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### MODEL OF THE PROBLEM

**Objective Function** 

$$\min f(x)$$

### **Constraints**

$$g(x) \leq 0$$

$$x_j \in \{0,1\} \ \forall j \in J$$

$$f: \mathbb{R}^n \to \mathbb{R}$$

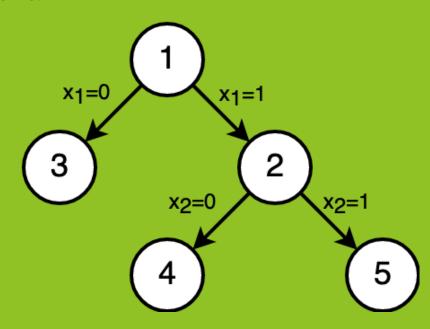
$$g: \mathbb{R}^n \to \mathbb{R}^m$$

$$J \subseteq N \coloneqq \{1 \dots n\}$$

### TYPICAL APPROACH

**STRATEGY USED:** Best Bound First

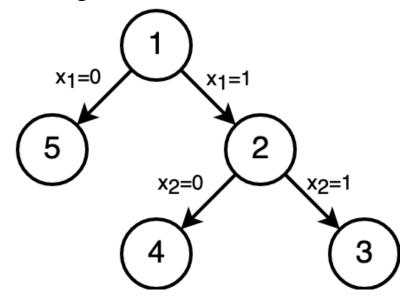
**DRAWBACK:** search tends to be initially trapped in the upper part of the tree, where lower bounds are smaller, which is counterproductive in terms of probability of updating the incumbent.



# PROXIMITY SEARCH APPROACH

### **MIXED STRATEGY:**

- 1. Selecting the next node to elaborate (e.g. using best-bound first)
- 2. Making a sequence of diving branching steps by visiting the nodes at increasing distance from the root.



### STEP 1

Finding an initial feasible solution  $\tilde{x}$  by:

- An ad-hoc heuristic
- Running a black-box MIP solver

### STEP 2

Adding an explicit cutoff constraint ( $\theta > 0$ )

$$f(x) \le f(\tilde{x}) - \theta$$

### STEP 3

Replacing the objective function by a proximity one (e.g. the Hamming distance):

$$\Delta(x,\tilde{x}) \coloneqq \sum_{j \in J: \tilde{x}_j = 0} x_j + \sum_{j \in J: \tilde{x}_j = 1} (1 - x_j)$$

$$\tilde{x} = (101001)$$
  
 $x = (011101)$   
 $\Delta(x, \tilde{x}) = 3$ 

### STEP 4

Run the black-box MIP solver to hopefully find a new incumbent  $x^*$  with  $f(x^*) \le f(\tilde{x}) - \theta$ 

### STEP 5

The new solution  $x^*$  is possibly improved by solving a convex problem where all the binary variables have been fixed to their value in  $x^*$ 

### STEP 6

Recenter  $\Delta(x, \cdot)$  by setting  $\tilde{x} := x^*$ , and/or update  $\theta$ 

### THE ALGORITHM

### Proximity Search:

- 1. let  $\tilde{x}$  be the initial heuristic feasible solution to refine; repeat
- 2. explicitly add the *cutoff constraint*  $f(x) \leq f(\tilde{x}) \theta$  to the MIP model;
- 3. replace f(x) by the "proximity" objective function  $\Delta(x, \tilde{x})$ ;
- 4. run the MIP solver on the new model until a termination condition is reached, and let  $x^*$  be the best feasible solution found ( $x^*$  empty if none); if  $x^*$  is nonempty and  $J \subset N$  then
- 5. refine  $x^*$  by solving the convex program  $x^* := \operatorname{argmin}\{f(x) : g(x) \leq 0, x_j = x_j^* \, \forall j \in J\}$

#### end

6. recenter  $\Delta(x,\cdot)$  by setting  $\tilde{x} := x^*$ , and/or update  $\theta$  until an overall termination condition is reached;

# Proximity Search Variants



### **Variants**

Possible alternative implementations of the basic algorithm:

Proximity
Search without
recentering

Proximity objective function never changes

Proximity
Search with
recentering

Proximity objective function changes when a new incumbent is found

Proximity
Search with
an incumbent

Each new incumbent triggers the MIP solver internal refinement heuristic

# Proximity Search without recentering (proxy\_norec)

### **Features**

- Each new incumbent  $\bar{x}$  is declared unfeasible
- New cutoff constraint:  $f(x) \le f(\bar{x}) - \theta$

