Pyro: A Spatial-Temporal Big-Data Storage System





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Applications

- A huge amount of geo-tagged events are generated and stored in real-time.
 - Tweets, Photos
 - Taxi locations
 - Smartphone User Traces
- Query ask for events within a given time range and geographic area: geometry query.

Challenges

- Efficiently store and retrieve Spatial-temporal data
- Achieve Scalability
- Handle dynamic workload hotspot

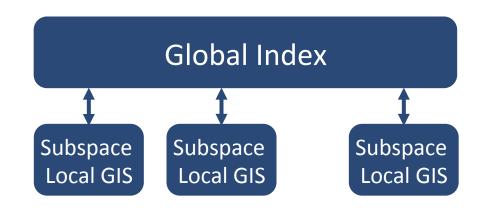


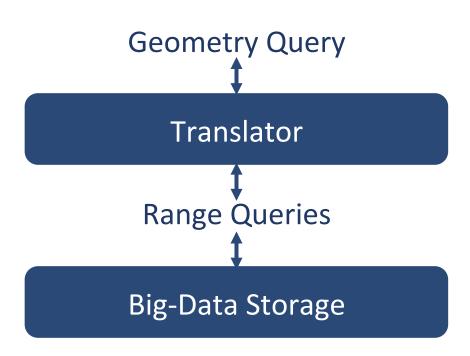
Prior Approaches

- Make Geographic Information Systems (GIS) scalable
- Make Big-Data storage system understand spatial-temporal workload

Contributions

- Pyro is the first holistic solution specifically designed for Spatial-Temporal Applications.
 - Internally understands
 Spatial-Temporal data and query
 - Aggregatively optimizes IO
 - Manages data replicas to mitigate workload hotspots





Background

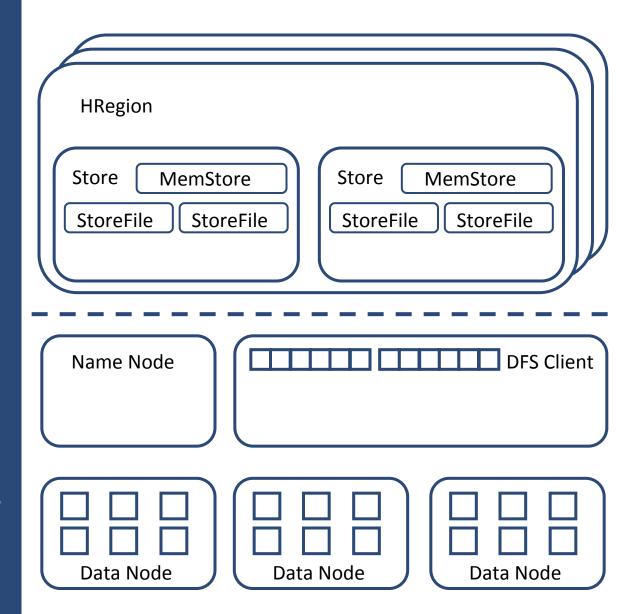
• Hbase

- The table is horizontally divided into HRegions.
- Each HRegion is vertically divided into stores, one store per column family.
- Data is first cached in the MemStore, and then flushed into a StoreFile when the size threshold is reached.

HDFS

- The Name Node manages file system namespaces.
- Data Nodes store data chunks
- DFS Client exposes APIs.

