

Plurality Points and Condorcet Points in Euclidean Space

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Outline

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- 3 Necessary lemmata by case
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Informal description

We can represent opinion standpoints by points in space. Goal: find a (spacial) equilibrium between these quantified “opinions”.

- Voters: Multiset in Euclidean space (\mathbb{R}^d with ℓ_2 norm)
- Plurality point: Closer to at least as many voters as any other point
- Condorcet point: No other point is closer to an absolute majority of voters

Notation

Definition (Voters)

Voters: $V = \{v_1, \dots, v_n\} \subseteq \mathbb{R}^d$ (a multiset, in general). The components are denoted as $v_i = (v_i^{(1)}, \dots, v_i^{(d)})$.

Definition (Multidimensional medians)

Informally: at least half the voters are on each side, per dimension:

$$\begin{aligned} \mathcal{M}_V = \{ \theta : & |\{v \in V : \forall j : v^{(j)} \leq \theta^{(j)}\}| \geq n/2 \\ & \wedge |\{v \in V : \forall j : v^{(j)} \geq \theta^{(j)}\}| \geq n/2 \} \end{aligned}$$

Notation

Definition (Median corners)

Let $V^{(i)}$ be the i -th coordinate of points in V , and $x_l^{(i)}$ and $x_h^{(i)}$ denote the $\lceil n/2 \rceil$ -th and $\lceil (n+1)/2 \rceil$ -th smallest numbers in $V^{(i)}$. Then the set

$$C = \{(x_{t_1}^{(i)}, \dots, x_{t_d}^{(i)}) : t_1, \dots, t_d \in \{l, h\}\}$$

collects all corners of \mathcal{M}_V . Clearly, $|C| \leq 2^d$.

Definition (Half spaces)

Given a line L , we can partition $\mathbb{R}^d = L^+ \cup L \cup L^-$, where L^+ and L^- are the open half spaces separated by L . We define:

$$V_L = L \cap V, \quad V_L^+ = L^+ \cap V, \quad V_L^- = L^- \cap V$$

$$n_L = |V_L|, \quad n_L^+ = |V_L^+|, \quad n_L^- = |V_L^-|$$

Formal definition

Definition (Preference)

- v_i prefers θ_1 to θ_2 if $d(v_i, \theta_1) < d(v_i, \theta_2)$
- $[\theta_1 \succ \theta_2] := \{v \in V : d(v, \theta_1) < d(v, \theta_2)\}$, and accordingly for other relations

Definition (Plurality point (PP))

$\Delta \in \mathbb{R}^d$ is a plurality point iff

$$\forall \theta \in \mathbb{R}^d : |[\theta \succ \Delta]| \leq |[\Delta \succ \theta]|$$

Definition (Condorcet point (CP))

$\Delta \in \mathbb{R}^d$ is a plurality point iff

$$\forall \theta \in \mathbb{R}^d : |[\theta \succ \Delta]| \leq n/2$$

Equivalence of PPs and CPs

Lemma (Equivalence)

In \mathbb{R}^d , a point is a plurality point iff it is a Condorcet point.

Proof.

- 1 A PP is a CP by definition.
- 2
 - Assume x is a CP, but not a PP, so $\exists y : |[x \succ y]| < |[y \succ x]|$.
 - Let z be the midpoint of x, y . Then for each $c \in [y \succeq x]$,
 $d(c, z) < d(c, x)$.
 - Thus, $|[z \succ x]| > n/2$, which is a contradiction.
 - Therefore, a CP is also a PP.

From now on, we constrain ourselves to PPs in \mathbb{R}^2 .

Idea behind algorithm

The presented algorithm works by successively cutting down the space of possible PPs via a series of case distinctions, until only $\mathcal{O}(1)$ candidate points remain, which can be checked exhaustively. For this purpose, the following series of lemmata are employed.

Furthermore:

Definition

Δ_V denotes the set of all PPs of V

Collinear case

Lemma (Collinear points)

If all voters in V are collinear, then

- 1 $\Delta_V = \mathcal{M}_V$, and
- 2 $|\Delta_V| \geq 1$.

Proof.

- 1 has been shown by Hansen & Thisse (1981) for CCs in \mathbb{R}^2 . By the equivalence lemma, this holds also for PPs.
- 2 By definition, $\mathcal{M}_V \neq \emptyset$. When $|V|$ is even, we may have $|\mathcal{M}_V| = |\Delta_V| > 1$.

“Tukey depth condition”

This necessary and sufficient condition is used to decide whether a candidate point is a valid PP.

Lemma

In \mathbb{R}^2 , Δ is a plurality point iff for any line L through Δ , $n_L^+ \leq n/2$ and $n_L^- \leq n/2$.

