

# Digital Systems Design and Laboratory

## [ 1. Number Systems and Conversion ]

Chung-Wei Lin

[cwlin@csie.ntu.edu.tw](mailto:cwlin@csie.ntu.edu.tw)

CSIE Department

National Taiwan University

# Outline

- ❑ **Digital Systems and Switching Circuits**

- ❑ Number Systems and Conversion

- ❑ Binary Arithmetic

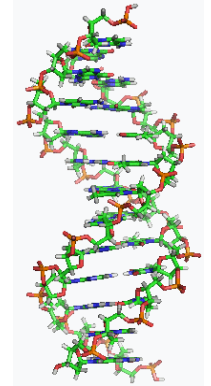
- ❑ Representation of Negative Numbers

- ❑ Binary Codes

# Historical Digital Systems

能传达的信息来代表意义:

- ☐ Abacus
- ☐ Braille
- ☐ DNA
- ☐ Flag semaphore
- ☐ International maritime signal flags
- ☐ Morse code



Source: Wikipedia



INTERNATIONAL FLAGS AND PENNANTS					
ALPHABET FLAGS			NUMERAL PENNANTS		
Alfa		Kilo		Uniform	
Bravo		Lima		Victor	
Charlie		Mike		Whiskey	
Delta		November		X-ray	
Echo		Oscar		Yankee	
Foxtrot		Papa		Zulu	
Golf		Quebec		SUBSTITUTES	
Hotel		Romeo		1st Substitute	
India		Sierra		2nd Substitute	
Juliett		Tango		3rd Substitute	
				CODE (Answering Pennant or Second Point)	

## International Morse Code

1. A dash is equal to three dots.
2. The space between parts of the same letter is equal to one dot.
3. The space between two letters is equal to three dots.
4. The space between two words is equal to five dots.

A		U	
B		V	
C		W	
D		X	
E		Y	
F		Z	
G			
H			
I			
J			
K		1	
L		2	
M		3	
N		4	
O		5	
P		6	
Q		7	
R		8	
S		9	
T		0	

# Digital vs. Analog

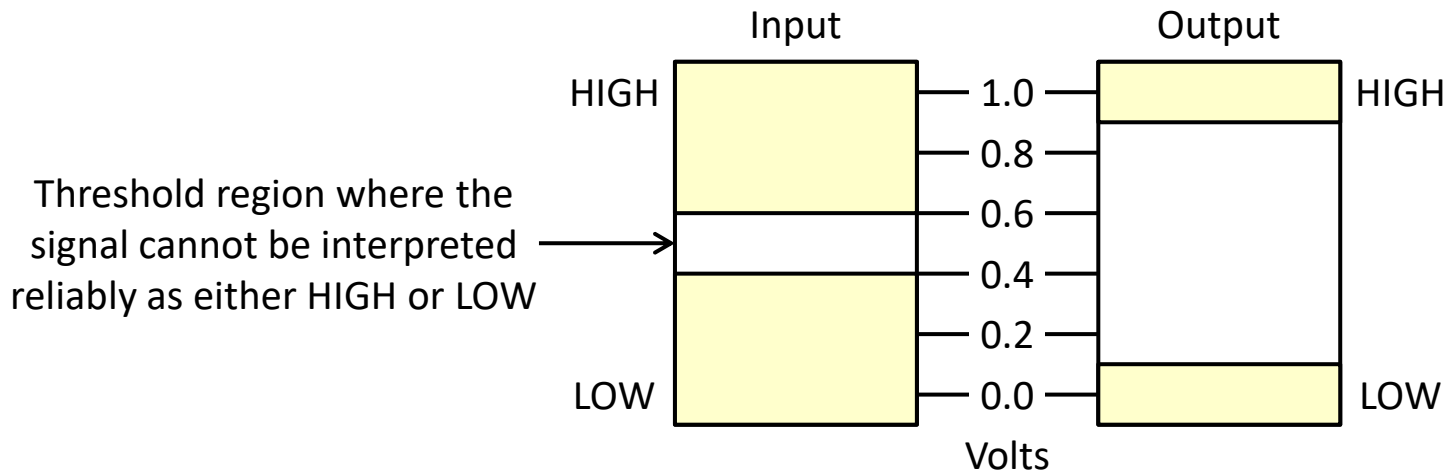
除了0/1以外的值  
也可能有意义。

## ❑ The physical quantities or signals in

➤ A digital system assumes only discrete values

- Example: 0V and +1V
- Greater accuracy and reliability (why?)

