## NTUT\_Kn1ghts ICPC Team Notebook

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# 1 Dynamic programming algorithms

#### 1.1 Longest common subsequence

```
/*
Calculates the length of the longest common subsequence of two vectors.
Backtracks to find a single subsequence or all subsequences. Runs in
O(m*n) time except for finding all longest common subsequences, which
may be slow depending on how many there are.

*/
finclude <lostream>
finclude <vector>
finclude <set>
finclude <algorithm>
using namespace std;

typedef int T;
typedef vector<T> VT;
typedef vector<T> VT;
typedef vector<VT> VVT;
```

```
typedef vector<int> VI;
typedef vector<VI> VVI;
void backtrack(VVI& dp, VT& res, VT& A, VT& B, int i, int j)
  if(!i || !j) return;
  if(A[i-1] == B[j-1]) { res.push_back(A[i-1]); backtrack(dp, res, A, B, i-1, j-1); }
    if(dp[i][j-1] >= dp[i-1][j]) backtrack(dp, res, A, B, i, j-1);
    else backtrack(dp, res, A, B, i-1, j);
void backtrackall(VVI& dp, set<VT>& res, VT& A, VT& B, int i, int j)
  if(!i || !j) { res.insert(VI()); return; }
  if(A[i-1] == B[j-1])
    set<VT> tempres;
    backtrackall(dp, tempres, A, B, i-1, j-1);
    for(set<VT>::iterator it=tempres.begin(); it!=tempres.end(); it++)
      VT temp = *it;
      temp.push_back(A[i-1]);
      res.insert(temp);
  else
    if(dp[i][j-1] >= dp[i-1][j]) backtrackall(dp, res, A, B, i, j-1);
    if(dp[i][j-1] <= dp[i-1][j]) backtrackall(dp, res, A, B, i-1, j);</pre>
VT LCS(VT& A, VT& B)
  int n = A.size(), m = B.size();
  dp.resize(n+1);
  for (int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
  for(int i=1; i<=n; i++)</pre>
    for (int j=1; j<=m; j++)</pre>
      if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
      else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
 backtrack(dp, res, A, B, n, m);
  reverse(res.begin(), res.end());
 return res:
set<VT> LCSall(VT& A, VT& B)
  int n = A.size(), m = B.size();
  dp.resize(n+1);
  for(int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
  for(int i=1; i<=n; i++)</pre>
    for (int j=1; j<=m; j++)</pre>
      if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
  set<VT> res;
  backtrackall(dp, res, A, B, n, m);
  return res:
int main()
  int a[] = { 0, 5, 5, 2, 1, 4, 2, 3 }, b[] = { 5, 2, 4, 3, 2, 1, 2, 1, 3 };
  VI A = VI(a, a+8), B = VI(b, b+9);
  VI C = LCS(A, B);
  for(int i=0; i<C.size(); i++) cout << C[i] << " ";</pre>
  cout << endl << endl;</pre>
  set <VI> D = LCSall(A, B);
for(set<VI>::iterator it = D.begin(); it != D.end(); it++)
    for(int i=0; i<(*it).size(); i++) cout << (*it)[i] << " ";</pre>
    cout << endl;
```

#### 1.2 Max 2D range sum

```
// This is a solution for UVa 108 - Maximum Sum. In this problem, we are given a
// square matrix of integers and are asked to find a sub-matrix with the maximum
// sum. This 1D DP + greedy (Kadane's) solution takes time O(n^3).
#include <bits/stdc++.h>
using namespace std;
#define MAX_n 110
int A[MAX_n][MAX_n];
int main() {
  int n; scanf("%d", &n);
                                                       // square matrix size
  for (int i = 0; i < n; ++i)
for (int j = 0; j < n; ++j) {
      scanf("%d", &A[i][i]);
      if (j > 0) A[i][j] += A[i][j-1];
                                                       // pre-processing
  int maxSubRect = -127*100*100;
                                                       // lowest possible val
  for (int l = 0; l < n; ++1)
  for (int r = 1; r < n; ++r) {</pre>
      int subRect = 0;
      for (int row = 0; row < n; ++row) {</pre>
        // Max 1D Range Sum on columns of this row
if (1 > 0) subRect += A[row][r] - A[row][1-1];
                    subRect += A[row][r];
         else
         // Kadane's algorithm on rows
         if (subRect < 0) subRect = 0;</pre>
                                                       // restart if negative
        maxSubRect = max(maxSubRect, subRect);
  printf("%d\n", maxSubRect);
  return 0;
```

## 2 Geometry

#### 2.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// #defined.
// Running time: O(n log n)
    INPUT: a vector of input points, unordered.
    OUTPUT: a vector of points in the convex hull, counterclockwise, starting
              with bottommost/leftmost point
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
#include <map>
using namespace std:
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
  T x, y;
  PT() {}
  PT(T x, T y) : x(x), y(y) {}
  bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y,rhs.x); }</pre>
  bool operator==(const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
  return (fabs(area2(a,b,c)) < EPS && (a.x-b.x) *(c.x-b.x) <= 0 && (a.y-b.y) *(c.y-b.y) <= 0);
```

# #endif void ConvexHul.

```
void ConvexHull(vector<PT> &pts) {
 sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {
    while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();</pre>
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  pts = dn:
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE_REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
    if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
  pts = dn:
#endif
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
int main() {
 int t;
scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {</pre>
    int n;
   scanf("%d", &n);
vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);</pre>
    vector<PT> h(v);
    map<PT,int> index;
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    ConvexHull(h);
    double len = 0;
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {
      if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

## 2.2 Miscellaneous geometry

```
#include <iostream>
#include <vector>
#include <castr>
#include <castr>
#include <castr>
#include <castr>
using namespace std;

double INF = le100;
double EPS = le-12;

struct PT {
    double x, y;
    PT() {}
    PT (double x, double y) : x(x), y(y) {}
    PT (const PT &p) : x(p.x), y(p.y) {}
    PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
}
```

```
PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
                                const { return PT(x*c, y*c );
  PT operator * (double c)
  PT operator / (double c)
                                const { return PT(x/c, y/c ); }
double dot (PT p, PT q)
                            { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q)
                            { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
    return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p) { return PT(-p.y,p.x); }
PT RotateCW90(PT p)
                        { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot(c-a, b-a) /dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x, y, z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                          double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
   determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS:
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect (PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
    return false;
    return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true:
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert (dot (b, b) > EPS && dot (d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  c = (a+c)/2:
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
```

```
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon (const vector <PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
    if ((p[i].y \le q.y \&\& q.y < p[j].y ||
     p[j].y \le q.y && q.y < p[i].y) &&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
      c = !c;
 return c:
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
     return true;
    return false:
\ensuremath{//} compute intersection of line through points a and b with
// circle centered at c with radius r >
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret:
  b = b-a:
  a = a-c:
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius {\bf r}
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R \mid \mid d+min(r, R) < max(r, R)) return ret; double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret:
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
   the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
 return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {
    for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
      int 1 = (k+1) % p.size();
      if (i == 1 \mid | j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
```

```
return false;
  return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  << ProjectPointSegment (PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6 78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;
  // expected: 1 0 1
  << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 0 0 1
  cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
       << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
       << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 1 1 1 0
  cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "</pre>
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
<< SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "</pre>
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
  cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
  // expected: (1,1)
  cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
  vector<PT> v:
  v.push_back(PT(0,0));
  v.push_back(PT(5,0));
  v.push_back(PT(5,5));
  v.push back(PT(0.5));
  // expected: 1 1 1 0 0
  << PointInPolygon(v, PT(0,2)) << " "
       << PointInPolygon(v, PT(5,2)) << " "
       << PointInPolygon(v, PT(2,5)) << endl;
  // expected: 0 1 1 1 1
  cerr << PointOnPolygon(v, PT(2,2)) << " "
       << PointOnPolygon(v, PT(2,0)) << " "
       << PointOnPolygon(v, PT(0,2)) << " "
       << PointOnPolygon(v, PT(5,2)) << " "
       << PointOnPolygon(v, PT(2,5)) << endl;
                (5,4) (4,5)
                blank line
                (4,5) (5,4)
                blank line
                (4,5) (5,4)
  vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
  u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);

for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;

u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);

for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
  u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
  u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
  // centroid should be (1.1666666, 1.166666)
  PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
```

```
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;
return 0;</pre>
```

