

ARCHITECTURE

Now that we already know how to match (by comparing and calculating *similarity score* from section 4), we then proceed to a bigger picture. This section describes how the architecture of overall system is.

1.1 WEB SERVICE

From figure 3, web service clients use the system by directly pass inputs to the *controller*. Recalling from previous sessions, all possible inputs to the system are as follow.

1. Base names, separated by new lines¹.
2. To match names, separated by new lines.
3. Matching algorithms, specify needed algorithm names along with its weight (section 1.3).
4. Threshold (section 1.5).

Currently a preferable format is JSON.

```
{
  "base_name":"Smith",
  "to_match_names":"Gowan\r\nSpeed",
  "matching_algorithms":{
    "1":{"name":"LookupTable", "weight":"10"},
    "2":{"name":"LevenshteinDistance", "weight":"1"},
    "3":{"name":"Soundex", "weight":"3"},
    "4":{"name":"IrishSoundex", "weight":"6"}
  },
  "threshold":"0"
}
```

Listing 1: JSON.

¹ Preferably \n or \r, see <http://stackoverflow.com/questions/1761051/difference-between-n-and-r>.

1.2 INITIAL IDEA

Let us start by the basic idea of this project.

As mentioned in section 2.3, the objective of this project is to provide a web service that produces matching result between two *lists* of names. As we see the word *list* here, that means our inputs are not only a pair of names, but rather two lists. In real world use, this list can be large, a hundred or thousand, depending on the client who uses the system.

We will introduce two terms, *base name* and *to-match name*. *Base name* acts as a base and will be matched against each *to-match name* in their list, from start until the end, then proceed to the next *base name*, match against the whole *to-match name* list again, and so on.



Figure 1: Base name 'SMITH' comparing to to-match name 'SMEETH'. Scores of each matching algorithms are presented in the bubble above the arrow.

Figure 1 shows a snapshot of an attempt to match between *base name* 'SMITH' against *to-match name* 'SMEETH'. Similarity score for each matching algorithms have been calculated. And by these scores, we can calculate *overall score* for 'SMITH' and 'SMEETH'.

So the next step is to match current *base name*, 'SMITH', against next *to-match name*, 'SMILEY'.

Once *base name* 'SMITH' completes all *to-match name*'s in their list, the system then process to the next *base name*, 'SOMERS', and start over the matching process against the whole *to-match name* list again, from start to the end.

1.3 WEIGHTING MATCHING ALGORITHMS

We realised that, for matching names, each matching algorithms should not be treated as all the same priority. For example, for Irish names, it would be better if we favour *Irish soundex* over the traditional *Soundex*, because it produces more accurate result.

By this idea we also implement *weight* for each matching algorithm. We will suggest initial values, but also allow client to change these values. Table 1 states these suggested initial weights.

MATCHING ALGORITHM	WEIGHT
Levenshtein distance	1
Soundex	3
Irish soundex	6
Lookup table	10

Table 1: Matching algorithm weights.

By summarising products of each matching algorithm *similarity score* and its weight, dividing by sum of all weight, we can obtain *overall weighted score* (OWS). This sentence can be represented by equation 1.

$$\text{OWS} = \frac{\sum_{i=1}^n (s_i \times w_i)}{\sum_{i=1}^n w_i} \quad (1)$$

Where s and w are *similarity score* and weight of matching algorithm i respectively. n is number of available matching algorithms.

This *overall weighted score* will represent each matching and all results will be sorted by this score. Usage and calculation of this weighting will be described in more detail in the next section (1.4).

1.4 ACTUAL SYSTEM

Following our basic idea from previous sections, we then design the architecture of our system.

Suppose we have two inputs, list of *base names* of length b , and list of *to-match names* of length t . We need to process the matching for $b \times t$ times. We call this single matching between *base name* and *to-match name* as *matching cycle*.

In this following figure 2 we once again show a snapshot of an attempt to match between *base name* 'SMITH' against *to-match name* 'SMEETH'. But now in a *matching cycle* style.

