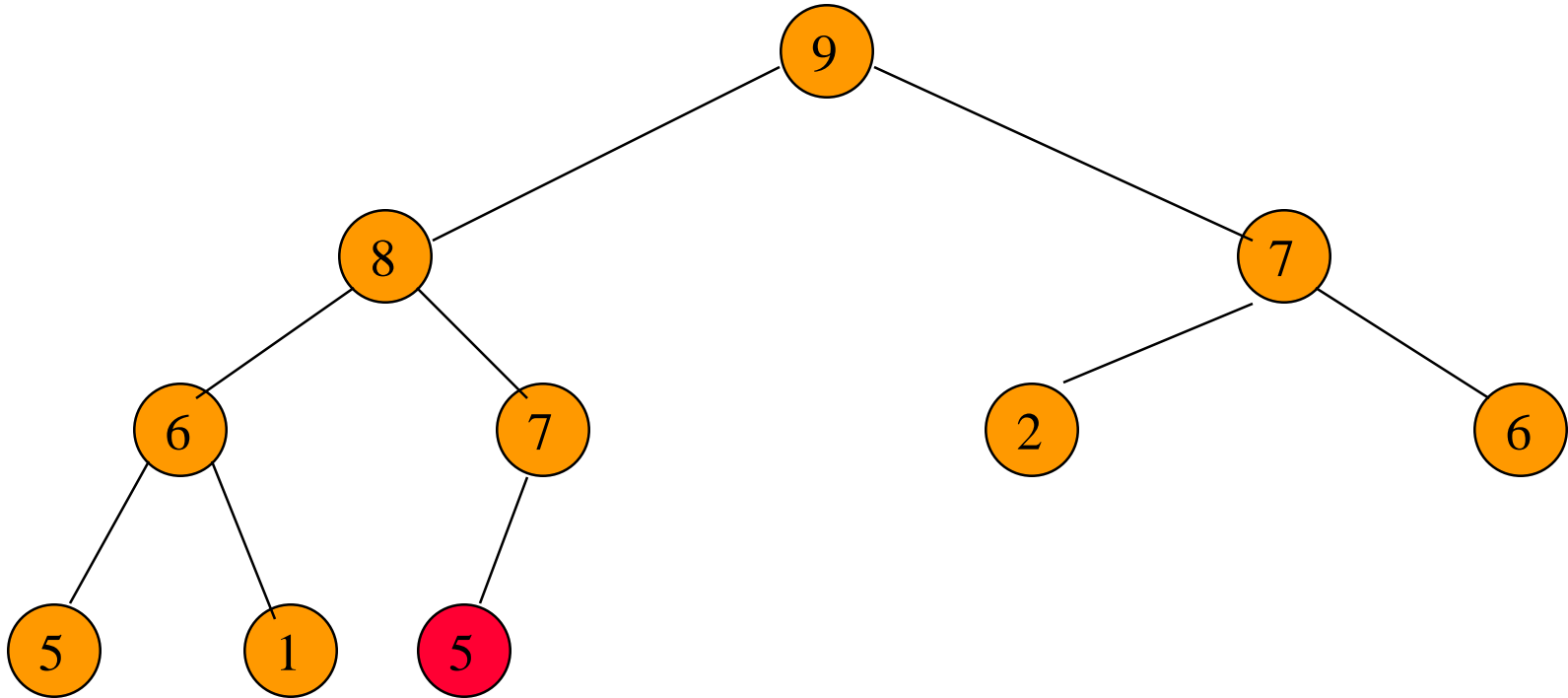


# Heap Operations

Data structures

Spring 2017

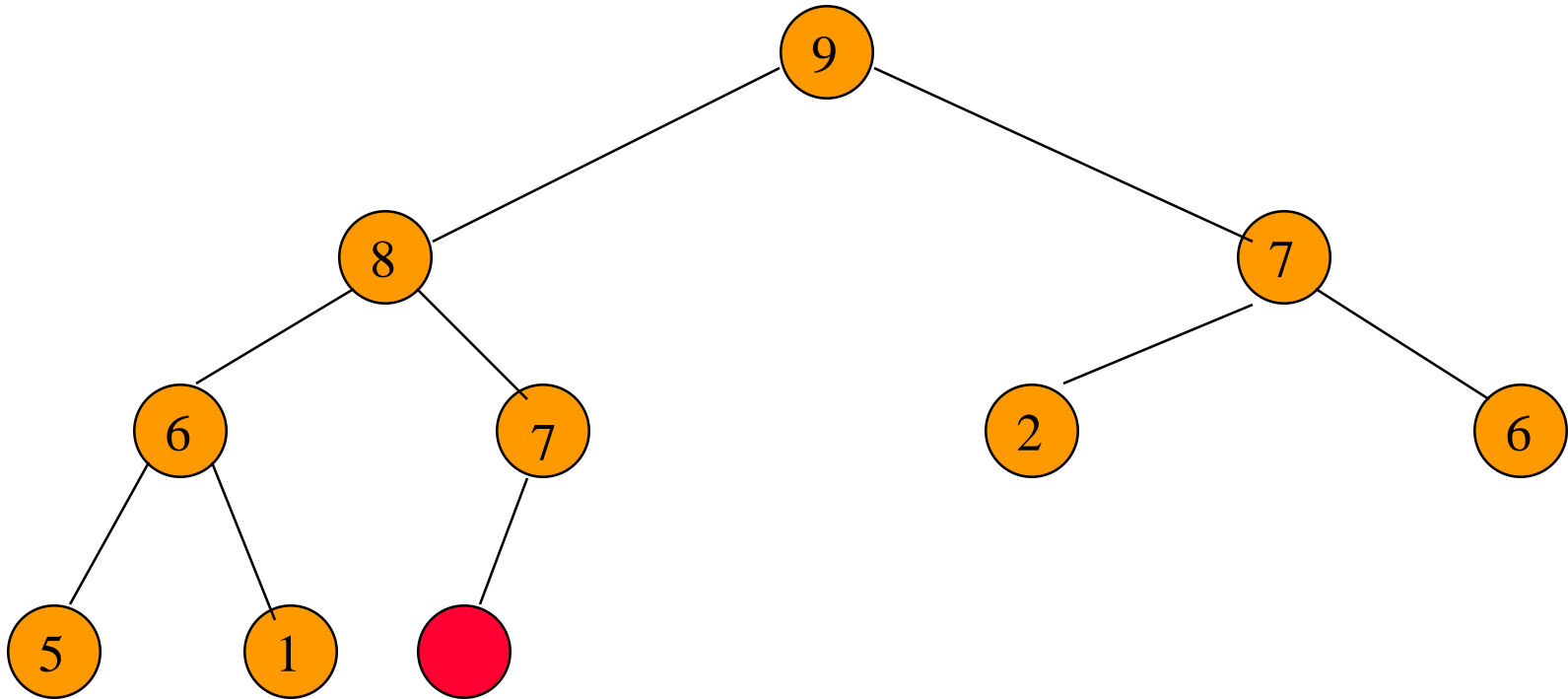
# Putting An Element Into A Max Heap



Complete binary tree with 10 nodes.

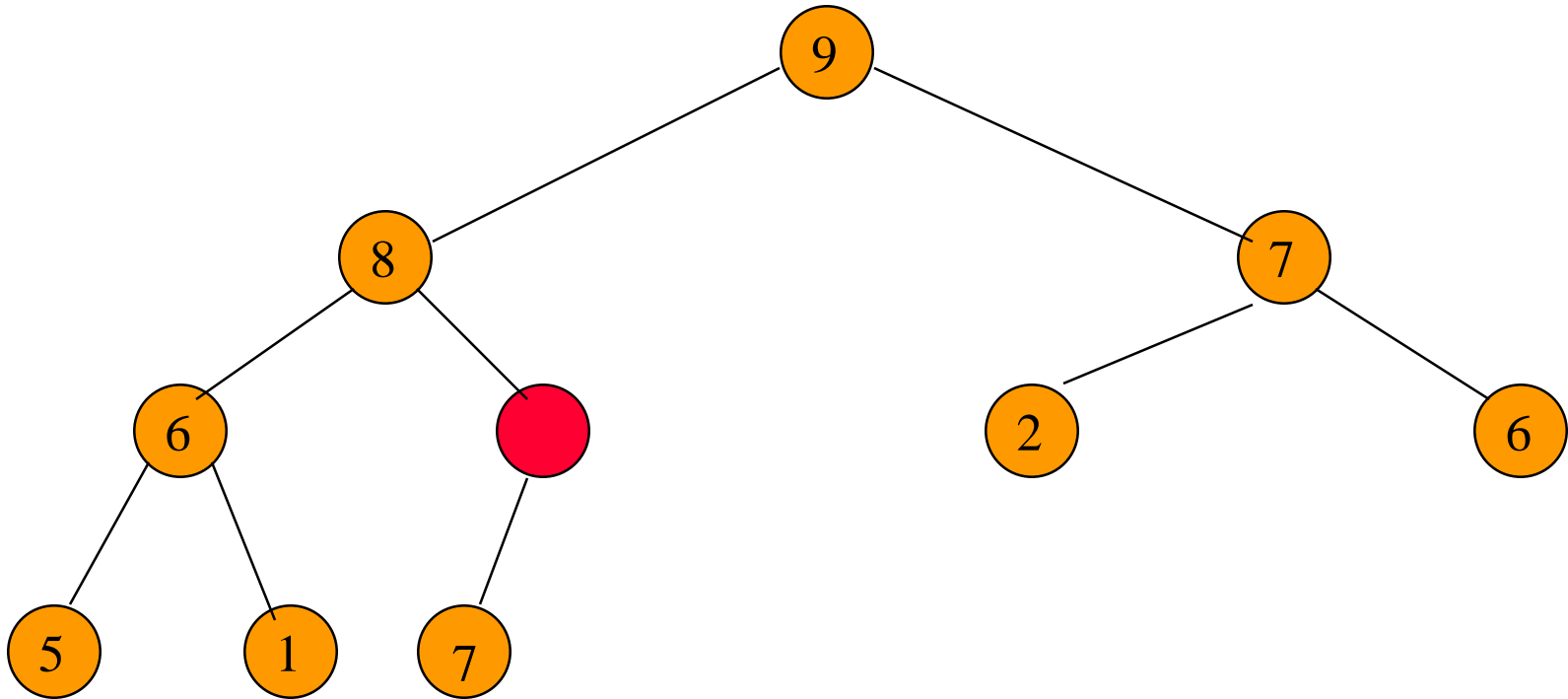
New element is 5.

# Putting An Element Into A Max Heap



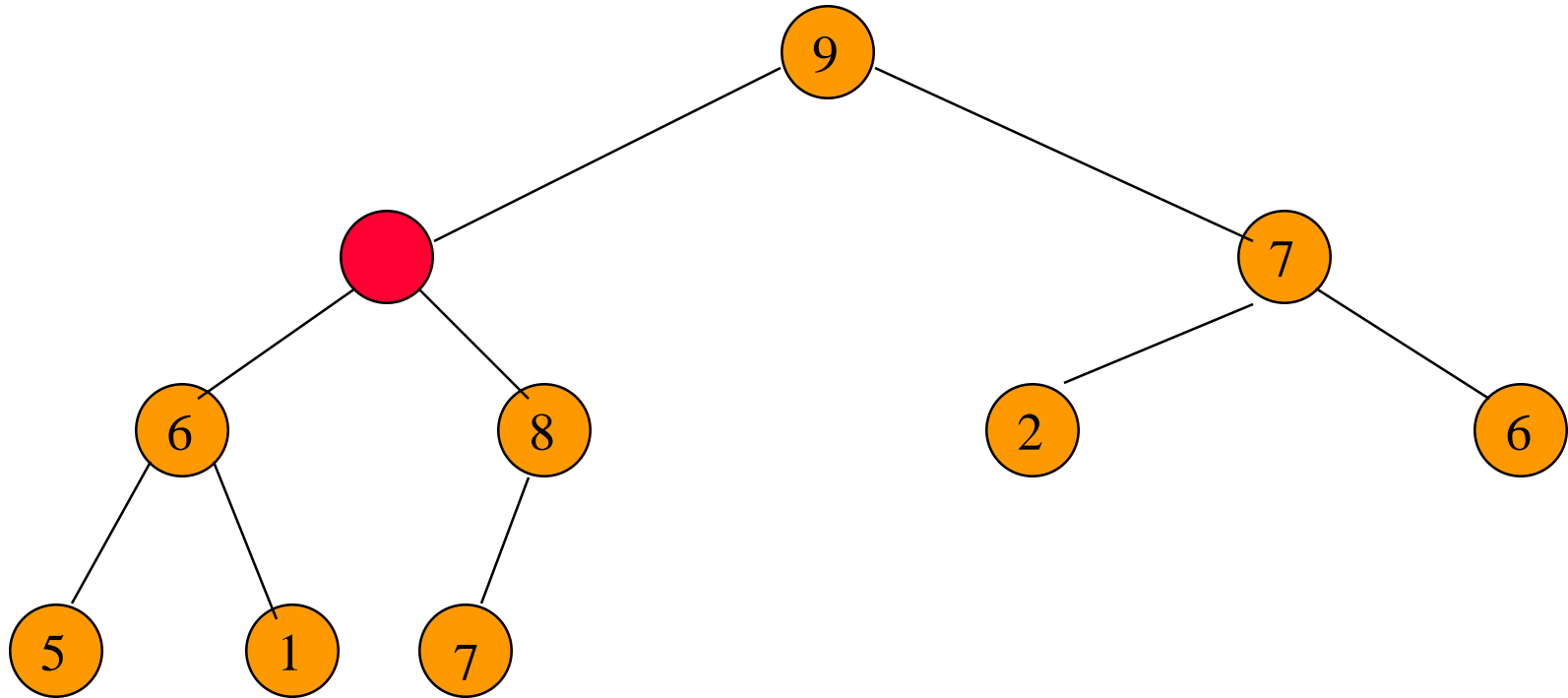
New element is 20.

# Putting An Element Into A Max Heap



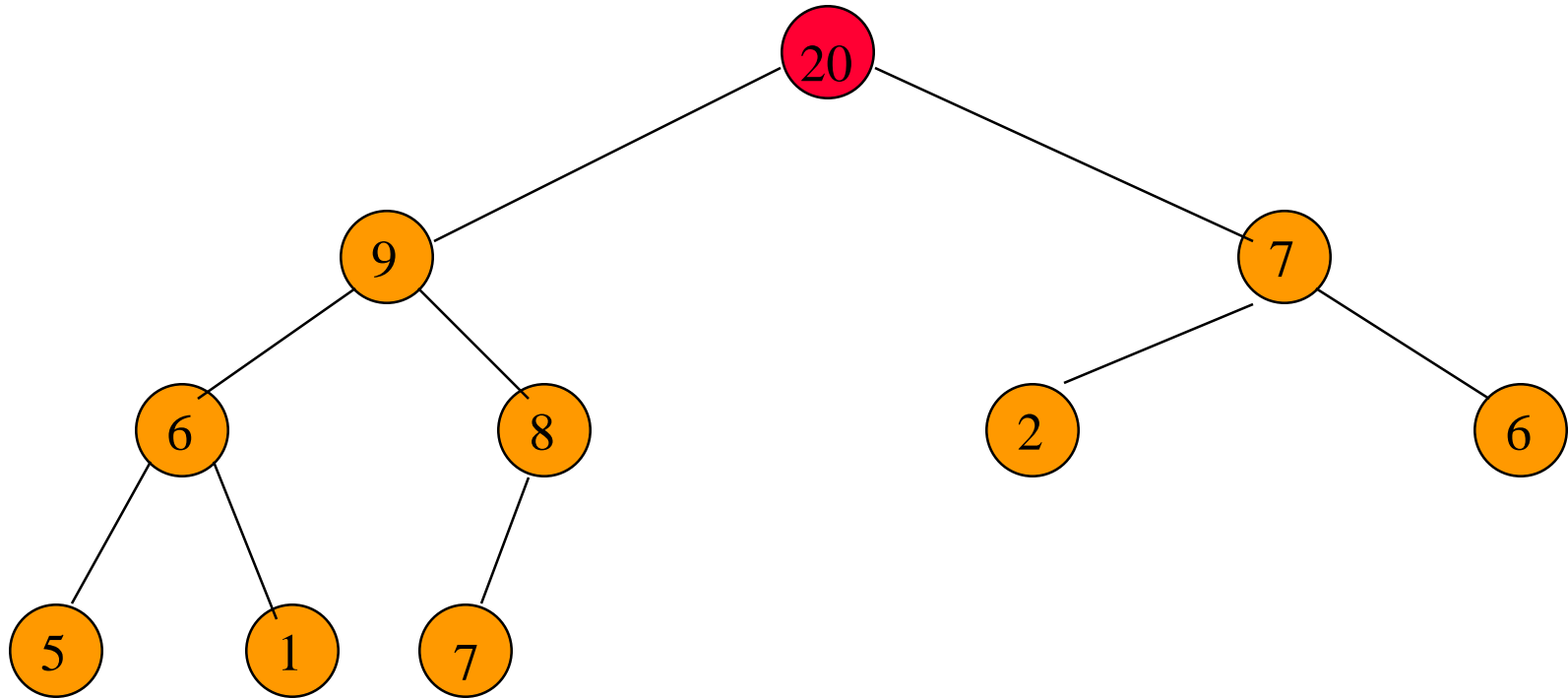
New element is 20.

