

NTUT_Knights ICPC Team Notebook

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1 Dynamic programming algorithms

1.1 Longest common subsequence

```

typedef vector<T> VT;
typedef vector<VT> VVT;

typedef vector<int> VI;
typedef vector<VI> VVI;

void backtrack(VVI& dp, VT& res, VT& A, VT& B, int i, int j)
{
    if(!i || !j) return;
    if(A[i-1] == B[j-1]) { res.push_back(A[i-1]); backtrack(dp, res, A, B, i-1, j-1); }
    else
    {
        if(dp[i][j-1] >= dp[i-1][j]) backtrack(dp, res, A, B, i, j-1);
        else backtrack(dp, res, A, B, i-1, j);
    }
}

void backtrackall(VVI& dp, set<VT>& res, VT& A, VT& B, int i, int j)
{
    if(!i || !j) { res.insert(VI()); return; }
    if(A[i-1] == B[j-1])
    {
        set<VT> tempres;
        backtrackall(dp, tempres, A, B, i-1, j-1);
        for(set<VT>::iterator it=tempres.begin(); it!=tempres.end(); it++)
        {
            VT temp = *it;
            temp.push_back(A[i-1]);
            res.insert(temp);
        }
    }
    else
    {
        if(dp[i][j-1] >= dp[i-1][j]) backtrackall(dp, res, A, B, i, j-1);
        if(dp[i][j-1] <= dp[i-1][j]) backtrackall(dp, res, A, B, i-1, j);
    }
}

VT LCS(VT& A, VT& B)
{
    VVI dp;
    int n = A.size(), m = B.size();
    dp.resize(n+1);
    for(int i=0; i<=n; i++) dp[i].resize(m+1, 0);

    for(int i=1; i<=n; i++)
        for(int j=1; j<=m; j++)
        {
            if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
            else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
        }

    VT res;
    backtrack(dp, res, A, B, n, m);
    reverse(res.begin(), res.end());
    return res;
}

set<VT> LCSall(VT& A, VT& B)
{
    VVI dp;
    int n = A.size(), m = B.size();
    dp.resize(n+1);
    for(int i=0; i<=n; i++) dp[i].resize(m+1, 0);
    for(int i=1; i<=n; i++)
        for(int j=1; j<=m; j++)
        {
            if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
            else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
        }

    set<VT> res;
    backtrackall(dp, res, A, B, n, m);
    return res;
}

int main()
{
    int a[] = { 0, 5, 5, 2, 1, 4, 2, 3 }, b[] = { 5, 2, 4, 3, 2, 1, 2, 1, 3 };
    VI A = VI(a, a+8), B = VI(b, b+9);
    VI C = LCS(A, B);

    for(int i=0; i<C.size(); i++) cout << C[i] << " ";
    cout << endl << endl;

    set<VI> D = LCSall(A, B);
    for(set<VI>::iterator it = D.begin(); it != D.end(); it++)
    {
        for(int i=0; i<(*it).size(); i++) cout << (*it)[i] << " ";
        cout << endl;
    }
}

```

```

/*
Calculates the length of the longest common subsequence of two vectors.
Backtracks to find a single subsequence or all subsequences. Runs in
O(m*n) time except for finding all longest common subsequences, which
may be slow depending on how many there are.
*/

```

```

#include <iostream>
#include <vector>
#include <set>
#include <algorithm>

using namespace std;

typedef int T;

```

1.2 Max 2D range sum

// This is a solution for UVa 108 - Maximum Sum. In this problem, we are given a square matrix of integers and are asked to find a sub-matrix with the maximum sum. This 1D DP + greedy (Kadane's) solution takes time $O(n^3)$.

```
#include <bits/stdc++.h>
using namespace std;

#define MAX_n 110

int A[MAX_n][MAX_n];

int main() {
    int n; scanf("%d", &n);           // square matrix size
    for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j) {
            scanf("%d", &A[i][j]);
            if (j > 0) A[i][j] += A[i][j-1]; // pre-processing
        }
    int maxSubRect = -127*100*100;      // lowest possible val
    for (int l = 0; l < n; ++l)
        for (int r = l; r < n; ++r) {
            int subRect = 0;
            for (int row = 0; row < n; ++row) {
                // Max 1D Range Sum on columns of this row
                if (l > 0) subRect += A[row][r] - A[row][l-1];
                else subRect += A[row][r];
                // Kadane's algorithm on rows
                if (subRect < 0) subRect = 0; // restart if negative
                maxSubRect = max(maxSubRect, subRect);
            }
        }
    printf("%d\n", maxSubRect);
    return 0;
}
```

1.3 0-1 Knapsack

// This is a solution for UVa 10130 - SuperSale. This is a 0-1 Knapsack variant solved using top-down DP $O(n*W)$.

```
#include <bits/stdc++.h>
using namespace std;

const int MAX_N = 1010;
const int MAX_W = 40;

int N, V[MAX_N], W[MAX_N], memo[MAX_N][MAX_W];

int dp(int id, int remW) {
    if ((id == N) || (remW == 0)) return 0; // two base cases
    int &ans = memo[id][remW];
    if (ans != -1) return ans; // computed before
    if (W[id] > remW) return ans = dp(id+1, remW); // no choice, skip
    return ans = max(dp(id+1, remW), // has choice, skip
                    V[id]+dp(id+1, remW-W[id])); // or take
}

int main() {
    int T; scanf("%d", &T);
    while (T--) {
        memset(memo, -1, sizeof memo);
        scanf("%d", &N);
        for (int i = 0; i < N; ++i)
            scanf("%d %d", &V[i], &W[i]);
        int ans = 0;
        int G; scanf("%d", &G);
        while (G--) {
            int MW; scanf("%d", &MW);
            ans += dp(0, MW);
        }
        printf("%d\n", ans);
    }
    return 0;
}
```

1.4 Traveling salesman

// This is a solution for UVa 10496 - Collecting Beepers. The problem is a variant of the Traveling Salesman Problem (TSP): Given n cities and their pairwise distances in the form of a matrix 'dist' of size $n * n$, compute the minimum cost of making a tour that starts from any city s , goes through all the other $n - 1$ cities exactly once, and finally returns to the city s . In this case, the salesman is Karel in a 2D world who can only move along the x and y axis. The cities are beepers whose coordinates are given, from which pairwise distances can be calculated.

INPUT: The first line is the number of test cases. The first line of each test case is world's size (x -size and y -size). Next is the starting position of Karel. Next is the number of beepers. Next are the beepers' x - and y -coordinates.

OUTPUT: For each test case, output the minimum distance to move from Karel's starting position to each of the beepers and back to the starting position.

```
#include <bits/stdc++.h>
using namespace std;

#define LSONe(S) ((S) & -(S))

const int MAX_n = 11;

int dist[MAX_n][MAX_n], memo[MAX_n][1<<(MAX_n-1)]; // Karel + max 10 beepers

int dp(int u, int mask) {
    if (mask == 0) return dist[u][0]; // mask = free coordinates
    int &ans = memo[u][mask]; // close the loop
    if (ans != -1) return ans; // computed before
    ans = 2000000000;
    int m = mask;
    while (m) {
        int two_pow_v = LSONe(m); // up to  $O(n)$ 
        int v = __builtin_ctz(two_pow_v)+1; // but this is fast
        ans = min(ans, dist[u][v] + dp(v, mask^two_pow_v)); // offset v by +1
        m -= two_pow_v;
    }
    return ans;
}

int main() {
    int TC; scanf("%d", &TC);
    while (TC--) {
        int xsize, ysize; scanf("%d %d", &xsize, &ysize); // these two values are not used
        int x[MAX_n], y[MAX_n];
        scanf("%d %d", &x[0], &y[0]);
        int n; scanf("%d", &n); ++n; // include Karel
        for (int i = 1; i < n; ++i) // Karel is at index 0
            scanf("%d %d", &x[i], &y[i]);
        for (int i = 0; i < n; ++i) // build distance table
            for (int j = i; j < n; ++j)
                dist[i][j] = dist[j][i] = abs(x[i]-x[j]) + abs(y[i]-y[j]); // Manhattan distance
        memset(memo, -1, sizeof memo);
        printf("The shortest path has length %d\n", dp(0, (1<<(n-1))-1)); // DP-TSP
    }
    return 0;
}
```

2 Geometry

2.1 Convex hull

// Compute the 2D convex hull of a set of points using the monotone chain algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is #defined.

// Running time: $O(n \log n)$

INPUT: a vector of input points, unordered.