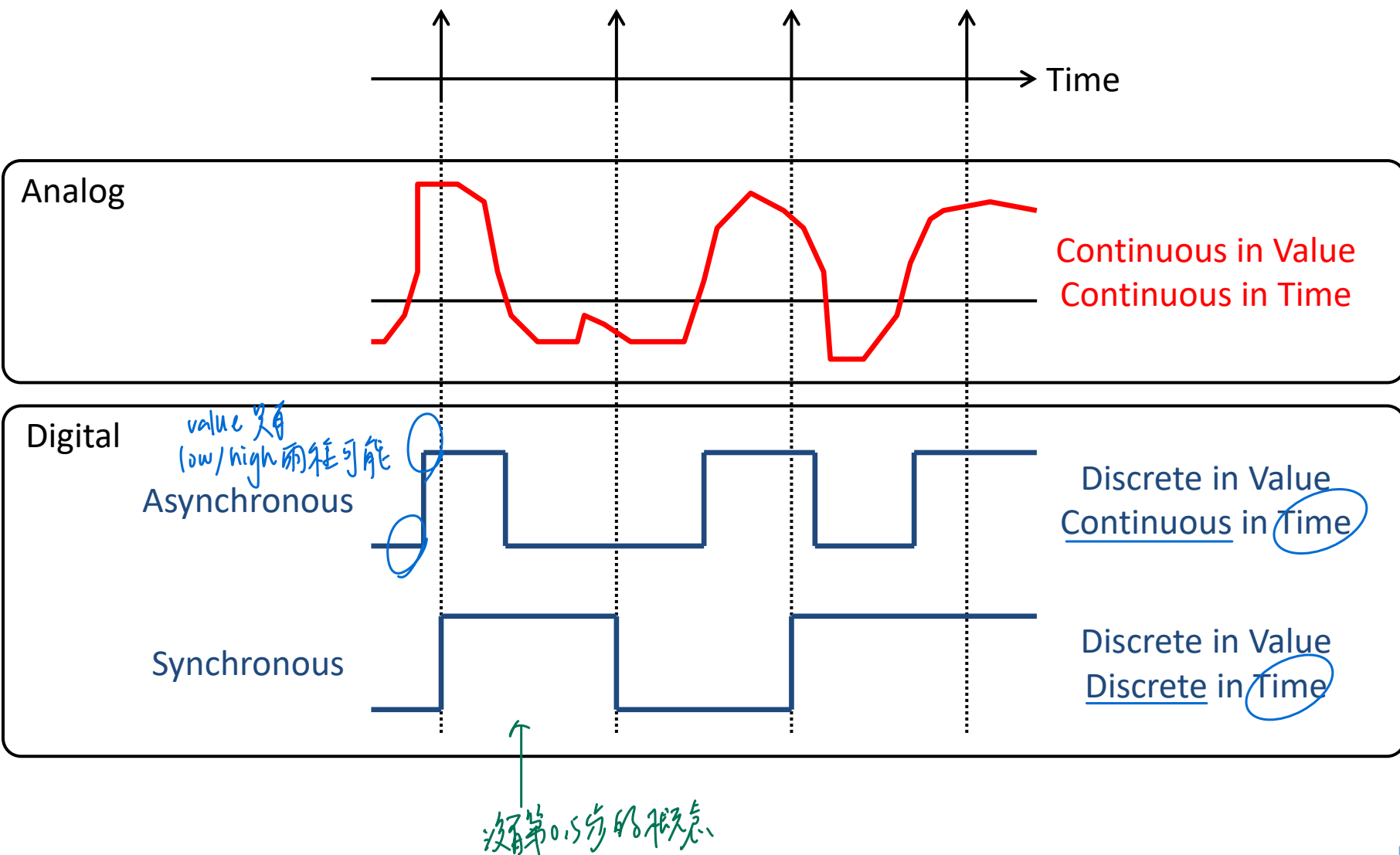
↓  
值. 时间  
continuous



➤ An analog system varies continuously over a specified range

- Example: any value between 0V to +1V

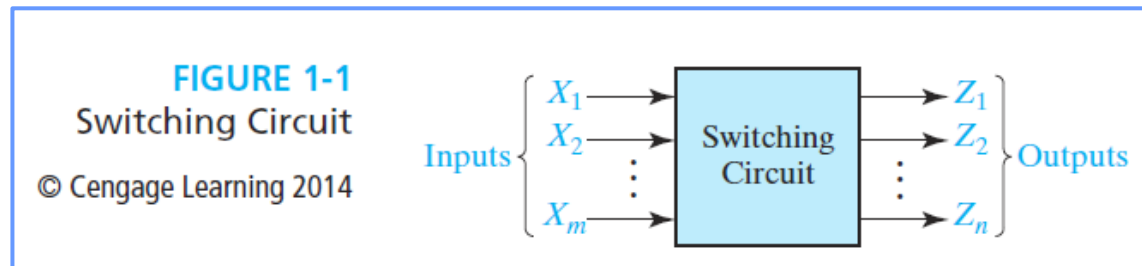
# Signal Examples over Time



# Digital Systems and Switching Circuits

## □ Subsystems of a digital system take the form of a switching circuit which has discrete inputs and outputs

- Switching devices are generally **two-state** devices
  - i.e., output can assume only **two** different discrete values
- It is natural to use **binary** numbers internally in digital systems



## □ Two types of switching circuits

- **Combinational circuits:** outputs depend only on present inputs
  - Memoryless
- **Sequential circuits:** outputs depend on both present and past inputs
  - In general, sequential circuits = combinational circuits + memory

有記憶功能 (M1後)

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- ❑ Digital Systems and Switching Circuits
- ❑ **Number Systems and Conversion**
- ❑ Binary Arithmetic
- ❑ Representation of Negative Numbers
- ❑ Binary Codes

# Number Systems (1/2)

□ Positional notation: each digit is multiplied by an appropriate power of base depending on its position in the number

➤ The point separates the positive and negative powers of base

- Example: decimal (base 10) numbers

$$-953.78_{10} = 9 \times 10^2 + 5 \times 10^1 + 3 \times 10^0 + 7 \times 10^{-1} + 8 \times 10^{-2}$$

➤ A positive number N with base R (positive integer,  $R > 1$ ):