# Putting An Element Into A Max Heap



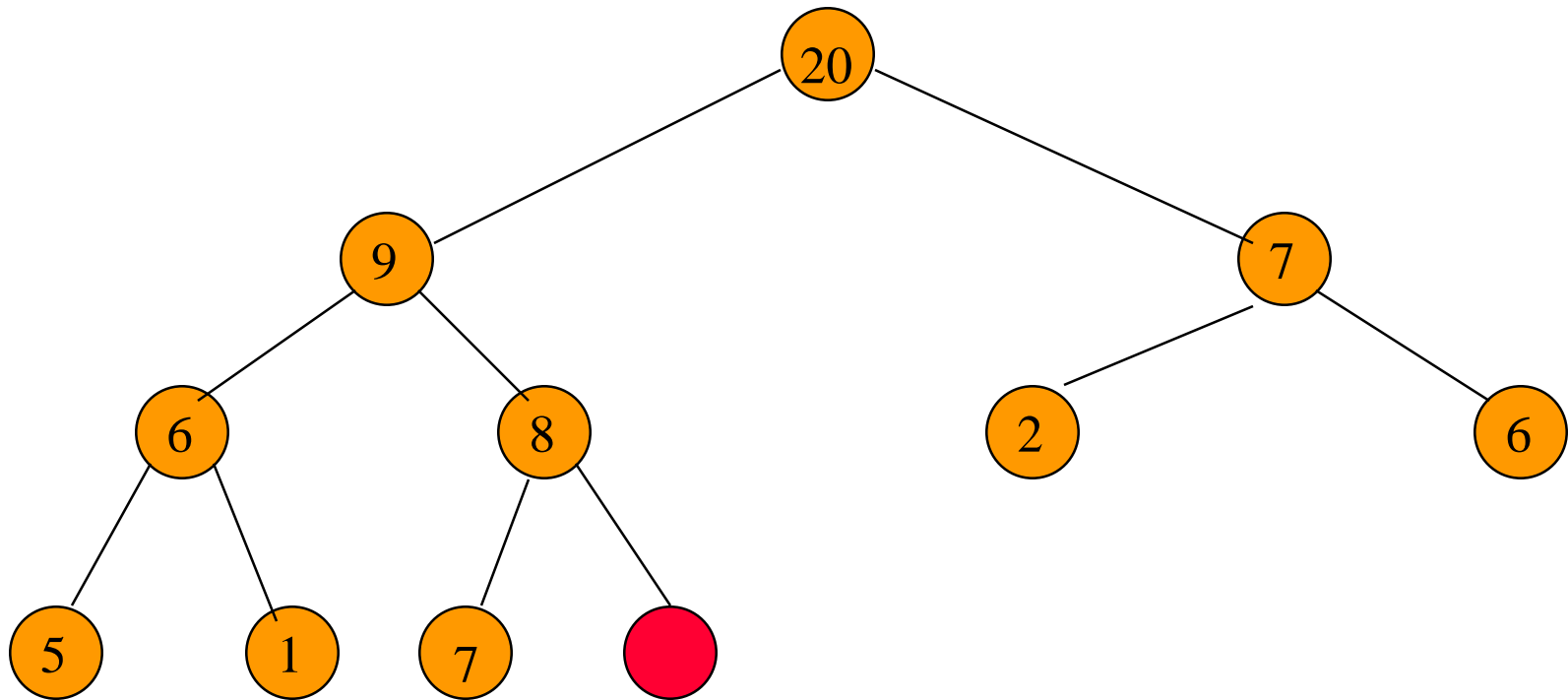
New element is 20.

# Putting An Element Into A Max Heap



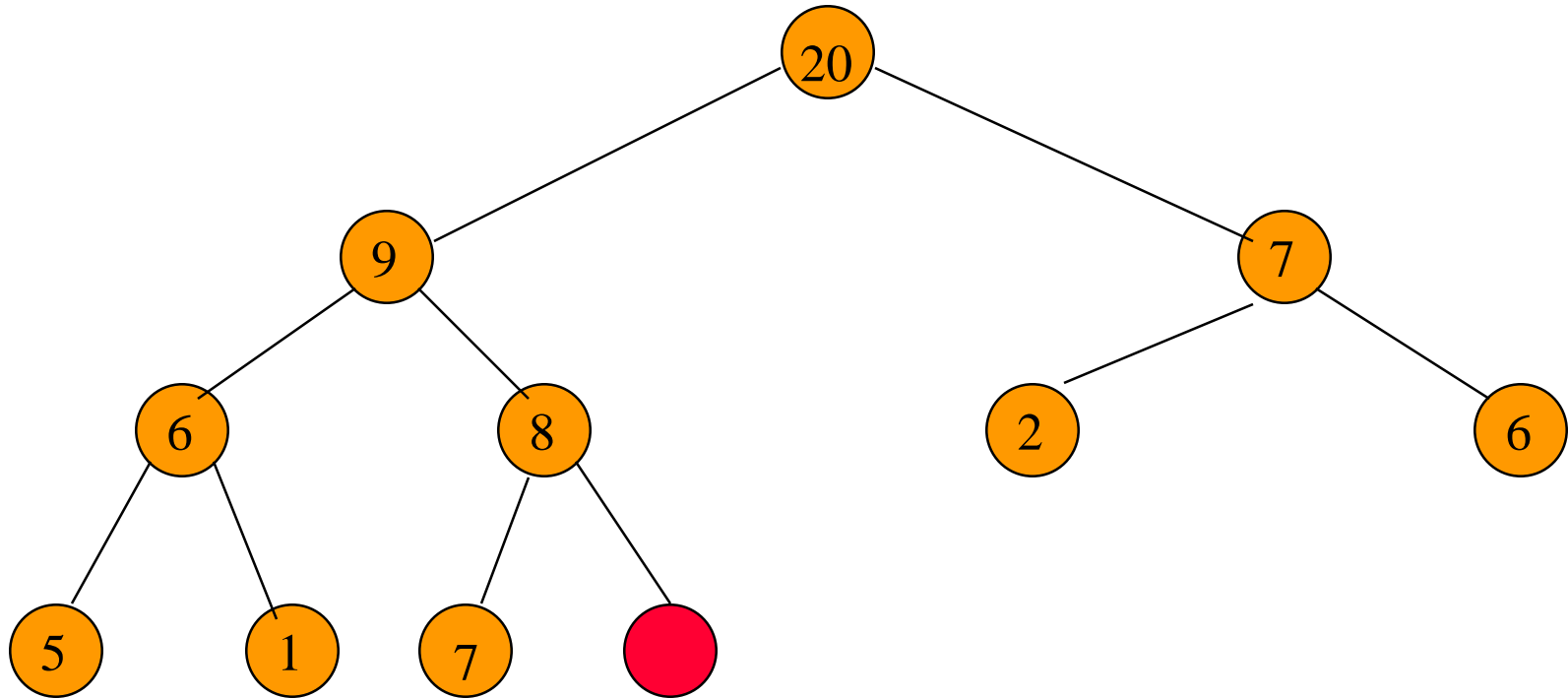
New element is 20.

# Putting An Element Into A Max Heap



Complete binary tree with **11** nodes.

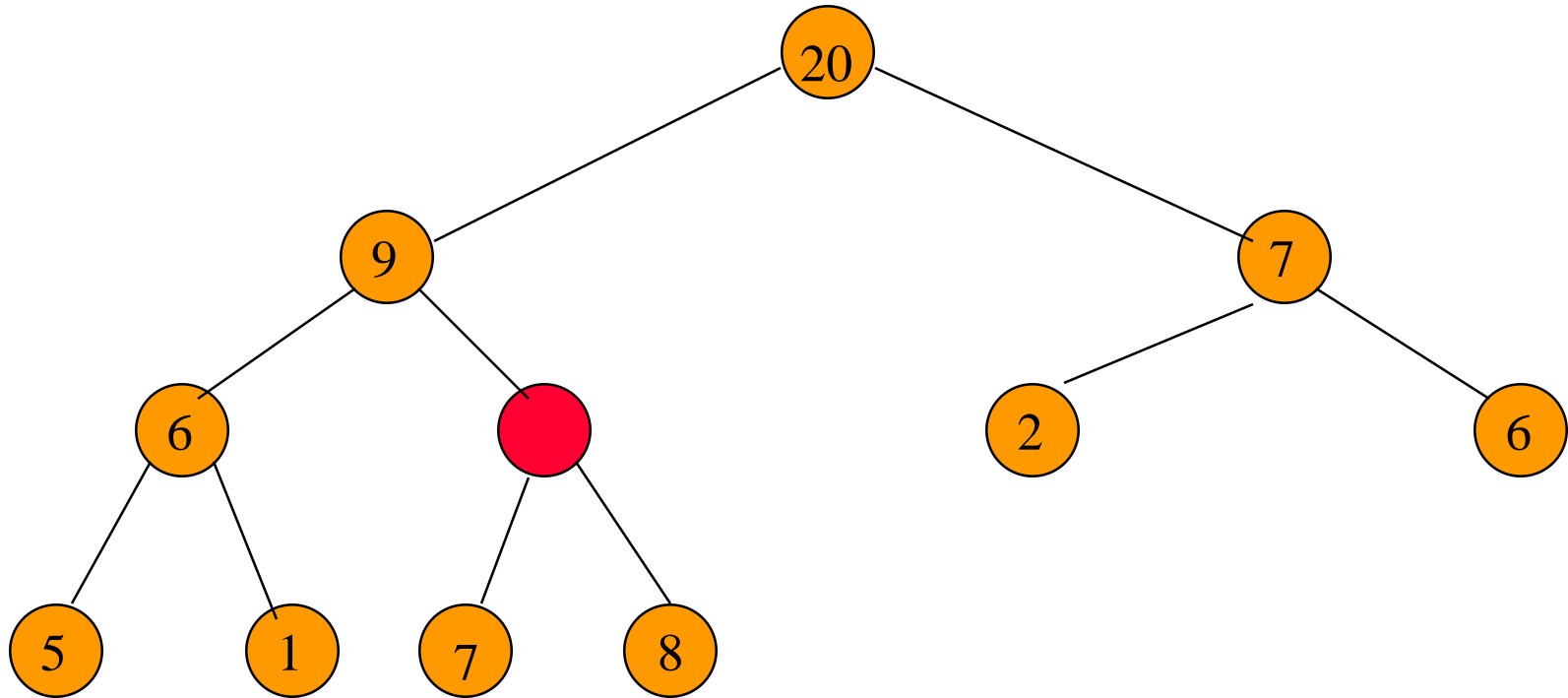
# Putting An Element Into A Max Heap



New element is 15.

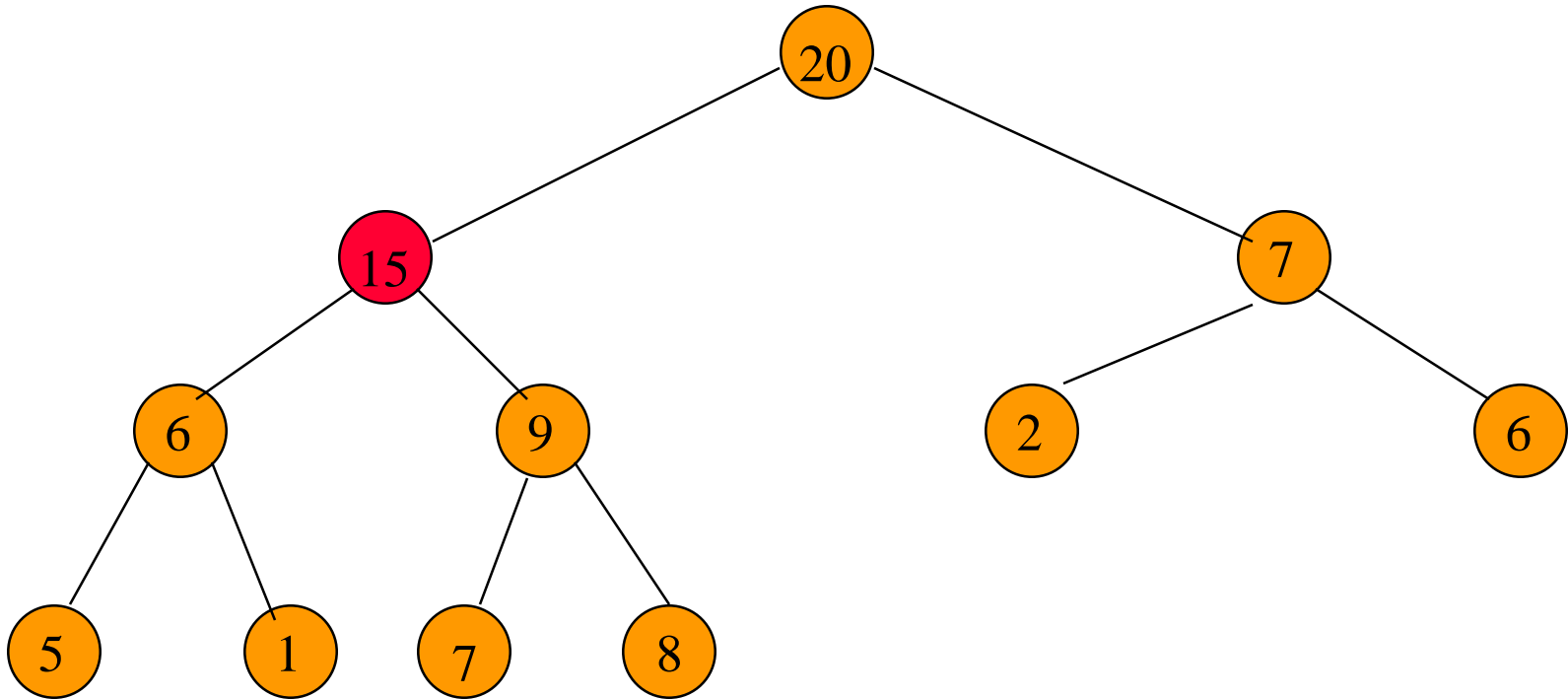


# Putting An Element Into A Max Heap



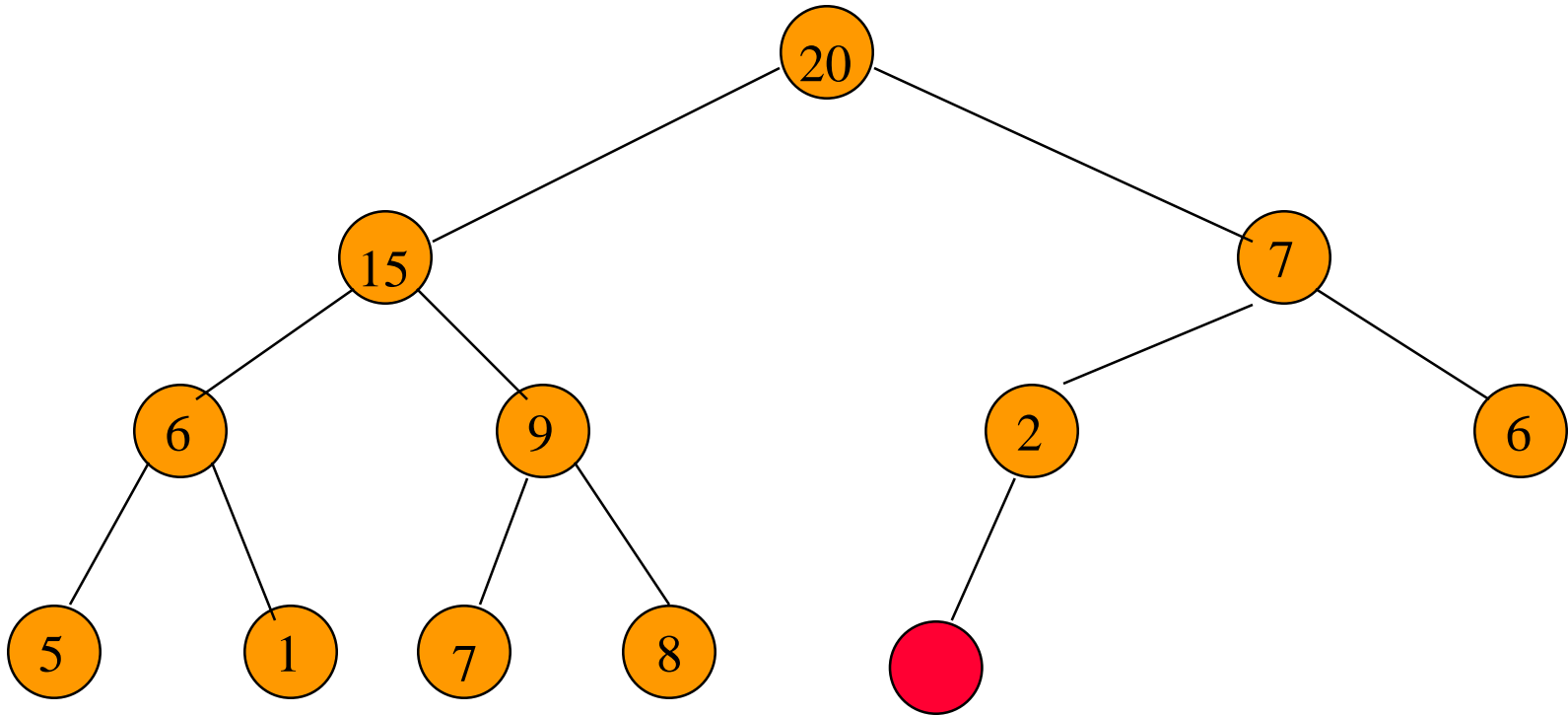
New element is 15.

# Putting An Element Into A Max Heap



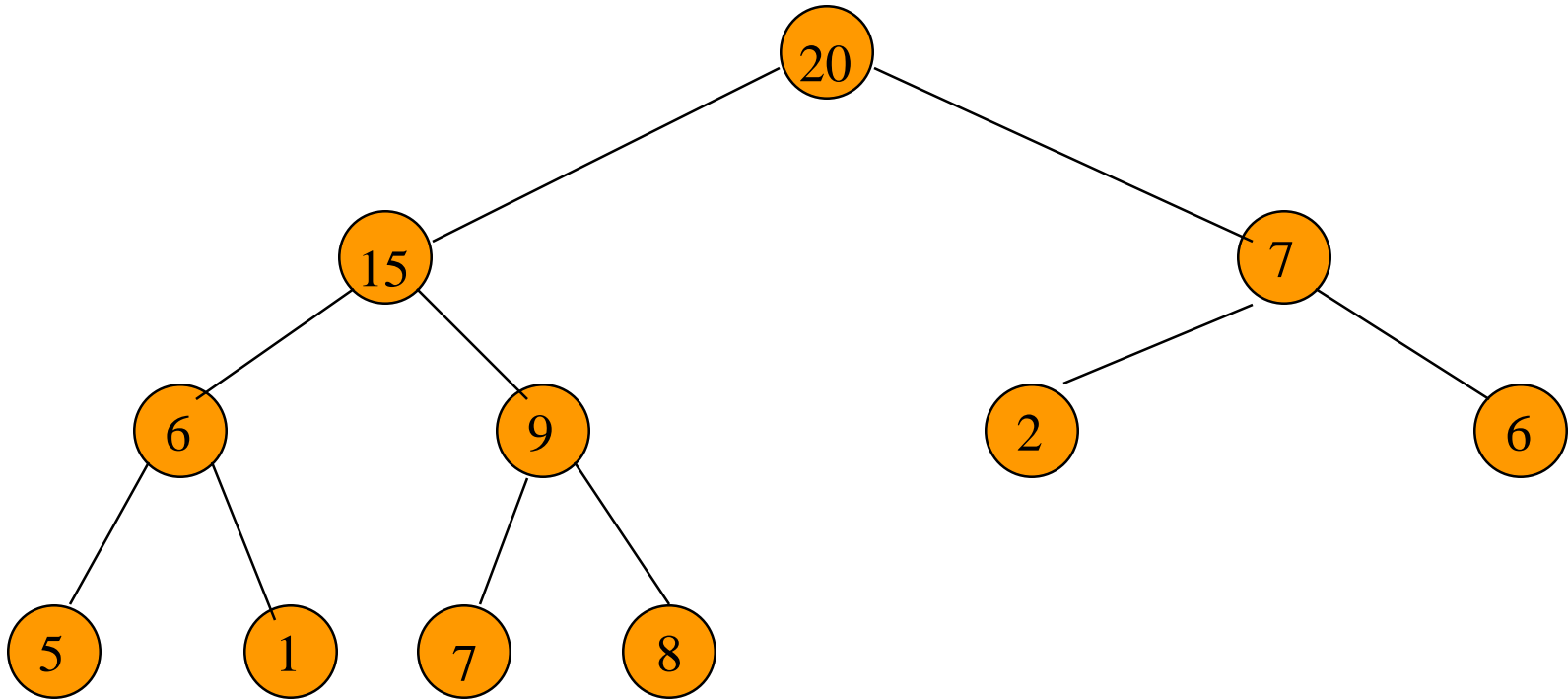
New element is 15.