**Drawbacks** 

The MIP incumbent is never updated explicitly

- No propagation scheme
- No variable-fixing scheme
- No refinement heuristic

# Proximity Search without recentering: algorithm

### Proximity Search, no recentering (proxy\_norec):

- 1. let  $\tilde{x}$  be the initial heuristic feasible solution to refine;
- 2. explicitly add the *cutoff constraint*  $f(x) \le f(\tilde{x}) \theta$  to the MIP model;
- 3. replace f(x) by the "proximity" objective function  $\Delta(x, \tilde{x})$ ;
- 4. run the MIP solver on the new model until a termination condition is reached or no more feasible solutions are found;

### **Callback** (when a new incumbent is discovered):

- 5. Let  $\bar{x}$  be the discovered incumbent, record  $\bar{x}$ ;
- 6. Introduce the *new cutoff constraint*  $f(x) \le f(\bar{x}) \theta$ ;
- 7. refinement of the current incumbent;

# Proximity Search with recentering (proxy\_rec)

### **Features**

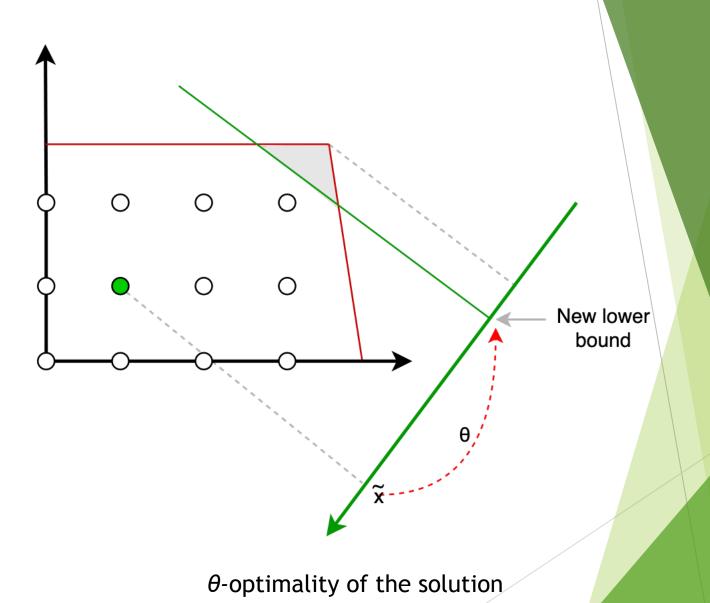
- Each new incumbent is the center of the proximity objective function  $\Delta(x, \cdot)$
- MIP solver as a Black Box

### **Drawbacks**

The MIP incumbent is never updated \_\_\_\_\_ explicitly

- No propagation scheme
- No variable-fixing scheme
- No refinement heuristic

Proximity Search with recentering: property



# Proximity Search with recentering: algorithm

### Proximity Search, with recentering (proxy\_rec):

- 1. let  $\tilde{x}$  be the initial heuristic feasible solution to refine; repeat
- 2. explicitly add the *cutoff constraint*  $f(x) \le f(\tilde{x}) \theta$  to the MIP model;
- 3. replace f(x) by the "proximity" objective function  $\Delta(x, \tilde{x})$ ;
- 4. run the MIP solver on the new model until the first feasible solution is found and let it be  $x^*$  ( $x^*$  empty if none);
- 5. refinement of the current incumbent  $x^*$ ;
- 6. recenter  $\Delta(x, \bullet)$  by setting  $\tilde{x} := x^*$ ;

until an overall termination condition is reached;

# Proximity Search with an incumbent (proxy\_incum)

### **Features**

New proximity objective function:

$$\Delta(x, \cdot) + Mz \qquad M > 0$$

Soft cutoff constraints:

$$f(x) \le f(\tilde{x}) - \theta + z$$
  $z \ge 0$ 

**Drawbacks** 

A too aggressive (large)  $\theta$  would make this approach useless:

$$\theta = f(\tilde{x}) \longrightarrow z = f(x)$$

# Proximity Search with an incumbent: algorithm

### Proximity Search, with an incumbent (proxy\_incum):

- 1. let  $\tilde{x}$  be the initial heuristic feasible solution to refine; **repeat**
- 2. explicitly add the soft cutoff constraint  $f(x) \le f(\tilde{x}) \theta + z$  to the MIP model;
- 3. replace f(x) by the "proximity" objective function  $\Delta(x, \tilde{x}) + Mz$  (M > 0);
- 4. run the MIP solver on the new model until the first feasible solution is found and let it be  $x^*$  ( $x^*$  empty if none);
- 5. refinement of the current incumbent  $x^*$ ;
- 6. recenter  $\Delta(x, \cdot)$  by setting  $\tilde{x} := x^*$ ; until an overall termination condition is reached:

## Trivial example (1)

Let's consider the Proximity Search with recentering applicated on a trivial MIP problem:

min 
$$5x_1 + 3x_2 + 4x_3 + 2x_4$$
  
 $x_1 + x_2 \ge 1$   
 $x_4 \ge 7x_1$   
 $x_4 \ge 8x_2$   
 $x_i \in \{0;1\} \ i=1,2,3$   
 $x_4 \ge 0$ 

Suppose to have the initial (feasible) solution  $\tilde{x} = (1,1,1,8)$ Set  $\theta = 2$ 

Use the Hamming Distance as "proximity" objective function:

$$\Delta(x,\,\tilde{x})=3-x_1-x_2-x_3$$

The objective function in  $\tilde{x}$  has value:  $f(\tilde{x}) = 28$ 

## Trivial example (2)

The model after Step 1 and 2 of the algorithm:

$$\min \frac{3 - x_1 - x_2 - x_3}{3 - x_1 - x_2 - x_3} \quad \text{Hamming Distance}$$

$$5x_1 + 3x_2 + 4x_3 + 2x_4 \le 28 - 2 \quad \text{Cutoff constraint}$$

$$x_1 + x_2 \ge 1$$

$$x_4 \ge 7x_1$$

$$x_4 \ge 8x_2$$

$$x_i \in \{0;1\} \ \ i = 1,2,3$$

$$x_4 \ge 0$$

$$\Delta(x, \tilde{x}) = 1 : (0, 1, 1, 8) \quad -> \quad f(x^*) = 23 \le 26 \quad -> \quad \text{feasible solution}$$

$$(1, 0, 1, 8)$$

$$(1, 1, 0, 8)$$

 $x^*$ = (0, 1, 1, 8) is the new incumbent, it will be the new center  $\tilde{x}$ 

## Trivial example (3)

The model, recentered, after Step 1 and 2 of the algorithm:

min 
$$x_1 + 2 - x_2 - x_3$$
 Recentered Hamming Distance  $5x_1 + 3x_2 + 4x_3 + 2x_4 \le 26$  Previous cutoff constraint  $5x_1 + 3x_2 + 4x_3 + 2x_4 \le 23 - 2$  New cutoff constraint  $x_1 + x_2 \ge 1$   $x_4 \ge 7x_1$   $x_4 \ge 8x_2$   $x_i \in \{0;1\} \ i=1,2,3$   $x_4 \ge 0$ 

# Performances and comparisons



## Related approaches in the literature

The ideas from which the Proximity Search takes its success are not new in the literature:

Replacing constraints with penalizations

Augmented Lagrangian (Hestenes, Powell, 1969) Replacing the objective function with a distance-related function

Parametric Branch and Bound Scheme (Glover, 1978)

Feasability Pump (Fischetti, Glover, Lodi, 2005 - in the interpretation of Boland, 2012) Using a cutoff constraint

Parametric Tabu Search (Glover, 2006)

## Related approaches in the literature

The ideas from which the Proximity Search takes its success are not new in the literature:

Solving subproblems with fixed variables

Relaxation Induced Neighborhood Search (Danna, Rothberg, Le Pape, 2005) Inserting a distance constraint

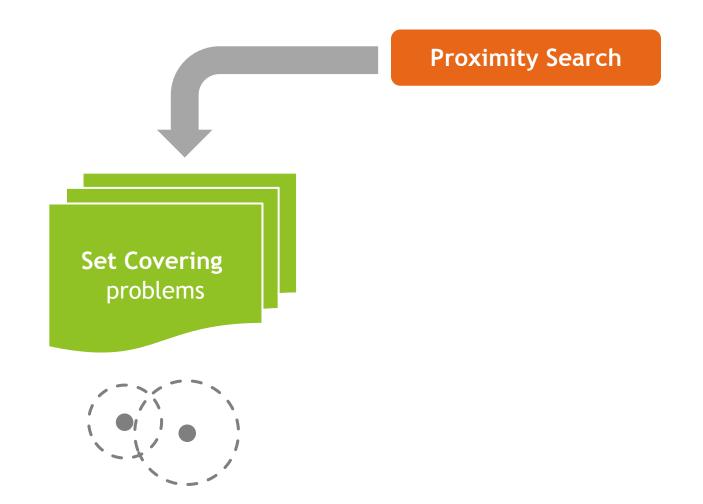
Local Branching (Fischetti, Lodi, 2003)

So, is **Proximity Search** a good choice?