$$\begin{aligned} N &= (a_4 a_3 a_2 a_1 a_0 . a_{-1} a_{-2} a_{-3})_R \\ &= a_4 R^4 + a_3 R^3 + a_2 R^2 + a_1 R^1 + a_0 R^0 + a_{-1} R^{-1} + a_{-2} R^{-2} + a_{-3} R^{-3} \end{aligned}$$

- Base is also called radix
- Base is indicated as subscript

➤ Why do people use the decimal number system?



# Number Systems (2/2)

## □ Examples

### ➤ Decimal (base 10) numbers

- $953.78_{10} = 9 \times 10^2 + 5 \times 10^1 + 3 \times 10^0 + 7 \times 10^{-1} + 8 \times 10^{-2}$

### ➤ Binary (base 2) numbers

- $1011.11_2 = 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 + 1 \times 2^{-1} + 1 \times 2^{-2}$   
 $= 11.75_{10}$

### ➤ Octal (base 8) numbers

- $147.3_8 = 1 \times 8^2 + 4 \times 8^1 + 7 \times 8^0 + 3 \times 8^{-1}$   
 $= 103.375_{10}$

### ➤ Hexadecimal (base 16) numbers

- Digits: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F
- $A2F_{16} = 10 \times 16^2 + 2 \times 16^1 + 15 \times 16^0$   
 $= 2607_{10}$

# Conversion of Decimal Integer

## □ Convert a decimal integer to base R using division

$$\text{➤ } N = (a_n a_{n-1} \dots a_2 a_1 a_0)_R = a_n R^n + a_{n-1} R^{n-1} + \dots + a_3 R^3 + a_2 R^2 + a_1 R^1 + a_0$$

$$\text{➤ } N/R = a_n R^{n-1} + a_{n-1} R^{n-2} + \dots + a_3 R^2 + a_2 R^1 + a_1 = Q_1 \quad \text{rem.} = a_0$$

$$\text{➤ } Q_1/R = a_n R^{n-2} + a_{n-1} R^{n-3} + \dots + a_3 R^1 + a_2 = Q_2 \quad \text{rem.} = a_1$$

$$\text{➤ } Q_2/R = a_n R^{n-3} + a_{n-1} R^{n-4} + \dots + a_3 = Q_3 \quad \text{rem.} = a_2$$