# Complexity Of Put



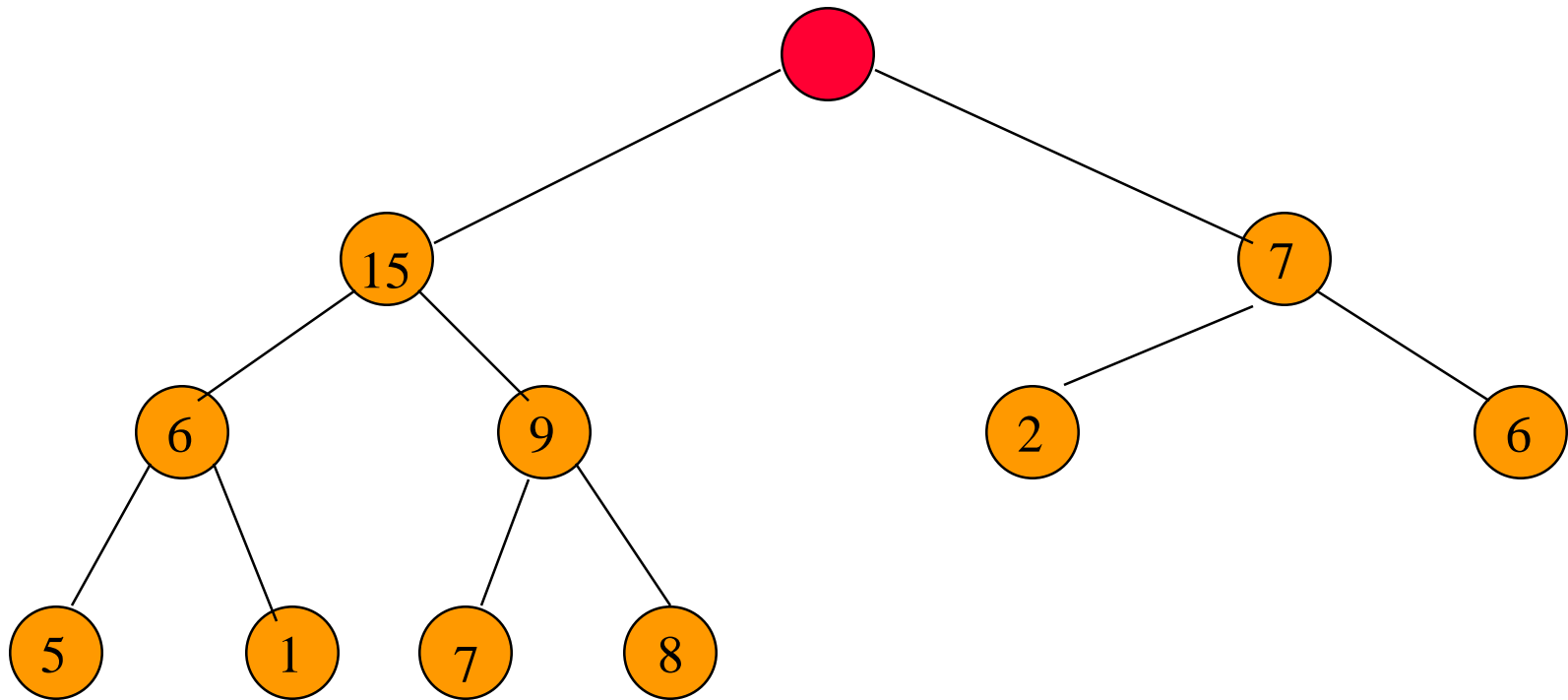
Complexity is  $O(\log n)$ , where  $n$  is heap size.

# Removing The Max Element



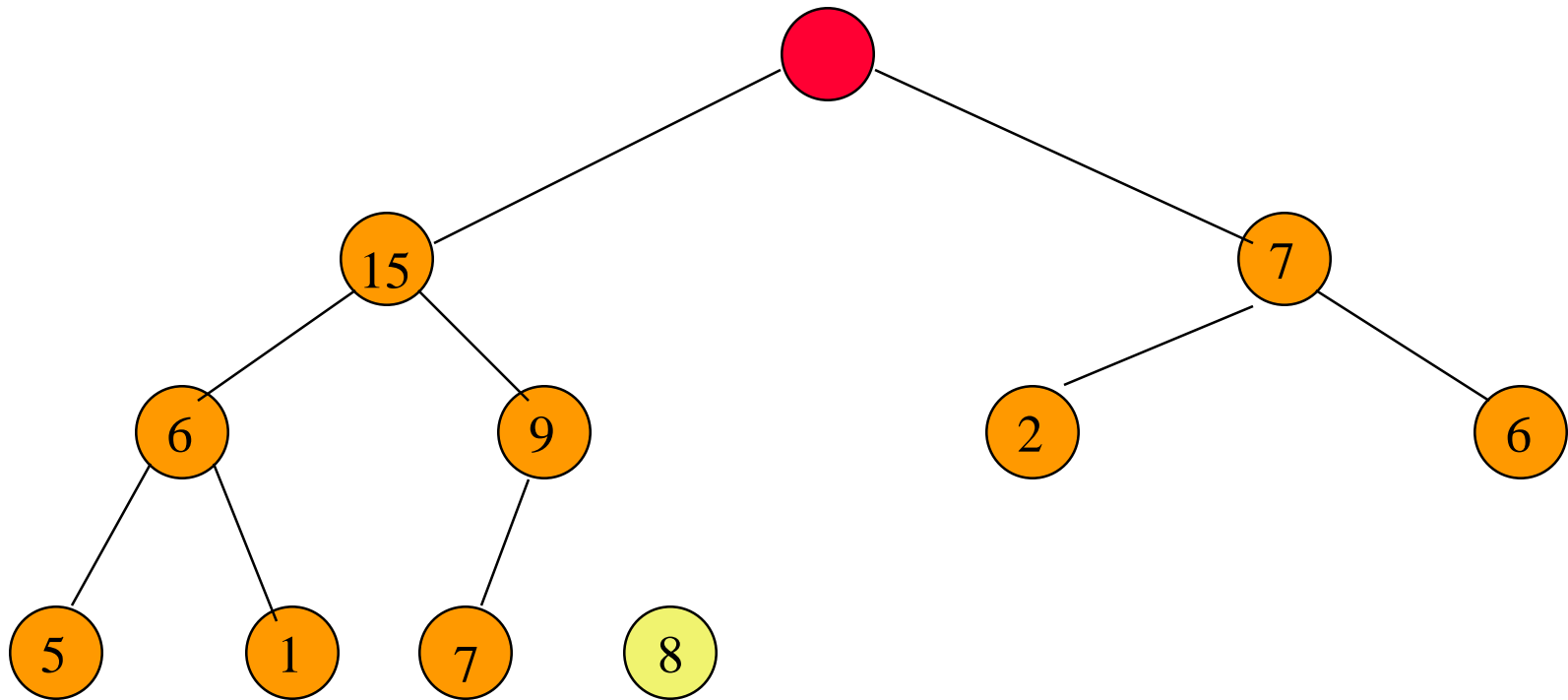
Max element is in the root.

# Removing The Max Element



After max element is removed.

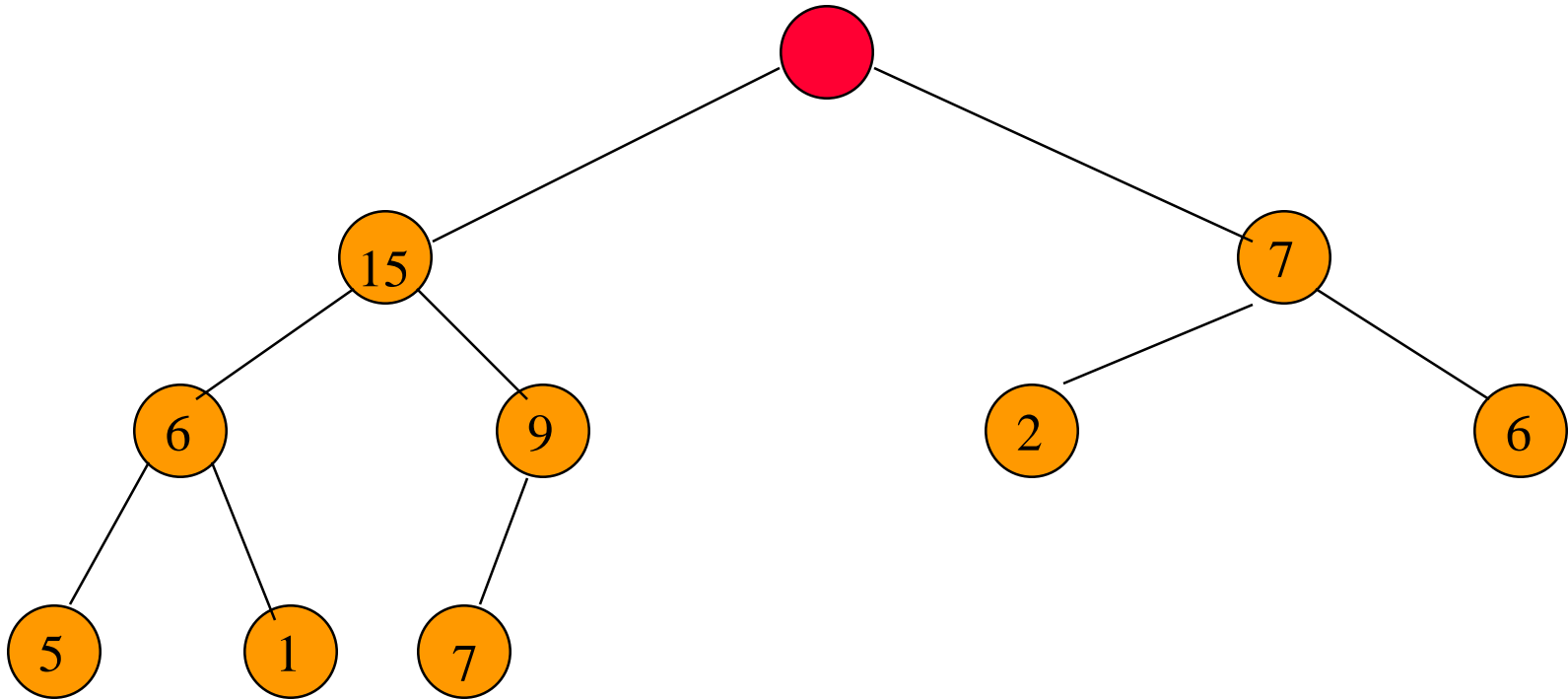
# Removing The Max Element



Heap with 10 nodes.

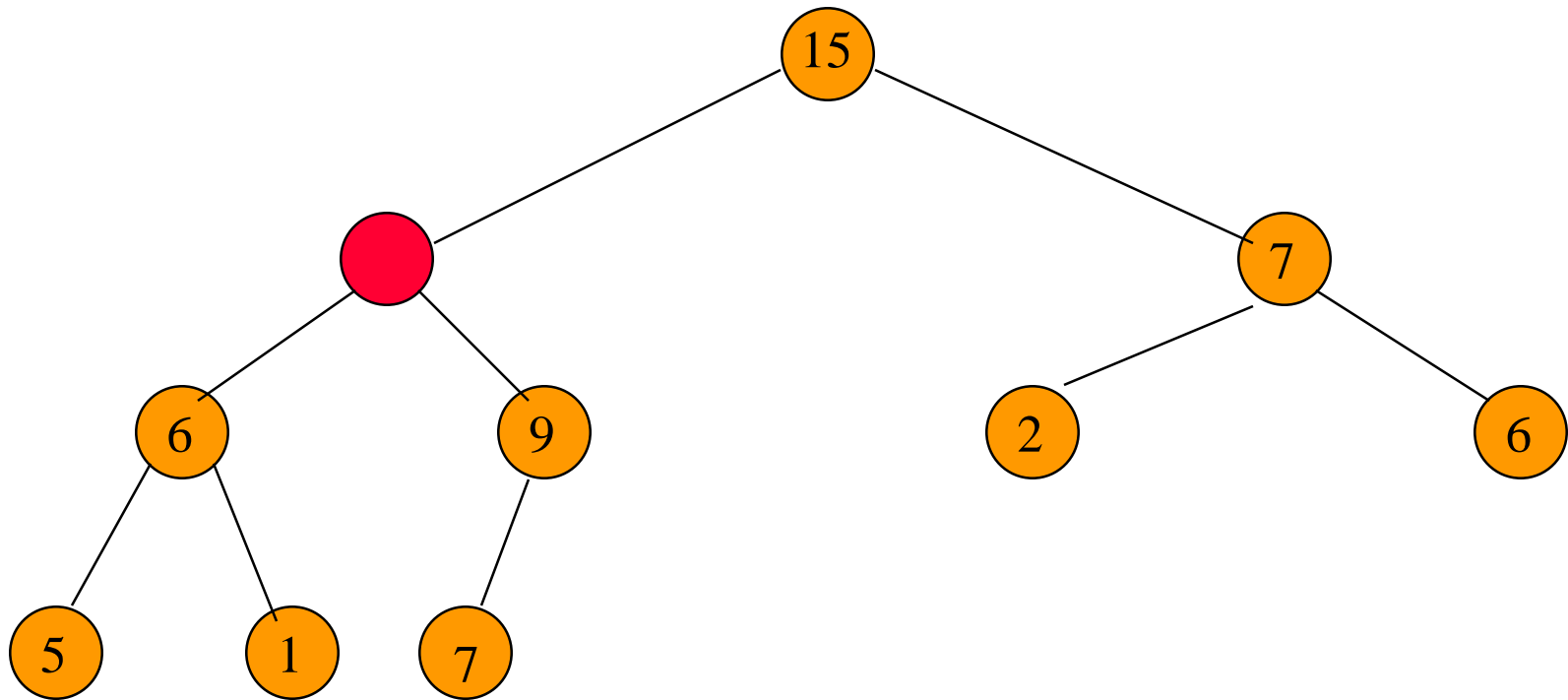
Reinsert 8 into the heap.

# Removing The Max Element



Reinsert **8** into the heap.

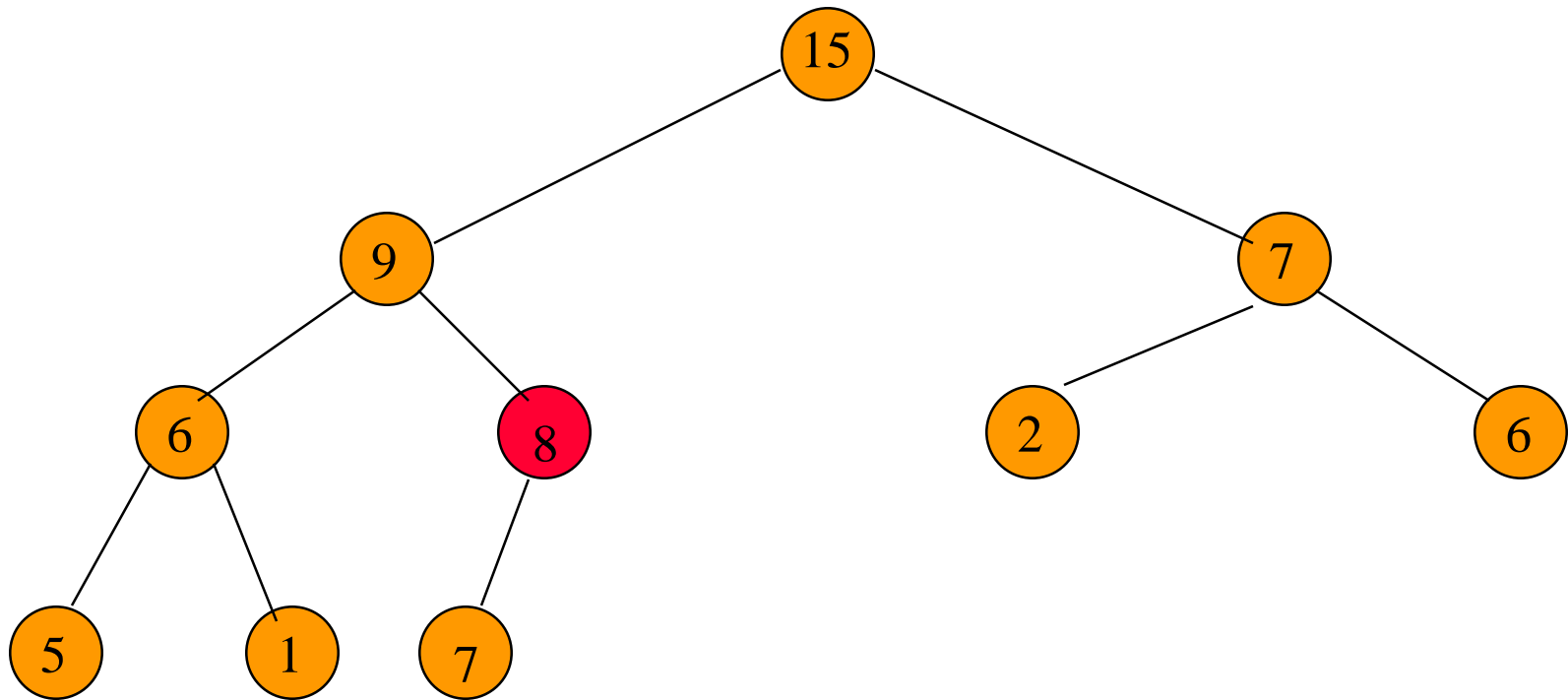
# Removing The Max Element



Reinsert **8** into the heap.