OUTPUT: a vector of points in the convex hull, counterclockwise, starting with bottommost/leftmost point

```
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
// BEGIN CUT
#include <map>
// END CUT
```

```
using namespace std;
```

```
#define REMOVE_REDUNDANT
```

```
typedef double T;
```

```

const T EPS = 1e-7;
struct PT {
    T x, y;
    PT() {}
    PT(T x, T y) : x(x), y(y) {}
    bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y, rhs.x); }
    bool operator==(const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y, rhs.x); }
};

T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }

#ifdef REMOVE_REDUNDANT
bool between(const PT &a, const PT &b, const PT &c) {
    return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 && (a.y-b.y)*(c.y-b.y) <= 0);
}
#endif

void ConvexHull(vector<PT> &pts) {
    sort(pts.begin(), pts.end());
    pts.erase(unique(pts.begin(), pts.end(), pts.end()));
    vector<PT> up, dn;
    for (int i = 0; i < pts.size(); i++) {
        while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
        while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();
        up.push_back(pts[i]);
        dn.push_back(pts[i]);
    }
    pts = dn;
    for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);

#ifdef REMOVE_REDUNDANT
    if (pts.size() <= 2) return;
    dn.clear();
    dn.push_back(pts[0]);
    dn.push_back(pts[1]);
    for (int i = 2; i < pts.size(); i++) {
        if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();
        dn.push_back(pts[i]);
    }
    if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
        dn[0] = dn.back();
        dn.pop_back();
    }
    pts = dn;
#endif
}

// BEGIN CUT
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)

int main() {
    int t;
    scanf("%d", &t);
    for (int caseno = 0; caseno < t; caseno++) {
        int n;
        scanf("%d", &n);
        vector<PT> v(n);
        for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);
        vector<PT> h(v);
        map<PT,int> index;
        for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
        ConvexHull(h);

        double len = 0;
        for (int i = 0; i < h.size(); i++) {
            double dx = h[i].x - h[(i+1)%h.size()].x;
            double dy = h[i].y - h[(i+1)%h.size()].y;
            len += sqrt(dx*dx+dy*dy);
        }

        if (caseno > 0) printf("\n");
        printf("%.2f\n", len);
        for (int i = 0; i < h.size(); i++) {
            if (i > 0) printf(" ");
            printf("%d", index[h[i]]);
        }
        printf("\n");
    }
}

// END CUT

```

```

#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>

using namespace std;

double INF = 1e100;
double EPS = 1e-12;

struct PT {
    double x, y;
    PT() {}
    PT(double x, double y) : x(x), y(y) {}
    PT(const PT &p) : x(p.x), y(p.y) {}
    PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
    PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
    PT operator * (double c) const { return PT(x*c, y*c); }
    PT operator / (double c) const { return PT(x/c, y/c); }
};

double dot(PT p, PT q) { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {
    return os << "(" << p.x << ", " << p.y << ")";
}

// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCW90(PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
    return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
}

// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
    return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
}

// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
    double r = dot(b-a,b-a);
    if (fabs(r) < EPS) return a;
    r = dot(c-a, b-a)/r;
    if (r < 0) return a;
    if (r > 1) return b;
    return a + (b-a)*r;
}

// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
    return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
}

// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
    double a, double b, double c, double d)
{
    return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
}

// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
    return fabs(cross(b-a, c-d)) < EPS;
}

bool LinesCollinear(PT a, PT b, PT c, PT d) {
    return LinesParallel(a, b, c, d)
        && fabs(cross(a-b, a-c)) < EPS
        && fabs(cross(c-d, c-a)) < EPS;
}

// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
    if (LinesCollinear(a, b, c, d)) {
        if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
            dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
        if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
            return false;
        return true;
    }
    if (cross(d-a, b-a) + cross(c-a, b-a) > 0) return false;
    if (cross(a-c, d-c) + cross(b-c, d-c) > 0) return false;
    return true;
}

// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if

```

2.2 Miscellaneous geometry

// C++ routines for computational geometry.

```

// segments intersect first
PT ComputelineIntersection(PT a, PT b, PT c, PT d) {
    b=b-a; d=d-c; c=c-a;
    assert(dot(b, b) > EPS && dot(d, d) > EPS);
    return a + b*cross(c, d)/cross(b, d);
}

// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
    b=(a+b)/2;
    c=(a+c)/2;
    return ComputelineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
}

// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
    bool c = 0;
    for (int i = 0; i < p.size(); i++){
        int j = (i+1)%p.size();
        if ((p[i].y <= q.y && q.y < p[j].y ||
            p[j].y <= q.y && q.y < p[i].y) &&
            q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
            c = !c;
    }
    return c;
}

// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
    for (int i = 0; i < p.size(); i++)
        if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)
            return true;
    return false;
}

// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
    vector<PT> ret;
    b = b-a;
    a = a-c;
    double A = dot(b, b);
    double B = dot(a, b);
    double C = dot(a, a) - r*r;
    double D = B*B - A*C;
    if (D < -EPS) return ret;
    ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
    if (D > EPS)
        ret.push_back(c+a+b*(-B-sqrt(D))/A);
    return ret;
}

// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
    vector<PT> ret;
    double d = sqrt(dist2(a, b));
    if (d > r+R || d+min(r, R) < max(r, R)) return ret;
    double x = (d*d-R*R+r*r)/(2*d);
    double y = sqrt(r*r-x*x);
    PT v = (b-a)/d;
    ret.push_back(a+v*x + RotateCCW90(v)*y);
    if (y > 0)
        ret.push_back(a+v*x - RotateCCW90(v)*y);
    return ret;
}

// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
    double area = 0;
    for(int i = 0; i < p.size(); i++) {
        int j = (i+1) % p.size();
        area += p[i].x*p[j].y - p[j].x*p[i].y;
    }
    return area / 2.0;
}

double ComputeArea(const vector<PT> &p) {
    return fabs(ComputeSignedArea(p));
}

PT ComputeCentroid(const vector<PT> &p) {
    PT c(0,0);

```

```

double scale = 6.0 * ComputeSignedArea(p);
for (int i = 0; i < p.size(); i++){
    int j = (i+1) % p.size();
    c = c + (p[i].x*p[j].y - p[j].x*p[i].y);
}
return c / scale;
}

// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
    for (int i = 0; i < p.size(); i++) {
        for (int k = i+1; k < p.size(); k++) {
            int j = (i+1) % p.size();
            int l = (k+1) % p.size();
            if (i == 1 || j == k) continue;
            if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
                return false;
        }
    }
    return true;
}

int main() {
    // expected: (-5,2)
    cerr << RotateCCW90(PT(2,5)) << endl;

    // expected: (5,-2)
    cerr << RotateCW90(PT(2,5)) << endl;

    // expected: (-5,2)
    cerr << RotateCCW(PT(2,5), M_PI/2) << endl;

    // expected: (5,2)
    cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;

    // expected: (5,2) (7.5,3) (2.5,1)
    cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "
        << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
        << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;

    // expected: 6.78903
    cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;

    // expected: 1 0 1
    cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
        << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
        << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;

    // expected: 0 0 1
    cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
        << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
        << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;

    // expected: 1 1 1 0
    cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
        << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
        << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
        << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;

    // expected: (1,2)
    cerr << ComputelineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;

    // expected: (1,1)
    cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;

    vector<PT> v;
    v.push_back(PT(0,0));
    v.push_back(PT(5,0));
    v.push_back(PT(5,5));
    v.push_back(PT(0,5));

    // expected: 1 1 1 0 0
    cerr << PointInPolygon(v, PT(2,2)) << " "
        << PointInPolygon(v, PT(2,0)) << " "
        << PointInPolygon(v, PT(0,2)) << " "
        << PointInPolygon(v, PT(5,2)) << " "
        << PointInPolygon(v, PT(2,5)) << endl;

    // expected: 0 1 1 1 1
    cerr << PointOnPolygon(v, PT(2,2)) << " "
        << PointOnPolygon(v, PT(2,0)) << " "
        << PointOnPolygon(v, PT(0,2)) << " "
        << PointOnPolygon(v, PT(5,2)) << " "
        << PointOnPolygon(v, PT(2,5)) << endl;

    // expected: (1,6)
    // (5,4) (4,5)
    // blank line
    // (4,5) (5,4)
    // blank line
    // (4,5) (5,4)

```

```

vector<PT> u = CircleLineIntersection(P(0,6), P(2,6), P(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleLineIntersection(P(0,9), P(9,0), P(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(P(1,1), P(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(P(1,1), P(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(P(1,1), P(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(P(1,1), P(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;

// area should be 5.0
// centroid should be (1.1666666, 1.1666666)
PT pa[] = { P(0,0), P(5,0), P(1,1), P(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;

return 0;
}

```

3 Numerical algorithms

3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```

// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.

```

```

#include <iostream>
#include <vector>
#include <algorithm>

using namespace std;

typedef vector<int> VI;
typedef pair<int, int> PII;

// return a % b (positive value)
int mod(int a, int b) {
    return ((a%b) + b) % b;
}

// computes gcd(a,b)
int gcd(int a, int b) {
    while (b) { int t = a%b; a = b; b = t; }
    return a;
}

// computes lcm(a,b)
int lcm(int a, int b) {
    return a / gcd(a, b) * b;
}

// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
{
    int ret = 1;
    while (b)
    {
        if (b & 1) ret = mod(ret*a, m);
        a = mod(a*a, m);
        b >>= 1;
    }
    return ret;
}

// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
    int xx = y = 0;
    int yy = x = 1;
    while (b) {
        int q = a / b;
        int t = b; b = a%b; a = t;
        t = xx; xx = x - q*xx; x = t;
        t = yy; yy = y - q*yy; y = t;
    }
    return a;
}

```

```

// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
    int x, y;
    VI ret;
    int g = extended_euclid(a, n, x, y);
    if (!(b%g)) {
        x = mod(x*(b / g), n);
        for (int i = 0; i < g; i++)
            ret.push_back(mod(x + i*(n / g), n));
    }
    return ret;
}

