# NTUT\_Kn1ghts ICPC Team Notebook

#### **Contents**

1	Dynamic programming algorithms 1			
	1.1	Longest common subsequence	1	
	1.2	Max 2D range sum	2	
	1.3	0-1 Knapsack	2	
	1.4	Traveling salesman	2	
2	Geometry			
	2.1	Convex hull	2	
	2.2	Miscellaneous geometry	3	
3	Nun	nerical algorithms	5	
	3.1	Number theory (modular, Chinese remainder, linear Diophantine)	5	
	3.2	Systems of linear equations, matrix inverse, determinant	6	
	3.3	Reduced row echelon form, matrix rank	6	
	3.4	Fast Fourier transform	7	
	3.5	Simplex algorithm	7	
4	Graj	ph algorithms	3	
	4.1	Bellman-Ford shortest paths with negative edge weights	8	
	4.2	Fast Dijkstra's algorithm	9	
	4.3	Strongly connected components	9	
	4.4	Eulerian path	9	
	4.5	Kruskal's algorithm	9	
	4.6	Minimum spanning trees	O	
5	Data	a structures 10	)	
	5.1	Suffix array	0	
	5.2	Binary Indexed (Fenwick) Tree	1	
	5.3	Union-find disjoint sets	1	
	5.4	KD-tree	1	
	5.5	Splay tree	3	
	5.6	Segment tree	3	
	5.7	Lowest common ancestor	4	
6	Misc	cellaneous 15	5	
	6.1	Dijkstra's algorithm using struct	5	
	6.2	Longest increasing subsequence	5	
	6.3	Knuth-Morris-Pratt	6	
	6.4	Dates	6	
	6.5	Prime numbers	6	
	6.6	Latitude/longitude		
	6.7	C++ input/output	7	
	6.8	Random STL stuff		
	6.9	Regular expressions (Java)	7	

# 1 Dynamic programming algorithms

#### 1.1 Longest common subsequence

Calculates the length of the longest common subsequence of two vectors. Backtracks to find a single subsequence or all subsequences. Runs in O(m\*n) time except for finding all longest common subsequences, which may be slow depending on how many there are.

\*/

#include <iostream>
#include <set>
#include <alcording
#include <

```
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
void backtrack(VVI& dp, VT& res, VT& A, VT& B, int i, int j)
  if(!i || !j) return;
  if(A[i-1] == B[j-1]) { res.push_back(A[i-1]); backtrack(dp, res, A, B, i-1, j-1); }
  else
    if(dp[i][j-1] >= dp[i-1][j]) backtrack(dp, res, A, B, i, j-1);
    else backtrack(dp, res, A, B, i-1, j);
void backtrackall(VVI& dp, set<VT>& res, VT& A, VT& B, int i, int j)
  if(!i || !j) { res.insert(VI()); return; }
  if(A[i-1] == B[j-1])
    set<VT> tempres;
    backtrackall(dp, tempres, A, B, i-1, j-1);
    for(set<VT>::iterator it=tempres.begin(); it!=tempres.end(); it++)
      VT temp = *it:
      temp.push_back(A[i-1]);
      res.insert(temp);
    \label{eq:if} \textbf{if}(dp[i][j-1] >= dp[i-1][j]) \ backtrackall(dp, res, A, B, i, j-1);
    if(dp[i][j-1] \le dp[i-1][j]) backtrackall(dp, res, A, B, i-1, j);
VT LCS(VT& A, VT& B)
 int n = A.size(), m = B.size();
dp.resize(n+1);
  for (int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
  for(int i=1; i<=n; i++)</pre>
    for (int j=1; j<=m; j++)</pre>
      if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
  backtrack(dp, res, A, B, n, m);
  reverse(res.begin(), res.end());
  return res:
set<VT> LCSall(VT& A, VT& B)
  int n = A.size(), m = B.size();
  dp.resize(n+1);
  for(int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
  for(int i=1; i<=n; i++)</pre>
    for (int j=1; j<=m; j++)</pre>
      if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
  set<VT> res;
  backtrackall(dp, res, A, B, n, m);
  return res;
  int a[] = { 0, 5, 5, 2, 1, 4, 2, 3 }, b[] = { 5, 2, 4, 3, 2, 1, 2, 1, 3 };
  VI A = VI(a, a+8), B = VI(b, b+9);
  VI C = LCS(A, B);
  for(int i=0; i<C.size(); i++) cout << C[i] << " ";</pre>
  cout << endl << endl:
  set <VI> D = LCSall(A, B);
  for(set<VI>::iterator it = D.begin(); it != D.end(); it++)
    for(int i=0; i<(*it).size(); i++) cout << (*it)[i] << " ";</pre>
```

#### 1.2 Max 2D range sum

```
// This is a solution for UVa 108 - Maximum Sum. In this problem, we are given a
// square matrix of integers and are asked to find a sub-matrix with the maximum
// sum. This 1D DP + greedy (Kadane's) solution takes time O(n^3).
#include <bits/stdc++.h>
using namespace std;
#define MAX n 110
int A[MAX n][MAX n];
int main() {
 int n; scanf("%d", &n);
                                                  // square matrix size
  for (int i = 0; i < n; ++i)
   for (int j = 0; j < n; ++j) {
      scanf("%d", &A[i][j]);
      if (j > 0) A[i][j] += A[i][j-1];
                                                  // pre-processing
  int maxSubRect = -127 * 100 * 100;
                                                  // lowest possible val
  for (int 1 = 0; 1 < n; ++1)
    for (int r = 1; r < n; ++r) {
     int subRect = 0:
      for (int row = 0; row < n; ++row) {</pre>
        // Max 1D Range Sum on columns of this row
        if (1 > 0) subRect += A[row][r] - A[row][1-1];
        else      subRect += A[row][r];
         // Kadane's algorithm on rows
        if (subRect < 0) subRect = 0;</pre>
                                                  // restart if negative
        maxSubRect = max(maxSubRect, subRect);
  printf("%d\n", maxSubRect);
  return 0:
```

### 1.3 0-1 Knapsack

```
// This is a solution for UVa 10130 - SuperSale. This is a 0-1 Knapsack variant
// solved using top-down DP O(n*W).
#include <bits/stdc++.h>
using namespace std;
const int MAX_N = 1010;
const int MAX_W = 40;
int N, V[MAX_N], W[MAX_N], memo[MAX_N][MAX_W];
int dp(int id, int remW) {
 if ((id == N) || (remW == 0)) return 0;
                                                 // two base cases
  int &ans = memo[id][remW];
  if (ans != -1) return ans;
                                                 // computed before
  if (W[id] > remW) return ans = dp(id+1, remW); // no choice, skip
  return ans = max(dp(id+1, remW),
                                                // has choice, skip
                   V[id]+dp(id+1, remW-W[id])); // or take
int main() {
  int T; scanf("%d", &T);
  while (T--) {
   memset(memo, -1, sizeof memo);
    scanf("%d", &N);
    for (int i = 0; i < N; ++i)
     scanf("%d %d", &V[i], &W[i]);
    int ans = 0;
   int G; scanf("%d", &G);
    while (G--) {
     int MW: scanf("%d", &MW);
      ans += dp(0, MW):
    printf("%d\n", ans);
  return 0:
```