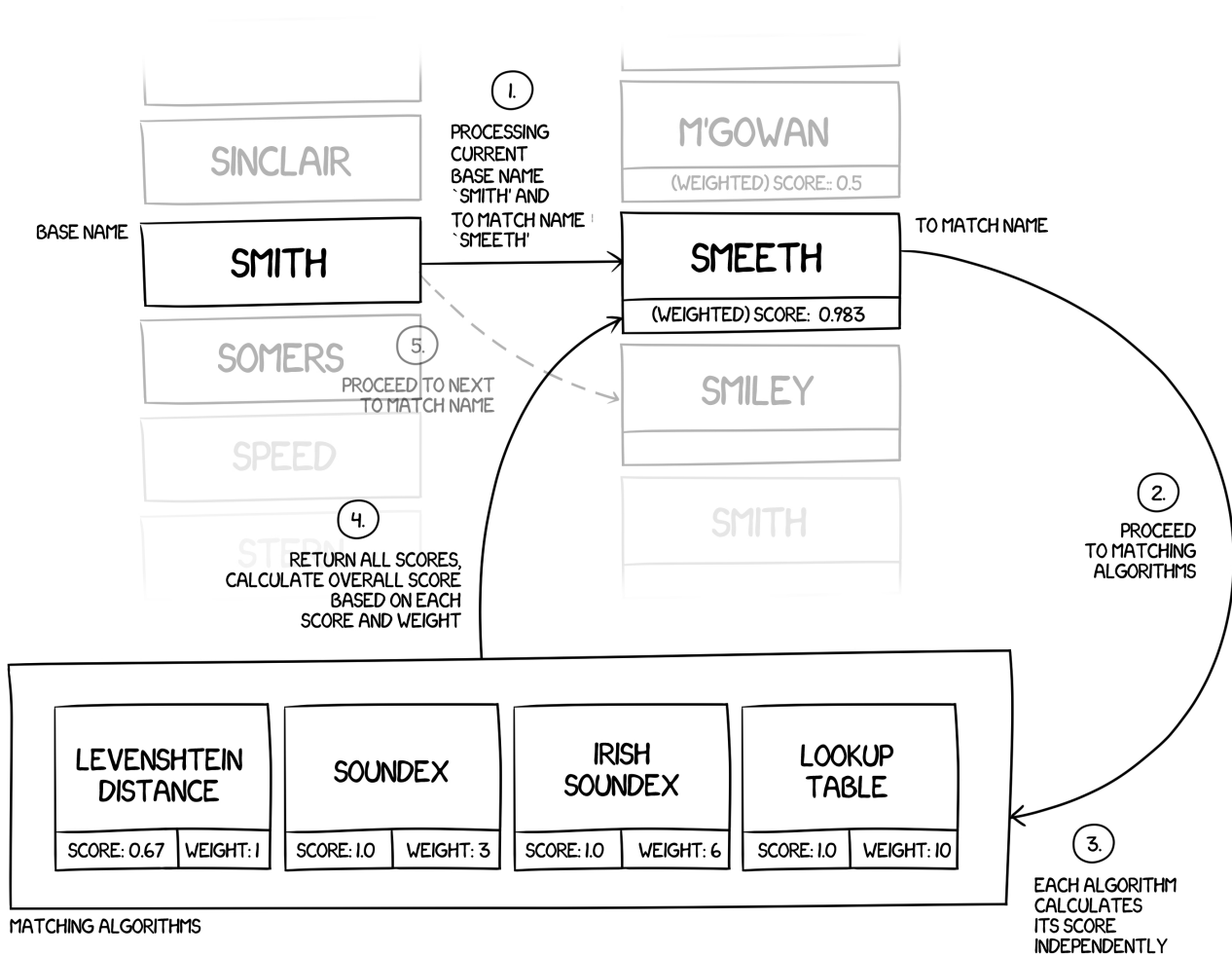


Figure 2: One matching cycle.

One matching cycle consists of 5 steps as shown in figure 2.

1. Processing current *base name* 'SMITH' and *to-match name* 'SMEETH'
2. Proceed to matching algorithms.
3. Each algorithm calculates its score independently.
 - a) Levenshtein distance between 'SMITH' and 'SMEETH' is 2 (1 substitution of I to E and 1 insertion of E). So *similarity score* is $(6 - 2) \div 6 = 0.667$ where 6 comes from the length of the longer string, SMEETH.
Weight of this algorithm is 1.
 $\therefore \text{WEIGHTED SCORE} = 0.667 \times 1 = 0.667$

- b) Soundex of both 'SMITH' and 'SMEETH' are S530 so their *similarity score* is 1.0.

Weight of this algorithm is 3.

$$\therefore \text{WEIGHTED SCORE} = 1.0 \times 3 = 3.0$$

- c) Irish soundex of both 'SMITH' and 'SMEETH' are S530 so their *similarity score* is 1.0.

