**Title: Algorithm Efficiency and Sorting** 

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Assignment: 1

# Question 1

## a) Insertion Sort (According to slide number 37)

4 | 8 3 7 6 2 1 5 Original List

48 | 3 7 6 2 1 5 After pass 1

3 4 8 | <mark>7</mark> 6 2 1 5 After pass 2

3 4 7 8 | <mark>6</mark> 2 1 5 After pass 3

3 4 6 7 8 | 2 1 5 After pass 4

234678 | 15 After pass 5

1234678 | <mark>5</mark> After pass 6

12345678 After pass 7

# b) Selection Sort (According to slide number 15)

48376215 Initial Array

453**7**621**| 8** After 1<sup>st</sup> swap

4531<mark>6</mark>2**78** After 2<sup>st</sup> swap

4 5 3 1 2 | 6 7 8 After 3st swap

4 2 3 1 | **5 6 7 8** After 4<sup>st</sup> swap

1 2 3 | 4 5 6 7 8 After 5<sup>st</sup> swap

1 2 | 3 4 5 6 7 8 After 6st swap

**1 | 2345678** After 7<sup>st</sup> swap

# c) Bubble Sort (According to slide number 42)

#### Pass 1

48376215 | Initial Array

48376215

43<mark>87</mark>6215

437<mark>86</mark>215

4376<mark>82</mark>15

43762<mark>81</mark>5

437621<mark>85</mark>

4376215**|8** 

## Pass 2

<mark>43</mark>76215**8** 

3 <mark>4 7</mark> 6 2 1 5 **8** 

34<mark>76</mark>215**8** 

346<mark>72</mark>15**8** 

3 4 6 2 <mark>7 1</mark> 5 **8** 

3 4 6 2 1 <mark>7 5</mark> **8** 

346215 **| 78** 

## Pass 3

**34**6215**78** 

3 <mark>4 6</mark> 2 1 5 **7 8** 

34<mark>62</mark>15**78** 

3 4 2 <mark>6 1</mark> 5 **7 8** 

3 4 2 1 <mark>6 5 **7 8**</mark>

34215**|678** 

## Pass 4

34215**678** 

3 <mark>4 2</mark> 1 5 **6 7 8** 

3 2 <mark>4 1</mark> 5 **6 7 8** 

3 2 1 <mark>4 5</mark> **6 7 8** 

3214 | 5678

## Pass 5

**32**14**5678** 

2 3 1 4 5 6 7 8

21<mark>34</mark>5678

213 | 45678

#### Pass 6

**21**3**45678** 

12345678

12 | 345678

## Pass 7

12 | 345678

12345678

## d) Merge Sort

Initial array: 48376215

mergesort(array, 0, 7)

mergesort(array, 0, 3)

mergesort(array, 0, 1)

mergesort(array, 0, 0)

mergesort(array, 1, 1)

merge(array, 0, 0, 1)

48376215

mergesort(array, 2, 3)

mergesort(array, 2, 2)

mergesort(array, 3, 3)

merge(array, 2, 2, 3)

48376215

merge(array, 0, 1, 3)

34786215

mergesort(array, 4, 7)

mergesort(array, 4, 5)

mergesort(array, 4, 4)

mergesort(array, 5, 5)

merge(array, 4, 4, 5)

34782615

mergesort(array, 6, 7)

mergesort(array, 6, 6)

mergesort(array, 7, 7)

merge(array, 6, 6, 7)

34782615

merge(array, 4, 5, 7)

34781256

merge(array, 0, 3, 7)

12345678

## e) Quick Sort

Initial array: 48376215

quicksort(array, 0, 7)

partition(array, 0, 7)

13246875

quicksort(array, 0, 2)

partition(array, 0, 2)

13246875

quicksort(array, 0, -1)

quicksort(array, 1, 2)

partition(array, 1, 2)

12346875

quicksort(array, 1, 1)

quicksort(array, 3, 2)

quicksort(array, 4, 7)

partition(array, 4, 7)

