4190.101 **Discrete Mathematics**

Chapter 9 Relations

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Equivalence Relations

Section 9.5

Section Summary

- Equivalence Relations
- Equivalence Classes
- Equivalence Classes and Partitions

Equivalence Relations

- Definition 1: A relation on a set A is called an equivalence relation if it is reflexive, symmetric, and transitive.
- Definition 2: Two elements a and b that are related by an equivalence relation are called equivalent. The notation a ~ b is often used to denote that a and b are equivalent elements with respect to a particular equivalence relation.

Strings

- **Example**: Suppose that R is the relation on the set of strings of English letters such that a R b if and only if l(a) = l(b), where l(x) is the length of the string x. Is R an equivalence relation?
- Solution: Show that all of the properties of an equivalence relation hold.
 - Reflexivity: Because I(a) = I(a), it follows that a R a for all strings a.
 - Symmetry: Suppose that a R b. Since I(a) = I(b), I(b) = I(a) also holds and b R a.
 - Transitivity: Suppose that a R b and b R c. Since l(a) = l(b), and l(b) = l(c), l(a) = l(c) also holds and a R c.

Congruence Modulo m

- **Example**: Let m be an integer with m>1. Show that the relation $R = \{(a,b) \mid a \equiv b \pmod{m}\}$
 - is an equivalence relation on the set of integers.
- **Solution**: Recall that $a \equiv b \pmod{m}$ if and only if m divides a-b.
 - Reflexivity: $a \equiv a \pmod{m}$ since a a = 0 is divisible by m since $0 = 0 \cdot m$.
 - Symmetry: Suppose that $a \equiv b \pmod{m}$. Then a b is divisible by m, and so a b = km, where k is an integer. It follows that b a = (-k) m, so $b \equiv a \pmod{m}$.
 - Transitivity: Suppose that $a \equiv b \pmod{m}$ and $b \equiv c \pmod{m}$. Then m divides both a b and b c. Hence, there are integers k and l with a b = km and b c = lm. We obtain by adding the equations:

$$a-c = (a-b) + (b-c) = km + lm = (k+l) m.$$

Therefore, $a \equiv c \pmod{m}$.

Divides

- **Example**: Show that the "divides" relation on the set of positive integers is not an equivalence relation.
- **Solution**: The properties of reflexivity, and transitivity do hold, but there relation is not transitive. Hence, "divides" is not an equivalence relation.
 - Reflexivity: a | a for all a.
 - Not Symmetric: For example, 2 | 4, but 4 / 2. Hence, the relation is not symmetric.
 - Transitivity: Suppose that a divides b and b divides c. Then there are positive integers k and l such that b = ak and c = bl. Hence, c = a(kl), so a divides c. Therefore, the relation is transitive.

Equivalence Classes

- **Definition 3**: Let R be an equivalence relation on a set A. The set of all elements that are related to an element a of A is called the equivalence class of a. The equivalence class of a with respect to R is denoted by $[a]_R$.
- Note that $[a]_R = \{s \mid (a,s) \subseteq R\}.$
- When only one relation is under consideration, we can write [a], without the subscript R, for this equivalence class.
- If $b \in [a]_R$, then b is called a representative of this equivalence class. Any element of a class can be a representative of the class.
- The equivalence classes of the relation congruence modulo m are called the *congruence classes modulo* m. The congruence class of an integer a modulo m is denoted by $[a]_m$, so $[a]_m = \{..., a-2m, a-m, a, a+m, a+2m, ... \}$. For example,

$$[0]_4 = \{..., -8, -4, 0, 4, 8, ...\}$$

$$[1]_4 = \{..., -7, -3, 1, 5, 9, ...\}$$

$$[2]_4 = \{..., -6, -2, 2, 6, 10, ...\}$$

$$[3]_4 = \{..., -5, -1, 3, 7, 11, ...\}$$

Equivalence Classes and Partitions

- **Theorem 1**: let *R* be an equivalence relation on a set *A*. These statements for elements *a* and *b* of *A* are equivalent:
 - (i) a R b(ii) [a] = [b](iii) $[a] \cap [b] \neq \emptyset$
- **Proof:** We show that (i) implies (ii). Assume that a R b. Now suppose that $c \in [a]$. Then a R c. Because a R b and a R c is symmetric, a R c and a R c is transitive and a R c and a R c, it follows that a R c and a R c it follows that a R c and a R c it follows that a R c and a R c it follows that a R c and a R c it follows that a R c and a R c it follows that a R c and a R c it follows that a R c and a R c it follows that a R c is transitive and a R c and a R c it follows that a R c is transitive and a R c in the following th

(see text for proof that (ii) implies (iii) and (iii) implies (i))

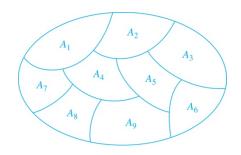
Partition of a Set

• **Definition**: A *partition* of a set S is a collection of disjoint nonempty subsets of S that have S as their union. In other words, the collection of subsets A_i , where $i \in I$ (where I is an index set), forms a partition of S if and only if

$$-A_i \neq \emptyset$$
 for $i \in I$,

$$-A_i \cap A_j = \emptyset$$
 when $i \neq j$,

$$-$$
 and $\bigcup_{i \in I} A_i = S$



A Partition of a Set

An Equivalence Relation Partitions a Set

• Let R be an equivalence relation on a set A. The union of all the equivalence classes of R is all of A, since an element a of A is in its own equivalence class $[a]_R$. In other words,

$$\bigcup_{a \in A} [a]_R = A$$

- From Theorem 1, it follows that these equivalence classes are either equal or disjoint, so $[a]_R \cap [b]_R = \emptyset$ when $[a]_R \neq [b]_R$.
- Therefore, the equivalence classes form a partition of A, because they split A into disjoint subsets.