### 1.4 Traveling salesman

```
// This is a solution for UVa 10496 - Collecting Beepers. The problem is a
// variant of the Traveling Salesman Problem (TSP): Given n cities and their
// pairwise distances in the form of a matrix 'dist' of size n \star n, compute the
// minimum cost of making a tour that starts from any city s, goes through all
// the other n - 1 cities exactly once, and finally returns to the city s. In
// this case, the salesman is Karel in a 2D world who can only move along the
// x and y axis. The cities are beepers whose coordinates are given, from which
// pairwise distances can be calculated. Algorithm takes time O(2^n \, \star \, n^2).
// INPUT: The first line is the number of test cases. The first line of each
// test case is world's size (x-size and y-size). Next is the starting position
// of Karel. Next is the number of beepers. Next are the beepers' x- and y-
// coordinates.
// OUTPUT: For each test case, output the minimum distance to move from Karel's
// starting position to each of the beepers and back to the starting position.
#include <bits/stdc++.h>
using namespace std;
#define LSOne(S) ((S) & -(S))
const int MAX n = 11:
int dist[MAX_n][MAX_n], memo[MAX_n][1<<(MAX_n-1)]; // Karel + max 10 beepers</pre>
int dp(int u, int mask) {
                                                  // mask = free coordinates
  if (mask == 0) return dist[u][0];
                                                  // close the loop
  int &ans = memo[u][mask];
  if (ans != -1) return ans;
                                                  // computed before
  ans = 2000000000;
  int m = mask;
  while (m) {
                                                  // up to O(n)
   int two_pow_v = LSOne(m);
                                                  // but this is fast
    int v = __builtin_ctz(two_pow_v)+1;
                                                  // offset v by +1
    ans = min(ans, dist[u][v] + dp(v, mask^two_pow_v)); // keep the min
    m -= two_pow_v;
  return ans;
int main() {
  int TC; scanf("%d", &TC);
  while (TC--) {
   int xsize, ysize; scanf("%d %d", &xsize, &ysize); // these two values are not used int x[MAX_n], y[MAX_n]; scanf("%d %d", &x[0], &y[0]); int n; scanf("%d", &n); ++n; // include Karel
    for (int i = 1; i < n; ++i)
                                                  // Karel is at index 0
      scanf("%d %d", &x[i], &y[i]);
    memset (memo, -1, sizeof memo);
    printf("The shortest path has length d^n, dp(0, (1 << (n-1))-1)); // DP-TSP
  return 0:
```

# 2 Geometry

#### 2.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// #defined.
//
// Running time: O(n log n)
//
// INPUT: a vector of input points, unordered.
// OUTPUT: a vector of points in the convex hull, counterclockwise, starting
// with bottommost/leftmost point

#include <csstdio>
#include <cassert>
#include <algorithm>
#include <cmath>
// BEGIN CUT
#include <map>
// END CUT
```

```
#define REMOVE REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
  Т х, у;
  PT() {}
  PT(T x, T y) : x(x), y(y) {}
  bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y,rhs.x); }</pre>
 bool operator==(const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2 (PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE_REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
  return (fabs(area2(a,b,c)) < EPS && (a.x-b.x) \star (c.x-b.x) <= 0 && (a.y-b.y) \star (c.y-b.y) <= 0);
#endif
void ConvexHull(vector<PT> &pts) {
  sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {
    while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();</pre>
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  pts = dn;
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE_REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {
   if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();</pre>
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
  pts = dn:
#endif
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
int main() {
  int t:
  scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {</pre>
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);</pre>
    vector<PT> h(v);
    map<PT, int> index;
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    ConvexHull(h):
    double len = 0:
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
    printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {
  if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
```

using namespace std;

#### 2.2 Miscellaneous geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100:
double EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT (double x, double y) : x(x), y(y) {}
  PT (const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y);
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y);
  PT operator * (double c)
                                 const { return PT(x*c, y*c );
  PT operator / (double c)
                                 const { return PT(x/c, y/c ); }
double dot(PT p, PT q)
                             { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x+q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
    return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCW90 (PT p)
                        { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
 return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                            double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
    dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
      return false;
    return true;
```

```
if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
 b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// \ {\tt integer} \ {\tt arithmetic} \ {\tt by} \ {\tt taking} \ {\tt care} \ {\tt of} \ {\tt the} \ {\tt division} \ {\tt appropriately}
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||
      p[j].y \le q.y && q.y < p[i].y) &&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
      return true;
    return false:
// compute intersection of line through points a and b with
// circle centered at c with radius r >
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
  h = h-a:
  a = a-c:
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r * r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
// compute intersection of circle centered at a with radius \boldsymbol{r}
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R || d+min(r, R) < max(r, R)) return ret;
double x = (d*d-R*R*r*r)/(2*d);</pre>
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or // counterclockwise fashion. Note that the centroid is often known as
   the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
```

```
double ComputeArea(const vector<PT> &p) {
  return fabs (ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++){}
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
 // tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++)</pre>
    for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
int l = (k+1) % p.size();
      if (i == 1 \mid \mid j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false:
  return true:
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "</pre>
       << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "</pre>
        << ProjectPointSegment (PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
  // expected: 1 0 1
  cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
<< LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;</pre>
  // expected: 0 0 1
  cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
       << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
        << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
        << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
        << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
        << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
  // expected: (1,2)
  cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
  // expected: (1,1)
  cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
  vector<PT> v;
  v.push_back(PT(0,0));
  v.push_back(PT(5,0));
  v.push_back(PT(5,5));
  v.push_back(PT(0,5));
  // expected: 1 1 1 0 0
  cerr << PointInPolygon(v, PT(2,2)) << " '
       << PointInPolygon(v, PT(2,0)) << " "
        << PointInPolygon(v, PT(0,2)) << " "
       << PointInPolygon(v, PT(5,2)) << " "
<< PointInPolygon(v, PT(2,5)) << endl;</pre>
  // expected: 0 1 1 1 1
  cerr << PointOnPolygon(v, PT(2,2)) << " "
        << PointOnPolygon(v, PT(2,0)) << " "
        << PointOnPolygon(v, PT(0,2)) << " "
        << PointOnPolygon(v, PT(5,2)) << " "
```

```
<< PointOnPolygon(v, PT(2,5)) << endl;
                  (5,4) (4,5)
                  blank line
                  (4,5) (5,4)
                  blank line
                  (4,5) (5,4)
ior (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << end];
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << end];
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << end];
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << end];</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl:
return 0:
```

## 3 Numerical algorithms

# 3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
       return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
// computes lcm(a,b)
int lcm(int a, int b)
       return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1;
        return ret;
// returns q = qcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
        int yy = x = 1;
```

```
while (b) {
                  int q = a / b;
                  int t = b; b = a%b; a = t;
                  t = xx; xx = x - q*xx; x = t;
                  t = yy; yy = y - q * yy; y = t;
        return a;
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
        int x, y;
        VI ret;
        int g = extended_euclid(a, n, x, y);
        if (! (b%g)) {
    x = mod(x*(b / g), n);
                  for (int i = 0; i < q; i++)
                          ret.push_back(mod(x + i*(n / g), n));
        return ret;
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
if (g > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g);
// Chinese remainder theorem: find z such that
// z \circledast m[i] = r[i] for all i. Note that the solution is // unique modulo M = lcm_i (m[i]). Return (z, M). On // failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
        PII ret = make_pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {</pre>
                  ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
                  if (ret.second == -1) break;
        return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
        if (!a && !b)
                  if (c) return false;
                  x = 0; y = 0;
                  return true:
        if (!a)
                 if (c % b) return false;
                  x = 0; y = c / b;
                  return true;
        if (!b)
                 if (c % a) return false;
                  x = c / a; v = 0;
                  return true;
        int g = gcd(a, b);
        if (c % g) return false;
        x = c / g * mod_inverse(a / g, b / g);
         y = (c - a * x) / b;
         return true;
int main() {
        // expected: 2
        cout << gcd(14, 30) << endl;
         // expected: 2 -2 1
        int x, y;
int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
```

// expected: 95 451

# 3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
    (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
// Running time: O(n^3)
              a[][] = an nxn matrix
              b[][] = an nxm matrix
// OUTPUT: X
                     = an nxm matrix (stored in b[][])
              A^{-1} = an nxn matrix (stored in a[][])
              returns determinant of a[1[1
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T det = 1;
  for (int i = 0; i < n; i++) {</pre>
    int pj = -1, pk = -1;
for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
      for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
    if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre>
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
if (pj != pk) det *= -1;
irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c;
    for (int p = 0; p < m; p++) b[pk][p] *= c;
for (int p = 0; p < n; p++) if (p != pk) {
       c = a[p][pk];
      a[p][pk] = 0;
       for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
       for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
```

```
return det;
int main() {
  const int n = 4;
  const int m = 2;
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \}; double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);</pre>
  double det = GaussJordan(a, b);
   // expected: 60
  cout << "Determinant: " << det << endl;
   // expected: -0.233333 0.166667 0.133333 0.0666667
                   0.166667 0.166667 0.333333 -0.333333
                   0.233333 0.833333 -0.133333 -0.0666667
                   0 05 -0 75 -0 1 0 2
  cout << "Inverse: " << endl;</pre>
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++)
  cout << a[i][j] << ' ';</pre>
     cout << endl:
   // expected: 1.63333 1.3
                  -0.166667 0.5
                   2.36667 1.7
                   -1.85 -1.35
  cout << "Solution: " << endl;</pre>
  for (int i = 0; i < n; i++) {</pre>
    for (int j = 0; j < m; j++)
cout << b[i][j] << ' ';</pre>
     cout << endl;
```

