The WACC Language Specification

Second Year Computing Laboratory
Department of Computing
Imperial College London

1 What is WACC?

WACC (pronounced "whack") is a simple variant on the While family of languages encountered in many program reasoning/verification courses (in particular in the Models of Computation course taught to our 2nd year undergraduates). It features all of the common language constructs you would expect of a While-like language, such as program variables, simple expressions, conditional branching, looping and no-ops. It also features a rich set of extra constructs, such as simple types, functions, arrays and basic tuple creation on the heap.

The WACC language is intended to help unify the material taught in our more theoretical courses (such as Models of Computation) with the material taught in our more practical courses (such as Compilers). The core of the language should be simple enough to reason about and the extensions should pose some interesting challenges and design choices for anyone implementing it.

2 WACC Language Syntax

We give the syntax of the WACC language in Backus-Naur Form (BNF) extended with some basic regular expression notation that simplifies the presentation:

- $(x\text{-}y)$ stands for 'range', meaning any value from x to y inclusive;

- $(x)?$ stands for 'optional', meaning that x can occur zero or one times;

- $(x)+$ stands for 'repeatable', meaning that x can occur one or more times;

- $(x)^*$ stands for 'optional and repeatable', meaning that x can occur zero or more times.

2.1 BNF

$\langle program \rangle$	::=	'begin' $\langle func \rangle^*$ $\langle stat \rangle$ 'end'
$\langle func \rangle$	::=	$\langle type \rangle$ $\langle ident \rangle$ '(' $\langle param\text{-}list \rangle?$ ')' 'is' $\langle stat \rangle$ 'end'
$\langle param\text{-}list \rangle$	::=	$\langle param \rangle$ (',' $\langle param \rangle$)*
$\langle param \rangle$	::=	$\langle type \rangle$ $\langle ident \rangle$
$\langle stat \rangle$	::=	'skip'
	|	$\langle type \rangle$ $\langle ident \rangle$ '=' $\langle assign\text{-}rhs \rangle$
	|	$\langle assign\text{-}lhs \rangle$ '=' $\langle assign\text{-}rhs \rangle$
	|	'read' $\langle assign\text{-}lhs \rangle$
	|	'free' $\langle expr \rangle$
	|	'return' $\langle expr \rangle$
	|	'exit' $\langle expr \rangle$
	|	'print' $\langle expr \rangle$
	|	'println' $\langle expr \rangle$

$$| \quad \text{`if'} \langle expr \rangle \text{`then'} \langle stat \rangle \text{`else'} \langle stat \rangle \text{`fi'}$$
$$| \quad \text{`while'} \langle expr \rangle \text{`do'} \langle stat \rangle \text{`done'}$$
$$| \quad \text{`begin'} \langle stat \rangle \text{`end'}$$
$$| \quad \langle stat \rangle \text{`;'} \langle stat \rangle$$

$\langle assign\text{-}lhs \rangle$	::=	$\langle ident \rangle$
	|	$\langle array\text{-}elem \rangle$
	|	$\langle pair\text{-}elem \rangle$

$\langle assign\text{-}rhs \rangle$	::=	$\langle expr \rangle$
	|	$\langle array\text{-}liter \rangle$
	|	`newpair` ‘(’ $\langle expr \rangle$ ‘,’ $\langle expr \rangle$ ‘)’
	|	$\langle pair\text{-}elem \rangle$
	|	`call` $\langle ident \rangle$ ‘(’ $\langle arg\text{-}list \rangle$? ‘)’

$\langle arg\text{-}list \rangle$	::=	$\langle expr \rangle$ (‘,’ $\langle expr \rangle$)*

$\langle pair\text{-}elem \rangle$	::=	`fst` $\langle expr \rangle$
	|	`snd` $\langle expr \rangle$

$\langle type \rangle$	::=	$\langle base\text{-}type \rangle$
	|	$\langle array\text{-}type \rangle$
	|	$\langle pair\text{-}type \rangle$

$\langle base\text{-}type \rangle$	::=	`int`
	|	`bool`
	|	`char`
	|	`string`

$\langle array\text{-}type \rangle$	::=	$\langle type \rangle$ ‘[’ ‘]’