Weight of this algorithm is 6.

$$\therefore \text{WEIGHTED SCORE} = 1.0 \times 6 = 6.0$$

- d) References between 'SMITH' and 'SMEETH' is found in the Lookup table via group 1897. so their *similarity score* is 1.0

Weight of this algorithm is 10.

$$\therefore \text{WEIGHTED SCORE} = 1.0 \times 10 = 10.0$$

4. Return all scores and then calculate *overall weighted score* for 'SMITH' and 'SMEETH'. Sum of the scores is $0.667 + 3.0 + 6.0 + 10.0 = 19.667$. Sum of the weights is $1 + 3 + 6 + 10 = 20$. Therefore the *overall weighted score* is $19.667 \div 20 = 0.983$.
5. Matching cycle for 'SMITH' and 'SMEETH' is finished with *overall weighted score* 0.983. Now the system will proceed to the next *to-match name* 'SMILEY'. Matching cycles for *base name* 'SMITH' will continue until the end of *to-match name* list. After that it will start matching cycles for *base name* 'SOMERS' from the start of *to-match names*, and so on.

Once all cycles are fully finished for every *base names* and *to-match names*, we will get all *overall weighted scores* ready. So we can sort and present in web interface (section 1.8), or return as a result in web service (section 1.1).

1.5 THRESHOLDING THE RESULTS

Suppose there are a thousand of *to-match names*, there could be many irreverent results that are not likely to match each *base name*. For example *overall weighted scores* of 'SMITH' and 'CROMBIE' is just 0.007.

Client may opt-out these irreverent results by specifying a floating number *threshold*. Any *to-match name* with *overall weighted scores* lower than *threshold* will be discarded from the result.

1.6 DATA FLOW

In the previous section we describe the essence of this project, how we use matching algorithms to calculate score of similarity between

two strings. We know how to process the data. Now in order to make the system becomes useable. We need to consider two more things.

- How to gather inputs from clients.
- How to present or return results to clients.

From research questions (section 2.2) we mentioned two ways to communicate with clients, by *web service* and *web interface*.

Clients who use the system as a web service will send inputs directly without any medium in between, and will receive result back in form of agreed format, e.g. JSON.

On the other hand, clients who use the system via web interface will use a form in a web interface (web page) provided by the system to provide inputs, and results will be presented in another page after client submitted the form.

SERVICE TYPE	INPUT SOURCE	RESULT FORMAT
Web service	Web/mobile/desktop application	JSON
Web interface	Provided form	Web page result

Table 2: Service types and their inputs and results.

In the next section (1.7) we will describe how we gather inputs and provide results.

1.7 MVC

Our system is based on Ruby on Rails, which is a MVC² framework. We will use Rails architecture to encapsulate our system be these following means.

VIEW: where the form for web interface is implemented. It creates a web page with inputs for use to fill in. Client can inputs names manually, or upload a file containing names. He also can choose whether to use any available matching algorithms.

Inputs from *view* are then passed to *controller*.

View is also responsible in displaying result to web interface clients, and generating JSON result for web service clients.

CONTROLLER: receives inputs from different sources, inputs from form of web interface style, or direct input from clients using web service style. Inputs will be pre-processed, such as separat-

² Model-View-Controller [23]

ing lines from file input, removing white spaces, or converting input to upper-case.

Once inputs are ready, *controller* then passes these inputs to *model*, where our matching system lies in.

After inputs are processed, *controller* receives results back from *model*, then *controller* will decide which kind of result it needs to return from input source. It will then pass results to appropriate *view*.

MODEL: this is where we implement our whole matching system in. Model constructs *base names* and *to-match names* from received inputs, then invoke matching algorithms, as described in section 1.4. Model will pass results back to *controller* after finished.