#### 3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT: a[][] = an nxm matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
             returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
 int n = a.size();
  int m = a[0].size();
  int r = 0;
  for (int c = 0; c < m && r < n; c++) {
    int j = r;
    for (int i = r + 1; i < n; i++)
     if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
for (int i = 0; i < n; i++) if (i != r) {</pre>
      T t = a[i][c];
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];</pre>
```

```
return r;
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
    {16, 2, 3, 13},
    { 5, 11, 10, 8},
    { 9, 7, 6, 12},
    { 4, 14, 15, 1},
    {13, 21, 21, 13}};
  VVT a(n);
  for (int i = 0; i < n; i++)
   a[i] = VT(A[i], A[i] + m);
  int rank = rref(a);
  // expected: 3
  cout << "Rank: " << rank << endl;
  // expected: 1 0 0 1
                0 1 0 3
                 0 0 1 -3
                 0 0 0 3 108626-15
                0 0 0 2 22045e-15
  cout << "rref: " << endl:
  for (int i = 0; i < 5; i++) {
  for (int j = 0; j < 4; j++)
    cout << a[i][j] << ' ';</pre>
    cout << endl;
```

#### 3.4 Fast Fourier transform

```
#include <cassert>
#include <cstdio>
#include <cmath>
struct cpx
  cpx(){}
  cpx(double aa):a(aa),b(0){}
  cpx(double aa, double bb):a(aa),b(bb){}
  double a;
  double b;
  double modsq(void) const
    return a * a + b * b;
  cpx bar(void) const
    return cpx(a, -b);
1:
cpx operator +(cpx a, cpx b)
  return cpx(a.a + b.a, a.b + b.b);
cpx operator *(cpx a, cpx b)
  return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
cpx operator / (cpx a, cpx b)
  cpx r = a * b.bar():
  return cpx(r.a / b.modsq(), r.b / b.modsq());
cpx EXP (double theta)
  return cpx(cos(theta), sin(theta));
const double two_pi = 4 * acos(0);
// in:
            input array
// out:
            output array
// step: {SET TO 1} (used internally)
// size: length of the input/output {MUST BE A POWER OF 2}
           either plus or minus one (direction of the FFT)
// dir:
/// RESULT: \operatorname{out}[k] = \sum_{j=0}^{n} \{\operatorname{size} - 1\} \operatorname{in}[j] * \exp(\operatorname{dir} * 2\operatorname{pi} * i * j * k / \operatorname{size})
void FFT(cpx *in, cpx *out, int step, int size, int dir)
```

```
if(size < 1) return;</pre>
  if(size == 1)
    out[0] = in[0];
    return;
  FFT(in, out, step * 2, size / 2, dir);
  FFT(in + step, out + size / 2, step \star 2, size / 2, dir);
  for(int i = 0; i < size / 2; i++)
    cpx even = out[i];
   cpx odd = out[i + size / 2];
out[i] = even + EXP(dir * two pi * i / size) * odd;
    out[i + size / 2] = even + EXP(dir * two_pi * (i + size / 2) / size) * odd;
// Usage:
// f[0...N-1] and g[0..N-1] are numbers
// Want to compute the convolution h, defined by
// h[n] = sum \ of \ f[k]g[n-k] \ (k = 0, ..., N-1).
// Here, the index is cyclic; f[-1] = f[N-1], f[-2] = f[N-2], etc.
// Let F[0...N-1] be FFT(f), and similarly, define G and H.
// The convolution theorem says H[n] = F[n]G[n] (element-wise product).
// To compute h[] in O(N log N) time, do the following:

    Compute F and G (pass dir = 1 as the argument).

    2. Get H by element-wise multiplying F and G.
    3. Get h by taking the inverse FFT (use dir = -1 as the argument)
        and *dividing by N*. DO NOT FORGET THIS SCALING FACTOR.
int main (void)
 printf("If rows come in identical pairs, then everything works.\n");
  cpx \ a[8] = \{0, 1, cpx(1,3), cpx(0,5), 1, 0, 2, 0\};
  cpx b[8] = \{1, cpx(0,-2), cpx(0,1), 3, -1, -3, 1, -2\};
  cpx B[8];
  FFT(a, A, 1, 8, 1);
  FFT(b, B, 1, 8, 1);
  for(int i = 0; i < 8; i++)
    printf("%7.21f%7.21f", A[i].a, A[i].b);
  printf("\n");
  for(int i = 0; i < 8; i++)
    cpx Ai(0,0);
    for (int j = 0; j < 8; j++)
      Ai = Ai + a[j] * EXP(j * i * two_pi / 8);
    printf("%7.21f%7.21f", Ai.a, Ai.b);
  printf("\n");
  cpx AB[8];
  for(int i = 0; i < 8; i++)
AB[i] = A[i] * B[i];
  cpx aconvb[8];
  FFT(AB, aconvb, 1, 8, -1);
  for (int i = 0; i < 8; i++)
    aconvb[i] = aconvb[i] / 8;
  for (int i = 0; i < 8; i++)
    printf("%7.21f%7.21f", aconvb[i].a, aconvb[i].b);
  printf("\n");
  for (int i = 0; i < 8; i++)
    cpx aconvbi(0,0);
    for (int j = 0; j < 8; j++)
      aconvbi = aconvbi + a[j] * b[(8 + i - j) % 8];
    printf("%7.21f%7.21f", aconvbi.a, aconvbi.b);
 printf("\n");
  return 0;
```

## 3.5 Simplex algorithm

// Two-phase simplex algorithm for solving linear programs of the form  $^{\prime\prime}$ 

```
maximize
        subject to Ax <= b
                       x >= 0
// INPUT: A -- an m x n matrix
           b -- an m-dimensional vector
            c -- an n-dimensional vector
            x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
            above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m. n:
   \mbox{LPSolver} \mbox{(const VVD \&A, const VD \&b, const VD \&c)} \ : \\
    m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
    for (int i=0, i < m, i++) for (int j=0, j < n, j++) D[i][j] = A[i][j]; for (int i=0, i < m, i++) { B[i] = n + i; D[i][n] = -1, D[i][n + 1] = b[i]; } for (int j=0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
    N[n] = -1; D[m + 1][n] = 1;
  void Pivot(int r, int s) {
    double inv = 1.0 / D[r][s];
for (int i = 0; i < m + 2; i++) if (i != r)</pre>
    for (int i = 0; i < m + 2; i++) if (j != s)
   D[i][j] -= D[r][j] * D[i][s] * inv;
for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
for (int i = 0; i < m + 2; i++) if (j != s) D[i][s] *= -inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
  bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
    while (true) {
      int s = -1;
      for (int j = 0; j <= n; j++) {
  if (phase == 2 && N[j] == -1) continue;</pre>
         if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] && N[j] < N[s]) s = j;
       if (D[x][s] > -EPS) return true;
       int r = -1;
       for (int i = 0; i < m; i++) {
         if (D[i][s] < EPS) continue;</pre>
         if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
            if (r == -1) return false:
      Pivot(r, s);
  DOUBLE Solve(VD &x) {
    for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n + 1] < -EPS) {
       Pivot(r, n);
       if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -numeric_limits<DOUBLE>::infinity();
       for (int i = 0; i < m; i++) if (B[i] == -1) {
         int s = -1;
         int s = 1,
for (int j = 0; j <= n; j++)
   if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i][s] && N[j] < N[s]) s = j;</pre>
         Pivot(i, s):
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
     x = VD(n);
     for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return D[m][n + 1];
```