$\langle pair\text{-}type \rangle$	::=	`pair` ‘(’ $\langle pair\text{-}elem\text{-}type \rangle$ ‘,’ $\langle pair\text{-}elem\text{-}type \rangle$ ‘)’

$\langle pair\text{-}elem\text{-}type \rangle$	::=	$\langle base\text{-}type \rangle$
	|	$\langle array\text{-}type \rangle$
	|	`pair`

$\langle expr \rangle$	::=	$\langle int\text{-}liter \rangle$
	|	$\langle bool\text{-}liter \rangle$
	|	$\langle char\text{-}liter \rangle$
	|	$\langle str\text{-}liter \rangle$
	|	$\langle pair\text{-}liter \rangle$
	|	$\langle ident \rangle$
	|	$\langle array\text{-}elem \rangle$
	|	$\langle unary\text{-}oper \rangle \langle expr \rangle$
	|	$\langle expr \rangle \langle binary\text{-}oper \rangle \langle expr \rangle$
	|	‘(’ $\langle expr \rangle$ ‘)’

$\langle unary\text{-}oper \rangle$	::=	‘!’ | ‘-’ | `len` | `ord` | `chr`

$\langle binary\text{-}oper \rangle$	::=	‘*’ | ‘/’ | ‘%’ | ‘+’ | ‘-’ | ‘>’ | ‘>=’ | ‘<’ | ‘<=’ | ‘==’ | ‘!=’ | ‘&&’ | ‘||’

$\langle ident \rangle$	::=	(‘_’ | ‘a’-‘z’ | ‘A’-‘Z’) (‘_’ | ‘a’-‘z’ | ‘A’-‘Z’ | ‘0’-‘9’)*

$\langle array\text{-}elem \rangle$	::=	$\langle ident \rangle$ (‘[’ $\langle expr \rangle$ ‘]’)+

$\langle \textit{int-liter} \rangle$	::=	$\langle \textit{int-sign} \rangle$? $\langle \textit{digit} \rangle +$
$\langle \textit{digit} \rangle$	::=	('0'-'9')
$\langle \textit{int-sign} \rangle$	::=	'+' | '-'
$\langle \textit{bool-liter} \rangle$	::=	'true' | 'false'
$\langle \textit{char-liter} \rangle$	::=	''' $\langle \textit{character} \rangle$ '''
$\langle \textit{str-liter} \rangle$	::=	'"' $\langle \textit{character} \rangle$* '"'
$\langle \textit{character} \rangle$	::=	any-ASCII-character-except-'\'-'''-'"'
	|	'\' $\langle \textit{escaped-char} \rangle$
$\langle \textit{escaped-char} \rangle$	::=	'0' | 'b' | 't' | 'n' | 'f' | 'r' | '"' | ''' | '\'
$\langle \textit{array-liter} \rangle$	::=	'[' ($\langle \textit{expr} \rangle$ (',' $\langle \textit{expr} \rangle$)*)? ']'
$\langle \textit{pair-liter} \rangle$	::=	'null'
$\langle \textit{comment} \rangle$	::=	'#' (any-character-except-EOL)* $\langle \textit{EOL} \rangle$

NB: There is an additional constraint on the syntax of function definitions, that every execution path through the body of the function must end with either a 'return' statement or an 'exit' statement.

3 WACC Language Semantics

We now go through each of the language components and explain their behaviour and purpose in more detail.

3.1 Types

With the exception of nested pairs (see below), the WACC language is both statically and strongly typed.

- Static, in that, once declared, the type of a variable is fixed for the duration of the program.
- Strong, in that the compiler should not coerce between types.

There is also no explicit typecasting.

Basic Types: The basic types in the WACC language are:

- int: The Integer type. The WACC language supports integers in the inclusive range $[-2^{31}, 2^{31}-1]$.
- bool: The Boolean type (true or false).
- char: The Character type. The WACC language supports only the ASCII characters.
- string: The String type. The WACC language treats strings as a list of characters.

We write 'T' to denote an arbitrary type.

Arrays: As well as the basic types given above, the WACC language also supports the array type. We write 'T[]' to denote an array whose elements are of type 'T'. Note that 'T' can be of any type, including another array type, which allows for nested arrays. Arrays are allocated on the heap. As well as their elements, each array also tracks its length, which is set when it is created.