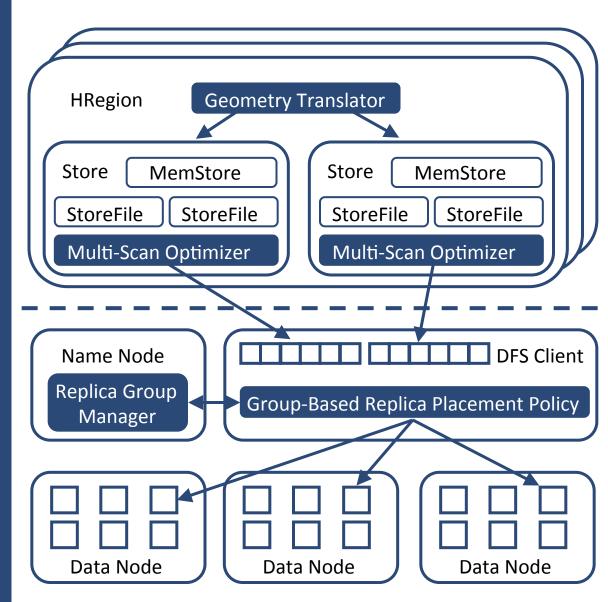
HBase



- Geometry Translator
 - Encoding spatial-temporal information into row keys, and translating geometry queries into range scans

- Multi-Scan Optimizer
 - Aggregatively optimizing all range scans of the same geometry query

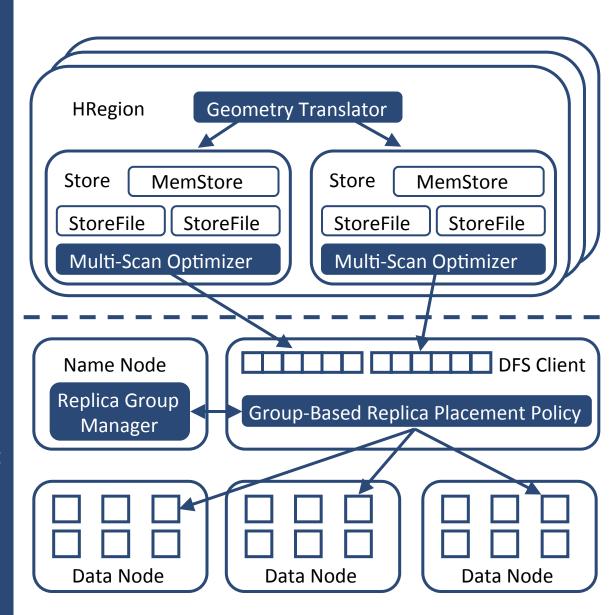
- Group-Based Replica Placement
 - Improves data locality during workload dynamics.



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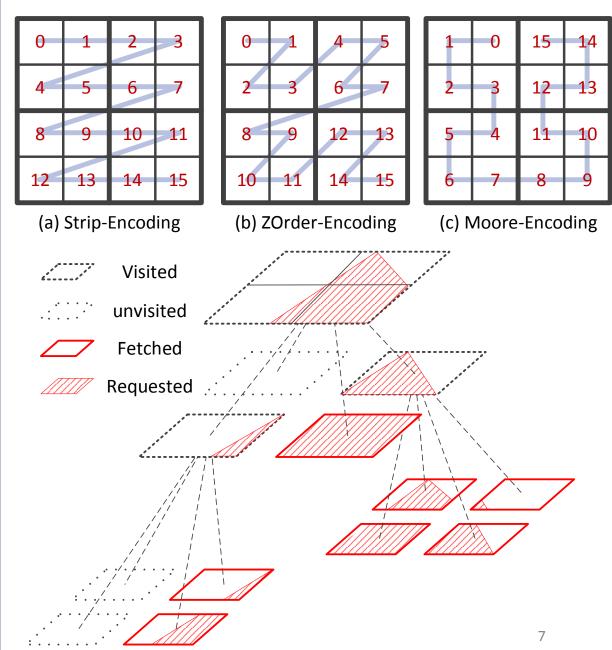
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Geometry Translator

- The space is recursively divided into tiles using a quad-tree
- Using a space filling curve (Z, Moore, Hilbert, etc.) to encode tiles