// computes b such that ab = 1 (mod n), returns -1 on failure
int mod_inverse(int a, int n) {
    int x, y;
    int g = extended_euclid(a, n, x, y);
    if (g > 1) return -1;
    return mod(x, n);
}

// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
    int s, t;
    int g = extended_euclid(m1, m2, s, t);
    if (r1%g != r2%g) return make_pair(0, -1);
    return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g);
}

// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
    PII ret = make_pair(r[0], m[0]);
    for (int i = 1; i < m.size(); i++) {
        ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
        if (ret.second == -1) break;
    }
    return ret;
}

// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
    if (!(a && b))
    {
        if (c) return false;
        x = 0; y = 0;
        return true;
    }
    if (!a)
    {
        if (c % b) return false;
        x = 0; y = c / b;
        return true;
    }
    if (!b)
    {
        if (c % a) return false;
        x = c / a; y = 0;
        return true;
    }
    int g = gcd(a, b);
    if (c % g) return false;
    x = c / g + mod_inverse(a / g, b / g);
    y = (c - a*x) / b;
    return true;
}

int main() {
    // expected: 2
    cout << gcd(14, 30) << endl;

    // expected: 2 -2 1
    int x, y;
    int g = extended_euclid(14, 30, x, y);
    cout << g << " " << x << " " << y << endl;

    // expected: 95 451
    VI sols = modular_linear_equation_solver(14, 30, 100);
    for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";
    cout << endl;

    // expected: 8
    cout << mod_inverse(8, 9) << endl;

    // expected: 23 105

```

```
//          11 12
PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
cout << ret.first << " " << ret.second << endl;
ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second << endl;

// expected: 5 -15
if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
cout << x << " " << y << endl;
return 0;
}
```

3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
//
// Uses:
// (1) solving systems of linear equations (AX=B)
// (2) inverting matrices (AX=I)
// (3) computing determinants of square matrices
//
// Running time: O(n^3)
//
// INPUT:   a[][] = an nxn matrix
//          b[][] = an nxm matrix
//
// OUTPUT:  X      = an nxm matrix (stored in b[][])
//          A^-1    = an nxn matrix (stored in a[][])
//          returns determinant of a[][]

#include <iostream>
#include <vector>
#include <cmath>

using namespace std;

const double EPS = 1e-10;

typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;

T GaussJordan(VVT &a, VVT &b) {
    const int n = a.size();
    const int m = b[0].size();
    VI irow(n), icol(n), ipiv(n);
    T det = 1;

    for (int i = 0; i < n; i++) {
        int pj = -1, pk = -1;
        for (int j = 0; j < n; j++) if (!ipiv[j])
            for (int k = 0; k < n; k++) if (!ipiv[k])
                if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
        if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }
        ipiv[pj]++;
        swap(a[pj], a[pk]);
        swap(b[pj], b[pk]);
        if (pj != pk) det *= -1;
        irow[i] = pj;
        icol[i] = pk;

        T c = 1.0 / a[pk][pk];
        det *= a[pk][pk];
        a[pk][pk] = 1.0;
        for (int p = 0; p < n; p++) a[pk][p] *= c;
        for (int p = 0; p < m; p++) b[pk][p] *= c;
        for (int p = 0; p < n; p++) if (p != pk) {
            c = a[p][pk];
            a[p][pk] = 0;
            for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
            for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
        }

        for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
            for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
        }
    }

    return det;
}

int main() {
    const int n = 4;
    const int m = 2;
}
```

```
double A[n][n] = { {1,2,3,4},{1,0,1,0},{5,3,2,4},{6,1,4,6} };
double B[n][m] = { {1,2},{4,3},{5,6},{8,7} };
VVT a(n), b(n);
for (int i = 0; i < n; i++) {
    a[i] = VT(A[i], A[i] + n);
    b[i] = VT(B[i], B[i] + m);
}

double det = GaussJordan(a, b);

// expected: 60
cout << "Determinant: " << det << endl;

// expected: -0.233333 0.166667 0.133333 0.066667
//          0.166667 0.166667 0.333333 -0.333333
//          0.233333 0.833333 -0.133333 -0.066667
//          0.05 -0.75 -0.1 0.2
cout << "Inverse: " << endl;
for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
        cout << a[i][j] << ' ';
    }
    cout << endl;
}

// expected: 1.63333 1.3
//          -0.166667 0.5
//          2.36667 1.7
//          -1.85 -1.35
cout << "Solution: " << endl;
for (int i = 0; i < n; i++) {
    for (int j = 0; j < m; j++) {
        cout << b[i][j] << ' ';
    }
    cout << endl;
}
}
```

3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
//
// Running time: O(n^3)
//
// INPUT:   a[][] = an nxm matrix
//
// OUTPUT:  rref[][] = an nxm matrix (stored in a[][])
//          returns rank of a[][]

#include <iostream>
#include <vector>
#include <cmath>

using namespace std;

const double EPSILON = 1e-10;

typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;

int rref(VVT &a) {
    int n = a.size();
    int m = a[0].size();
    int r = 0;
    for (int c = 0; c < m && r < n; c++) {
        int j = r;
        for (int i = r + 1; i < n; i++)
            if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
        if (fabs(a[j][c]) < EPSILON) continue;
        swap(a[j], a[r]);

        T s = 1.0 / a[r][c];
        for (int j = 0; j < m; j++) a[r][j] *= s;
        for (int i = 0; i < n; i++) if (i != r) {
            T t = a[i][c];
            for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];
        }
        r++;
    }
    return r;
}

int main() {
    const int n = 5, m = 4;
    double A[n][m] = {
        {16, 2, 3, 13},
    }
}
```

```

    { 5, 11, 10, 8},
    { 9, 7, 6, 12},
    { 4, 14, 15, 1},
    {13, 21, 21, 13}};
VVT a(n);
for (int i = 0; i < n; i++)
    a[i] = VT(A[i], A[i] + m);

int rank = rref(a);

// expected: 3
cout << "Rank: " << rank << endl;

// expected: 1 0 0 1
//           0 1 0 3
//           0 0 1 -3
//           0 0 3.10862e-15
//           0 0 2.22045e-15
cout << "rref: " << endl;
for (int i = 0; i < 5; i++) {
    for (int j = 0; j < 4; j++)
        cout << a[i][j] << ' ';
    cout << endl;
}
}

```

3.4 Fast Fourier transform

```

#include <cassert>
#include <cstdio>
#include <cmath>

struct cpx
{
    cpx() {}
    cpx(double aa):a(aa),b(0){}
    cpx(double aa, double bb):a(aa),b(bb){}
    double a;
    double b;
    double modsq(void) const
    {
        return a * a + b * b;
    }
    cpx bar(void) const
    {
        return cpx(a, -b);
    }
};

cpx operator +(cpx a, cpx b)
{
    return cpx(a.a + b.a, a.b + b.b);
}

cpx operator *(cpx a, cpx b)
{
    return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
}

cpx operator /(cpx a, cpx b)
{
    cpx r = a * b.bar();
    return cpx(r.a / b.modsq(), r.b / b.modsq());
}

cpx EXP(double theta)
{
    return cpx(cos(theta), sin(theta));
}

const double two_pi = 4 * acos(0);

// in:    input array
// out:    output array
// step:   {SET TO 1} (used internally)
// size:   length of the input/output (MUST BE A POWER OF 2)
// dir:    either plus or minus one (direction of the FFT)
// RESULT: out[k] = \sum_{j=0}^{size-1} in[j] * exp(dir * 2pi * i * j * k / size)
void FFT(cpx *in, cpx *out, int step, int size, int dir)
{
    if (size < 1) return;
    if (size == 1)
    {
        out[0] = in[0];
        return;
    }
    FFT(in, out, step * 2, size / 2, dir);
    FFT(in + step, out + size / 2, step * 2, size / 2, dir);
}

```

```

for(int i = 0 ; i < size / 2 ; i++)
{
    cpx even = out[i];
    cpx odd = out[i + size / 2];
    out[i] = even + EXP(dir * two_pi * i / size) * odd;
    out[i + size / 2] = even + EXP(dir * two_pi * (i + size / 2) / size) * odd;
}

// Usage:
// f[0...N-1] and g[0...N-1] are numbers
// Want to compute the convolution h, defined by
// h[n] = sum of f[k]g[n-k] (k = 0, ..., N-1).
// Here, the index is cyclic; f[-1] = f[N-1], f[-2] = f[N-2], etc.
// Let F[0...N-1] be FFT(f), and similarly, define G and H.
// The convolution theorem says H[n] = F[n]G[n] (element-wise product).
// To compute h[] in O(N log N) time, do the following:
// 1. Compute F and G (pass dir = 1 as the argument).
// 2. Get H by element-wise multiplying F and G.
// 3. Get h by taking the inverse FFT (use dir = -1 as the argument)
//    and *dividing by N*. DO NOT FORGET THIS SCALING FACTOR.

