

November 24, 2020 at 13:23

**1. Intro.** Given the specification of a filomino puzzle in *stdin*, this program outputs DLX data for the problem of finding all solutions.

The specification consists of  $m$  lines of  $n$  entries each. An entry is either ‘.’ or a digit from 1 to 9 or **a** to **f**.

A solution means that all ‘.’ entries are replaced by digits. Every maximal rookwise connected set of cells labeled  $d$  must be a  $d$ -omino.

The maximum digit in the solution will be the maximum digit specified. (For example, the program will make no attempt to fit pentominoes into the blank cells, if all of the specified digits are less than 5.)

The main interest in this program is its method for finding all feasible  $d$ -ominoes that cover a given entry  $d$ : They must not be adjacent to a  $d$  that’s not included. The algorithm used here is an instructive generalization of Algorithm R in exercise 7.2.2–75 of *The Art of Computer Programming*.

```
#define maxn 16      /* at most 16 (or I'll have to go beyond hex) */
#define maxd 16      /* digits of the solution must be less than this */
#define bufsize 80
#define pack(i,j) (((i)+1) << 8) + (j) + 1
#define unpack(ij) icoord = ((ij) >> 8) - 1, jcoord = ((ij) & #ff) - 1
#define board(i,j) brd[pack(i,j)]
#define panic(message)
    { fprintf(stderr, "%s: %s", message, buf); exit(-1); }

#include <stdio.h>
#include <stdlib.h>
char buf[bufsize];
int brd[pack(maxn, maxn)]; /* the given pattern */
int dmax; /* the maximum digit seen */
⟨Global data structures for Algorithm R 18⟩;
⟨Subroutines 19⟩;

main()
{
    register int a, d, i, j, k, l, m, n, p, q, s, t, u, v, di, dj, icoord, jcoord;
    ⟨Read the input into board 2⟩;
    ⟨Print the item-name line 3⟩;
    for (d = 1; d ≤ dmax; d++) ⟨Print all the options for d-ominoes 4⟩;
}
```

2.  $\langle$  Read the input into *board* 2  $\rangle \equiv$ 

```

printf("l_filomino-dlx:\n");
for (i = n = t = 0; i ≤ maxn; i++) {
    if (!fgets(buf, bufsize, stdin)) break;
    printf("l_%s", buf);
    for (j = k = 0; ; j++, k++) {
        if (buf[k] == '\n') break;
        if (buf[k] == '.') continue;
        if (buf[k] ≥ '1' ∧ buf[k] ≤ '9') board(i, j) = buf[k] - '0', t++;
        else if (buf[k] ≥ 'a' ∧ buf[k] ≤ 'f') board(i, j) = buf[k] - 'a' + 10, t++;
        else panic("illegal_entry");
        if (board(i, j) > dmax) dmax = board(i, j);
    }
    if (j > n) n = j; /* short rows are extended with '.'s */
}
if (i > maxn) panic("too_many_rows");
m = i;
for (i = 0; i < m; i++) board(i, -1) = board(i, n) = -1; /* frame the board */
for (j = 0; j < n; j++) board(-1, j) = board(m, j) = -1;
fprintf(stderr, "OK, I've read %d clues, for a %dx%d board.\n", t, dmax, m, n);
mm = m, nn = n;

```

This code is used in section 1.

3. There are primary items  $ij$  for  $0 \leq i < m$  and  $0 \leq j < n$ . They represent the cells to be filled.

There are secondary items  $h dij$  for each boundary edge of a  $d$ -omino between  $(i, j - 1)$  and  $(i, j)$ , for  $0 \leq i < m$  and  $1 \leq j < n$ . Similarly, secondary items  $v dij$  for  $1 \leq i < m$  and  $0 \leq j < n$  are for boundaries between  $(i - 1, j)$  and  $(i, j)$  in the vertical dimension.

 $\langle$  Print the item-name line 3  $\rangle \equiv$ 

```

for (i = 0; i < m; i++)
    for (j = 0; j < n; j++) printf("%x%x", i, j);
printf("|");
for (i = 0; i < m; i++)
    for (j = 1; j < n; j++)
        for (d = 1; d ≤ dmax; d++) printf("_h%x%x%x", d, i, j);
for (i = 1; i < m; i++)
    for (j = 0; j < n; j++)
        for (d = 1; d ≤ dmax; d++) printf("_v%x%x%x", d, i, j);
printf("\n");

```

This code is used in section 1.