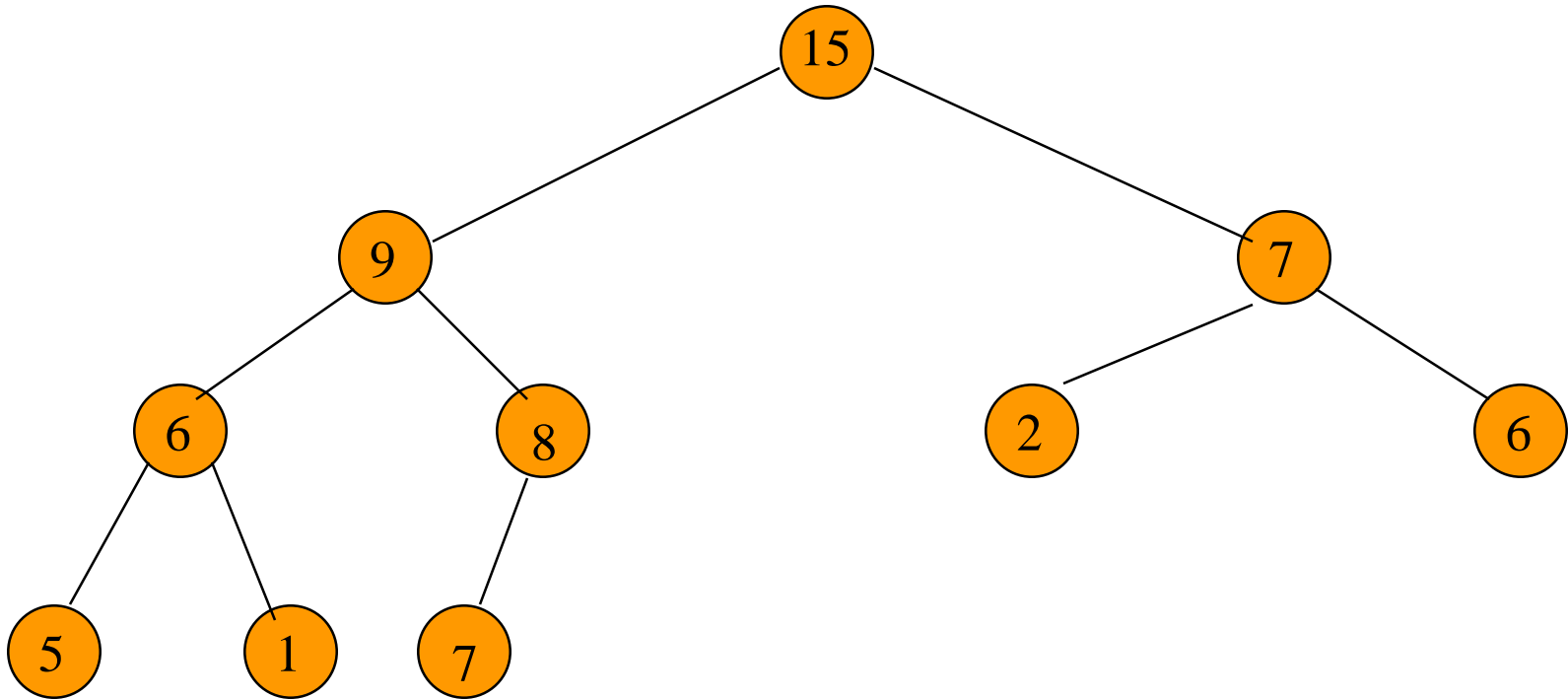


# Removing The Max Element



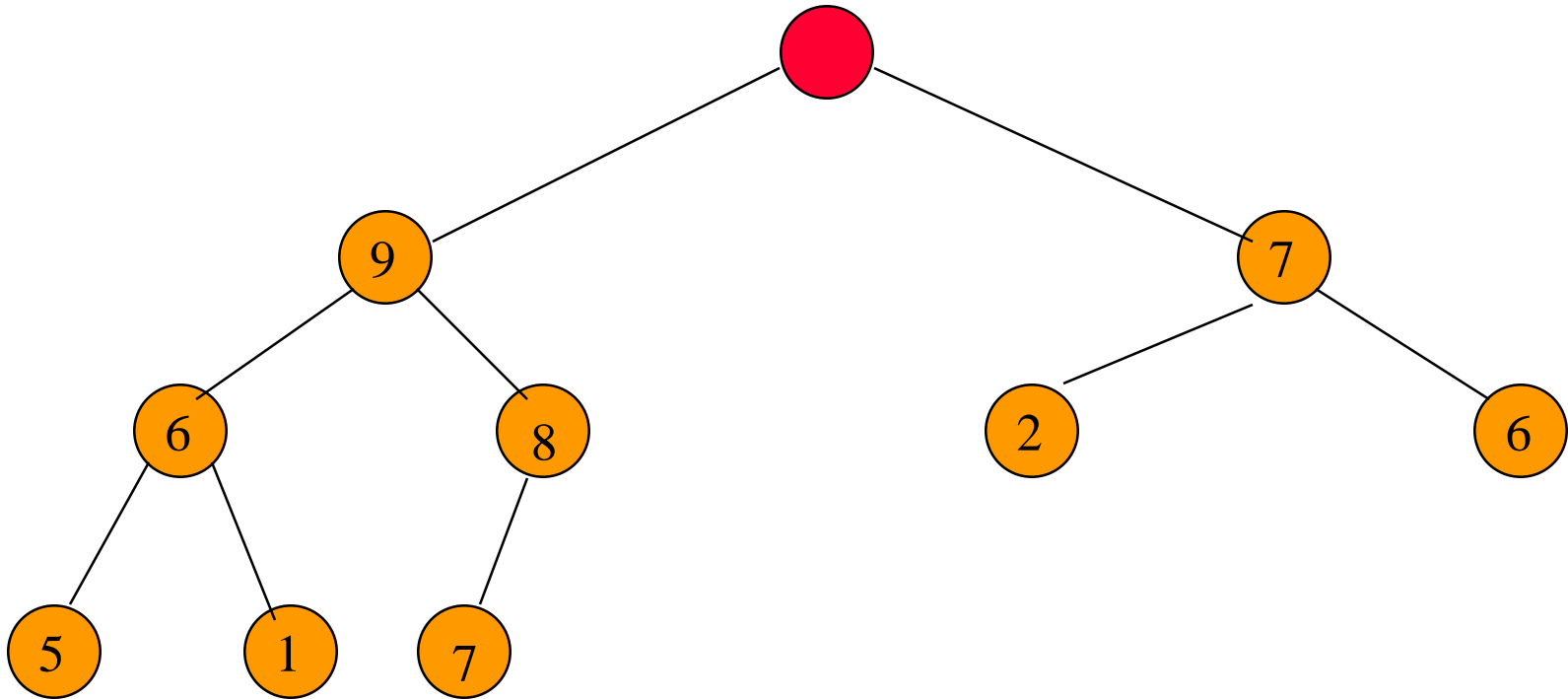
Reinsert **8** into the heap.

# Removing The Max Element



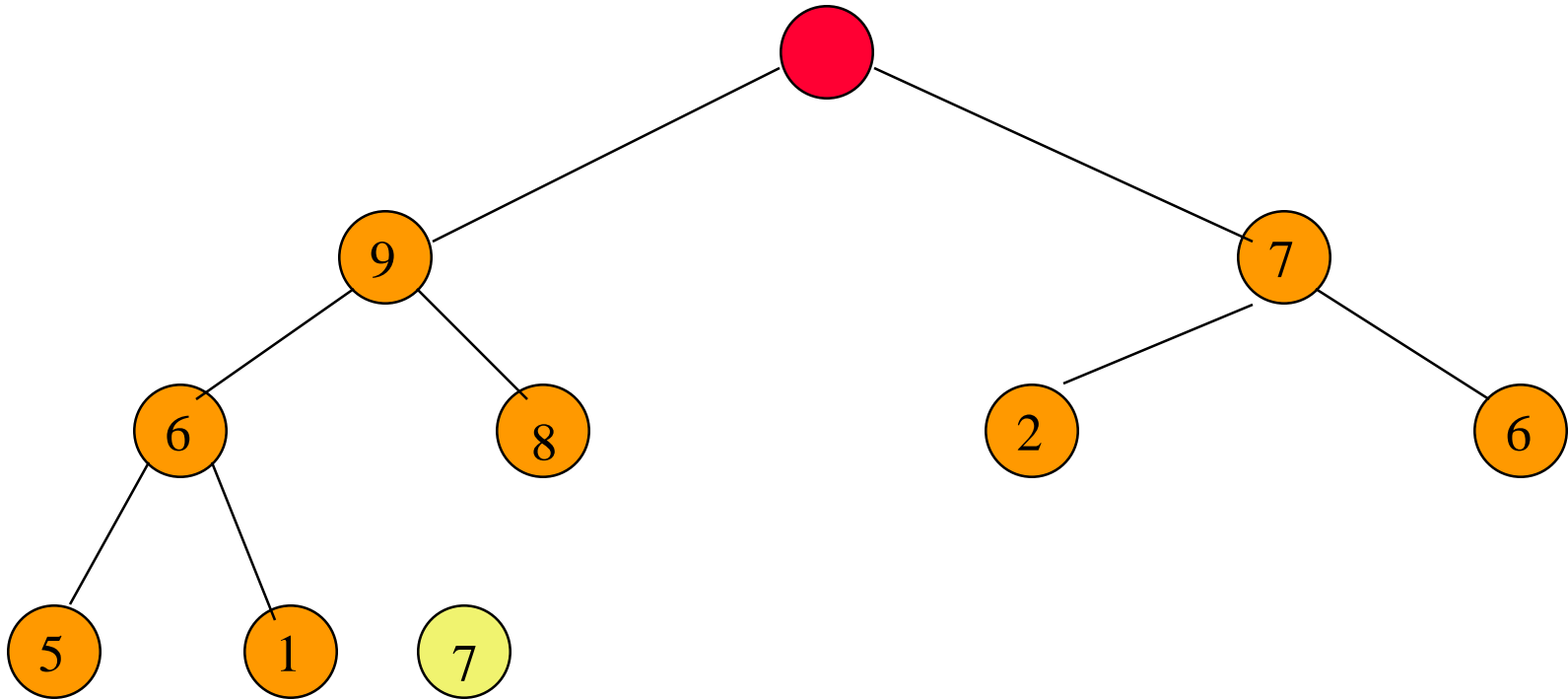
Max element is 15.

# Removing The Max Element



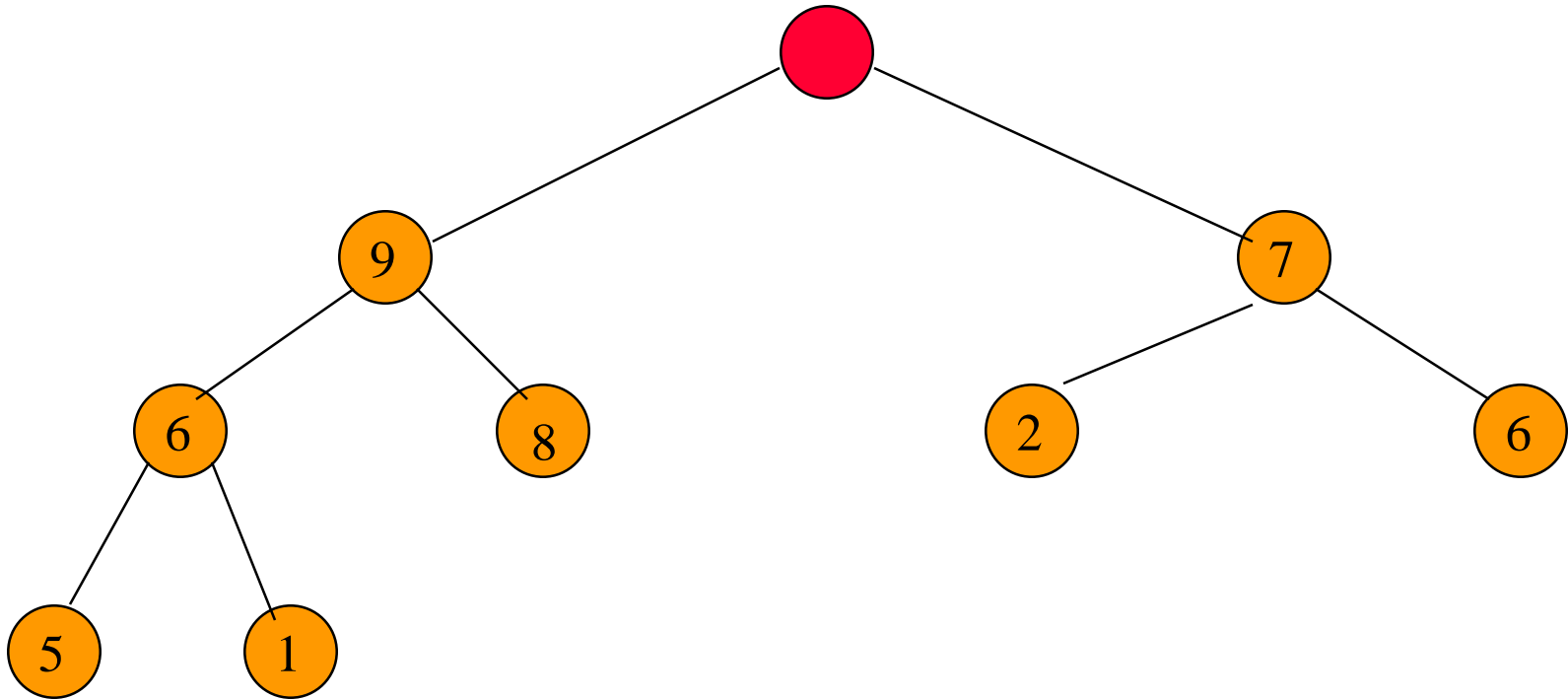
After max element is removed.

# Removing The Max Element



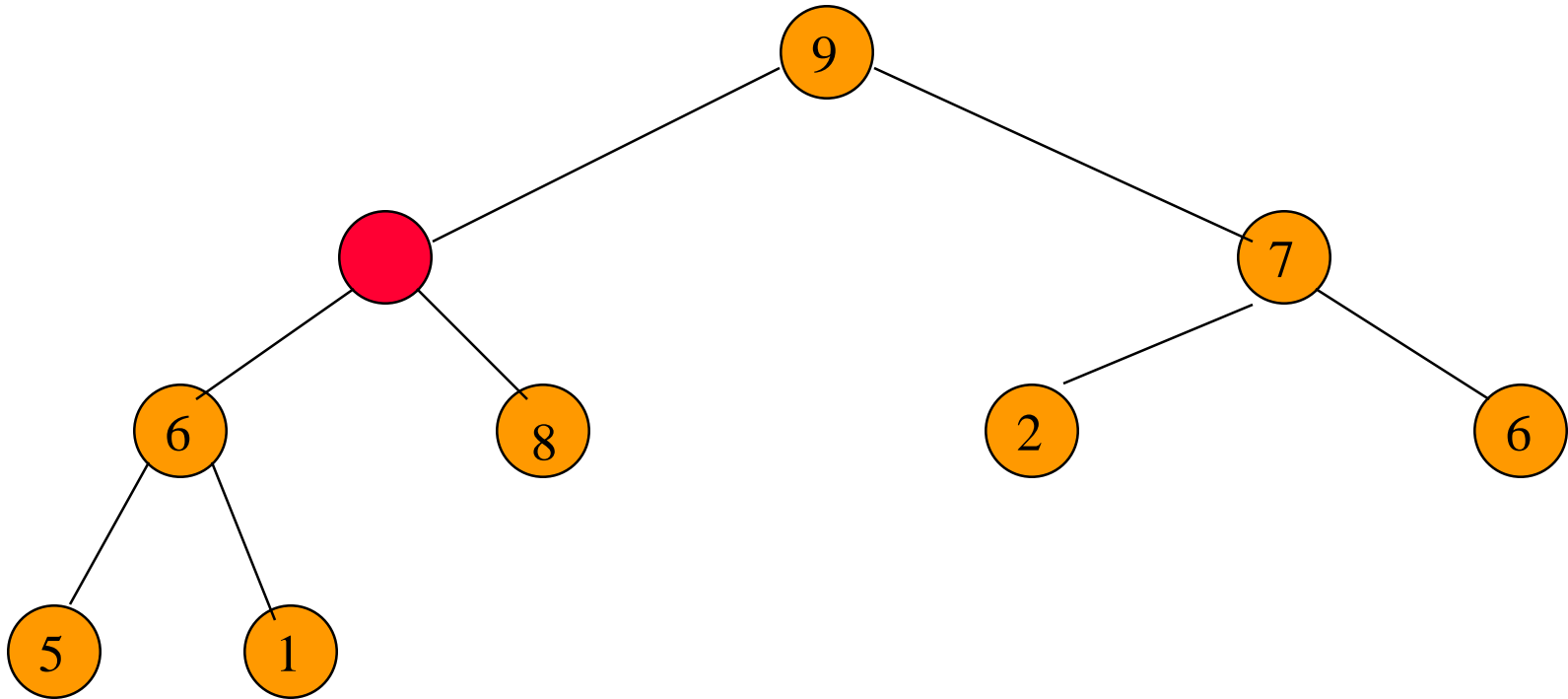
Heap with 9 nodes.

# Removing The Max Element



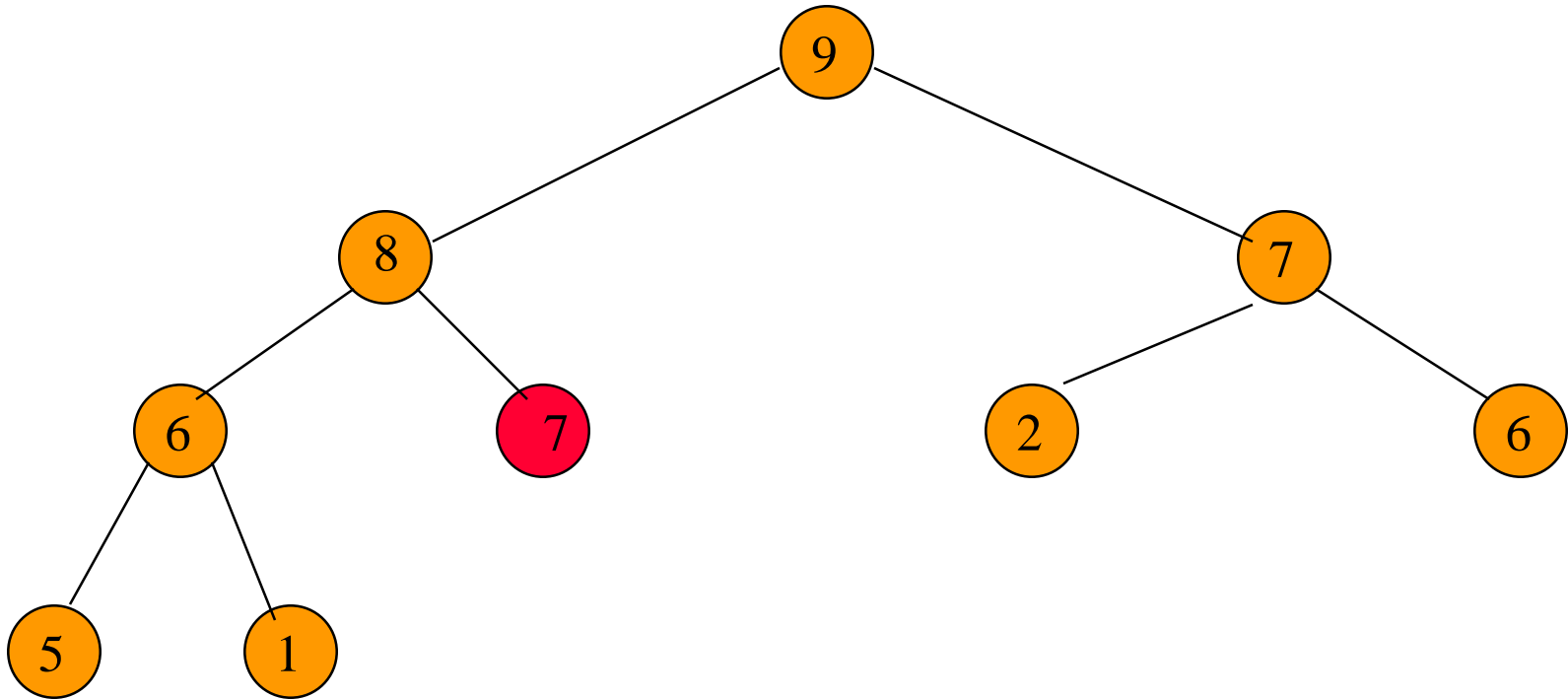
Reinsert **7**.