```
};
int main() {
  const int m = 4;
  const int n = 3;
  DOUBLE _A[m][n] =
    { 6, -1, 0 },
     \{-1, -5, 0\},
    { 1, 5, 1 },
    \{-1, -5, -1\}
  DOUBLE _b[m] = { 10, -4, 5, -5 };
  DOUBLE _c[n] = { 1, -1, 0 };
  VD b(\underline{b}, \underline{b} + m);
   VD c(_c, _c + n);
  for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);</pre>
  LPSolver solver (A, b, c);
  DOUBLE value = solver.Solve(x):
  cerr << "VALUE: " << value << endl; // VALUE: 1.29032
cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1</pre>
  for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
  cerr << endl:
  return 0:
```

# 4 Graph algorithms

# 4.1 Bellman-Ford shortest paths with negative edge weights

```
// This function runs the Bellman-Ford algorithm for single source
// shortest paths with negative edge weights. The function returns
// false if a negative weight cycle is detected. Otherwise, the
// function returns true and dist[i] is the length of the shortest
// path from start to i.
// Running time: O(|V|^3)
    INPUT: start, w[i][j] = cost \ of \ edge \ from \ i \ to \ j
    OUTPUT: dist[i] = min weight path from start to i
             prev[i] = previous node on the best path from the
                        start node
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool BellmanFord (const VVT &w, VT &dist, VI &prev, int start) {
 int n = w.size();
  prev = VI(n, -1);
  dist = VT(n, 1000000000);
  dist[start] = 0;
  for (int k = 0; k < n; k++) {
   if (k == n-1) return false;
dist[j] = dist[i] + w[i][j];
         prev[j] = i;
  return true;
```

### 4.2 Fast Dijkstra's algorithm

```
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
// Running time: O(|E| log |V|)
#include <queue>
#include <cstdio>
using namespace std;
const int INF = 2000000000;
typedef pair<int, int> PII;
int main() {
        int N, s, t;
scanf("%d%d%d", &N, &s, &t);
        vector<vector<PII> > edges(N);
        for (int i = 0; i < N; i++) {
                int M;
                scanf("%d", &M);
for (int j = 0; j < M; j++) {
                         int vertex, dist;
                         scanf("%d%d", &vertex, &dist);
                         edges[i].push_back(make_pair(dist, vertex)); // note order of arguments here
         // use priority queue in which top element has the "smallest" priority
        priority_queue<PII, vector<PII>, greater<PII> > Q;
        vector<int> dist(N, INF), dad(N, -1);
        Q.push(make_pair(0, s));
dist[s] = 0;
        while (!Q.empty()) {
                PII p = Q.top();
                Q.pop();
                int here = p.second;
                if (here == t) break;
                if (dist[here] != p.first) continue;
                for (vector<PII>::iterator it = edges[here].begin(); it != edges[here].end(); it++) {
                         if (dist[here] + it->first < dist[it->second]) {
                                 dist[it->second] = dist[here] + it->first;
                                 dad[it->second] = here;
                                 Q.push(make_pair(dist[it->second], it->second));
        printf("%d\n", dist[t]);
        if (dist[t] < INF)</pre>
                for (int i = t; i != -1; i = dad[i])
                         printf("%d%c", i, (i == s ? '\n' : ' '));
Sample input:
5 0 4
2 1 2 3 1
2 2 4 4 5
3 1 4 3 3 4 1
20123
2 1 5 2 1
Expected:
4 2 3 0
```

# 4.3 Strongly connected components

```
#include<memory.h>
struct edge{int e, nxt;};
int V, E;
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group_cnt, group_num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill_forward(int x)
{
   int i;
}
```

```
v[x]=true;
for(i=sp[x];i;i=e[i].nxt) if(!v[e[i].e]) fill_forward(e[i].e);
stk[+stk[0]]=x;
}
void fill_backward(int x)
{
   int i;
   v[x]=false;
   group_num[x]=group_cnt;
   for(i=sp[x];i:i=er[i].nxt) if(v[er[i].e]) fill_backward(er[i].e);
}
void add_edge(int v1, int v2) //add edge v1->v2
{
   e [++E].e=v2; e [E].nxt=sp [v1]; sp [v1]=E;
   er[ E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
}
void SCC()
{
   int i;
   stk[0]=0;
   memset (v, false, sizeof(v));
   for(i=1;i<=v;i++) if(!v[i]) fill_forward(i);
   group_cnt=0;
   for(i=stk[0];i>=1;i--) if(v[stk[i]]) (group_cnt++; fill_backward(stk[i]);)
for(i=stk[0];i>=1;i--) if(v[stk[i]]) (group_cnt++; fill_backward(stk[i]);)
```

#### 4.4 Eulerian path

```
struct Edge;
typedef list<Edge>::iterator iter;
struct Edge
        int next vertex:
        iter reverse edge:
        Edge(int next_vertex)
                :next_vertex(next_vertex)
};
const int max_vertices = ;
int num_vertices;
list<Edge> adj[max_vertices];
                                        // adjacency list
vector<int> path;
void find_path(int v)
        while(adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().reverse_edge);
                adj[v].pop_front();
        path.push_back(v);
void add edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse_edge = itb;
        itb->reverse_edge = ita;
```

## 4.5 Kruskal's algorithm

```
// This is a solution for UVa 11747 - Heavy Cycle Edges. Abridged problem
// description: Undirected graph. Generate Minimum Spanning Tree (MST).
#include <bits/stdc++.h>
#define LOCAL
#define 11 long long
using namespace std;
int parent[1020];
```

```
struct edge {
    11 n1, n2, w;
} node[25020];
int compare(edge A, edge B) {
    return A.w < B.w;
int find_root(int a) {
    if (a != parent[a])
        return parent[a] = find_root(parent[a]);
    return a;
int main() {
#ifdef LOCAL
    freopen("in1.txt", "r", stdin);
freopen("out.txt", "w", stdout);
#endif
    int n, m, p_n1, p_n2; // parent_n1 , parent_n2
    vector<int> hce;  //heavy edge circle
    while (cin >> n >> m && n + m != 0) {
        for (int i = 0; i < m; i++) {</pre>
            cin >> node[i].n1 >> node[i].n2 >> node[i].w;
        for (int i = 0; i < n; i++)</pre>
            parent[i] = i:
        sort(node, node + m, compare);
        hce.clear():
        for (int i = 0; i < m; i++) {
            p_n1 = find_root(node[i].n1);
p_n2 = find_root(node[i].n2);
             if (p_n1 != p_n2)
                parent[p_n2] = p_n1;
            else
                hce.push_back(node[i].w);
             //debug
             /**<
             for(int i = 0 ; i < n ; i++)
                cout << parent[i] << ' ';
             cout << '\n';
         sort(hce.begin(), hce.end());
        if (hce.size()) {
             for (int i = 0; i < hce.size() - 1; i++)</pre>
                cout << hce[i] << ' ';
             cout << hce[hce.size() - 1];</pre>
        else
            cout << "forest":
        cout << '\n';
    return 0:
```

## 4.6 Minimum spanning trees

```
// This function runs Prim's algorithm for constructing minimum
// weight spanning trees.
// Running time: O(|V|^2)
    INPUT: w[i][j] = cost \ of \ edge \ from \ i \ to \ j
              NOTE: Make sure that w[i][j] is nonnegative and
              symmetric. Missing edges should be given -1
              weight.
    OUTPUT: edges = list of pair<int,int> in minimum spanning tree
             return total weight of tree
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std:
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
```

```
typedef vector<VI> VVI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
T Prim (const VVT &w, VPII &edges) {
  int n = w.size();
  VI found (n);
  VI prev (n, -1);
  VT dist (n, 1000000000);
  int here = 0;
  dist[here] = 0;
  while (here !=-1) {
    found[here] = true;
    int best = -1:
    for (int k = 0; k < n; k++) if (!found[k]) {</pre>
       if (w[here][k] != -1 && dist[k] > w[here][k]){
         dist[k] = w[here][k];
         prev[k] = here;
       if (best == -1 || dist[k] < dist[best]) best = k;</pre>
    here = best:
  T tot_weight = 0;
  for (int i = 0; i < n; i++) if (prev[i] != -1) {
  edges.push_back (make_pair (prev[i], i));</pre>
    tot_weight += w[prev[i]][i];
  return tot_weight;
int main(){
  int ww[5][5] = {
    {0, 400, 400, 300, 600},
     {400, 0, 3, -1, 7},
    {400, 3, 0, 2, 0},
    {300, -1, 2, 0, 5},
    {600, 7, 0, 5, 0}
  VVT w(5, VT(5));
  for (int i = 0; i < 5; i++)
for (int j = 0; j < 5; j++)
w[i][j] = ww[i][j];
  // expected: 305
                 0 3
  VPII edges:
  cout << Prim (w, edges) << endl;</pre>
  for (int i = 0; i < edges.size(); i++)
  cout << edges[i].first << " " << edges[i].second << endl;</pre>
```

## 5 Data structures

### 5.1 Suffix array