4.  $\langle$  Print all the options for  $d$ -ominoes 4  $\rangle \equiv$ 

```

{
    for (di = 0; di < m; di++)
        for (dj = 0; dj < n; dj++)  $\langle$  Print the options for  $d$ -ominoes starting at  $(di, dj)$  5  $\rangle$ ;
}

```

This code is used in section 1.

**5.** Now comes the interesting part. I assume the reader is familiar with Algorithm R in the solution to exercise 7.2.2–75. But we add a new twist: A *forced move* is made to a  $d$ -cell if we’ve chosen a vertex adjacent to it. The first vertex ( $v_0$ ) is also considered to be forced.

Since I’m not operating with a general graph, the **ARCS** and **NEXT** aspects of Algorithm R are replaced with a simple scheme: Codes 1, 2, 3, 4 are used respectively for north, west, east, and south. In other words, the operation ‘ $a \leftarrow \text{ARCS}(v)$ ’ is changed to ‘ $a \leftarrow 1$ ’; ‘ $a \leftarrow \text{NEXT}(a)$ ’ is changed to ‘ $a \leftarrow a + 1$ ’; ‘ $a = \Lambda$ ?’ becomes ‘ $a = 5$ ?’ The vertex  $\text{TIP}(a)$  is the cell north, west, east, or south of  $v$ , depending on  $a$ .

A forced move at level  $l$  is indicated by  $a_l = 0$ .

If cell  $(di, dj)$  is not already filled, we fill it with a  $d$ -mino that uses only unfilled cells and doesn’t come next to a  $d$ -cell.

```

⟨Print the options for  $d$ -ominoes starting at  $(di, dj)$  5⟩ ≡
{
   $u = \text{pack}(di, dj)$ ;
  if ( $\neg \text{board}(di, dj)$ ) {
    for ( $q = 1$ ;  $q \leq 4$ ;  $q++$ )
      if ( $\text{brd}[u + \text{dir}[q]] \equiv d$ ) break;    /* next to  $d$  */
    if ( $q \leq 4$ ) continue;
    forcing = 0;
  } else if ( $\text{board}(di, dj) \neq d$ ) continue;
  else forcing = 1;
  ⟨Do step R1 6⟩;
  ⟨Do step R2 7⟩;
  ⟨Do step R3 11⟩;
  ⟨Do step R4 12⟩;
  ⟨Do step R5 13⟩;
  ⟨Do step R6 15⟩;
  ⟨Do step R7 17⟩;
  done: checktags();
}

```

This code is used in section 4.

```

6.  ⟨Do step R1 6⟩ ≡
r1:  /* initialize */
    for ( $i = 0$ ;  $i < m$ ;  $i++$ )
      for ( $j = 0$ ;  $j < n$ ;  $j++$ ) tag[pack( $i, j$ )] = 0;
   $v = vv[0] = u$ , tag[ $v$ ] = 1;
   $i = ii[0] = 0$ ,  $a = aa[0] = 0$ ,  $l = 1$ ;

```

This code is used in section 5.

**7.** At the beginning of step R2, we’ve just chosen the vertex  $u$ , which is  $vv[l - 1]$ . If  $l > 1$ , it’s a vertex adjacent to  $v = vv[i]$  in direction  $a$ , where  $i = ii[l - 1]$  and  $a = aa[l - 1]$ .

```

⟨Do step R2 7⟩ ≡
r2:  /* enter level  $l$  */
    if (forcing) ⟨Make forced choices of all  $d$ -cells adjacent to  $u$ ; but goto r7 if there’s a problem 8⟩;
    if ( $l \equiv d$ ) {
      ⟨Print an option for the current  $d$ -omino 9⟩;
      ⟨Undo the latest forced moves 10⟩;
    }

```

This code is used in section 5.

**8.** Ye olde depth-first search.

If forcing, we backtrack if the  $d$ -omino gets too big, or if we're forced to choose a  $d$ -cell whose options have already been considered.

If not forcing, we backtrack if we're next to a  $d$ -cell, or if solutions for this cell have already been considered.