12345678

quicksort(array, 4, 4)

quicksort(array, 6, 7)

partition(array, 6, 7)

12345678

quicksort(array, 6, 5)

quicksort(array, 7, 7)

## **Question 2**

## Mergesort

$$T(1) = \Theta(1)$$

$$T(n) = 2T(n/2) + \Theta(n)$$

$$T(n) = 2(2T(n/4) + \Theta(n/2)) + \Theta(n)$$

$$T(n) = 4T(n/4) + 2\Theta(n/2) + \Theta(n)$$

$$T(n) = 4(2T(n/8) + \Theta(n/4)) + 2\Theta(n/2) + \Theta(n)$$

$$T(n) = 8T(n/8) + 4\Theta(n/4) + 2\Theta(n/2) + \Theta(n)$$

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$$T(n) = 2^k \, T(n/2^k) + 2^{k-1} \, \Theta(n/2^{k-1}) + 2^{k-2} \Theta(n/2^{k-2}) + \ldots + 2\Theta(n/2) + \Theta(n)$$

when k = logn

$$T(n) = nT(n/n) + (n/2)\Theta(n/(n/2)) + (n/4)\Theta(n/(n/4)) + ... + 2\Theta(n/2) + \Theta(n)$$

$$T(n) = nT(1) + \Theta(n) + \Theta(n) + \dots + \Theta(n) + \Theta(n)$$

$$T(n) = n\Theta(1) + k\Theta(n)$$

plugging the value of k into the equation

$$T(n) = n\Theta(1) + logn\Theta(n)$$

Answer: Θ(nlogn)

# Quicksort

$$T(1) = \Theta(1)$$

$$T(n) = T(n-1) + \Theta(n)$$

$$T(n) = T(n-2) + \Theta(n-1) + \Theta(n)$$

$$T(n) = T(n-3) + \Theta(n-2) + \Theta(n-1) + \Theta(n)$$

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$$T(n) = T(n-k) + \Theta(n-k+1) + \Theta(n-k+2) + ... + \Theta(n-2) + \Theta(n-1) + \Theta(n)$$

when k = n - 1

$$T(n) = T(1) + \Theta(2) + \Theta(3) + ... + \Theta(n-2) + \Theta(n-1) + \Theta(n)$$

$$T(n) = \Theta(1) + \Theta(2) + \Theta(3) + ... + \Theta(n-2) + \Theta(n-1) + \Theta(n)$$