➤ Continue until ...

$$\text{➤ } Q_i/R = 0 \quad \text{rem.} = a_n$$

➤ Example: convert **53<sub>10</sub>** to binary

$$2 \overline{) 53}$$

$$2 \overline{) 26} \dots \text{remainder} = 1 = a_0 \quad (\text{LSB})$$

$$2 \overline{) 13} \dots \text{remainder} = 0 = a_1$$

$$2 \overline{) 6} \dots \text{remainder} = 1 = a_2$$

$$2 \overline{) 3} \dots \text{remainder} = 0 = a_3$$

$$2 \overline{) 1} \dots \text{remainder} = 1 = a_4$$

$$0 \dots \text{remainder} = 1 = a_5 \quad (\text{MSB})$$

$$\mathbf{53_{10} = 110101_2}$$

一定会有1 否则加不出53



小数点前那位.

# Conversion of Decimal Fraction (1/2)

## □ Convert a decimal fraction to base R using multiplication

$$\text{➤ } F = (.a_{-1}a_{-2}a_{-3}\dots a_{-m})_R = a_{-1}R^{-1} + a_{-2}R^{-2} + a_{-3}R^{-3} + \dots + a_{-m}R^{-m}$$

$$\text{➤ } FR = a_{-1} + a_{-2}R^{-1} + a_{-3}R^{-2} + \dots + a_{-m}R^{-m+1} = a_{-1} + F_1$$

$$\text{➤ } F_1R = a_{-2} + a_{-3}R^{-1} + \dots + a_{-m}R^{-m+2} = a_{-2} + F_2$$

$$\text{➤ } F_2R = a_{-3} + \dots + a_{-m}R^{-m+3} = a_{-3} + F_3$$

➤ Continue until  $F_i = 0$  or ... (next slide)

➤ Example: convert **.375<sub>10</sub>** to binary

$$\begin{array}{r}
 .375 \\
 \times 2 \\
 \hline
 (0) .750 \quad a_{-1} = 0 \quad (\text{MSB}) \\
 \times 2 \\
 \hline
 (1) .500 \quad a_{-2} = 1 \\
 \times 2 \\
 \hline
 (1) .000 \quad a_{-3} = 1 \quad (\text{LSB})
 \end{array}$$

$\frac{1}{4} + \frac{1}{8} = 0.375$

$\frac{1}{4} =$   
 $\frac{1}{8} =$

從小數點往後排.

**.375<sub>10</sub> = .011<sub>2</sub>**

# Conversion of Decimal Fraction (2/2)

## □ Sometimes, the result is a repeating fraction

➤ Example: convert  $.7_{10}$  to binary

急用  $\frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots$   
加出 0.7

$$\begin{array}{r} .7 \\ \times 2 \\ \hline (1) .4 \\ \times 2 \\ \hline (0) .8 \\ \times 2 \\ \hline (1) .6 \\ \times 2 \\ \hline (1) .2 \\ \times 2 \\ \hline (0) .4 \end{array} \leftarrow \begin{array}{l} \text{Repeating} \end{array}$$

$$.7_{10} = .1 \underline{0110} \underline{0110} \dots_2$$

无限多项等比级数.

# Conversion between Two Bases (1/2)

## □ Convert between two bases $R_1$ and $R_2$ other than decimal

➤ Base  $R_1 \rightarrow$  base 10  $\rightarrow$  base  $R_2$

➤ Example: convert **231.3**<sub>4</sub> to base 7

$\bullet \quad 231.3_4 = 2 \times 4^2 + 3 \times 4^1 + 1 \times 4^0 + 3 \times 4^{-1} = 45.75_{10}$

$\begin{array}{c} \uparrow \uparrow \uparrow \uparrow \\ 4^2 \ 4^1 \ 4^0 \ 4^{-1} \end{array}$

$$\begin{array}{r} 7 \overline{) 45} \\ 7 \overline{) 6} \end{array} \dots \text{remainder} = 3$$

$$\begin{array}{r} 0 \dots \text{remainder} = 6 \end{array}$$

$45_{10} = 63_7$

$$\begin{array}{r} \phantom{.}75 \\ \times \phantom{.}7 \\ \hline (5) .25 \\ \times \phantom{.}7 \\ \hline (1) .75 \end{array}$$

