### **NAME**

```
simp — simplistic programming language
```

#### **SYNOPSIS**

```
simp
simp [-i] -e string [arg ...]
simp [-i] -p string [arg ...]
simp [-i] ile [arg ...]
```

#### DESCRIPTION

The **simp** utility is an interpreter for *Simp*, a minimalist (yet expressive), statically scoped, dynamically typed, properly tail-recursive, general-purpose LISP-1 programming language in which imperative, funcional and message-passing programming styles (to name a few) may be conveniently expressed. It reads expressions from a file, string, or standard input, and evaluates them (possibly writing to standard output the result of the evaluation, if in the REPL mode).

The options are as follows:

```
-e string
```

Read expressions from string but do not write the resulting evaluation into standard output.

**-i** Enter the interactive REPL mode after evaluating expressions from a string or file. This flag is set by default if no argument is given.

```
-p string
```

Read expressions from string and write the resulting evaluation into standard output.

In the first synopsis form, expressions are interpreted interactively in a REPL (read-eval-print loop). Expressions are read from standard input (usually a terminal) and written to standard output.

In the second synopsis form (with  $-\mathbf{e}$ ), expressions are read from string and the resulting evaluation is not written into standard output. Unless the flag  $-\mathbf{i}$  is given, no interactive REPL occurs after evaluating the string.

In the third synopsis form (with  $-\mathbf{p}$ ), expressions are read from string and the resulting evaluation is written into standard output. Unless the flag  $-\mathbf{i}$  is given, no interactive REPL occurs after evaluating the string.

In the fourth synopsis form (with an argument and neither  $-\mathbf{e}$  or  $-\mathbf{i}$ ), expressions are read from file and evaluated without outputing the result. If file is an hyphen (-), read expressions from the standard input instead. Unless the flag  $-\mathbf{i}$  is given, no interactive REPL occurs after evaluating the file.

**simp** (like most dialects of Lisp) employs a fully parenthesized prefix notation for programs and other data known as *S-expression* (or "sexp" for short). However, **simp** implements *S-expressions* as **vectors** rather than cons cells (aka pairs). In fact, it is the vector, not the pair, the primary mean of data combination. In **simp**, lists can be implemented as a chain of vectors, in which the last element of each vector is a pointer to the next vector.

### LANGUAGE

A **simp** program is a sequence of expressions, possibly alternated with comments and whitespace. The kind of value that an expression returns is specified below after the  $\Rightarrow$  symbol. An expression can be a special form, a symbol, a literal or a procedure combination. The initial **simp** syntax environment starts out with a few builtin syntactical forms; and its regular environment starts out with a number of variables bound to builtin procedures. They are listed below.

**simp** (like most Lisps) is a dynamically typed programming language. Types are associated with objects rather than with variables. (Statically typed languages, by contrast, associate types with variables and expressions as well with values). Objects are self-typing; the primitive data type of each object is implicit in the object itself. Primitive data types are list below, together with the builtin procedures that manipulate them.

**simp** is a statically scoped language with block structure. To each place where an identifier is bound in a program there corresponds a *region* of the program text within which the binding is visible. Every mention of an identifier refers to the binding of the identifier that established the innermost of the regions containing the use.

Expressions can be evaluated, yielding a value. An expression can evaluate to itself (a so called *self-evaluating* expression). An evaluation can do more than just compute things, such as mutate an object, or perform input/output. Such additional processes are called *side effects*. An expression can also signal an error, in which case the evaluation algorithm is interrupted, even when at a deep level of recursion.

### **Syntax Sugars**

 $\OBJECT \Rightarrow OBJECT$ 

Any object prefixed with a backslash is equivalent to quoting the object. For example,  $\setminus$  () is expanded to the special form (quote ()) while parsing.

#### **Non-vector Forms**

 $SYMBOL \Rightarrow OBJECT$ 

Return the value that the given symbol is bound to in the current environment. SYMBOL must be a symbol already bound in the current environment; otherwise an error is signaled.