Strings vs. Arrays: In the WACC language, arrays of characters may be treated as strings for the purposes of type-checking, but strings may not be treated as arrays. This is a one-way typing relaxation.

Pairs: Pairs are allocated on the heap and contain two elements that can be of any type. We write 'pair(T_1, T_2)' to donate a pair whose first element is of type 'T_1' and second element is of type 'T_2' (these need not be the same). Note that if either 'T_1' or 'T_2' is a pair type, we do not write the type of the sub-elements. For example, a pair whose first element is an integer and whose second element is a pair of characters is written as 'pair(int, pair)' and not as 'pair(int, pair(char, char))'. It is obvious that we lose some typing information in this way. Moreover, due to the loss of type information for nested pairs, it is possible to subtly coerce between types.

3.2 Program Scopes

The WACC language includes explicit scoping. Various statements introduce new program scopes, which have an effect on the visibility of program variables.

Whenever a new variable is declared it is added to the current program scope. When a program scope is exited, every variable created within that scope is destroyed. This means that variables are not accessible by statements outside the scope of their creation, although they are accessible in child scopes.

The main, or global scope is created at the start of a WACC program and is exited at the end of the program. Functions can only be created at the beginning of this global scope, but they may be called from within any child scope.

We will see that several other program constructs, including functions, while loops and conditional branches, introduce new program scopes during their execution.

3.3 Programs

A WACC program ⟨*program*⟩ consists of zero or more function definitions followed by the body of the main function. The whole program is written between the 'begin' and 'end' tokens, denoting the main or global program scope. A WACC file (extension .wacc) only ever contains a single WACC program.

3.4 Function Definitions

A function definition ⟨*func*⟩ consists of a return type, a function name and zero or more typed parameters followed by the function's body. A function's body, which is denoted by the 'is' and 'end' tokens, is executed in its own scope containing only the parameters passed into the function. As mentioned above, any execution path through the body of the function must end with either a 'return' statement, whose expression type must match the function's return type, or an 'exit' statement.

Functions can only be defined at the beginning of the global scope, before the body of the main function. Functions may, however, be both recursive and mutually recursive.

3.5 Statements

A statement ⟨*stat*⟩ consists of: a no-op, a variable definition, an assignment, an input read, a memory free, a function return, an exit call, a print command, a conditional branch, a while loop, a scope introduction or the sequential composition of two statements.

We discuss each of these in more detail below.

No-op Statements: A no-op statement 'skip' does not do anything. It is used where a statement is expected but we do not want to do anything. For example, in an if statement where we want to have an empty else clause.

Variable Declaration Statements: A variable declaration statement creates a new program variable in the current scope setting its static type and initial value. The statement must be given a valid WACC type ⟨*type*⟩, a variable name ⟨*ident*⟩ and an initial assignment value ⟨*assign-rhs*⟩.

Variable names must not clash with another variable name which resides in the same scope, function names or string literals such as 'begin', 'end', 'skip', or 'if' etc.

They can consist of a letter or underscore followed by any number of alphanumeric characters or underscores or none.

The initial assignment to a variable follows all of the assignment restrictions discussed in detail in the assignment statement section below.

A variable must be declared before it is accessed.

Any attempt to access an undeclared variable results in a semantic error: `"Variable ID is not defined in this scope"`.

Additionally, every use of a variable must match the type assigned to that variable when it was declared.

A variable can only be accessed within the scope of its declaration (or any child scope) and it is destroyed when exiting this scope. Variables must be unique within their scope, so a variable cannot be redefined within the same scope. Variables can, however, be redefined within a child scope. In this case the redefined variable will have the new value in the child scope and any of its child scopes. Redefined variables are permitted to have a different type than the instances outside of its scope. For example:

```
1  begin
2    int x = 1 ;
3    begin
4      bool x = true ;
5    end ;
6  end
```

compiles successfully.