⟨ Make forced choices of all  $d$ -cells adjacent to  $u$ ; but **goto** *r7* if there's a problem *8* ⟩  $\equiv$

```

for ( $stack[0] = u, s = 1; s;$ ) {
     $u = stack[--s];$ 
    for ( $q = 1; q \leq 4; q++$ ) {
         $t = u + dir[q];$ 
        if ( $brd[t] \neq d$ ) continue;    /* not a  $d$ -cell */
        if ( $tag[t]$ ) continue;        /* we've already chosen this  $d$ -cell */
        if ( $t < vv[0]$ ) goto r7;      /* it came earlier than  $(di, dj)$  */
        if ( $l \equiv d$ ) goto r7;        /* we've already got  $d$  vertices */
         $aa[l] = 0, vv[l++] = t, tag[t] = 1, stack[s++] = t;$     /* forced move to  $t$  */
    }
}

```

This code is used in section 7.

**9.** OK, we've got a viable  $d$ -omino to pass to the output.

⟨ Print an option for the current  $d$ -omino *9* ⟩  $\equiv$

```

{
     $curstamp++;$ 
    for ( $p = 0; p < d; p++$ ) {
         $unpack(vv[p]);$ 
         $printf("\%x\%x", icoord, jcoord);$ 
         $stamp[vv[p]] = curstamp;$ 
    }
    for ( $p = 0; p < d; p++$ ) {
         $unpack(vv[p]);$ 
        for ( $q = 1; q \leq 4; q++$ )
            if ( $stamp[vv[p] + dir[q]] \neq curstamp$ ) {    /* boundary edge detected */
                switch ( $q$ ) {
                    case 1: if ( $icoord$ )  $printf("\%x\%x\%x", d, icoord, jcoord);$  break;
                    case 2: if ( $jcoord$ )  $printf("\%h\%x\%x\%x", d, icoord, jcoord);$  break;
                    case 3: if ( $jcoord < n - 1$ )  $printf("\%h\%x\%x\%x", d, icoord, jcoord + 1);$  break;
                    case 4: if ( $icoord < m - 1$ )  $printf("\%v\%x\%x\%x", d, icoord + 1, jcoord);$  break;
                }
            }
    }
     $printf("\n");$ 
}

```

This code is used in section 7.

**10.** ⟨ Undo the latest forced moves *10* ⟩  $\equiv$ 

```

for ( $l--; aa[l] \equiv 0; l--$ ) {
    if ( $l \equiv 0$ ) goto done;
     $tag[vv[l]] = 0;$ 
}

```

This code is used in sections 7 and 17.

**11.**  $\langle \text{Do step R3 11} \rangle \equiv$   
 $r3:$   $\quad /* \text{ advance } a */$   
 $\quad a++;$

This code is used in section 5.

**12.**  $\langle \text{Do step R4 12} \rangle \equiv$   
 $r4:$   $\quad /* \text{ done with level? } */$   
 $\quad \text{if } (a \neq 5) \text{ goto } r5;$   
 $\quad \text{if } (i \equiv l - 1) \text{ goto } r6;$   
 $\quad v = vv[+i], a = 1;$

This code is used in section 5.

**13.**  $\langle \text{Do step R5 13} \rangle \equiv$   
 $r5:$   $\quad /* \text{ try } a */$   
 $\quad u = v + dir[a];$   
 $\quad \text{if } (brd[u]) \text{ goto } r3; \quad /* \text{ not really a neighbor of } v */$   
 $\quad tag[u]++;$   
 $\quad \text{if } (tag[u] > 1) \text{ goto } r3; \quad /* \text{ already chosen } */$   
 $\quad \text{if } (\neg forcing) \langle \text{If } u \text{ was already handled, or if it's adjacent to a } d\text{-cell, goto } r3 \text{ 14} \rangle;$   
 $\quad ii[l] = i, aa[l] = a, vv[l] = u, l++;$   
 $\quad \text{goto } r2;$

This code is used in section 5.

**14.**  $\langle \text{If } u \text{ was already handled, or if it's adjacent to a } d\text{-cell, goto } r3 \text{ 14} \rangle \equiv$   
 $\{$   
 $\quad \text{if } (u < vv[0]) \text{ goto } r3; \quad /* \text{ it's earlier than } (di, dj) */$   
 $\quad \text{if } (brd[u]) \text{ goto } r3; \quad /* \text{ not a blank cell } */$   
 $\quad \text{for } (q = 1; q \leq 4; q++)$   
 $\quad \quad \text{if } (brd[u + dir[q]] \equiv d) \text{ goto } r3;$   
 $\}$

This code is used in section 13.