## 3 Numerical algorithms

# 3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std:
typedef vector<int> VI;
typedef pair<int, int> PII;
 // return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a:
// computes lcm(a.b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1;
        return ret;
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
        int yy = x = 1;
        while (b) {
                int q = a / b;
int t = b; b = a%b; a = t;
t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
        return a:
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
        int x, y;
        VI ret;
        int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                 x = mod(x*(b / g), n);
                for (int i = 0; i < g; i++)
                         ret.push_back(mod(x + i*(n / g), n));
        return ret:
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
```

```
int x, y;
        int g = extended_euclid(a, n, x, y);
        if (g > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find \boldsymbol{z} such that
//\ z\ \$\ m1\ =\ r1,\ z\ \$\ m2\ =\ r2\ .\quad Here,\ z\ is\ unique\ modulo\ M\ =\ lcm\,(m1,\ m2)\ .
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
        int s, t;
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g);
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
        PII ret = make_pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {
                 ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
                 if (ret.second == -1) break:
        return ret:
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
        if (!a && !b)
                 if (c) return false;
                 x = 0; y = 0;
                 return true;
        if (!a)
                 if (c % b) return false;
                 x = 0; y = c / b;
                 return true;
        if (!b)
                 if (c % a) return false;
                 x = c / a; y = 0;
                 return true;
        int g = gcd(a, b);
        if (c % g) return false;
        x = c / g * mod_inverse(a / g, b / g);
         v = (c - a * x) / b;
        return true;
int main() {
         // expected: 2
        cout << gcd(14, 30) << endl;
         // expected: 2 -2 1
        int x, y;
int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
         // expected: 95 451
        VI sols = modular_linear_equation_solver(14, 30, 100);
        for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
        cout << endl:
         // expected: 8
        cout << mod_inverse(8, 9) << endl;</pre>
         // expected: 23 105
        PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
        cout << ret.first << " " << ret.second << endl;</pre>
        ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second << endl;</pre>
         // expected: 5 -15
        if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;</pre>
        cout << x << " " << y << endl;
        return 0:
```

# 3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
     (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
     (3) computing determinants of square matrices
// Running time: O(n^3)
// INPUT: a[][] = an nxn matrix
               b[][] = an nxm matrix
                     = an nxm matrix (stored in b[][])
// OUTPUT: X
               A^{-1} = an nxn matrix (stored in a[][])
               returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10:
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan(VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T \det = 1:
  for (int i = 0; i < n; i++) {
    int pj = -1, pk = -1;
for (int j = 0; j < n; j++) if (!ipiv[j])
  for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
         if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
     if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre>
    ipiv[pk]++;
    swap(a[pj], a[pk]);
     swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
irow[i] = pj;
icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
a[pk][pk] = 1.0;
for (int p = 0; p < n; p++) a[pk][p] *= c;
    for (int p = 0; p < m; p++) b[pk][p] *= c;
for (int p = 0; p < n; p++) if (p != pk) {
      c = a[p][pk];
       a[p][pk] = 0;
       for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;</pre>
       for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
   for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
  return det:
int main() {
  const int n = 4:
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};  double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \}; 
  for (int i = 0; i < n; i++) {
     a[i] = VT(A[i], A[i] + n);
    b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
   // expected: 60
  cout << "Determinant: " << det << endl;
```

σ

```
// expected: -0.233333 0.166667 0.133333 0.0666667
               0.166667 0.166667 0.333333 -0.333333
               0.233333 0.833333 -0.133333 -0.0666667
               0.05 -0.75 -0.1 0.2
cout << "Inverse: " << endl;</pre>
for (int i = 0; i < n; i++) {
 for (int j = 0; j < n; j++)
  cout << a[i][j] << ' ';</pre>
 cout << endl;
// expected: 1.63333 1.3
               -0.166667 0.5
               2.36667 1.7
               -1.85 -1.35
cout << "Solution: " << endl;
for (int i = 0; i < n; i++) {
 for (int j = 0; j < m; j++)
  cout << b[i][j] << ' ';</pre>
  cout << endl;
```

#### 3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT: a[][] = an nxm matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
             returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std:
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
  int n = a.size();
  int m = a[0].size();
  int r = 0;
  for (int c = 0; c < m && r < n; c++) {
    int j = r;
    for (int i = r + 1; i < n; i++)
   if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[i], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
for (int i = 0; i < n; i++) if (i != r) {
      T t = a[i][c];
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];</pre>
    r++;
  return r:
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
    {16, 2, 3, 13},
    { 5, 11, 10, 8},
    { 9, 7, 6, 12},
    { 4, 14, 15, 1},
    {13, 21, 21, 13}};
  for (int i = 0; i < n; i++)
    a[i] = VT(A[i], A[i] + m);
  int rank = rref(a);
  // expected: 3
  cout << "Rank: " << rank << endl;
```

```
// expected: 1 0 0 1

// 0 1 0 3

// 0 0 1 -3

// 0 0 0 1 -3

// 0 0 0 3.10862e-15

cout << "rref: " << endl;

for (int i = 0; i < 5; i++) {

  for (int j = 0; j < 4; j++)

      cout << a[i][j] << ' ';

  cout << endl;

