

CS112 Primality Tester

N:

K:

Fermat Result: 701 **is prime** with probability 0.9687500000000000

MR Result: 701 **is prime** with probability 0.9990234375000000

CS112 Primality Tester

N:

K:

Fermat Result: 1105 **is prime** with probability 0.8750000000000000

MR Result: 1105 **is not prime**

```
def mod_exp(x, y, N):
    # You will need to implement this function and change the return value.

    if y == 0:
        return 1

    # Recursively call mod_exp function with y divided by 2 rounding down.
    z = mod_exp(x, (y//2), N)

    if y % 2 == 0: # when y is even
        return (z**2) % N
    else:
        return (x*(z**2)) % N

def fprobability(k):
    # You will need to implement this function and change the return value.
    # Fermat's test returns YES when N is not prime: p(fail) <= (1/2^k), so
    p(success) = 1 - p(fail)
    return 1 - 1/(2**k)

def mprobability(k):
    # You will need to implement this function and change the return value.
    # Miller-Rabin test passes when N is not prime : p(fail) <= (1/4^k), so
    p(success) = 1 - p(fail)
    return 1 - 1/(4**k)

def run_fermat(N,k):
    # You will need to implement this function and change the return value, which
    # should be
    # either 'prime' or 'composite'.
    #
    # To generate random values for a, you will most likely want to use
    # random.randint(low,hi) which gives a random integer between low and
    # hi, inclusive.

    # If N is 2, it is prime
    if N == 2:
        return 'prime'

    for i in range(1, k+1):
        # Create random number to test
```

```

        random_number = random.randint(1, N - 1)
        if mod_exp(random_number, N-1, N) != 1:
            return 'composite'
    return 'prime'

```

```
def run_miller_rabin(N,k):
```

```

    # You will need to implement this function and change the return value, which
    should be

```

```

    # either 'prime' or 'composite'.
    #

```

```

    # To generate random values for a, you will most likely want to use
    # random.randint(low,hi) which gives a random integer between low and
    # hi, inclusive.

```

```

    # Return 'prime' when N is 2

```

```

    if N == 2:
        return 'prime'

```

```

    # K tests

```

```

    for i in range(1, k+1):
        # Pick a random number 1 <= a < N
        random_number = random.randint(1, N-1)

```

```

        # Initial test
        if mod_exp(random_number, N - 1, N) != 1:
            return 'composite'

```

```

        # Initial test passed, compute the sequence and find the first value that
        is not equal to 1

```

```

        y = (N - 1) / 2
        compute_sequence = True
        while compute_sequence:
            value = mod_exp(random_number, y, N)
            if value != 1:
                if (value + 1) != N: # value is not -1
                    return 'composite'
                else:
                    break # Passed the test, move on

```

```

        if y % 2 != 0:
            compute_sequence = False # Can't square root anymore
        else:
            y = y / 2

```

```

    return 'prime'

```

3. For functions that calculate probabilities which are `fprobability()` and `mprobability()`, I would say that they have time complexity of $O(1)$ which is constant time. Since the number of bits in k is very tiny compared to the number of bits n in N , we can say that they have time complexity of $O(1)$ respect to n . Also, they just calculate the probabilities and return the resulting value, so both have space complexity of $O(1)$.

For `mod_exp()`, the modular operation has time complexity of $O(n^2)$, and multiplication and division also take $O(n^2)$ time, and the function is calling itself n times recursively. Therefore, the time complexity of this function would be $O(n^3)$. For space complexity, z is getting return value and it needs $O(n)$ bits. Since the function recursively calls itself N times, the space complexity would be $O(n^2)$.

For `run_fermat()` function, it calls `mod_exp()` function k times unless it breaks out of the loop in the middle. So, the time complexity would be $O(kn^3)$. We store random number in the loop, but continuously call `mod_exp()`, so it would have the space complexity of $O(n^2)$.

For `run_miller_rabin()` function, it calls `mod_exp()` in the for loop but use another loop inside to continue to square root of the number. Therefore, the time complexity for this function would be $O(n^4)$. Since we keep calling `mod_exp()`, the space complexity would also be $O(n^2)$.

4. `fprobability()` calculates the probability of success from running the `run_fermat()`. The probability that the testing primality function would say yes when N is not prime is less than or equal to $(1/2)^k$. Therefore, the probability that the algorithm would give the correct result would be $1 - (1/2)^k$. Similarly, the probability that `run_miller_rabin()` would give wrong results is less than or equal to $(1/4)^k$. Therefore, $p(\text{success})$ would be $1 - (1/4)^k$.