$$T(n) = \Theta(n(n+1)/2)$$

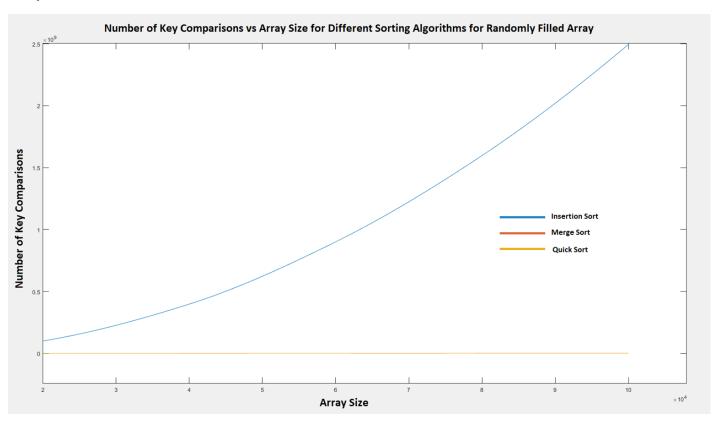
$$T(n) = \Theta(n^2/2) + \Theta(n/2)$$

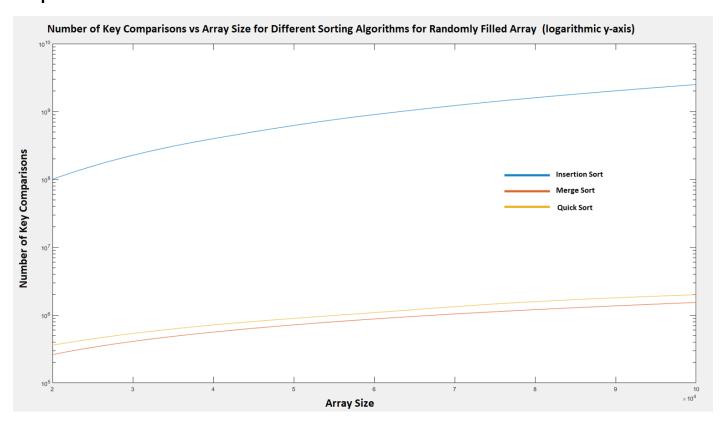
Answer: Θ(n²)

## **Question 3**

In the following graph, merge sort and quick sort curves look as if they overlap so I also plotted the same graph with logarithmic y-axis.

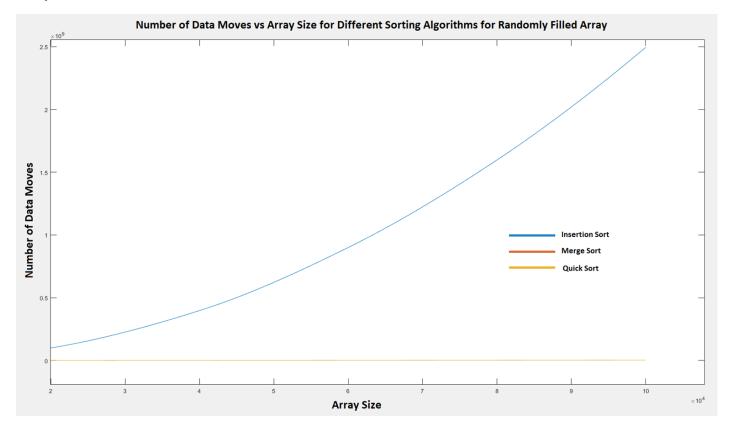
## Graph 1



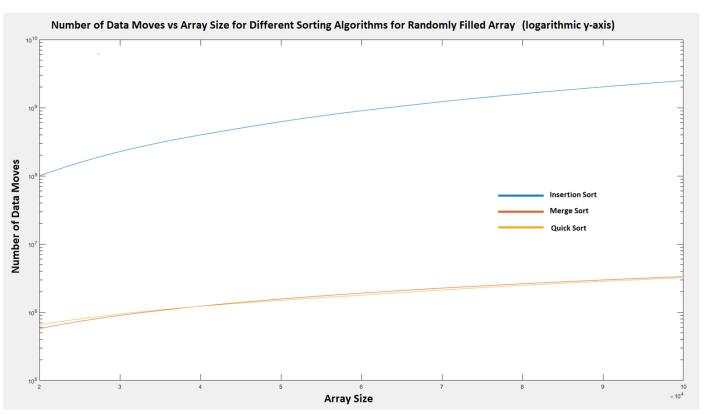


In the following graph, merge sort and quick sort curves look as if they overlap so I also plotted the same graph with logarithmic y-axis.