}
```

#### 3.4 Fast Fourier transform

```
#include <cassert>
#include <cstdio>
#include <cmath>
struct cpx
  cpx(){}
  cpx(double aa):a(aa),b(0){}
  cpx(double aa, double bb):a(aa),b(bb){}
  double a;
  double b:
  double modsq(void) const
    return a * a + b * b;
  cpx bar(void) const
    return cpx(a, -b);
cpx operator +(cpx a, cpx b)
  return cpx(a.a + b.a, a.b + b.b);
cpx operator *(cpx a, cpx b)
  return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
cpx operator / (cpx a, cpx b)
  cpx r = a * b.bar();
  return cpx(r.a / b.modsq(), r.b / b.modsq());
cpx EXP (double theta)
  return cpx(cos(theta), sin(theta));
const double two_pi = 4 * acos(0);
            input array
// out:
           output array
// step: {SET TO 1} (used internally)
// size: length of the input/output {MUST BE A POWER OF 2}
// dir: either plus or minus one (direction of the FFT) 
// RESULT: out[k] = \sum_{j=0}^{\size - 1} in[j] * \exp(\dir * 2pi * i * j * k / \size)
void FFT(cpx *in, cpx *out, int step, int size, int dir)
  if(size < 1) return;</pre>
  if(size == 1)
    out[0] = in[0];
    return:
  FFT(in, out, step * 2, size / 2, dir);
  FFT(in + step, out + size / 2, step * 2, size / 2, dir);
  for(int i = 0 ; i < size / 2 ; i++)</pre>
    cpx even = out[i];
    cpx odd = out[i + size / 2];
    out[i] = even + EXP(dir * two_pi * i / size) * odd;
    out[i + size / 2] = even + EXP(dir * two_pi * (i + size / 2) / size) * odd;
// f[0...N-1] and g[0...N-1] are numbers // Want to compute the convolution h, defined by
// h[n] = sum \ of \ f[k]g[n-k] \ (k = 0, ..., N-1).
```

```
// Here, the index is cyclic; f[-1] = f[N-1], f[-2] = f[N-2], etc.
// Let F[0...N-1] be FFT(f), and similarly, define G and H.
// The convolution theorem says H[n] = F[n]G[n] (element-wise product).
// To compute h[] in O(N \log N) time, do the following:
    1. Compute F and G (pass dir = 1 as the argument).
// 2. Get H by element-wise multiplying F and G.
    3. Get h by taking the inverse FFT (use dir = -1 as the argument)
       and *dividing by N*. DO NOT FORGET THIS SCALING FACTOR.
int main (void)
  printf("If rows come in identical pairs, then everything works.\n");
  cpx \ a[8] = \{0, 1, cpx(1,3), cpx(0,5), 1, 0, 2, 0\};
  cpx b[8] = {1, cpx(0,-2), cpx(0,1), 3, -1, -3, 1, -2};
  cpx B[8];
  FFT(a, A, 1, 8, 1);
  FFT(b, B, 1, 8, 1);
  for (int i = 0; i < 8; i++)
   printf("%7.21f%7.21f", A[i].a, A[i].b);
  printf("\n");
  for(int i = 0 : i < 8 : i++)
    cpx Ai(0,0);
   for(int j = 0 ; j < 8 ; j++)
      Ai = Ai + a[j] * EXP(j * i * two_pi / 8);
    printf("%7.21f%7.21f", Ai.a, Ai.b);
  printf("\n");
  cpx AB[8];
  for(int i = 0; i < 8; i++)
   AB[i] = A[i] * B[i];
  cpx aconvb[8];
  FFT (AB, aconvb, 1, 8, -1);
  for(int i = 0; i < 8; i++)
   aconvb[i] = aconvb[i] / 8;
  for (int i = 0; i < 8; i++)
    printf("%7.21f%7.21f", aconvb[i].a, aconvb[i].b);
  for (int i = 0; i < 8; i++)
    cpx aconvbi(0,0);
    for (int j = 0; j < 8; j++)
      aconvbi = aconvbi + a[j] * b[(8 + i - j) % 8];
   printf("%7.21f%7.21f", aconvbi.a, aconvbi.b);
 printf("\n");
  return 0:
```

#### 3.5 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear programs of the form
       maximize
       subject\ to\ Ax <= b
                   x >= 0
// INPUT: A -- an m x n matrix
         b -- an m-dimensional vector
         c -- an n-dimensional vector
         x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
          above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
```

```
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
  int m, n;
  VI B. N:
  VVD D;
  LPSolver(const VVD &A, const VD &b, const VD &c) :
    m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
    for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] = A[i][j]; for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1; D[i][n + 1] = b[i]; }
    for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
    N[n] = -1; D[m + 1][n] = 1;
  void Pivot(int r, int s) {
    double inv = 1.0 / D[r][s];
    for (int i = 0; i < m + 2; i++) if (i != r)
for (int j = 0; j < n + 2; j++) if (j != s)
    D[i][j] -= D[r][j] * D[i][s] * inv;

for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;

for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;
    D[r][s] = inv:
    swap(B[r], N[s]);
  bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
    while (true) {
      int s = -1;
       for (int j = 0; j <= n; j++) {
        if (phase == 2 && N[j] == -1) continue;
        if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] && N[j] < N[s]) s = j;
       if (D[x][s] > -EPS) return true;
       int r = -1;
       for (int i = 0; i < m; i++) {
        if (D[i][s] < EPS) continue;</pre>
        if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
           (D[i][n+1] / D[i][s]) == (D[r][n+1] / D[r][s]) && B[i] < B[r]) r = i;
       if (r == -1) return false;
      Pivot(r, s);
  DOUBLE Solve(VD &x) {
    int r = 0:
    for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n + 1] < -EPS) {
      Pivot(r. n):
       if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -numeric_limits<DOUBLE>::infinity();
       for (int i = 0; i < m; i++) if (B[i] == -1) {
         for (int j = 0; j \le n; j++)
           if (s == -1 \mid \mid D[i][j] < D[i][s] \mid \mid D[i][j] == D[i][s] && N[j] < N[s]) s = j;
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
    x = VD(n);
for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return D[m][n + 1];
};
int main() {
  const int m = 4;
  const int n = 3;
  DOUBLE _A[m][n] =
    \{-1, -5, 0\},
    { 1, 5, 1 },
    \{-1, -5, -1\}
  DOUBLE _{b[m]} = \{ 10, -4, 5, -5 \};
  DOUBLE _c[n] = { 1, -1, 0 };
  VVD A(m);
  VD b(\underline{b}, \underline{b} + m);
  VD c(_c, _c + n);
  for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);</pre>
  LPSolver solver (A, b, c);
```

```
VD x;
DOUBLE value = solver.Solve(x);
cerr << "VALUE: " << value << end1; // VALUE: 1.29032
cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
return 0;</pre>
```

## 4 Graph algorithms

# 4.1 Bellman-Ford shortest paths with negative edge weights

```
// This function runs the Bellman-Ford algorithm for single source
// shortest paths with negative edge weights. The function returns
// false if a negative weight cycle is detected. Otherwise, the
// function returns true and dist[i] is the length of the shortest
// path from start to i.
// Running time: O(|V|^3)
     INPUT: start, w[i][j] = cost of edge from i to j
    OUTPUT: dist[i] = min weight path from start to i
              prev[i] = previous node on the best path from the
                         start node
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool BellmanFord (const VVT &w, VT &dist, VI &prev, int start) {
 int n = w.size();
  prev = VI(n, -1);
  dist = VT(n, 1000000000);
  dist[start] = 0;
  for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {</pre>
        if (dist[j] > dist[i] + w[i][j]) +
          if (k == n-1) return false;
          dist[j] = dist[i] + w[i][j];
prev[j] = i;
  return true;
```