 $LITERAL \Rightarrow OBJECT$ 

Return the given literal expression. LITERAL can be any expression, except a symbol or a vector

# **Special Forms**

```
(and EXPRESSION ...) ⇒ OBJECT
```

Evaluate each expression in turn and return false when one evaluate to false; return a true value otherwise.

(define SYMBOL EXPRESSION) ⇒ VOID

Bind the given symbol to the given expression in the current environment.

(defmacro SYMBOL SYMBOL ... EXPRESSION) ⇒ VOID

Bind the given symbol to the given expression in the current syntax environment.

(defmacro... SYMBOL SYMBOL ... SYMBOL EXPRESSION) ⇒ VOID

Same as **defmacro**, but creates a variadic argument macro, in which the last given symbol is bound to the remaining actual unevaluated arguments.

(defun SYMBOL SYMBOL ... EXPRESSION) ⇒ VOID

Bind, in the current environment, the given first symbol to the closure in which the given following symbols are bound to the actual arguments in the given body expression. This is a combination of the **define** and **lambda** syntactical forms.

 $(\mathbf{defun}\dots SYMBL\ SYMBOL\ \dots\ SYMBOL\ EXPRESSION) \Rightarrow VOID$ 

Same as **defun**, but in the closure, the last symbol is bound to a vector containing the remaining actual arguments. This is a combination of the **define** and **lambda...** syntactical forms.

(**do** EXPRESSION ...)  $\Rightarrow$  OBJECT

Evaluate each given expression and return the result of the last evaluation.

(false)

Return the false boolean constant.

(if CONDITION ... THEN ... [ELSE])  $\Rightarrow$  OBJECT

Evaluate each *CONDITION* expression in turn and, if it returns a true value, return the immediatelly following *THEN* expression. If all condition-then pairs have been tested, return the optional *ELSE* condition, if it exists.

(lambda SYMBOL ... EXPRESSION) ⇒ PROCEDURE

Return a closure in which the given symbols are bound to the actual arguments in the given body expression.

(lambda... SYMBOL ... SYMBOL EXPRESSION) ⇒ PROCEDURE

(The form is literally "lambda" three dots). Return a variadic closure in which the given symbols but the last are bound to the actual arguments in the given body expression. The last symbol is bound to a vector containing the remaining actual arguments.

(let [SYMBOL EXPRESSION] ... EXPRESSION)  $\Rightarrow$  OBJECT

Evaluate the final given expression in an environment in which each given symbol is bound to the given expression after it.

(or EXPRESSION ...)  $\Rightarrow OBJECT$ 

Evaluate each expression in turn and return true when one evaluate to true; return a false value otherwise.

(quote EXPRESSION) ⇒ OBJECT

Return the given expression itself, without evaluating it.

(redefine SYMBOL OBJECT) ⇒ VOID

Rebound the given symbol to the given object in the current environment. It is an error if the symbol is not already bound in the current environment.

(true) ⇒ TRUE

Return the true boolean constant.

### Values

There are two means of comparing objects. Two objects are the *same* if they are bound to the same thing or location. Two objects are *equivalent* if the values stored in their locations are equivalent by some predicate.

The notation  $(+26) \Rightarrow 8$  means that the expression (+26) evaluates to the value represented by 8. The notation  $(+26) \equiv (+44)$  means that both expressions (+26) and (+44) evaluate to the same value.

(same? OBJECT ...) ⇒ BOOLEAN

Return whether the given objects are the same.

### **Numbers**

Numbers are self-evaluating objects that represent an integer or real value. **simp** currently only implements fixed-size integer and double-sized floating point data types. The procedures below only apply to integers. The external representation of a number is a number literal.

(+ NUMBER ...) ⇒ NUMBER

Return the sum of the given numbers.

 $(* NUMBER ...) \Rightarrow NUMBER$ 

Return the product of the given numbers.

(- NUMBER NUMBER ...) ⇒ NUMBER

Return the difference of the given numbers.