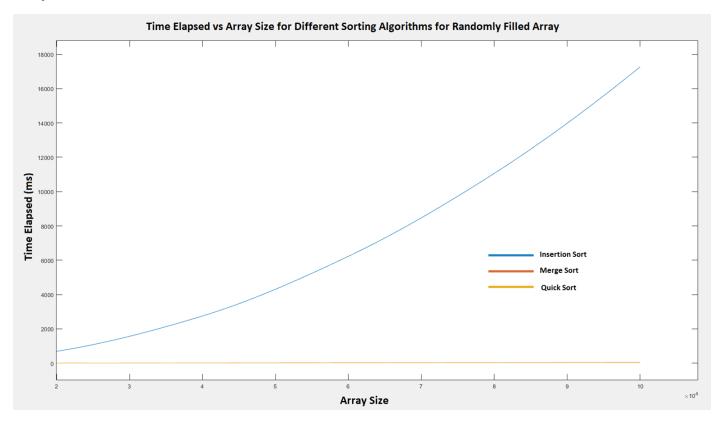
**Graph 3** 



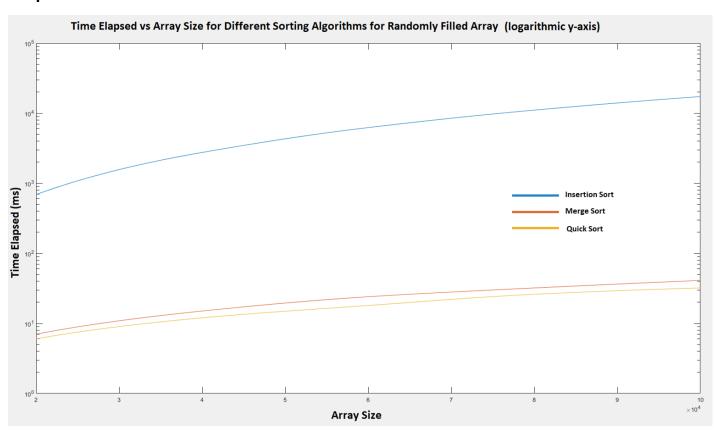
**Graph 4** 



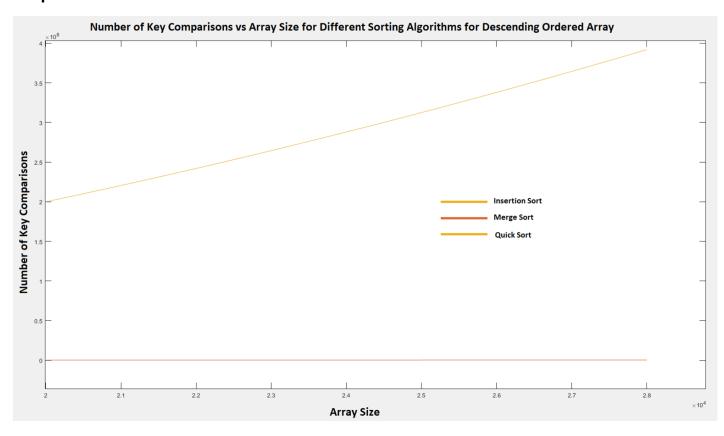
In the following graph, merge sort and quick sort curves look as if they overlap so I also plotted the same graph with logarithmic y-axis.

