Computation of matrix functions with fully automatic Schur-Parlett and Rational Krylov methods.

Roberto Zanotto

Matrix Functions

 $\operatorname{sqrt}(A), \exp(A), f(A) \dots$ where A is a square matrix. Useful generalization of scalar function $f: \mathbb{C} \to \mathbb{C}$

Examples of applications:

- > Stiff differential equations;
- > Nuclear magnetic resonance;
- > Control theory.

Mathematical definition of f(A) can be obtained through polynomial interpolation.

Matrix polynomials:

$$p(x) = x^3 + 4x - 7$$
, $p(A) = A^3 + 4A - 7I$

Interesting property: value of p(A) is fully determined by the values of p on the spectrum of A:

$$p(A) = q(A) \Leftrightarrow p^{(j)}(\lambda) = q^{(j)}(\lambda), \quad \forall \lambda \in \Lambda(A)$$

By extending this property to matrix functions, we get a definition of f(A):

$$f(A) \coloneqq p(A)$$

where p takes the same values of f on the spectrum of A, i.e.

$$p^{(j)}(\lambda) = f^{(j)}(\lambda), \quad \forall \lambda \in \Lambda(A)$$

Computation of Matrix Functions

Hermite interpolation with Vandermonde matrix is unstable and the resulting method is slow.

There are specialized methods (scaling and squaring for expm and logm, Björk-Hammarling for sqrtm).

We want an algorithm that is both:

- > Generic working for any function f and matrix A;
- > **Automatic** requiring no other information form the user (such as derivatives of *f* or other data structures).

Automatic Differentiation

Dual numbers can be used to compute derivatives in an automatic, efficient and stable manner:

$$x + y\varepsilon$$
 with $\varepsilon \neq 0$, $\varepsilon^2 = 0$
 $f(x + y\varepsilon) = f(x) + f'(x)y\varepsilon \Rightarrow f'(x) = \mathrm{Eps}(f(x + \varepsilon))$

We use the Julia language (a "new Matlab"), whose typesystem allows for pain-free use of dual numbers on userdefined functions.

Dual numbers can be generalized for higher-order derivatives (TaylorSeries.jl).

Automatic Differentiation - example

$$f(x) = (x - 3)^{8} \qquad \text{We can compute } f'(5) \text{ with}$$
 function $f(x)$ $\qquad f(5 + \varepsilon)$: $\qquad x -= 3 \qquad (5 + \varepsilon) - 3 = 2 + \varepsilon$ for $i = 1:3 \qquad (2 + \varepsilon)(2 + \varepsilon) = 4 + 4\varepsilon$ $\qquad x *= x \qquad (4 + 4\varepsilon)(4 + 4\varepsilon) = 16 + 32\varepsilon$ end $\qquad (16 + 32\varepsilon)(16 + 32\varepsilon) = 256 + 1024\varepsilon$ return $\qquad \text{The result } f(5 + \varepsilon) = 256 + 1024\varepsilon$ means end $\qquad f(5) = 256 \text{ and } f'(5) = 1024$

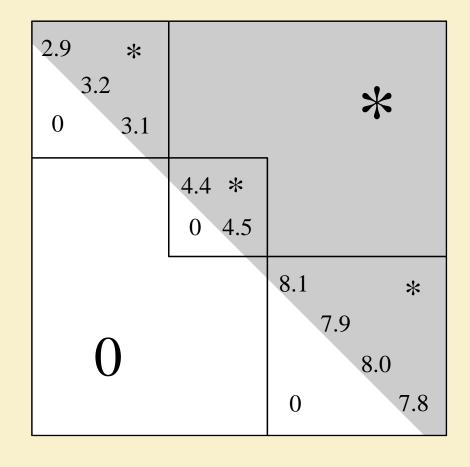
Schur-Parlett – dense matrix functions

Proposed by Higham and Davies in 2003.

$$A = QTQ^* \Rightarrow f(A) = Qf(T)Q^*$$

Need to compute F = f(T):

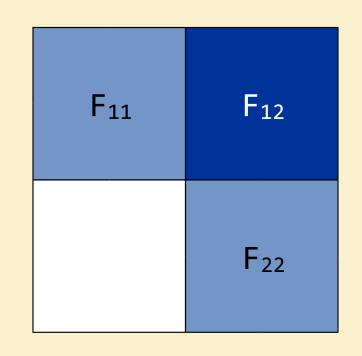
- 1. Group eigenvalues in blocks by proximity;
- 2. Compute $f(T_{ii})$ of diagonal blocks with Taylor;
- 3. Use the Parlett recurrence to reconstruct the upper part.