```
// Suffix array construction in O(L log^2 L) time. Routine for // computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
             of substring s[i...L-1] in the list of sorted suffixes.
             That is, if we take the inverse of the permutation suffix[],
             we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
struct SuffixArray {
  const int L;
  string s;
  vector<vector<int> > P;
  vector<pair<int,int>,int> > M;
```

```
SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);</pre>
    for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
      P.push_back(vector<int>(L, 0));
       for (int i = 0; i < L; i++)
        M[i] = make\_pair(make\_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
       sort(M.begin(), M.end());
          P[level] [M[i].second] = (i > 0 \& M[i].first == M[i-1].first) ? P[level] [M[i-1].second] : i; 
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
      if (P[k][i] == P[k][j]) {
        i += 1 << k;
          j += 1 << k;
         len += 1 << k:
    return len:
1:
// BEGIN CUT
   The following code solves UVA problem 11512: GATTACA.
#ifdef TESTING
int main() {
  int T;
  cin >> T;
  for (int caseno = 0; caseno < T; caseno++) {</pre>
    cin >> s;
    SuffixArray array(s);
   sutinxArray array(s);
vector<int> v = array.GetSuffixArray();
int bestlen = -1, bestpos = -1, bestcount = 0;
for (int i = 0; i < s.length(); i++) {
  int len = 0, count = 0;
  for (int j = i+1; j < s.length(); j++) {
    int l = array.LongestCommonPrefix(i, j);
    if() = length() = length(); j++</pre>
         if (1 >= len) {
           if (1 > len) count = 2; else count++;
       if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen) > s.substr(i, len)) {
         bestlen = len;
         bestcount = count:
         bestpos = i:
    if (bestlen == 0) {
      cout << "No repetitions found!" << endl;</pre>
      cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;</pre>
#else
// END CUT
int main() {
  // bobocel is the O'th suffix
  // obocel is the 5'th suffix
       bocel is the 1'st suffix
        ocel is the 6'th suffix
          cel is the 2'nd suffix
          el is the 3'rd suffix
            1 is the 4'th suffix
  SuffixArray suffix("bobocel");
  vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
  for (int i = 0; i < v.size(); i++) cout << v[i] << " ";</pre>
  cout << endl:
  cout << suffix.LongestCommonPrefix(0, 2) << endl:
#endif
// END CUT
```

#### 5.2 Binary Indexed (Fenwick) Tree

```
#include <iostream>
using namespace std;
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
  while (x \le N)
   tree[x] += v;
   x += (x & -x);
// get cumulative sum up to and including x
int get(int x) {
  while(x) {
   res += tree[x];
   x -= (x & -x);
  return res;
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
  int idx = 0, mask = N;
  while (mask && idx < N)
    int t = idx + mask;
    if(x >= tree[t]) {
     idx = t;
      x -= tree[t];
    mask >>= 1;
  return idx:
```

#### 5.3 Union-find disjoint sets

```
#include <iostream>
#include <vector>
using namespace std;
struct UnionFind {
    vector\into C;
    UnionFind(\int n) : C(n) { for (\int i = 0; i < n; i++) C[i] = i; }
    int find(\int x) { return (C[x] == x) ? x : C[x] = find(C[x]); }
    void merge(\int x, \int y) { C[find(x)] = find(y); }
};
int main() {
    int n = 5;
    UnionFind uf(n);
    uf.merge(0, 2);
    uf.merge(1, 0);
    uf.merge(3, 4);
    for (\int i = 0; i < n; i++) cout << i << " " << uf.find(i) << endl;
    return 0;
}</pre>
```

## 5.4 KD-tree

```
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
    ntype x, y;
    point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
bool operator==(const point &a, const point &b)
    return a.x == b.x && a.y == b.y;
// sorts points on x-coordinate
bool on_x(const point &a, const point &b)
    return a.x < b.x:
  sorts points on v-coordinate
bool on_y(const point &a, const point &b)
    return a.y < b.y;
// squared distance between points
ntype pdist2(const point &a, const point &b)
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
// bounding box for a set of points
struct bbox
    ntype x0, x1, y0, y1;
    bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
    // computes bounding box from a bunch of points
    void compute(const vector<point> &v) {
        for (int i = 0; i < v.size(); ++i) {</pre>
            x0 = min(x0, v[i].x); x1 = max(x1, v[i].x);
            y0 = min(y0, v[i].y); y1 = max(y1, v[i].y);
    // squared distance between a point and this bbox, 0 if inside
    ntype distance(const point &p) {
        if (p.x < x0) {
            if (p.y < y0)
                                return pdist2(point(x0, y0), p);
            else if (p.y > y1) return pdist2(point(x0, y1), p);
                                return pdist2(point(x0, p.y), p);
        else if (p.x > x1) {
           if (p.y < y0)
                                return pdist2(point(x1, y0), p);
            else if (p.y > y1) return pdist2(point(x1, y1), p);
            else
                                return pdist2(point(x1, p.y), p);
        else (
                                return pdist2(point(p.x, y0), p);
           if (p.v < v0)
            else if (p.y > y1) return pdist2(point(p.x, y1), p);
            else
                                return 0;
// stores a single node of the kd-tree, either internal or leaf
struct kdnode
    bool leaf:
                    // true if this is a leaf node (has one point)
    point pt;
                    // the single point of this is a leaf
    bbox bound:
                   // bounding box for set of points in children
    kdnode *first, *second; // two children of this kd-node
    kdnode() : leaf(false), first(0), second(0) {}
    "kdnode() { if (first) delete first; if (second) delete second; }
    // intersect a point with this node (returns squared distance)
    ntype intersect (const point &p) {
```

```
return bound.distance(p);
    // recursively builds a kd-tree from a given cloud of points
    void construct(vector<point> &vp)
        // compute bounding box for points at this node
        bound.compute(vp);
        // if we're down to one point, then we're a leaf node
        if (vp.size() == 1) {
            leaf = true:
            pt = vp[0];
        else
               split on x if the bbox is wider than high (not best heuristic...)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on_x);
            // otherwise split on y-coordinate
            else
                sort(vp.begin(), vp.end(), on_y);
            // divide by taking half the array for each child
            // (not best performance if many duplicates in the middle)
            int half = vp.size()/2;
            vector<point> vl(vp.begin(), vp.begin()+half);
            vector<point> vr(vp.begin()+half, vp.end());
            first = new kdnode(); first->construct(v1);
            second = new kdnode(); second->construct(vr);
   }
};
// simple kd-tree class to hold the tree and handle queries
struct kdtree
    kdnode *root;
    // constructs a kd-tree from a points (copied here, as it sorts them)
    kdtree(const vector<point> &vp) {
        vector<point> v(vp.begin(), vp.end());
        root = new kdnode():
        root =>construct (v):
    "kdtree() { delete root; }
    // recursive search method returns squared distance to nearest point
    ntype search(kdnode *node, const point &p)
        if (node->leaf) {
            // commented special case tells a point not to find itself
              if (p == node->pt) return sentry;
              PISP
                return pdist2(p, node->pt);
        ntype bfirst = node->first->intersect(p);
        ntype bsecond = node->second->intersect(p);
        // choose the side with the closest bounding box to search first
           (note that the other side is also searched if needed)
        if (bfirst < bsecond) {</pre>
            ntype best = search(node->first, p);
            if (bsecond < best)</pre>
                best = min(best, search(node->second, p));
            return best;
        else (
            ntype best = search(node->second, p);
            if (bfirst < best)</pre>
                best = min(best, search(node->first, p));
            return best;
    // squared distance to the nearest
    ntype nearest (const point &p) {
        return search (root, p);
};
// some basic test code here
int main()
    // generate some random points for a kd-tree
    vector<point> vp;
    for (int i = 0; i < 100000; ++i) {
        vp.push_back(point(rand()%100000, rand()%100000));
    kdtree tree(vp);
```