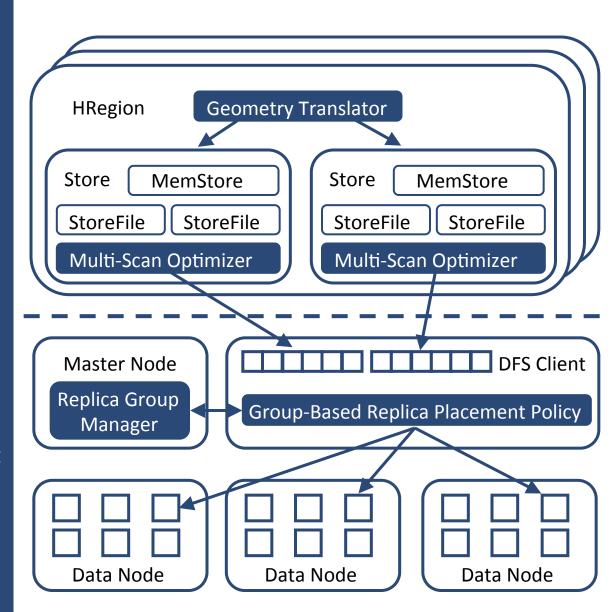
- Use the same quad-tree to calculate the tiles that intersect with the geometry
- Tiles then turns into range scans.



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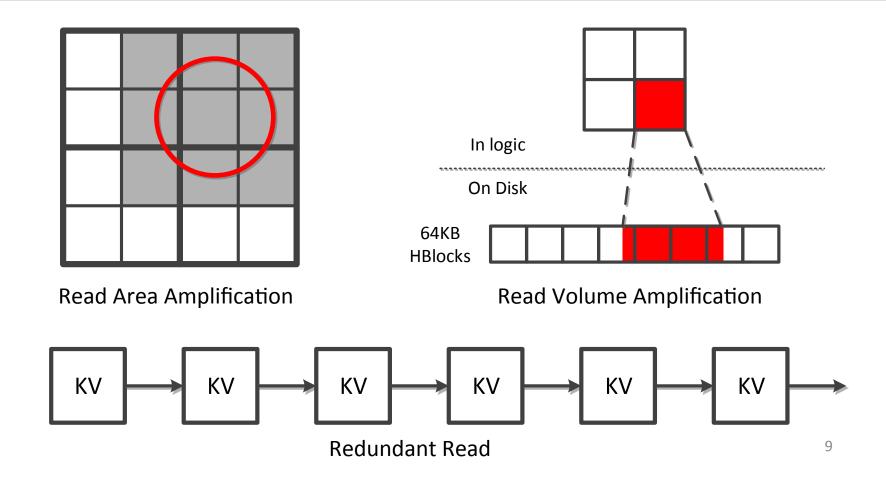
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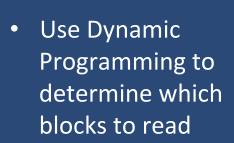
Multi-Scan Optimizer: Read Amplification

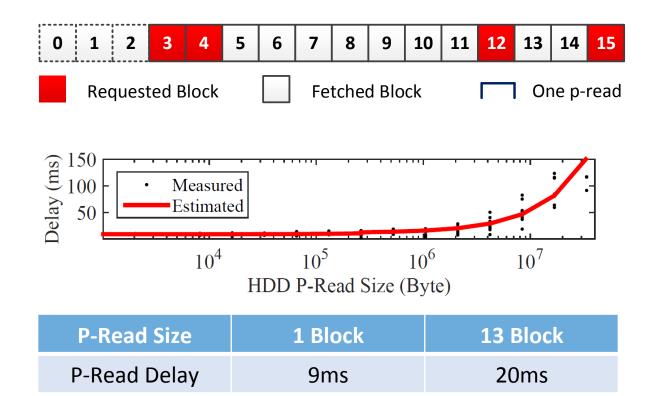
- A Geometry query may translate into a large number of range scans.
- These range scans usually force the underlying system to fetch more data or repeatedly go through the same data structure.



Multi-Scan Optimizer: Use Small Tile and HBlocks

- Keep tile size and block size small, and aggregatively optimize range scans.
- Profile P-Read delay vs size.





