$\S 1$ HULLTR INTRODUCTION 1

1. Introduction. This is a hastily written implementation of treehull, using treaps to guarantee good average access time.

```
format Graph int
                           /* gb_graph defines the Graph type and a few others */
  format Vertex int
  format Arc int
  format Area int
#include "gb_graph.h"
#include "gb_miles.h"
#include "gb_rand.h"
  \langle Type declarations _3\rangle
  int n = 128;
  (Global variables 2)
  (Procedures 9)
  main()
  {
    ⟨Local variables 7⟩
    Graph *g = miles(128, 0, 0, 0, 0, 0, 0);
    mems = ccs = 0;
    \langle Find convex hull of g 10\rangle;
    printf("Total_of_wd_mems_and_wd_calls_on_ccw.\n", mems, ccs);
  }
     I'm instrumenting this in a simple way.
\#define o mems ++
#define oo mems += 2
\#define ooo mems +=3
\langle \text{Global variables 2} \rangle \equiv
  int mems;
                /* memory accesses */
  int ccs;
              /* calls on ccw */
See also section 5.
```

This code is used in section 1.

2 Data structures hulltr §3

3. Data structures. For now, each vertex is represented by two coordinates stored in the utility fields x.I and y.I. I'm also putting a serial number into z.I, so that I can check whether different algorithms generate identical hulls.

We use separate nodes for the current convex hull. These nodes have a bunch of fields: p-vert points to the vertex; p-succ and p-pred point to next and previous nodes in a circular list; p-left and p-right point to left and right children in a tree that's superimposed on the list; p-parent is present too, it points to the parent node; p-prio is the priority if we are implementing the tree as a treap.

The *head* node has the root of the tree in its *right* field, and it represents the special vertex that isn't in the tree.

```
\langle \text{Type declarations 3} \rangle \equiv
  typedef struct node_struct {
     struct vertex_struct *vert;
     struct node_struct *succ, *pred, *left, *right, *parent;
     long prio;
  } node;
This code is used in section 1.
     \langle \text{Initialize the array of nodes 4} \rangle \equiv
  head = (\mathbf{node} *) \ gb\_alloc((g\neg n) * \mathbf{sizeof}(\mathbf{node}), working\_storage);
  if (head \equiv \Lambda) return (1);
                                      /* fixthis */
  next\_node = head;
This code is used in section 6.
      \langle \text{Global variables } 2 \rangle + \equiv
  \mathbf{node} *head:
                       /* beginning of the hull data structure */
  node *next\_node;
                             /* first unused slot in that array */
  Area working_storage;
                             /* used to disambiguate entries with equal coordinates */
  int serial\_no = 1;
```

6. We assume that the vertices have been given to us in a GraphBase-type graph. The algorithm begins with a trivial hull that contains only the first two vertices.

```
\langle Initialize the data structures _{6}\rangle \equiv
   init_area(working_storage);
   (Initialize the array of nodes 4);
   o, u = g \rightarrow vertices;
   v = u + 1;
   u \rightarrow z.I = 0;
   v \rightarrow z.I = 1;
   p = ++ next\_node;
   ooo, head \neg succ = head \neg pred = head \neg right = p;
   oo, p \neg succ = p \neg pred = head;
   o, p \neg parent = head;
   oo, p \neg left = p \neg right = \Lambda;
   gb\_init\_rand(110);
   o, p \rightarrow prio = gb\_next\_rand();
   o, head \neg vert = u;
   o, p \neg vert = v;
   next\_node ++;
   if (n < 150) printf("Beginning_with_(%s; %s)\n", u \rightarrow name, v \rightarrow name);
This code is used in section 10.
```