```
// query some points
for (int i = 0; i < 10; ++i) {
    point q(rand()*100000, rand()*100000);
        cout << "Closest squared distance to (" << q.x << ", " << q.y << ")"
        << " is " << tree.nearest(q) << endl;
}
return 0;
}</pre>
```

### 5.5 Splay tree

```
#include <cstdio>
#include <algorithm>
using namespace std;
const int N_MAX = 130010;
const int oo = 0x3f3f3f3f3f;
struct Node
  Node *ch[2], *pre;
  int val, size;
  bool isTurned:
} nodePool[N_MAX], *null, *root;
Node *allocNode(int val)
  static int freePos = 0;
  Node *x = &nodePool[freePos ++];
  x->val = val, x->isTurned = false;
  x->ch[0] = x->ch[1] = x->pre = null;
  x->size = 1;
  return x:
inline void update(Node *x)
  x->size = x->ch[0]->size + x->ch[1]->size + 1;
inline void makeTurned(Node *x)
  if(x == null)
    return;
  swap(x->ch[0], x->ch[1]);
x->isTurned ^= 1;
inline void pushDown(Node *x)
  if(x->isTurned)
    makeTurned(x->ch[0]);
    makeTurned(x->ch[1]);
    x->isTurned ^= 1;
inline void rotate(Node *x, int c)
  Node *y = x->pre;
  x->pre = y->pre;
if(y->pre != null)
    y->pre->ch[y == y->pre->ch[1]] = x;
   y->ch[!c] = x->ch[c];
  if (x->ch[c] != null)
   x->ch[c]->pre = y;
  x->ch[c] = y, y->pre = x;
  update(y);
  if(y == root)
    root = x;
void splay(Node *x, Node *p)
  while (x->pre != p)
    if(x->pre->pre == p)
      rotate(x, x == x->pre->ch[0]);
    else
      Node \star y = x->pre, \star z = y->pre;
      if(y == z \rightarrow ch[0])
```

```
if(x == y->ch[0])
          rotate(y, 1), rotate(x, 1);
          rotate(x, 0), rotate(x, 1);
        if(x == y->ch[1])
          rotate(y, 0), rotate(x, 0);
          rotate(x, 1), rotate(x, 0);
  update(x);
void select(int k, Node *fa)
  Node *now = root;
  while (1)
    pushDown(now);
    int tmp = now->ch[0]->size + 1;
if(tmp == k)
     break:
    else if(tmp < k)</pre>
     now = now \rightarrow ch[1], k \rightarrow tmp;
    else
      now = now -> ch[0];
  splay(now, fa);
Node *makeTree(Node *p, int 1, int r)
  if(1 > r)
   return null;
  int mid = (1 + r) / 2;
  Node *x = allocNode(mid);
  x->pre = p;
 x->ch[0] = makeTree(x, 1, mid - 1);
x->ch[1] = makeTree(x, mid + 1, r);
  update(x);
 return x;
int main()
  null = allocNode(0);
  null->size = 0;
  root = allocNode(0);
  root->ch[1] = allocNode(oo);
  root->ch[1]->pre = root;
  update(root);
  scanf("%d%d", &n, &m);
  root->ch[1]->ch[0] = makeTree(root->ch[1], 1, n);
  splay(root->ch[1]->ch[0], null);
  while (m --)
    int a, b;
    scanf("%d%d", &a, &b);
    a ++, b ++;
    select(a - 1, null);
    select(b + 1, root);
    makeTurned(root->ch[1]->ch[0]);
  for(int i = 1; i <= n; i ++)</pre>
    select(i + 1, null);
    printf("%d ", root->val);
```

# 5.6 Segment tree

```
// This is a solution for UVa 12299 - RMQ with Shifts. In this problem, we are
// given an array of integers A and a list of queries of 2 types:
1. query(L, R) returns the minimum value in A[L...R].
// 2. shift(i_1, i_2, ..., i_k) performs a left "circular shift" of A[i_1],
// A[i_2], ..., A[i_k]. For example, if A = {6, 2, 4, 8, 5, 1, 4}, then
shift(2, 4, 5, 7) yields 6, 8, 4, 5, 4, 1, 2. After that, shift(1, 2)
yields 8, 6, 4, 5, 4, 1, 2.
```

```
#include <bits/stdc++.h>
#define LOCAL
#define Lson(x) ((x << 1) + 1)
#define Rson(x) ((x << 1) + 2)
#define INF 99999999
using namespace std;
const int N = 100005:
int shift[35], num[N], len_shift;
string strLine;
struct Node {
    int left, right, Min_Value;
} node[4 * N];
void build(int left, int right, int x = 0) {
   node[x].left = left;
    node[x].right = right;
    if (left == right) {
        node[x].Min_Value = num[left];
        return:
    int mid = (left + right) / 2;
    //cout << mid << '\n' ;
    //cout << x << ' ' << node[x].left << ' ' << node[x].right << ' ' << '\n';
    build(left, mid, Lson(x));
    build(mid + 1, right, Rson(x));
    node[x].Min_Value = min(node[Lson(x)].Min_Value, node[Rson(x)].Min_Value);
void handle() {
    len_shift = 0;
   shift[len_shift] = shift[len_shift] * 10 + (int)(strLine[i] - '0');
        else (
            shift[++len shift] = 0;
    //finaly char is ')' , so len_shift is right
    sort(shift, shift + len_shift);
    //debug
    for(int i = 0; i < len_shift; i++)
    cout << shift[i] << ' ';</pre>
    cout << '\n' ;
int query(int left, int right, int x = 0) {
    if (node[x].left >= left && node[x].right <= right)</pre>
       return node[x].Min_Value;
    int mid = (node[x].left + node[x].right) / 2;
    int ans = INF;
    //cout << x << ' ' << node[x].left << ' ' << node[x].right << ' ' << node[x].Min_Value << '\n';
    if (left <= mid)</pre>
        ans = min(ans, query(left, right, Lson(x)));
    if (mid < right)</pre>
       ans = min(ans, query(left, right, Rson(x)));
    return ans:
void set_num(int position, int value, int x = 0) {
    if (node[x].left == position && node[x].right == position) {
        node[x].Min_Value = value;
    int mid = (node[x].left + node[x].right) / 2;
    if (position <= mid)</pre>
        set_num(position, value, Lson(x));
    if (mid < position)</pre>
        set_num(position, value, Rson(x));
    node[x].Min_Value = min(node[Lson(x)].Min_Value, node[Rson(x)].Min_Value);
int main() {
    int n, q, intTemp;
    ios::sync_with_stdio(0);
#ifdef LOCAL
    freopen("out.txt", "w", stdout);
    freopen("in1.txt", "r", stdin);
```

```
#endif // LOCAL
    cin >> n >> q;
    for (int i = 1; i <= n; i++)
        cin >> num[i];
    build(1, n);
    //debug
    for(int i = 0; i < 13; i++){
    cout << node[i].left << ' ' << node[i].right << ' ' << node[i].Min_Value << '\n';</pre>
    return 0 ;
    while (q--) {
        cin >> strLine;
        if (strLine[0] == 'q') {
             handle();
             cout << query(shift[0], shift[1]) << '\n';</pre>
        } else if (strLine[0] == 's') {
             handle();
             intTemp = num[shift[0]];
             for (int i = 1; i < len_shift; i++) {</pre>
                 set_num(shift[i - 1], num[shift[i]]);
num[shift[i - 1]] = num[shift[i]];
             num[shift[len_shift - 1]] = intTemp;
             set_num(shift[len_shift - 1], intTemp);
             //debua
             //cout << intTemp << ' ' << shift[len_shift-1] << '\n';
             //for(int i = 1; i <= n; i++)
                 cout << num[i] << ' ';
    return 0;
```

#### 5.7 Lowest common ancestor