# Removing The Max Element



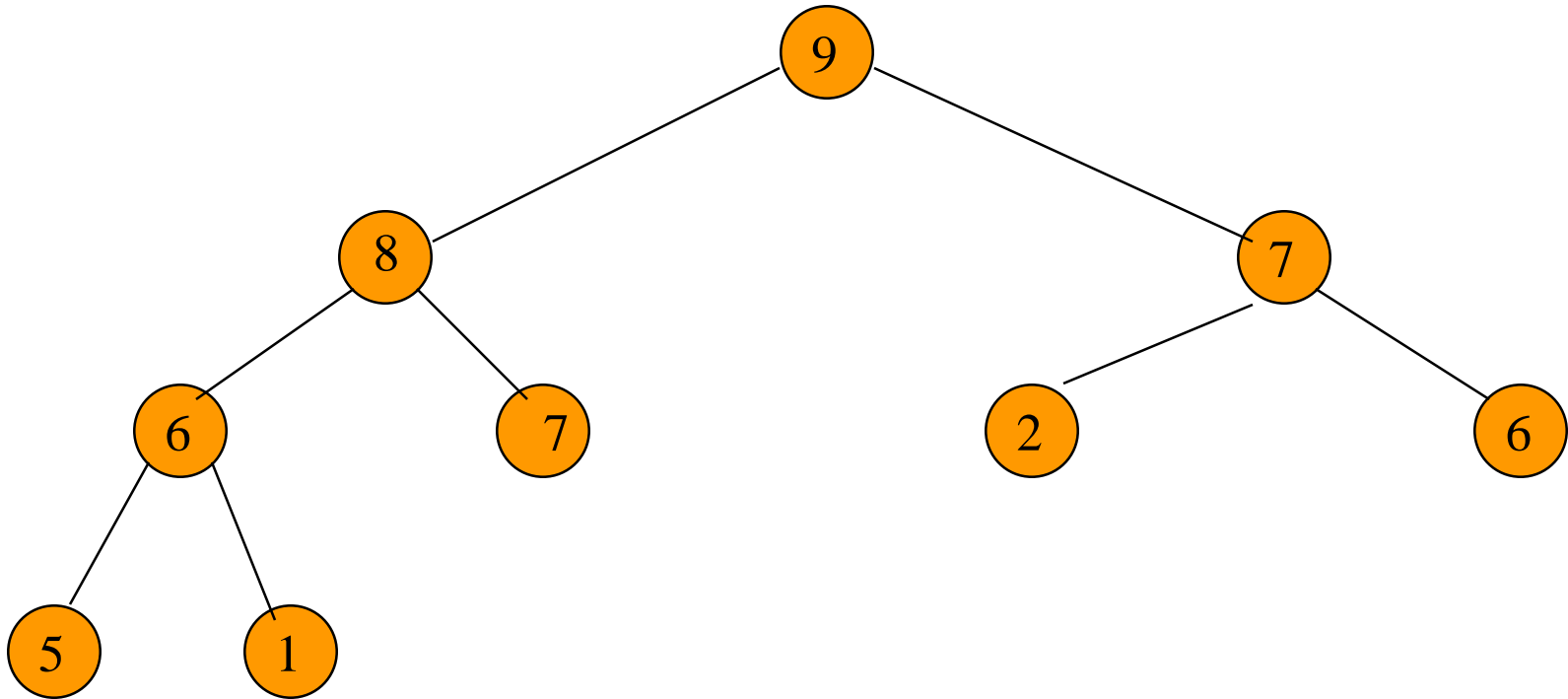
Reinsert **7**.

# Removing The Max Element



Reinsert **7**.

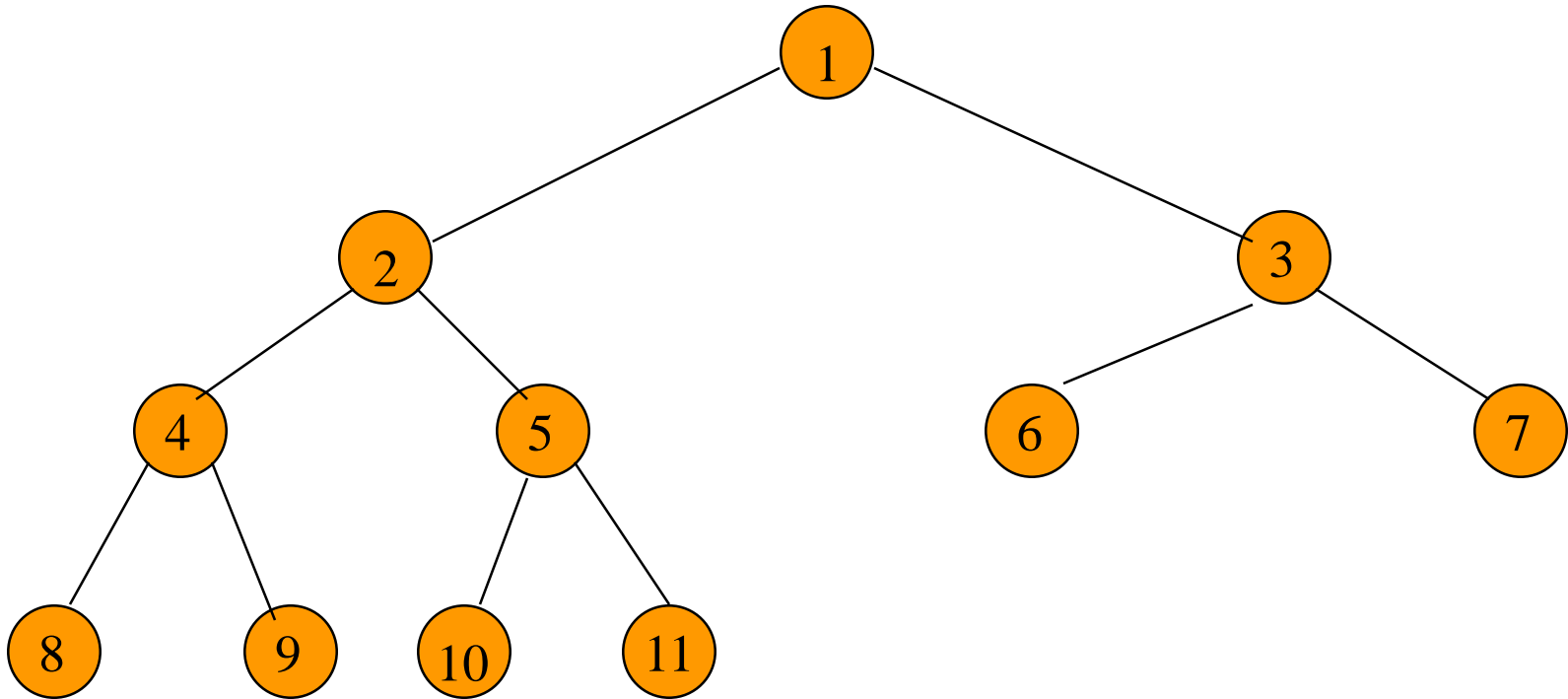
# Complexity Of Remove Max Element



Complexity is  $O(\log n)$ .

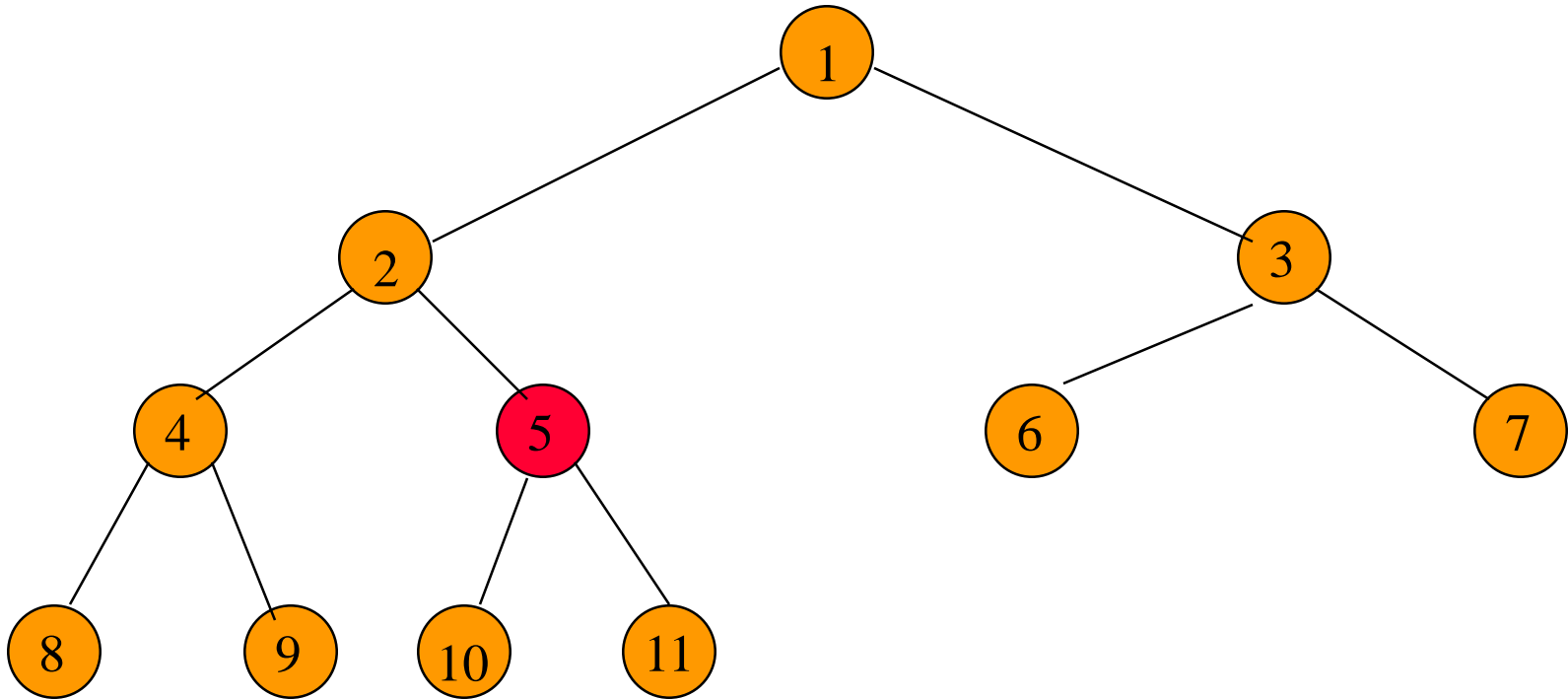


# Initializing A Max Heap



input array =  $[-, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11]$

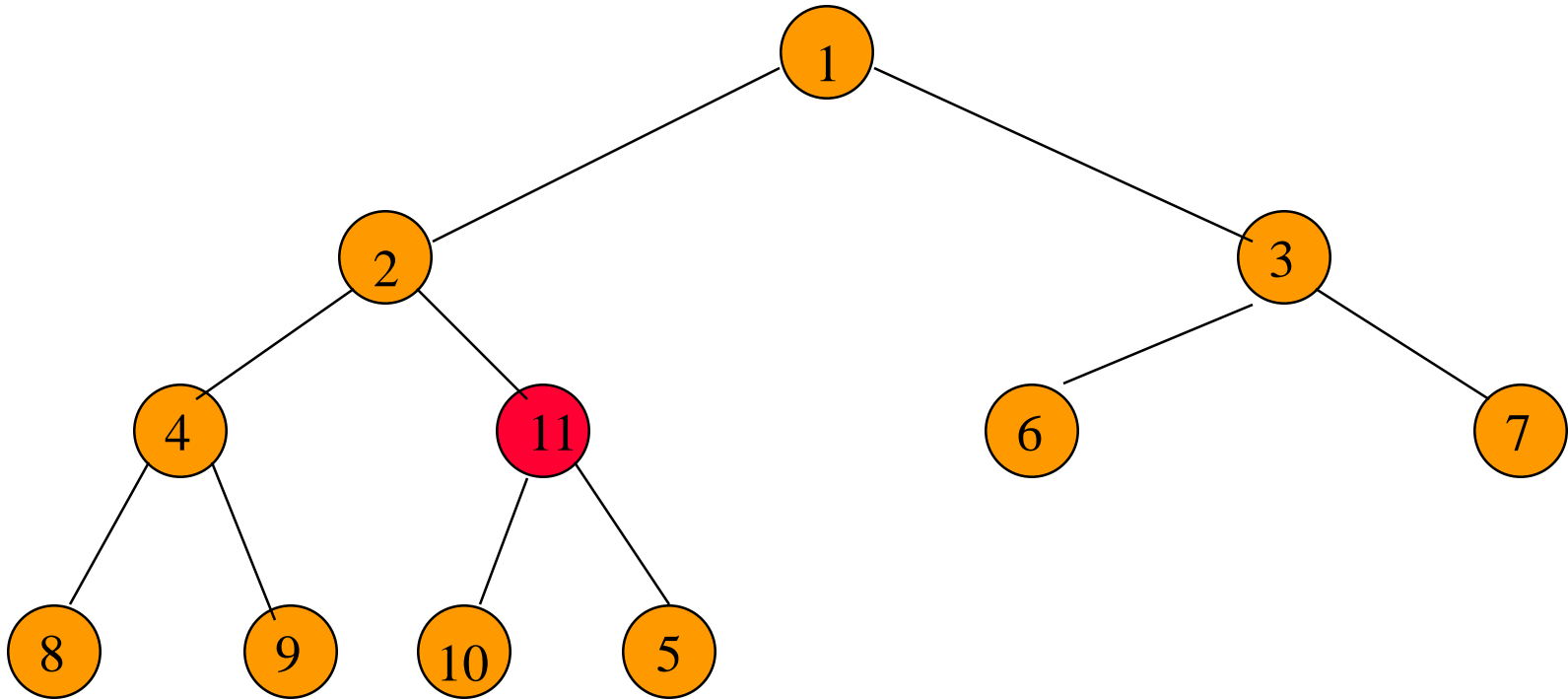
# Initializing A Max Heap



Start at rightmost array position that has a child.

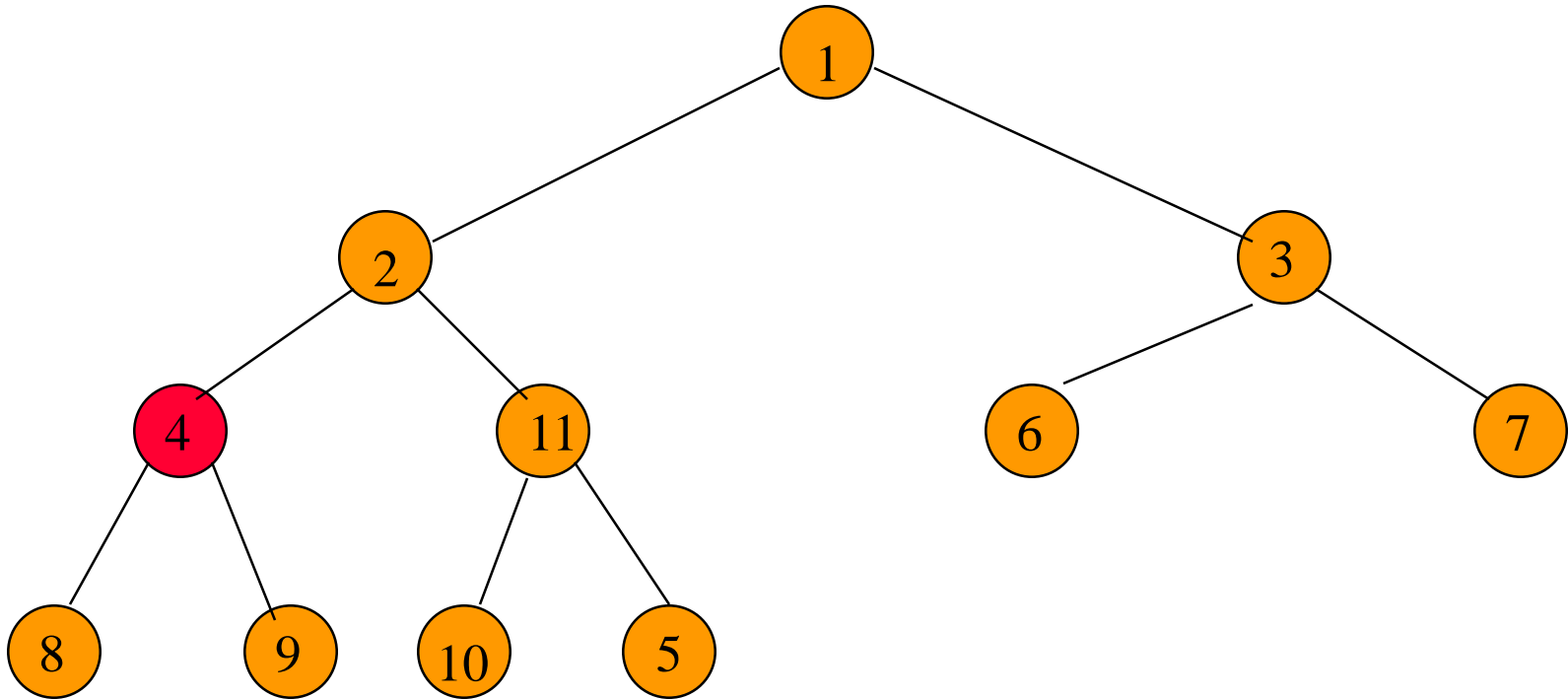
Index is  $n/2$ .