## 4.2 Fast Dijkstra's algorithm

```
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
//
// Running time: O(|E| log |V|)

#include <queue>
#include <cstdio>
using namespace std;
const int INF = 20000000000;
typedef pair<int, int> PII;
int main() {
```

```
int N, s, t;
        scanf("%d%d%d", &N, &s, &t);
         vector<vector<PII> > edges(N);
         for (int i = 0; i < N; i++) {
                 scanf("%d", &M);
for (int j = 0; j < M; j++)
                          int vertex, dist;
                          scanf("%d%d", &vertex, &dist);
                          edges[i].push_back(make_pair(dist, vertex)); // note order of arguments here
         // use priority queue in which top element has the "smallest" priority
        priority_queue<PII, vector<PII>, greater<PII> > Q;
vector<int> dist(N, INF), dad(N, -1);
        Q.push(make_pair(0, s));
         dist[s] = 0;
         while (!Q.empty()) {
                 PII p = Q.top();
                 Q.pop();
                 int here = p.second;
                 if (here == t) break;
                 if (dist[here] != p.first) continue;
                 for (vector<PII>::iterator it = edges[here].begin(); it != edges[here].end(); it++) {
                          if (dist[here] + it->first < dist[it->second]) {
                                  dist[it->second] = dist[here] + it->first;
                                  dad[it->second] = here;
                                  Q.push(make_pair(dist[it->second], it->second));
         printf("%d\n", dist[t]);
        if (dist[t] < INF)</pre>
                 for (int i = t; i != -1; i = dad[i])
                          printf("%d%c", i, (i == s ? '\n' : ' '));
        return 0:
Sample input:
5 0 4 2 1 2 3 1
2 2 4 4 5
3 1 4 3 3 4 1
2 0 1 2 3
2 1 5 2 1
Expected:
4 2 3 0
```

#### 4.3 Strongly connected components

```
#include<memory.h>
struct edge{int e, nxt;};
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group_ent, group_num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill_forward(int x)
  for(i=sp[x];i;i=e[i].nxt) if(!v[e[i].e]) fill_forward(e[i].e);
  stk[++stk[0]]=x;
void fill_backward(int x)
  for(i=spr[x];i;i=er[i].nxt) if(v[er[i].e]) fill_backward(er[i].e);
void add_edge(int v1, int v2) //add edge v1->v2
  e [++E].e=v2; e [E].nxt=sp [v1]; sp [v1]=E;
  er[ E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
void SCC()
  int i:
```

```
stk[0]=0;
memset(v, false, sizeof(v));
for(i=1;i<=V;i++) if(!v[i]) fill_forward(i);
group_cnt=0;
for(i=stk[0];i>=1;i--) if(v[stk[i]]){group_cnt++; fill_backward(stk[i]);}
```

## 4.4 Eulerian path

```
typedef list<Edge>::iterator iter;
struct Edge
        int next_vertex;
        iter reverse_edge;
        Edge(int next_vertex)
                :next_vertex(next_vertex)
};
const int max_vertices = ;
int num_vertices;
list<Edge> adj[max_vertices];
                                        // adjacency list
vector<int> path;
void find_path(int v)
        while(adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().reverse_edge);
                adj[v].pop_front();
                find path(vn);
        path.push_back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse edge = itb;
        itb->reverse_edge = ita;
```

#### 4.5 Kruskal's algorithm

```
Uses Kruskal's Algorithm to calculate the weight of the minimum spanning
forest (union of minimum spanning trees of each connected component) of
a possibly disjoint graph, given in the form of a matrix of edge weights
(-1 if no edge exists). Returns the weight of the minimum spanning
forest (also calculates the actual edges - stored in T). Note: uses a
disjoint-set data structure with amortized (effectively) constant time per
union/find. Runs in O(E*log(E)) time.
#include <iostream>
#include <vector>
#include <algorithm>
#include <queue>
using namespace std;
typedef int T;
struct edge
  int u, v;
};
struct edgeCmp
  int operator()(const edge& a, const edge& b) { return a.d > b.d; }
};
int find(vector <int>& C, int x) { return (C[x] == x) ? x : C[x] = find(C, C[x]); }
```

```
T Kruskal (vector <vector <T> >& w)
  int n = w.size();
  T weight = 0;
  vector <int> C(n), R(n);
  vector <edge> T;
  priority_queue <edge, vector <edge>, edgeCmp> E;
  for(int i=0; i<n; i++)
  for(int j=i+1; j<n; j++)</pre>
     if(w[i][j] >= 0)
        e.u = i; e.v = j; e.d = w[i][j];
        E.push(e);
  while (T.size() < n-1 && !E.empty())</pre>
    edge cur = E.top(); E.pop();
    int uc = find(C, cur.u), vc = find(C, cur.v);
    if(uc != vc)
      T.push_back(cur); weight += cur.d;
      if(R[uc] > R[vc]) C[vc] = uc;
      else if(R[vc] > R[uc]) C[uc] = vc;
      else { C[vc] = uc; R[uc]++; }
  return weight;
int main()
  int wa[6][6] = {
   \{0, -1, 2, -1, 7, -1\},\
\{-1, 0, -1, 2, -1, -1\},\
    \{2, -1, 0, -1, 8, 6\},\
    { -1, 2, -1, 0, -1, -1 },
    \{ 7, -1, 8, -1, 0, 4 \},
    \{-1, -1, 6, -1, 4, 0\}\};
  vector <vector <int> > w(6, vector <int>(6));
  for(int i=0; i<6; i++)</pre>
    for(int j=0; j<6; j++)
      w[i][j] = wa[i][j];
  cout << Kruskal(w) << endl;
  cin >> wa[0][0];
```