Repeating

$.75_{10} = .\underline{51}_7$

$\bullet \quad 45.75_{10} = \mathbf{63.51}_7$

# Conversion between Two Bases (2/2)

## ❑ Convert between binary and octal/hexadecimal by inspection

- Start at the binary point
- Divide bits into groups of three/four
  - **Add 0's if necessary**
- Replace each group by an octal/hexadecimal digit

## ❑ Binary to octal *每个 digit 框起来 (3 digit 可表 0~7)*

*二转八* ➤  $1001101.010111_2 = \underline{001} \ \underline{001} \ \underline{101} . \underline{010} \ \underline{111}_2$   
 $= 115.27_8$

## ❑ Binary to hexadecimal *每个 digit 框起来*

*二转16* ➤  $\overset{\text{不够补0}}{0} \underline{1001101} . \overset{\text{不够补0}}{010111}_2 = \underline{0100} \ \underline{1101} . \underline{0101} \ \underline{1100}_2$   
 $= 4D.5C_{16}$   
*↑ 从小数点往左框*      *↑ 从小数点往右框*

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- ❑ Digital Systems and Switching Circuits
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- ❑ **Binary Arithmetic**
- ❑ Representation of Negative Numbers
- ❑ Binary Codes

# Addition

## □ Addition table

➤  $0 + 0 = 0$

➤  $0 + 1 = 1$

➤  $1 + 0 = 1$

➤  $1 + 1 = 0$  (and carry 1 to the next column)

## □ Example: add $13_{10}$ and $11_{10}$ in binary

$$\begin{array}{r} 13_{10} = \quad 1101 \\ 11_{10} = + \quad 1011 \\ \hline 11000 = 24_{10} \end{array}$$

1111 ← Carries



# Subtraction

## Subtraction table

- $0 - 0 = 0$
- $1 - 0 = 1$
- $1 - 1 = 0$
- $0 - 1 = 1$  (and borrow 1 from the next column)
  - Borrow 1 from the next column = subtract 1 at the next column and add 2 at the current column

❑ Example: subtract  $19_{10}$  and  $29_{10}$  in binary

$$\begin{array}{r} 29_{10} = 11101 \\ 19_{10} = - \underline{10011} \\ \hline 01010 = 10_{10} \end{array}$$

# Multiplication

## ❑ Multiplication table

➤  $0 \times 0 = 0$

➤  $0 \times 1 = 0$

➤  $1 \times 0 = 0$

➤  $1 \times 1 = 1$

## ❑ Example: multiply $13_{10}$ and $11_{10}$ in binary

$$\begin{array}{r} 13_{10} = \quad \quad 1101 \\ 11_{10} = \times \quad \quad 1011 \\ \hline \quad \quad \quad 1101 \\ \quad \quad \quad 1101 \\ \quad \quad \quad 0000 \\ \quad \quad 1101 \\ \hline 10001111 = 143_{10} \end{array}$$

# Division

❑ Similar to (but easier than) decimal division

❑ Example: divide  $145_{10}$  and  $11_{10}$  in binary

$$\begin{array}{r}
 11_{10} = 1011 \overline{) 10010001} \\
 \underline{1011} \phantom{0000} \\
 1110 \phantom{000} \\
 \underline{1011} \phantom{000} \\
 1101 \phantom{00} \\
 \underline{1011} \phantom{00} \\
 10 \phantom{00} = 2_{10}
 \end{array}$$

又看够不够除  
 不够→看下一位  
 (101) → 10010  
 够  
 ⇒ 上面商1 (x)  
 做减法  
 10010  
 - 1011  
 -----

1101 = 13<sub>10</sub>  
 10010001 = 145<sub>10</sub>

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# Negative Numbers

□  $n$  = word length = number of bits

□ Sign and Magnitude (SM)

