

8)

pseudocode

q8 (A, m) {

$t[m][\log_2 m + 1]$

for ($i = 0, m-1$) {

$t[i][0] \leftarrow i$

}

for ($i = 1, \log_2 m$) {

for ($j = 0, m-2^i$) {

$a \leftarrow t[j][i-1]$

$b \leftarrow t[j+2^{i-1}][i-1]$

if ($A[a] > A[b]$) {

$t[j][i] \leftarrow a$

} else if ($A[a] \leq A[b]$) {

$t[j][i] \leftarrow b$

}

}

}

return t ;

}

q8-2 (A, n) {

$l \leftarrow 0, r \leftarrow m-1$

while ($l < r$) {

$mid = \frac{l+r}{2}$

if ($A[mid] \geq n$) {

$r \leftarrow mid$

} else {

$l \leftarrow mid + 1$

}

}

return r ;

}

```

qs-u(A, n) {
    l ← 0, r ← n-1
    while (l ≤ r) {
        mid =  $\frac{l+r}{2}$ 
        if (A[mid] > x) {
            r ← mid;
        } else {
            l ← mid + 1
        }
    }
    return r;
}

```

algorithm: - we will first sort A and then store unique values in B and frequency of each unique value in C (index)

we define 2 points for upper & lower bound for array, using them we can find the no. of elements in A b/w (a, b].

now we create a 2D matrix, matrix element (i][j] store index of max/min occurring element from our array C

now find indices such that $B(i) \geq a$ & $B(j) \leq b$

~~the~~ answer can be found in table.

time complexity: \rightarrow as this $\rightarrow O(n \log n)$