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TRANSITIVE CLOSURE OF A GRAPH

Given a directed graph, find out if a vertex j is reachable from another vertex i for all vertex pairs (i, j) in the given graph. Being reachable mean that there is a path from vertex i to j . The reachability matrix is called transitive closure of a graph.

The all-pairs shortest path problem is the determination of the shortest graph distances between every pair of vertices in a given graph. The problem can be solved using applications of Dijkstra's algorithm or all at once using the Floyd-Warshall algorithm.

The Floyd-Warshall algorithm is an algorithm for finding shortest paths in a weighted graph with positive or negative edge weights (but with no negative cycles).

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FLYOD-WARSHALL ALGORITHM

The Floyd-Warshall algorithm is an algorithm for finding shortest paths in a weighted graph with positive or negative edge weights (but with no negative cycles). A single execution of the algorithm will find the lengths (summed weights) of the shortest paths between all pairs of vertices.

The Floyd-Warshall algorithm is an example of dynamic programming, and was published in its currently recognized form by Robert Floyd in 1962. However, it is essentially the same as algorithms previously published by Bernard Roy in 1959 and also by Stephen Warshall in 1962 for finding the transitive closure of a graph, and is closely related to Kleene's algorithm (published in 1956) for converting a deterministic finite automaton into a regular expression. The modern formulation of the algorithm as three nested for-loops was first described by Peter Ingerman, also in 1962.

The algorithm is also known as Floyd's algorithm, the Roy-Warshall algorithm, the Roy-Floyd algorithm, or the WFI algorithm.

20 PATH RECONSTRUCTION using WFI Algorithm

The Floyd-Warshall algorithm typically only provides the lengths of the paths between all pairs of vertices. With simple modifications, it is possible to create a method to reconstruct the actual path between any two endpoint vertices. While one may be inclined to store the actual path from each vertex to each other vertex, this is not necessary, and in fact, is very costly in terms of memory. Instead, the shortest-path tree can be calculated for each node in $\Theta(|E|)$ time using $\Theta(|V|)$ memory to store each tree which allows us to efficiently reconstruct a path from any two connected vertices.

Let,

$dist[][]$ be a $|V| \times |V|$ array of minimum distances initialized to INF
 $next[][]$ be a $|V| \times |V|$ array of vertex indices initialized to NULL

procedure FloydWarshallWithPathReconstruction(G)

 for each edge (u, v)

$dist[u][v] \leftarrow w(u, v)$ // the weight of the edge (u, v)

$next[u][v] \leftarrow v$

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        for k from 1 to |V| // standard Floyd-Warshall implementation
            for i from 1 to |V|
                for j from 1 to |V|
                    if dist[i][j] > dist[i][k] + dist[k][j] then
                        dist[i][j] ← dist[i][k] + dist[k][j]
                        next[i][j] ← next[i][k]

40 procedure Path(u, v)
    if next[u][v] = NULL then
        return []

    path = [u]

    while u ≠ v
        u ← next[u][v]
        path.append(u)

50 return path

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#include <stdio.h>

#define NF 786
#define MX 10

void showMatrix(int graph[][MX], int vertices, const char *text){
60     int i, j;

    printf("\n\t%s is....\n", text);
    printf("\t\t u|v |");
    for(i=0; i < vertices; i++)
        printf("%4d ", i);
    printf("\n");

    printf("\t\t-----");
    for(i=0; i < vertices; i++)
70         printf("-----");
    printf("\n");

    for(i=0; i < vertices; i++){
        printf("\t\t%5d |", i);
        for(j = 0; j < vertices; j++)
            printf("%4d ", graph[i][j]);
        printf("\n");
    }
    //printf("\n");

80 }

void tcFloydWarshall (int graph[][MX], int vertices){
    int dist[MX][MX], i, j, k;
    int next[MX][MX];

    for (i = 0; i < vertices; i++){
        for (j = 0; j < vertices; j++){
90             dist[i][j] = graph[i][j];

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        next[i][j] = j;
    }
}

printf("\nAt Initialization ...");
showMatrix(dist, vertices, "Cost Matrix");
showMatrix(next, vertices, "Path Matrix");

for (k = 0; k < vertices; k++){
100     printf("\nHit Enter key to Proceed ...");
        getc(stdin);
        getc(stdin);

        printf("ITERATION %2d\n", k+1);
        for (i = 0; i < vertices; i++){
            for (j = 0; j < vertices; j++){
                if (dist[i][j] > dist[i][k] + dist[k][j]){
                    dist[i][j] = dist[i][k] + dist[k][j];
                    next[i][j] = next[i][k];
110                }
            }
        }
        sleep(3);

        showMatrix(dist, vertices, "Cost Matrix");
        showMatrix(next, vertices, "Path Matrix");
    }
}

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void initMatrix(int graph[][MX]){
    int i, j, weight;
    for(i=0; i < MX; i++)
        for(j = 0; j < MX; j++)
            graph[i][j] = NF;
}

