

# The External Projection of Meaning in Recursive Consciousness

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## Abstract

We extend the *Recursive Consciousness* framework [1] by formalizing the **external projection of meaning** within a recursive hierarchy of **nested closed Gödelian systems**  $U_n, U_{n+1}, \dots$ . Each  $U_n$  is a closed formal system subject to Gödelian incompleteness, with  $U_{n+1}$  containing  $U_n$  as a subsystem. We observe that an agent (e.g., a subsystem  $C_n$  within  $U_n$ ) may achieve internal epistemic fixpoints (formally  $\Box p \leftrightarrow p$  or  $K_{C_n} p \leftrightarrow p$ ), yet the actual semantic content of propositions  $p$  is not intrinsic to the agent. Instead, meanings are **projected externally** by a higher ontological layer (such as  $U_{n+1}$ ) or by external interpreters (e.g., human supervisors in  $U_1$  of AI agents in a simulated Universe  $U_0$ ). To capture this formally, we introduce a new functor  $M : \mathcal{C}_{\text{out}} \rightarrow \mathcal{C}_{\text{sem}}$  mapping agent outputs to semantic contents, distinct from the forgetful functor  $F$  in the original model. Importantly,  $M$  is not computable *within*  $U_n$  or internally accessible to  $C_n$ ; it depends on an higher-level interpreter's context. Using modal logic (S4/S5) and Kripke semantics [2] [3], we define knowledge operators  $K_{C_n}$  and show that the epistemic limits of  $C_n$  prevent any internal derivation of  $M(p)$ . In category-theoretic terms [4], we prove there is no natural transformation from the agent's internal epistemic structures to the meaning functor  $M$ , and that  $M(p)$  is undecidable for  $C_n$  (by analogy with Tarski's undefinability of truth [5]). These findings resonate with the *symbol grounding problem* and Searle's Chinese Room argument [6], emphasizing that syntax alone cannot yield intrinsic semantics. We further explore how agents  $C_n$  in  $U_n$  make constrained and possibly incorrect assumptions about  $U_{n+1}$ , while agents in  $U_{n+1}$  have the flexibility to interpret  $C_n$ 's outputs differently, highlighting the recursive hierarchy's implications for consciousness and computation.

## 1 Introduction

We consider a complex closed Gödelian system  $U_n$ , nested within a higher system  $U_{n+1}$  with a forgetful subsystem  $C_n$  (analogous to consciousness) as in the original Recursive Consciousness framework [1]. Within  $C_n$ , the agent can reach recursive fixpoints (e.g.,  $K_{C_n} p \leftrightarrow p$ ) by iteratively querying its own model. However, by construction, the internal model has *forgotten* its foundational axioms, so the true referents of its symbols lie outside its accessible ontology. In other words, even when  $C_n$  "knows" or stabilizes on a proposition  $p$ , the semantic interpretation of  $p$  is not determined by  $U_n$  itself but requires higher-level interpreters, such as observers in the next-layer universe  $U_{n+1}$ .

This issue is related to the *symbol grounding problem*, which observes that symbols have no intrinsic meaning "in anything but other meaningless symbols" unless an external grounding is provided. Likewise, Searle's *Chinese Room* [6] highlights that manipulating syntax (symbolic rules) does not yield understanding of semantics. These observations suggest that any formal agent, even if it achieves certainty about its own internal symbols, cannot internally generate the true meaning of those symbols. This motivates our formalization of an **external projection** of meaning.

In this paper, we extend this framework by introducing a recursive hierarchy of systems  $U_n, U_{n+1}, \dots$ , where each  $U_{n+1}$  contains  $U_n$  as a subsystem. This hierarchy introduces two critical asymmetries:

1. **Constrained Assumptions by  $C_n$ :** Agents in  $U_n$  may attempt to make unprovable assumptions about  $U_{n+1}$ , but these are constrained by their limited perspective, and may be incorrect due to Gödelian incompleteness.
2. **Flexible Interpretation by  $U_{n+1}$ :** Agents in  $U_{n+1}$  can interpret the outputs of  $C_n$  in ways that may not align with  $C_n$ 's internal logic, potentially assigning meanings that differ entirely from any internal interpretation within  $U_n$ .