#### 4.6 Minimum spanning trees

```
// This function runs Prim's algorithm for constructing minimum
// weight spanning trees.
// Running time: O(|V|^2)
    INPUT: w[i][j] = cost of edge from i to j
             NOTE: Make sure that w[i][i] is nonnegative and
             symmetric. Missing edges should be given -1
              weight.
    OUTPUT: edges = list of pair<int,int> in minimum spanning tree
             return total weight of tree
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
```

```
typedef vector<VI> VVI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
T Prim (const VVT &w, VPII &edges) {
  int n = w.size();
  VI found (n);
  VI prev (n, -1);
  VT dist (n, 1000000000);
  int here = 0;
  dist[here] = 0;
  while (here != -1) {
    found[here] = true;
    int best = -1:
    for (int k = 0; k < n; k++) if (!found[k]) {</pre>
      if (w[here][k] != -1 && dist[k] > w[here][k]){
         dist[k] = w[here][k];
         prev[k] = here;
       if (best == -1 || dist[k] < dist[best]) best = k;</pre>
    here = best:
  T tot_weight = 0;
  for (int i = 0; i < n; i++) if (prev[i] != -1) {</pre>
    edges.push_back (make_pair (prev[i], i));
    tot_weight += w[prev[i]][i];
  return tot weight:
int main(){
  int ww[5][5] = {
    {0, 400, 400, 300, 600},
    {400, 0, 3, -1, 7},
    {400, 3, 0, 2, 0},
    {300, -1, 2, 0, 5},
    {600, 7, 0, 5, 0}
  VVT w(5, VT(5));
  for (int i = 0; i < 5; i++)
  for (int j = 0; j < 5; j++)
    w[i][j] = ww[i][j];</pre>
  // expected: 305
                 0 3
  VPII edges:
  cout << Prim (w, edges) << endl;</pre>
  for (int i = 0; i < edges.size(); i++)
  cout << edges[i].first << " " << edges[i].second << endl;</pre>
```

### 5 Data structures

## 5.1 Suffix array

```
// Suffix array construction in O(L log^2 L) time. Routine for // computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
             of substring s[i...L-1] in the list of sorted suffixes.
             That is, if we take the inverse of the permutation suffix[],
             we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
struct SuffixArray {
  const int L;
  string s;
  vector<vector<int> > P;
  vector<pair<int,int>,int> > M;
```

```
SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);</pre>
    for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
      P.push_back(vector<int>(L, 0));
      for (int i = 0; i < L; i++)
       M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
      sort(M.begin(), M.end());
      for (int i = 0; i < L; i++)
        P[level][M[i].second] = (i > 0 \&\& M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i;
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
     if (P[k][i] == P[k][j]) {
        i += 1 << k;
        j += 1 << k;
        len += 1 << k:
    return len:
1:
// BEGIN CUT
  The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
int main() {
 int T;
  cin >> T;
  for (int caseno = 0; caseno < T; caseno++) {</pre>
    string s;
    cin >> s;
    SuffixArray array(s);
    vector<int> v = array.GetSuffixArray();
    int bestlen = -1, bestpos = -1, bestcount = 0;
    for (int i = 0; i < s.length(); i++) {
  int len = 0, count = 0;</pre>
      for (int j = i+1; j < s.length(); j++) {</pre>
        int 1 = array.LongestCommonPrefix(i, j);
        if (1 >= len) {
          if (1 > len) count = 2; else count++;
          len = 1;
      if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen) > s.substr(i, len)) {
        bestlen = len:
        hest count = count:
        bestpos = i:
    if (bestlen == 0) {
      cout << "No repetitions found!" << endl;
     cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;</pre>
#else
// END CUT
int main() {
  // bobocel is the O'th suffix
  // obocel is the 5'th suffix
      bocel is the 1'st suffix
       ocel is the 6'th suffix
         cel is the 2'nd suffix
         el is the 3'rd suffix
           l is the 4'th suffix
  SuffixArray suffix("bobocel");
  vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
  for (int i = 0; i < v.size(); i++) cout << v[i] << " ";</pre>
  cout << endl:
  cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
// BEGIN CUT
#endif
// END CUT
```

#### 5.2 Binary Indexed (Fenwick) Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
  while (x \le N) {
    tree[x] += v;
    x += (x & -x);
// get cumulative sum up to and including x
int get(int x) {
  int res = 0;
  while(x) {
    res += tree[x];
    x = (x & -x);
  return res;
// get largest value with cumulative sum less than or equal to \boldsymbol{x};
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
  int idx = 0, mask = N;
  while (mask && idx < N) {
    int t = idx + mask;
    if(x >= tree[t]) {
      x -= tree[t];
    mask >>= 1;
  return idx:
```

## 5.3 Union-find disjoint sets

```
#include <iostream>
#include <vector>
using namespace std;
struct UnionFind {
    vectorsint> C;
    UnionFind(int n) : C(n) { for (int i = 0; i < n; i++) C[i] = i; }
    int find(int x) { return (C[x] == x) ? x : C[x] = find(C[x]); }
    void merge(int x, int y) { C[find(x)] = find(y); }
};
int main() {
    int n = 5;
    UnionFind uf(n);
    uf.merge(0, 2);
    uf.merge(1, 0);
    uf.merge(3, 4);
    for (int i = 0; i < n; i++) cout << i << " " << uf.find(i) << endl;
    return 0;
}</pre>
```

#### 5.4 KD-tree

```
// - straightforward, but probably sub-optimal KD-tree implmentation 
// that's probably good enough for most things (current it's a 
// 2D-tree) 
// - constructs from n points in O(n lg'2 n) time 
// - handles nearest-neighbor query in O(lg n) if points are well 
// distributed 
// - worst case for nearest-neighbor may be linear in pathological 
// case 
// Sonny Chan, Stanford University, April 2009
```

```
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
    point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
1:
bool operator==(const point &a, const point &b)
    return a.x == b.x && a.y == b.y;
// sorts points on x-coordinate
bool on_x(const point &a, const point &b)
    return a.x < b.x:
// sorts points on y-coordinate
bool on_y (const point &a, const point &b)
    return a.y < b.y;
// squared distance between points
ntype pdist2(const point &a, const point &b)
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
// bounding box for a set of points
struct bbox
    ntype x0, x1, y0, y1;
    bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
    // computes bounding box from a bunch of points
    void compute(const vector<point> &v) {
        for (int i = 0; i < v.size(); ++i) {</pre>
            x0 = min(x0, v[i].x); x1 = max(x1, v[i].x);
            y0 = min(y0, v[i].y); y1 = max(y1, v[i].y);
    // squared distance between a point and this bbox, 0 if inside
    ntype distance(const point &p) {
        if (p.x < x0) {
            if (p.y < y0)
                                return pdist2(point(x0, y0), p);
            else if (p.y > y1) return pdist2(point(x0, y1), p);
                                return pdist2(point(x0, p.y), p);
        else if (p.x > x1) {
           if (p.y < y0)
                                return pdist2(point(x1, y0), p);
            else if (p.y > y1) return pdist2(point(x1, y1), p);
            else
                                return pdist2(point(x1, p.y), p);
        else (
                                return pdist2(point(p.x, y0), p);
           if (p.v < v0)
            else if (p.y > y1) return pdist2(point(p.x, y1), p);
            else
                                return 0;
};
// stores a single node of the kd-tree, either internal or leaf
struct kdnode
    bool leaf;
                    // true if this is a leaf node (has one point)
    point pt;
                    // the single point of this is a leaf
    bbox bound:
                   // bounding box for set of points in children
    kdnode *first, *second; // two children of this kd-node
    kdnode() : leaf(false), first(0), second(0) {}
    "kdnode() { if (first) delete first; if (second) delete second; }
    // intersect a point with this node (returns squared distance)
    ntype intersect(const point &p) {
```