 $\S 7$ HULLTR DATA STRUCTURES 3

We'll probably need a bunch of local variables to do elementary operations on data structures. $\langle \text{Local variables } 7 \rangle \equiv$ Vertex *u, *v, *vv, *w; **node** *p, *pp, *q, *qq, *qqq, *r, *rr, *s, *ss, *tt, **par, **ppar, *prepar; int replaced; /* will be nonzero if we've just replaced a hull element */ This code is used in section 1. Here's a routine I used when debugging (in fact I should have written it sooner than I did). \langle Verify the integrity of the data structures $\rangle \equiv$ p = head;count = 0;**do** { count ++: $p \rightarrow prio = (p \rightarrow prio \& \#ffff0000) + count;$ if $(p \rightarrow succ \rightarrow pred \neq p)$ printf("succ/pred_failure_at_%s!\n", $p \rightarrow vert \rightarrow name$); if $(p \rightarrow left \neq \Lambda \land p \rightarrow left \rightarrow parent \neq p)$ printf ("parent/lchild_failure_lat_\%s!\n", $p \rightarrow vert \rightarrow name$); if $(p \rightarrow right \neq \Lambda \land p \rightarrow right \rightarrow parent \neq p)$ $printf("parent/rchild_failure_lat_\%s!\n", p \rightarrow vert \rightarrow name);$ $p = p \rightarrow succ;$ } while $(p \neq head)$; count = 1; $inorder(head \neg right);$ This code is used in section 10. 9. $\langle \text{Procedures 9} \rangle \equiv$ int count; inorder(p)node *p; **if** (p) { $inorder(p \rightarrow left);$ if $((p\rightarrow prio \& \#ffff) \neq ++count)$ { $printf("tree_lnode_l%d_lis_lmissing_lat_l%d:_l%s!\n", count, p-prio \& #ffff, p-vert-name);$ $count = p \rightarrow prio \& \#ffff;$ $inorder(p \rightarrow right);$ } See also sections 14, 16, and 19. This code is used in section 1.

4 HULL UPDATING HULLTR §10

10. Hull updating. The main loop of the algorithm updates the data structure incrementally by adding one new vertex at a time. If the new vertex lies outside the current convex hull, we put it into the cycle and possibly delete some vertices that were previously part of the hull.

```
 \langle \text{ Find convex hull of } g \mid 10 \rangle \equiv \\ \langle \text{ Initialize the data structures } 6 \rangle; \\ \text{ for } (oo, vv = g \neg vertices + 2; vv < g \neg vertices + g \neg n; vv ++) \ \{ \\ vv \neg z.I = ++ serial\_no; \\ o, q = head \neg pred; \\ replaced = 0; \\ o, u = head \neg vert; \\ \text{ if } (o, ccw(vv, u, q \neg vert)) \ \langle \text{ Do Case 1 12} \rangle \\ \text{ else } \langle \text{ Do Case 2 17} \rangle; \\ \langle \text{ Verify the integrity of the data structures 8} \rangle; \\ \} \\ \langle \text{ Print the convex hull 11} \rangle; \\ \text{This code is used in section 1.}
```

11. Let me do the easy part first, since it's bedtime and I can worry about the rest tomorrow.

```
 \begin{split} &\langle \operatorname{Print \ the \ convex \ hull \ 11} \rangle \equiv \\ &p = head; \\ &printf ("\operatorname{The_{\sqcup}convex\_hull_{\sqcup}is:\n"}); \\ &\mathbf{do} \ \{ \\ &printf ("_{\sqcup \sqcup} \% \mathsf{s} \backslash \mathsf{n"}, p \text{-}vert \text{-}name); \\ &p = p \text{-}succ; \\ &\} \ \mathbf{while} \ (p \neq head); \end{split}
```

This code is used in section 10.

