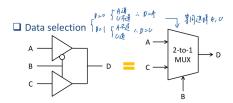
# Digital System Design and Lab: HW3

Lo Chun, Chou R13922136

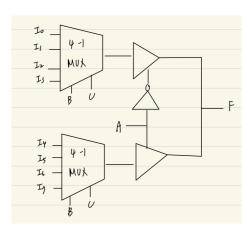
April 25, 2025

# 1

By lecture slide LEC-09 p.10-11, we knew that using two three-state buffers with one inverter could do data selection, and is equivalent to a 2-to-1 MUX:

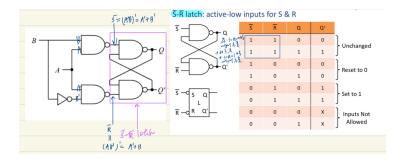


So, we can use the two 4-to-1 MUXs and this setting to implement the 8-to-1 MUX as follows:



# $\mathbf{2}$

First, observe that part of the given circuit is a  $\bar{S} - \bar{R}$  latch:



with the inputs:

$$\bar{S} = (AB)' = A' + B'$$

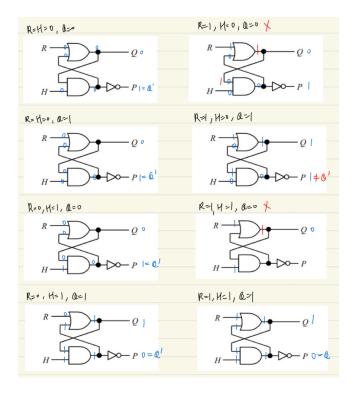
$$\bar{R} = (AB')' = A' + B$$

Thus, we can form the truth table of this latch by checking the values of  $\bar{S},\,\bar{R}$  and Q in the table above:

A	В	$\bar{S}$	$\bar{R}$	Q	$Q^+$
0	0	1	1	0	0
0	0	1	1	1	1
0	1	1	1	0	0
0	1	1	1	1	1
1	0	1	0	0	0
1	0	1	0	1	0
1	1	0	1	0	1
1	1	0	1	1	1

3

The following is the different cases of the latch:



(1)

From the above cases, we can see that when R=1 and H=0,  $P=1\neq Q'=0.$  Therefore, we should not let:

$$R=1$$
 and  $H=0$ 

(2)

The next-state table is shown below:

R	Н	Q	Q+
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	X
1	0	1	1
1	1	0	X
1	1	1	1

And we can construct the K-map as follows:

HOR	D	I	
0 0	0	X	
o	0	J	
l	1		
ا ۵	0	*	

Which would give us the characteristic equation:

$$Q^+ = R + H \cdot Q$$

### 4

(1)

```
lookahead adder, gate level modeling modify the input/output of module
wire 60, 61, 62, P0, P1, P2; // 6 is the generate, P is the propagate, both 3 bits // Note: cannot use: assign 60 = A[0] & B[0]; since we are required to use gate level moduling // Check LAB-01.pdf p.33 for the formula
       AND g_0(G0, A[0], B[0]);
AND g_1(G1, A[1], B[1]);
AND g_2(G2, A[2], B[2]);
       OR or_0(P0, A[0], B[0]);
OR or_1(P1, A[1], B[1]);
OR or_2(P2, A[2], B[2]);
       // aim: Then calculate the carry f
wire Cin_0, Cin_1, Cin_2, Cout_2;
       // explain: the following equations are using notations in LAB-01.pdf p.33
       // SUBCOM:
wire p0c0;
AND int11(p0c0, P0, Cin_0);
OR c_1(Cin_1, G0, p0c0);
       // subaim: C_2 = G_1 + P_1 & C_1
wire g6p1, c6p6p1;
AND int22(c6p6p1, 60, P1);
AND int22(c6p6p1, p8c0, P1);
// note: we add 0 since we only need to OR 3 terms, but we need to choose from OR and OR4
CR4 c_2(Cin_2, G1, g6p1, c6p6p1, 1'b0);
        // subaim: C_3 = G_2 + P_2 & C_2
       // suboth: L3 = 0.2 + P.2 x (2)
wire glp2, g0p1p2, c0p0p1p2;
AND int31(glp2, G1, P2);
AND int32(g0p1p2, g0p1, P2);
AND int33(c0p0p1p2, c0p0p1, P2);
OR4 c_3(Cout_2, G2, g1p2, g0p1p2, c0p0p1p2);
         // aim: Calculate the result of each full adder
// note: the carries are stored in Di (dummy), since we've already precomputed them, and these Di would not be used later
       // note: the corries are stored in Di
wire D0, D1, D2;
FA fa_6(D0, S[0], A[0], B[0], Cin_0);
FA fa_1(D1, S[1], A[1], B[1], Cin_1);
FA fa_2(D2, S[2], A[2], B[2], Cin_2);
```

**(2)** 

# 5

I use Surfer instead of GTKWave to present the waveform:



# 6

We can find the maximum delay and one of the transition from the attached terminal output screenshots.

(1)

The maximum delay is 23 ticks on transition 000+000+0 --> 000+111+1 lab1.v:70: \$finish called at 3276800000 (1ps) 3276800 / 000 / 111 / 1 / 1000 / 1111

**(2)** 

The maximum delay is 20 ticks on transition 000+000+0 --> 000+011+1 lab1.v:70: \$finish called at 3276800000 (1ps) 3276800 / 000 / 011 / 1 / 0100 / z111

7