int main(void)
{
    printf("If rows come in identical pairs, then everything works.\n");

    cpx a[8] = {0, 1, cpx(1,3), cpx(0,5), 1, 0, 2, 0};
    cpx b[8] = {1, cpx(0,-2), cpx(0,1), 3, -1, -3, 1, -2};
    cpx A[8];
    cpx B[8];
    FFT(a, A, 1, 8, 1);
    FFT(b, B, 1, 8, 1);

    for(int i = 0 ; i < 8 ; i++)
    {
        printf("%7.21f%7.21f", A[i].a, A[i].b);
    }
    printf("\n");
    for(int i = 0 ; i < 8 ; i++)
    {
        cpx Ai(0,0);
        for(int j = 0 ; j < 8 ; j++)
        {
            Ai = Ai + a[j] * EXP(j * i * two_pi / 8);
        }
        printf("%7.21f%7.21f", Ai.a, Ai.b);
    }
    printf("\n");

    cpx AB[8];
    for(int i = 0 ; i < 8 ; i++)
        AB[i] = A[i] * B[i];
    cpx aconvb[8];
    FFT(AB, aconvb, 1, 8, -1);
    for(int i = 0 ; i < 8 ; i++)
        aconvb[i] = aconvb[i] / 8;
    for(int i = 0 ; i < 8 ; i++)
    {
        printf("%7.21f%7.21f", aconvb[i].a, aconvb[i].b);
    }
    printf("\n");
    for(int i = 0 ; i < 8 ; i++)
    {
        cpx aconvbi(0,0);
        for(int j = 0 ; j < 8 ; j++)
        {
            aconvbi = aconvbi + a[j] * b[(8 + i - j) % 8];
        }
        printf("%7.21f%7.21f", aconvbi.a, aconvbi.b);
    }
    printf("\n");

    return 0;
}

```

3.5 Simplex algorithm

```

// Two-phase simplex algorithm for solving linear programs of the form
//
//      maximize      c^T x
//      subject to    Ax <= b
//                   x >= 0
//
// INPUT: A -- an m x n matrix
//        b -- an m-dimensional vector
//        c -- an n-dimensional vector
//        x -- a vector where the optimal solution will be stored

```

```
//
// OUTPUT: value of the optimal solution (infinity if unbounded
//         above, nan if infeasible)
//
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).

#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>

using namespace std;

typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;

const DOUBLE EPS = 1e-9;

struct LPSolver {
    int m, n;
    VI B, N;
    VVD D;

    LPSolver(const VVD &A, const VD &b, const VD &c) :
        m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
        for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] = A[i][j];
        for (int i = 0; i < m; i++) { B[i] = n + 1; D[i][n] = -1; D[i][n + 1] = b[i]; }
        for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
        N[n] = -1; D[m + 1][n] = 1;
    }

    void Pivot(int r, int s) {
        double inv = 1.0 / D[r][s];
        for (int i = 0; i < m + 2; i++) if (i != r)
            for (int j = 0; j < n + 2; j++) if (j != s)
                D[i][j] -= D[r][j] * D[i][s] * inv;
        for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
        for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;
        D[r][s] = inv;
        swap(B[r], N[s]);
    }

    bool Simplex(int phase) {
        int x = phase == 1 ? m + 1 : m;
        while (true) {
            int s = -1;
            for (int j = 0; j <= n; j++) {
                if (phase == 2 && N[j] == -1) continue;
                if (s == -1 || D[x][j] < D[x][s] || D[x][j] == D[x][s] && N[j] < N[s]) s = j;
            }
            if (D[x][s] > -EPS) return true;
            int r = -1;
            for (int i = 0; i < m; i++) {
                if (D[i][s] < EPS) continue;
                if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
                    (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s]) && B[i] < B[r]) r = i;
            }
            if (r == -1) return false;
            Pivot(r, s);
        }
    }

    DOUBLE Solve(VD &x) {
        int r = 0;
        for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
        if (D[r][n + 1] < -EPS) {
            Pivot(r, n);
            if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return numeric_limits<DOUBLE>::infinity();
            for (int i = 0; i < m; i++) if (B[i] == -1) {
                int s = -1;
                for (int j = 0; j <= n; j++)
                    if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i][s] && N[j] < N[s]) s = j;
                Pivot(i, s);
            }
        }
        if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
        x = VD(n);
        for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
        return D[m][n + 1];
    }
};

int main() {
    const int m = 4;
    const int n = 3;
    DOUBLE _A[m][n] = {
        { 6, -1, 0 },

```

```
{ -1, -5, 0 },
{ 1, 5, 1 },
{ -1, -5, -1 }
};
DOUBLE _b[m] = { 10, -4, 5, -5 };
DOUBLE _c[n] = { 1, -1, 0 };

VVD A(m);
VD b(_b, _b + m);
VD c(_c, _c + n);
for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);

LPSolver solver(A, b, c);
VD x;
DOUBLE value = solver.Solve(x);

cerr << "VALUE: " << value << endl; // VALUE: 1.29032
cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
cerr << endl;
return 0;
}
```

4 Graph algorithms

4.1 Bellman-Ford shortest paths with negative edge weights

```
// This function runs the Bellman-Ford algorithm for single source
// shortest paths with negative edge weights. The function returns
// false if a negative weight cycle is detected. Otherwise, the
// function returns true and dist[i] is the length of the shortest
// path from start to i.
//
// Running time: O(|V|^3)
//
// INPUT: start, w[i][j] = cost of edge from i to j
// OUTPUT: dist[i] = min weight path from start to i
//         prev[i] = previous node on the best path from the
//         start node
```

```
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>

using namespace std;

typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;

typedef vector<int> VI;
typedef vector<VI> VVI;

bool BellmanFord(const VVT &w, VI &dist, VI &prev, int start) {
    int n = w.size();
    prev = VI(n, -1);
    dist = VT(n, 1000000000);
    dist[start] = 0;

    for (int k = 0; k < n; k++) {
        for (int i = 0; i < n; i++) {
            for (int j = 0; j < n; j++) {
                if (dist[j] > dist[i] + w[i][j]) {
                    if (k == n - 1) return false;
                    dist[j] = dist[i] + w[i][j];
                    prev[j] = i;
                }
            }
        }
    }

    return true;
}
```

4.2 Fast Dijkstra's algorithm


```
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
//
// Running time:  $O(|E| \log |V|)$ 

#include <queue>
#include <cstdio>

using namespace std;
const int INF = 2000000000;
typedef pair<int, int> PII;

int main() {
    int N, s, t;
    scanf("%d%d%d", &N, &s, &t);
    vector<vector<PII>> edges(N);
    for (int i = 0; i < N; i++) {
        int M;
        scanf("%d", &M);
        for (int j = 0; j < M; j++) {
            int vertex, dist;
            scanf("%d%d", &vertex, &dist);
            edges[i].push_back(make_pair(dist, vertex)); // note order of arguments here
        }
    }

    // use priority queue in which top element has the "smallest" priority
    priority_queue<PII, vector<PII>, greater<PII>> Q;
    vector<int> dist(N, INF), dad(N, -1);
    Q.push(make_pair(0, s));
    dist[s] = 0;
    while (!Q.empty()) {
        PII p = Q.top();
        Q.pop();
        int here = p.second;
        if (here == t) break;
        if (dist[here] != p.first) continue;

        for (vector<PII>::iterator it = edges[here].begin(); it != edges[here].end(); it++) {
            if (dist[here] + it->first < dist[it->second]) {
                dist[it->second] = dist[here] + it->first;
                dad[it->second] = here;
                Q.push(make_pair(dist[it->second], it->second));
            }
        }
    }

    printf("%d\n", dist[t]);
    if (dist[t] < INF)
        for (int i = t; i != -1; i = dad[i])
            printf("%d%c", i, (i == s ? '\n' : ' '));

    return 0;
}

/*
Sample input:
5 0 4
2 1 2 3 1
2 2 4 4 5
3 1 4 3 3 4 1
2 0 1 2 3
2 1 5 2 1

Expected:
5
4 2 3 0
*/
```

4.3 Strongly connected components

```
#include <memory.h>
struct edge{int e, nxt;};
int V, E;
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group_cnt, group_num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill_forward(int x)
{
    int i;
    v[x]=true;
    for(i=sp[x];i=i=e[i].nxt; i!=e[i].e) fill_forward(e[i].e);
    stk[++stk[0]]=x;
}
void fill_backward(int x)
```

```
{
    int i;
    v[x]=false;
    group_num[x]=group_cnt;
    for(i=spr[x];i=i=er[i].nxt; i!=er[i].e) fill_backward(er[i].e);
}
void add_edge(int v1, int v2) //add edge v1->v2
{
    e[++E].e=v2; e[E].nxt=sp[v1]; sp[v1]=E;
    er[E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
}
void SCC()
{
    int i;
    stk[0]=0;
    memset(v, false, sizeof(v));
    for(i=1;i<=V;i++) if(!v[i]) fill_forward(i);
    group_cnt=0;
    for(i=stk[0];i>=1;i--) if(v[stk[i]]){group_cnt++; fill_backward(stk[i]);}
}
```

4.4 Eulerian path

```
struct Edge;
typedef list<Edge>::iterator iter;

struct Edge
{
    int next_vertex;
    iter reverse_edge;