 $\S12$ HULLTR HULL UPDATING \S

12. In Case 1 we don't need the tree structure since we've already found that the new vertex is outside the hull at the tree root position.

```
\langle \text{ Do Case 1 } 12 \rangle \equiv
   \{ qqq = head;
      while (1) {
         o, r = qqq \rightarrow succ;
         if (r \equiv q) break;
                                       /* can't eliminate any more */
         if (oo, ccw(vv, qqq \rightarrow vert, r \rightarrow vert)) break;
         \langle \text{ Delete or replace } qqq \text{ from the hull } 15 \rangle;
         qqq = r;
      qq = qqq;
      qqq = q;
      while (1) {
         o, r = qqq \rightarrow pred;
         if (r \equiv qq) break;
         if (oo, ccw(vv, r \rightarrow vert, qqq \rightarrow vert)) break;
         \langle \text{ Delete or replace } qqq \text{ from the hull } 15 \rangle;
         qqq = r;
      q = qqq;
     if (\neg replaced) (Insert vv at the right of the tree 13);
      if (n < 150) printf("New_hull_usequence_u(%s;_u%s;_u%s)\n", q-vert-name, vv-name, qq-vert-name);
   }
This code is used in section 10.
         At this point q \equiv head \neg pred is the tree's rightmost node.
\langle \text{Insert } vv \text{ at the right of the tree } 13 \rangle \equiv
      tt = next\_node ++;
      o, tt \rightarrow vert = vv;
      o, tt \neg succ = head;
      o, tt \neg pred = q;
      o, head \neg pred = tt;
      o, q \rightarrow succ = tt;
      oo, tt \neg left = tt \neg right = \Lambda;
      o, tt \neg prio = gb\_next\_rand();
      if (n < 150) \ printf("(Inserting_\%s_\at_\right_\of_\text{tree}, \prio=\%d)\n", vv\undername, tt\underprio);
      if (o, tt \rightarrow prio < q \rightarrow prio) rotup(q, \&(q \rightarrow right), tt, tt \rightarrow prio);
      else {
                 /* easy case, no rotation necessary */
         o, tt \neg parent = q;
         o, q \rightarrow right = tt;
      }
   }
This code is used in section 12.
```

6 HULL UPDATING HULLTR §14

14. The link from parent to child hasn't been set when the priorities indicate necessary rotation.

```
\langle \text{Procedures } 9 \rangle + \equiv
   rotup(p, pp, q, qp)
        node *p;
                         /* parent of inserted node */
                             /* link field in parent */
        node **pp;
                         /* inserted node */
        node *q;
        long qp;
                         /* its priority */
   { \mathbf{node} *pr, **ppr;}
                               /* grandparent */
     \mathbf{node} *qq;
                        /* child who is reparented */
     while (1) {
        o, pr = p \rightarrow parent;
        if (o, pr \neg right \equiv p) ppr = \&(pr \neg right);
        else ppr = \&(pr \neg left);
        if (pp \equiv \&(p \rightarrow right)) {
                                         /* we should rotate left */
          if (n < 150) printf("...(rotating_left)\n");
           o, qq = q \rightarrow left;
           o, q \rightarrow left = p;
           o, p \rightarrow parent = q;
           o, p \rightarrow right = qq;
           if (qq \neq \Lambda) o, qq \neg parent = p;
        }
        else { /* we should rotate right */
           if (n < 150) printf("...(rotating_right)\n");
           o, qq = q \rightarrow right;
           o, q \rightarrow right = p;
           o, p \neg parent = q;
           o, p \rightarrow left = qq;
           if (qq \neq \Lambda) o, qq \neg parent = p;
        if (o, qp \ge pr \neg prio) break;
        p = pr;
        pp = ppr;
     o, q \neg parent = pr;
     o, *ppr = q;
```