Once the child scope is exited all variables defined in the child scope are destroyed and any redefined variables are returned to their original value before redefinition. However, any reassignment inside the child scope which occurs **before** the variable redefinition still persists afterwards. For example:

```
1  begin
2    int x = 1 ;
3    begin
4      x = 2 ;
5      int x = 5 ;
6      println x
7    end ;
8    println x
9  end
```

The first **println** (line 6) outputs **5** whereas the the second **println** (line 8) outputs **2**. Notice here that x is reassigned to be **2** inside the child scope (line 4) and still persists outside the child scope because the reassignment occurs before the redefinition. However, the `int` variable x storing the value **5** is destroyed upon exiting the child scope and can no longer be accessed. Conversely:

```
1  begin
2    int x = 1 ;
3    begin
4      int x = 5 ;
5      x = 2 ;
6      println x
7    end ;
8    println x
9  end
```

In this instance, lines **4** and **5** have been swapped so that the variable redefinition occurs before the reassignment. The outputs upon executing this program are **2** and **1** respectively. The reassignment in this case is associated with the variable **x** declared inside the same scope which is eventually exited and destroyed, whereas the previous reassignment was associated with the variable **x** in the parent scope.

> **Gap 1.** *Complete the above paragraph* (**5 marks**)

Assignment Statements: An assignment statement updates its target (the left-hand side of the '='') with a new value (the right-hand side of the '='). The target of an assignment can be either a program variable, an array element or a pair element. The assignment value can be one of five possible types: an expression, an array literal, a function call, a pair constructor or a pair element.

- If the assignment value is an expression ⟨*expr*⟩ then the expression's type must match that of the target. The expression is then evaluated and the resulting value is copied into the target.

- If the assignment value is an array literal ⟨*array-liter*⟩ then the value array's element type must match that of the target array, but their lengths can be different. The value array is then allocated on the heap with each element initialised to the given value. After that, the reference of the value array is copied to the reference of the target array. For more details on array literals, see the expressions section.

- If the assignment value is a function call 'call' then the function's return value type must match that of the target. The number and types of the function's arguments must also match the function's definition. A function is called by its name and its arguments are passed by value (for int, bool and char types) or by reference (for strings, arrays and pairs). When called, the function's body is executed in a new scope, not related to the current scope. The only declared variables are the function's parameters, whose types are set by the function definition and whose values are set by the function call's arguments. When the execution of the function body terminates, the function's return value is then copied into the assignment target.

- If the assignment value is pair constructor 'newpair' then the target must be of type 'pair(T$_1$, T$_2$)'. The pair constructor is passed two expressions that must match T$_1$ and T$_2$ respectively. A 'newpair' assignment allocates enough memory on the heap to store the pair structure and its elements. It then initialises each element of the pair using the evaluation of the first expression for the first element and the evaluation of the second expression for the second element. Pairs, in the WACC language, are always used by reference, so a reference to the pair is copied into the target, rather than the actual content of the pair.

- If the assignment value is a pair element ⟨*pair-elem*⟩ then the expression passed to the pair element must be of type 'pair', must not be a 'null' pair literal and must match the type of the first (or second) element of the target pair when using the 'fst' (or 'snd') keyword. The pair expression is evaluated to obtain a reference to a pair and this is dereferenced to find the corresponding pair element, which is then copied into the target. Any attempt to dereference a 'null' reference should generate a runtime error.

Read Statements: A read statement 'read' is a special assignment statement that takes its value from the standard input and writes it to its argument. Just like a general assignment statement, a read statement can target a program variable, an array element or a pair element. However, the read statement can only handle character or integer input.

The read statement determines how it will interpret the value from the standard input based on the type of the target. For example, if the target is of type 'int' then it will convert the input string into an integer (truncated to fit within the minimum/maximum range of the 'int' type). If the standard input value is incompatible with the type of the target, then this value will not be consumed from the standard input stream. Instead, the program will continue, leaving the target's value unchanged.

Memory Free Statements: A memory free statement 'free' is used to free the heap memory allocated for a pair or array and its immediate content. The statement is given an expression that must be of type 'pair(T$_1$, T$_2$)' or 'T[]' (for some T, T$_1$, T$_2$). The expression must evaluate to a valid reference to a pair or array, otherwise a segmentation fault will occur at runtime.

If the reference is valid, then the memory for each element of the pair/array is freed, so long as the element is not a reference to another pair or another array (i.e. free is not recursive). Then the memory that stores the pair/array itself is also freed.

