

RSA Encryption

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RSA is a public key encryption system which mainly uses 3 components:

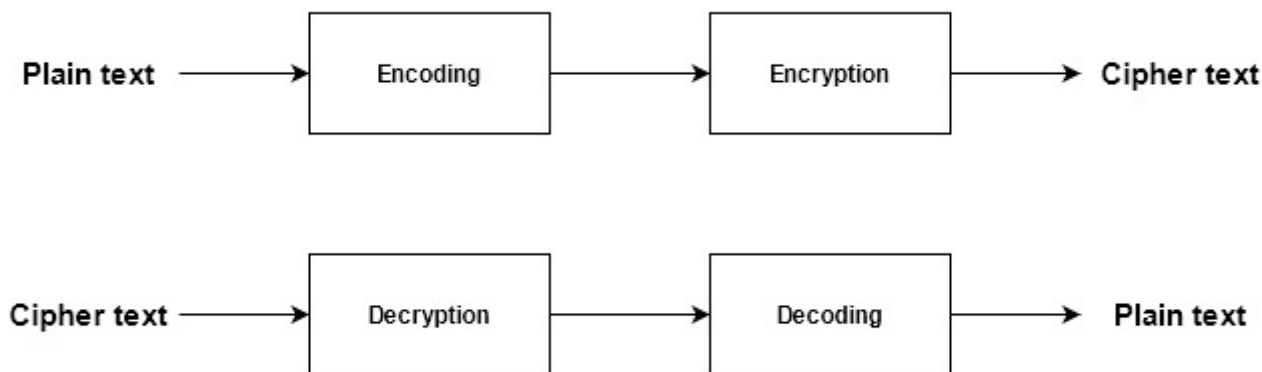
- **e** : the public key
- **d** : the private key
- **n** : a very large number created during the key generation process for modulus operations

This project uses the RSA algorithm to encrypt a series of encoded text using the publicly known parameters **e** and **n**. One of the potential attacks on RSA based encryption is when **n** is small, it can be factorized to retrieve the parameters used to generate it, and thus calculate **d**, the decryption key.

In this project, we use **n = 31313**, which is relatively small and can easily be factorized to calculate **d** and compromise the encryption system. The calculation of **d** is implemented in the class constructor in `RSA/RSA.cpp`.

To speed up the process of encryption and decryption, we use the fast-exponentiation algorithm with the time complexity $O(b)$, where b is the number of bits in the exponent.

The encryption-decryption process implemented in this project follows the following pipeline



Encoding and decoding

The encoding process is defined by the function `encode()` in `RSA/RSA.cpp`. The characters A-Z are assigned a value 0-25 (in order) and a 3 letter word is encoded using the following pattern:

$$\text{encoded value} = \text{character1} * 26^2 + \text{character2} * 26^1 + \text{character 3} * 26^0$$

The encoding only occurs for words with 3 characters.

The decoding process is the inverse of encoding where the 3 characters are extracted from the *encoded value*. This is defined by `decode()` in `RSA/RSA.cpp`.

Calculation of d

To calculate the private key d , we use the factorization method to obtain the two prime numbers used to generate n . It should be noted that n is a product of two large prime numbers p and q . So, if these are calculated, we can find $\phi(n)$ which is a product of $p-1$ and $q-1$. Finally, d is calculated by obtaining the inverse of e in $\text{mod}(\phi(n))$.

Encryption and decryption

A plain-text encoding P is encrypted by calculating $P^e \text{ mod}(n)$. A cipher-text C is decrypted by calculating $C^d \text{ mod}(n)$.