# Joint Scheduling of URLLC and eMBB Traffic in 5G Wireless Networks

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#### System Model

 Saturated system [1]: Each eMBB user has infinite amount of data to be served.

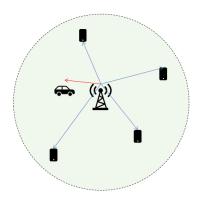


Figure: System model

# System Framework

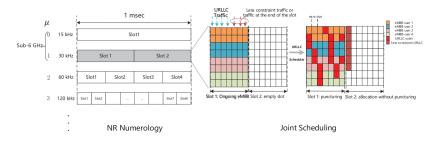


Figure: System framework [2]

#### **Problem Statement**

- URLLC puncturing/superposition commonly used scheduling procedure: allocate *subchannels* for eMBB users in the given time slot first, then schedule *resources* for URLLC users in the respective time minislots accordingly.
- Multiplexing conducting eMBB resource allocation using only eMBB channel state information (CSI) – has been broadly discussed in the literature.
- However, joint scheduling of eMBB and URLLC traffics has not been well studied.
- That is, in addition to eMBB CSI, could we also leverage URLLC information (e.g. observed behavior of URLLC demands, any others?) to schedule eMBB?

#### Examples

- Discussion with Yung-Ching senpai: Reduce eMBB retransmissions of data damaged by URLLC puncturing/superposition.
  - Approach 1: Reserve resources for URLLC traffic.
  - Approach 2: Pre-allocate resources for some URLLC users (less overhead compared to the first method, which increases URLLC reliability).
- However, their paper does not address the aforementioned problems, but aims to tackle two questions:
  - For linear loss model, would the traditional multiplexing procedure be optimal in the long-term?
  - For convex loss model, could we somehow incorporate URLLC demands into eMBB resource allocation formulation?

#### Solution

- The main idea is to perform mathematical analysis on eMBB time slot.
- This is done by their model 'proportion of an eMBB time slot's resource'.
- Linear, convex, and threshold loss models are discussed.

#### Linear Loss Model

A URLLC puncturing problem is formulated as follows:

maximize 
$$\sum_{u} ln(\bar{r_u})$$
  
subject to  $\sum_{u} \alpha_{u,n,l} \leq \frac{1}{l}, \quad \forall n, \forall l,$   
 $\alpha_{u,n,l} \in \left\{0, \frac{1}{l}\right\}, \quad \forall u, \forall n, \forall l,$   
 $\sum_{u} \sum_{l} \beta_{v_0,u,n,m,l} = D_{v_0,n,m}, \quad \forall n, \forall m,$   
 $\beta_{v_0,u,n,m,l} \leq \alpha_{u,n,m,l}, \quad \forall u, \forall n, \forall m, \forall l,$   
 $\beta_{v_0,u,n,m,l} \in \left\{0, \frac{1}{\mathfrak{m}l}\right\}, \quad \forall u, \forall n, \forall m, \forall l,$   
 $(1)$ 

where

$$\alpha_{u,n,m,l} = \frac{\alpha_{u,n,l}}{m}, \quad \forall u, \forall n, \forall m, \forall l$$
 (2)

The average rate is calculated as follows.

$$\bar{r}_u = \frac{r_{u,0} + r_{u,1} + \ldots + r_{u,n-1}}{n}$$
 (3)

$$=\frac{1}{\mathfrak{n}}\sum_{n}r_{u,n}\tag{4}$$

$$=\frac{1}{\mathfrak{n}}\sum_{n}\sum_{l}r_{u,n,l}\tag{5}$$

$$=\frac{1}{\mathfrak{n}}\sum_{n}\sum_{l}\left(\alpha_{u,n,l}-\beta_{v_0,u,n,l}\right)\mathfrak{r}_{u,n,l},\quad\forall u$$
 (6)

In this paper, subchannel allocation is neglected.

$$\bar{r}_u = \frac{1}{\mathfrak{n}} \sum_{n} r_{u,n} \tag{7}$$

$$=\frac{1}{\mathfrak{n}}\sum\left(\alpha_{u,n}-\beta_{v_0,u,n}\right)\mathfrak{r}_{u,n},\quad\forall u\tag{8}$$

• The first idea that comes to my mind is to give each eMBB user an equal share of URLLC demands.

$$\beta_{v_0,u,n,m} = \frac{D_{v_0,n,m}}{\mathfrak{u}}, \quad \forall u, \forall n, \forall m$$
 (9)

• However, this cannot guarantee

$$\beta_{v_0,u,n,m} \le \alpha_{u,n,m}, \quad \forall u, \forall n, \forall m.$$
 (10)

• It is then intuitive to give each eMBB user a proportionally equal share of URLLC demands.

