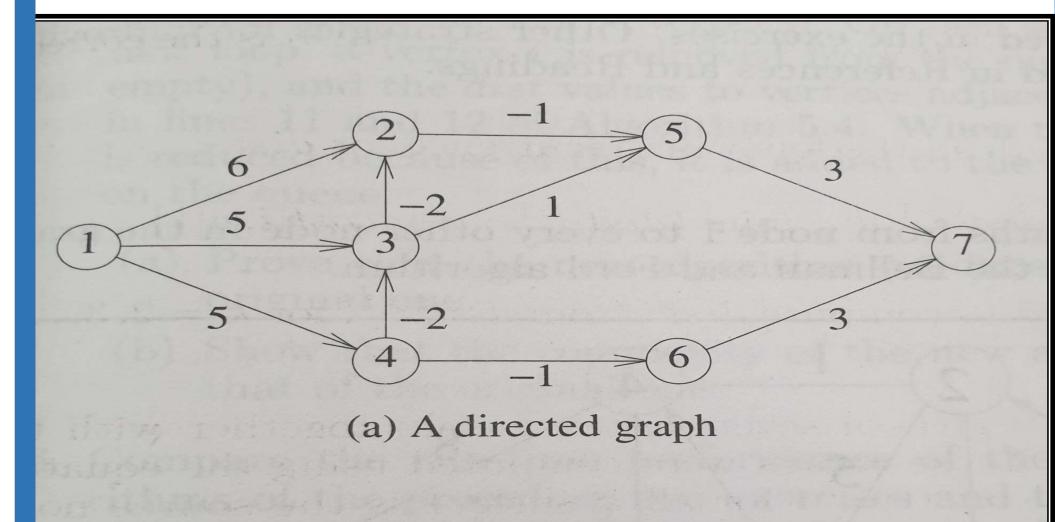
# Bellman Ford Algorithm

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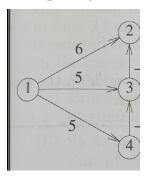
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- Single Source Shortest Path Algorithm
- Permits Graph with Negative Edges
- Significance of Negative Edge
- Negative edges signifies facts like: discount, free offer, schemes etc.
- There is possibility of cycle in a graph. [Problem?]
- Generally with negative edges present in graph, the cost of traversal can be zero or some times negative.
- Negative cost of traversal means, the customer is paid for travelling instead of paying for services.
- Example: Jio plans, Airline traffic etc

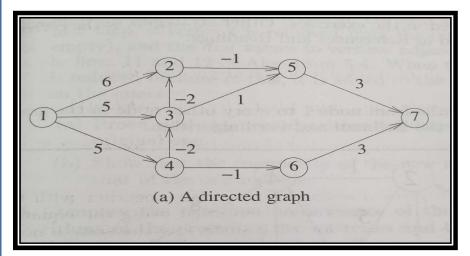


- Stage wise solution is designed
- Stage 1: Select Source Vertex : Given vertex: Vertex 1
- Stage 2: The path of length = 1 will be allowed.
- It means: computations for all the vertices directly reachable from source vertex can be selected as reached vertex [cost will be as per the graph]



K	1	2	3	4	5	6	7
1	0	6	5	5	00	00	00

- Stage 3: Path of length = 2
- All the vertices reachable in path of length=2, are selected and distances are computed as per the graph.



K	1	2	3	4	5	6	7
1	0	6	5	5	00	00	00
2	0	3	3	5	5	4	00

• All vertices except vertex-7 are reachable.

#### **Bellman Ford Algorithm: Formulation**

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Formula used for computing Shortest Path:
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dist^{k}[u] = min \{ dist^{k-1}[u], min \{ dist^{k-1}[i] + cost[i, u] \} \}
```

where "i" represents all the vertices except vertex "u" present in the graph. The vertex "i" should have connection to vertex "u" in the graph.

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Stage 4: Path of length = 3
(Sample Calculations)

dist<sup>3</sup>[u] = min { dist<sup>2</sup>[u] min { dist<sup>3</sup>[u] min { dist
```

 $dist^{3}[u] = min { dist^{2}[u], min { dist^{2}[i] + cost[i, u]}}$ 

**Consider Vertex: 2 as destination** 

How to avoid complexity of calculation

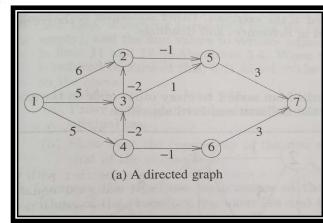
Refer the previous calculation

What is "i" [Either 1 or 3]

If i=1 then distance is 6

If i=3 then distance will be  $dist[3]+cost[3,2] \rightarrow$  previous computations

 $dist^{3}[2] = min \{ dist^{2}[2], min \{ dist^{2}[3] + cost[3,2] \} \} = 1$ 



• Final matrix for the given graph

K	1	2	3	4	5	6	7
1	0	6	5	5	$\infty$	$\infty$	$\infty$
2	0	3	3	5	5	4	$\infty$
3	0	1	3	5	2	4	7
4	0	1	3	5	0	4	5
5	0	1	3	5	0	4	3
6	0	1	3	5	0	4	3

## **Bellman Ford Algorithm: Home Work**

