

Title of Your Project

Compiler Construction '17 Final Report

First1 Last1 First2 Last2
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1. Introduction

We chose to do the Tail-call optimization lab. However general Tail-call elimination would be a complex task in JVM so we intended to start out by only implementing the sub problem of Tail-recursion optimization.

This allows us to focus solve the problem without any changes to the existing compiler stages. We believe that all we need to do is to introduce a new compiler stage for AST rewriting, and implement tail recursion elimination as an AST transformation from a tail recursive method to a method with a while loop.

This stage will have to be executed after the Type checking phase, as having the method symbols attached for on method call expressions would be helpful.

2. Examples

Benefits of tail recursion optimizations is the most obvious on method executions that would otherwise run out of stack space. Consider for example the following program.

It recursively increments *i* until *i* overflows to 0 and returns the result.

```
class OverflowPlease {  
  def makeZero(i: Int): Int = {  
    if (i == 0) {  
      i  
    } else {  
      this.makeZero(i + 1)  
    }  
  }  
}  
  
object Main extends App {  
  println(new OverflowPlease().makeZero(0))  
}
```

each recursion creates a new stack frame, so a stack that fits over 4 million frames would be required to execute

this program. But with tail recursion optimization we expect this program to successfully overflow to 0 only using a couple of stack frames.

This should also make the programs that do not run out of stack space faster since popping and pushing stack frames is a relatively slow procedure.

here is a similar example from the lab description that recursively gets the last element of a linked list

```
class Helper {  
  def last(a: List): List = {  
    if (a.hasNext()) {  
      this.last(a.next())  
    } else {  
      a  
    }  
  }  
}
```

2.1 Recursion on trees

Recursion on tree structures is interesting because in binary trees we need to recurse on the left and right subtrees, but only the one of the recursive calls may be in the tail position. Still removing one of the recursive calls is better than none, especially if the tree tends to be unbalanced.

this example has a function that tail recursively sums the contents of the tree.

and another one that sets all the contents of the tree to 0.

```
class TreeHelper {  
  def countNodes(a: Tree, acc: Int): List = {  
    val r: Tree = NULL;  
    val l: Tree = NULL;  
    val sum: Int = acc;  
    if (a == NULL) {  
      sum  
    } else {  

```

```

    r=a.getRight();
    l=a.getLeft();
    sum = this.countNodes(r, acc) + sum + 1;
    this.countNodes(l, sum)
  }
}

def cleanseTree(a: Tree): List = {
  val r:Tree = NULL;
  val l:Tree = NULL;
  if(!(a==NULL)){
    a.setValue(0);
    r=a.getRight();
    l=a.getLeft();
    this.cleanseTree(r);
    this.cleanseTree(l)
  }
}

```

2.2 corecursive tail calls

Here is toy example of the more general problem that we do not intend to handle. Here two methods `isEven(Int)` and `isOdd(Int)` use tail calls to each other to determine if the input method is odd or even.

```

class numbers {
  def isEven(i:Int):Boolean = {
    if(i==0) True
    else isOdd(i-1)
  }

  def isOdd(i:Int):Boolean = {
    if(i==1) True
    else isEven(i-1)
  }
}

```

3. Implementation

This is a very important section, you explain to us how you made it work.

3.1 Theoretical Background

If you are using theoretical concepts, explain them first in this subsection. Even if they come from the course (e.g., lattices), try to explain the essential points *in your own words*. Cite any reference work you used like this [Appel 2002]. This should convince us that you know the theory behind what you coded.

3.2 Implementation Details

Describe all non-obvious tricks you used. Tell us what you thought was hard and why. If it took you time to figure out the solution to a problem, it probably means it wasn't easy and you should definitely describe the solution in detail here. If you used what you think is a cool algorithm for some problem, tell us. Do not however spend time describing trivial things (we know what a tree traversal is, for instance).

After reading this section, we should be convinced that you knew what you were doing when you wrote your extension, and that you put some extra consideration for the harder parts.

4. Possible Extensions

Handling corecursive tail calls like in the even odd example in Section 2.2 would be the obvious next step.

The benefit is that we would get improved performance and avoid some stack overflows in some corecursive programs.

But this extension is a lot more challenging. The typical approach for call recursion elimination of simply jumping to the tail called method instead of popping and pushing stack frames would not be possible since every method in JVM has its own private address space.

One possible approach to eliminate corecursive tail calls that would build well on work in this project would be to implement a AST transformation that rewrites corecursive methods into a single big method that can simulate all of the methods forming the corecursion. After that our tail recursion elimination AST transformation can be applied to the AST generated by the method merger.

References

A. W. Appel. *Modern Compiler Implementation in Java*. Cambridge University Press, 2nd edition, 2002.