```
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
                                          // children[i] contains the children of node i
int A[max_nodes][log_max_nodes+1];
                                          // A[i][j] is the 2^j-th ancestor of node i, or -1 if that
      ancestor does not exist
int L[max_nodes];
                                          // L[i] is the distance between node i and the root
// floor of the binary logarithm of n
int lb(unsigned int n)
    if(n==0)
        return -1;
    int. p = 0:
    if (n >= 1<<16) { n >>= 16; p += 16; }
    if (n >= 1<< 8) { n >>= 8; p += 8; }
if (n >= 1<< 4) { n >>= 4; p += 4; }
    if (n >= 1<< 2) { n >>= 2; p += 2; }
    if (n >= 1<< 1) {
    return p;
void DFS(int i, int 1)
    L[i] = 1;
for(int j = 0; j < children[i].size(); j++)
        DFS(children[i][j], 1+1);
int LCA(int p, int q)
     // ensure node p is at least as deep as node q
    if(L[p] < L[q])
        swap(p, q);
    // "binary search" for the ancestor of node p situated on the same level as q
    for(int i = log_num_nodes; i >= 0; i--)
        if(L[p] - (1<<i) >= L[q])
            p = A[p][i];
    if(p == q)
        return p;
    // "binary search" for the LCA
    for(int i = log num nodes; i >= 0; i--)
        if (A[p][i] != -1 && A[p][i] != A[q][i])
```

```
p = A[p][i];
             q = A[q][i];
    return A[p][0];
int main(int argc,char* argv[])
     // read num_nodes, the total number of nodes
    log_num_nodes=1b(num_nodes);
    for(int i = 0; i < num_nodes; i++)</pre>
        // read p, the parent of node i or -1 if node i is the root
        if(p != -1)
             children[p].push_back(i);
        else
             root = i;
     // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)
    for(int i = 0; i < num_nodes; i++)</pre>
             if(A[i][j-1] != -1)
                 A[i][j] = A[A[i][j-1]][j-1];
             else
                 A[i][j] = -1;
     // precompute L
    DFS (root, 0);
    return 0;
```

### 6 Miscellaneous

## 6.1 Dijkstra's algorithm using struct

```
// This is a solution for UVa 929 - Number Maze. Abridged problem statement:
// What is the shortest path from the maze's top-left to bottom-right corner?
#include <bits/stdc++.h>
#define LOCAL
#define INF 99999999
using namespace std;
int intMap[1010][1010] = {}, intValue[1010][1010] = {};
int m, n;
struct Node {
    int x, y, v;
    void read(int _x, int _y, int _v) {
        x = x;
         y = \underline{y};
    bool operator<(const Node &a) const {</pre>
         return v > a.v:
| nodNode:
void print_map() {
   for (int i = 1; i <= n; i++) {</pre>
        for (int j = 1; j <= m; j++) {
   if (intValue[i][j] == 99999999)</pre>
                  cout << 'r' << ' ';
                  cout << intValue[i][j] << ' ';
         cout << '\n';
    cout << '\n';
void bfs() {
    int x, y, intDirection[4][2] = {-1, 0, 0, 1, 1, 0, 0, -1};
int intDx, intDy;
    Node nodTemp;
    priority_queue<Node> deqNode;
```

```
nodTemp.read(1, 1, 0);
    deqNode.push (nodTemp);
    while (deqNode.size())
        x = deqNode.top().x;
        y = deqNode.top().y;
        deqNode.pop();
        for (int i = 0; i < 4; i++) {
            intDx = intDirection[i][0] + x;
            intDy = intDirection[i][1] + y;
            //cout << intDx << ' ' << intDy << ' ' << intValue[x][y] + intMap[intDx][intDy] << ' ' <<
                  i << '\n' :
            if (intValue[x][y] + intMap[intDx][intDy] < intValue[intDx][intDy]) {</pre>
                intValue[intDx][intDy] = intValue[x][y] + intMap[intDx][intDy];
                 nodTemp.read(intDx, intDy, intValue[intDx][intDy]);
                 deqNode.push(nodTemp);
        //print_map() ;
int main() {
#ifdef LOCAL
    freopen("in1.txt", "r", stdin);
    freopen("out.txt", "w", stdout);
#endif
    ios::sync_with_stdio(false);
    int intCase;
    cin >> intCase;
    while (intCase--) {
        cin >> n >> m;
        for (int i = 1; i <= n; i++) {
            for (int j = 1; j <= m; j++) {
    cin >> intMap[i][j];
                intValue[i][j] = INF;
        for (int i = 1; i <= n; i++) {
            intValue[i][0] = 0;
            intValue[i][m + 1] = 0;
            intMap[i][0] = INF + 1;
            intMap[i][m + 1] = INF + 1;
        for (int i = 1; i <= m; i++) {
            intValue[0][i] = 0;
            intValue[n + 1][i] = 0;
intMap[0][i] = INF + 1;
            intMap[n + 1][i] = INF + 1;
        intValue[1][1] = intMap[1][1];
        //cout << intValue[1][1] << '\n';
        cout << intValue[n][m] << '\n';</pre>
```

# 6.2 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
//
// Running time: O(n log n)
//
// INPUT: a vector of integers
// OUTPUT: a vector containing the longest increasing subsequence
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
```

```
VI LongestIncreasingSubsequence(VI v) {
  VPII best;
  VI dad(v.size(), -1);
  for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASNG
    PII item = make_pair(v[i], 0);
    VPII::iterator it = lower_bound(best.begin(), best.end(), item);
    item.second = i;
#else
    PII item = make_pair(v[i], i);
    VPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
    if (it == best.end()) {
     dad[i] = (best.size() == 0 ? -1 : best.back().second);
      best_push_back(item);
     dad[i] = it == best.begin() ? -1 : prev(it)->second;
      *it = item;
  VI ret:
  for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push_back(v[i]);
  reverse(ret.begin(), ret.end());
  return ret:
```

#### 6.3 Knuth-Morris-Pratt

```
Finds all occurrences of the pattern string p within the
text string t. Running time is O(n + m), where n and m
are the lengths of p and t, respecitvely.
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildPi(string& p, VI& pi)
  pi = VI(p.length());
  int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
    while (k >= -1 &  p[k+1] != p[i])
     k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP (string& t, string& p)
 VI pi;
  buildPi(p, pi);
  int k = -1;
for(int i = 0; i < t.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != t[i])
     k = (k == -1) ? -2 : pi[k];
    if(k == p.length() - 1) {
     // p matches t[i-m+1, ..., i]
      cout << "matched at index " << i-k << ": ";
      cout << t.substr(i-k, p.length()) << endl;</pre>
     k = (k == -1) ? -2 : pi[k];
  return 0:
int main()
  string a = "AABAACAADAABAABA", b = "AABA";
  KMP(a, b); // expected matches at: 0, 9, 12
  return 0;
```

#### 6.4 Dates

```
// Routines for performing computations on dates. In these routines,
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
#include <iostream>
#include <string>
using namespace std;
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y) {
  return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
// converts integer (Julian day number) to Gregorian date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
 int x, n, i, j;
  x = id + 68569;
 n = 4 * x / 146097;
  x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
  x = 1461 * i / 4 - 31;
  j = 80 * x / 2447;
  d = x - 2447 * j / 80;
  x = j / 11;

m = j + 2 - 12 * x;
  y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
 return dayOfWeek[jd % 7];
int main (int argc, char **argv) {
  int jd = dateToInt (3, 24, 2004);
  int m, d, y;
intToDate (jd, m, d, y);
  string day = intToDay (jd);
  // expected output:
  // 2453089
// 3/24/2004
// Wed
  cout << jd << endl
    << m << "/" << d << "/" << y << endl
    << day << endl:
```