**15.**  $\langle \text{Do step R6 15} \rangle \equiv$   
 $r6:$   $\quad /* \text{ backtrack } */$   
 $\quad \langle \text{Undo previous forced moves 16} \rangle;$   
 $\quad \text{for } (i = ii[l], k = i + 1; k \leq l; k++) \{$   
 $\quad \quad t = vv[k];$   
 $\quad \quad \text{for } (q = 1; q \leq 4; q++)$   
 $\quad \quad \quad \text{if } (brd[t + dir[q]] \equiv 0) tag[t + dir[q]]--; \quad /* \text{ untag the neighbors of } vv[k] */$   
 $\quad \quad \}$   
 $\quad \text{for } (a = aa[l] + 1, v = vv[i]; a \leq 4; a++)$   
 $\quad \quad \text{if } (brd[v + dir[a]] \equiv 0) tag[v + dir[a]]--; \quad /* \text{ untag late neighbors of } vv[i] */$   
 $\quad \quad a = aa[l];$   
 $\quad \quad \text{goto } r3;$

This code is used in section 5.

16.  $\langle \text{Undo previous forced moves 16} \rangle \equiv$   
**for** ( $l \leftarrow$ ;  $aa[l] \equiv 0$ ;  $l \leftarrow$ ) {  
     **if** ( $l \equiv 0$ ) **goto** *done*;  
      $t = vv[l]$ ;  
     **for** ( $q = 1$ ;  $q \leq 4$ ;  $q++$ )  
         **if** ( $brd[t + dir[q]] \equiv 0$ )  $tag[t + dir[q]] \leftarrow$ ;     /\* untag the neighbors of  $vv[l]$  \*/  
          $tag[t] = 0$ ;  
**}**

This code is used in section 15.

17.  $\langle \text{Do step R7 17} \rangle \equiv$   
 $r7$ :     /\* recover from bad forcing \*/  
      $\langle \text{Undo the latest forced moves 10} \rangle$ ;  
      $i = ii[l]$ ,  $v = vv[i]$ ,  $a = aa[l]$ ;  
     **goto**  $r3$ ;

This code is used in section 5.

18.  $\langle \text{Global data structures for Algorithm R 18} \rangle \equiv$   
**int** *forcing*;  
**int**  $dir[5] = \{0, -(1 \ll 8), -1, 1, 1 \ll 8\}$ ;  
**int**  $tag[pack(maxn, maxn)]$ ;  
**int**  $vv[maxd]$ ,  $aa[maxd]$ ,  $ii[maxd]$ ,  $stack[maxd]$ ;     /\* state variables \*/  
**int** *curstamp*;  
**int**  $stamp[pack(maxn, maxn)]$ ;  
**int**  $mm, nn$ ;

This code is used in section 1.

19.  $\langle \text{Subroutines 19} \rangle \equiv$   
**void** *debug*(**char** \**message*)  
   {  
     *fprintf*(*stderr*, "%s!\n", *message*);  
   }

See also sections 20 and 21.

This code is used in section 1.

20. Here's a handy routine for debugging the tricky parts.

$\langle \text{Subroutines 19} \rangle + \equiv$   
**void** *showtags*(**void**)  
   {  
     **register** **int**  $i, j$ ;  
     **for** ( $i = 0$ ;  $i < mm$ ;  $i++$ )  
         **for** ( $j = 0$ ;  $j < nn$ ;  $j++$ )  
             **if** ( $tag[pack(i, j)]$ ) *printf*("%x%x:%d\n",  $i, j, tag[pack(i, j)]$ );  
   }

21.  $\langle \text{Subroutines 19} \rangle + \equiv$

```

void checktags(void)
{
    register int i, j, q;
    for (i = 0; i < mm; i++)
        for (j = 0; j < nn; j++)
            if (tag[pack(i, j)] {
                if (pack(i, j)  $\equiv$  vv[0]) continue;
                for (q = 1; q  $\leq$  4; q++)
                    if (pack(i, j)  $\equiv$  vv[0] + dir[q]) break;
                if (q  $\leq$  4) continue;
                debug("bad_␣tag");
            }
}