# Initializing A Max Heap

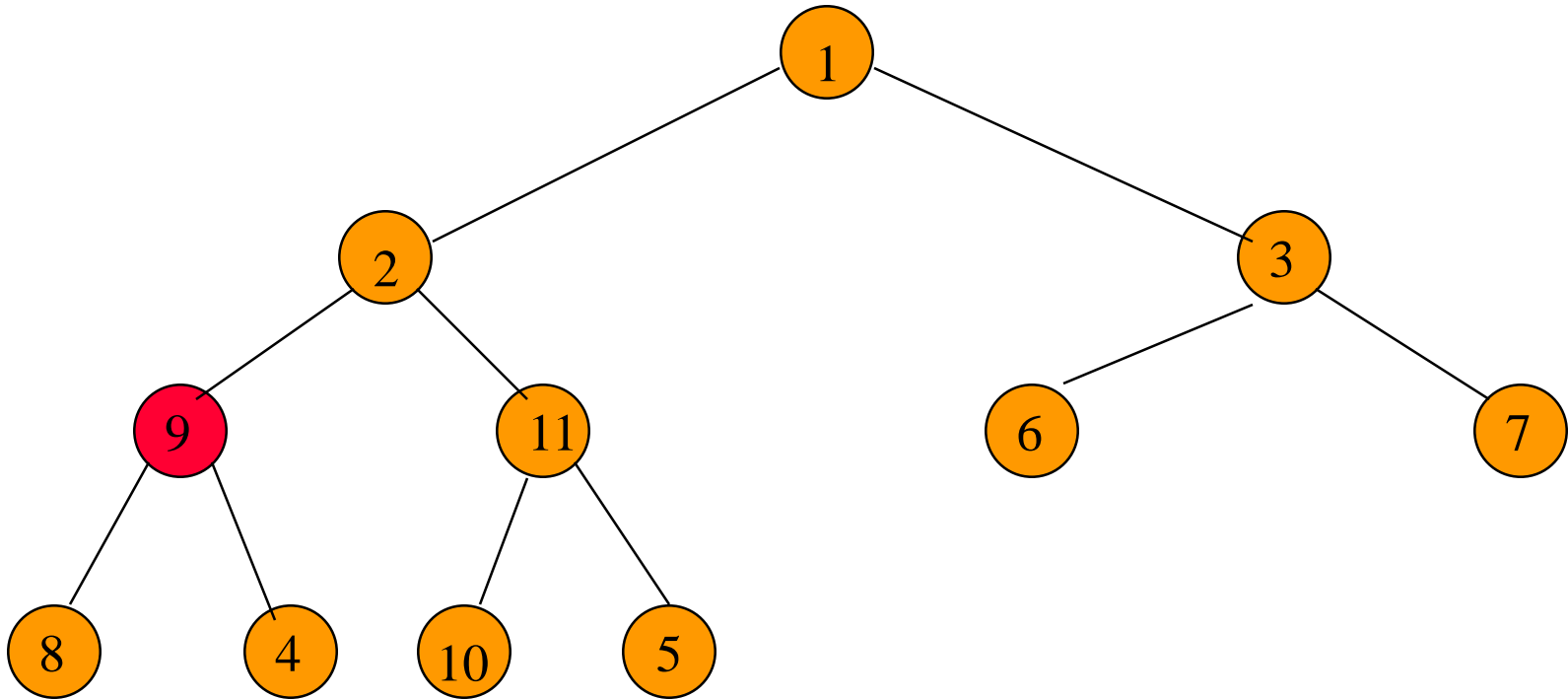


Move to next lower array position.