(/ NUMBER NUMBER ...) ⇒ NUMBER

Return the ratio of the given numbers.

 $(= NUMBER ...) \Rightarrow BOOLEAN$ 

```
Return whether the given numbers are equal.
```

```
(< NUMBER ...) \Rightarrow BOOLEAN
```

Return whether the given numbers are monotonically increasing.

(> NUMBER ...) ⇒ BOOLEAN

Return whether the given numbers are monotonically decreasing.

```
(<= NUMBER ...) ⇒ BOOLEAN
```

Return whether the given numbers are monotonically nondecreasing.

 $(>= NUMBER ...) \Rightarrow BOOLEAN$ 

Return whether the given numbers are monotonically nonincreasing.

```
(abs NUMBER) ⇒ NUMBER
```

Return the absolute, non negative, value of the given number.

```
(number? OBJECT) \Rightarrow BOOLEAN
```

Return whether the given object is a number.

# **Symbols**

Symbols are non-self-evaluating objects holding an interned string of characters. A symbol evaluates to the value bound to its homonymous variable in the current environment. The external representation of a symbol is an identifier.

```
(symbol? OBJECT) ⇒ BOOLEAN
```

Return whether the given object is a symbol object.

#### **Booleans**

Boolean are one of two distinct unique self-evaluating objects: the true and false objects. These objects have no read external representation (therefore they cannot be created by the **read** procedure).

Booleans can be used to control the evaluation of conditional procedures. **simp** interprets the false boolean object as a logical false, and any other object (including the true boolean object itself) as a logical true.

```
(boolean? OBJECT) \Rightarrow BOOLEAN
```

Return whether the given object is a boolean object.

```
(false? OBJECT) ⇒ BOOLEAN
```

Return whether the given object is the false object.

```
(not OBJECT) \Rightarrow BOOLEAN
```

Return the true object if the given object is the false object. Return the false object otherwise.

```
(true? OBJECT) ⇒ BOOLEAN
```

Return whether the given object is not the false object.

# **Bytes**

Bytes are self-evaluating objects representing an integer in the range from 0 to 255 inclusive. The external representation of a byte is a character literal.

```
(byte? OBJECT) ⇒ BOOLEAN
```

Return whether the given object is a byte object.

# **Strings**

Strings (or bytevectors) are self-evaluating objects bound to a homogeneous sequence of zero or more locations in memory, each one holding exactly a byte and indexed from zero. A *byte* is an exact integer in the range from 0 to 255 inclusive. A string is typically more space-efficient than a vector containing the same values. The external representation of a string is a string literal.

```
(empty? OBJECT) ⇒ BOOLEAN
```

Return whether the given object is the string object with zero elements.

```
(string BYTE ...) ⇒ STRING
```

Return a newly allocated string containing the given bytes.

(string-<? STRING STRING ...) ⇒ BOOLEAN

Return whether the given strings are lexicographically sorted in increasing order.

(string→>? STRING STRING ...) ⇒ BOOLEAN

Return whether the given strings are lexicographically sorted in decreasing order.

(string-<=? STRING STRING ...) ⇒ BOOLEAN

Return whether the given strings are lexicographically sorted in nondecreasing order.

(string->=? STRING STRING ...) ⇒ BOOLEAN

Return whether the given strings are lexicographically sorted in nonincreasing order.

 $(\texttt{string-concat} \ \textit{STRING} \ \ldots) \ \Rightarrow \ \textit{STRING}$ 

Return a newly allocated string whose elements form the concatenation of the given strings.

(string-copy! STRING STRING) ⇒ VOID

Copy all elements of the second string into the first one. The length of the second string must be lesser or equal to the first's, or an error will be signaled

(string-clone STRING) ⇒ STRING

Return a newly allocated string with size and elements the same as the given string.

(string-length STRING) ⇒ NUMBER

Return the number of elements of the given string.

(string-alloc NUMBER) ⇒ STRING

Return a newly allocated string of the given size filled with the zero byte.