    Edge(int next_vertex)
        :next_vertex(next_vertex)
    { }
};

const int max_vertices = ;
int num_vertices;
list<Edge> adj[max_vertices]; // adjacency list

vector<int> path;

void find_path(int v)
{
    while(adj[v].size() > 0)
    {
        int vn = adj[v].front().next_vertex;
        adj[vn].erase(adj[v].front().reverse_edge);
        adj[v].pop_front();
        find_path(vn);
    }
    path.push_back(v);
}

void add_edge(int a, int b)
{
    adj[a].push_front(Edge(b));
    iter ita = adj[a].begin();
    adj[b].push_front(Edge(a));
    iter itb = adj[b].begin();
    ita->reverse_edge = itb;
    itb->reverse_edge = ita;
}
```

4.5 Kruskal's algorithm

```
/*
Uses Kruskal's Algorithm to calculate the weight of the minimum spanning
forest (union of minimum spanning trees of each connected component) of
a possibly disjoint graph, given in the form of a matrix of edge weights
(-1 if no edge exists). Returns the weight of the minimum spanning
forest (also calculates the actual edges - stored in T). Note: uses a
disjoint-set data structure with amortized (effectively) constant time per
union/find. Runs in  $O(E \cdot \log(E))$  time.
*/

#include <iostream>
#include <vector>
#include <algorithm>
#include <queue>

using namespace std;
```

```

typedef int T;

struct edge
{
    int u, v;
    T d;
};

struct edgeCmp
{
    int operator()(const edge& a, const edge& b) { return a.d > b.d; }
};

int find(vector<int>& C, int x) { return (C[x] == x) ? x : C[x] = find(C, C[x]); }

T Kruskal(vector<vector<T>> &w)
{
    int n = w.size();
    T weight = 0;

    vector<int> C(n), R(n);
    for(int i=0; i<n; i++) { C[i] = i; R[i] = 0; }

    vector<edge> T;
    priority_queue<edge, vector<edge>, edgeCmp> E;

    for(int i=0; i<n; i++)
        for(int j=i+1; j<n; j++)
            if(w[i][j] >= 0)
            {
                edge e;
                e.u = i; e.v = j; e.d = w[i][j];
                E.push(e);
            }

    while(T.size() < n-1 && !E.empty())
    {
        edge cur = E.top(); E.pop();

        int uc = find(C, cur.u), vc = find(C, cur.v);
        if(uc != vc)
        {
            T.push_back(cur); weight += cur.d;

            if(R[uc] > R[vc]) C[vc] = uc;
            else if(R[vc] > R[uc]) C[uc] = vc;
            else { C[vc] = uc; R[uc]++; }
        }
    }

    return weight;
}

int main()
{
    int wa[6][6] = {
        { 0, -1, 2, -1, 7, -1 },
        { -1, 0, -1, 2, -1, -1 },
        { 2, -1, 0, -1, 8, 6 },
        { -1, 2, -1, 0, -1, -1 },
        { 7, -1, 8, -1, 0, 4 },
        { -1, -1, 6, -1, 4, 0 };

    vector<vector<int>> w(6, vector<int>(6));

    for(int i=0; i<6; i++)
        for(int j=0; j<6; j++)
            w[i][j] = wa[i][j];

    cout << Kruskal(w) << endl;
    cin >> wa[0][0];
}

```

4.6 Minimum spanning trees

```

// This function runs Prim's algorithm for constructing minimum
// weight spanning trees.
//
// Running time: O(|V|^2)
//
// INPUT:   w[i][j] = cost of edge from i to j
//
// NOTE: Make sure that w[i][j] is nonnegative and
// symmetric. Missing edges should be given -1
// weight.
//
// OUTPUT:  edges = list of pair<int,int> in minimum spanning tree

```

```

// return total weight of tree

#include <iostream>
#include <queue>
#include <cmath>
#include <vector>

using namespace std;

typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;

typedef vector<int> VI;
typedef vector<VI> VVI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;

T Prim(const VVT &w, VPII &edges) {
    int n = w.size();
    VI found(n);
    VI prev(n, -1);
    VT dist(n, 1000000000);
    int here = 0;
    dist[here] = 0;

    while (here != -1) {
        found[here] = true;
        int best = -1;
        for (int k = 0; k < n; k++) if (!found[k]) {
            if (w[here][k] != -1 && dist[k] > w[here][k]) {
                dist[k] = w[here][k];
                prev[k] = here;
            }
            if (best == -1 || dist[k] < dist[best]) best = k;
        }
        here = best;
    }

    T tot_weight = 0;
    for (int i = 0; i < n; i++) if (prev[i] != -1) {
        edges.push_back(make_pair(prev[i], i));
        tot_weight += w[prev[i]][i];
    }

    return tot_weight;
}

int main() {
    int ww[5][5] = {
        { 0, 400, 400, 300, 600 },
        { 400, 0, 3, -1, 7 },
        { 400, 3, 0, 2, 0 },
        { 300, -1, 2, 0, 5 },
        { 600, 7, 0, 5, 0 }
    };
    VVT w(5, VT(5));
    for (int i = 0; i < 5; i++)
        for (int j = 0; j < 5; j++)
            w[i][j] = ww[i][j];

    // expected: 305
    // 2 1
    // 3 2
    // 0 3
    // 2 4

    VPII edges;
    cout << Prim(w, edges) << endl;
    for (int i = 0; i < edges.size(); i++)
        cout << edges[i].first << " " << edges[i].second << endl;
}

```

5 Data structures

5.1 Suffix array

```

// Suffix array construction in O(L log^2 L) time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
//
// INPUT:   string s
//
// OUTPUT:  array suffix[] such that suffix[i] = index (from 0 to L-1)
// of substring s[i...L-1] in the list of sorted suffixes.
// That is, if we take the inverse of the permutation suffix[],

```

```
// we get the actual suffix array.

#include <vector>
#include <iostream>
#include <string>

using namespace std;

struct SuffixArray {
    const int L;
    string s;
    vector<vector<int>> > P;
    vector<pair<pair<int,int>,int>> > M;

    SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L) {
        for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
        for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
            P.push_back(vector<int>(L, 0));
            for (int i = 0; i < L; i++)
                M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
            sort(M.begin(), M.end());
            for (int i = 0; i < L; i++)
                P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i;
        }
    }

    vector<int> GetSuffixArray() { return P.back(); }

    // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
    int LongestCommonPrefix(int i, int j) {
        int len = 0;
        if (i == j) return L - i;
        for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
            if (P[k][i] == P[k][j]) {
                i += 1 << k;
                j += 1 << k;
                len += 1 << k;
            }
        }
        return len;
    }
};

// BEGIN CUT
// The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
int main() {
    int T;
    cin >> T;
    for (int caseno = 0; caseno < T; caseno++) {
        string s;
        cin >> s;
        SuffixArray array(s);
        vector<int> v = array.GetSuffixArray();
        int bestlen = -1, bestpos = -1, bestcount = 0;
        for (int i = 0; i < s.length(); i++) {
            int len = 0, count = 0;
            for (int j = i+1; j < s.length(); j++) {
                int l = array.LongestCommonPrefix(i, j);
                if (l >= len) {
                    if (l > len) count = 2; else count++;
                    len = l;
                }
            }
            if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen) > s.substr(i, len)) {
                bestlen = len;
                bestcount = count;
                bestpos = i;
            }
        }
        if (bestlen == 0) {
            cout << "No repetitions found!" << endl;
        } else {
            cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;
        }
    }
}

#else
// END CUT
int main() {

    // bobocel is the 0'th suffix
    // obocel is the 5'th suffix
    // bocel is the 1'st suffix
    // ocel is the 6'th suffix
    // cel is the 2'nd suffix
    // el is the 3'rd suffix
    // l is the 4'th suffix
    SuffixArray suffix("bobocel");
    vector<int> v = suffix.GetSuffixArray();
```

```
// Expected output: 0 5 1 6 2 3 4
//
for (int i = 0; i < v.size(); i++) cout << v[i] << " ";
cout << endl;
cout << suffix.LongestCommonPrefix(0, 2) << endl;
}
// BEGIN CUT
#endif
// END CUT
```

5.2 Binary Indexed (Fenwick) Tree

```
#include <iostream>
using namespace std;

#define LOGSZ 17

int tree[(1<<LOGSZ)+1];
int N = (1<<LOGSZ);

// add v to value at x
void set(int x, int v) {
    while(x <= N) {
        tree[x] += v;
        x += (x & -x);
    }
}

// get cumulative sum up to and including x
int get(int x) {
    int res = 0;
    while(x) {
        res += tree[x];
        x -= (x & -x);
    }
    return res;
}

// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
    int idx = 0, mask = N;
    while(mask && idx < N) {
        int t = idx + mask;
        if(x >= tree[t]) {
            idx = t;
            x -= tree[t];
        }
        mask >>= 1;
    }
    return idx;
}
```

5.3 Union-find disjoint sets

```
#include <iostream>
#include <vector>
using namespace std;
struct UnionFind {
    vector<int> C;
    UnionFind(int n) : C(n) { for (int i = 0; i < n; i++) C[i] = i; }
    int find(int x) { return (C[x] == x) ? x : C[x] = find(C[x]); }
    void merge(int x, int y) { C[find(x)] = find(y); }
};
int main()
{
    int n = 5;
    UnionFind uf(n);
    uf.merge(0, 2);
    uf.merge(1, 0);
    uf.merge(3, 4);
    for (int i = 0; i < n; i++) cout << i << " " << uf.find(i) << endl;
    return 0;
}
```

5.4 KD-tree