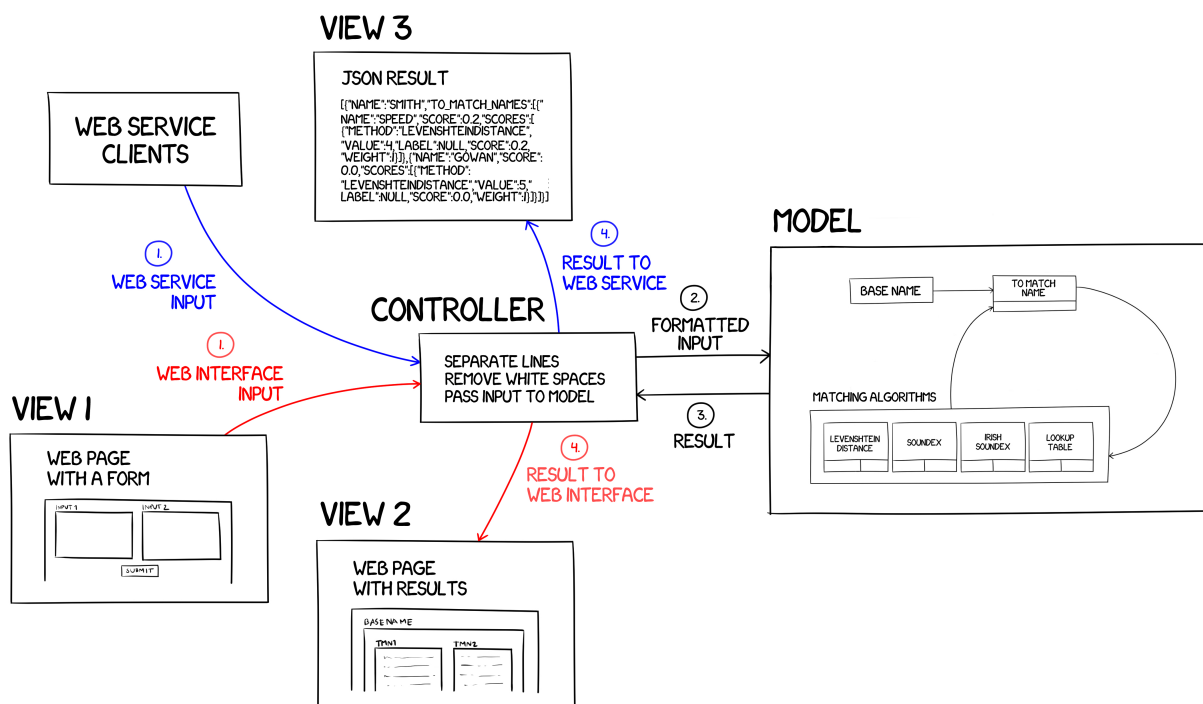


Figure 3: Data flow from web service client and web interface client.

INTRODUCTION

This project contains some backgrounds which are not part of computer science. Here introduces adequate information in order to help getting start.

2.1 MOTIVATION

The first civil registration in Ireland was performed on 1864 [1]. Before that time census materials were mostly lost or incomplete. So genealogical researches need to rely on parish records and also some ‘census substitute’ documents, such as land ownership and tenancy records¹.

However, for these documents, each of them usually does not contain enough information to indentify individuals. Some of them contains name and address, whereas others might contain only name. In order to fulfil missing information of one individual that scattered among many documents, *Record linkage* is one method to do so.

Record linkage uses a person's name as a basis to link that person's information between many documents. Together with other coherent attributes to ensure the link is correct, a more complete information about that person can be obtained.

In addition, this is not only just for one person. We can assume the relationship of the person to others that might be close to, and apply the information to those people as well. For example, if we know that there is a record that is believed to consist of people from the same area in each page [3] (but no area or address is mentioned, or some is missing in the page). And we can find one or more person's addresses in that page by using record linkage. We might be able to apply those addresses to all people in that page as well.

Apparently linking or matching person's name is important in the process. Unfortunately, in the 19th century, in Ireland, there was no standard of the spelling of names, handwriting could be difficult to read and contractions or abbreviations were often used. Many people were not literate, so they asked literate people to write their names. This way even names with the same pronunciation and for the same individual could be written in many different ways, depending on who wrote them.

*“Record linkage is used in historical research, social studies, marketing, administration and government as well as in genealogy”
– Winstanley [2]
section 2.2*

¹ [2] section 1.1

In addition to the various ways of spelling one's name, people from this time also often use Irish names which equivalent to modern names, for example, Irish version of *Smith* could be *Gowan*. There are also some Irish prefixes like *O'*, *M'*, *Mac*, etc. When combined together this would result in *O'Gowan* or *M'Gowen*, and so on.

An example list of possible equivalent Irish names of *Smith* could be as follow.