Schur-Parlett – Parlett recurrence

Parlett recurrence made recursive and cache-oblivious.

$F_{\mathtt{i}\mathtt{i}}$	Fik	$F_{\mathtt{i}\mathtt{j}}$	
		F _{kj}	
		F _{jj}	



$$T_{ii}F_{ij} - F_{ij}T_{jj} = F_{ii}T_{ij} - T_{ij}F_{jj}$$

$$+ \sum_{k=i+1}^{j-1} (F_{ik}T_{kj} - T_{ik}F_{kj})$$

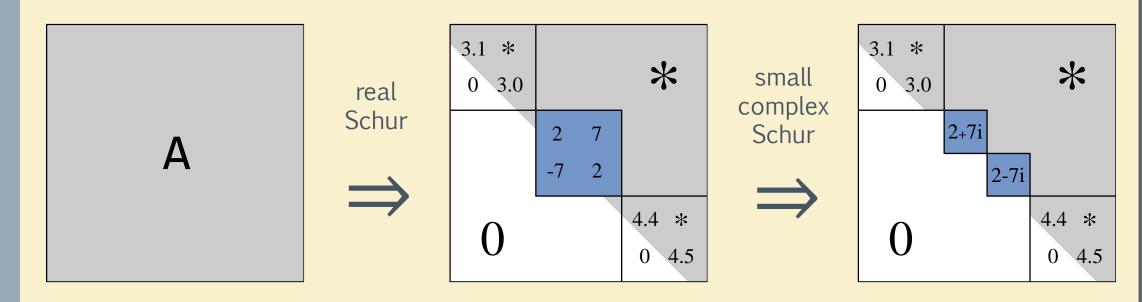
$$T_{11}F_{12} - F_{12}T_{22} = F_{11}T_{12} - T_{12}F_{22}$$

~3x speedup for n = 2500

Problem: conjugated eigenvalues with big imaginary part must go in different blocks, even with real *A*.

Original solution: do everything in complex arithmetic.

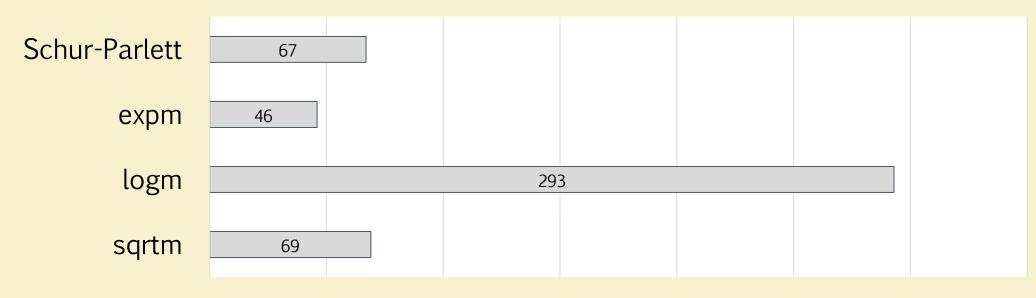
Our solution: complex Schur factorization can be "delayed" and done on small blocks:



Also allows for Parlett recurrence with real Sylvester equations, $\sim 2x$ speedup for n = 2500 on whole Schur-Parlett.

Schur-Parlett – performance results

The whole procedure typically spends $\sim 2/3$ of the time doing A's Schur decomposition (varies depending on the eigenvalues' distribution).



■ Execution time (s) on randn(2500, 2500)

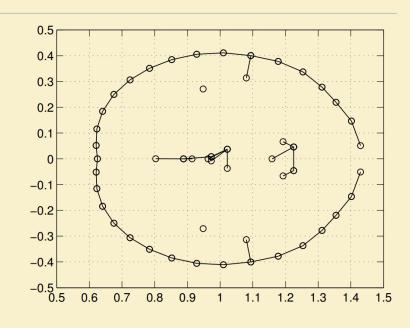
Schur-Parlett – numerical accuracy

	Specialized method	Schur-Parlett		
ехр	1.0e-14	1.5e-13		
log	4.0e-14	4.1e-14		
sqrt	4.3e-14	4.5e-14		

Relative error versus Matlab's Symbolic Toolbox (on 50x50 random matrix).

Boils down to how well the eigenvalues can be clustered.

Can behave badly (7e-4) with "snake" eigenvalues (as shown in original paper, experiment 4)



Code comparison for the matrix cube root $A^{1/3}$

Matlab

```
function r = cuberoot(x, k)
     c = 1.0;
     for i = 0:k-1
           c = c*(1/3 - i);
     end
     r = c*x^{(1/3 - k)};
end
funm(A, @cuberoot)
```

Our Julia implementation

```
schurparlett((x) \rightarrow x^{(1/3)}, A)
```

Rational Krylov - sparse matrix functions

(as described by Güttel in 2013)

Goal: compute f(A)b for a sparse A.

