

# Error Codes Part III: CRC Code Performance

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Receiver  $\frac{R(x)}{G(x)} = \frac{T(x) + E(x)}{G(x)} \Rightarrow \frac{\cancel{E(x)}}{\cancel{G(x)}} ?$

$\frac{E(x)}{G(x)} ?$

$$\begin{array}{r} T(x) = 10011010 \\ E(x) = 01000001 \\ R(x) = 11011011 \end{array}$$

Single Bit Error

$$\sum (x^i) = 00010000$$

$$\sum (x^i) = x^i \quad \text{power of 2}$$

If  $G(x) \geq 2$  terms (not power of 2)  
remainder will be produced

$$\begin{array}{r} 32 \\ \hline 172+4 \\ \hline 7 \end{array} \rightarrow \text{remainder}$$

$$\left. \begin{array}{l} G(x) \text{ has } 2 \\ \text{or more terms} \end{array} \right\}$$

Double Bit Error  $E(x) = 0010001000$

$$E(x) = x^i + x^j \quad i \neq j$$

$$= x^j (x^{i-j} + 1)$$

$$= \underbrace{(x^j)}_{G(x)} \underbrace{(x^k + 1)}_{G(x)}$$

$$k = i - j$$

So if  $G(x) \rightarrow G(x)$  does not divide  $x^k + 1$

example:  $x^{18} + x^{14} + 1$  will not

divide  $x^k + 1$  evenly for  $k$  up to 32, 78

$$\begin{array}{r} \underline{G(x)} \\ G(x) \end{array}$$

there's no polynomial with an odd # terms  
divisible evenly by  $(x+1)$   $\rightarrow$  So pick  
generator polynomials like  $(x+1)^p$  a  
factor.  $\rightarrow$  catch odd # errors

Proof

Assume  $\Sigma(x)$  has odd # terms  
and  $u$  divisible by  $x+1$

$$\Sigma(x) = (x+1)Q(x)$$

$$\text{Let } x=1$$

$$\begin{aligned} \Sigma(1) &= (1+1)Q(1) \\ &= 0 \cdot Q(1) = 0 \end{aligned}$$

But if  $\Sigma(x)$  has odd # terms,

$$\left. x^5 + x^4 + 1 \right|_{x=1} = 1+1+1 = 1$$

$$\begin{array}{r|l} A+B & A+B \\ \hline 0 & 0 \\ 0 & 1 \\ \hline 1 & 1 \\ 1 & 0 \\ \hline \end{array}$$

$$\Sigma(1) = 1$$

Burst Error

$$E(x) = 00111111000$$

$$E(x) = x^i (x^{k-i} + \dots + 1)$$

$$\text{If } G(x) \Rightarrow x^0 = 1 \quad \text{G(x)}$$

then  $x^i$  is not a factor  
so remainder is not 0

## Burst Error

If burst  $\leq r+1$  (where  $r$  is  
 # check bits) including a term  $x^r$   
 +1 ( $x^0$ ) in GCD will catch all

$$\text{burst} \leq r$$

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If burst  $> r+1$  / several short bursts

$$P(\text{bad message undetected}) = \left(\frac{1}{2}\right)^{r-1}$$



Interpolated Sums

$$x^{12} + x^{11} + x^3 + x^2 + x + 1$$

$$x^{16} + x^{15} + x^2 + 1$$

$$x^{16} + x^{12} + x^5 + 1$$

APX  
protocols

✓ All have +1

✓ All have (x+1) as a factor

✓ All have 2 terms