(string-get STRING NUMBER) ⇒ BYTE

Return the i-th element of the given string.

(string-set! STRING NUMBER BYTE) ⇒ STRING

Set the i-th element of the given string to the given byte.

(string-slice STRING [NUMBER [NUMBER]]) ⇒ STRING

Return the given string indexed from the first given number (default 0) with the second given number (default maximum possible) elements. For example, (string-slice "abcde" 1

3) returns the string "bcd". Both strings point to the same memory, so changing an element in the resulting string also changes the corresponding element in the original string.

(string? OBJECT) ⇒ BOOLEAN

Return whether the given object is a string object.

#### Vectors

Vectors are non-self-evaluating objects bound to a heterogeneous sequence of zero or more locations in memory, each one holding exactly an object of arbitrary type, and indexed from zero. The external representation of a number is a number literal.

A vector with zero element is called a *nil*. A vector with one element is called a *box*. A vector with two elements is called a *pair*. A vector with a number of n elements is called an *n-uple*.

More complex data structures, such as linked lists, hash tables, trees, and records (to name a few), can be implemented in terms of vectors.

(alloc NUMBER) ⇒ VECTOR

Return a newly allocated vector of the given size filled with the null object.

(car VECTOR) ⇒ OBJECT

Return the first element of the given vector. (car v) is equivalent to (get v 0).

 $(cdr \ VECTOR) \Rightarrow VECTOR$ 

Return the given vector indexed from the second element. (cdr v) is equivalent to (slice v 1 (- (length v) 1)).

(clone VECTOR)  $\Rightarrow$  VECTOR

Return a newly allocated vector with size and elements the same as the given vector.

(concat VECTOR ...) ⇒ VECTOR

Return a newly allocated vector whose elements form the concatenation of the given vectors.

(copy! VECTOR VECTOR) ⇒ VOID

Copy all elements of the second vector into the first one. The length of the second vector must be lesser or equal to the first's.

(equiv? VECTOR ...) ⇒ BOOLEAN

Return whether the given vectors have the same length and the same elements.

(get VECTOR NUMBER) ⇒ OBJECT

Return the i-th element of the given vector.

 $(length VECTOR) \Rightarrow NUMBER$ 

Return the number of elements of the given vector.

 $(member PROCEDURE OBJECT VECTOR) \Rightarrow OBJECT$ 

If the given vector contains an object for which the given binary predicate procedure passes for the given object, return the slice of the given vector starting from such object until the end; return the false object otherwise. For example, (member  $> 3 \setminus (0\ 2\ 4\ 6\ 8)$ ) returns the slice (4 6 8), which begins from the first object larger than 3.

(null? OBJECT)  $\Rightarrow$  BOOLEAN

Return whether the given object is the vector object with zero elements.

(reverse VECTOR) ⇒ VECTOR

Return a newly allocated vector with the same size and elements as the given vector, but in reverse order.

(reverse! VECTOR) ⇒ VECTOR

Reverse the elements of the given vector in place, and return it.

(set! VECTOR NUMBER OBJECT) ⇒ VECTOR

Set the i-th element of the given vector to the given object.

(slice VECTOR [NUMBER [NUMBER]]) ⇒ VECTOR

Return the given vector indexed from the first given number (default 0) with the second given number (default maximum possible) elements. For example, (slice (vector 'a' 'b' 'c' 'd' 'e') 1 3) returns the vector ('b' 'c' 'd'). Both vectors point to the same memory, so changing an element in the resulting vector also changes the corresponding element in the original vector.

(string->vector STRING) ⇒ VECTOR

Return a newly allocated vector filled with the bytes of the given string.

(vector OBJECT ...) ⇒ VECTOR

Return a newly allocated vector containing the given objects.

(vector? OBJECT) ⇒ BOOLEAN

Return whether the given object is a vector object.

### **Procedures**

Procedures are self-evaluating objects that represent either a builtin procedure or a closure created with the **lambda** syntactical form.

```
(apply PROCEDURE OBJECT ... VECTOR) ⇒ OBJECT
```

Return the result of the application of the given procedure to all the given objects and the objects in the given vector.