```
// -----
// A straightforward, but probably sub-optimal KD-tree implementation
```

```
// that's probably good enough for most things (current it's a
// 2D-tree)
//
// - constructs from n points in O(n lg^2 n) time
// - handles nearest-neighbor query in O(lg n) if points are well
//   distributed
// - worst case for nearest-neighbor may be linear in pathological
//   case
//
// Sonny Chan, Stanford University, April 2009
// -----

#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>

using namespace std;

// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();

// point structure for 2D-tree, can be extended to 3D
struct point {
    ntype x, y;
    point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
};

bool operator==(const point &a, const point &b)
{
    return a.x == b.x && a.y == b.y;
}

// sorts points on x-coordinate
bool on_x(const point &a, const point &b)
{
    return a.x < b.x;
}

// sorts points on y-coordinate
bool on_y(const point &a, const point &b)
{
    return a.y < b.y;
}

// squared distance between points
ntype pdist2(const point &a, const point &b)
{
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
}

// bounding box for a set of points
struct bbox
{
    ntype x0, x1, y0, y1;

    bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}

    // computes bounding box from a bunch of points
    void compute(const vector<point> &v) {
        for (int i = 0; i < v.size(); ++i) {
            x0 = min(x0, v[i].x);    x1 = max(x1, v[i].x);
            y0 = min(y0, v[i].y);    y1 = max(y1, v[i].y);
        }
    }

    // squared distance between a point and this bbox, 0 if inside
    ntype distance(const point &p) {
        if (p.x < x0) {
            if (p.y < y0)    return pdist2(point(x0, y0), p);
            else if (p.y > y1) return pdist2(point(x0, y1), p);
            else            return pdist2(point(x0, p.y), p);
        }
        else if (p.x > x1) {
            if (p.y < y0)    return pdist2(point(x1, y0), p);
            else if (p.y > y1) return pdist2(point(x1, y1), p);
            else            return pdist2(point(x1, p.y), p);
        }
        else {
            if (p.y < y0)    return pdist2(point(p.x, y0), p);
            else if (p.y > y1) return pdist2(point(p.x, y1), p);
            else            return 0;
        }
    }
};

// stores a single node of the kd-tree, either internal or leaf
struct kdnnode
{
    bool leaf;        // true if this is a leaf node (has one point)
```

```
    point pt;        // the single point of this is a leaf
    bbox bound;       // bounding box for set of points in children

    kdnnode *first, *second; // two children of this kd-node

    kdnnode() : leaf(false), first(0), second(0) {}
    ~kdnnode() { if (first) delete first; if (second) delete second; }

    // intersect a point with this node (returns squared distance)
    ntype intersect(const point &p) {
        return bound.distance(p);
    }

    // recursively builds a kd-tree from a given cloud of points
    void construct(vector<point> &vp)
    {
        // compute bounding box for points at this node
        bound.compute(vp);

        // if we're down to one point, then we're a leaf node
        if (vp.size() == 1) {
            leaf = true;
            pt = vp[0];
        }
        else {
            // split on x if the bbox is wider than high (not best heuristic...)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on_x);
            // otherwise split on y-coordinate
            else
                sort(vp.begin(), vp.end(), on_y);

            // divide by taking half the array for each child
            // (not best performance if many duplicates in the middle)
            int half = vp.size()/2;
            vector<point> vl(vp.begin(), vp.begin()+half);
            vector<point> vr(vp.begin()+half, vp.end());
            first = new kdnnode();    first->construct(vl);
            second = new kdnnode();   second->construct(vr);
        }
    }
};

// simple kd-tree class to hold the tree and handle queries
struct kdtree
{
    kdnnode *root;

    // constructs a kd-tree from a points (copied here, as it sorts them)
    kdtree(const vector<point> &vp) {
        vector<point> v(vp.begin(), vp.end());
        root = new kdnnode();
        root->construct(v);
    }

    ~kdtree() { delete root; }

    // recursive search method returns squared distance to nearest point
    ntype search(kdnnode *node, const point &p)
    {
        if (node->leaf) {
            // commented special case tells a point not to find itself
            // if (p == node->pt) return sentry;
            // else
            return pdist2(p, node->pt);
        }

        ntype bfirst = node->first->intersect(p);
        ntype bsecond = node->second->intersect(p);

        // choose the side with the closest bounding box to search first
        // (note that the other side is also searched if needed)
        if (bfirst < bsecond) {
            ntype best = search(node->first, p);
            if (bsecond < best)
                best = min(best, search(node->second, p));
            return best;
        }
        else {
            ntype best = search(node->second, p);
            if (bfirst < best)
                best = min(best, search(node->first, p));
            return best;
        }
    }

    // squared distance to the nearest
    ntype nearest(const point &p) {
        return search(root, p);
    }
};

// -----
```

```
// some basic test code here

int main()
{
    // generate some random points for a kd-tree
    vector<point> vp;
    for (int i = 0; i < 100000; ++i) {
        vp.push_back(point(rand()%100000, rand()%100000));
    }
    kdtree tree(vp);

    // query some points
    for (int i = 0; i < 10; ++i) {
        point q(rand()%100000, rand()%100000);
        cout << "Closest squared distance to (" << q.x << ", " << q.y << ") "
              << " is " << tree.nearest(q) << endl;
    }

    return 0;
}

// -----
```

5.5 Splay tree

```
#include <cstdio>
#include <algorithm>
using namespace std;

const int N_MAX = 130010;
const int oo = 0x3f3f3f3f;
struct Node
{
    Node *ch[2], *pre;
    int val, size;
    bool isTurned;
} nodePool[N_MAX], *null, *root;

Node *allocNode(int val)
{
    static int freePos = 0;
    Node *x = &nodePool[freePos++];
    x->val = val, x->isTurned = false;
    x->ch[0] = x->ch[1] = x->pre = null;
    x->size = 1;
    return x;
}

inline void update(Node *x)
{
    x->size = x->ch[0]->size + x->ch[1]->size + 1;
}

inline void makeTurned(Node *x)
{
    if (x == null)
        return;
    swap(x->ch[0], x->ch[1]);
    x->isTurned ^= 1;
}

inline void pushDown(Node *x)
{
    if (x->isTurned)
    {
        makeTurned(x->ch[0]);
        makeTurned(x->ch[1]);
        x->isTurned ^= 1;
    }
}

inline void rotate(Node *x, int c)
{
    Node *y = x->pre;
    x->pre = y->pre;
    if (y->pre != null)
        y->pre->ch[y == y->pre->ch[1]] = x;
    y->ch[!c] = x->ch[c];
    if (x->ch[c] != null)
        x->ch[c]->pre = y;
    x->ch[c] = y, y->pre = x;
    update(y);
    if (y == root)
        root = x;
}

void splay(Node *x, Node *p)
```

```
{
    while (x->pre != p)
    {
        if (x->pre->pre == p)
            rotate(x, x == x->pre->ch[0]);
        else
        {
            Node *y = x->pre, *z = y->pre;
            if (y == z->ch[0])
            {
                if (x == y->ch[0])
                    rotate(y, 1), rotate(x, 1);
                else
                    rotate(x, 0), rotate(x, 1);
            }
            else
            {
                if (x == y->ch[1])
                    rotate(y, 0), rotate(x, 0);
                else
                    rotate(x, 1), rotate(x, 0);
            }
        }
    }
    update(x);
}

void select(int k, Node *fa)
{
    Node *now = root;
    while (1)
    {
        pushDown(now);
        int tmp = now->ch[0]->size + 1;
        if (tmp == k)
            break;
        else if (tmp < k)
            now = now->ch[1], k -= tmp;
        else
            now = now->ch[0];
    }
    splay(now, fa);
}

Node *makeTree(Node *p, int l, int r)
{
    if (l > r)
        return null;
    int mid = (l + r) / 2;
    Node *x = allocNode(mid);
    x->pre = p;
    x->ch[0] = makeTree(x, l, mid - 1);
    x->ch[1] = makeTree(x, mid + 1, r);
    update(x);
    return x;
}

int main()
{
    int n, m;
    null = allocNode(0);
    null->size = 0;
    root = allocNode(0);
    root->ch[1] = allocNode(oo);
    root->ch[1]->pre = root;
    update(root);

    scanf("%d%d", &n, &m);
    root->ch[1]->ch[0] = makeTree(root->ch[1], 1, n);
    splay(root->ch[1]->ch[0], null);

    while (m --)
    {
        int a, b;
        scanf("%d%d", &a, &b);
        a ++, b ++;
        select(a - 1, null);
        select(b + 1, root);
        makeTurned(root->ch[1]->ch[0]);
    }

    for (int i = 1; i <= n; i ++ )
    {
        select(i + 1, null);
        printf("%d ", root->val);
    }
}
```

5.6 Segment tree

```
// This is a solution for UVa 12299 - RMQ with Shifts. In this problem, we are
// given an array of integers A and a list of queries of 2 types:
// 1. query(L, R) returns the minimum value in A[L...R].
// 2. shift(i_1, i_2, ..., i_k) performs a left "circular shift" of A[i_1],
//    A[i_2], ..., A[i_k]. For example, if A = {6, 2, 4, 8, 5, 1, 4}, then
//    shift(2, 4, 5, 7) yields 6, 8, 4, 5, 4, 1, 2. After that, shift(1, 2)
//    yields 8, 6, 4, 5, 4, 1, 2.

