



Stony Brook University



ESE/CSE 346

Probability Problems (Part 1)

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Electrical & Computer Engineering



$$(a) \quad p = \frac{20}{60} = \frac{1}{3} \Rightarrow \text{prob a circuit busy } \left(\frac{1}{3}\right)$$

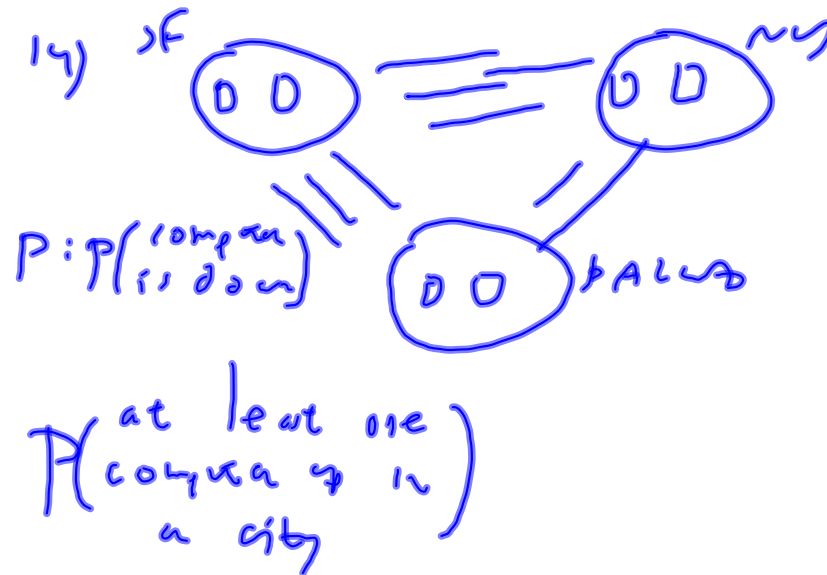
$$\begin{aligned}
 (b) \quad & \cancel{P} P(\text{at least one free circuit}) & \left| \begin{array}{l} q = 1 - p \\ = P(\text{free}) \end{array} \right. \\
 & = P(\text{find a free circuit}) \\
 & = 1 - P(0 \text{ circuits free}) \\
 & = 1 - (3) q^0 (1 - q)^3 = 1 - (1 - q)^3 = 1 - p^3
 \end{aligned}$$

(c)

$$P(\text{one circuit free}) = \binom{3}{1} \underbrace{p}_{\text{circuit free}} (1-p)^2$$

(d)

$$\begin{aligned} \# \text{ busy circuits} &= 1 \cdot P(1 \text{ busy}) + 2 \cdot P(2 \text{ busy}) + 3 \cdot P(3 \text{ busy}) \\ &= 1 \cdot \binom{3}{1} p (1-p)^2 + 2 \cdot \binom{3}{2} p^2 (1-p) + 3 \cdot \binom{3}{3} p^3 \underbrace{(1-p)^0}_{=1} \\ &= 3p(1-p)^2 + 6p^2(1-p) + 3p^3 \end{aligned}$$

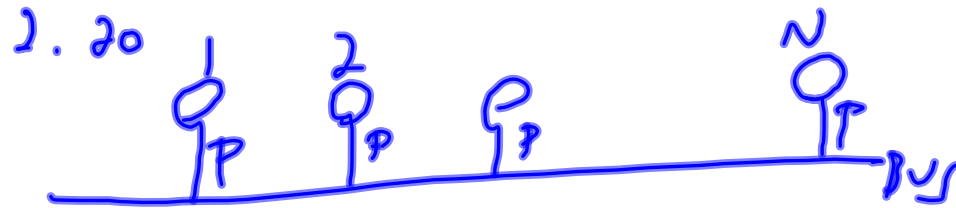


p : prob a computer down

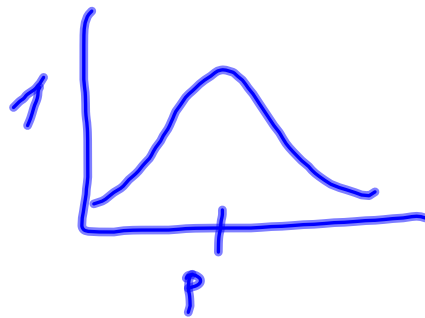
$$P\left(\begin{array}{l} \text{at least one} \\ \text{computer up in} \\ \text{a group} \end{array}\right) = 1 - P\left(\begin{array}{l} \text{both computers} \\ \text{down in} \\ \text{a group} \end{array}\right)$$

$$= 1 - p^2$$

$$P\left(\begin{array}{l} \text{at least one} \\ \text{computer up} \\ \text{in each group} \end{array}\right) = (1 - p^2)^3$$