(for-each PROCEDURE VECTOR VECTOR  $\dots$ )  $\Rightarrow$  VOID

Apply the given procedure element-wise to the elements of the given vectors, in order. It is an error if the given procedure does not accept as many arguments as there are vectors. The procedure **for-each** is like **map** but calls procedure for its side effects rather than for its value.

(map PROCEDURE VECTOR VECTOR ...) ⇒ VECTOR

Return the vector of results of applications of the given procedure element-wise to the elements of the given vectors, in order. It is an error if the given procedure does not accept as many argu-

ments as there are vectors.

```
(string-for-each PROCEDURE STRING STRING ...) ⇒ VOID
```

Apply the given procedure element-wise to the bytes of the given strings, in order. It is an error if the given procedure does not accept as many arguments as there are strings. The procedure **string-for-each** is like **string-map** but calls procedure for its side effects rather than for its value.

```
(string-map PROCEDURE STRING STRING ...) ⇒ STRING
```

Return the string of results of applications of the given procedure element-wise to the bytes of the given string, in order. It is an error if the given procedure does not accept as many arguments as there are strings, or if the procedure does not return a byte.

```
(procedure? OBJECT) ⇒ BOOLEAN
```

Return whether the given object is a procedure object.

#### **Ports**

Ports are self-evaluating objects representing input or output devices. A port object has no read external representation. The printed external representation of a port is unique for each port object, but unpredictable.

Ports can be input or output ports; and can be used to read/write bytes from/into files or strings. Ports can be closed. When a port is closed, no further input/output operation is permited on that port. Input/output operation can be buffered, and closing a port flushes the buffer.

```
Operations on ports are listed below.
```

```
(port? OBJECT) ⇒ BOOLEAN
```

Return whether the given object is a port object.

 $(stderr) \Rightarrow PORT$ 

Return the standard error port.

 $(stdin) \Rightarrow PORT$ 

Return the standard input port.

(stdout) ⇒ PORT

Return the standard output port.

# Input/Output

The *read external representation* of an object is a representation of an object as a sequence of characters. This representation is not necessarily unique; and certain objects may have no read external representation at all. For example, the integer 28 has the read external representations "28", "+28", and "0x1C", among others. Closures and port objects, for example, have no read external representation. The read external representation of an object is parsed by the **read** procedure.

The *printed external representation* of an object is a representation of an object as a sequence of characters. This representation is unique; and all objects have one. For example, the printed external representation of the integer 28 is the sequence of characters "28", and the printed external representation of a vector consisting of the integers 8 and 13 is the sequence of characters "(8 13)". For most data types, an object's printed external representation is also a valid read external representation for the object. The printed external representation of an object is generated by the **write** procedure.

The *pretty-printed external representation* of an object is a representation of an object as a sequence of characters. This representation is not necessarily unique; and all objects have one. The pretty-printed external representation is similar to the regular printed external representation, except that it is not guaranteed that a pretty-printed external representation for an object is a valid read external representation. In other words, while the printed external representation of an object can be read by the interpreter, the pretty-printed external representation of an object is made for user consumption. The pretty-printed

external representation of an object is generated by the display procedure.

(display OBJECT [PORT]) ⇒ VOID

Write the given object into the given port (standard output, by default) in user-readable form.

(eof? OBJECT) ⇒ BOOLEAN

Return whether the given object is the end-of-file object.

 $(newline [PORT]) \Rightarrow VOID$ 

Write a newline into the given port (standard output, by default).

(read [PORT]) ⇒ OBJECT

Read an object from the given port in its read external representation. It returns false on read error and the end-of-object file when the port reached a end.

(write OBJECT [PORT]) ⇒ VOID

Write the given object into the given port (standard output, by default) in its printed external representation.

# **Environments**

Evaluation is subject to the current environment (or the environment given to the eval special form).