➤ 1-bit sign +  $(n-1)$ -bit magnitude

• Example:  $3_{10} = 0011$  and  $-3 = \underline{1011} = 3_{10}$

➤ Common for people but awkward for computers

第 1 bit      剩下的 bit(s)  
[ 0 正      magnitude  
  1 負

□ 1's complement → 所有 bit 翻過來。

➤ Complement  $N$  bits, i.e.,  $\bar{N} = (2^n - 1) - N$

• Example:  $3 = 0011$  and  $\bar{3} = 1100$

□ 2's complement → 1's comp + 1

➤ Complement  $N$  bits and then add 1, i.e.,  $N^* = 2^n - N = \bar{N} + 1$

➤ Or complement all bits from MSB to the left of the rightmost 1

• Example:  $3 = 0011$  and  $3^* = 1101$

# Signed Binary Integers

**TABLE 1-1**  
Signed Binary Integers (word length:  $n = 4$ )  
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可以有4位来描述数字。

$+N$	Positive Integers (all systems)	$-N$	Sign and Magnitude	Negative Integers 2's Complement $N^*$	1's Complement $\bar{N}$
+0	0000	-0	1000	—	1111
+1	0001	-1	1001	1111	1110
+2	0010	-2	1010	1110	1101
+3	0011	-3	1011	1101	1100
+4	0100	-4	1100	1100	1011
+5	0101	-5	1101	1011	1010
+6	0110	-6	1110	1010	1001
+7	0111	-7	1111	1001	1000
		-8		1000	—

SM, 1's comp  $\rightarrow$  有  $\pm 0$

sign 1: 负, 0: 正  
 $[-2^{n-1}+1, 2^{n-1}-1]$   $[-2^{n-1}, 2^{n-1}-1]$  只有 2's comp 可以表示 -8  $[-2^{n-1}, 2^{n-1}-1]$

❑ For word length  $n = 4$ , there are  $2^4$  different permutations

➤ SM and  $\bar{N}$ :  $[-7, \dots, -0, +0, \dots, +7]$ , i.e.,  $[-2^{n-1}+1, 2^{n-1}-1]$

➤  $N^*$ :  $[-8, \dots, +0, \dots, +7]$ , i.e.,  $[-2^{n-1}, 2^{n-1}-1]$

❑ Always view the first bit as the sign bit

❑ Exercise: what is  $1110_2$ ?

# Addition of 2's Complement Numbers (1/2)

⚡ 注意: 如果没有 overflow

无论是正+负 / 负+负 ... 任何情形  
即可直接算

## Steps

- Add just as if all numbers are positive
- Ignore the carry, if any, from the sign bit

## Cases (assume $A > 0$ , $B > 0$ , and word length = $n$ )

- Case 1:  $A + B$  and  $|A + B| < 2^{n-1}$  → Correct (n=4)  
→ 这个 4 bit 可表示的值
- Case 2:  $A + B$  and  $|A + B| \geq 2^{n-1}$  → Wrong (overflow)
- Case 3:  $A - B$  and  $A < B$  → Correct
- Case 4:  $-A + B$  and  $A \leq B$  → Correct (ignore the carry)
- Case 5:  $-A - B$  and  $|A + B| \leq 2^{n-1}$  → Correct (ignore the carry)
- Case 6:  $-A - B$  and  $|A + B| > 2^{n-1}$  → Wrong (overflow) x's comp: 如果带进位去 → 直接省略

Case 1	Case 2	Case 3	Case 4	Case 5	Case 6
+3    0011	+5 <u>0</u> 101	+5    0101	-5       1011	-3       1101	-5       1011
+4    0100	+6 <u>0</u> 110	-6    1010	+6       0110	-4       1100	-6       1010
+7    0111	<u>1</u> 011	-1    1111	+1    (1) 0001	-7    (1) 1001	(1) <u>0</u> 101

\* 检查是否有 overflow → 看第一位是否有特殊情况 (0+0 应该为 0 → 但变 1 (负) 才代表有 overflow)

# Addition of 2's Complement Numbers (2/2)

## ❑ Why to ignore the carry, i.e., subtract $2^n$ ?

➤ Add(-A, +B) where  $B > A$  省0的bit

- $A^* + B = (2^n - A) + B = 2^n + (B - A)$

➤ Add(-A, -B) where  $A + B \leq 2^{n-1}$

- $A^* + B^* = (2^n - A) + (2^n - B) = 2^n + 2^n - (A + B) = 2^n + (A + B)^*$

## ❑ How to detect overflow?

➤ Check the sign

- (+) + (+) becomes (-)