$$\beta_{\nu_0,u,n,m} = \frac{\alpha_{u,n,m}}{\frac{1}{\mathfrak{m}}} D_{\nu_0,n,m} \tag{11}$$

$$=\alpha_{u,n}D_{v_0,n,m} \tag{12}$$

Such an allocation guarantees URLLC constraints:

$$\sum_{u} \beta_{v_0, u, n, m} = D_{v_0, n, m}, \quad \forall n, \forall m,$$

$$\beta_{v_0, u, n, m} \leq \alpha_{u, n, m}, \quad \forall u, \forall n, \forall m.$$
(13)

$$\beta_{v_0,u,n,m} \le \alpha_{u,n,m}, \quad \forall u, \forall n, \forall m.$$
 (14)

URLLC allocation in eMBB time slot can then be derived:

$$\beta_{\mathbf{v}_0,\mathbf{u},\mathbf{n}} = \sum_{\mathbf{m}} \beta_{\mathbf{v}_0,\mathbf{u},\mathbf{n},\mathbf{m}} \tag{15}$$

$$=\alpha_{u,n}\sum_{m}D_{v_0,n,m} \tag{16}$$

$$=\alpha_{u,n}D_{v_0,n} \tag{17}$$

#### Theorem 1

For a wireless system under the linear superposition/puncturing loss model we have that  $C = C^{LR}$ 

 By this theorem, the proportionally equal share policy is optimal.

$$\bar{r}_{u} = \frac{1}{\mathfrak{n}} \sum_{n} \left( \alpha_{u,n} - \alpha_{u,n} D_{v_{0},n} \right) \mathfrak{r}_{u,n} \tag{18}$$

$$=\frac{1}{\mathfrak{n}}\sum_{n}\left(1-D_{\nu_{0},n}\right)\alpha_{u,n}\mathfrak{r}_{u,n},\quad\forall u$$
 (19)

 Do note that URLLC allocation disappears from the objective function of Problem 1, and URLLC demands are constants.

• This problem then can be solved with gradient algorithm [1].

maximize 
$$\sum_{u} \ln(\bar{r}_{u})$$
subject to 
$$\sum_{u} \alpha_{u,n,l} \leq \frac{1}{\mathfrak{l}}, \quad \forall n, \forall l, \qquad (20)$$

$$\alpha_{u,n,l} \in \{0, \frac{1}{\mathfrak{l}}\}, \quad \forall u, \forall n, \forall l$$

 However, this paper does not execute two problem transformations required to apply the above algorithm.

#### Convex Loss Model

Recall that

$$\bar{r}_u = \frac{1}{\mathfrak{n}} \sum_n r_{u,n} \tag{21}$$

$$=\frac{1}{\mathfrak{n}}\sum_{n}\left(\alpha_{u,n}-\alpha_{u,n}f(x)\right)\mathfrak{r}_{u,n}\tag{22}$$

• where  $f(\cdot)$  is a convex function, and

$$x = \frac{\beta_{v_0, u, n}}{\alpha_{u, n}} \tag{23}$$

Thus

$$\bar{r}_{u} = \frac{1}{n} \sum_{n} \left( \alpha_{u,n} - \alpha_{u,n} f\left(\frac{\sum_{m} \beta_{v_{0},u,n,m}}{\alpha_{u,n}}\right) \right) \mathfrak{r}_{u,n}, \quad \forall u \quad (24)$$

• They first fluidize the bandwidth:

maximize 
$$\sum_{u} ln(\bar{r}_{u})$$
  
subject to  $\sum_{u} \alpha_{u,n} = 1 \quad \forall n,$   
 $\alpha_{u,n} \geq 0 \quad \forall u, \forall n,$  (25)  
 $\sum_{u} \beta_{v_{0},u,n,m} = D_{v_{0},n,m} \forall n, \forall m,$   
 $\beta_{v_{0},u,n,m} \leq \alpha_{u,n,m} \ \forall u, \forall n, \forall m,$   
 $\beta_{v_{0},u,n,m} \geq 0 \quad \forall u, \forall n, \forall m$ 

where

$$\alpha_{u,n,m} = \frac{\alpha_{u,n}}{m}, \quad \forall u, \forall n, \forall m$$
 (26)

- I can prove the objective function is concave (by perspective function), and that this is a convex optimization problem.
- Nevertheless, how can we allocate eMBB resources for the current eMBB time slot if the problem depends on unknown URLLC demands in the future?