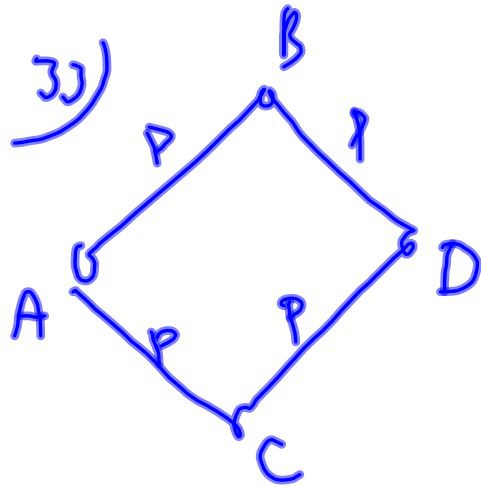
$$\begin{array}{c} | \quad | \quad | \quad | \\ \text{slot} \quad \text{slot} \quad \text{slot} \quad \text{time} \rightarrow \\ 1 \quad 2 \quad 3 \end{array} \quad \left| \begin{array}{l} \text{Thruput} = P(\text{all stations receive}) \\ = \underline{N p (1-p)^{N-1}} \end{array} \right.$$



$$\frac{d \text{Thruput}}{d p} = 0$$

$$\frac{1}{N} N = 1$$

$$p_{\text{optimal}} = \frac{1}{N}$$



p : $p(\text{link available})$

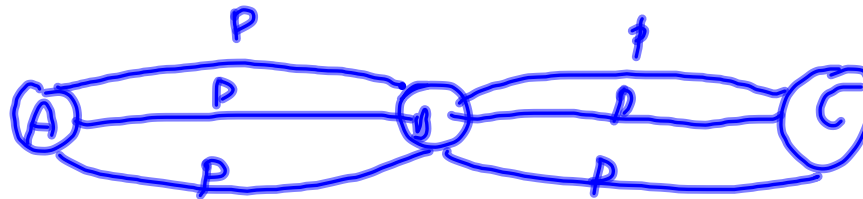
(a) $\mathcal{L} = p^2 = p(\text{AD path available})$

(a)

(b)
$$P\left[\begin{array}{c} \text{at least one} \\ \text{A to D path} \\ \text{available} \end{array}\right] = 1 - P\left[\begin{array}{c} \text{no AD} \\ \text{path available} \end{array}\right]$$

$$= 1 - (1 - p^2)^2$$

9)



p : link in use

(c)

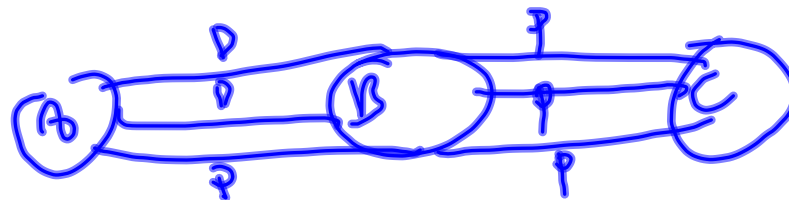
$\# \text{ of links} =$

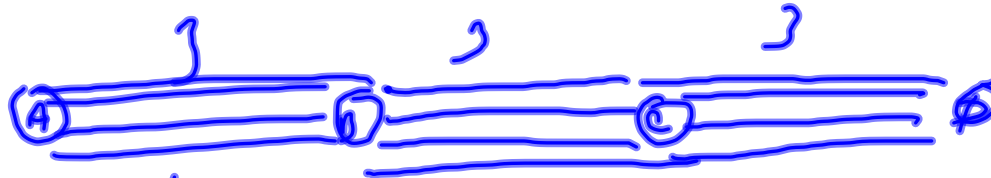
$$\sum_{n=1}^6 n p^n (1-p)^{6-n} \binom{6}{n}$$

(b)

p is low

$$\begin{aligned}
 P\left[\begin{array}{l} \text{at least one} \\ \text{idle path from} \\ A \text{ to } C \end{array}\right] &= \left(P\left[\begin{array}{l} \text{at least 1 idle} \\ \text{link from } A \text{ to } B \end{array}\right] \right)^2 \\
 &= \left(1 - P\left[\begin{array}{l} \text{no idle link} \\ \text{from } A \text{ to } B \end{array}\right] \right)^2 = (1 - p^2)^2
 \end{aligned}$$





p : prob link busy

(2) $P(A \rightarrow D)$
 $P(\text{all blocked})$

$P(\text{at least one link pair is completely blocked})$

$= 1 - P(\text{at least one link between each adjacent pair})$

$= 1 - \left(P(\text{at least one link between an adjacent pair of nodes}) \right)^3$

$= 1 - (1 - p^3)^3$