```

**22. Index.**

*a*: [1](#).  
*aa*: [6](#), [7](#), [8](#), [10](#), [13](#), [15](#), [16](#), [17](#), [18](#).  
*board*: [1](#), [2](#), [5](#).  
*brd*: [1](#), [5](#), [8](#), [13](#), [14](#), [15](#), [16](#).  
*buf*: [1](#), [2](#).  
*bufsize*: [1](#), [2](#).  
*checktags*: [5](#), [21](#).  
*curstamp*: [9](#), [18](#).  
*d*: [1](#).  
*debug*: [19](#), [21](#).  
*di*: [1](#), [4](#), [5](#), [8](#), [14](#).  
*dir*: [5](#), [8](#), [9](#), [13](#), [14](#), [15](#), [16](#), [18](#), [21](#).  
*dj*: [1](#), [4](#), [5](#), [8](#), [14](#).  
*dmax*: [1](#), [2](#), [3](#).  
*done*: [5](#), [10](#), [16](#).  
*exit*: [1](#).  
*fgets*: [2](#).  
*forcing*: [5](#), [7](#), [13](#), [18](#).  
*fprintf*: [1](#), [2](#), [19](#).  
*i*: [1](#), [20](#), [21](#).  
*icoord*: [1](#), [9](#).  
*ii*: [6](#), [7](#), [13](#), [15](#), [17](#), [18](#).  
*ij*: [1](#).  
*j*: [1](#), [20](#), [21](#).  
*jcoord*: [1](#), [9](#).  
*k*: [1](#).  
*l*: [1](#).  
*m*: [1](#).  
*main*: [1](#).  
*maxd*: [1](#), [18](#).  
*maxn*: [1](#), [2](#), [18](#).  
*message*: [1](#), [19](#).  
*mm*: [2](#), [18](#), [20](#), [21](#).  
*n*: [1](#).  
*nn*: [2](#), [18](#), [20](#), [21](#).  
*p*: [1](#).  
*pack*: [1](#), [5](#), [6](#), [18](#), [20](#), [21](#).  
*panic*: [1](#), [2](#).  
*printf*: [2](#), [3](#), [9](#), [20](#).  
*q*: [1](#), [21](#).  
*r1*: [6](#).  
*r2*: [7](#), [13](#).  
*r3*: [11](#), [13](#), [14](#), [15](#), [17](#).  
*r4*: [12](#).  
*r5*: [12](#), [13](#).  
*r6*: [12](#), [15](#).  
*r7*: [8](#), [17](#).  
*s*: [1](#).  
*showtags*: [20](#).  
*stack*: [8](#), [18](#).  
*stamp*: [9](#), [18](#).

*stderr*: [1](#), [2](#), [19](#).  
*stdin*: [1](#), [2](#).  
*t*: [1](#).  
*tag*: [6](#), [8](#), [10](#), [13](#), [15](#), [16](#), [18](#), [20](#), [21](#).  
*u*: [1](#).  
*unpack*: [1](#), [9](#).  
*v*: [1](#).  
*vv*: [6](#), [7](#), [8](#), [9](#), [10](#), [12](#), [13](#), [14](#), [15](#), [16](#), [17](#), [18](#), [21](#).



- ⟨ Do step R1 6 ⟩    Used in section 5.
- ⟨ Do step R2 7 ⟩    Used in section 5.
- ⟨ Do step R3 11 ⟩    Used in section 5.
- ⟨ Do step R4 12 ⟩    Used in section 5.
- ⟨ Do step R5 13 ⟩    Used in section 5.
- ⟨ Do step R6 15 ⟩    Used in section 5.
- ⟨ Do step R7 17 ⟩    Used in section 5.
- ⟨ Global data structures for Algorithm R 18 ⟩    Used in section 1.
- ⟨ If  $u$  was already handled, or if it's adjacent to a  $d$ -cell, **goto**  $r3$  14 ⟩    Used in section 13.
- ⟨ Make forced choices of all  $d$ -cells adjacent to  $u$ ; but **goto**  $r7$  if there's a problem 8 ⟩    Used in section 7.
- ⟨ Print all the options for  $d$ -ominoes 4 ⟩    Used in section 1.
- ⟨ Print an option for the current  $d$ -omino 9 ⟩    Used in section 7.
- ⟨ Print the item-name line 3 ⟩    Used in section 1.
- ⟨ Print the options for  $d$ -ominoes starting at  $(di, dj)$  5 ⟩    Used in section 4.
- ⟨ Read the input into *board* 2 ⟩    Used in section 1.
- ⟨ Subroutines 19, 20, 21 ⟩    Used in section 1.
- ⟨ Undo previous forced moves 16 ⟩    Used in section 15.
- ⟨ Undo the latest forced moves 10 ⟩    Used in sections 7 and 17.

# FILOMINO-DLX

	Section	Page
Intro .....	<a href="#">1</a>	1
Index .....	<a href="#">22</a>	8