

Graphs

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- Vertices + edges (weighted?)
- Directed vs. undirected (bidirectional)
- Complete graph
 - o Clique: $O(n^2)$ edges
- Path: $v \rightarrow v$
 - o Length = #edges
 - o Simple = no repeated vertices
 - o Cycle = circular path (simple cycle)
- Connected: no discrete components

Problems

- Shortest path, Travelling salesman, Topological sort

Graphs

- $G = (V, E, w)$ // vertices, edges, weight function

Adjacency

- $\text{adj}(v)$ = neighbours/ successors of v
- Adjacency matrix - $O(N^2)$ memory, not suitable for sparse graph
 - o Undirected = symmetrical matrix
- Adjacency list - $O(V + E)$, store edges of each vertex

BFS Traversal

- Source $\rightarrow i \rightarrow i + 1$
- Not unique
- Gives a BFS tree (root = source)
 - o Vertex label = distance from source
- Connected & undirected graph: BFS visits all vertices
 - o Disconnected: call BFS k times
- Recursively put neighbours into queue
- $O(V + E)$

DFS Traversal

- Use stack to remember where to back-track to
- $O(V + E)$

Topological Sort

- Directed Acyclic Graph (DAG)
- In-degree/ out-degree: #edges in/ out
- Recursively put vertices with in-degree 0 into queue
 - o Dequeue: remove vertex + edges from graph
- Use adjacency list
 - o When queued, update in-degree of neighbours
- $O(V + E)$