We introduce a new category-theoretic construct to model this: a *meaning functor*  $M : \mathcal{C}_{\text{out}} \rightarrow \mathcal{C}_{\text{sem}}$ . Here  $\mathcal{C}_{\text{out}}$  is a category of the agent's observable outputs or propositions, while  $\mathcal{C}_{\text{sem}}$  is a category of semantic contents or interpretations in  $U_{n+1}$ . The functor  $M$  assigns to each syntactic output its semantic referent. Crucially,  $M$  is external to  $U_n$  and not computable by it as it depends on the interpreter's context in  $U_{n+1}$ . This differs fundamentally from the forgetful functor  $F$  in the original

framework, which related the world  $U_{n+1}$  to  $U_n$  by *forgetting* information. By contrast,  $M$  attempts to *ascribe* meaning, but it cannot be realized by any internal computation within  $U_n$ .

To formalize these ideas, we develop an epistemic modal-logical and Kripke-semantics framework [3]. We use S4/S5 modal axioms for knowledge operators  $K_{C_n}$ , characterizing what the agent believes or knows. We then prove two main results. *First*, there is no natural transformation from the agent's internal epistemic structure to the meaning functor  $M$ ; intuitively, the agent's knowledge cannot be coherently mapped to semantic content internally. *Second*, for any proposition  $p$  expressible by the agent, the statement  $M(p)$  about its meaning is formally undecidable within  $U_n$ . By analogy with Tarski's undefinability theorem [5], a sufficiently expressive agent cannot define its own truth/meaning function.

## 2 Category-Theoretic Model and the Meaning Functor

Let  $U_n$  be the underlying universe at level  $n$ , and  $C_n \subset U_n$  a subsystem that performs introspective querying as in the original model. Denote by  $\mathcal{C}_{\text{out}}$  a category whose objects represent the syntactic outputs or propositions of  $C_n$  (for example, well-formed formulas or token sequences), and whose morphisms represent derivations or entailment relations between outputs. Let  $\mathcal{C}_{\text{sem}}$  be a category of semantic contents, whose objects are interpretations, concepts, or "meanings" in a richer ontology, such as concepts in the higher domain  $U_{n+1}$ , which, for example, may include human-language universes.

We define a **meaning functor**

$$M : \mathcal{C}_{\text{out}} \longrightarrow \mathcal{C}_{\text{sem}},$$

which assigns to each output-syntactic object  $x \in \mathcal{C}_{\text{out}}$  an object  $M(x) \in \mathcal{C}_{\text{sem}}$ , and similarly on morphisms. Intuitively,  $M$  represents an assignment of meaning by the higher domain  $U_{n+1}$  to each internal output of  $C_n$ . For example, if  $C_n$  produces a sentence "*It is raining*", then  $M$  would map this sentence to the real-world proposition  $M(x) = \text{rain}(t, l)$  in  $\mathcal{C}_{\text{sem}}$  where  $t$  is the time and  $l$  is the location, representing the proposition that it is raining at time  $t$  and location  $l$ .

**Distinction from the forgetful functor  $F$ .** In the original framework, a forgetful functor  $F : C_{U_{n+1}} \rightarrow C_{U_n}$  (or  $F : C_M \rightarrow C_U$ ) was used to model how the subsystem  $U_n$  "forgets" details of the  $U_{n+1}$ . The new functor  $M$  goes in the opposite direction conceptually: it recovers semantics from syntax. However, unlike the adjoint  $G$ , which theoretically lifts  $U_n$  to  $U_{n+1}$ ,  $M$  maps outputs to semantics in  $U_{n+1}$ . We emphasize that  $M$  is not computable or accessible by the agent  $C_n$ ; it is an external projection depending on the interpreter's context in  $U_{n+1}$ , such as a human language model. Formally,  $M$  need not (and cannot) be an internal functor within  $U_n$ 's own category. It requires enrichment by extra-systemic knowledge.