§15 HULLTR HULL UPDATING 7

15. Nodes don't need to be recycled.

```
\langle Delete or replace qqq from the hull 15\rangle \equiv
  if (replaced) {
      o, pp = qqq \neg pred;
      o, tt = qqq \neg succ;
      o, pp \neg succ = tt;
      o,\, tt \neg pred \,=\, pp\,;
      o, prepar = qqq \neg parent;
      if (o, prepar \rightarrow right \equiv qqq) par = \&(prepar \rightarrow right);
      else par = \&(prepar \rightarrow left);
      o, pp = qqq \rightarrow left;
      if (o, (ss = qqq \neg right) \equiv \Lambda) {
        if (n < 150) printf("(Deleting_\'\'s_\from_\tree,_\case_1)\n", qqq\(\frac{1}{2}vert\);
        o,*par = pp;
        if (pp \neq \Lambda) o, pp \neg parent = prepar;
      else if (pp \equiv \Lambda) {
        if (n < 150) printf("(Deleting_\%s_\from_\text{tree},_\case_\2)\n", <math>qqq \rightarrow vert \rightarrow name);
        o,*par = ss;
        o, ss \neg parent = prepar;
      else {
        if (n < 150) printf("(Deleting_\%s\_from\_tree,\_hard\_case)\n\", qqq^\neg vert^\neg name);
         oo, deldown(prepar, par, pp, ss, pp \neg prio, ss \neg prio);
   else {
      o, qqq \rightarrow vert = vv;
      replaced = 1;
This code is used in sections 12 and 17.
```

8 HULL UPDATING HULLTR $\S16$

```
\langle \text{Procedures } 9 \rangle + \equiv
16.
  deldown(p, pp, ql, qr, qlp, qrp)
       node *p; /* parent of deleted node */
                        /* link field in that parent */
       node **pp;
       node *ql, *qr; /* children of deleted node */
       int qlp, qrp; /* their priorities */
  { node *qq; /* grandchild of deleted node */
     if (qlp < qrp) {
       if (n < 150) printf("...(moving_left_lchild_up)\n");
       o, ql \rightarrow parent = p;
       o, *pp = ql;
       o,\,qq\,=\,ql {\rightarrow} right;
                                                                                  /* tail recursion */
       if (qq \neq \Lambda) o, deldown(ql, \&(ql \neg right), qq, qr, qq \neg prio, qrp);
       else {
          o, ql \rightarrow right = qr;
          o, qr \rightarrow parent = ql;
       }
     else {
       if (n < 150) printf("...(moving_right_child_up)\n");
       o, qr \rightarrow parent = p;
       o, *pp = qr;
       o, qq = qr \rightarrow left;
       if (qq \neq \Lambda) o, deldown(qr, \&(qr \rightarrow left), ql, qq, qlp, qq \rightarrow prio); /* tail recursion */
       else {
          o, qr \rightarrow left = ql;
          o, ql \neg parent = qr;
  }
```

§17 HULLTR

```
\langle \text{ Do Case 2 } 17 \rangle \equiv
17.
  \{ o, qq = head \neg right; \}
      while (1) {
        if (qq \equiv q \lor (o, ccw(u, vv, qq \rightarrow vert))) {
           o, r = qq \rightarrow left;
           if (r \equiv \Lambda) {
              preppar = qq;
              o, ppar = \&(qq \rightarrow left);
              break;
         }
        else {
           o, r = qq \neg right;
           if (r \equiv \Lambda) {
              preppar = qq;
              o, ppar = \&(qq \neg right);
              o, qq = qq \neg succ;
              break;
         qq = r;
     if (o, (r = qq \neg pred) \equiv head \lor (oo, ccw(vv, qq \neg vert, r \neg vert))) {
        if (r \neq head) {
           while (1) {
               qqq = r;
              o, r = qqq \rightarrow pred;
              if (r \equiv head) break;
              if (oo, ccw(vv, r \rightarrow vert, qqq \rightarrow vert)) break;
               \langle \text{ Delete or replace } qqq \text{ from the hull } 15 \rangle;
           r = qqq;
         }
         qqq = qq;
         while (1) {
           if (qqq \equiv q) break;
           oo, rr = qqq \neg succ;
           if (oo, ccw(vv, qqq \rightarrow vert, rr \rightarrow vert)) break;
           \langle Delete or replace qqq from the hull 15\rangle;
            qqq = rr;
        if (\neg replaced) (Insert vv in tree, linked by ppar 18);
        if (n < 150)
           printf("New_hull_sequence_h(%s;_ks;_ks))^n, r-vert-name, vv-name, qqq-vert-name);
```

9

HULL UPDATING

This code is used in section 10.