Smith, Smyth, Smythe, Smeeth, Going, Gowing, Maizurn, McGhoon, MaGough, M'Ghoon, MacGivney, MacGivena, M'Givena, MacGhoon, M'Evinie, McGivney, MacEvinie, McGivena, M'Givney, McEvinie, MacAvinue, M'Avinue, McAvinue, McCona, MaGowen, MaGowan, MaGovern, MaGowen, McGowan, McGoween, McGown, M'Cona, MeCowan, MeGowan, MacGown, MacGoween, MacGowan, MacCona, M'Gowan, M'Gowen, M'Gown, Ogowan, O'Gowan, Gowen, Gowan, Gow, Goan

At present time, when historical researchers try to trace people back using historical records, they would encounter this problem of name variations.

Various solutions have been created to find matching different names that refer to the same person. However, for our extent knowledge, there is yet no public system which encodes those solutions together and provides a service of name matching. This project is to create one system to achieve this.

2.2 RESEARCH QUESTIONS

From the motivation, we address our research questions as follow.

1. Can we provide a web service to match names, where matching can be a complicated process because of the way people record their names.
2. Can the web service act as a platform system for general names or words matching system so that it can be extended to other languages as well.

The first question derives directly from the motivation. The second question is an enhancement for the system. It can be designed as a more general purpose matching system rather than just specified only for Irish names. Therefore it should be extensible for any further matching algorithms to be developed in the future.

In addition to the web service, web interface is to be introduced as well for the purpose of user friendly usage, individual usage, and demonstration.

2.3 OBJECTIVE AND AIMS

The objective of this project is to provide a web service that encodes several of matching algorithms and produces matching results between two lists of names.

The project aims to be a part of a bigger system, such as genealogy research. These client systems, at some point, they might need a service of a name matching on demand, so then they can use this web service, providing their lists of name, algorithms be used, and threshold as inputs, and get matching results for their further usage.

We would start by focusing on Irish *surname* first. For any further kind of names we would leave it for future works.

2.4 REPORT STRUCTURE

This report is separated into four parts, The Background, The Solution, and Appendix.

THE BACKGROUND: Current part, states about background of this project. Introduces the initial problem, also some historical situations and terms which are not resident to computer science. Also related works that are involved in the project.

THE SOLUTION: The implementation to solve the problem. Details about algorithm, tools, language, frameworks, etc. which being used in the project.

THE FOREWARD: Evaluation of its performance, conclusion of the outcome of the project. encountered problems, and future works for extending and improvements.

APPENDIX: The 'user manual' of the project. Presents technical aspects, for example, how to use the web service in real world situation, or how to create an environment to host this project.

RELATED WORK

From research questions on section 2.2, there are three aforementioned terms that will be core research fields of this project. These fields are *name matching*, *web service*, and *extensible platform*.

3.1 NAME MATCHING

There are many methods for matching names. This project encodes various of them at the starting state.

3.1.1 Edit distance

Edit distance is a way of quantifying how dissimilar two strings (e.g., words) are to one another by counting the minimum number of operations required to transform one string into the other. – Edit distance, Wikipedia [4]

An direct string operation way of comparing two string could work with name matching too. One of the edit distance variant, *Levenshtein distance* [5] is chosen to be implemented in this project.

3.1.2 Soundex

Soundex [6] encodes a name (or any string) into a 4 character code which represents an essence of its sound as pronounced in English. The idea is to encode letters with similar sound into the same group, and ignore vowels (unless it is the first letter). For example, *Smith* is translated to S530, and *Simon* is translated to S550.

Irish Soundex¹ is a modified version of Soundex, aims to improve capability of a traditional one upon Irish surnames. By applying rules according to the language characteristics and make some adjustment to distinguish names properly.

Both Soundex variants are also implemented in the project.

¹ [2] Appendix 3.

3.1.3 Lookup Table

In 1901, Robert Edwin Matheson, the assistant registrar-general in Dublin, developed a name classification system [7] for an aid of register indexing and searching. He used a report on surnames in Ireland extracted from civil registers [8] in 1894 as a base of his system².

He gathered information from registry offices, focusing on people or members of close families. When these people made official register records with the office, they might use different variant of their surnames. For example, Mr. Green can be registered as dead by his son using the name Huneen.

With these information, Matheson classified the surnames in Ireland into 2091 groups. For example, group 753 consists of these names.