Function Return Statements: A return statement can only be present in the body of a non-main function and is used to return a value from that function. The type of the expression given to the return statement must match the return type of the function. Once the return statement is executed, the function is immediately exited.

Expression Type	Behaviour	Example Expression	Example Output
int	Output the integer converted to a decimal string.	10	"10"
bool	Output "true" if the boolean is true and "false" otherwise.	false	"false"
char	Output a single-character string.	'c'	"c"
string or char[]	Output the string, otherwise convert the char array to a string and then output	['a','b','c']	"abc"
Other Array Types	Outputs the address of the array as a string	[0, 1, 2, 3, 4, 5]	0x22150
pair	Outputs the base address of the pair on the heap	newpair(a, b)[1]	0x22150

Table 1: The behaviour of the print statements for each type of expression.

Exit Statements: The exit statement takes the form : 'exit' $\langle expr \rangle$, and is used to terminate the program with exit code e where e is the evaluated expression of $\langle expr \rangle$. $\langle expr \rangle$ can either be an integer or an expression that evaluates to an integer e. When the exit statement is executed, it displays 'The exit code is e' then immediately terminates. If no exit statement exists in the program, the default exit code is 0.

> **Gap 2.** *define exit statements* **(2 marks)**

Print Statements: There are two types of print command in the WACC language. The 'print' command takes an expression and prints the result of its evaluation to the standard output. The 'println' command is similar, but additionally prints out a new line afterwards.

The output representation of each expression evaluation depends on the type of the expression. The behaviour of the print statements for each type of expression is shown in Table 1, along with some example cases.

> **Gap 3.** *Fill in Table* 1 **(3 marks)**

Conditional Branch Statements: A conditional branch statement 'if' evaluates an expression and determines which program path to follow. The statement is given a condition expression, that must be of type 'bool', and two body statements, one for the 'then' branch and one for the 'else' branch.

If the condition evaluates to 'true', then the 'then' body statement is executed. Otherwise, the 'else' body statement is executed. Each of the program branches is executed in its own scope, which are denoted by the 'then' and 'else' tokens and the 'else' and 'fi' tokens, respectively.

While Loop Statements: A while loop statement takes the form 'while' $\langle expr \rangle$ 'do' $\langle stat \rangle$ 'done'. A while loop repeatedly evaluates a body statement, while its condition expression evaluates to true. The statement is given a condition expression $\langle expr \rangle$, that must be of type bool once evaluated, and a body statement $\langle stat \rangle$.

The condition is always checked before executing the body statement. If the condition evaluates to true, then the body statement is executed. This is repeated until the condition evaluates to false. If the condition initially evaluates to false, the body statement will never be executed and the program will jump to the next statement below the body. The body statement is executed in its own scope, which

---

[1]This is not exactly an expression because it can only appear on the right hand side of an assignment. However, it gives the best example here.

is denoted by the 'do' and 'done' token. Since the body has its own scope, any new variables declared within the body are destroyed upon exiting the while loop. A while loop will run indefinitely if the variant does not make any progress, slowly towards a `false` condition expression. For example :

```
 1  begin
 2    int x = 3 ;
 3    int y = 7 ;
 4    print "initial value of x: " ;
 5    println x ;
 6    while y > 0 do
 7      x = x + 1 ;
 8      y = y - 1
 9    done ;
10    print "final value of x: " ;
11    println x
12  end
```

In the program above, the variable y is decremented in every iteration of the loop which causes it to eventually be less than 0 evaluating the while loop condition expression to `false` and finally exiting the loop and eventually terminating the program.

Gap 4. *Define/describe while loop statements*   (**6 marks**)

Scoping Statements: A scoping statement introduces a new program scope, which is denoted by the 'begin' and 'end' tokens.

Sequential Composition: A sequential composition is a statement $\langle stat \rangle$ followed by the symbol ; and another statement $\langle stat \rangle$. A sequential composition s1 ; s2 first executes s1, then s2. If we start with initial program state **h**, then executing s1 will give us an intermediate state **h1**. Then with program state **h1** we execute s2 giving our final state **h2**.

