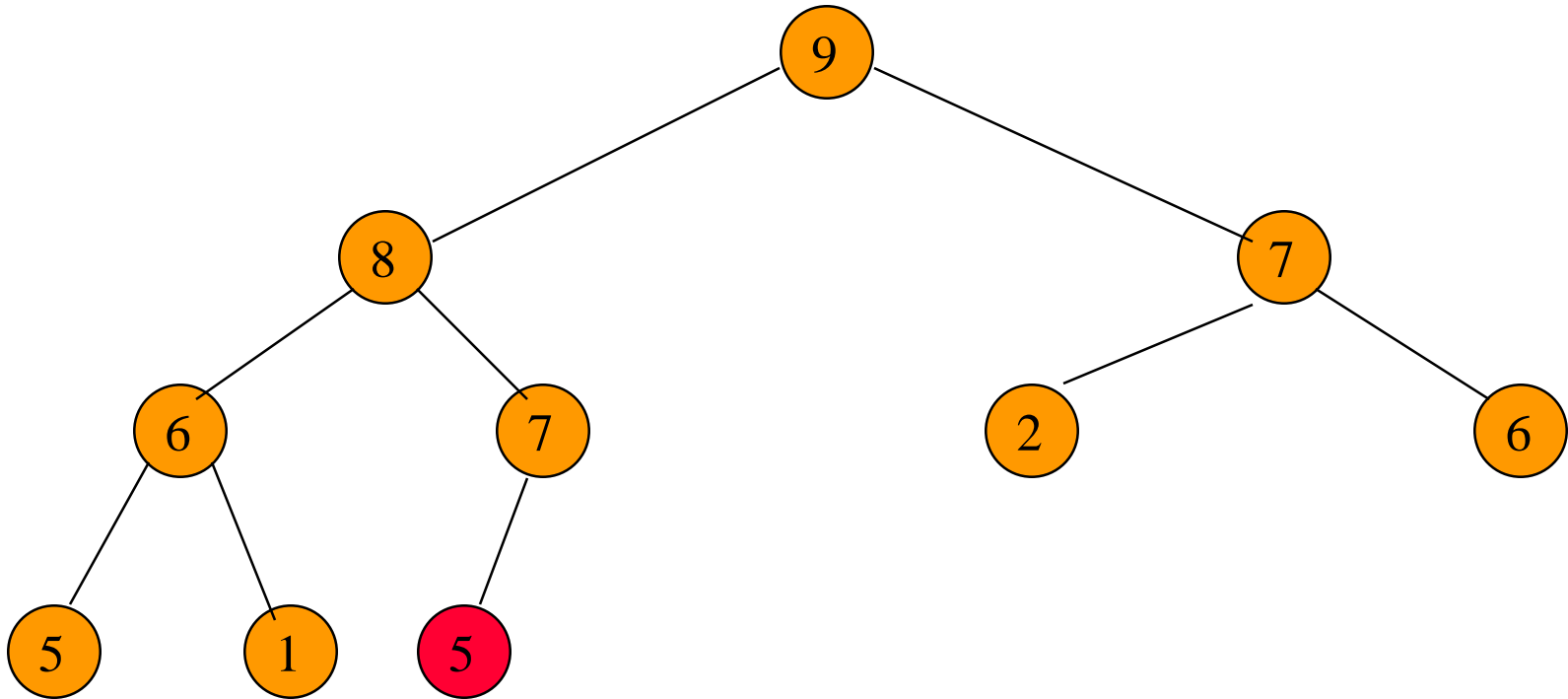


Heap Operations

Data structures

Fall 2018

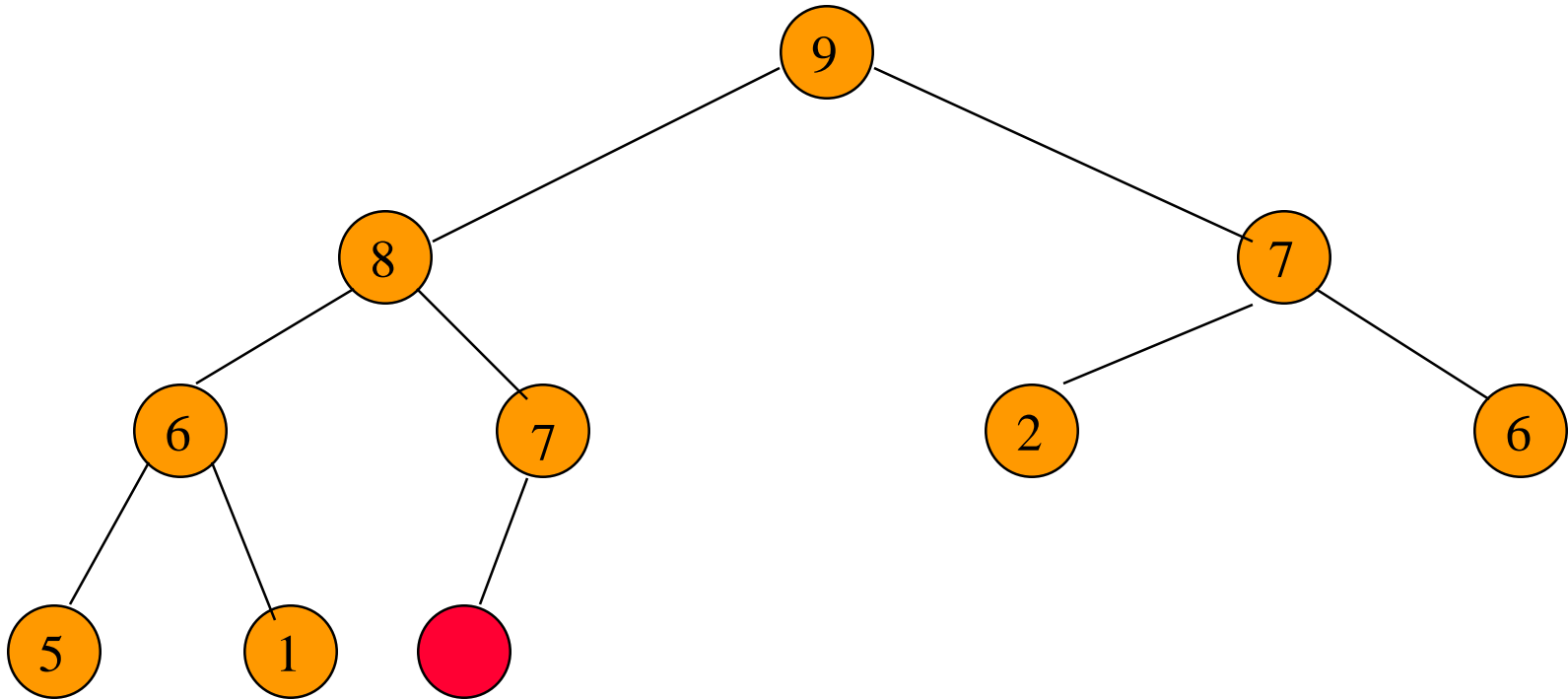
Putting An Element Into A Max Heap