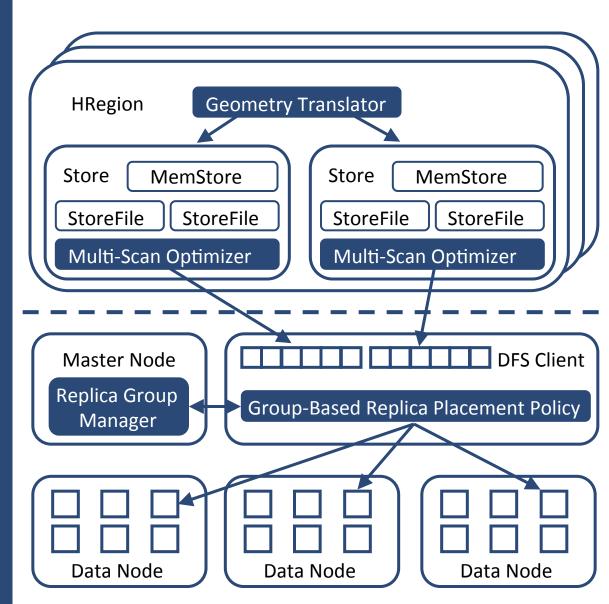
Adaptive Aggregation Algorithm:

$$S[i] = \min\{S[j-1] + E(j,i) | 1 \le j \le i\}$$

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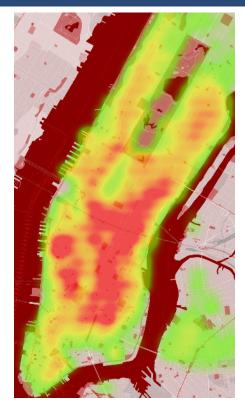
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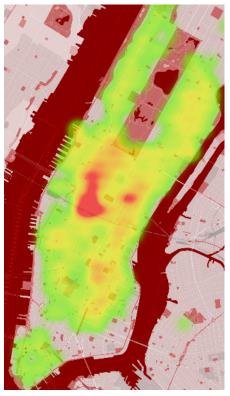


Group-Based Replica Placement

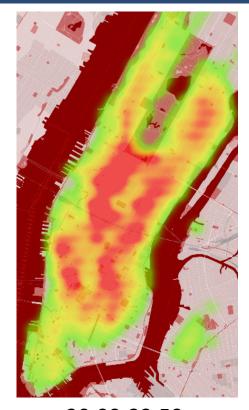
- Each HRegion handles a range of row keys, that corresponds to a subarea in the space.
- Spatial-temporal applications naturally create dynamic workload hotspots within small areas that may overwhelm corresponding HRegion servers.



20:00-23:59 Dec 31, 2012



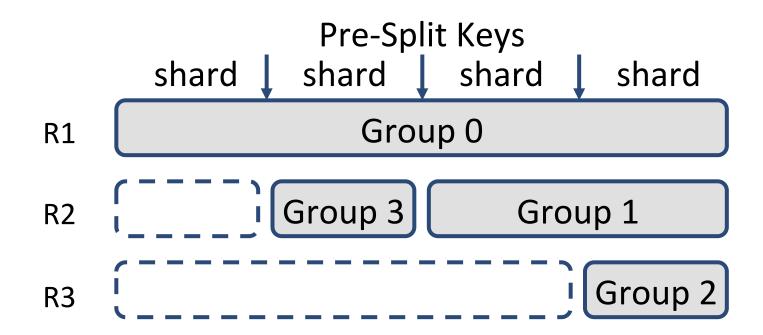
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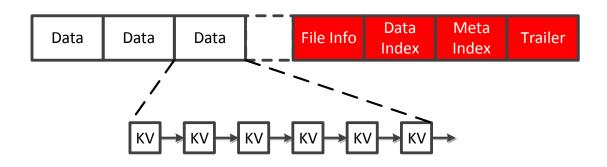
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Group-Based Replica Placement Policy

- A HRegion can split to input multiple daughter HRegions, and these daughter HRegions can be moved into other machines to mitigate workload hotspot.
- HRegions usually co-locate with HDFS datanodes that allows read/write data locality. Splitting may destroy data locality.
- Pyro employs group-based replica placement to achieve data locality.