An identifier may name a location in memory where a value can be stored. Such identifier is called a *variable* and is said to be *bound* to that location. The set of all visible bindings in effect at some point in a program is known as the *environment* in effect at that point. The value stored in the location to which a variable is bound is called the variable's *value*. By abuse of terminology, the variable is sometimes said to name the value or to be bound to the value. This is not quite accurate, but confusion rarely results from this practice.

An environment is a sequence of *frames*. Each frame is a table (possible empty) of *bindings*, which associate variable names with their corresponding location in memory. Frames are structured in a singly linked list: each frame can point to either nothing or to a parent tree. The *value of a variable* with respect to an environment is the value on the location given by the binding of the variable in the first frame in the environment that contains a binding for that variable. If no frame in the sequence specifies a binding for the variable, then the variable is said to be *unbound* in the environment.

(environment [ENVIRONMENT]) ⇒ ENVIRONMENT

Return a newly allocated environment using the given environment as parent environment. If no environment is given, return the empty environment.

(environment? OBJECT) ⇒ BOOLEAN

Return whether the given object is an environment object.

(eval OBJECT ENVIRONMENT) ⇒ OBJECT

Evaluate the given object in the given environment and return the result of the evaluation.

#### FORMAL SYNTAX

This section provides a formal syntax for **simp** written in an extended Backus-Naur form (BNF). Non-terminals are written between angle braces (<...>). A terminal symbol is written between double quotation marks ("...").

The following extensions to BNF are used to make the description more consise:

- <thing>\* means zero or more occurrences of <thing>.
- <thing>+ means one or more occurrences of <thing>.
- <thing>? means zero or one occurrence of <thing>.

The BNF is augmented with the concepts of character classes and character ranges. A character class is expressed between square braces and colons ([:...:]) and denotes a named set of characters. A character range is a set of characters and/or character classes between square braces ([...]) and denotes any character in the set or in the classes. For example, ([abc[:delimiter:]]) means an

**a**, or **b**, or **c** character, or a character in the [:delimiter:] class. The notion of character range is augmented as follows.

- The character has the same special meaning in a character range it has in ERE. For example, [0-9] is the same as [0123456789] (which is the same as [:decimal:]).
- The ^ character has the same special meaning in a character range it has in ERE. For example, [^abc] means any character but a, b, or c.
- The opening bracket [ may occur anywhere in a character range.
- The closing bracket ] may occur only as the first character in a character range

Unprintable and hard-to-type characters are represented in the same escape notation used in string literals. For example,  $\n$  is the newline.

# **Alphabet**

The alphabet for this grammar is all the 256 bytes that can be read from a file augmented with the end-of-file indicator.

The character classes are defined as follows.

```
[:space:] ← [\f\n\r\t\v]
[:binary:] ← [0-1]
[:octal:] ← [0-7]
[:decimal:] ← [0-9]
[:hexadecimal:] ← [0-9A-Fa-f]
[:delimiter:] ← [][()#[:eof:][:space:]]
```

The end-of-file indicator, in special, is represented by the special class [:eof:].

The backslash character (\), the double-quote character ("), and the single-quote character ('), which have special meanings and thus would need to be escaped, are represented by the special classes [:slash:], [:double-quote:], and [:single-quote:], respectively.

The special character class [:anything:] represents any character in the alphabet.

# Tokens

A **token** is the lexical element used to compose well formed expressions. Some characters, known as **delimiters**, have special meaning during the program parsing, because certain tokens require a delimiter to occur after them. A token is defined as follows:

The end-of-file is the token that terminates a program. It is actually not a character, but is interpreted as if it were.

```
<end-of-file> ← [:eof:]
```

Single-character tokens are the following:

```
<left-paren> ← "("
<right-paren> ← ")"
```

A character literal is composed by one character element between single quotes. A string literal is composed by zero or more character elements between double quotes. A character element is any character other than a double quote or a backslash or an escaped character. Chracter literals and string literals are used to represent characters and strings (alsow known as bytevectors) respectively. The single and the double quotation mark characters that terminates a character and a string are themselves delimiters.