10 HULL UPDATING HULLTR §18

```
\langle Insert vv in tree, linked by ppar 18 \rangle \equiv
18.
      tt = next\_node ++;
     o, tt \neg vert = vv;
     o, tt \neg succ = qq;
      o, tt \neg pred = r;
      o,\,qq \neg pred \,=\, tt;
      o, r \rightarrow succ = tt;
      oo, tt \rightarrow left = tt \rightarrow right = \Lambda;
      o, tt \rightarrow prio = gb\_next\_rand();
     if (n < 150) printf("(Inserting_\%s_\at_\bottom_\of_\tree, \prio=\%d)\n", vv\rightarrow name, tt\rightarrow prio);
      if (o, tt \neg prio < preppar \neg prio) rotup(preppar, ppar, tt, tt \neg prio);
      else { /* easy case, no rotation needed */
        o, tt \neg parent = preppar;
        o,*ppar = tt;
   }
```

This code is used in section 17.

19. Determinants. I need code for the primitive function *ccw*. Floating-point arithmetic suffices for my purposes.

We want to evaluate the determinant

```
ccw(u, v, w) = \begin{vmatrix} u(x) & u(y) & 1 \\ v(x) & v(y) & 1 \\ w(x) & w(y) & 1 \end{vmatrix} = \begin{vmatrix} u(x) - w(x) & u(y) - w(y) \\ v(x) - w(x) & v(y) - w(y) \end{vmatrix}.
```

```
\langle \text{Procedures } 9 \rangle + \equiv
   int ccw(u, v, w)
            Vertex *u, *v, *w;
    { register double wx = (double) w \rightarrow x.I, wy = (double) w \rightarrow y.I;}
        register double det = ((double) \ u \rightarrow x.I - wx) * ((double) \ v \rightarrow y.I - wy) - ((double)
               u \rightarrow y.I - wy) * ((\mathbf{double}) \ v \rightarrow x.I - wx);
        Vertex *uu = u, *vv = v, *ww = w, *t;
        if (det \equiv 0) {
            det = 1;
           \text{if } (u \neg x.I > v \neg x.I \lor (u \neg x.I \equiv v \neg x.I \land (u \neg y.I > v \neg y.I \lor (u \neg y.I \equiv v \neg y.I \land u \neg z.I > v \neg z.I)))) \  \, \{ v \neg x.I > v \neg x.I \land (u \neg y.I \equiv v \neg x.I \land (u \neg y.I \rightarrow v \neg x.I)) \} \  \, \{ v \neg x.I > v \neg x.I \land (u \neg x.I \equiv v \neg x.I) \land (u \neg y.I \rightarrow v \neg x.I) \} \  \, \} 
               t = u; u = v; v = t; det = -det;
           if (v \rightarrow x.I > w \rightarrow x.I \lor (v \rightarrow x.I \equiv w \rightarrow x.I \land (v \rightarrow y.I > w \rightarrow y.I \lor (v \rightarrow y.I \equiv w \rightarrow y.I \land v \rightarrow z.I > w \rightarrow z.I)))) {
               t = v; v = w; w = t; det = -det;
           \text{if } (u \neg x.I > v \neg x.I \lor (u \neg x.I \equiv v \neg x.I \land (u \neg y.I > v \neg y.I \lor (u \neg y.I \equiv v \neg y.I \land u \neg z.I < v \neg z.I)))) \  \, \{ v \neg x.I > v \neg x.I \lor (u \neg x.I \equiv v \neg x.I \land (u \neg y.I \Rightarrow v \neg x.I)) \} \  \, \{ v \neg x.I > v \neg x.I \lor (u \neg x.I \equiv v \neg x.I \land (u \neg x.I \Rightarrow v \neg x.I)) \} \  \, \{ v \neg x.I > v \neg x.I \lor (u \neg x.I \equiv v \neg x.I) \land (u \neg x.I \equiv v \neg x.I) \} \  \, \} 
               det = -det;
            }
        if (n < 150)
            printf("cc(%s; u%s; u%s) uisu%s n", uu - name, vv - name, ww - name, det > 0? "true": "false");
        ccs++;
        return (det > 0);
Area: 5.
                                                                                                  n: 1.
ccs: 1, 2, 19.
                                                                                                  name: 6, 8, 9, 11, 12, 13, 15, 17, 18, 19.
ccw: 2, 10, 12, 17, <u>19</u>.
                                                                                                  next_node: 4, 5, 6, 13, 18.
count: 8, 9.
                                                                                                  node: 3, 4, 5, 7, 9, 14, 16.
deldown: 15, \underline{16}.
                                                                                                  node_struct: 3.
det: \underline{19}.
                                                                                                  o: \underline{2}.
g: \underline{1}.
                                                                                                  oo: 2, 6, 10, 12, 13, 15, 17, 18.
qb\_alloc: 4.
                                                                                                  ooo: 2, 6.
gb\_graph: 1.
                                                                                                 p: \quad \underline{7}, \ \underline{9}, \ \underline{14}, \ \underline{16}.
gb\_init\_rand: 6.
                                                                                                  par: 7, 15.
gb\_next\_rand: 6, 13, 18.
                                                                                                  parent: \underline{3}, 6, 8, 13, 14, 15, 16, 18.
Graph: 1.
                                                                                                  pp: \frac{7}{14}, \frac{14}{15}, \frac{16}{16}.
head: 3, 4, \underline{5}, 6, 8, 10, 11, 12, 13, 17.
                                                                                                  ppar: 7, 17, 18.
init\_area: 6.
                                                                                                  ppr: \underline{14}.
inorder: 8, 9.
                                                                                                  pr: 14.
left: 3, 6, 8, 9, 13, 14, 15, 16, 17, 18.
                                                                                                  pred: 3, 6, 8, 10, 12, 13, 15, 17, 18.
main: \underline{1}.
                                                                                                  prepar: \underline{7}, 15.
mems: 1, 2.
                                                                                                  preppar: \underline{7}, 17, 18.
miles: 1.
                                                                                                  printf: 1, 6, 8, 9, 11, 12, 13, 14, 15, 16, 17, 18, 19.
```