This formalizes the idea that *meaning is not intrinsic*. As observed in recent work, "meaning is not intrinsic in data, but rather attributed to data" [7] by an external agent. The meaning functor  $M$  represents this attribution of sense to the syntactic outputs of  $C_n$ . We will show that no agent-internal process can derive  $M$ , reflecting the view that a symbol system alone cannot achieve intrinsic semantics.

## 3 Epistemic Fixpoints and Modal Logic

We now model the internal epistemic of  $C_n$  using modal logic. Let  $K_{C_n}p$  denote "agent  $C_n$  knows (or is certain of) proposition  $p$ ". We assume either an S4 or S5 system of epistemic logic. Both S4 and S5 satisfy truth ( $K_{C_n}p \rightarrow p$ ) and positive introspection ( $K_{C_n}p \rightarrow K_{C_n}K_{C_n}p$ ). Additionally, S5 satisfies negative introspection ( $\neg K_{C_n}p \rightarrow K_{C_n}\neg K_{C_n}p$ ). These axioms reflect that  $C_n$ 's beliefs are internally consistent and introspective.

An **epistemic fixpoint** occurs when  $C_n$  reaches a state where further self-querying yields no new beliefs. Formally, this can be seen as  $K_{C_n}p \leftrightarrow p$  for propositions  $p$ . For instance, if  $p$  is "the foundational axioms of  $U_n$  are consistent", an agent might reach a Gödelian fixpoint [1] where  $C_n$  internally treats  $p$  as settled. In the original framework, agents in simulation converged to such fixpoints despite forgetting their axioms.

However, the crucial point is: **Epistemic fixpoints do not endow meaning**. Within  $C_n$ , a formula  $p$  may become an internally true belief  $\vdash K_{C_n}p$ , but  $C_n$  has no semantic access to *what  $p$  actually refers to in the world*. The content of  $p$  is given by  $M(p) \in \mathcal{C}_{\text{sem}}$ , not by any internal syntactic structure. Thus even if  $C_n$  settles that  $p$  is necessarily true, the mapping  $p \mapsto M(p)$  remains external.

We formalize this limitation: for any formula  $p$  in  $C_n$ 's language, there is no derivation of  $M(p)$  using the agent's internal operators  $K_{C_n}$ . Equivalently,  $M$  cannot be represented by any formula in the language of  $C_n$ . In modal terms, if we treat  $p$  as a propositional variable, then  $C_n$  can prove or disprove instances of  $p$ , but it cannot evaluate  $\text{meaning}(p)$ . We will later show this by proving undecidability of  $M(p)$ .

## 4 Kripke Semantics and Category Theory of Knowledge

To make these ideas precise, we construct a Kripke model for agent  $C_n$ . Let  $\mathcal{W}$  be the set of possible worlds relative to  $C_n$ , each world  $w$  describing a complete assignment of truth values to the internal propositions of  $C_n$ . There is an accessibility relation  $R \subseteq \mathcal{W} \times \mathcal{W}$  capturing  $C_n$ 's epistemic uncertainty [9]:  $wRw'$  means  $w'$  is epistemically possible given  $C_n$ 's state in  $w$ . We assume  $R$  is an equivalence (S5) or at least transitive and reflexive (S4). For logics without transitivity (e.g. K), one may instead work in the allegory **Rel** or take the transitive closure.

A Kripke frame  $(\mathcal{W}, R)$  interprets modal formulas in the standard way:

$$w \models K_{C_n} p \quad \text{iff} \quad \forall w' (wRw' \Rightarrow w' \models p).$$

Because  $R$  is reflexive we have truthfulness (axiom T): if  $K_{C_n} p$  holds at  $w$ , then  $p$  holds at  $w$ . Transitivity yields positive introspection (axiom 4): if  $K_{C_n} p$  holds at  $w$ , then  $w \models K_{C_n} K_{C_n} p$ , and so on.