Gap 5. *Define/describe sequential composition*
i.e. `<stat> ; <stat>`   (**2 marks**)

3.6   Expressions

An expression $\langle expr \rangle$ consists of a literal (integer, boolean, character, string or pair), a variable, an array element, a unary expression, a binary expression or an expression enclosed by parenthesis.

We discuss the meaning of each of these expressions in more detail below.

The expressions of the WACC language have been chosen to be side-effect free. A side-effect free expression does not modify the state of the program. This means that, if the program state is not changed by other code, evaluating an expression will always give the same result.

Gap 6. *Define side-effect free expressions*   (**1 mark**)

Integer Literals: An integer literal $\langle int\text{-}liter \rangle$ consists of a sequence of decimal digits. Optionally, the sequence can be preceded by a '+' or a '-' symbol.

Boolean Literals: A boolean literal $\langle bool\text{-}liter \rangle$ is either 'true' or 'false'.

Character Literals: A character literal $\langle char\text{-}liter \rangle$ is a single ASCII character between two ''' symbols. A '\' can be used to escape the character that immediately follows the '\'. The meaning of each escaped character is shown in Table 2.

Gap 7. *Fill in Table* 2   (**2 marks**)

Representation	ASCII Value	Description	Symbol
\0	0x00	null terminator	NUL
\b	0x08	backspace	BS
\t	0x09	tab	TAB
\n	0x0a	new line	NL
\f	0x0c	new page	FF
\r	0x0d	carriage return	CR
\"	0x22	double quote	"
\'	0x27	single quote	'
\\	0x5c	backslash	\

Table 2: The escaped-characters available in the WACC language.

Operator	Argument Type	Return Type	Meaning
'!'	bool	bool	Logical Not
'-'	int	int	Negation
'len'	T[]	int	Array Length
'ord'	char	int	ASCII Value Of Char
'chr'	int	char	ASCII Char From Value

Table 3: The unary operators of the WACC language with their types and meanings.

String Literals: A string literal ⟨*str-liter*⟩ is a sequence of characters between two '"' symbols. Each character in the string literal can be escaped in the same way as in character literal.

Pair Literals: The only pair literal ⟨*pair-liter*⟩ is 'null' which represents a reference that does not point to any pair. The 'null' pair literal can match the type of any pair. To see how pairs are created, read the 'newpair' case of the assignment statement.

Array Literals: Array literals cannot occur directly in expressions, but they do occur in the WACC language as assignment values. An array literal starts with a '[' token and ends with a ']' token. The elements of the array (zero or more) are given between these brackets and are separated by ',' tokens. All elements of an array must be of exactly the same type, so the type of any non-empty array literal can be statically determined. If, however, an array literal is empty, we allow it to be of any array type. For example, the array '[]' can be of type 'int[]','bool[]', 'char[]', etc... depending on the context, but the array '[1]' must be of type 'int[]'.

Variables: When a variable expression ⟨*ident*⟩ is evaluated it returns the value of that variable. If the variable is of type 'T' then the return type of the expression is also 'T'.

Array Elements: An array element expression evaluates to return an element from an array. The expression consists of two sub-expressions, the first of which must be of type 'T[]' and the second of which must be of type 'int'. The return type of the overall expression is 'T'.

The first expression is evaluated to find an array 'a' and the second is evaluated to find an index 'i'. The overall expression returns the element at the index 'i' of array 'a', that is, 'a[i]'.

If the array has length l then the index 'i' must be between 0 and $(l-1)$, otherwise the expression will generate a runtime error.

Unary Operators: A unary operator ⟨*unary-oper*⟩ has a single sub-expression. The unary operators available in the WACC language are shown in Table 3. All unary operators have the same precedence, they are evaluated from right to left.