12 DETERMINANTS HULLTR §19

```
prio: 3, 6, 8, 9, 13, 14, 15, 16, 18.
q: \frac{7}{14}.
ql: \underline{16}.
qlp: \underline{16}.
qp: \underline{14}.
qq: \quad \underline{7}, \ 12, \ \underline{14}, \ \underline{16}, \ 17, \ 18.
qqq: \ \ \underline{7}, \ 12, \ 15, \ 17.
qr: \underline{16}.
qrp: \underline{16}.
r: \underline{7}.
replaced: 7, 10, 12, 15, 17.
right: 3, 6, 8, 9, 13, 14, 15, 16, 17, 18.
rotup: 13, \ \underline{14}, \ 18.
rr: \quad \underline{7}, \quad 17.
s: \underline{7}.
serial\_no: \underline{5}, 10.
ss: \frac{7}{2}, 15.
succ: 3, 6, 8, 11, 12, 13, 15, 17, 18.
t: \underline{19}.
tt: <u>7</u>, 13, 15, 18.
u: \ \underline{7}, \ \underline{19}.
uu: \underline{19}.
v: \underline{7}, \underline{19}.
vert: \ \underline{3}, 6, 8, 9, 10, 11, 12, 13, 15, 17, 18.
Vertex: 7, 19.
vertex\_struct: 3.
vertices: 6, 10.
vv: \quad \underline{7}, \ 10, \ 12, \ 13, \ 15, \ 17, \ 18, \ \underline{19}.
w: \ \ \underline{7}, \ \underline{19}.
working_storage: 4, \underline{5}, 6.
ww: \underline{19}.
wx: \underline{19}.
wy: \underline{19}.
```

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