

# Simulation and High-Performance Computing

## Part 4: Higher-order Timestepping Methods

Steffen Börm

Christian-Albrechts-Universität zu Kiel

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# Reminder: Explicit timestepping methods

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# General Runge-Kutta method

**Observation:** All single-step methods introduced in this lecture so far share a common form:

$$\tilde{y}(t + \delta) = y(t) + \delta \sum_{i=1}^n b_i k_i,$$
$$k_i = f \left( t + \delta c_i, y(t) + \delta \sum_{j=1}^n a_{ij} k_j \right) \quad \text{for all } i \in [1 : n].$$

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- **Explicit method** if  $a_{ij} = 0$  for  $j \geq i$ ,
- **Semi-implicit method** if  $a_{ij} = 0$  for  $j > i$ ,
- **Implicit method** otherwise.

# Butcher tableaux

Idea: Collect all coefficients in a simple representation.

$$\begin{array}{c|ccc} c_1 & a_{11} & \cdots & a_{1n} \\ \vdots & \vdots & \ddots & \vdots \\ c_n & a_{n1} & \cdots & a_{nn} \\ \hline & b_1 & \cdots & b_n \end{array}$$



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Runge and Crank-Nicolson:

$$\begin{array}{c|cc} 0 & 0 & \\ 1/2 & 1/2 & 0 \\ \hline & 0 & 1 \end{array}$$

$$\begin{array}{c|cc} 0 & 0 & \\ 1 & 1/2 & 1/2 \\ \hline & 1/2 & 1/2 \end{array}$$

# Classic fourth-order Runge-Kutta

Idea: Based on Simpson's quadrature.

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# Classic Runge-Kutta: Implementation

**Approach:** Store approximated derivatives  $k_1, \dots, k_4$  in auxiliary variables.

```
dx1 = v;  
dv1 = -c/m * x;  
  
dx2 = v + 0.5 * delta * dv1;  
dv2 = -c/m * (x + 0.5 * delta * dx1);  
  
dx3 = v + 0.5 * delta * dv2;  
dv3 = -c/m * (x + 0.5 * delta * dx2);  
  
dx4 = v + delta * dv3;  
dv4 = -c/m * (x + delta * dx3);  
  
x += delta * (dx1 + 2.0 * dx2 + 2.0 * dx3 + dx4) / 6.0;  
v += delta * (dv1 + 2.0 * dv2 + 2.0 * dv3 + dv4) / 6.0;
```

# Experiment: Crank-Nicolson vs Runge-Kutta

**Approach:** Start at  $t = 0$ , perform successive timesteps to reach  $t = 10$ .

$\delta$	Crank-Nic		Runge-Kutta	
	error	ratio	error	ratio
1/2	9.2 <sub>-2</sub>		8.1 <sub>-4</sub>	
1/4	2.7 <sub>-2</sub>	3.4	1.2 <sub>-4</sub>	6.9
1/8	7.0 <sub>-3</sub>	3.9	9.2 <sub>-6</sub>	12.6
1/16	1.8 <sub>-3</sub>	4.0	6.4 <sub>-7</sub>	14.5
1/32	4.4 <sub>-4</sub>	4.0	4.1 <sub>-8</sub>	15.3
1/64	1.1 <sub>-4</sub>	4.0	2.6 <sub>-9</sub>	15.7
1/128	2.8 <sub>-5</sub>	4.0	1.7 <sub>-10</sub>	15.8
1/256	6.9 <sub>-6</sub>	4.0	1.1 <sub>-11</sub>	15.9
1/512	1.7 <sub>-6</sub>	4.0	6.6 <sub>-13</sub>	15.9

**Observation:** Classic Runge-Kutta is indeed of fourth order, since the error behaves like  $\delta^4$ .



# Multistep methods

**Problem:** Higher-order Runge-Kutta methods are computationally expensive.

**Idea:** Re-use results computed in previous steps in order to save time.

**Approach:** We let  $t_i := t_0 + \delta i$ .

An  $m$ -step method computes  $y(t_{i+1})$  based on  $y(t_i), \dots, y(t_{i-m+1})$  (and maybe additional data corresponding to these previous states).

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**Example:** Leapfrog method,  $y(t + \delta)$  depends on  $y(t)$  and  $y(t + \frac{\delta}{2})$ .

# Adams-Bashforth method

Idea: Fundamental theorem of calculus states

$$y(t_{i+1}) = y(t_i) + \int_{t_i}^{t_{i+1}} y'(s) ds.$$

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Approximate the interval by replacing  $y'$  with the interpolating polynomial

$$p(s) = \sum_{j=0}^m y'(t_{i-j}) \ell_{i,j}(s)$$

in the points  $t_i, t_{i-1}, \dots, t_{i-m}$  with Lagrange polynomials  $\ell_{i,0}, \dots, \ell_{i,m}$ .

