# Homework 2 Report

Yige Hu and Zhiting Zhu

# 1 P1

```
1: function REC_PSUM(a, x_0, b, n)
       if (n == 1) then
2:
 3:
          s(0) = x_0; return; end;
 4:
       end if
       x = zeros(n/2, 1);
 5:
       a\_new = zeros(n/2 - 1, 1);
 6:
 7:
       x(0) = x_0;
 8:
       parfor i = 1 : n do
          x(i) = b(i);
9:
10:
       end parfor
       parfor i = 0 : n/2 - 1 do
11:
          y(i) = x(2*i)*a(2*i+1) + x(2*i+1);
12:
          if (i!=0) then
13:
              a\_new(i) = a(2*i)*a(2*i+1);
14:
          end if
15:
       end parfor
16:
       c = \text{REC\_PSUM}(a\_new, y(0), y[1:n/2-1], n/2);
17:
       s(0) = x_0;
18:
       parfor i = 1 : n - 1 do
19:
          if isOdd(i) then
20:
21:
              s(i) = c(i/2);
22:
              s(i) = c((i-1)/2) * a(i) + x(i);
23:
24:
          end if
       end parfor
25:
       return s;
26:
27: end function
```

# 2 P2

## 2.1 Algorithm

```
1: function SCAN(x, n, l)

2: step = ceil(log_2(n))

3: temp = n >> 1
```

```
offset = 1
 4:
       parfor i = 0 : n/2 - 1 do
 5:
          for j = i; j < temp; j+ = nthreads do
 6:
              indx2 = offset * (2 * i + 2) - 1
 7:
              indx1 = offset * (2 * i + 1) - 1
 8:
              x(indx2) = x(indx1) + x(indx2)
9:
          end for
10:
11:
          offset* = 2
          temp = temp >> 1
12:
13:
       end parfor
       temp = 2
14:
       offset >>= 1
15:
16:
       parfor i = 1 : n/2 - 1 do
          offset >>= 1
17:
          for j = i; j < temp; j+ = nthreads do
18:
              indx2 = offset * (2 * i + 1) - 1
19:
              indx1 = offset * 2 * i - 1
20:
21:
              x(indx2) = x(indx1) + x(indx2)
22:
          end for
          temp* = 2
23:
       end parfor
24:
25: end function
```

#### 2.2 Result

| Wall Clock Time(us) | Number of threads |           |            |
|---------------------|-------------------|-----------|------------|
| Length of Arrary    | sequential        | 6 threads | 12 threads |
| 1M                  | 15679             | 15500     | 38192      |
| 10M                 | 156797.9          | 212012.5  | 160871     |
| 100M                | 730794.8          | 1513714   | 1262623.5  |
| 1B                  | 7305516.5         | 14843186  | 12431315.5 |

Table 1: Wall clock execution time for different array size with different number of threads for 1D vectors

| Wall Clock Time(us) | Number of threads |           |            |
|---------------------|-------------------|-----------|------------|
| Length of Arrary    | sequential        | 6 threads | 12 threads |
| 1M                  | 20525.5           | 79923.5   | 146187     |
| 10M                 | 247284.5          | 539063    | 375131.5   |
| 100M                | 2046770           | 4615023.5 | 3381959    |

Table 2: Wall clock execution time for different array size with different number of threads for 4D vectors

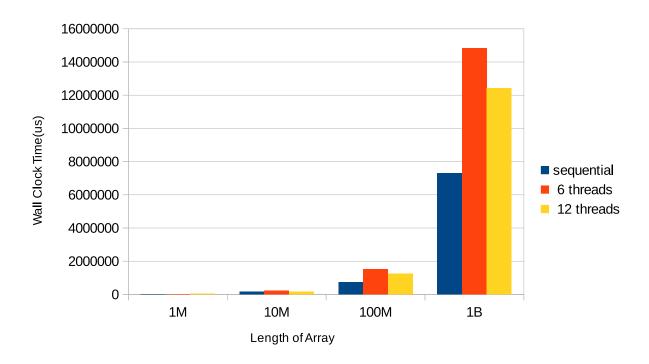


Figure 1: Wall clock execution time for different array size with different number of threads for 1D vectors