- (-) + (-) becomes (+)



# Addition of 1's Complement Numbers

## ❑ End-around carry

- Add just as if all numbers are positive
- Add the carry out back to the rightmost bit

Case 1	Case 2	Case 3	Case 4	Case 5	Case 6
$\begin{array}{r} +3 \quad 0011 \\ +4 \quad 0100 \\ \hline +7 \quad 0111 \end{array}$	$\begin{array}{r} +5 \quad \color{red}0101 \\ +6 \quad \color{red}0110 \\ \hline \quad \color{red}1011 \end{array}$	$\begin{array}{r} +5 \quad 0101 \\ -6 \quad 1001 \\ \hline -1 \quad 1110 \end{array}$	$\begin{array}{r} -5 \quad 1010 \\ +6 \quad 0110 \\ \hline +1 \quad (1)0000 \\ \quad \quad 1 \\ \hline \quad \quad 0001 \end{array}$	$\begin{array}{r} -3 \quad 1100 \\ -4 \quad 1011 \\ \hline -7 \quad (1)0111 \\ \quad \quad 1 \\ \hline \quad \quad 1000 \end{array}$	$\begin{array}{r} -5 \quad \color{red}1010 \\ -6 \quad \color{red}1001 \\ \hline \quad (1)\color{red}0011 \\ \quad \quad 1 \\ \hline \quad \quad 0100 \end{array}$

## ❑ How to detect overflow?

- Check the sign
  - (+) + (+) becomes (-)
  - (-) + (-) becomes (+)

一样可检测溢出1 bit

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# Decimal Digits to Binary Codes

## ❑ Input/output interface generally uses decimal digits

- How to code decimal digits using binary codes?
- Choose 10 elements from 16 binary numbers of 4 bits
- Binary-Coded-Decimal (BCD)

• Example: 937.25 → 1001 0011 0111 . 0010 0101

螢幕上 937.25 如何顯示 ↗  
(非靜態二進位的問題)

For error  
checking

For analog  
quantity

**TABLE 1-2**  
Binary Codes for  
Decimal Digits

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Decimal Digit	8-4-2-1 Code (BCD)	6-3-1-1 Code	Excess-3 Code	2-out-of-5 Code	Gray Code
0	0000	0000	0011	00011	0000
1	0001	0001	0100	00101	0001
2	0010	0011	0101	00110	0011
3	0011	0100	0110	01001	0010
4	0100	0101	0111	01010	0110
5	0101	0111	1000	01100	1110
6	0110	1000	1001	10001	1010
7	0111	1001	1010	10010	1011
8	1000	1011	1011	10100	1001
9	1001	1100	1100	11000	1000

4 bit 分列代碼  
6, 3, 1, 1

(BCD+3)

Only 1 bit  
difference for two  
successive digits

# Warning: Conversion or Coding?

## ❑ Do NOT mix up

- Conversion of a decimal number to a binary number
- Coding a decimal digit with a binary code

## ❑ Example

➤ Conversion:  $13_{10} = 1101_2$

➤ Coding:  $13 = 0001\ 0011$

↑  
二进制表示 1, 3



# Text to Binary Codes

## ❑ ASCII

- American Standard Code for Information Interchange
- Developed from telegraph code
- English alphanumeric symbols
- 7 bits
- 94 printable characters are numbered  $32_{10}$  to  $126_{10}$