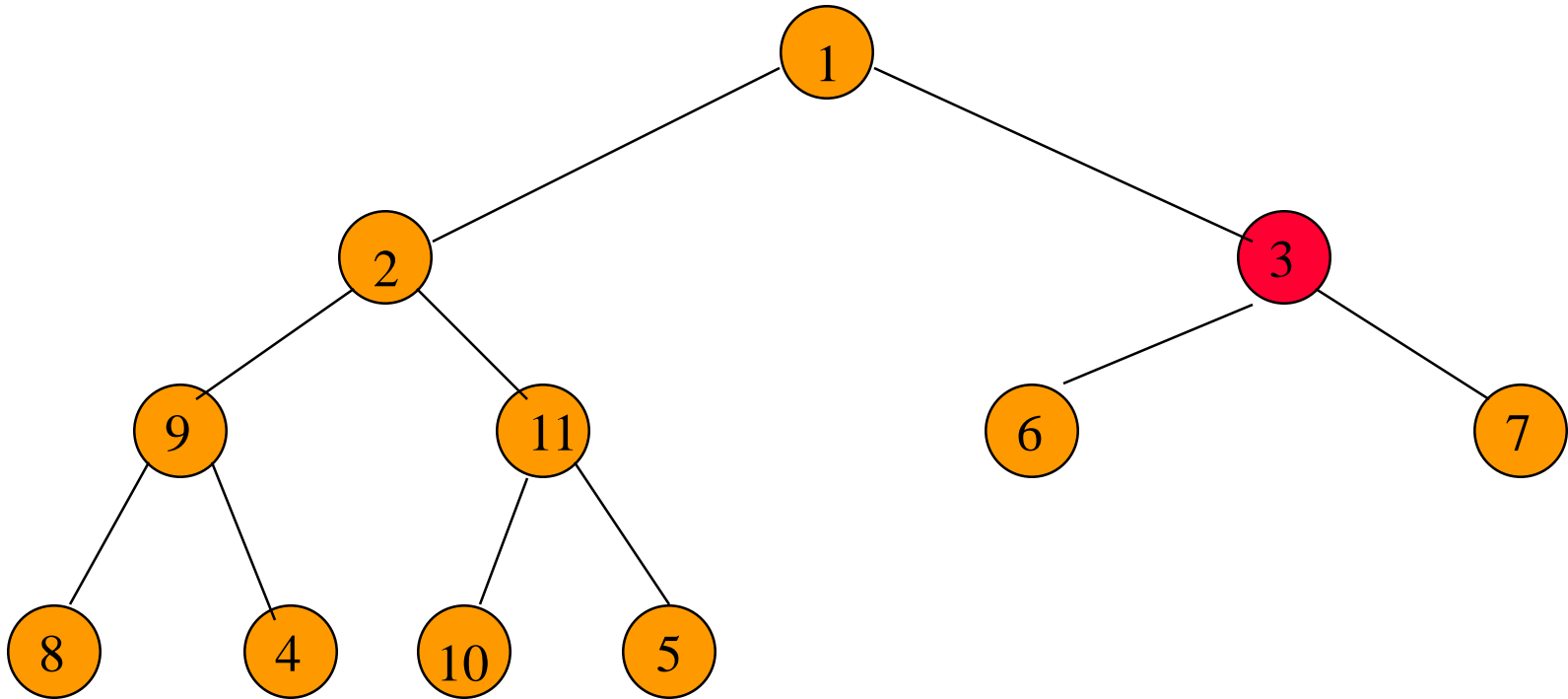
# Initializing A Max Heap



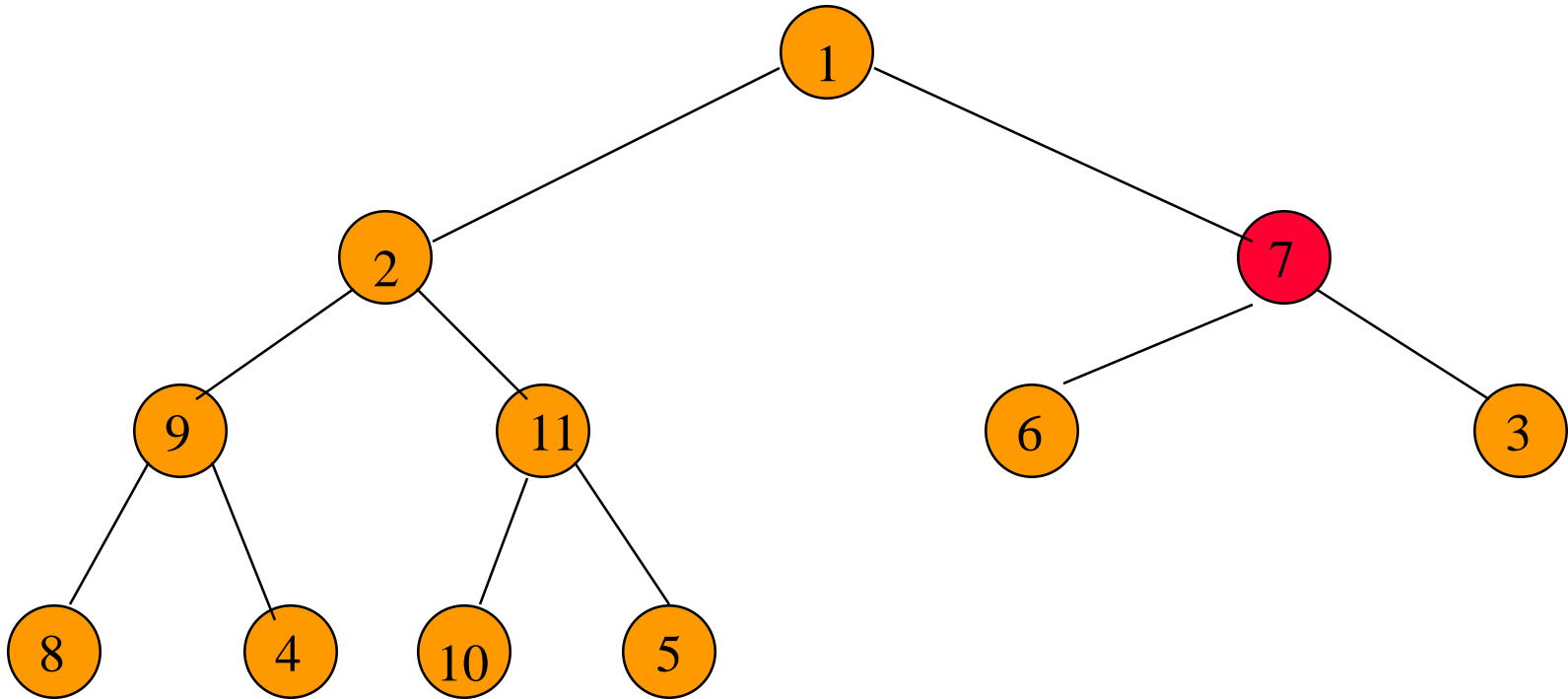
# Initializing A Max Heap



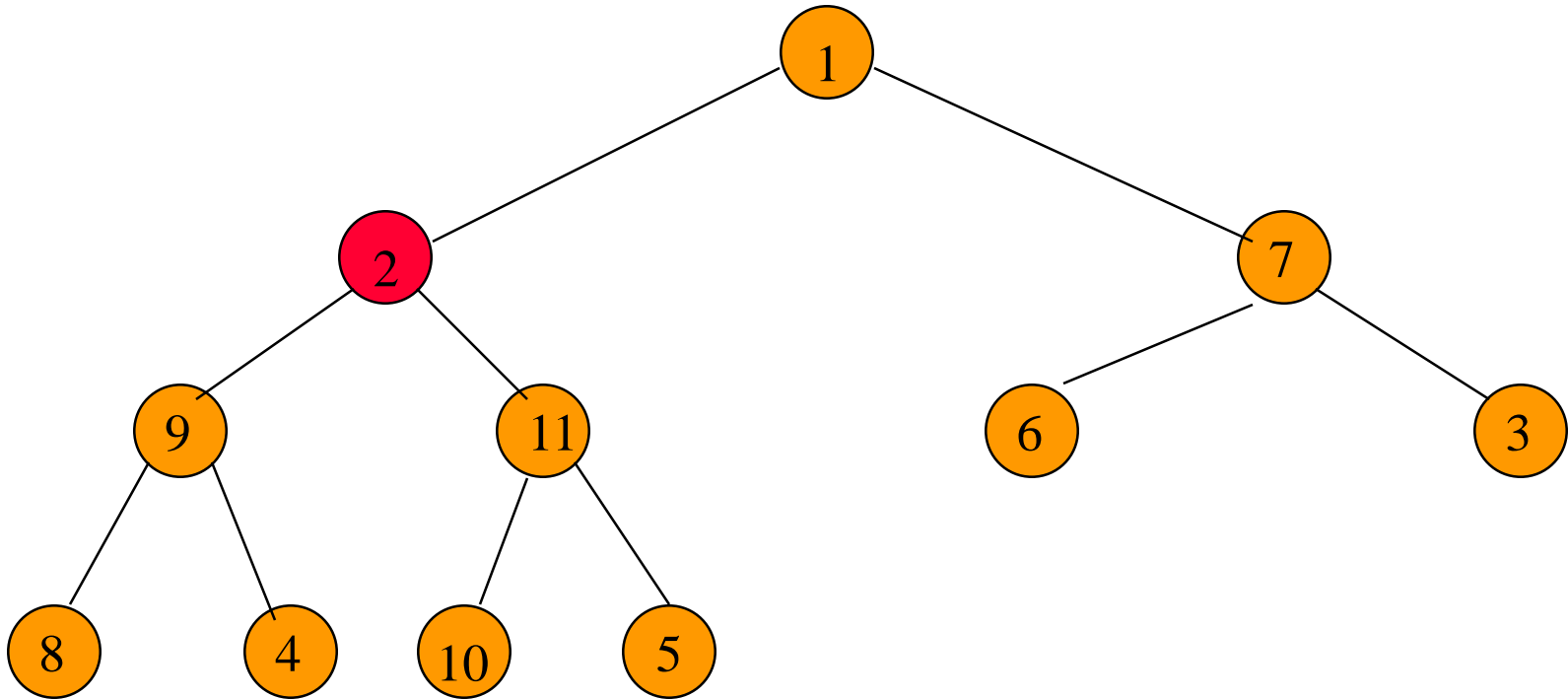
# Initializing A Max Heap



# Initializing A Max Heap

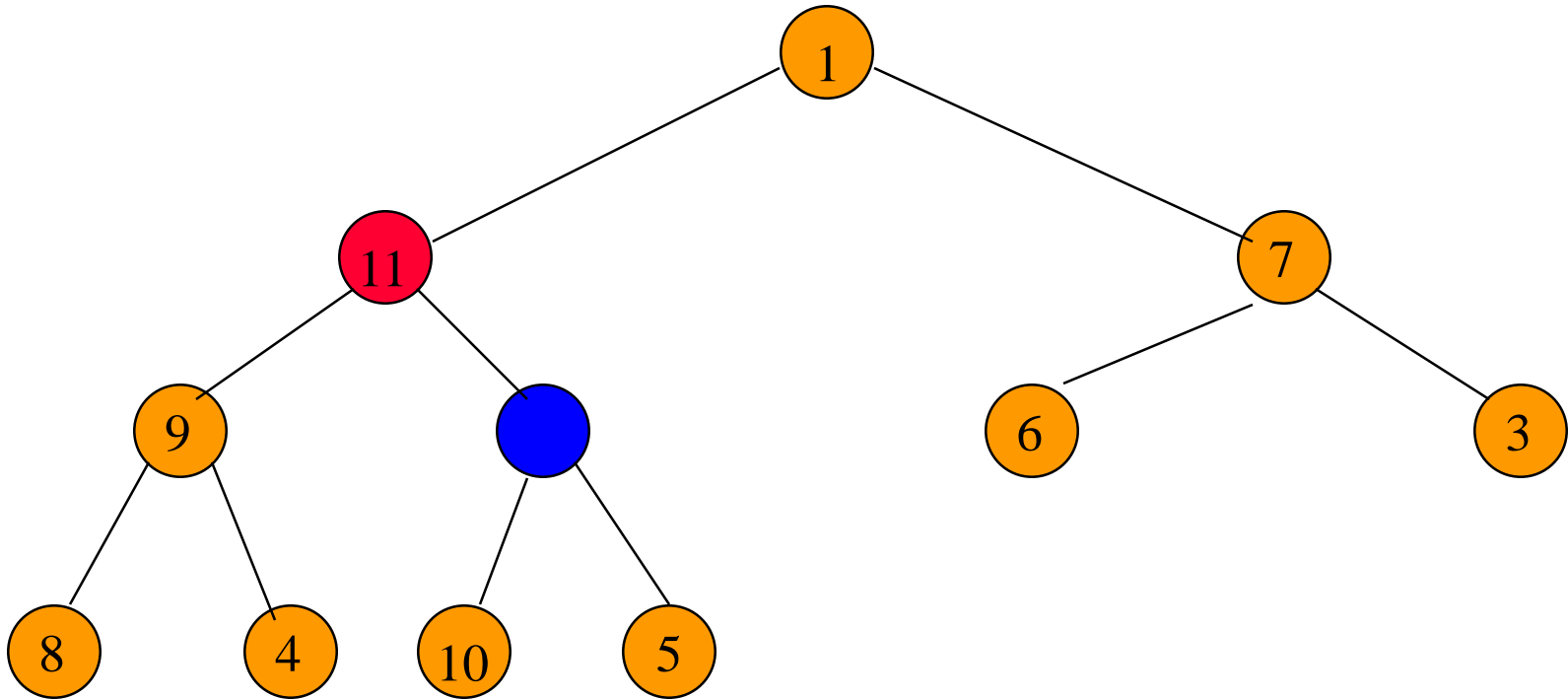


# Initializing A Max Heap



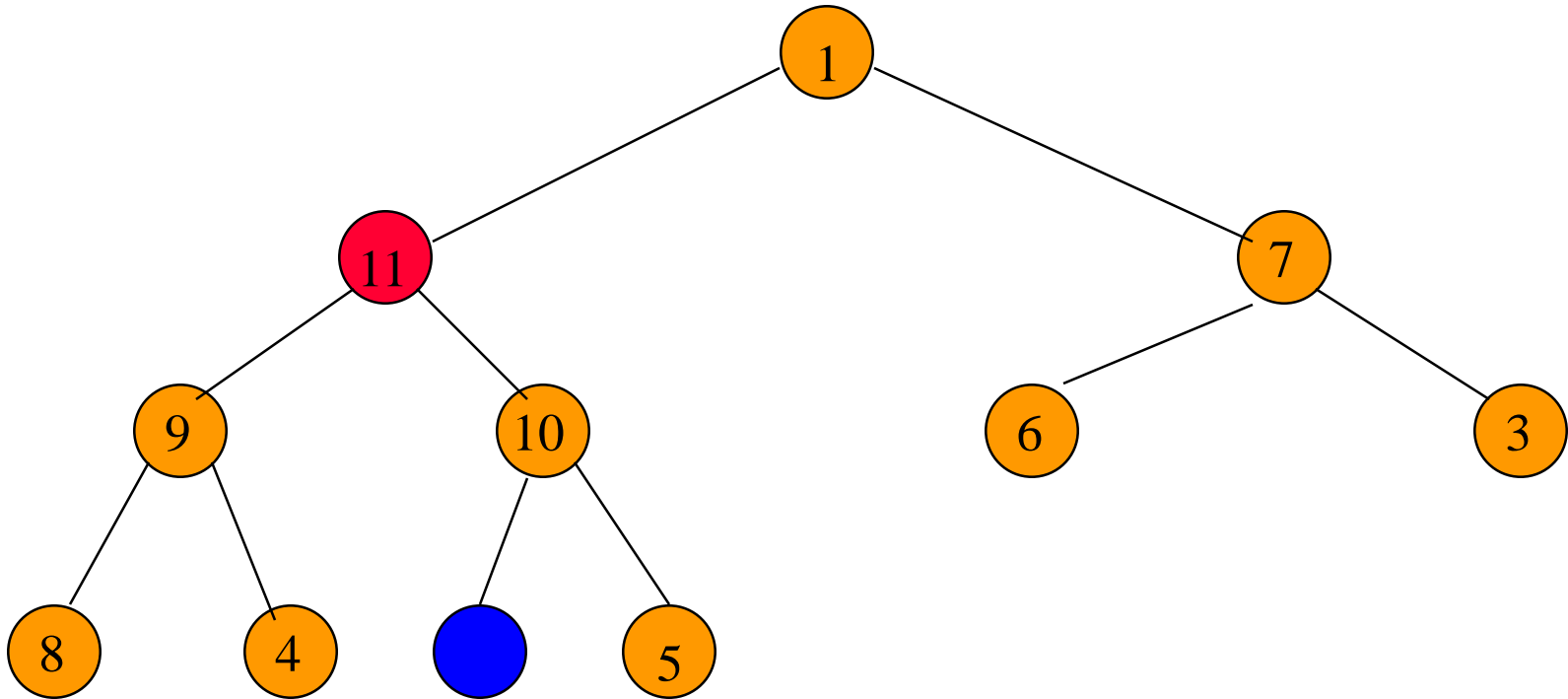


# Initializing A Max Heap



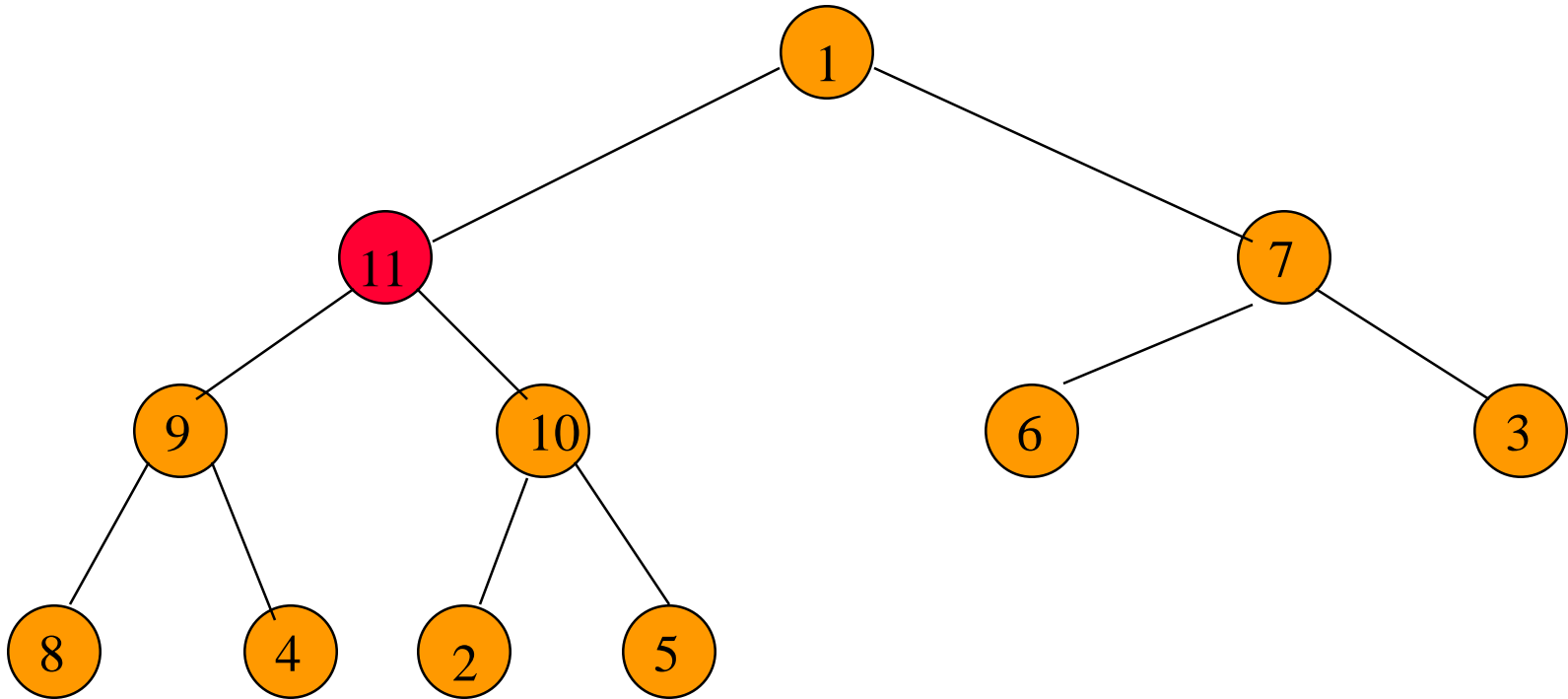
Find a home for 2.

# Initializing A Max Heap



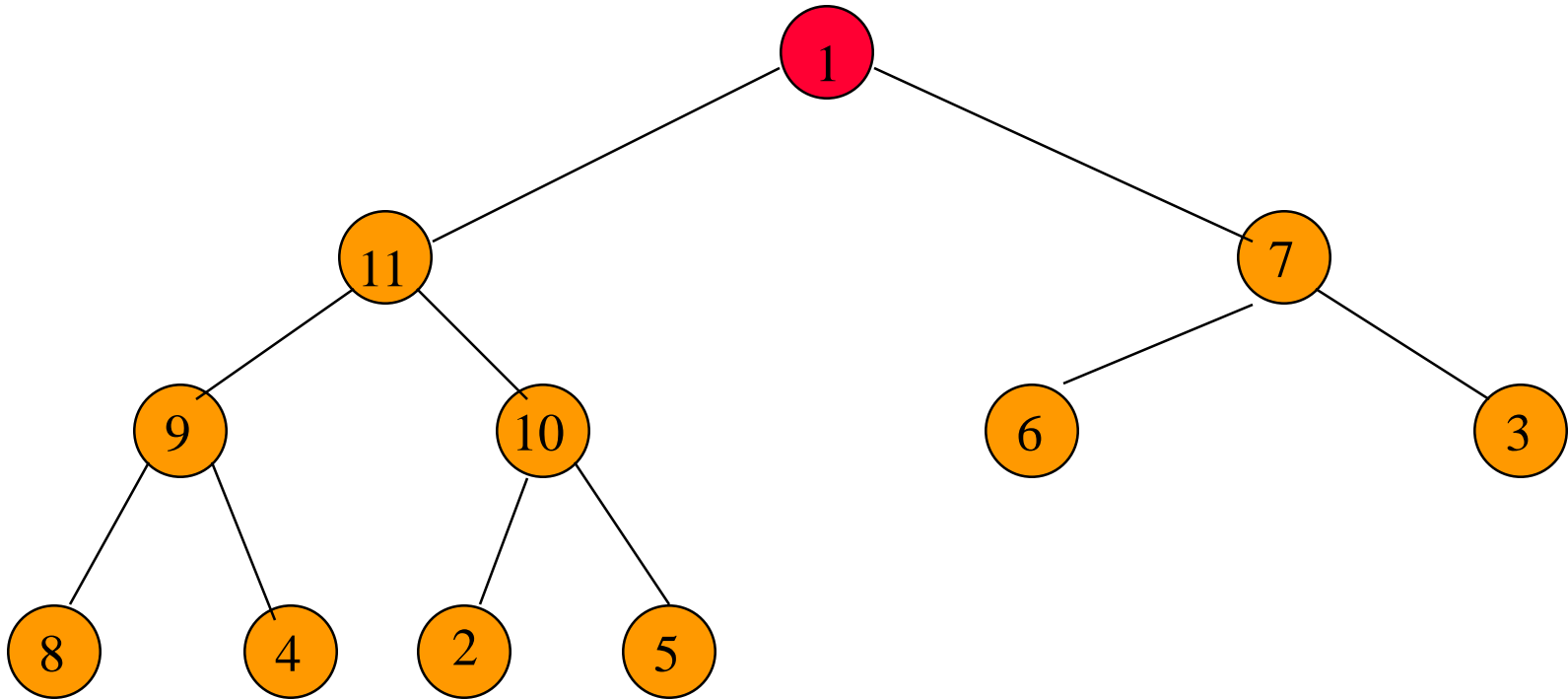
Find a home for 2.

# Initializing A Max Heap



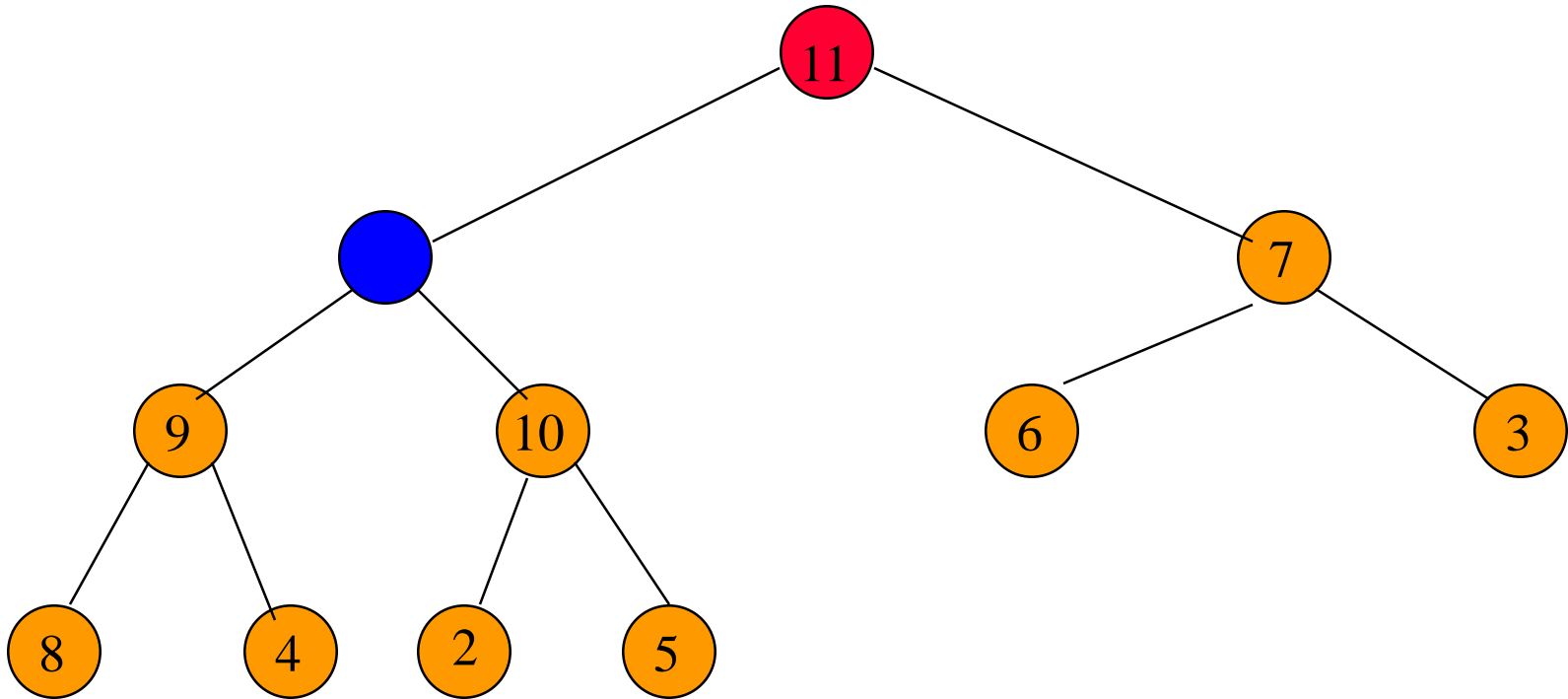
Done, move to next lower array position.

# Initializing A Max Heap



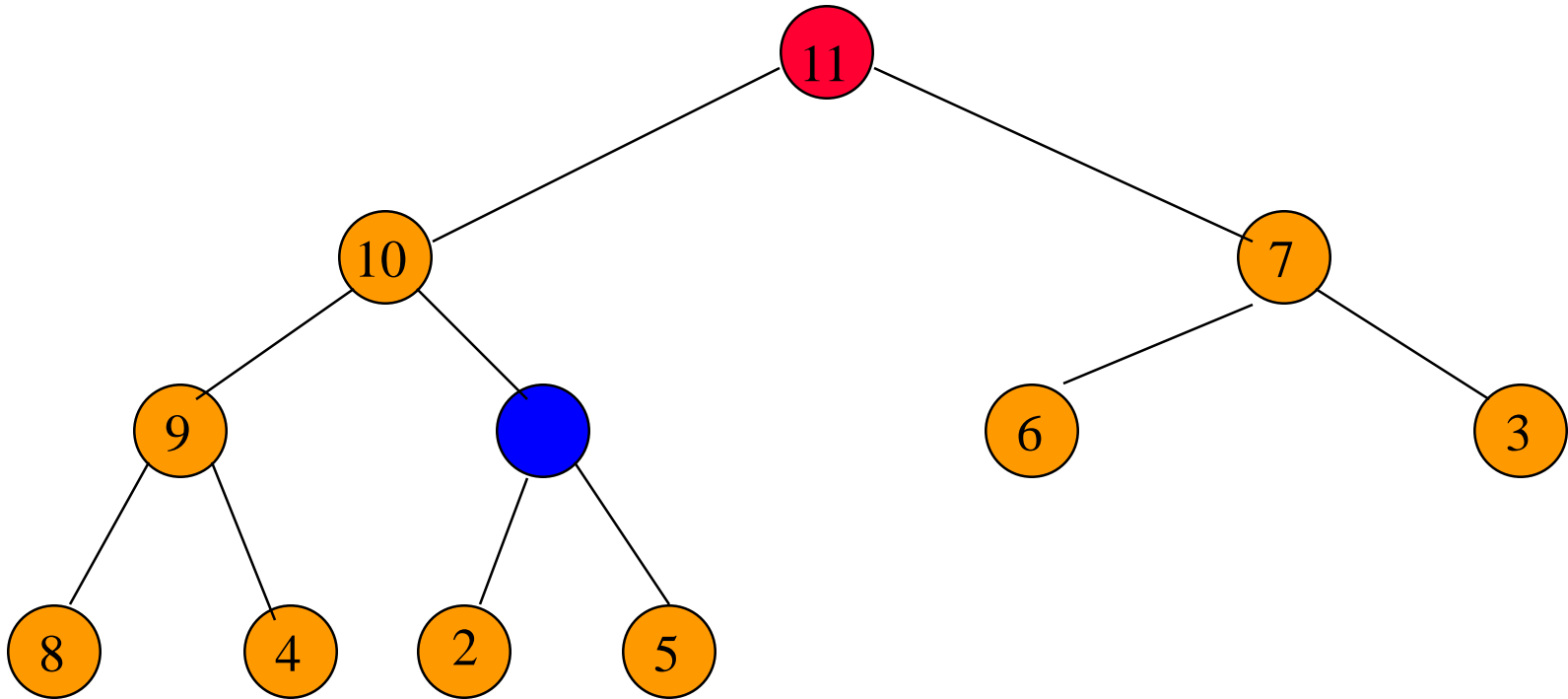
Find home for 1.

# Initializing A Max Heap



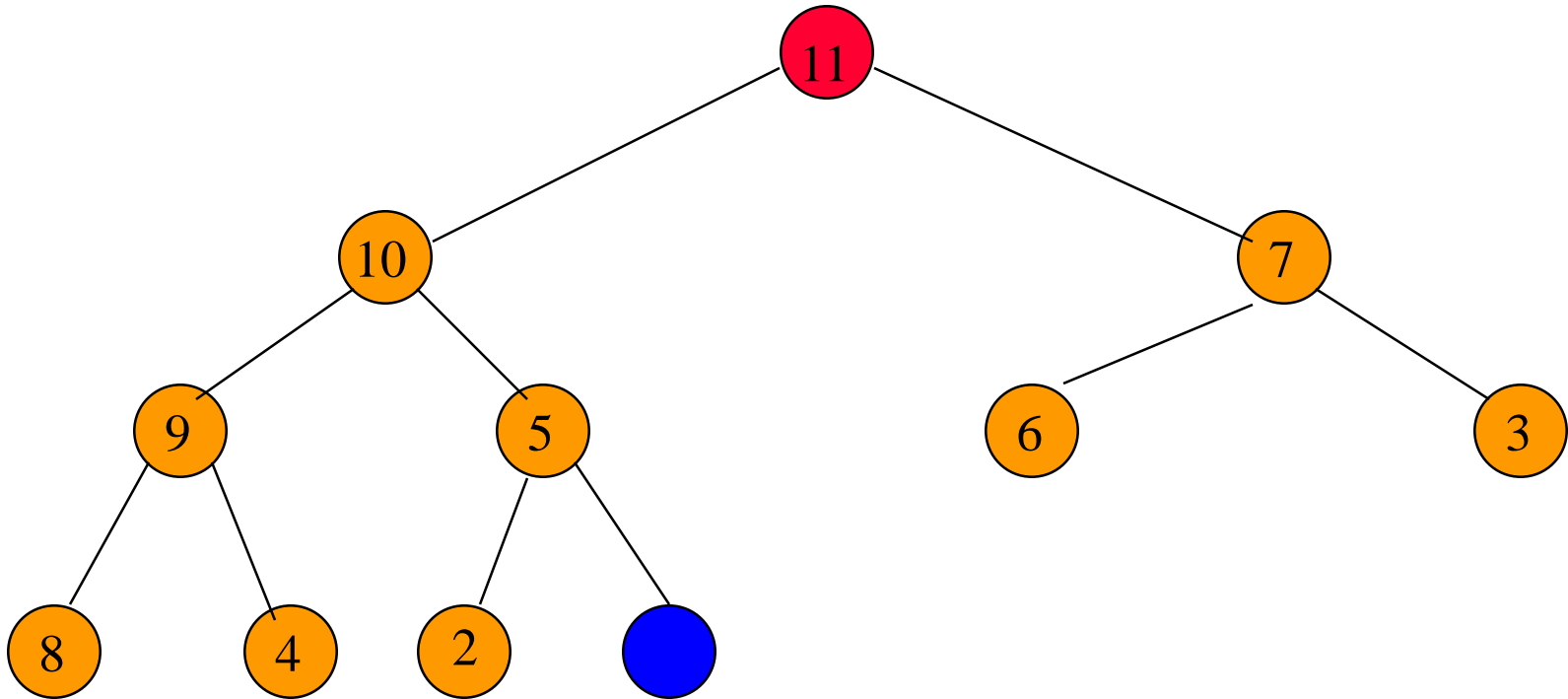
Find home for 1.