An Equivalence Relation Partitions a Set (continued)

- **Theorem 2**: Let R be an equivalence relation on a set S. Then the equivalence classes of R form a partition of S. Conversely, given a partition $\{A_i \mid i \in I\}$ of the set S, there is an equivalence relation R that has the sets A_i , $i \in I$, as its equivalence classes.
- Proof: We have already shown the first part of the theorem.
- For the second part, assume that {A_i | i ∈ I} is a partition of S. Let R be the relation on S consisting of the pairs (x, y) where x and y belong to the same subset A_i in the partition. We must show that R satisfies the properties of an equivalence relation.
 - Reflexivity: For every $a \in S$, $(a,a) \in R$, because a is in the same subset as itself.
 - Symmetry: If $(a,b) \in R$, then b and a are in the same subset of the partition, so $(b,a) \in R$.
 - Transitivity: If $(a,b) \in R$ and $(b,c) \in R$, then a and b are in the same subset of the partition, as are b and c. Since the subsets are disjoint and b belongs to both, the two subsets of the partition must be identical. Therefore, $(a,c) \in R$ since a and c belong to the same subset of the partition.

Partial Orderings

Section 9.6

Section Summary

- Partial Orderings and Partially-ordered Sets
- Lexicographic Orderings
- Hasse Diagrams
- Lattices (not currently in overheads)
- Topological Sorting (not currently in overheads)

Partial Orderings

- Definition 1: A relation R on a set S is called a partial ordering, or partial order, if it is reflexive, antisymmetric, and transitive.
 - A set together with a partial ordering R is called a partially ordered set, or poset, and is denoted by (S, R). Members of S are called elements of the poset.

Partial Orderings (continued)

- Example 1: Show that the "greater than or equal" relation (≥) is a partial ordering on the set of integers.
 - Reflexivity: $a \ge a$ for every integer a.
 - Antisymmetry: If $a \ge b$ and $b \ge a$, then a = b.
 - Transitivity: If $a \ge b$ and $b \ge c$, then $a \ge c$.

These properties all follow from the order axioms for the integers. (See Appendix 1).

Partial Orderings (continued)

- Example 2: Show that the divisibility relation
 () is a partial ordering on the set of integers.
 - Reflexivity: a | a for all integers a. (see Example 9 in Section 9.1)
 - Antisymmetry: If a and b are positive integers with $a \mid b$ and $b \mid a$, then a = b. (see Example 12 in Section 9.1)
 - Transitivity: Suppose that a divides b and b divides c. Then there are positive integers k and l such that b = ak and c = bl. Hence, c = a(kl), so a divides c. Therefore, the relation is transitive.
- (**Z**⁺, |) is a poset.

Partial Orderings (continued)

- Example 3: Show that the inclusion relation (\subseteq) is a partial ordering on the power set of a set S.
 - Reflexivity: $A \subseteq A$ whenever A is a subset of S.
 - Antisymmetry: If A and B are positive integers with $A \subseteq B$ and $B \subseteq A$, then A = B.
 - Transitivity: If $A \subseteq B$ and $B \subseteq C$, then $A \subseteq C$.

The properties all follow from the definition of set inclusion.

Comparability

• **Definition 2:** The elements a and b of a poset (S, \leq) are comparable if either $a \leq b$ or $b \leq a$. When a and b are elements of S so that neither $a \leq b$ nor $b \leq a$, then a and b are called incomparable.

The symbol ≤ is used to denote the relation in any poset.

- **Definition 3:** If (S, \leq) is a poset and every two elements of S are comparable, S is called a *totally ordered* or *linearly ordered set*, and \leq is called a *total order* or a *linear order*. A totally ordered set is also called a *chain*.
- **Definition 4:** (S, \leq) is a well-ordered set if it is a poset such that \leq is a total ordering and every nonempty subset of S has a least element.

Lexicographic Order

• **Definition:** Given two posets (A_1, \leq_1) and (A_2, \leq_2) , the lexicographic ordering on $A_1 \times A_2$ is defined by specifying that (a_1, a_2) is less than (b_1, b_2) , that is,

$$(a_1, a_2) \prec (b_1, b_2),$$

either if $a_1 \prec_1 b_1$ or if $a_1 = b_1$ and $a_2 \prec_2 b_2$.

- **Example:** Consider strings of lowercase English letters. A lexicographic ordering can be defined using the ordering of the letters in the alphabet. This is the same ordering as that used in dictionaries.
 - discreet \prec discrete, because these strings differ in the seventh position and $e \prec t$.
 - discreet
 < discreetness, because the first eight letters agree, but the second string is longer.