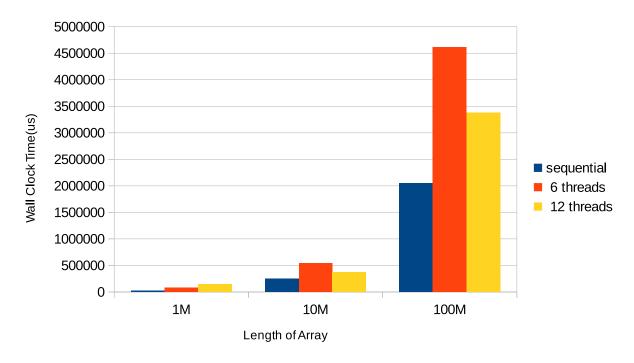


Figure 2: Timing measurements for different array size with different number of threads for 4D vectors

From Figure 1 and 2 and table 1 and 2, the parallel execution runs slower than the sequential execution. The speed up of doubling the number of threads is less than 2. The results are out of our expectation and we have consider two reasons that may be related to this result. First, the cost of administering threads

by OpenMP is high. Second, the input array is shared between multiple threads. During execution, several threads may try to write to this array at different locations which will incur frequent cache coherence for different processors as each thread is pinned to a different processor.

## 3 P3

## 3.1 Algorithm

```
1: function SEARCH(comm, keys, num_keys, sub_arr, arr_size, num_threads, rank)
       low\_pos = arr\_size * rank
 3:
       high\_pos = low\_pos + arr\_size - 1
       low = 0
 4:
       high = arr\_size - 1
 5:
       for i = 0; i < num\_keys; i + + do
 6:
           low = 0
 7:
 8:
           high = arr\_size - 1
           position = low \\
 9:
           k = keys[i]
10:
           if k < sub\_arr[high] \land k > sub\_arr[low] then
11:
              if arr\_size - 1 \le num\_threads then
12:
                  parfor i = 1 : num\_threads do
13:
                     if sub\_arr[low + i] \le k then
14:
                         position = low
15:
                     end if
16:
                  end parfor
17:
              else
18:
                  len = (arr\_size - 1)/(num\_threads + 1)
19:
                  position = low
20:
                  while len! = 1 do
21:
                     parfor i = i : num\_threads do
22:
                         left = low + i * len
23:
                         right = (left + len - 1) > high?high: (left + len - 1)
24:
                         if sub\_arr[left] == k then
25:
                            position = left
26:
                            len = 1
27:
                         else if sub\_arr[right] == k then
28:
                            position = right
29:
                            len = 1
30:
                         else
31:
32:
                             if sub\_arr[left] < k \land sub\_arr[right] > k then
                                len = (len - 1)/(num\_threads + 1)
33:
34:
                                low = left
                                high = right
35:
                                position = low
36:
                             end if
37:
                         end if
38:
                     end parfor
39:
40:
                  end while
              end if
41:
              position = position + low\_pos
42:
              Output position
43:
           else if k == sub\_arr[high] then
44:
45:
              position = high + low\_pos
              Output position
46:
```

```
\begin{array}{lll} 47: & \textbf{else} \\ 48: & \textbf{if } k == sub\_arr[low] \textbf{ then} \\ 49: & position = low + low\_pos \\ 50: & \text{Output position} \\ 51: & \textbf{end if} \\ 52: & \textbf{end if} \\ 53: & \textbf{end for} \\ 54: & \textbf{end function} \end{array}
```

## 3.2 Result

| Length 1 core | 1 socket | 1 node | 2 node |
|---------------|----------|--------|--------|
|---------------|----------|--------|--------|

Table 3: Wall clock execution time for different configurations to search array with different length

We have sussessfully run the program with 16 keys but for large keyspace ( $2^{20}$  keys) our program hangs. We could not figure out the reason for hanging but we suspect that we do not have enough memory to support the full execution.