130 int createGraph(int graph[][MX]){
    int i, j, vCnt=0, weight;
    int u, v, vertices, type;

    printf("\n\tGraph Creation [Undirected/Directed]...\n");

    printf("\t\tType of Graph [0: UnDirected] := ");
    scanf("%d", &type);

    if(type != 0)
140         type = 1;

    do{
        printf("\t\tHow Many Vertices [upto %d vertices]?? ", MX);
        scanf("%d", &vertices);
    }while(vertices < 1 || vertices > MX);

    printf("\n");
}

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printf("\n\tVertices starts at 0 and terminates at %d\n", vertices-1);
printf("\t\tVertex ID of -1 terminates Input\n");
printf("\n\tEnter Existing Edges in the Graph\n\n");
printf("\t\t-----\n");
printf("\t\tEdge#      'u'      'v'      Cost      Remark\n");
printf("\t\t-----\n");

do{
    do{
        printf("\t\t   %2d          ", vCnt+1);
        scanf("%d%d%d", &u, &v, &weight );

        if(u == v && weight < 0)
            printf("\t\t\t\t\t\t\t\t\t\tNegative Cycle\n");
        while(u == v && (u != -1 && v != -1) && weight < 0);

        printf(" \t ");

        if((u != -1 || v != -1) && u < vertices && v < vertices){
            if(graph[u][v] == NF){
                if(type)
                    graph[u][v] = weight;
                else
                    graph[u][v] = graph[v][u] = weight;

                printf("\t\t\t\t\t\t\t\t\t\tEdge Taken\n");
            } else
                printf("\t\t\t\t\t\t\t\t\t\tEdge Exists\n");

        }else
            printf("\t\t\t\t\t\t\t\t\t\tInvalid Edge\n");

        vCnt++;

    }while(u != -1 || v != -1);

    for(i = 0; i < MX; i++)
        if(graph[i][i] == NF)
            graph[i][i] = 0;

    return vertices;
}

int printPathWFI(int dist[][MX], int next[][MX], int u, int v){
    int pathCost = NF;

    return pathCost;
}

int main(){

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int graph[MX][MX];

int vertices;
initMatrix(graph);

vertices = createGraph(graph);

printf("\nGraph with %2d vertices ...\n", vertices);
showMatrix(graph, vertices, "Adjacency Matrix");
tcFloydWarshall(graph, vertices);

return 0;
}

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/** ----- EXECUTION TRAIL -----

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Graph Creation [Undirected/Directed]...
Type of Graph [0: UnDirected] := 4
How Many Vertices [upto 10 vertices]?? 4

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Vertices starts at 0 and terminates at 3
Vertex ID of -1 terminates Input

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Enter Existing Edges in the Graph

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Edge#	'u'	'v'	Cost	Remark
1	0	2	3	Edge Taken
2	2	2	1	Edge Taken
3	2	3	2	Edge Taken
4	2	1	-2	Edge Taken
5	3	2	4	Edge Taken
6	1	0	4	Edge Taken
7	1	3	3	Edge Taken
8	-1	-1	12	Invalid Edge

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Graph with 4 vertices ...

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Adjacency Matrix is....

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u v	0	1	2	3
0	0	786	3	786
1	4	0	786	3
2	786	-2	1	2
3	786	786	4	0

```

At Initialization ...

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Cost Matrix is....

u v	0	1	2	3
0	0	786	3	786
1	4	0	786	3
2	786	-2	1	2
3	786	786	4	0

Path Matrix is....

u v	0	1	2	3
0	0	1	2	3
1	0	1	2	3
2	0	1	2	3
3	0	1	2	3

ITERATION 1

Cost Matrix is....

u v	0	1	2	3
0	0	786	3	786
1	4	0	7	3
2	786	-2	1	2
3	786	786	4	0

Path Matrix is....

u v	0	1	2	3
0	0	1	2	3
1	0	1	0	3
2	0	1	2	3
3	0	1	2	3

ITERATION 2

Cost Matrix is....

u v	0	1	2	3
0	0	786	3	786
1	4	0	7	3
2	2	-2	1	1
3	786	786	4	0

Path Matrix is....

u v	0	1	2	3
0	0	1	2	3
1	0	1	0	3
2	1	1	2	1
3	0	1	2	3

ITERATION 3

Cost Matrix is....

u v	0	1	2	3
0	0	1	3	4
1	4	0	7	3
2	2	-2	1	1
3	6	2	4	0

Path Matrix is....

u v	0	1	2	3
0	0	2	2	2
1	0	1	0	3
2	1	1	2	1
3	2	2	2	3

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ITERATION 4

Cost Matrix is....

u v	0	1	2	3
0	0	1	3	4
1	4	0	7	3
2	2	-2	1	1
3	6	2	4	0

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Path Matrix is....

u v	0	1	2	3
0	0	2	2	2
1	0	1	0	3
2	1	1	2	1
3	2	2	2	3

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