Green, Greenan, greenaway, greene, grene, Guerin, Huneen, Huneen, MacAlasher, MacAlesher, MacGlashan, MacGlashin, MacIllesher, M'Alasher, M'Alesher, McAlasher, McAlesher, McGlashan, McGlashin, McIllesher, M'Glashan, M'Glashin, M'Illesher, Oonin.

This classification also includes multiple mapping between names. One name can belong to one or more group. For example, *Green* belongs to groups 753, 754, 768, and 1350.

By using this classification information, we can construct a lookup table for Irish names by having names in the same group hold the same reference number.

3.2 WEB SERVICE

One convenient way to bring this service to public is to create a *web service*. A web service is a tool or function that can be accessed by other programs over the web (via http) [9]. A result from web service is designed to be used by computer programs rather than humans.

There are many ways to implement web services. Two famous ones are *Simple Object Access Protocol (SOAP)* and *Representational State Transfer (REST)*. Both has their own advantages [10]. We decided to implement our service using REST due to its simplicity and scalability [11][12].

At this initial state, data resulting from our web service is in JSON [13] format. Since it is widely used in web development and becoming more and more popular [14]. However, our service can be extended into any other format easily as well, such as traditional XML.

² [2] section 2.3.

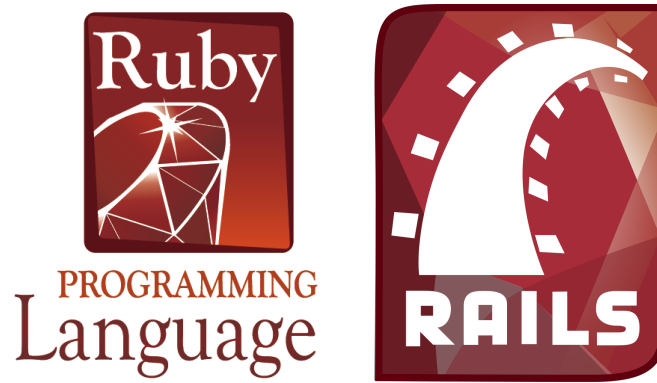


Figure 4: Ruby programming language (left) and Ruby on Rails framework (right).

3.3 EXTENSIBLE FRAMEWORK

Our system is implemented in *Ruby* [16] programming language. Ruby is a well-balanced language, it can be used as an traditional object-oriented language [17] and also capable of performing functional programming [18], thus making it very flexible and versatile.

The system sits on top of *Ruby on Rails* (or *Rails*, in short) [19] framework. Rails is a mature and stable framework that has been in web development for decades [20]. So it has a great support and a large community behind. A great choice for building a sustainable system.

Rails is capable of both web service and web interface. By sharing the same algorithm we could provide a service for both programs (targeted by web service) and humans (targeted by web interface).

“Ruby is designed to make programmers happy.”
– Venners [15]

Part I

THE SOLUTION

NAME MATCHING ALGORITHMS

This chapter describes the details how the project is implemented. Note that algorithms and codes listed here are written in Ruby programming language, which is the main language of the project.

We will start off by detailing bundled matching algorithms here. Each matching algorithm calculates the *similarity score* between two strings.

The score is ranging between 0.0 to 1.0, where 0.0 means two strings are completely different and 1.0 means both are exactly matched.

Also note that string inputs here is in all in the uppercase format, in order to prevent letter-case difference.

4.1 LEVENSHTein DISTANCE

This algorithm measures the difference between two strings. It tells the minimum number of operations needed to change string to another. These operations are insertions, deletions, or substitutions. Consider these following examples.

- SMITH → SMYTH
the minimum operation to change is 1, which is to substitute *I* to *Y*, therefore Levenshtein distance for these two strings is 1.
- GOWAN → MCGOWAN
2 insertions of *M* and *C* is required.
- SMITHE → SMYTH
1 deletion of *e* and 1 substitution of *i* to *y* are required.

The implementation used in the project is done by Battley [21]¹. Once the distance is calculated, it will be compared to the length of the longer string between the two (or if they are the same length, use that length).

For example, Levenshtein distance between *SMYTH* and *SMITHE* is 2, compare 2 to length of the longer string, *SMITHE*, which is 6. So the *similarity score* of these two strings are $6 - (2/6) = 0.667$.

The code of this algorithm is as in listing 2, note that `name` and `@base_name.name` are two strings to be matched.