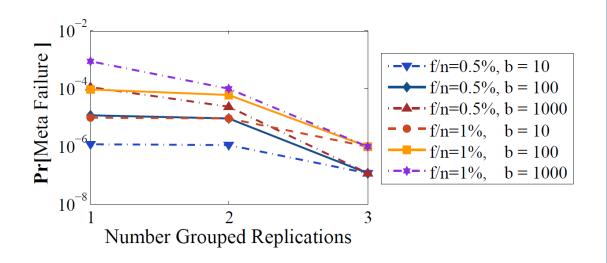


Group-Based Replica Placement Asymmetry



n: # of servers, f: # of failed servers,

g: # of groups, b: # of DFS blocks in the file



The asymmetry in replica groups caters HFile format: meta data locates at the end of the Hfile.

Meta blocks: minimize the probability of losing any DFS block

Data blocks: minimize the expectation of the number of unavailable DFS blocks.

Evaluation

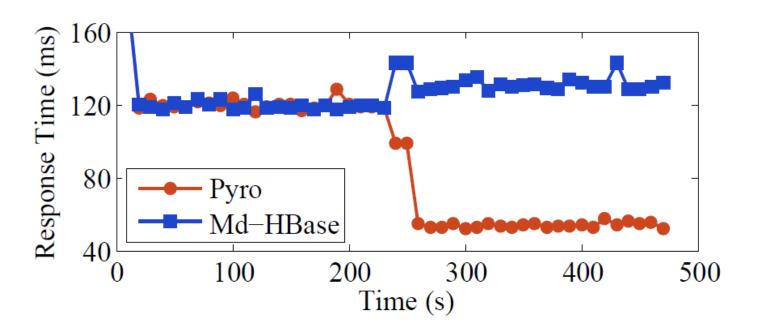
- Open data: ~700,000,000 NYC taxi trips from 2010 to 2013.
 - https://publish.illinois.edu/dbwork/open-data/

- Experimenting on an 80-server cluster:
 - 1 PyroDFS namenode, 30 datanodes
 - 1 PyroDB master, 3 ZooKeeper nodes, 30 co-located HRegion servers.
 - Remaining nodes generate workload and log latency.

- Compare with Md-HBase
 - Md-HBase adds an translation layer above Hbase, and uses Z-order encoding.

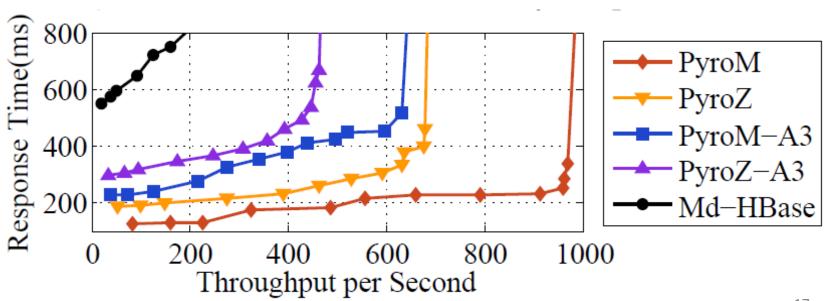
Evaluation

- Manually splitting a Pyro region vs Manually splitting a Md-HBase region.
 - To make the evaluation fair, this evaluation submits range scans rather than geometry query into two systems. In this case, both geometry translator and multi-scan optimizer in Pyro are disabled.
 - Both systems use Z-order encoding algorithm



Evaluation

- Throughput measurement of 100m X 100m rectangle geometry.
 - PyroM: Pyro using Moore encoding
 - PyroZ: Pyro using Zorder encoding
 - PyroM A3: PyroM, disabled adaptive aggregation algorithm
 - PyroZ A3: PyroZ, disabled adaptive aggregation algorithm



Thank you Q&A