

## 1 Generally Useful Maths

### Trig Properties

$$\begin{aligned}\sin^2 x + \cos^2 x &= 1 & \frac{d}{dx} \sin x &= \cos x \\ \tan x &= \frac{\sin x}{\cos x} & \frac{d}{dx} \cos x &= -\sin x \\ \frac{d}{dx} \tan x &= \sec^2 x & \frac{d}{dx} \cot x &= -\csc^2 x \\ \frac{d}{dx} \arcsin x &= \frac{1}{\sqrt{1-x^2}} & \frac{d}{dx} \arccos x &= \frac{-1}{\sqrt{1-x^2}} \\ \frac{d}{dx} \arctan x &= \frac{1}{1+x^2}\end{aligned}$$

### Log & Exp Properties

$$\begin{aligned}\frac{x}{dx} b^x &= b^x \ln x & \log\left(\frac{1}{x}\right) &= -\log x \\ \log_a x &= \frac{\log_b x}{\log_a b} & \frac{d}{dx} e^{ax} &= a e^{ax} \\ \frac{d}{dx} a^x &= a^x \ln a & \frac{d}{dx} \ln x &= \frac{1}{x} \\ x^0 &= 1 & x^n \cdot x^m &= x^{n+m} & x^{-n} &= \frac{1}{x^n} \\ \log_a x^n &= n \log_a x & \frac{e^{-nx}}{e^x} &= e^{-(n+1)x} \\ \log_a\left(\frac{x}{y}\right) &= \log_a x - \log_a y \\ \log_a(xy) &= \log_a x + \log_a y\end{aligned}$$

### Useful Series

$$\begin{aligned}r^0 + r^1 + r^2 + r^3 &= \frac{r^n - 1}{r - 1} \\ \text{for an alternating series the following} & \\ \text{will work to start:} & \\ \sum_{n=0}^{\infty} (-1)^n \text{ or } \sum_{n=0}^{\infty} (-1)^{n+1} &\end{aligned}$$

### In Class Terminology

$$\begin{aligned}\text{the relative error formula: } & \frac{|x - \hat{x}|}{x} \\ \text{these were repesd strangely in class:} & \\ x' = f(t, x) & \quad x(2) = 1 \rightarrow t = 2, x = 1 \\ \text{If } x'' = xx' & \text{ then } x''' = xx'' + x'x'\end{aligned}$$

## 2 Base Conversion

### Decimal to Binary

For this simply find the place of the largest binary number that (of the form  $2^n$ ) that is within the number. Successivley subtract these numbers while keeping track of their place to generate the binary number.

### Binary to Decimal

For this notice that each place in the decimal number has a corresponding power of 2. If the decimal number has a floating point then the power is negative counting from zero. This generates a sum of the form:  
 $2^n + \dots + 2^2 + 2^1 + 2^{-1} + 2^{-2} + \dots + 2^{-m}$   
 Where  $n$  is the most significant digit and  $m$  is the least. The  $2^{-1}$  term is the begining of the floating point numbers.

### Binary to Octal

Simply follow the table:  $000 \rightarrow 0$   
 $001 \rightarrow 1$   $002 \rightarrow 2$   $003 \rightarrow 3$   
 $004 \rightarrow 4$   $005 \rightarrow 5$   $006 \rightarrow 6$   $007 \rightarrow 7$

## Binary to Hex

This identical to the Octal method, the Hex symbols range from 0 to  $F$  and binary from 0000 to 1111. Simply count up un binary and there is a simple conversion.

### One & Two's Complement

### 3 IEEE Floating Points

#### Definitions

$s$  = signed bit,  $c$  = based exponent,  $F$  = fraction. The general form for this is  $(-1)^s \cdot 2^{c-127} \cdot 1.F$ , for both  $|s| = 1$   
 For single precision:  $|c| = 8$ ,  $|F| = 23$   
 For double precision:  $|c| = 11$ ,  $|F| = 52$

### Converting to IEEE Format

A number will have the form  $D_n \dots D_1 D_0.F_0 F_1 \dots F_m$ , to start we need to shift the values left (normalize) so that the number is now of the form:  $D_n.F_0 F_1 \dots F_{m+(n-1)} \cdot 10^{n-1}$ .