#include <bits/stdc++.h>

#define LEFT_CHILD(x) ((x << 1) + 1)
#define RIGHT_CHILD(x) ((x << 1) + 2)
#define INF 999999999
#define MAX_N 100005

using namespace std;

int queryParams[35], numbers[MAX_N], queryParamsCount;
string inputQuery;

struct node_t {
    int left, right, minValue;
} nodes[4 * MAX_N];

void build(int left, int right, int x = 0) {
    nodes[x].left = left;
    nodes[x].right = right;
    if (left == right) {
        nodes[x].minValue = numbers[left];
        return;
    }
    int mid = (left + right) / 2;
    build(left, mid, LEFT_CHILD(x));
    build(mid + 1, right, RIGHT_CHILD(x));
    nodes[x].minValue = min(nodes[LEFT_CHILD(x)].minValue, nodes[RIGHT_CHILD(x)].minValue);
}

void parseInputQuery() {
    queryParamsCount = 0;
    queryParams[queryParamsCount] = 0;
    // Parse input query starting from after opening parenthesis i.e. index 6.
    for (int i = 6; i < inputQuery.length(); i++) {
        if ('0' <= inputQuery[i] && inputQuery[i] <= '9') {
            queryParams[queryParamsCount] = queryParams[queryParamsCount] * 10 + (int)(inputQuery[i] - '0');
        }
        else {
            queryParams[++queryParamsCount] = 0; // next query parameter!
        }
    }
}

int query(int left, int right, int x = 0) {
    if (left <= nodes[x].left && nodes[x].right <= right) {
        return nodes[x].minValue;
    }
    int mid = (nodes[x].left + nodes[x].right) / 2;
    int ans = INF;
    if (left <= mid) {
        ans = min(ans, query(left, right, LEFT_CHILD(x)));
    }
    if (mid < right) {
        ans = min(ans, query(left, right, RIGHT_CHILD(x)));
    }
    return ans;
}

void setValue(int position, int newValue, int x = 0) {
    if (nodes[x].left == position && nodes[x].right == position) {
        nodes[x].minValue = newValue;
        return;
    }
    int mid = (nodes[x].left + nodes[x].right) / 2;
    if (position <= mid) {
        setValue(position, newValue, LEFT_CHILD(x));
    }
    if (mid < position) {
        setValue(position, newValue, RIGHT_CHILD(x));
    }
    nodes[x].minValue = min(nodes[LEFT_CHILD(x)].minValue, nodes[RIGHT_CHILD(x)].minValue);
}

int main()
{
    int n, q, placeholder;
    cin >> n >> q;
```

```
for (int i = 1; i <= n; i++) {
    cin >> numbers[i];
}
build(1, n);
while (q--) {
    cin >> inputQuery;
    if (inputQuery[0] == 'q') { // perform operation query.
        parseInputQuery();
        cout << query(queryParams[0], queryParams[1]) << '\n';
    }
    else if (inputQuery[0] == 's') { // perform operation shift.
        parseInputQuery();
        placeholder = numbers[queryParams[0]];
        for (int i = 1; i < queryParamsCount; i++) {
            setValue(queryParams[i - 1], numbers[queryParams[i]]);
            numbers[queryParams[i - 1]] = numbers[queryParams[i]];
        }
        setValue(queryParams[queryParamsCount - 1], placeholder);
        numbers[queryParams[queryParamsCount - 1]] = placeholder;
    }
}
return 0;
}
```

5.7 Lowest common ancestor

```
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;

vector<int> children[max_nodes]; // children[i] contains the children of node i
int A[max_nodes][log_max_nodes+1]; // A[i][j] is the 2^j-th ancestor of node i, or -1 if that
// ancestor does not exist
int L[max_nodes]; // L[i] is the distance between node i and the root

// floor of the binary logarithm of n
int lb(unsigned int n)
{
    if (n==0)
        return -1;
    int p = 0;
    if (n >= 1<<16) { n >>= 16; p += 16; }
    if (n >= 1<<8) { n >>= 8; p += 8; }
    if (n >= 1<<4) { n >>= 4; p += 4; }
    if (n >= 1<<2) { n >>= 2; p += 2; }
    if (n >= 1<<1) { p += 1; }
    return p;
}

void DFS(int i, int l)
{
    L[i] = l;
    for (int j = 0; j < children[i].size(); j++)
        DFS(children[i][j], l+1);
}

int LCA(int p, int q)
{
    // ensure node p is at least as deep as node q
    if (L[p] < L[q])
        swap(p, q);

    // "binary search" for the ancestor of node p situated on the same level as q
    for (int i = log_num_nodes; i >= 0; i--)
        if (L[p] - (1<<i) >= L[q])
            p = A[p][i];

    if (p == q)
        return p;

    // "binary search" for the LCA
    for (int i = log_num_nodes; i >= 0; i--)
        if (A[p][i] != -1 && A[p][i] != A[q][i])
        {
            p = A[p][i];
            q = A[q][i];
        }

    return A[p][0];
}

int main(int argc, char* argv[])
{
    // read num_nodes, the total number of nodes
    log_num_nodes = lb(num_nodes);

    for (int i = 0; i < num_nodes; i++)
    {
        int p;
```

```

// read p, the parent of node i or -1 if node i is the root
A[i][0] = p;
if (p != -1)
    children[p].push_back(i);
else
    root = i;
}

// precompute A using dynamic programming
for (int j = 1; j <= log_num_nodes; j++)
    for (int i = 0; i < num_nodes; i++)
        if (A[i][j-1] != -1)
            A[i][j] = A[A[i][j-1]][j-1];
        else
            A[i][j] = -1;

// precompute L
DFS(root, 0);

return 0;
}

```

6 Miscellaneous

6.1 Longest increasing subsequence

```

// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
//
// Running time: O(n log n)
//
// INPUT: a vector of integers
// OUTPUT: a vector containing the longest increasing subsequence

#include <iostream>
#include <vector>
#include <algorithm>

using namespace std;

typedef vector<int> VI;
typedef pair<int,int> PII;
typedef vector<PII> VPPII;

#define STRICTLY_INCREASNG

VI LongestIncreasingSubsequence (VI v) {
    VPPII best;
    VI dad(v.size(), -1);

    for (int i = 0; i < v.size(); i++) {
#ifdef STRICTLY_INCREASNG
        PII item = make_pair(v[i], 0);
        VPPII::iterator it = lower_bound(best.begin(), best.end(), item);
        item.second = i;
#else
        PII item = make_pair(v[i], i);
        VPPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
        if (it == best.end()) {
            dad[i] = (best.size() == 0 ? -1 : best.back().second);
            best.push_back(item);
        } else {
            dad[i] = it == best.begin() ? -1 : prev(it)->second;
            *it = item;
        }
    }

    VI ret;
    for (int i = best.back().second; i >= 0; i = dad[i])
        ret.push_back(v[i]);
    reverse(ret.begin(), ret.end());
    return ret;
}

```

6.2 Knuth-Morris-Pratt

```
/*
```

Finds all occurrences of the pattern string *p* within the text string *t*. Running time is $O(n + m)$, where *n* and *m* are the lengths of *p* and *t*, respectively.

```

*/

#include <iostream>
#include <string>
#include <vector>

using namespace std;

typedef vector<int> VI;

void buildPi(string& p, VI& pi)
{
    pi = VI(p.length());
    int k = -2;
    for (int i = 0; i < p.length(); i++) {
        while (k >= -1 && p[k+1] != p[i])
            k = (k == -1) ? -2 : pi[k];
        pi[i] = ++k;
    }
}

int KMP (string& t, string& p)
{
    VI pi;
    buildPi(p, pi);
    int k = -1;
    for (int i = 0; i < t.length(); i++) {
        while (k >= -1 && p[k+1] != t[i])
            k = (k == -1) ? -2 : pi[k];
        k++;
        if (k == p.length() - 1) {
            // p matches t[i-m+1, ..., i]
            cout << "matched at index " << i-k << " : ";
            cout << t.substr(i-k, p.length()) << endl;
            k = (k == -1) ? -2 : pi[k];
        }
    }
    return 0;
}

int main()
{
    string a = "AABAACAADAABAABA", b = "AABA";
    KMP(a, b); // expected matches at: 0, 9, 12
    return 0;
}

```

6.3 Dates

```

// Routines for performing computations on dates. In these routines,
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
// integers.