A number literal begins with an optional signal and is followed by the number body. A delimiter must occur after a numer literal.

An identifier is any sequence of non-delimiter characters that does not form another type of token. A delimiter character must occur after an identifier.

# Escape sequences

Within a string literal, sequences of characters beginning with a backslash ( $\$ ) are called **escape sequences** and represent bytes other than the characters themselves. Most escape sequences represent a single byte, but some forms may represent more than one byte. An invalid escape sequence is equivalent to the character after the backslash; for example, the string literal " $\$ " does not contain a valid escape sequence, so it is equivalent to " $\$ ". The valid escape sequences are as follows:

```
\a
        Alarm (U+0007).
        Backspace (U+0008).
/b
\t.
        Horizontal tab (U+0009).
\n
        Line feed (U+000A).
\v
        Vertical tab (U+000B).
\f
        Form feed (U+000C).
        Carriage return (U+000D).
\r
\e
        Escape character (U+001B).
\"
        Double quote (U+0022).
```

```
Name Byte whose value is the 1-, 2-, or 3-digit octal number num.

Name Byte whose value is the 1- or 2-digit hexadecimal number num.

Name Bytes encoding, in UTF-8, the 4-digit hexadecimal number num.

Name Bytes encoding, in UTF-8, the 8-digit hexadecimal number num.
```

# Intertoken space

Tokens are separated by intertoken space, which includes both whitespace and comments. Intertoken space is used for improved readability, and as necessary to separate tokens from each other.

Whitespace can occur between any two tokens, but not within a token. Whitespace occurring inside a string literal is significant.

Comments are annotations in the source code and are treated exactly like whitespace. A hash character (#) outside a string literal indicates the start of a comment. The comment continues to the end of the line on which the hash character appears.

# Read external representation

The following is a simplification of the syntax of a read external representation. This syntax is not complete, because intertoken-space may occur on either side of any token (but not within a token).

### **Program**

A **simp** program is a sequence of characters forming whitespace, comments, and tokens. The tokens in a program must form syntactically well formed expressions.

### FORMAL SEMANTICS

I have no idea what a formal semantics is or does.

#### **EXAMPLES**

[TODO]

#### **SEE ALSO**

Harold Abelson, Gerald Jay Sussman, and Julie Sussman, *Structure and Interpretation of Computer Programms*, *The MIT Press*, 1996.

#### **STANDARDS**

The **simp** programming language is compliant with nothing, as it has not been standardised. It was influenced by the Scheme LISP dialect.

The syntax for comments and number literals breaks the usual LISP tradition, and are influenced by shell script comments and C constants, respectively.

Parts of this manual (especially at the **DESCRIPTION** section) were blatantly stolen from *Revised Report on the Algorithmic Language Scheme*.

#### HISTORY

The **simp** programming language was developed as a personal playground for programming language theory, motivated by the reading of the Wizard Book (Abelson & Sussman). It first appeared as a C implementation in 2022.

#### **AUTHORS**

The **simp** programming language was designed by Lucas de Sena <lucas AT seninha DOT org>.

# **BUGS**

The **simp** programming language implemented in **simp** is not complete, and may not conform to this manpage.

This manual page is also not complete, as the language is only informally specified, and may change significantly from one release to the other.

This manual uses the terms "string" and "bytevector" interchangeably, as both refer to the same **simp** data structure. Note that "string" and "string literal" refer to different concepts; the former is a data type, while the latter is a token type.

This manual avoids to use the word "character" to refer to the elements of a string. This manual uses the word "character" to refer solely to the units that compose tokens read by the parser. Strings in **simp** can possibly contain no valid character (in the sense of a UTF-8 encoded codepoint). This manual uses the term "byte" instead to refer to the elements of a string.

There's no "character" data type, either in the C sense of a "byte", or in the sense of a UTF-8 encoded codepoint. A single byte can be represented as a one-element string. A UTF-8 encoded codepoint can be represented as a string containing the encoding bytes.