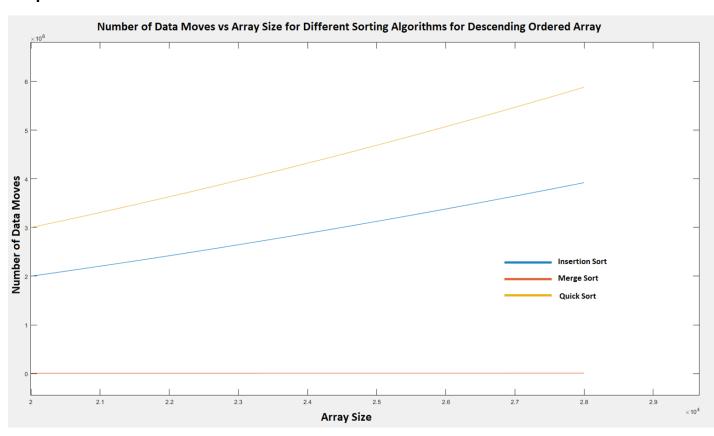


Graph 6

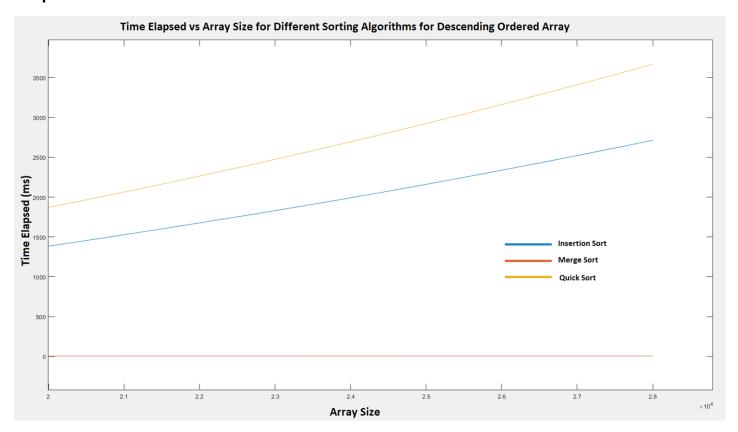


# Graph 7



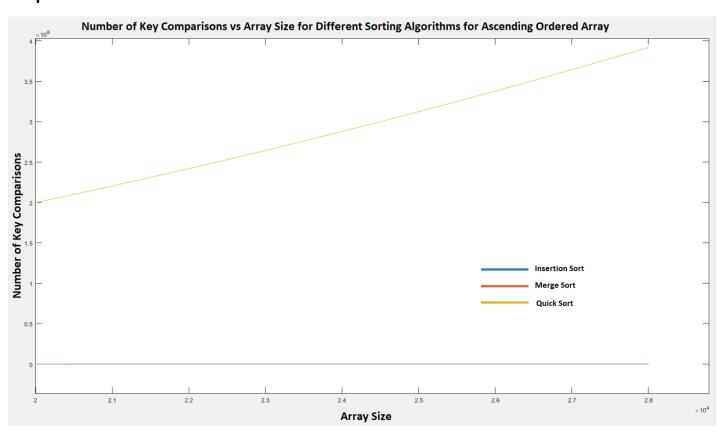


## Graph 9

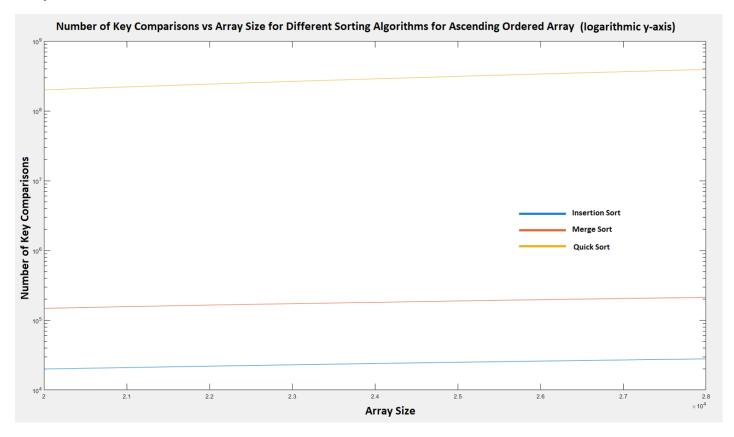


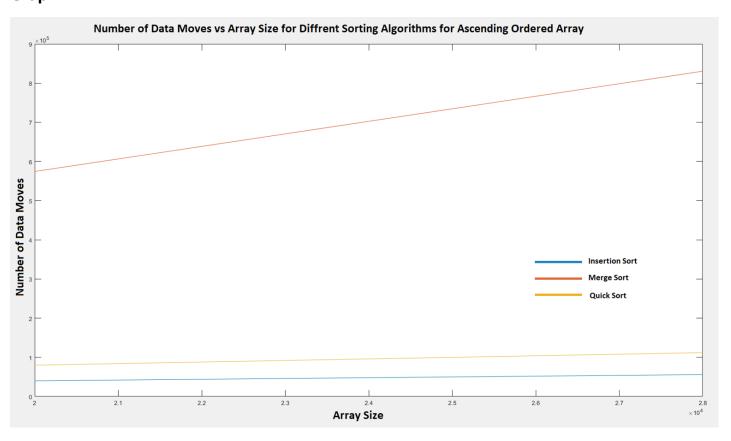
In the following graph, merge sort and insertion sort curves look as if they overlap so I also plotted the same graph with logarithmic y-axis.

