beta E \left\{ v_z(z', e') \cdot (1+r) \mid e \right\} \\ - \bar{\lambda} + \underline{\lambda} = 0 \quad (5) \end{aligned}$$

KKT Complementary slackness:

$$\bar{\lambda} \geq 0, \quad \bar{\lambda}(\hat{a}' - z) = 0 \quad (6)$$

$$\underline{\lambda} \geq 0, \quad \underline{\lambda} \hat{a}' = 0 \quad (7)$$

Assume $u(\cdot)$ s.t. we always have $\hat{a}' < z$

$\Leftrightarrow \bar{\lambda} = 0$. [i.e., one never optimal choose $c=0$!]

So we focus on $\underline{\lambda} \geq 0$.

- $\underline{\lambda} > 0 \Leftrightarrow \hat{a}' = 0$ (ad-hoc BL binding)

$$\lambda = 0 \Leftrightarrow \hat{a}' > 0 \quad (\text{not})$$

(4), (5) and (7) together imply:

$$u_c(c) = \beta(1+r) \mathbb{E}\{u_c(c') | e\} \quad \text{if} \left(\begin{array}{c} \hat{a}' > 0 \\ \Updownarrow \\ \lambda = 0 \end{array} \right)$$

$$u_c(c) > \beta(1+r) \mathbb{E}\{u_c(c') | e\} \quad \text{if} \left(\begin{array}{c} \lambda > 0 \\ \Updownarrow \\ \hat{a}' = 0 \end{array} \right)$$

As seen on slide 19