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$$\begin{aligned} y(t_{i+1}) &\approx y(t_i) + \int_{t_i}^{t_{i+1}} p(s) ds = y(t_i) + \sum_{j=0}^m y'(t_{i-j}) \underbrace{\int_{t_i}^{t_{i+1}} \ell_{i,j}(s) ds}_{=: a_{ij}} \\ &= y(t_i) + \sum_{j=0}^m a_{ij} f(t_{i-j}, y(t_{i-j})). \end{aligned}$$

# Adams-Bashforth coefficients

Equidistant points  $t_i = t_0 + \delta i$  imply

$$\ell_{i,j}(s) = \prod_{\substack{k=0 \\ k \neq j}}^m \frac{s - t_{i-k}}{t_{i-j} - t_{i-k}} = \prod_{\substack{k=0 \\ k \neq j}}^m \frac{s - t_i + \delta k}{\delta(k - j)}$$

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Coefficients given by

$$a_{ij} = \int_{t_i}^{t_{i+1}} \ell_{i,j}(s) ds$$



# Adams-Bashforth coefficients

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$$\ell_{i,j}(s) = \prod_{\substack{k=0 \\ k \neq j}}^m \frac{s - t_{i-k}}{t_{i-j} - t_{i-k}} = \prod_{\substack{k=0 \\ k \neq j}}^m \frac{s - t_i + \delta k}{\delta(k - j)} = \prod_{\substack{k=0 \\ k \neq j}}^m \frac{\hat{s} + k}{k - j}$$

with  $s = \delta \hat{s} + t_i$ .

Coefficients given by

$$a_{ij} = \int_{t_i}^{t_{i+1}} \ell_{i,j}(s) ds = \delta \int_0^1 \ell_{i,j}(\delta \hat{s} + t_i) d\hat{s}$$

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The same coefficients are used in all timesteps.

# Computing the weights

**Idea:** Monomials  $p_i(t) = t^i$  have to be integrated exactly,

$$\sum_{j=0}^m w_j p_i(-j) = \int_0^1 p_i(s) ds.$$

**Example:** Interpolation points  $0, -1, -2, -3$ , weights have to solve

$$\begin{array}{ll} w_0 + w_1 + w_2 + w_3 = 1, & (p_0(t) = 1) \\ -w_1 - 2w_2 - 3w_3 = 1/2, & (p_1(t) = t) \\ w_1 + 4w_2 + 9w_3 = 1/3, & (p_2(t) = t^2) \\ -w_1 - 8w_2 - 27w_3 = 1/4. & (p_3(t) = t^3) \end{array}$$

The solution is  $w_0 = \frac{55}{24}$ ,  $w_1 = -\frac{59}{24}$ ,  $w_2 = \frac{37}{24}$ ,  $w_3 = -\frac{9}{24}$ .

# Adams-Bashforth algorithm

**Idea:** Store  $y_i := \tilde{y}(t_i)$  and  $f_i := f(t_i, \tilde{y}(t_i))$ .

$$y_{i+1} := y_i + \delta \sum_{j=0}^m w_j f_{i-j},$$

$$f_{i+1} := f(t_{i+1}, y_{i+1}).$$

**Observation:** We require only one evaluation of  $f$  per step.

**Storage:** We have to store the derivatives  $f_i, \dots, f_{i-m}$ .

Older derivatives can be cyclically overwritten.

**Problem:** We need approximations for the first  $m + 1$  states before we can start the Adams-Bashforth algorithm.

# Adams-Bashforth: Implementation

**Approach:** Store previous states in arrays  $x$ ,  $v$  and previous derivatives in arrays  $dx$  and  $dv$ .

```
for(i=3; i<n; i++) {  
    x[(i+1)%4] = x[i%4] + delta * (w0 * dx[i%4]  
                                   + w1 * dx[(i-1)%4]  
                                   + w2 * dx[(i-2)%4]  
                                   + w3 * dx[(i-3)%4]);  
    v[(i+1)%4] = v[i%4] + delta * (w0 * dv[i%4]  
                                   + w1 * dv[(i-1)%4]  
                                   + w2 * dv[(i-2)%4]  
                                   + w3 * dv[(i-3)%4]);  
  
    dx[(i+1)%4] = v[(i+1)%4];  
    dv[(i+1)%4] = -c/m * x[(i+1)%4];  
}
```

# Experiment: Runge-Kutta vs Adams-Bashforth

**Approach:** Start at  $t = 0$ , perform successive timesteps to reach  $t = 10$ .

$\delta$	Runge-Kutta		Adams-Bash	
	error	ratio	error	ratio
1/2	8.1 <sub>-4</sub>		2.0 <sub>-2</sub>	
1/4	1.2 <sub>-4</sub>	6.9	2.3 <sub>-3</sub>	8.7
1/8	9.2 <sub>-6</sub>	12.6	3.0 <sub>-4</sub>	7.5
1/16	6.4 <sub>-7</sub>	14.5	2.4 <sub>-5</sub>	12.7
1/32	4.1 <sub>-8</sub>	15.3	1.7 <sub>-6</sub>	14.5
1/64	2.6 <sub>-9</sub>	15.7	1.1 <sub>-7</sub>	15.3
1/128	1.7 <sub>-10</sub>	15.8	6.9 <sub>-9</sub>	15.7
1/256	1.1 <sub>-11</sub>	15.9	4.4 <sub>-10</sub>	15.8
1/512	6.6 <sub>-13</sub>	15.9	2.7 <sub>-11</sub>	15.9

**Observation:** Both methods of fourth order.

# Summary

Runge-Kutta methods use multiple intermediate derivatives.

$$k_i := f \left( t + \delta c_i, y(t) + \delta \sum_{j=1}^s a_{ij} k_j \right) \quad \text{for all } i \in [1 : s],$$

$$\tilde{y}(t + \delta) := y(t) + \delta \sum_{i=1}^s b_i k_i$$

Example: Classic Runge-Kutta method reaches fourth-order accuracy.

Multistep methods re-use previous states and derivatives.

$$\tilde{y}(t + \delta) = y(t) + \delta \sum_{j=0}^m w_j f(t - \delta j, \tilde{y}(t - \delta j)).$$

Example: Adams-Bashforth methods of any order can be constructed by solving a system of linear equations.