Approximate f(A)b with $r_m(A)b$ where $r_m = p_{m-1}/q_{m-1}$

Denominator is factored as $q_{m-1}(z) = \prod_{j=1}^{m-1} (1 - z/\xi_j)$ with poles ξ_i provided by the user.

Rational Krylov space is defined as:

$$Q_m(A,b) = q_{m-1}(A)^{-1} span\{b, Ab, ..., A^{m-1}b\}$$

Obtained with Ruhe's rational Arnoldi algorithm:

$$v_1 = b/||b||$$
, $v_{j+1} = orthonormalize((I - A/\xi_j)^{-1}Av_j)$

 $r_m(A)b$ is computed by projecting A into the Krylov space:

$$A_m = V^*AV$$
, $r_m(A)b = Vf(A_m)V^*b$

Rational Krylov – approximation accuracy

Depends on two factors:

- How well f can be approximated on A's spectrum by a rational function, hopefully with low degree (ill-posed problem);
- \rightarrow How well we choose the poles ξ_i .

We use the AAA algorithm for rational approximation to find good poles automatically.

AAA algorithm for rational approximation

Proposed by Nakatsukasa, Sète, Trefethen in 2017.

Input: function samples (real or complex).

Output: rational barycentric function of type (m-1, m-1):

$$r(z) = \frac{n(z)}{d(z)} = \sum_{j=1}^{m} \frac{w_j f_j}{z - z_j} / \sum_{j=1}^{m} \frac{w_j}{z - z_j}$$

$$r(z_j) = \infty/\infty$$
 but $\lim_{z \to z_j} r(z) = f_j$

Support/interpolation points (z_j) are chosen incrementally from samples in a greedy way, to avoid instabilities:

next z_i is chosen where $|f(z_i) - r(z_i)|$ is maximized.

After a new support point is found, weights are recomputed to minimize the approximation error:

$$f(z) \approx \frac{n(z)}{d(z)} \to minimize||f(z)d(z) - n(z)||, \qquad z \in Z^{(m)}$$

Is a least squares problem solvable with SVD:

$$minimize||Aw||, ||w|| = 1$$

When approximation error is small, we are done.

Poles can be then retrieved by solving a generalized eigenvalue problem, with accuracy up to machine precision.

Rational Krylov + AAA

- 1. Sample f on the 0-centered disk of radius norm(A);
- 2. Find the poles of *f* with AAA;
- 3. Use the poles for the rational Arnoldi approximation.

Computational costs:

Rational Krylov: m times sparse linear system + orthogonalization: O(m(L + mN))

AAA: m times SVD: $O(m(m^2M))$

Rational Krylov + AAA, Experiment 1

Matrices are from the SuiteSparse Matrix Collection. exp(A)b with Krylov is compared with dense expm:

Problem type	name	cond(A)	cond (exp, A)	# poles	error	# poles	error
2D/3D problem	jagmesh3	1168	7.0e0	9	1.5e-6	15	6.0e-14
Fluid dynamics	sherman4	2178	1.2e2	11	6.6e-6	19	6.7e-10
Structural problem	can_1072	2.0e34	3.4e1	11	7.3e-6	23	1.7e-14
Directed graph	SmaGri	Inf	7.8e1	9	2.6e-7	21	7.3e-13

Rational Krylov + AAA, Experiment 2

sqrt(A)b (positive definite A) against dense sqrtm:

Problem type	name	min eig(A)	norm(A)	# poles	error	# poles	error
Circuit simulatio	rajat19	1.7e-1	3.9e1	20	5.8e-6	45	3.7e-14
Structural problem	nos3	1.8e-2	7.7e2	50	8.8e-4	100	5.2e-9
Power network	685_bus	6.2e-2	2.6e4	50	1.5e-3	95	8.0e-5
Electro magnetic	mhd 1280b	1.5e-11	8.0e1	50	2.5e-2	50	8.8e-3

Custom sampling for AAA, more dense near 0, brings error from 10^{-3} to 10^{-12} (matrix "685_bus").

Future Work

Different algorithms could be used for the evaluation of the diagonal blocks in Schur-Parlett (replacing Taylor).

An AAA solution could solve some of Schur-Parlett's instability issues and lift the dependency on automatic differentiation, bringing the automatic algorithm to other languages such as Matlab and C (more research needed).

Try it out at: https://github.com/robzan8/MatFun.jl

Questions?