### Example

Converting the number  $-42.125$  to binary floating point with single precision:  
 $0.125 \cdot 2 = [0].25 \rightarrow 0.25 \cdot 2 = [0].5 \rightarrow 0.5 \cdot 2 = [1].0$  Thus the fractional part is: 0.001, the nonfractional part is:  
 $\frac{1}{32} \frac{0}{16} \frac{1}{8} \frac{0}{4} \frac{1}{2} \frac{0}{1} \rightarrow 101010$   
 The full value, 101010.001, when normalized is:  $1.01010001 \cdot 10^5$ . We find the  $c$  term with  $2^{c-127} = 2^5 \rightarrow c = 132$  and we already have  $F = .01010001$  This gives us the following number with single precision:  
 $[1][10000100][01010001\dots b_{23}]$

## 4 Loss of Significance

### Loss of Precision Theorem

The general form of the theorem is as follows:

$x$  and  $y$  are floating point numbers such that  $x > y > 0$ , the theorem states that given:  $2^{-p} \leq 1 - \frac{y}{x} \leq 2^{-q}$   
 there are at most  $p$  and at least  $q$  digits lost in the subtraction  $x - y$ .

practically speaking, we view the equation as:  $E(x) = f(x) - g(x)$ . If we notice this approaches 0 we have a concern of loss of precision at that point. to find that point we typically view the max loss acceptable as 1, so we set the euqation to  $\frac{g(x)}{f(x)} = \frac{1}{2}$ . We find the  $x = z$  values that cause the  $\frac{1}{2}$  flip and use a Taylor method there and use the normal formula elsewhere. We're just avoiding the loss of precision as  $x \rightarrow z$ .

### Rationalizing Numerators

In some cases we want to rationalize a numerator to avoid a loss of significance. The general form form for radicals in a

denominator is:

$$\sqrt[k]{x^n} + r + c \cdot \frac{\sqrt[k]{x^{n+r-c}}}{\sqrt[k]{x^{n+r-c}}} = \frac{x^n + r - 2c}{\sqrt[k]{x^{n+r-c}}}$$

## 5 Taylor, Maclaurin, & Euler

### Taylor Series

The Taylor series is a sum of derivatives of increasing order that equate to a function. The formula for the Taylor series of  $f(x)$  evaluated at  $a$  is:

$$\begin{aligned}f(x) &= \sum_{n=0}^{\infty} \frac{f^{(n)}(a)}{n!} (x - a)^n = f(a) + \\ f'(a)(x - a) &+ \frac{f''(a)}{2!} (x - a)^2 + \frac{f'''(a)}{3!} (x - a)^3\end{aligned}$$

### Taylor's Method for ODEs

This method takes advantage of the previously mentioned series. here this is some step size  $h$  that we take from some  $f(x)$  value. This is the Initial Value Problem (IVP).

$$\begin{aligned}f(x + h) &= f(x) + f'(x)h + \frac{1}{2!} f''(x)h^2 + \\ \frac{1}{3!} f'''(x)h^3 &+ \dots\end{aligned}$$

### Maclaurin Series

The Maclaurin series is just the Taylor series at the special case where  $x = 0$ . This gives the following:

$$\begin{aligned}f(x) &= f(0) + f'(0) + \frac{x^2}{2!} f''(0) + \frac{x^3}{3!} f'''(0) + \\ \frac{x^4}{4!} f''''(0) &+ \dots = \sum_{n=0}^{\infty} \frac{f^{(n)}(0)}{n!} x^n\end{aligned}$$

### Euler's Method for ODEs

This method is just a Taylor series of order 1 with the same step term  $h$ , though many steps can be taken:

$$f(x + h) = f(x) + f'(x)h$$

### Error Terms

We note that Taylor's theorem in terms of  $x + h$  is:

$$f(x + h) = \sum_{k=0}^n \frac{f^{(k)}(x)}{k!} h^k + E_{n+1}$$

Thus, error terms are of the form:

$$E_{n+1} = \frac{f^{(n+1)}(\xi)}{(n+1)!} h^{n+1}$$

It pays off to look at the term more specifically for the problem. A lot of times the error term takes the form  $\frac{n^2}{n!}$  or  $\frac{n^2}{n}$ .

It is important to note that we only care about the  $0.5 \cdot 10^n$  if our desired accuracy is to the  $n$ th decimal. Thus we set  $E_{n+1} < 0.5 \cdot 10^n$

### Accuracy to Postition Example

Imagine that we have an error term it will be of the following form:

$$E_{n+1} = \frac{f^{(n+1)}(\xi)}{(n+1)!} h^{n+1} \text{ Let } E_{n+1} = \frac{x^{n+1}}{n!}$$

## 6 Runge-Kutta Methods

### RK4

This is the 4th order (RK4) Runge-Kutta method for the Initial Value Problem

(IVP):

$$x(t + h) = x(t) \frac{1}{6} (K_1 + 2K_2 + 2K_3 + K_4)$$

where the following are values of  $K_n$ :

$$K_1 = hf(t, x)$$

$$K_2 = hf\left(t + \frac{1}{2}h, x + \frac{1}{2}K_1\right)$$

$$K_3 = hf\left(t + \frac{1}{2}h, x + \frac{1}{2}K_2\right)$$

$$K_4 = hf\left(t + h, x + K_3\right)$$

first, the  $K_n$  values are calculated in succession. They the  $K_n$  values are filled into the first formula above.