Assume minislot-homogeneous policy:

$$\beta_{\nu_0,u,n,m} = \frac{\frac{\gamma_{\nu_0,u,n}}{m}}{\frac{1}{m}} D_{\nu_0,n,m}$$
 (27)

$$= \gamma_{\nu_0, u, n} D_{\nu_0, n, m}, \quad \forall u, \forall n, \forall m, \tag{28}$$

where

$$\gamma_{\nu_0,u,n} \ge 0, \quad \forall u, \forall n$$
 (29)

$$\gamma_{\nu_0,u,n} \ge 0, \quad \forall u, \forall n$$

$$\sum_{u} \gamma_{\nu_0,u,n} = 1, \quad \forall n.$$
(29)

Thus

$$\sum_{n} \beta_{\nu_0, u, n, m} = D_{\nu_0, n, m}, \quad \forall n, \forall m$$
 (31)

If

$$\frac{\alpha_{u,n}}{\gamma_{v_0,u,n}} \ge 1, \quad \forall u, \forall n \tag{32}$$

• then combining with  $\mathfrak{m}D_{v_0,n,m} \leq 1, \forall n, \forall m$ ,

$$\frac{\alpha_{u,n}}{\gamma_{v_0,u,n}} \ge \mathfrak{m} D_{v_0,n,m}, \quad \forall u, \forall n, \forall m$$
 (33)

which materializes

$$\beta_{v_0,u,n,m} \le \alpha_{u,n,m}, \quad \forall u, \forall n, \forall m$$
 (34)

However, this implies that

$$\gamma_{\nu_0, u, n} = \alpha_{u, n}, \quad \forall u, \forall n \tag{35}$$

• Proof by contradiction: Assume  $\exists u: \exists n: \gamma_{v_0,u,n} > \alpha_{u,n}$ , then at that n:

$$\sum_{u} \gamma_{\nu_0, u, n} > \sum_{u} \alpha_{u, n} \tag{36}$$

$$1 > 1 \tag{37}$$

• But such an assumption materializes

$$\sum_{u} \gamma_{v_0, u, n} = \sum_{u} \alpha_{u, n} = 1, \quad \forall n$$
 (38)

• They introduce the notation of  $(1 - \delta)$  factor to argue away some fundamental issues.

- Such an assignment eliminates the URLLC constraints.
- The average rate is then

$$\bar{r}_{u} = \frac{1}{\mathfrak{u}} \sum_{n} \left( \alpha_{u,n} - \alpha_{u,n} f\left(\frac{\gamma_{v_{0},u,n} D_{v_{0},n}}{\alpha_{u,n}}\right) \right) \mathfrak{r}_{u,n}, \quad \forall u \quad (39)$$

• Finally, they hide the  $D_{v_0,n}$  from their problem formulation and state that it is a convex optimization problem.

$$\begin{aligned} \max_{\phi^{s},\gamma^{s}} & \sum_{u \in \mathcal{U}} U'_{u}\left(\overline{r}_{u}(t-1)\right) g_{u}^{s}(\phi_{u}^{s},\gamma_{u}^{s}), & (9) \\ \text{s.t.} & \phi^{s} \geq (1-\delta)\,\gamma^{s}, & (10) \\ & \sum_{u \in \mathcal{U}} \phi_{u}^{s} = 1 \text{ and } \sum_{u \in \mathcal{U}} \gamma_{u}^{s} = 1, & (11) \end{aligned}$$

$$u \in \mathcal{U} \qquad \qquad u \in \mathcal{U}$$

$$\phi^s \in [0, 1]^{|\mathcal{U}|} \text{ and } \gamma^s \in [0, 1]^{|\mathcal{U}|}. \tag{12}$$

Figure: Problem

#### **Issues**

- Multiple URLLC users scenario is neglected.
- Subchannel-wise rate is neglected.
- Physical resource block allocation of eMBB (in linear loss model analysis and convex loss model) and URLLC (in linear and convex loss models) is neglected.
- Superposition interference in convex loss model is neglected.
- Bandwidth 'fluidizability' in convex loss model is not guaranteed.
- Homogeneity in convex loss model is not guaranteed.

#### Notes

- They incorrectly address Pareto optimality.
- They incorrectly reference gradient algorithm's moving average rate update function (twice).
- They model channel states, but I think it was sufficient to describe channel states using time slot index.
- They hide the formulation of linear loss model.

#### References

- Alexander L. Stolyar. "On the Asymptotic Optimality of the Gradient Scheduling Algorithm for Multiuser Throughput Allocation". In: Operations Research 53.1 (2005), pp. 12–25. DOI: 10.1287/opre.1040.0156.
- [2] Hao Yin, Lyutianyang Zhang, and Sumit Roy. "Multiplexing URLLC Traffic Within eMBB Services in 5G NR: Fair Scheduling". In: IEEE Transactions on Communications 69.2 (2021), pp. 1080–1093. DOI: 10.1109/TCOMM.2020.3035582.