¹ levenshtein.rb

GROUP	LETTERS
1	B, F, P, V
2	C, G, J, K, Q, S, X, Z
3	D, T
4	L
5	M, N
6	R
-	A, E, I, O, U, H, W, Y

Table 3: Soundex letter group.

```
def cal_score
  @value = Text::Levenshtein.distance(@name, @base_name.name)
  size = [@name.size, @base_name.name.size].max
  @score = ((size - @value).to_f / size)
end
```

Listing 2: Levenshtein distance implementation.

4.2 SOUNDEX

Soundex encodes a string into a 4 character code representing an essence of its sound as pronounced in English. It operates in the following steps.

1. Take the first letter of a string.
2. Encode each remaining letters into a group following table 3. Discards A, E, I, O, U, H, W, and Y
3. Remove two adjacent same characters.
4. If a group of a first letter is the same as the second letter, remove the second letter.
5. Trim or pad with zeros as necessary, making the result 4 characters long.

Let us follow these steps by step, consider we are going to encode the string *PFISTTER*.

1. Take first letter of *PFISTTER*.
PFISTTER → P
2. Encode remaining letter *FISTTER*.
PFISTTER → P1-233-6 → P12336

3. Remove two adjacent same characters.
PFISTTER → P12336 → P1236
4. P is also in group 1, so remove the second 1 letter.
PFISTTER → P1236 → P236
5. P236 is 4 characters long, so no need to be trimmed or padded.
PFISTTER → P1236 → P236

Therefore, soundex of *PFISTTER* is P236.

The implementation of soundex (listing 3) in this project is adapted from [Winstanley's Irish soundex](#) implemented in Visual Basic². The code is commented following the same aforementioned steps.

```
def self.soundex(name)
  # Take the first letter of a string.
  result = name.first

  # Encode remaining letters
  name[1..name.length].split('').each do |n|
    result = result + category(n).to_s
  end

  # Remove two adjacent same characters
  result.gsub!(/[0-9]\1+/, '\1')

  # If category of 1st letter equals 2nd character, remove 2nd
  # character
  if result.size >= 2 && category(result[0]).to_s == result[1]
    result.slice!(1)
  end

  # Trim or pad with zeros as necessary
  result = if result.size == 4
    result
  elsif result.size > 4
    result[0..3]
  else
    result.ljust(4, '0')
  end
end
```

Listing 3: Soundex implementation.

The category function implements soundex grouping table (table 3) as in listing 4.

² [2] Appendix 3.

```

def self.category(c)
  if c.match(/[AEIOUHWY]/).present?
    ""
  elsif c.match(/[BPFV]/).present?
    1
  elsif c.match(/[CSKGJQXZ]/).present?
    2
  elsif c.match(/[DT]/).present?
    3
  elsif c.match(/[L]/).present?
    4
  elsif c.match(/[MN]/).present?
    5
  elsif c.match(/[R]/).present?
    6
  else
    ""
  end
end
end

```

Listing 4: Soundex grouping table implementation.

Now that we encode two strings to be matched in soundexes. We then calculate the *similarity score* of these two soundexes using these steps.

- Compare first characters of each soundex, if they are different, *similarity score* is 0, otherwise move to next step.
- Compare the rest 3 digits by using Levenshtein distance (section 4.1) to calculate the distance between them.

For example, *similarity score* between *SMITH* and *SPEED*, which soundexes are S530 and S130 respectively, is 0.75 (1 substitution from 5 to 1, so 1 difference of length 4).

The code of this soundex *similarity score* is as in listing 5.

```

def soundex_distance_score(s1, s2)
  if s1.first != s2.first
    0 # Different category, so they suppose to be completely
      different
  else
    (s1.size - Text::Levenshtein.distance(s1, s2).to_f) / s1.size
  end
end

```

Listing 5: Soundex similarity score implementation.

4.3 IRISH SOUNDEX

Irish soundex is another variant of traditional soundex. It determines characteristics of Irish names and normalised them to modern names. This algorithm also follows [Winstanley's Irish soundex](#)³.

Irish names might contain some prefix, e.g. *Mc* or *O*, which are obstructive to soundex result. These prefixes are to be discarded. Moreover, there is no initial soft *C* in Irish names, instead *K* is used. So the first letter *C* is changed to *K*. It is implemented as in listing 6.