Graph 10



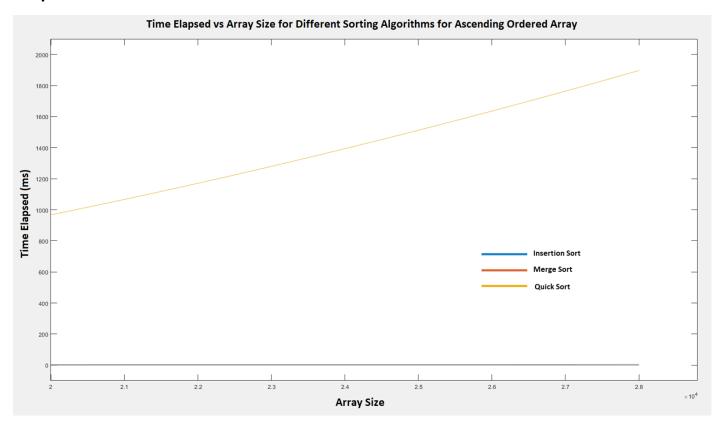
# Graph 11



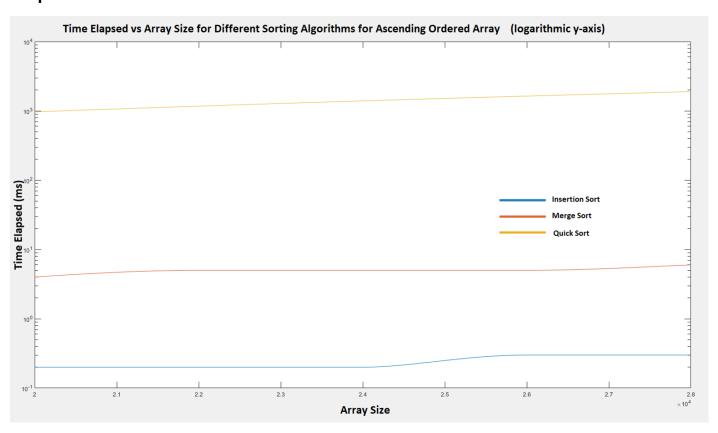


In the following graph, merge sort and insertion sort curves look as if they overlap so I also plotted the same graph with logarithmic y-axis.

Graph 13



Graph 14



# Sample Output

	(Randomly filled)				
	100232986 260952 359227	574464 655561	692.0 7.0 6.0		
	(Randomly filled)				
Insertion Sort Merge Sort Quick Sort	397895944 561980 716397		2751.0 15.0 12.0		
	(Randomly filled)				
Insertion Sort Merge Sort Quick Sort	900510601 877225 1088346		6228.0 24.0 18.0		
	(Randomly filled)				
Insertion Sort Merge Sort Quick Sort	1597421811 1203434 1580845	2617856	11060.0 32.0 26.0		
Array Size: 100000 (Randomly filled)					
Algorithm Insertion Sort Merge Sort Quick Sort	Key Comparisons 2494526740 1536068 1997379	2494626751 3337856	Time Elapsed 17268.0 41.0 32.0	(ms)	

Array Size: 20000	(Descending order)					
Insertion Sort Merge Sort Quick Sort	199990000 139216 199990000	574464 300079996	1384.0 4.0			
Array Size: 22000	(Descending order)					
_	241989000 154208	Data Moves 242032998 638464 363087996	5.0			
Array Size: 24000 (Descending order)						
_	287988000 170624	288035998	Time Elapsed (ms) 1991.0 6.0 2694.0			
Array Size: 26000	(Descending order)					
_	337987000	338038998 766464	Time Elapsed (ms) 2337.0 5.0 3162.0			
Array Size: 28000 (Descending order)						
Algorithm Insertion Sort Merge Sort Quick Sort	Key Comparisons 391986000 202512 391986000	Data Moves 392041998 830464 588111996	Time Elapsed (ms) 2713.0 6.0 3667.0			