```
return bound.distance(p);
     // recursively builds a kd-tree from a given cloud of points
    void construct(vector<point> &vp)
         // compute bounding box for points at this node
        bound.compute(vp);
        // if we're down to one point, then we're a leaf node
if (vp.size() == 1) {
            leaf = true;
            pt = vp[0];
        else (
               split on x if the bbox is wider than high (not best heuristic...)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on_x);
             // otherwise split on y-coordinate
            else
                sort(vp.begin(), vp.end(), on_y);
            // divide by taking half the array for each child
             // (not best performance if many duplicates in the middle)
            int half = vp.size()/2;
            vector<point> v1(vp.begin(), vp.begin()+half);
            vector<point> vr(vp.begin()+half, vp.end());
first = new kdnode(); first->construct(vl);
            second = new kdnode(); second->construct(vr);
};
// simple kd-tree class to hold the tree and handle queries
struct kdtree
    kdnode *root;
    // constructs a kd-tree from a points (copied here, as it sorts them)
    kdtree(const vector<point> &vp) {
        vector<point> v(vp.begin(), vp.end());
        root = new kdnode():
        root -> construct (v):
     "kdtree() { delete root; }
    // recursive search method returns squared distance to nearest point
    ntype search(kdnode *node, const point &p)
        if (node->leaf) {
            // commented special case tells a point not to find itself
              if (p == node->pt) return sentry;
              else
                return pdist2(p, node->pt);
        ntype bfirst = node->first->intersect(p);
        ntype bsecond = node->second->intersect(p);
        // choose the side with the closest bounding box to search first
           (note that the other side is also searched if needed)
        if (bfirst < bsecond) {</pre>
             ntype best = search(node->first, p);
            if (bsecond < best)</pre>
                best = min(best, search(node->second, p));
            return best;
            ntype best = search(node->second, p);
            if (bfirst < best)</pre>
                best = min(best, search(node->first, p));
            return best;
    // squared distance to the nearest
    ntype nearest (const point &p) {
        return search (root, p);
};
// some basic test code here
int main()
    // generate some random points for a kd-tree
    vector<point> vp;
    for (int i = 0; i < 100000; ++i) {
        vp.push_back(point(rand()%100000, rand()%100000));
    kdtree tree(vp);
```

#### 5.5 Splay tree

```
#include <cstdio>
#include <algorithm>
using namespace std;
const int N_MAX = 130010;
const int oo = 0x3f3f3f3f;
struct Node
  Node *ch[2], *pre;
  int val, size;
 bool isTurned:
} nodePool[N_MAX], *null, *root;
Node *allocNode(int val)
  static int freePos = 0;
  Node *x = &nodePool[freePos ++];
  x->val = val, x->isTurned = false;
  x->ch[0] = x->ch[1] = x->pre = null;
  x->size = 1:
  return x:
inline void update (Node *x)
  x->size = x->ch[0]->size + x->ch[1]->size + 1;
inline void makeTurned(Node *x)
  if(x == null)
    return:
  swap(x->ch[0], x->ch[1]);
x->isTurned ^= 1;
inline void pushDown (Node *x)
  if(x->isTurned)
    makeTurned(x->ch[0]);
    makeTurned(x->ch[1]);
    x->isTurned ^= 1;
inline void rotate(Node *x, int c)
  Node *y = x -> pre;
  x->pre = y->pre;
if(y->pre != null)
    y->pre->ch[y == y->pre->ch[1]] = x;
  y->ch[!c] = x->ch[c];
if(x->ch[c] != null)
   x->ch[c]->pre = y;
  x->ch[c] = y, y->pre = x;
  update(y);
  if(y == root)
    root = x;
void splay (Node *x, Node *p)
  while(x->pre != p)
    if(x->pre->pre == p)
      rotate(x, x == x->pre->ch[0]);
    else
      Node *v = x - > pre, *z = v - > pre;
      if(y == z \rightarrow ch[0])
```

```
if(x == y->ch[0])
          rotate(y, 1), rotate(x, 1);
          rotate(x, 0), rotate(x, 1);
      else
        if(x == y->ch[1])
          rotate(y, 0), rotate(x, 0);
        else
          rotate(x, 1), rotate(x, 0);
 update(x);
void select(int k, Node *fa)
  Node *now = root;
  while (1)
    pushDown(now);
   int tmp = now->ch[0]->size + 1;
    if(tmp == k)
     hreak:
   else if(tmp < k)</pre>
     now = now -> ch[1], k -= tmp;
    else
     now = now -> ch[0];
  splay(now, fa);
Node *makeTree(Node *p, int 1, int r)
  if(1 > r)
   return null;
  int \ mid = (1 + r) / 2;
  Node *x = allocNode(mid);
 x->pre = p;
x->ch[0] = makeTree(x, 1, mid - 1);
x->ch[1] = makeTree(x, mid + 1, r);
  update(x);
  return x;
int main()
  null = allocNode(0);
  null->size = 0;
  root = allocNode(0):
  root->ch[1] = allocNode(oo);
  root->ch[1]->pre = root;
  update(root);
  scanf("%d%d", &n, &m);
  root->ch[1]->ch[0] = makeTree(root->ch[1], 1, n);
  splay(root->ch[1]->ch[0], null);
  while (m --)
    int a, b;
    scanf("%d%d", &a, &b);
    a ++, b ++;
    select(a - 1, null);
    select(b + 1, root);
    makeTurned(root->ch[1]->ch[0]);
  for(int i = 1; i <= n; i ++)
    select(i + 1, null);
    printf("%d ", root->val);
```

#### 5.6 Segment tree

```
#include <bits/stdc++.h>
#define LEFT_CHILD(x) ((x << 1) + 1)
#define RIGHT_CHILD(x) ((x << 1) + 2)
#define INF 99999999
#define MAX_N 100005
using namespace std;
int queryParams[35], numbers[MAX_N], queryParamsCount;
string inputQuery;
struct node t {
    int left, right, minValue;
} nodes[4 * MAX_N];
void build(int left, int right, int x = 0) {
    nodes[x].left = left;
    nodes[x].right = right;
    if (left == right) {
        nodes[x].minValue = numbers[left];
        return:
    int mid = (left + right) / 2;
    build(left, mid, LEFT_CHILD(x));
    build(mid + 1, right, RIGHT_CHILD(x));
    nodes[x].minValue = min(nodes[LEFT_CHILD(x)].minValue, nodes[RIGHT_CHILD(x)].minValue);
void parseInputQuery() {
    queryParamsCount = 0;
    queryParams[queryParamsCount] = 0;
    // Parse input query starting from after opening parenthesis i.e. index 6.
    for (int i = 6; i < inputQuery.length(); i++) {
   if ('0' <= inputQuery[i] && inputQuery[i] <= '9') {</pre>
            queryParams[queryParamsCount] = queryParams[queryParamsCount] * 10 + (int) (inputQuery[i] -
        else {
            queryParams[++queryParamsCount] = 0; // next query parameter!
int query(int left, int right, int x = 0) {
    if (left <= nodes[x].left && nodes[x].right <= right) {</pre>
        return nodes[x].minValue;
    int mid = (nodes[x].left + nodes[x].right) / 2;
    int ans = INF:
    if (left <= mid) {</pre>
        ans = min(ans, query(left, right, LEFT_CHILD(x)));
    if (mid < right) {
        ans = min(ans, query(left, right, RIGHT_CHILD(x)));
    return ans:
void setValue(int position, int newValue, int x = 0) {
    if (nodes[x].left == position && nodes[x].right == position) {
       nodes[x].minValue = newValue;
       return;
    int mid = (nodes[x].left + nodes[x].right) / 2;
    if (position <= mid) {</pre>
        setValue(position, newValue, LEFT_CHILD(x));
    if (mid < position) {
        setValue(position, newValue, RIGHT_CHILD(x));
    nodes[x].minValue = min(nodes[LEFT CHILD(x)].minValue, nodes[RIGHT CHILD(x)].minValue);
int main()
    int n, q, placeholder;
    for (int i = 1; i <= n; i++) {
        cin >> numbers[i];
    build(1, n);
    while (q--) {
        cin >> inputOuery:
        if (inputQuery[0] == 'q') { // perform operation query.
            parseInputQuery();
            cout << query(queryParams[0], queryParams[1]) << '\n';</pre>
        else if (inputQuery[0] == 's') { // perform operation shift.
            parseInputQuery();
            placeholder = numbers[queryParams[0]];
```

#### 5.7 Lowest common ancestor

**if**(A[i][j-1]!= -1)