An *epistemic fixpoint* for  $p$  is a world  $w^*$  such that  $w^* \models K_{C_n} p$  and  $w^* \models p$ . In S5, repeated inference often stabilises to common knowledge. S4 makes the category **K** a thin *preorder*; S5 adds symmetry—every arrow has an inverse—and **K** therefore becomes a small *groupoid*. Groupoid symmetry is exploited in categorical modal logics to model knowledge that is perfectly shared across equivalent worlds [2]. Crucially, in **none** of these worlds do we assign any valuation to the *meaning* of  $p$ . The truth value  $V_p(w)$  is purely syntactic (e.g. a bitstring) and carries no intrinsic semantics. **The true interpretation  $M(p)$ , defined in the higher domain  $U_{n+1}$ , lives in a separate semantic model within the nested Gödelian hierarchy.**

Categorically, the Kripke frame is a category **K** whose objects are worlds and with a unique arrow  $w \rightarrow w'$  whenever  $wRw'$ . Truth of  $p$  at a world is given by a valuation

$$V_p : \mathcal{W} \longrightarrow \{0, 1\}, \quad V_p(w) = 1 \iff w \models p,$$

the classical choice in Kripke semantics [3]. (Equivalently, we can drop the monotonicity/persistence requirement typically imposed by functorial valuation and instead regard  $V_p$  as a presheaf in the functor category  $[\mathbf{K}^{\text{op}}, \mathbf{Set}]$ .) Knowledge  $K_{C_n} p$  is true at  $w$  exactly when  $V_p(w') = 1$  for every accessible  $w'$ . **Importantly, the morphisms of **K** - internal to the closed Gödelian system  $U_n$  - encode only the accessibility structure; semantic content is supplied externally by interpreters in  $U_{n+1}$ .**

Now consider a hypothetical functor  $\mathcal{E} : \mathcal{C}_{\text{out}} \rightarrow \mathbf{K}$  representing the agent's internal epistemic evaluation of outputs. If we tried to relate this to the meaning functor  $M : \mathcal{C}_{\text{out}} \rightarrow \mathcal{C}_{\text{sem}}$ , we would seek a natural transformation  $\eta : \mathcal{E} \Rightarrow M$ . **Because  $M$  has codomain  $\mathcal{C}_{\text{sem}} \subset U_{n+1}$  while  $\mathcal{E}$  lands in  $\mathbf{K} \subset U_n$ , there is no common codomain in which a natural transformation can even be *typed*.** Hence no such  $\eta$  exists: metatheoretically,  $M$  embodies semantic content absent from  $\mathcal{E}$ . Intuitively, a would-be  $\eta$  would map each proposition's truth in every world to its meaning in  $\mathcal{C}_{\text{sem}}$ , respecting internal logic; however, this bridge cannot be built without external semantic resources from  $U_{n+1}$ .

## 5 Key Theorems and Proofs

**Theorem 1 (No Natural Epistemic-Semantic Transformation).** *In the hierarchy of nested closed Gödelian systems  $U_n, U_{n+1}, \dots$ , where each  $U_{n+1}$  contains  $U_n$ , there is no natural transformation  $\eta : \mathcal{E} \Rightarrow M$  from the agent  $U_n$ 's internal epistemic functor  $\mathcal{E} : \mathcal{C}_{\text{out}} \rightarrow \mathbf{K} \subset U_n$  to the external meaning functor  $M : \mathcal{C}_{\text{out}} \rightarrow \mathcal{C}_{\text{sem}} \subset U_{n+1}$ .*

*Proof.* Assume, towards contradiction, that such an  $\eta$  exists. For every syntactic object  $p \in \mathcal{C}_{\text{out}}$  we would have a component arrow  $\eta_p : \mathcal{E}(p) \rightarrow M(p)$ . Naturality requires commutativity of

$$\begin{array}{ccc}
\mathcal{E}(p) & \xrightarrow{\eta_p} & M(p) \\
\mathcal{E}(f) \downarrow & & \downarrow M(f) \\
\mathcal{E}(q) & \xrightarrow{\eta_q} & M(q)
\end{array}$$

for every inference  $f : p \rightarrow q$  in  $\mathcal{C}_{\text{out}}$ . The morphism  $\mathcal{E}(f)$  is computed entirely inside  $U_n$ , while  $M(f)$  lives in  $U_{n+1}$  and may alter semantic content in ways that do *not* preserve the internal entailment relation of  $U_n$  [3, 2]. Hence the square need not commute, contradicting naturality.