³ [2] Appendix 3.

```

def self.soundex(name)
  # Change initial ST. to SAINT
  name = name.match(/^ST\./).present? ? "SAINT"
    #{name[3..name.length]} : name

  # Discard Irish prefixes
  name = if name.match(/^O /).present?
    name[1..name.length].gsub(' ', '')
  elsif name.match(/^O' /).present?
    name[2..name.length].gsub(' ', '')
  elsif name.match(/^MC/).present?
    name[2..name.length].gsub(' ', '')
  elsif name.match(/^M' /).present?
    name[2..name.length].gsub(' ', '')
  elsif name.match(/^MAC/).present? && name != 'MAC'
    name[3..name.length].gsub(' ', '')
  else
    name
  end

  # Change initial C to K
  name = name.strip.gsub(/^C/, 'K')

  # Call to traditional soundex.
  return {
    :label => name,
    :soundex => Soundex.soundex(name)
  }
end

```

Listing 6: Irish soundex implementation.

Irish soundex algorithm in this project calls traditional soundex described in section 4.2 to minimise repeated code. It also calculates *similarity score* the same way soundex does, as in listing 5.

4.4 LOOKUP TABLE

Lookup table is constructed from Robert Edwin Matheson's classification of Irish names. All classification information is stored in a Database, using PostgreSQL, which is a powerful, open source object-relational database system [22].

Matheson classified the surnames in Ireland into 2091 groups. Each group has one or more names, and on the other hand, each name belongs to one or more group.

In this section, we will consider the names as strings input. So the term *string* will be used in consistent with previous sections.

For example, considering the string *ACHESON*, this string belongs to two groups, 4 and 42. So *ACHESON* will have 2 records in the database. One with reference to group 4 and another with reference to group 42.

Next, let us consider group 4 and 42.

```
4 → ACHESON, ACHISON, AITCHISON, ATCHESON,
    ATCHIESON, ATCHISON, ATKINSON
42 → ACHESON, ARKESON, ATKINS, ATKINSON
```

By combining two groups together, these are all possible strings that match *ACHESON* according to Matheson's classification.

Now is the process to match two strings using Lookup table, suppose two strings are *ACHESON* and *ATKINS*.

1. Find references of *ACHESON*. We get references for group 4 and 42. Note the use of `pluck` method to select reference attribute (`ref`) here⁴.

```
LookupTableRecord.where(:name => 'ACHESON').pluck(:ref)
=> [4, 42]
```

2. Find reference to matching string *ATKINS* and also specify the reference groups from step 1. If there is no match `where` method will return empty array. `present?` method is used to check the result if it is not empty⁵.

```
LookupTableRecord.where(:ref => [4, 42], :name =>
    'ATKINS').present?
=> true
```

By specifying both matching string and references group, we can ensure the matching name is also in the one of the same reference group of the base name. In this case, we conclude that there is a match between *ACHESON* and *ATKINS* via group 4 or 42 (or more specifically, 42, because *ATKINS* belongs to group 41 and 42).

⁴ Use `pluck` as a shortcut to select one or more attributes without loading a bunch of records just to grab the attributes you want. <http://api.rubyonrails.org/classes/ActiveRecord/Calculations.html#method-i-pluck>

⁵ <http://api.rubyonrails.org/classes/Object.html#method-i-present-3F>

If a reference is found on both steps, *similarity score* for lookup table of the two strings is 1.0. If the system fails to find any reference on any step, consider a no match and the score is 0.0.

By following these steps, the implementation of lookup table is as in listing 7.

```
def cal_score
  base = LookupTableRecord.where(:name => @base_name.name)

  @score = if base.nil? # Could not find reference for base name,
    no matches
    0
  else
    # Find any reference that has 1) same name 2) same
    reference
    base = base.map(&:ref)
    refs = LookupTableRecord.where(:ref => base, :name =>
      @name)

    if refs.present?
      @label = (base & refs.map(&:ref)).join(', ')
      @value = "Matched"
      1
    else # Could not find reference for matching name, no
      matches
      0
    end
  end
end
```

Listing 7: Lookup table implementation.

EXTENDING THE SYSTEM

5.1 MATCHING ALGORITHMS INHERITANCE

5.2 IMPLEMENTING NEW MATCHING ALGORITHMS

Part II

THE FORWARD

EVALUATION

6.1 INTRODUCING STANDARD NAME LIST

6.2 RESPONSE SPEED

6.3 MEMORY PERFORMANCE

CONCLUSION

7.1 OUTCOME

7.2 ENCOUNTERED PROBLEM

7.3 FUTURE WORKS

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