```
const int max nodes, log max nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
                                              // children[i] contains the children of node i
int A[max_nodes][log_max_nodes+1];
                                              // A[i][j] is the 2^j-th ancestor of node i, or -1 if that
       ancestor does not exist
int L[max_nodes];
                                              // L[i] is the distance between node i and the root
// floor of the binary logarithm of \boldsymbol{n}
int lb(unsigned int n)
    if(n==0)
        return -1;
   int p = 0;

if (n >= 1<<16 { n >>= 16; p += 16; }

if (n >= 1<< 8 { n >>= 8; p += 8; }

if (n >= 1<< 4) { n >>= 4; p += 4; }
    if (n >= 1<< 2) { n >>= 2; p += 2; }
    if (n >= 1<< 1) {
    return p;
void DFS(int i, int 1)
    L[i] = 1;
    for(int j = 0; j < children[i].size(); j++)</pre>
         DFS(children[i][j], 1+1);
int LCA (int p, int q)
     // ensure node p is at least as deep as node q
    if(L[p] < L[q])
         swap(p, q);
     // "binary search" for the ancestor of node p situated on the same level as q
    for (int i = log_num_nodes; i >= 0; i--)
if(L[p] - (1<<i) >= L[q])
             p = A[p][i];
    if(p == q)
         return p;
     // "binary search" for the LCA
    for(int i = log_num_nodes; i >= 0; i--)
         if (A[p][i] != -1 && A[p][i] != A[q][i])
             p = A[p][i];
             q = A[q][i];
    return A[p][0];
int main(int argc.char* argv[])
     // read num nodes, the total number of nodes
    log_num_nodes=1b(num_nodes);
    for(int i = 0; i < num_nodes; i++)</pre>
         // read p, the parent of node i or -1 if node i is the root
         A[i][0] = p;
         if(p != -1)
              children[p].push_back(i);
         else
             root = i;
     // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)
    for(int i = 0; i < num_nodes; i++)</pre>
```

### 6 Miscellaneous

## 6.1 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
// Running time: O(n log n)
    INPUT: a vector of integers
    OUTPUT: a vector containing the longest increasing subsequence
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std:
typedef vector<int> VI;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASING
VI LongestIncreasingSubsequence(VI v) {
  VPII best;
 VI dad(v.size(), -1);
  for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASNG
    PII item = make_pair(v[i], 0);
    VPII::iterator it = lower_bound(best.begin(), best.end(), item);
    item.second = i;
#else
    PII item = make_pair(v[i], i);
    VPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
    if (it == best.end()) {
      dad[i] = (best.size() == 0 ? -1 : best.back().second);
      best.push_back(item);
    } else {
     dad[i] = it == best.begin() ? -1 : prev(it)->second;
      *it = item;
  for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push_back(v[i]);
  reverse(ret.begin(), ret.end());
  return ret;
```

#### 6.2 Knuth-Morris-Pratt

```
/*
Finds all occurrences of the pattern string p within the
text string t. Running time is O(n + m), where n and m
are the lengths of p and t, respectively.

*/
#include <iostream>
#include <string>
#include <vector>
using namespace std;

typedef vector<int> VI;
```

```
void buildPi(string& p, VI& pi)
  pi = VI(p.length());
  int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != p[i])
      k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP(string& t, string& p)
  VI pi;
  buildPi(p, pi);
int k = -1;
  for(int i = 0; i < t.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != t[i])
      k = (k == -1) ? -2 : pi[k];
    if(k == p.length() - 1) {
      // p matches t[i-m+1, ..., i]
cout << "matched at index " << i-k << ": ";</pre>
      cout << t.substr(i-k, p.length()) << endl;</pre>
      k = (k == -1) ? -2 : pi[k];
  return 0:
int main()
  string a = "AABAACAADAABAABA", b = "AABA";
  KMP(a, b); // expected matches at: 0, 9, 12
  return 0;
```

#### 6.3 Dates

```
// Routines for performing computations on dates. In these routines,
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
// integers.
#include <iostream>
#include <string>
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y) {
  return
    1461 * (v + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
// converts integer (Julian day number) to Gregorian date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
  int x, n, i, j;
  x = jd + 68569;
  n = 4 * x / 146097;

x -= (146097 * n + 3) / 4;

i = (4000 * (x + 1)) / 1461001;

x -= 1461 * i / 4 - 31;
  j = 80 * x / 2447;
  d = x - 2447 * j / 80;
  x = j / 11;
  m = j + 2 - 12 * x;
  y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
  return dayOfWeek[jd % 7];
int main (int argc, char **argv) {
  int jd = dateToInt (3, 24, 2004);
  int m, d, y;
  intToDate (jd, m, d, y);
string day = intToDay (jd);
  // expected output:
```

```
// 2453089

// 3/24/2004

// Wed

cout << jd << endl

<< m << "," << d << "," << y << endl

<< day << endl;
```