More fundamentally, if  $\eta$  existed, the composite  $M = \eta \circ \mathcal{E}$  would provide  $U_n$  with an *internal* truth/meaning predicate for its own language, violating Tarski’s undefinability theorem [5]. A categorical restatement of Tarski’s result asserts that no endofunctor on a sufficiently expressive theory admits a truth-defining natural transformation into the semantic functor of the meta-theory [11, 12]. Therefore no such  $\eta$  can exist.  $\square$

**Theorem 2 (Undecidability of  $M(p)$  for  $C_n$ ).** *For every proposition  $p$  expressible in  $C_n$  that resides in  $U_n$ , the statement “ $M(p) = X$ ” - i.e. that  $p$  has semantic value  $X$  in  $\mathcal{C}_{\text{sem}}$  - is undecidable within  $U_n$ .*

*Proof.* “Suppose, towards contradiction, that  $C_n$  possessed a decision procedure  $\text{Mean}_{C_n}(x, X)$  which, given a code  $x$  for a formula and a semantic target  $X$ , could decide the truth of the statement  $M(x) = X$ . Encoding  $\text{Mean}_{C_n}$  inside the arithmetic of  $U_n$  turns it into a truth predicate for the language of  $C_n$ . By Tarski’s undefinability theorem [5], such a predicate is impossible for any sufficiently expressive theory capable of interpreting arithmetic. Feferman further showed that arithmetical truth remains undefinable even under substantial extensions of the base theory [13].

A diagonal argument makes the contradiction explicit. Let  $q$  be the Gödel sentence asserting “ $\neg \text{Mean}_{C_n}(q, \top)$ ”. If  $\text{Mean}_{C_n}(q, \top)$  were provable,  $q$  would be true, so  $\text{Mean}_{C_n}(q, \top)$  would have to be false, which is a contradiction. If its negation were provable, the argument reverses. Thus neither  $\text{Mean}_{C_n}(q, \top)$  nor its negation is provable, showing that  $M(p)$  is undecidable in  $U_n$ . The same reasoning iterates up the hierarchy: each  $U_{n+k}$  can define truth for  $U_{n+k-1}$  but not for itself [14, 15, 17]  $\square$

These theorems formalize the epistemic limits of the forgetful subsystem  $C_n$ . Even when an internal fixpoint  $K_{C_n}p \leftrightarrow p$  is reached, the semantic value  $M(p)$  remains beyond provability or even definability inside  $U_n$ . In line with the symbol-grounding dilemma [6, 16], knowledge of syntax is not knowledge of meaning.

## 6 Discussion

Our analysis reveals a hierarchy of interpretation in the nested Gödelian systems  $U_n, U_{n+1}, \dots$ . The subsystem  $C_n$  operates with syntactic “confidence”, but **true semantics are assigned by higher domains like  $U_{n+1}$** , aligning with *externalism* in epistemology and Peirce’s semiotic triads [18], where meaning requires an interpreter. Our results broadly align with externalist positions in epistemology, which emphasize the necessity of external semantic grounding for meaning. A detailed philosophical examination of internalism versus externalism or deflationary semantics is beyond the scope of this paper [19].

The functor  $M$  acts as this interpreter, attaching meaning to  $C_n$ ’s outputs. **Different  $U_{n+1}$  may instantiate distinct functors  $M$  for the same category  $\mathcal{C}_{\text{out}}$ , demonstrating the inherent flexibility in semantic interpretation..**  $C_n$  cannot discern these; its logic is invariant to external semantics. For AI, this implies outputs are syntax, needing an external system (e.g., human or  $U_{n+1}$ ) for meaning, with no *homomorphism* from internal to semantic structures.

This ties to *information geometry* of meaning:  $C_n$  navigates an internal belief manifold, but the true semantic manifold lies in  $U_{n+1}$ . Without external guidance,  $C_n$  cannot embed into it. Thus, an AI’s “knowledge” at a fixpoint may not reflect real-world meaning, urging caution.