#include <iostream>
#include <string>

using namespace std;

string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};

// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y) {
    return
        1461 * (y + 4800 + (m - 14) / 12) / 4 +
        367 * (m - 2 - (m - 14) / 12 + 12) / 12 -
        3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
        d - 32075;
}

// converts integer (Julian day number) to Gregorian date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
    int x, n, i, j;

    x = jd + 68569;
    n = 4 * x / 146097;
    x -= (146097 * n + 3) / 4;
    i = (4000 * (x + 1)) / 1461001;
    x -= 1461 * i / 4 - 31;
    j = 80 * x / 2447;
    d = x - 2447 * j / 80;
    x = j / 11;
    m = j + 2 - 12 * x;
    y = 100 * (n - 49) + i + x;
}

```

```
// converts integer (Julian day number) to day of week
string intToDay (int jd){
    return dayOfWeek[jd % 7];
}

int main (int argc, char **argv){
    int jd = dateToInt (3, 24, 2004);
    int m, d, y;
    intToDate (jd, m, d, y);
    string day = intToDay (jd);

    // expected output:
    // 2453089
    // 3/24/2004
    // Wed
    cout << jd << endl
         << m << "/" << d << " " << y << endl
         << day << endl;
}
```

6.4 Prime numbers

```
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool isPrimeSlow (LL x)
{
    if (x<=1) return false;
    if (x<=3) return true;
    if (!(x%2) || !(x%3)) return false;
    LL s=(LL)(sqrt((double)(x))+EPS);
    for (LL i=5; i<=s; i+=6)
    {
        if (!(x%i) || !(x%(i+2))) return false;
    }
    return true;
}

// Primes less than 1000:
// 2 3 5 7 11 13 17 19 23 29 31 37
// 41 43 47 53 59 61 67 71 73 79 83 89
// 97 101 103 107 109 113 127 131 137 139 149 151
// 157 163 167 173 179 181 191 193 197 199 211 223
// 227 229 233 239 241 251 257 263 269 271 277 281
// 283 293 307 311 313 317 331 337 347 349 353 359
// 367 373 379 383 389 397 401 409 419 421 431 433
// 439 443 449 457 461 463 467 479 487 491 499 503
// 509 521 523 541 547 557 563 569 571 577 587 593
// 599 601 607 613 617 619 631 641 643 647 653 659
// 661 673 677 683 691 701 709 719 727 733 739 743
// 751 757 761 769 773 787 797 809 811 821 823 827
// 829 839 853 857 859 863 877 881 883 887 907 911
// 919 929 937 941 947 953 967 971 977 983 991 997

// Other primes:
// The largest prime smaller than 10 is 7.
// The largest prime smaller than 100 is 97.
// The largest prime smaller than 1000 is 997.
// The largest prime smaller than 10000 is 9973.
// The largest prime smaller than 100000 is 99991.
// The largest prime smaller than 1000000 is 999983.
// The largest prime smaller than 10000000 is 9999991.
// The largest prime smaller than 100000000 is 99999989.
// The largest prime smaller than 1000000000 is 999999937.
// The largest prime smaller than 10000000000 is 9999999967.
// The largest prime smaller than 100000000000 is 9999999997.
// The largest prime smaller than 1000000000000 is 99999999989.
// The largest prime smaller than 10000000000000 is 99999999997.
// The largest prime smaller than 100000000000000 is 99999999999.
// The largest prime smaller than 1000000000000000 is 999999999993.
// The largest prime smaller than 10000000000000000 is 9999999999989.
// The largest prime smaller than 100000000000000000 is 99999999999937.
// The largest prime smaller than 1000000000000000000 is 999999999999997.
// The largest prime smaller than 10000000000000000000 is 9999999999999989.
```

6.5 Latitude/longitude

```
/*
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
*/

#include <iostream>
```

```
#include <cmath>

using namespace std;

struct ll
{
    double r, lat, lon;
};

struct rect
{
    double x, y, z;
};

ll convert (rect& P)
{
    ll Q;
    Q.r = sqrt (P.x*P.x+P.y*P.y+P.z*P.z);
    Q.lat = 180/M_PI*asin (P.z/Q.r);
    Q.lon = 180/M_PI*acos (P.x/sqrt (P.x*P.x+P.y*P.y));

    return Q;
}

rect convert (ll& Q)
{
    rect P;
    P.x = Q.r*cos (Q.lon*M_PI/180)+cos (Q.lat*M_PI/180);
    P.y = Q.r*sin (Q.lon*M_PI/180)+cos (Q.lat*M_PI/180);
    P.z = Q.r*sin (Q.lat*M_PI/180);

    return P;
}

int main ()
{
    rect A;
    ll B;

    A.x = -1.0; A.y = 2.0; A.z = -3.0;

    B = convert (A);
    cout << B.r << " " << B.lat << " " << B.lon << endl;

    A = convert (B);
    cout << A.x << " " << A.y << " " << A.z << endl;
}
```

6.6 C++ input/output

```
#include <iostream>
#include <iomanip>

using namespace std;

int main ()
{
    // Output a specific number of digits past the decimal point,
    // in this case 5
    cout.setf (ios::fixed); cout << setprecision (5);
    cout << 100.0/7.0 << endl;
    cout.unsetf (ios::fixed);

    // Output the decimal point and trailing zeros
    cout.setf (ios::showpoint);
    cout << 100.0 << endl;
    cout.unsetf (ios::showpoint);

    // Output a '+' before positive values
    cout.setf (ios::showpos);
    cout << 100 << " " << -100 << endl;
    cout.unsetf (ios::showpos);

    // Output numerical values in hexadecimal
    cout << hex << 100 << " " << 1000 << " " << 10000 << dec << endl;
}
```

6.7 Random STL stuff

```
// Example for using stringstream and next_permutation

#include <algorithm>
#include <iostream>
```



```

#include <sstream>
#include <vector>

using namespace std;

int main(void) {
    vector<int> v;

    v.push_back(1);
    v.push_back(2);
    v.push_back(3);
    v.push_back(4);

    // Expected output: 1 2 3 4
    //                  1 2 4 3
    //                  ...
    //                  4 3 2 1
    do {
        ostringstream oss;
        oss << v[0] << " " << v[1] << " " << v[2] << " " << v[3];

        // for input from a string s,
        // istream iss(s);
        // iss >> variable;

        cout << oss.str() << endl;
    } while (next_permutation (v.begin(), v.end()));

    v.clear();

    v.push_back(1);
    v.push_back(2);
    v.push_back(1);
    v.push_back(3);

    // To use unique, first sort numbers. Then call
    // unique to place all the unique elements at the beginning
    // of the vector, and then use erase to remove the duplicate
    // elements.

    sort(v.begin(), v.end());
    v.erase(unique(v.begin(), v.end()), v.end());

    // Expected output: 1 2 3
    for (size_t i = 0; i < v.size(); i++)
        cout << v[i] << " ";
    cout << endl;
}

```

6.8 Regular expressions (Java)

```

// Code which demonstrates the use of Java's regular expression libraries.
// This is a solution for
//
// Loglan: a logical language
// http://acm.uva.es/p/v1/134.html
//
// In this problem, we are given a regular language, whose rules can be
// inferred directly from the code. For each sentence in the input, we must
// determine whether the sentence matches the regular expression or not. The
// code consists of (1) building the regular expression (which is fairly
// complex) and (2) using the regex to match sentences.

import java.util.*;
import java.util.regex.*;

public class Loglan {

    public static String BuildRegex () {

```

```

String space = " ";

String A = "[aeiou]";
String C = "[a-z&&[aeiou]]";
String MOD = "(g" + A + ")";
String BA = "(b" + A + ")";
String DA = "(d" + A + ")";
String LA = "(l" + A + ")";
String NAM = "([a-z]" + C + ")";
String PREDA = "((" + C + C + A + C + A + "|" + C + A + C + C + A + ")";

String predstring = "(" + PREDA + "(" + space + PREDA + ")*)";
String predname = "(" + LA + space + predstring + "|" + NAM + ")";
String preds = "(" + predstring + "(" + space + A + space + predstring + ")*)";
String predclaim = "(" + predname + space + BA + space + preds + "|" + DA + space +
    preds + ")";
String verbpred = "(" + MOD + space + predstring + ")";
String statement = "(" + predname + space + verbpred + space + predname + "|" +
    predname + space + verbpred + ")";
String sentence = "(" + statement + "|" + predclaim + ")";

return "^" + sentence + "$";
}

public static void main (String args[]){

    String regex = BuildRegex();
    Pattern pattern = Pattern.compile (regex);

    Scanner s = new Scanner(System.in);
    while (true) {

        // In this problem, each sentence consists of multiple lines, where the last
        // line is terminated by a period. The code below reads lines until
        // encountering a line whose final character is a '.'. Note the use of
        //
        // s.length() to get length of string
        // s.charAt () to extract characters from a Java string
        // s.trim() to remove whitespace from the beginning and end of Java string
        //
        // Other useful String manipulation methods include
        //
        // s.compareTo(t) < 0 if s < t, lexicographically
        // s.indexOf("apple") returns index of first occurrence of "apple" in s
        // s.lastIndexOf("apple") returns index of last occurrence of "apple" in s
        // s.replace(c,d) replaces occurrences of character c with d
        // s.startsWith("apple") returns (s.indexOf("apple") == 0)
        // s.toLowerCase() / s.toUpperCase() returns a new lower/uppercased string
        //
        // Integer.parseInt(s) converts s to an integer (32-bit)
        // Long.parseLong(s) converts s to a long (64-bit)
        // Double.parseDouble(s) converts s to a double

        String sentence = "";
        while (true){
            sentence = (sentence + " " + s.nextLine()).trim();
            if (sentence.equals("#")) return;
            if (sentence.charAt(sentence.length()-1) == '.' break;
        }

        // now, we remove the period, and match the regular expression

        String removed_period = sentence.substring(0, sentence.length()-1).trim();
        if (pattern.matcher (removed_period).find()){
            System.out.println ("Good");
        } else {
            System.out.println ("Bad!");
        }
    }
}

```