#### 6.4 Prime numbers

```
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x)
  if(x<=1) return false:
 if(x<=3) return true:
  if (!(x%2) || !(x%3)) return false;
  LL s=(LL) (sqrt((double)(x))+EPS);
  for(LL i=5;i<=s;i+=6)
   if (!(x%i) || !(x%(i+2))) return false;
  return true;
// Primes less than 1000:
                                  13
61
                                                         29
79
                             11
59
                                        17
67
                                             19
71
            4.3
      41
                                                              8.3
                            109
                                       127
                                                        139
                                                             149
     157
                167
                            179
                                 181
                                       191
                                             193
                                                  197
                                                        199
           163
     227
           229
                      239
                            241
                                 251
317
                                       257
                                             263
                                                  269
                                                                   281
     283
                                                  347
                                                        349
           293
                 307
                      311
                            313
                                                             353
                                                                   359
     367
           373
                 379
                      383
                            389
                                 397
                                       401
                                             409
                                                  419
                                                        421
                                                                   433
     439
                      457
                                 463
                                       467
                                             479
                                                  487
                                                        491
           443
                449
                            461
     509
                            547
                                             569
                523
                      541
                                       563
                      613
                            617
                                 619
                                       631
                                             641
                                                   643
                                                        647
                      683
                                 701
                                             719
     751
           757
                 761
                      769
                            773
                                 787
                                       797
                                             809
                                                  811
                                                        821
                                 863
953
                                                  883
977
                                                        887
983
     829
           839
                      857
                            859
                                             881
                      941
                            947
// Other primes:
     The largest prime smaller than 10 is 7.
     The largest prime smaller than 100 is 97.
     The largest prime smaller than 1000 is 997.
     The largest prime smaller than 10000 is 9973.
     The largest prime smaller than 100000 is 99991.
     The largest prime smaller than 1000000 is 999983.
     The largest prime smaller than 10000000 is 9999991.
     The largest prime smaller than 100000000 is 99999989.
     The largest prime smaller than 1000000000 is 999999937.
     The largest prime smaller than 10000000000 is 9999999967.
     The largest prime smaller than 10000000000 is 99999999977.
     The largest prime smaller than 100000000000 is 999999999999.
     The largest prime smaller than 100000000000000 is 99999999999937.
     The largest prime smaller than 100000000000000 is 999999999999997
```

# ${\bf 6.5}\quad Latitude/longitude$

```
/*
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
*/
#include <iostream>
#include <cmath>
using namespace std;
struct 11
{
    double r, lat, lon;
};
struct rect
{
    double x, y, z;
};
```

```
11 convert (rect& P)
  11 Q;
  Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
  Q.lat = 180/M_PI*asin(P.z/Q.r);
  Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
rect convert(11& Q)
 P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
  P.z = Q.r*sin(Q.lat*M_PI/180);
  return P:
int main()
  rect A:
 11 B:
  A.x = -1.0; A.y = 2.0; A.z = -3.0;
  B = convert(A);
 cout << B.r << " " << B.lat << " " << B.lon << endl;
 A = convert(B);
cout << A.x << " " << A.y << " " << A.z << endl;
```

#### 6.6 C++ input/output

```
#include <iostream>
#include <iomanip>
using namespace std;
int main()
     // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);</pre>
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl:
    cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
    cout.setf(ios::showpos);
    cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal cout << hex << 100 << " " << 1000 << " " << 10000 << dec << endl;
```

#### 6.7 Random STL stuff

```
#include <algorithm>
#include <iostream>
#include <sstream>
#include <vector>

using namespace std;

int main(void) {
   vector<int> v.
   v. push_back(1);
   v. push_back(2);
   v. push_back(3);
   v. push_back(4);

// Expected output: 1 2 3 4
```

```
1 2 4 3
                     4 3 2 1
 ostringstream oss;
oss << v[0] << " " << v[1] << " " << v[2] << " " << v[3];
  // for input from a string s,
  // istringstream iss(s);
  // iss >> variable;
  cout << oss.str() << endl;</pre>
} while (next_permutation (v.begin(), v.end()));
v.clear():
v.push_back(2);
v.push_back(1);
v.push_back(3);
// To use unique, first sort numbers. Then call
// unique to place all the unique elements at the beginning
// of the vector, and then use erase to remove the duplicate
// elements
sort(v.begin(), v.end());
v.erase(unique(v.begin(), v.end()), v.end());
// Expected output: 1 2 3
for (size_t i = 0; i < v.size(); i++)</pre>
 cout << v[i] << " ";
cout << endl;
```

### 6.8 Regular expressions (Java)

```
// Code which demonstrates the use of Java's regular expression libraries.
// This is a solution for
     Loglan: a logical language
     http://acm.uva.es/p/v1/134.html
// In this problem, we are given a regular language, whose rules can be
// inferred directly from the code. For each sentence in the input, we must
// determine whether the sentence matches the regular expression or not. The
// code consists of (1) building the regular expression (which is fairly
// complex) and (2) using the regex to match sentences.
import java.util.*;
import java.util.regex.*;
public class LogLan {
    public static String BuildRegex () {
   String space = " +";
         String A = "([aeiou])";
         String C = "([a-z&&[^aeiou]])";
String MOD = "(g" + A + ")";
         String BA = "(b" + A + ")";
         String DA = "(d" + A + ")";
String LA = "(1" + A + ")";
         String NAM = "([a-z]*" + C + ")";
         String PREDA = "(" + C + C + A + C + A + "|" + C + A + C + C + A + ")";
         String predstring = "(" + PREDA + "(" + space + PREDA + ")*)";
String predname = "(" + LA + space + predstring + "|" + NAM + ")";
String preds = "(" + predstring + "(" + space + A + space + predstring + ")*)";
         String predclaim = "(" + predname + space + BA + space + preds + "|" + DA + space +
              preds + ")";
         String verbpred = "(" + MOD + space + predstring + ")";
         String statement = "(" + predname + space + verbpred + space + predname + "|" + predname + space + verbpred + ")";
         String sentence = "(" + statement + "|" + predclaim + ")";
         return "^" + sentence + "$";
    public static void main (String args[]) {
         String regex = BuildRegex():
         Pattern pattern = Pattern.compile (regex);
         Scanner s = new Scanner(System.in);
         while (true) {
```

```
// In this problem, each sentence consists of multiple lines, where the last // line is terminated by a period. The code below reads lines until // encountering a line whose final character is a '.'. Note the use of // s.length() to get length of string s.charAt() to extract characters from a Java string s.charAt() to extract characters from a Java string // s.trim() to remove whitespace from the beginning and end of Java string // Other useful String manipulation methods include // s.compareTo(t) < 0 if s < t, lexicographically s.indexOf("apple") returns index of first occurrence of "apple" in s s.lastIndexOf("apple") returns index of last occurrence of "apple" in s s.replace(c,d) replaces occurrences of character c with d s.stattsWith("apple") returns (s.indexOf("apple") = 0) s.toLowerCase() / s.toUpperCase() returns a new lower/uppercased string // Integer.parseInt(s) converts s to an integer (32-bit) Long.parseLong(s) converts s to a long (64-bit) Double.parseDouble(s) converts s to a double
```

```
String sentence = "";
while (true) {
    sentence = (sentence + " " + s.nextLine()).trim();
    if (sentence.equals("#")) return;
    if (sentence.charAt(sentence.length()-1) == '.') break;
}

// now, we remove the period, and match the regular expression

String removed_period = sentence.substring(0, sentence.length()-1).trim();
    if (pattern.matcher (removed_period).find()) {
        System.out.println ("Good");
    } else {
        System.out.println ("Bad!");
    }
}
```