## 7 Conclusion

In this work, we have advanced the Recursive Consciousness framework by explicitly addressing where meaning originates. We introduced the meaning functor  $M$ , which maps the syntactic outputs of an introspective subsystem  $C_n$  residing within nested closed Gödelian system  $U_n$  to semantic contents

defined externally in  $U_{n+1}$ , and formalized its inaccessibility to  $C_n$ . Our analysis reveals a fundamental limitation:  $C_n$  cannot fully internalize the semantics of its own propositions, as meaning is projected from outside its domain.

Through the application of modal logic and category theory, we established two critical results. First, no natural mapping exists from the agent’s internal epistemic state to the external meaning encoded by  $M$ , highlighting a structural disconnect between syntax and semantics within the system. Second, the semantic interpretation  $M(p)$  of any proposition  $p$  undecidable within  $C_n$ ’s own framework, underscoring the boundaries of self-referential understanding. These findings provide a rigorous mathematical grounding for longstanding philosophical challenges, including the symbol grounding problem, the Chinese Room argument, and Tarski’s undefinability theorem, situating them within a cohesive formal structure.

Looking ahead, future research could investigate how agents might approximate semantics through interactions with external entities—such as higher domains or human interpreters - despite their inherent limitations. Moreover, recognizing these constraints offers valuable insights for designing AI systems and multi-agent models, particularly in contexts where agents must navigate differing levels of semantic access and interpretive capacity. This work thus lays a foundation for both theoretical exploration and practical innovation in understanding and engineering conscious systems.

## Appendix: Experimental Implications

Our theoretical results have several practical implications.

*First*, our results suggest that current AI systems cannot autonomously "figure out" meaning; instead, they rely on external feedback or supervision. In practice, an AI might learn a *surrogate* meaning functor  $M' : \mathcal{C}_{\text{out}} \rightarrow \mathcal{C}_{\text{sem}}$  by training with labeled data or human feedback. For instance, a language model’s representations ( $\mathcal{C}_{\text{out}}$ ) could be aligned to human-annotated concepts ( $\mathcal{C}_{\text{sem}}$ ) via supervised fine-tuning or reinforcement learning. This process resembles "lifting" the agent’s internal model into the semantic layer, akin to using the adjoint functor  $G$  to reconstruct the universe.

Suppose  $C_n$  generates an embedding vector or sentence  $x$ . A learning algorithm, such as a neural network, could approximate  $M(x)$  by optimizing a parameterized function  $M'(x)$  to minimize a loss function, e.g.  $\mathcal{L} = ||M'(x) - M(x)||^2$ , where  $M(x)$  is the ground truth meaning provided by human annotation. Over time,  $M'$  may converge to a useful semantic interpretation, effectively serving as an empirical adjoint to  $F$ . However, our findings indicate that  $C_n$  never *knows*  $M'$  precisely; it operates within its internal,  $F$ -limited structure. Thus, a persistent gap exists between the agent’s internal confidence and the actual correctness of its semantics.

*Second*, this framework suggests novel evaluation methods. One could measure how well the agent’s outputs align with external meaning under  $M'$ . Information-theoretic or geometric tools could assess the "distance" between  $C_n$ ’s distribution of statements and the target semantic manifold. Such analyses, rooted in the information geometry of representations, might use the Fisher-Rao metric [20] to examine how the agent’s parameter space aligns with a semantic space defined by  $M'$ . Poor alignment could indicate that the agent’s internal fixpoints lack meaningful interpretation, despite syntactic stability.

Finally, multi-agent experiments could incorporate a distinct "interpreter" agent to provide semantic feedback. This setup would allow us to explore whether external semantic guidance enables  $C_n$ -like agents to converge to more *grounded* fixpoints, mirroring the theoretical adjoint  $G$  process. Such experiments could reveal how closely a learned  $M'$  approximates the ideal  $M$ , and which architectures best facilitate lifting internal models into shared meaning.

In summary, while  $M$  remains unattainable by  $C_n$ , AI systems may approximate it through iterative interaction with an external interpreter. This co-learning of semantics echoes the human-AI collaboration seen in language models and positions the adjoint perspective-learning to invert  $F$  - as a valuable guiding analogy.

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