> **Gap 8.** *Fill in Table 3* (**2 marks**)

Opera- tor	Prece- dence	Argument 1 Type	Argument 2 Type	Return Type	Meaning
'*'	1	int	int	int	Multiply
'/'	1	int	int	int	Divide
'%'	1	int	int	int	Modulus
'+'	2	int	int	int	Plus
'-'	2	int	int	int	Minus
'>'	3	int/char	int/char	bool	Greater Than
'>='	3	int/char	int/char	bool	Greater Than or Equal
'<'	3	int/char	int/char	bool	Less Than
'<='	3	int/char	int/char	bool	Less Than or Equal
'=='	4	int/char/bool	int/char/bool	bool	Value Equality
'=='	4	string/T[]/pair	string/T[]/pair	bool	Referential Equality
'!='	4	int/char/bool	int/char/bool	bool	Value Inequality
'!='	4	string/T[]/pair	string/T[]/pair	bool	Referential Inequality
'&&'	5	bool	bool	bool	Logical And
'||'	6	bool	bool	bool	Logical Or

Table 4: The binary operators of the WACC language, with their types and meanings.

- The '!' operator performs a logical Not operation on the result of evaluating its sub-expression, returning 'true' if the sub-expression evaluates to 'false' and vice-versa.

- The '-' operator inverts the sign of the evaluation of its sub-expression.

- The 'len' operator returns the length of the array referenced by the evaluation of its sub-expression.

- The 'ord' operator returns the ASCII code for the evaluation of its sub-expression.

> **Gap 9.** *Define/describe the 'ord' operator* (**1 mark**)

- The 'chr' operator returns the character representation of the ASCII value of it's sub-expression.

> **Gap 10.** *Define/describe the 'chr' operator* (**1 mark**)

Binary Operators: A binary operator is used in in-fix style between two sub-expressions. The binary operators available in the WACC language are shown in Table 4. The operators have different precedences, as illustrated in the table, with 1 being the highest and 6 being the lowest.

> **Gap 11.** *Fill in Table* 4 (**2 marks**)

- The '*', '/', '%', '+' and '-' operators all have their standard mathematical behaviour, where integer underflow/overflow results in a runtime error. If the divisor of a division ('/') or modulus ('%') operator is evaluated to '0', then this also results in a runtime error. The result of a division operation is positive if both its dividend and divisor have the same sign, and negative otherwise. The result of a modulus operation has the same sign as its dividend.

- The '>', '>=', '<' and '<=' operators perform a comparison test on the evaluations of their sub expressions. They accept expressions of type 'int' or 'char', but both expressions must have exactly the same type. The result is 'true' if the comparison of the evaluated expressions is true. Otherwise, the result it 'false'.

- The '==' operator performs an equality test on the evaluations of its sub-expressions. It accepts any two expressions of exactly the same type. When applied to expressions of type 'int', 'bool' or 'char', the result is 'true' iff the values of the two arguments are the same. When applied to expressions of type 'string', 'T[]' or 'pair', the result is 'true' iff the two references point to the same object and they have exactly the same type. Otherwise, the result is 'false'.

- The '`!=`' operator returns the opposite result to the '`==`' operator.

- The '`&&`' operator performs a logical And operation on the result of evaluating its sub-expressions, returning '`true`' if both sub-expressions evaluate to '`true`' and '`false`' otherwise.

- The '`||`' operator performs a logical Or operation on the result of evaluating its sub-expressions, returning '`true`' if either sub-expression evaluates to '`true`' and '`false`' otherwise.

Parenthesis: We can introduce a pair of parenthesis around an expression to control its evaluation. The expression in a parenthesis is always evaluated first, regardless of the operator precedence.

3.7 Whitespace and Comments

Whitespace is used in the WACC language to delimit keywords and variables. For example, '`if a == 13`' denotes the start of an '`if`' statement with boolean condition '`a == 13`', whereas '`ifa == 13`' denotes a boolean expression comparing the variable '`ifa`' with the value '`13`'. Any other type of occurrence of whitespace is ignored by the compiler. Note, in particular, that the code indentation in the example programs has no meaning, it simply aids readability. Also note that whitespace inside a string or character literal is preserved by the compiler.

Comments are used in the WACC langauge to give the reader a helping hand when trying to make sense of the code. The syntax of commenting is the symbol `#` followed by any number of character except EOL and then ending with an EOL (starting a new line implies an EOL). For example, `#Hello This is a Comment!` The `#` symbol specifies the start of the comment and anything else after it will be treated as a comment regardless of whitespaces. To start writing code which the compiler will consider as input again requires starting on the next line.

Gap 12. *Define/describe comments* (**3 marks**)