# Initializing A Max Heap



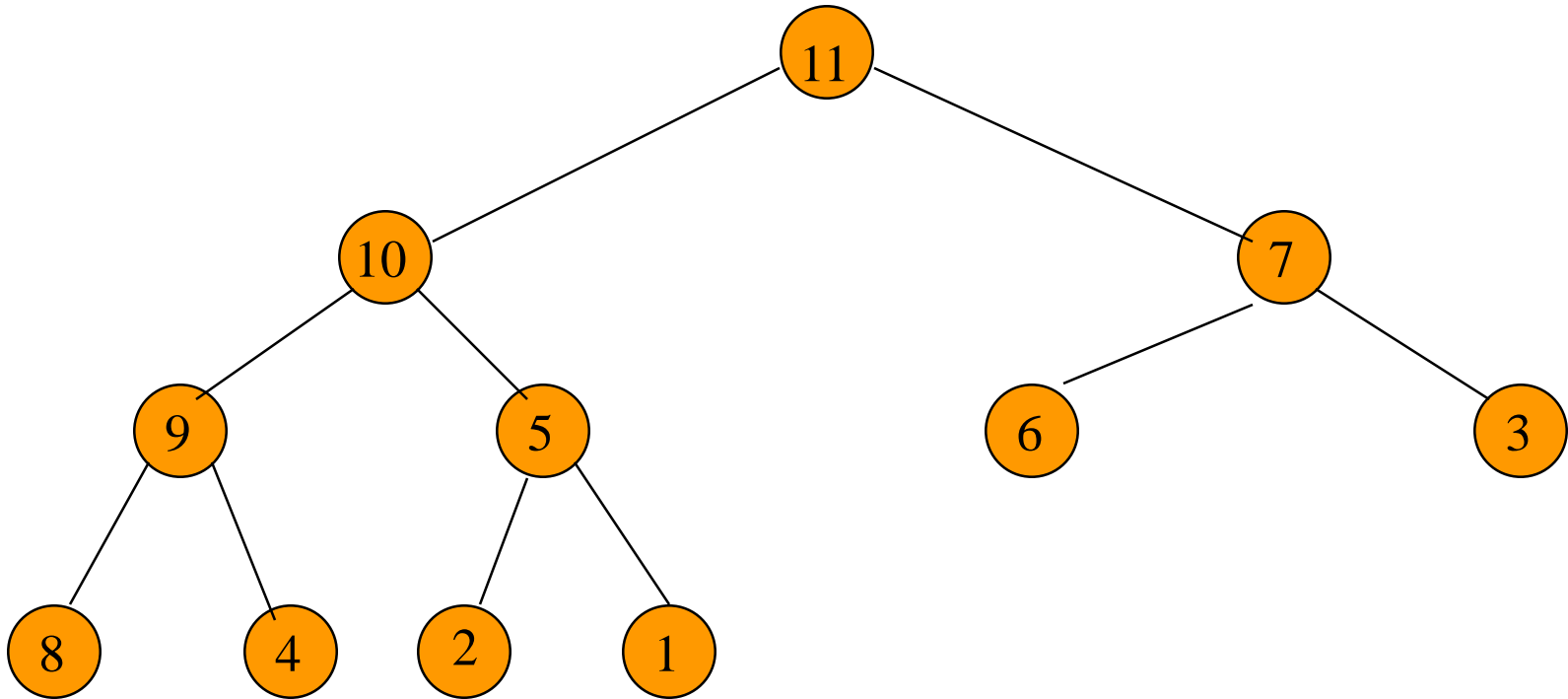
Find home for 1.

# Initializing A Max Heap



Find home for 1.

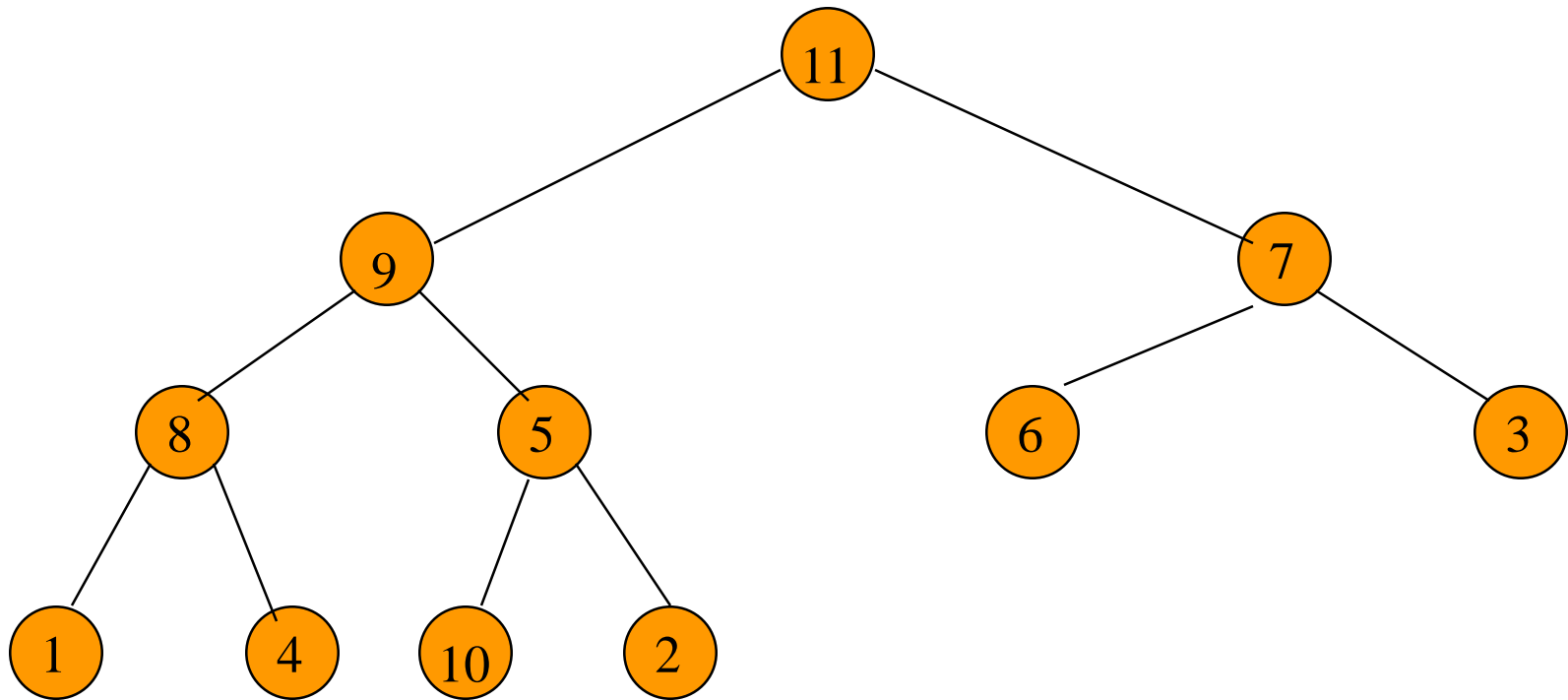
# Initializing A Max Heap



Done.



# Time Complexity



Height of heap =  $h$ .

Number of subtrees with root at level  $j$  is  $\leq 2^{j-1}$ .

Time for each subtree is  $O(h-j+1)$ .

# Complexity



Time for level  $j$  subtrees is  $\leq 2^{j-1}(h-j+1) = t(j)$ .

Total time is  $t(1) + t(2) + \dots + t(h-1) = O(n)$ .

# Huffman codes

# Fixed length codes

- Want to represent data as a sequence of 0's and 1's

For example: BACADAEAFABBAAGA

- A-000 B-001 C-010 D-011 E-100 F-101 G-110 H-111

001 000 010 000 011 000 100 000 101 000 001 001 000 000  
000 110 000 111

- This is a fixed length code.
- Can we make the sequence of 0's and 1's shorter ?

# Variable length codes

- A 0      B 100    C 1010    D 1011

E 1100 F 1101 G 1110    H 1111

100010100101101100011010100100000111001111

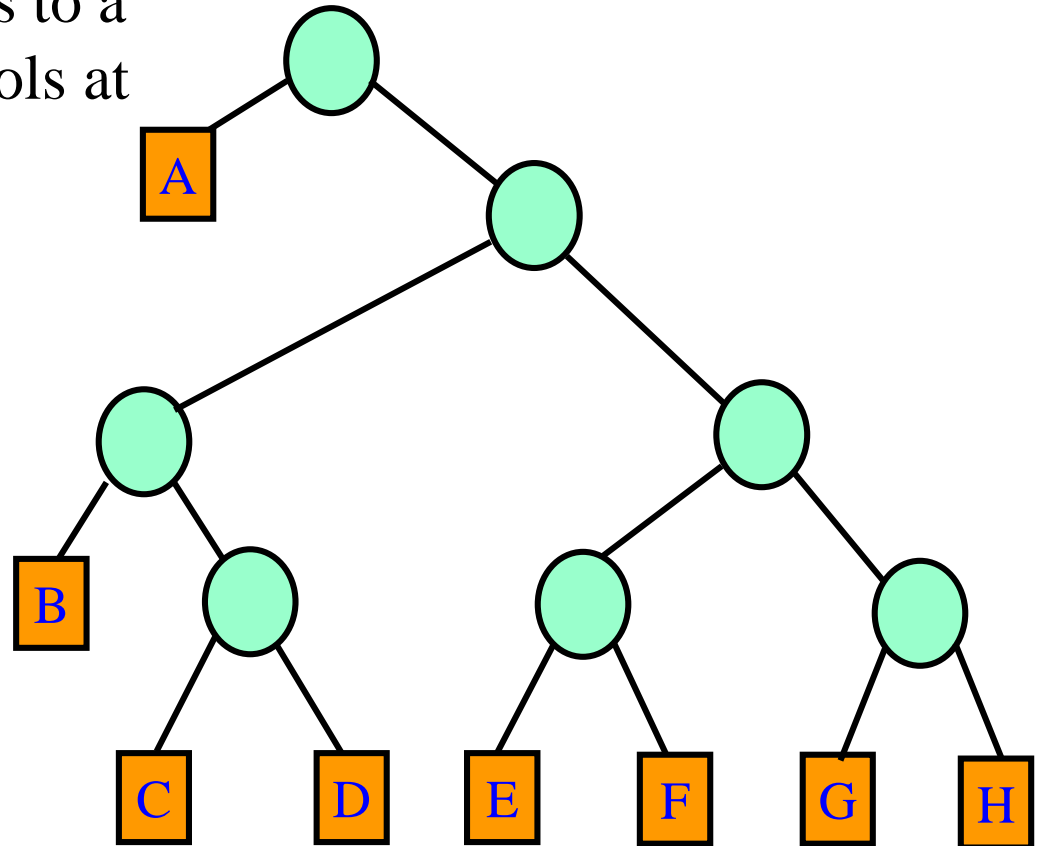
- This is a variable length code.
- How do we decode ?
- Use prefix codes: No codeword is a prefix of the other

# Representing prefix codes

A 0      B 100   C 1010   D 1011

E 1100   F 1101   G 1110   H 1111

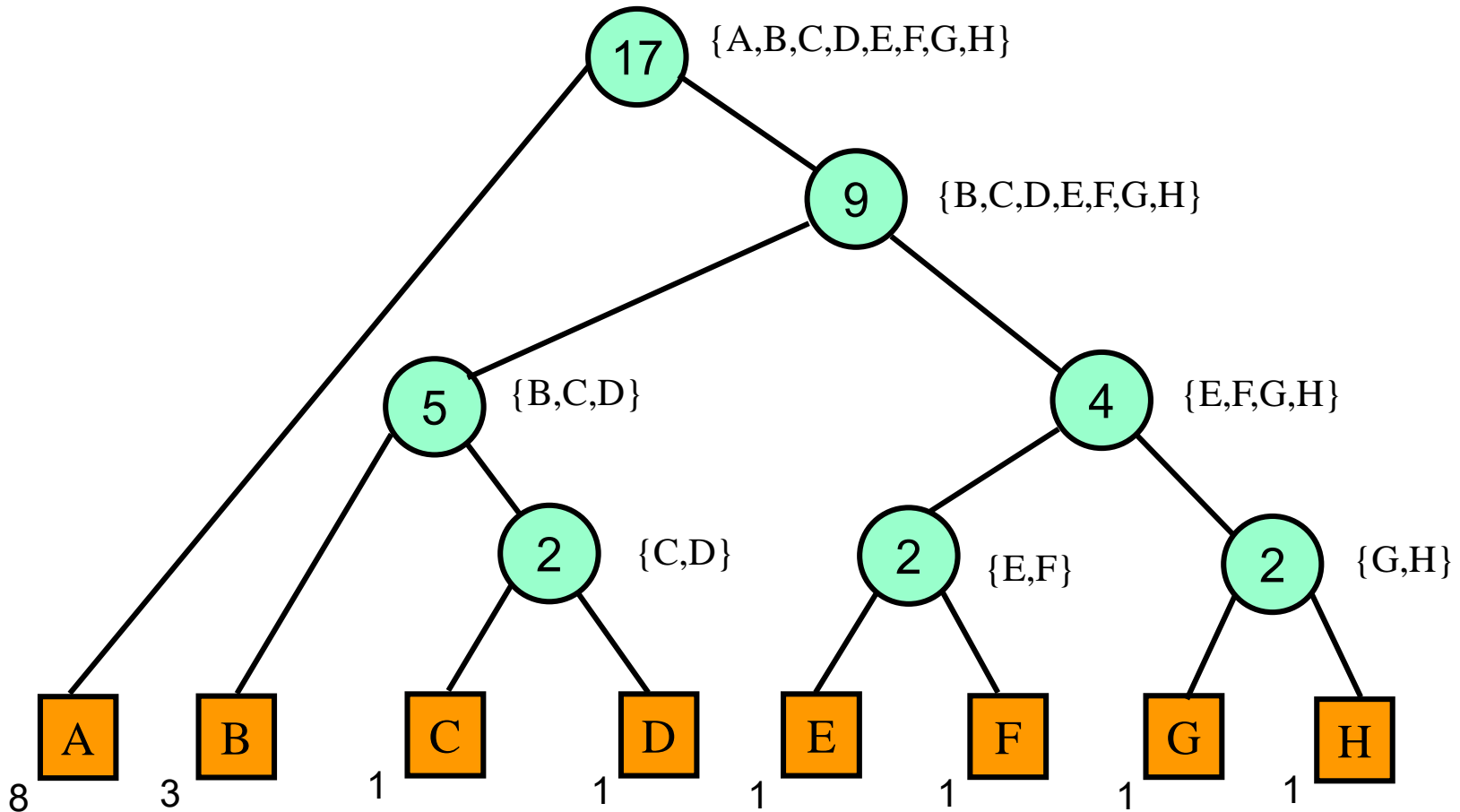
- A prefix code corresponds to a binary tree with the symbols at the **leaves** and vice versa.



# Construction Algorithm

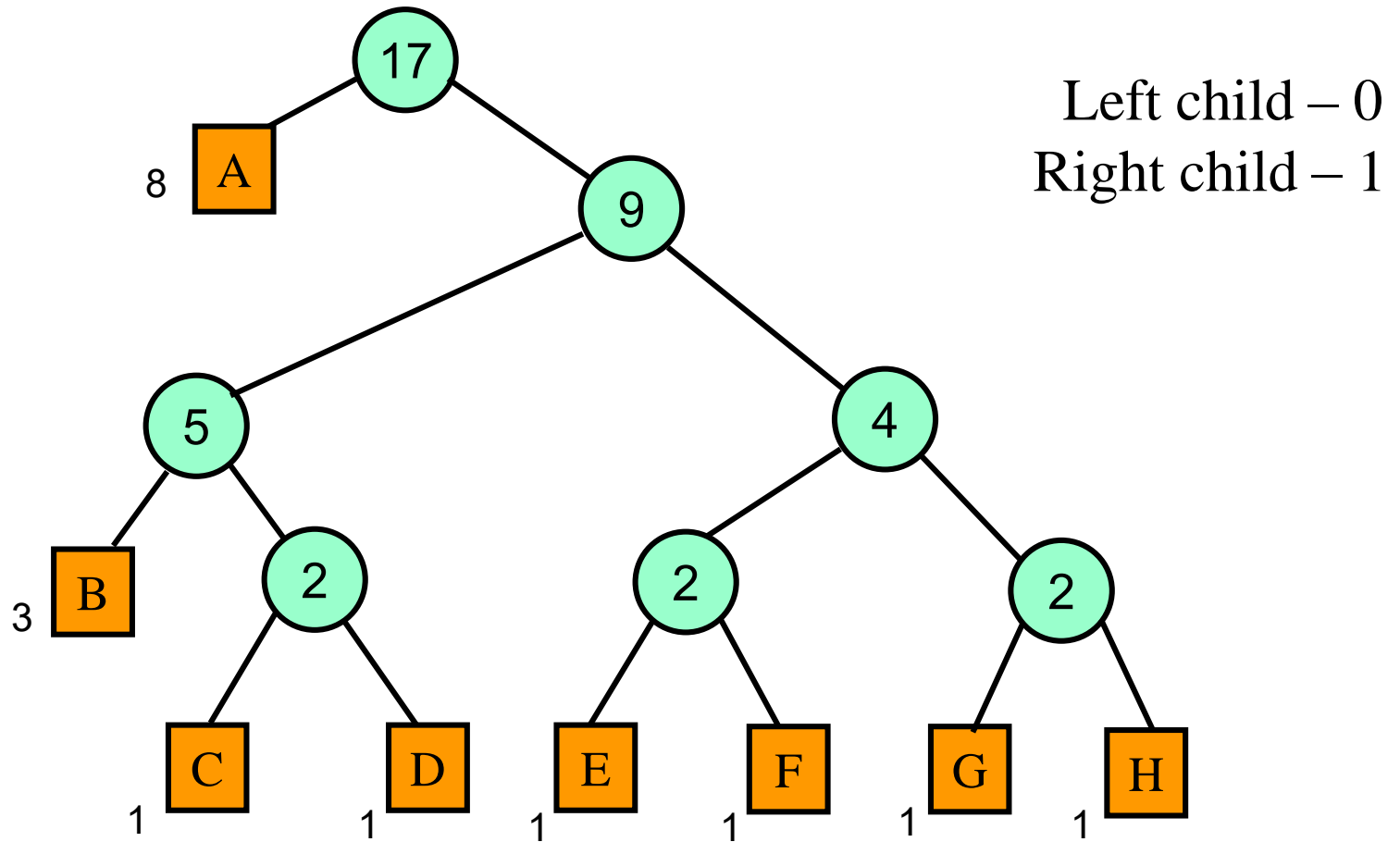
- Compute the character frequencies
- Insert the frequencies into a min heap
- Repeat
  - Delete the two minimum elements,  $f_1$  and  $f_2$  from the heap
  - insert  $f_1+f_2$  into the heap
- Complexity:  $O(n \log n)$ .

# Construction of Huffman tree





# Representation



$\text{code}(D) = 1011$ ;  $\text{code}(F) = 1101$