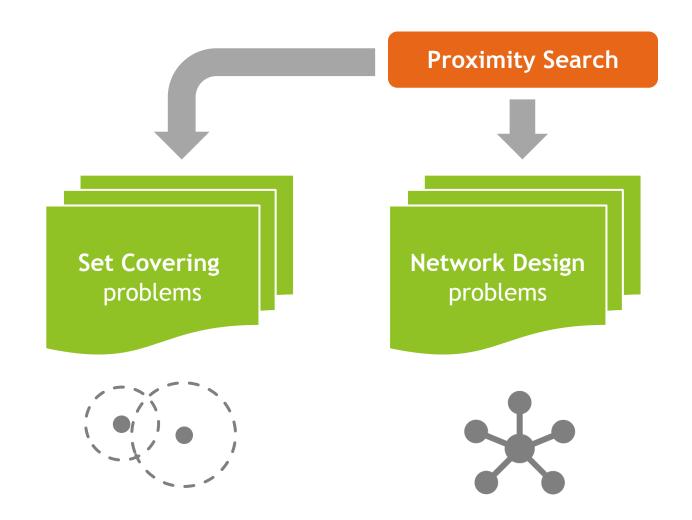
In the original article, Proximity Search has been **tested** on three different classes of problems:

**Proximity Search** 

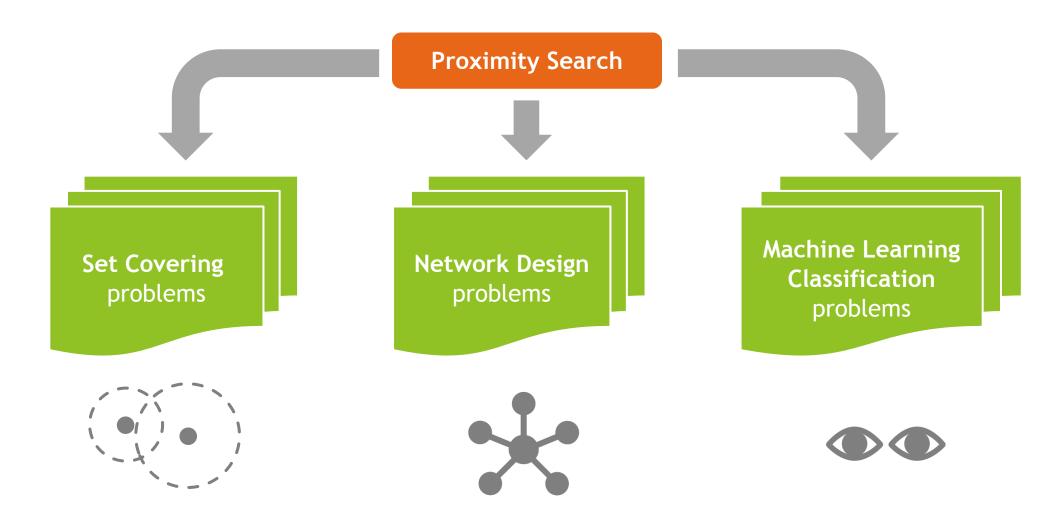
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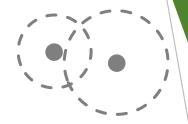
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## Set Covering Performances



|                        | time limit (s) |       |       |       |       |       |       |       |
|------------------------|----------------|-------|-------|-------|-------|-------|-------|-------|
|                        | 5              | 10    | 30    | 60    | 120   | 300   | 600   | 1,200 |
| Set covering instances |                |       |       |       |       |       |       | _     |
| proxy_norec            | 0.132          | 0.215 | 0.452 | 0.703 | 1.090 | 1.886 | 2.851 | 4.247 |
| cplex_def              | 0.178          | 0.310 | 0.698 | 1.121 | 1.753 | 2.880 | 4.108 | 5.775 |
| cplex_heu              | 0.174          | 0.305 | 0.703 | 1.113 | 1.671 | 2.697 | 3.774 | 5.086 |
| cplex_no_cuts          | 0.176          | 0.305 | 0.694 | 1.138 | 1.760 | 2.865 | 3.949 | 5.301 |
| cplex_gui_div          | 0.175          | 0.297 | 0.651 | 1.031 | 1.594 | 2.605 | 3.565 | 4.750 |
| proxy_incum            | 0.124          | 0.195 | 0.374 | 0.550 | 0.797 | 1.232 | 1.600 | 1.978 |
| proxy_rec              | 0.122          | 0.198 | 0.400 | 0.599 | 0.858 | 1.335 | 1.749 | 2.182 |
| locBra_orig            | 0.170          | 0.278 | 0.551 | 0.803 | 1.122 | 1.722 | 2.304 | 2.900 |
| locBra_aggr            | 0.121          | 0.192 | 0.376 | 0.561 | 0.773 | 1.157 | 1.533 | 1.974 |
| cplex_polish           | 0.181          | 0.298 | 0.596 | 0.876 | 1.251 | 1.895 | 2.498 | 3.252 |