It is equivalent checking if

$$\min_{L \in \mathcal{L}_\Delta} \{|V \cup \gamma| : \gamma \text{ is a closed halfspace separated by } L\} \geq n/2,$$

where \mathcal{L}_Δ is the space of lines through Δ . The quantity on the left is called *Tukey depth* of Δ with respect to V .

Non-collinear case & medians

Lemma (Non-collinear candidates)

If not all voters in V are collinear, then

- 1 $\Delta_V \subseteq \mathcal{M}_V$, and
- 2 $|\Delta_V| = 0$, or $|\Delta_V| = 1$.

So, a plurality point must be a median, but not vice versa. We can use this to cut down the searched space to the median; however, this is still infinite.

Δ is in \mathcal{M}_V

These lemmata reduce the possibilities in the non-collinear case further to $\mathcal{O}(1)$ candidate points.

Lemma ($\Delta \in V$)

If Δ is a PP of V , and $\Delta \in V$, then $\Delta \in C$, with $V \cup C \leq 2^d$.

Lemma ($\Delta \notin V$)

If $\Delta \notin V$, and L is a separating line through Δ (ie., $n_L = 0$), then Δ is the intersection of the internal tangents of the convex hulls of V_L^+ and V_L^- .

Subroutines

- $\text{SELECT}(S, n)$ returns the n -th smallest point in S . This can be calculated in $\mathcal{O}(|S|)$ time by the median-of-medians selection algorithm from Cormen et. al. Used for median calculation.
- $\text{VERIFYCANDIDATES}(C, V)$ returns all points from C whose Tukey depth in is greater or equal $|V|/2$ (ie., satisfying the condition in the “Tukey depth lemma” above). Tukey depth of one point can be computed in $\mathcal{O}(|V| \log |V|)$ time with an algorithm by Rousseeuw and Struyf (1998).

Subroutines

- `INTERNALTANGENTINTERSECTION(V, W)` returns the intersection of the internal tangents between the convex hulls of V and W . The tangents can be found by solving the linear program

$$\min_{k,d} F(k,d) = k, \quad \text{s.t.}$$

$$\forall v \in V : v_y \leq kv_x + d \quad \text{and}$$

$$\forall w \in W : w_y \geq kw_x + d,$$

and the opposite formulation. Linear programming with fixed dimension can be done in linear time in the number of constraints, by Megiddo (1984).

Pseudocode

```

1: procedure PLURALITYPOINT2D( $V$ )
2:   if all voters are collinear then
3:     return medians of  $V$ 
4:   else  $\triangleright$  Corners of medians
5:      $x_h = \text{SELECT}(V_x, \lceil (n+1)/2 \rceil)$ 
6:      $x_l = \text{SELECT}(V_x, \lceil n/2 \rceil)$ 
7:      $y_h = \text{SELECT}(V_y, \lceil (n+1)/2 \rceil)$ 
8:      $y_l = \text{SELECT}(V_y, \lceil n/2 \rceil)$ 
9:      $C = \{(x_h, y_h), (x_h, y_l), (x_l, y_h), (x_l, y_l)\}$ 

10:    if  $|C| = 1$  then  $\triangleright n$  odd, or  $C$  degenerate
11:      return  $\text{VERIFYCANDIDATES}(C, V)$ 
12:    else
13:       $P = \text{VERIFYCANDIDATES}(V \cap C, V)$ 
14:      if  $P \neq \emptyset$  then  $\triangleright \Delta \in V$ 
15:        return  $P$ 

16:    else
17:      if  $x_h = x_l$  then
18:         $V_a = \{v \in V : v^{(2)} \leq y_h\}$ 
19:         $V_b = \{v \in V : v^{(2)} \geq y_h\}$ 
20:      else
21:         $V_a = \{v \in V : v^{(1)} \leq x_h\}$ 
22:         $V_b = \{v \in V : v^{(1)} \geq x_h\}$ 
23:      end if

24:       $p = \text{INTERNALTANGENTINTERSECTION}(V_a, V_b)$ 
25:      return  $\text{VERIFYCANDIDATES}(\{p\}, V)$ 
26:    end if
27:  end if
28: end if
29: end procedure

```

Analysis

- Detecting collinearity (line 2): $\mathcal{O}(n)$
- Selecting median corners (lines 5–8): $\mathcal{O}(n)$
- Verifying candidates by calculating Tukey depth (lines 11, 13, 25): $\mathcal{O}(n \log n)$
- Computing internal tangent intersection (line 24): $\mathcal{O}(n)$
- Therefore: $\mathcal{O}(n \log n)$ overall complexity

Conclusion

- In Euclidean space, plurality points and Condorcet points are equivalent notions
- In \mathbb{R}^2 , they can be calculated in $\mathcal{O}(n \log n)$ time, using the presented algorithm
- This can be generalized to \mathbb{R}^d in $\mathcal{O}(n^{d-1} \log n)$ time