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Array Size: 20000	(Ascending order)				
Merge Sort Quick Sort		39998 574464 79996	Time Elapsed (ms) 0.2 4.0 968.0		
Array Size: 22000	(Ascending order)				
Insertion Sort Merge Sort	21999 165024				
Array Size: 24000	(Ascending order)				
Insertion Sort Merge Sort	23999 180608				
Array Size: 26000	(Ascending order)				
Insertion Sort Merge Sort	25999	51998 766464	Time Elapsed (ms) 0.3 5.0 1636.0		
Array Size: 28000 (Ascending order)					
Algorithm Insertion Sort Merge Sort Quick Sort	Key Comparisons 27999 212720 391986000	Data Moves 55998 830464 111996	Time Elapsed (ms) 0.3 6.0 1897.0		

#### **Question 4**

It can be seen from Graph 1 and Graph 2 that insertion sort requires more key comparisons than quick sort and quick sort requires more key comparisons than merge sort for randomly filled arrays.

It can be seen from Graph 3 and Graph 4 that insertion sort requires more data moves than both quick and merge sort for randomly filled arrays. Quick sort requires more data moves than merge sort up to some array size which is ~40000 . For array sizes which are greater ~40000 merge sort requires more data moves than quick sort.

Theoretical analysis suggests that while insertion sort algorithm has O(n²) time complexity, merge sort and quick sort algorithms have O(nlogn) time complexity on average case. These theoretical foundings are parallel to the foundings of my test program according to Graph 5 and Graph 6. Even though merge sort and quick sort have the same time complexity, quick sort has a better performance in practice in average case. One of the reasons for this is that quick sort is an inplace algorithm which does not require extra space like merge sort (Getting extra memory consumes time). Another reason could be that quick sort algorithm has a better cache locality since it frequently uses data that are close to each other.

It can be seen from Graph 7 that insertion sort and quick sort requires equal amount of key comparisons which is larger than the amount of key comparisons required by merge sort algorithm for arrays which are ordered in descending order.

It can be seen from Graph 8 that quick sort requires more data moves than insertion sort and insertion sort requires more data moves than merge sort for arrays which are ordered in descending order.

It can be seen from Graph 9 that quick sort requires more time than insertion sort and insertion sort requires more time than merge sort. From theoretical analysis it is known that quick sort algorithm has the worst case performance when the array is sorted in ascending or descending order (since the pivot will always be the largest or smallest). Theoretical analysis suggests that insertion sort has the time complexity of  $O(n^2)$  on worst and average cases and O(n) on the best case (array already sorted). Since merge sort always runs in  $O(n\log n)$  complexity, it makes sense that insertion sort requires more time than merge sort when the array is sorted in descending order. So it could be said that theoretical analysis is parallel to experiment results.

It can be seen from Graph 10 and Graph 11 that quick sort requires more key comparisons than merge sort and merge sort requires more key comparisons than insertion sort when the array is ordered in ascending order.

It can be seen from Graph 12 that merge sort requires more data moves than quick sort and quick sort requires more data moves than insertion sort when the array is ordered in ascending order.

Graph 13 and Graph 14 suggest that quick sort requires more time than merge sort and merge sort requires more time than insertion sort when the array is ordered in ascending order. It makes sense that quick sort takes a lot more time than others since when the array is ordered in ascending order quick sort performs in its worst case. Insertion sort runs in O(n) time when the array is already sorted, it outperforms merge sort and quick sort. Since merge sort always runs in O(nlogn) time, my experimental ranking of the required times by each of the sorting algorithm is parallel to theoretical analysis.

For array size ~33000, I started getting stack overflow error for arrays ordered in ascending or descending order because of quick sort on my computer so I limited the array size for arrays ordered in ascending or descending order to 28000.

I was having integer overflow problems with the data type int for move and comparison count so I changed the data type for move and comparison count to unsigned long long.