#### 6.5 Prime numbers

```
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x)
  if(x<=1) return false;</pre>
  if(x<=3) return true;</pre>
  if (!(x%2) || !(x%3)) return false;
  LL s=(LL)(sgrt((double)(x))+EPS);
  for(LL i=5;i<=s;i+=6)
    if (!(x%i) || !(x%(i+2))) return false;
  return true;
// Primes less than 1000:
                               11
                                     13
                                          17
                                                19
       41
             43
                   47
                        53
                               59
                                     61
                                           67
                                                 71
                                                             79
                                                                  83
                                                     137
       97
                 103
                       107
                             109
                                    113
                                         127
                                                131
                                                            139
                                                                 149
     157
                              179
           163
                  167
                        173
                                    181
                                          191
                                                193
                                                     197
                                                            199
                                                                  211
                                                                        223
     227
           229
                        239
                             241
                                    251
                                          257
                                                263
                                                     269
                                                                  277
                                                                        281
                                    317
     283
           293
                  307
                        311
                              313
                                          331
                                                     347
                                                            349
                                                                        359
                              389
                                    397
                                                409
      367
           373
                  379
                        383
                                          401
                                                      419
                                                            421
                                                                  431
                                                                        433
      439
           443
                  449
                        457
                              461
                                    463
                                          467
                                                479
                                                      487
                                                            491
                                                                  499
                                                                        503
                              547
                                    557
      509
                        541
                                                569
                                                      571
                                                            577
                                                                  587
                                                                        593
            521
                  523
                                          563
                              617
                                    619
                                                            647
      599
            601
                  607
                        613
                                          631
                                                641
                                                      643
                                                                  653
                                                                        659
            673
                        683
                              691
                                    701
                                          709
                                                719
                                                                  739
                                                                        743
      661
                  677
```

```
829 839 853 857 859 863 877 881 883 887
                 937 941 947 953 967 971 977 983 991
// Other primes:
     The largest prime smaller than 10 is 7.
     The largest prime smaller than 100 is 97.
     The largest prime smaller than 1000 is 997.
     The largest prime smaller than 10000 is 9973
     The largest prime smaller than 100000 is 99991
     The largest prime smaller than 1000000 is 999983.
     The largest prime smaller than 10000000 is 9999991. The largest prime smaller than 100000000 is 99999989.
     The largest prime smaller than 1000000000 is 999999937. The largest prime smaller than 10000000000 is 9999999967.
     The largest prime smaller than 10000000000 is 99999999977
     The largest prime smaller than 100000000000 is 999999999999.
     The largest prime smaller than 100000000000 is 999999999971.
     The largest prime smaller than 1000000000000 is 9999999999973
     The largest prime smaller than 10000000000000 is 999999999999999.
     The largest prime smaller than 100000000000000 is 99999999999937.
     The largest prime smaller than 1000000000000000 is 99999999999997.
```

### 6.6 Latitude/longitude

```
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
#include <iostream>
#include <cmath>
using namespace std;
struct 11
  double r, lat, lon;
};
struct rect
  double x, y, z;
};
11 convert(rect& P)
  11 Q;
  Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
  Q.lat = 180/M_PI*asin(P.z/Q.r);
  Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
  return 0;
rect convert(ll& Q)
  P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
  P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
  P.z = Q.r*sin(Q.lat*M_PI/180);
  return P:
int main()
  rect A:
  11 B:
 A.x = -1.0; A.y = 2.0; A.z = -3.0;
  B = convert(A);
cout << B.r << " " << B.lat << " " << B.lon << endl;</pre>
  A = convert(B);
cout << A.x << " " << A.y << " " << A.z << endl;
```

# 6.7 C++ input/output

```
#include <iomanip>
using namespace std;
int main()
{
    // Ouput a specific number of digits past the decimal point,
     // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);

    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf(ios::showpoint);

    // Output a '+' before positive values
    cout.setf(ios::showpos);
    cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal
    cout << hex << 100 << " " << 1000 << " " << 1000 << endl;
    cout << hex << 100 << " " << 1000 << endl;
    cout.unsetf(ios::showpos);</pre>
```

#### 6.8 Random STL stuff

```
// Example for using stringstreams and next_permutation
#include <algorithm>
#include <iostream>
#include <sstream>
#include <vector>
using namespace std;
int main(void) {
  vector<int> v;
  v.push_back(1);
  v.push_back(2);
  v.push back(3):
  v.push_back(4);
  // Expected output: 1 2 3 4
   ostringstream oss;
oss << v[0] << " " << v[1] << " " << v[2] << " " << v[3];
    // for input from a string s,
    // istringstream iss(s);
// iss >> variable;
    cout << oss.str() << endl;
  } while (next_permutation (v.begin(), v.end()));
  v.clear();
  v.push_back(1);
  v.push_back(2);
  v.push_back(1);
  v.push_back(3);
  // To use unique, first sort numbers. Then call
  // unique to place all the unique elements at the beginning
  // of the vector, and then use erase to remove the duplicate
  // elements.
  sort(v.begin(), v.end());
  v.erase(unique(v.begin(), v.end()), v.end());
  // Expected output: 1 2 3
  for (size_t i = 0; i < v.size(); i++)</pre>
   cout << v[i] << " ";
  cout << endl;</pre>
```

### 3.9 Regular expressions (Java)

```
// Code which demonstrates the use of Java's regular expression libraries.
// This is a solution for
      Loglan: a logical language
      http://acm.uva.es/p/v1/134.html
// In this problem, we are given a regular language, whose rules can be
\ensuremath{//} inferred directly from the code. For each sentence in the input, we must
// determine whether the sentence matches the regular expression or not. The
//\ {\rm code}\ {\rm consists}\ {\rm of}\ (1)\ {\rm building}\ {\rm the}\ {\rm regular}\ {\rm expression}\ ({\rm which}\ {\rm is}\ {\rm fairly}
// complex) and (2) using the regex to match sentences.
import java.util.*;
import java.util.regex.*;
public class LogLan {
      public static String BuildRegex () {
           String space = " +";
           String A = "([aeiou])";
String C = "([a-z&&[aeiou]])";
String MOD = "(g" + A + ")";
String BA = "(b" + A + ")";
String DA = "(d" + A + ")";
           String DA = '(1" + A + ")";

String LA = "(1" + A + ")";

String NAM = "([a-z]*" + C + ")";
           String PREDA = "(" + C + C + A + C + A + "|" + C + A + C + C + A + ")";
            String predstring = "(" + PREDA + "(" + space + PREDA + ")*)";
           String predname = "(" + LA + space + predstring + "|" + NAM + ")";

String preds = "(" + predstring + "(" + space + A + space + predstring + ")*)";

String predclaim = "(" + predname + space + BA + space + preds + "|" + DA + space +
                 preds + ")";
           preds + );
String verbpred = "(" + MOD + space + predstring + ")";
String statement = "(" + predname + space + verbpred + space + predname + "|" + predname + space + verbpred + ")";
String sentence = "(" + statement + "|" + predclaim + ")";
           return "^" + sentence + "$";
      public static void main (String args[]) {
            String regex = BuildRegex();
           Pattern pattern = Pattern.compile (regex);
```

```
Scanner s = new Scanner(System.in);
     while (true) {
         // In this problem, each sentence consists of multiple lines, where the last
         // line is terminated by a period. The code below reads lines until
         // encountering a line whose final character is a '.'. Note the use of
               s.length() to get length of string
               s.charAt() to extract characters from a Java string
               s.trim() to remove whitespace from the beginning and end of Java string
         // Other useful String manipulation methods include
               s.compareTo(t) < 0 if s < t, lexicographically
s.indexOf("apple") returns index of first occurrence of "apple" in s</pre>
               s.lastIndexOf("apple") returns index of last occurrence of "apple" in s
               s.replace(c,d) replaces occurrences of character c with d
               s.startsWith("apple) returns (s.indexOf("apple") == 0)
               s.toLowerCase() / s.toUpperCase() returns a new lower/uppercased string
               Integer.parseInt(s) converts s to an integer (32-bit)
               Long.parseLong(s) converts s to a long (64-bit)
               Double.parseDouble(s) converts s to a double
         String sentence = "";
         while (true) {
             sentence = (sentence + " " + s.nextLine()).trim();
             if (sentence.equals("#")) return;
             if (sentence.charAt(sentence.length()-1) == '.') break;
         // now, we remove the period, and match the regular expression
         String removed_period = sentence.substring(0, sentence.length()-1).trim();
         if (pattern.matcher (removed_period).find()){
             System.out.println ("Good");
         } else {
             System.out.println ("Bad!");
}
```