## ❑ Unicode

- <https://en.wikipedia.org/wiki/Unicode>

## ❑ UTF-8

- <https://en.wikipedia.org/wiki/UTF-8>

## ❑ Big-5

- Traditional Chinese characters
- <https://en.wikipedia.org/wiki/Big5>

ASCII Code													
Character						Character							
A <sub>6</sub>	A <sub>5</sub>	A <sub>4</sub>	A <sub>3</sub>	A <sub>2</sub>	A <sub>1</sub>	A <sub>0</sub>	A <sub>6</sub>	A <sub>5</sub>	A <sub>4</sub>	A <sub>3</sub>	A <sub>2</sub>	A <sub>1</sub>	A <sub>0</sub>
space	0	1	0	0	0	0	@	1	0	0	0	0	0
!	0	1	0	0	0	1	A	1	0	0	0	0	1
"	0	1	0	0	1	0	B	1	0	0	0	1	0
#	0	1	0	0	1	1	C	1	0	0	0	1	1
\$	0	1	0	1	0	0	D	1	0	0	1	0	0
%	0	1	0	1	0	1	E	1	0	0	1	0	1
&	0	1	0	1	1	0	F	1	0	0	1	1	0
'	0	1	0	1	1	1	G	1	0	0	1	1	1
(	0	1	1	0	0	0	H	1	0	1	0	0	0
)	0	1	1	0	0	1	I	1	0	1	0	0	1
*	0	1	1	0	1	0	J	1	0	1	1	0	0
+	0	1	1	0	1	1	K	1	0	1	1	0	1
,	0	1	1	1	0	0	L	1	0	1	1	0	0
-	0	1	1	1	0	1	M	1	0	1	1	0	1
.	0	1	1	1	1	0	N	1	0	1	1	1	0
/	0	1	1	1	1	1	O	1	0	1	1	1	1
0	1	0	1	0	0	0	P	1	0	1	0	0	0
1	1	0	1	0	0	1	Q	1	0	1	0	0	1
2	1	0	1	0	1	0	R	1	0	1	0	1	0
3	1	0	1	0	1	1	S	1	0	1	0	0	1
4	1	0	1	1	0	0	T	1	0	1	1	0	0
5	1	0	1	1	0	1	U	1	0	1	1	0	1
6	1	0	1	1	1	0	V	1	0	1	1	0	0
7	1	0	1	1	1	1	W	1	0	1	1	1	1
8	0	1	1	1	0	0	X	1	0	1	1	0	0
9	0	1	1	1	0	1	Y	1	0	1	1	0	1
:	0	1	1	1	0	1	Z	1	0	1	1	0	0
;	0	1	1	1	0	1	[	1	0	1	1	0	1
<	0	1	1	1	1	0	\	1	0	1	1	1	0
=	0	1	1	1	1	0	^	1	0	1	1	1	0
>	0	1	1	1	1	1	_	1	0	1	1	1	0
?	0	1	1	1	1	1	—	1	0	1	1	1	1
						Character							
,	1	1	0	0	0	0	,	1	1	0	0	0	0
-	1	1	0	0	0	1	a	1	1	0	0	0	1
.	1	1	0	0	1	0	b	1	1	0	0	1	0
/	1	1	0	0	1	1	c	1	1	0	0	1	1
0	1	1	0	1	0	0	d	1	1	0	1	0	0
1	1	1	0	1	0	1	e	1	1	0	1	0	1
2	1	1	0	1	1	0	f	1	1	0	1	1	0
3	1	1	0	1	1	1	g	1	1	0	1	1	1
4	1	1	1	0	0	0	h	1	1	0	1	0	0
5	1	1	1	0	0	1	i	1	1	0	1	0	1
6	1	1	1	0	1	0	j	1	1	0	1	0	1
7	1	1	1	0	1	1	k	1	1	0	1	1	1
8	1	1	1	1	0	0	l	1	1	0	1	1	0
9	1	1	1	1	0	1	m	1	1	0	1	1	0
:	1	1	1	1	1	0	n	1	1	0	1	1	1
;	1	1	1	1	1	1	o	1	1	0	1	1	1
<	1	1	1	1	0	0	p	1	1	1	0	0	0
=	1	1	1	1	0	0	q	1	1	1	0	0	1
>	1	1	1	1	0	1	r	1	1	1	0	0	1
?	1	1	1	1	0	1	s	1	1	1	0	0	1
	1	1	1	1	1	0	t	1	1	1	0	1	0
	1	1	1	1	1	1	u	1	1	1	0	1	0
	1	1	1	1	1	1	v	1	1	1	0	1	1
	1	1	1	1	1	1	w	1	1	1	0	1	1
	1	1	1	1	1	1	x	1	1	1	1	0	0
	1	1	1	1	1	1	y	1	1	1	1	0	0
	1	1	1	1	1	1	z	1	1	1	1	0	1
	1	1	1	1	1	1	{	1	1	1	1	0	1
	1	1	1	1	1	1		1	1	1	1	1	0
	1	1	1	1	1	1	}	1	1	1	1	1	0
	1	1	1	1	1	1	~	1	1	1	1	1	0
	1	1	1	1	1	1	delete	1	1	1	1	1	1

# Q&A