Comparison metric: geometric mean of **primal integral** (the lower, the better)

algorithms

## Set Covering Performances



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Proximity Search (three variants)

Comparison metric: geometric mean of **primal integral** (the lower, the better)

algorithms

## Set Covering Performances



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Proximity Search (three variants)

Local Branching (aggressive variant)

Comparison metric: geometric mean of **primal integral** (the lower, the better)

algorithms

## Comparison with Local Branching

### **Proximity Search**

 $\min \Delta(x, \tilde{x})$ 

$$f(x) \le f(\tilde{x}) - \theta$$

$$g(x) \le 0$$

Objective Function

Other constraints

### **Local Branching**

 $\min f(x)$ 

$$\Delta(x, \tilde{x}) \le k$$

$$g(x) \le 0$$

## Comparison with Local Branching

### **Proximity Search**

 $\min \Delta(x, \tilde{x})$ 

$$f(x) \le f(\tilde{x}) - \theta$$

$$g(x) \le 0$$

**Objective Function** 

Cutoff constraintNeighborhood constraint

Other constraints

**Local Branching** 

 $\min f(x)$ 

$$\Delta(x, \tilde{x}) \le k$$

$$g(x) \le 0$$

## Comparison with Local Branching

### **Proximity Search**

 $\min \Delta(x, \tilde{x})$ 

$$f(x) \le f(\tilde{x}) - \theta$$

$$g(x) \leq 0$$

Neighborha Cutoff constraint

Other constraints

### **Local Branching**

 $\min f(x)$ 

$$\Delta(x, \tilde{x}) \le k$$

$$g(x) \leq 0$$

The cutoff constraint does not exclude any improving solutions (for small values of  $\theta$ ).

For problems where a large improvement is possible for a small neighborhood, Local Branching is faster.

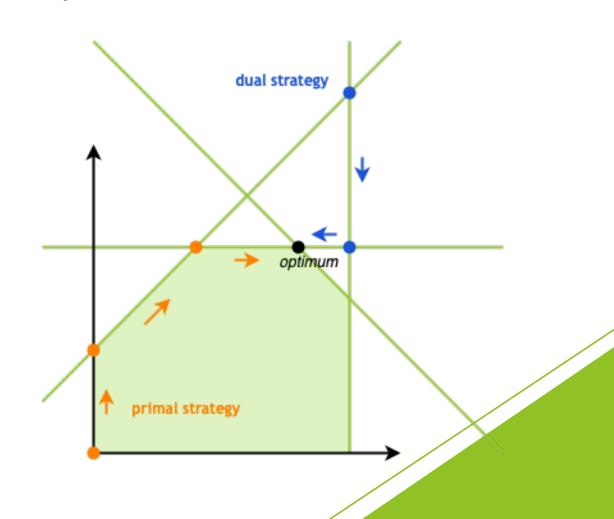
We can say that Proximity Search has a **primal nature**:

### **Primal Methods**

Produce a sequence of improved (feasible) solutions that eventually leads to an optimal one.

### **Dual Methods**

Eventually reach the optimal feasible solution through a sequence of more-than-optimal (infeasible) solutions.



We can say that Proximity Search has a **primal nature**:

**Primal Methods** 

**Dual Methods** 

We can say that Proximity Search has a **primal nature**:

**Primal Methods** 



Much more satisfactory in terms of "behavior as a **heuristic**".

**Dual Methods** 

A dual method stopped before the optimum is found, gives **no guarantee** to produce a feasible solution.

We can say that Proximity Search has a primal nature:

#### **Primal Methods**



Much more satisfactory in terms of "behavior as a **heuristic**".



Proximity Search can be trapped in a long series of small improvements.

### **Dual Methods**

A dual method stopped before the optimum is found, gives **no guarantee** to produce a feasible solution.



A more aggressive dual policy can produce less frequent but much larger improvements.



## Bibliography

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### Presentation adapted from:

Fischetti, M., Monaci, M.,

Proximity search for 0-1 mixed-integer convex programming,

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