§1 UNAVOIDABLE INTRO 1

November 24, 2020 at 13:25

1. Intro. A quickie to find a longest string that avoids the interesting set of "unavoidable" m-ary strings of length n constructed by Champarnaud, Hansel, and Perrin.

Their construction can be viewed as finding the minimum number of arcs to remove from the de Bruijn graph of (n-1)-tuples so that the resulting graph has no oriented cycles. (Because each n-letter string corresponds to an arc that must be avoided.)

This program constructs the graph and finds a longest path.

```
/* this many letters in the alphabet */
#define m 2
#define n 20
                       /* this many letters in each string */
                                      /* m^{n-1} */
#define space (1 \ll (n-1))
#include <stdio.h>
  char avoid[m*space];
                                 /* nonzero if the arc is removed */
                       /* outdegree, also used as pointer to next level */
  int deg[space];
  int link[space];
                        /* stack of vertices whose degree has dropped to zero */
                     /* staging area */
  int a[n+1];
                   /* the number of vertices on the current level */
  int count;
  int code;
                 /* an n-tuple represented in m-ary notation */
  main()
     register int d, j, k, l, q;
                           /* top of the linked stack */
     register int top;
     \langle \text{ Compute the } avoid \text{ and } deg \text{ tables } 2 \rangle;
     for (d = 0; count; d \leftrightarrow) {
       printf(\texttt{"Vertices}_{\square} \texttt{at}_{\square} \texttt{distance}_{\square} \texttt{%d:}_{\square} \texttt{%d} \texttt{`n"}, d, count);
       for (l = top, top = -1, count = 0; l \ge 0; l = link[l])
          (Decrease the degree of l's predecessors, and stack them if their degree drops to zero 5)
     ⟨Print out a longest path 6⟩;
     Algorithm 7.2.1.1F gives us the relevant prime powers here.
\langle Compute the avoid and deg tables 2\rangle \equiv
  for (j = 0; j < space; j++) deg[j] = m;
  count = d = 0;
  top = -1;
  for (j = n; j; j--) \ a[j] = 0;
  a[0] = -1, j = 1;
  while (1) {
    if (n \% j \equiv 0) (Generate an n-tuple to avoid 3);
     for (j = n; a[j] \equiv m - 1; j - -);
     if (j \equiv 0) break;
     a[j]++;
     for (k = j + 1; k \le n; k++) a[k] = a[k - j];
  printf("m=%d,_lm=%d:_lavoiding_lone_larc_lin_leach_lof_l%d_ldisjoint_lcycles\n", m, n, d);
```

This code is used in section 1.

2 INTRO UNAVOIDABLE §3

3. At this point $\lambda = a_1 \dots a_j$ is a prime string and $\alpha = a_1 \dots a_n = \lambda^{n/j}$. The crux of the Champarnaud/Hansel/Perrin method is to find the shortest prime μ such that α has the form $\mu^{\lfloor n/|\mu|\rfloor}\beta$, and to avoid the string $\beta \mu^{\lfloor n/|\mu|\rfloor}$.

We have $\mu = \lambda$ and $\beta = \epsilon$ if j < n. Otherwise we can use Duval's algorithm to discover all the prime prefixes of α , stopping when one of them has the desired form. (The resulting algorithm is quite pretty, if I do say so myself.)

```
\langle Generate an n-tuple to avoid _3\rangle \equiv
     d++;
     if (j < n) l = j, q = n;
        for (l = 1, k = 2; ; k++) {
             /* at this point a_1 \dots a_l is prime, and a_1 \dots a_{k-1} is its (k-1)-extension */
          if (a[k-l] < a[k]) {
             q = l * (\mathbf{int})(n/l);
             if (k > q) break;
             l=k;
             if (k \equiv n) break;
        }
     for (code = 0, k = q + 1; k < n; k++) code = m * code + a[k];
     for (k = 1; k \le q; k++) code = m * code + a[k];
     \langle \text{ Avoid the } n\text{-tuple encoded by } code \ 4 \rangle;
  }
This code is used in section 2.
     \langle Avoid the n-tuple encoded by code \langle \rangle \equiv
  avoid[code] = 1;
  q = code/m;
  deg[q]—;
  if (deg[q] \equiv 0) deg[q] = -1, link[q] = top, top = q, count ++;
This code is used in section 3.
5. \(\text{Decrease the degree of } l'\) predecessors, and stack them if their degree drops to zero 5\) \(\simega
  for (j = m - 1; j \ge 0; j - -) {
     k = l + j * space;
     if (\neg avoid[k]) {
       q = k/m;
        deg[q]--;
        if (deg[q] \equiv 0) deg[q] = l, link[q] = top, top = q, count ++;
This code is used in section 1.
```

 $\S6$ Unavoidable intro 3

6. Here I apologize for using a dirty trick: The current value of k happens to be the most recent value of l, a vertex with no predecessors.

This code is used in section 1.

4 INDEX UNAVOIDABLE §7

7. Index.

UNAVOIDABLE NAMES OF THE SECTIONS 5

UNAVOIDABLE

	Section	Pag	ge
Intro			1
Indov	7		/