Complete binary tree with 10 nodes.

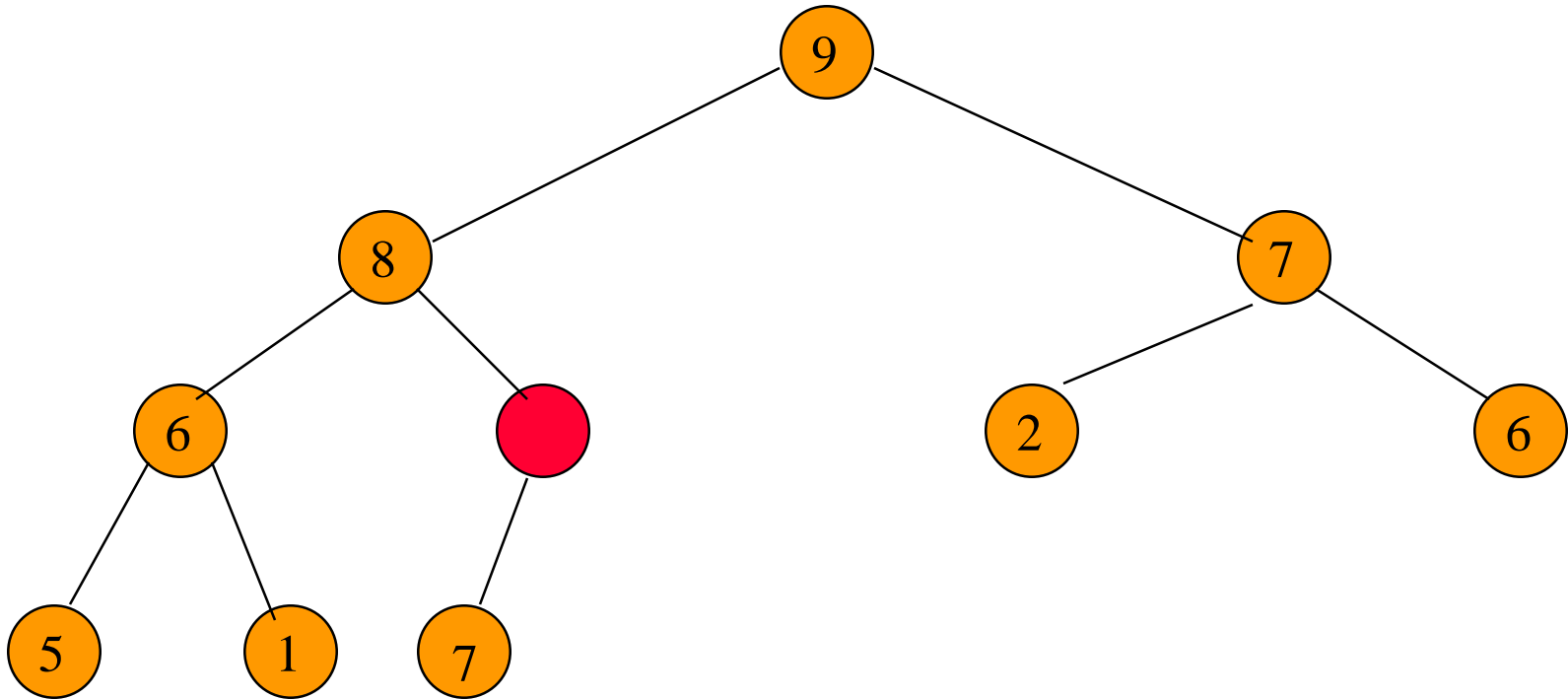
New element is 5.

Putting An Element Into A Max Heap



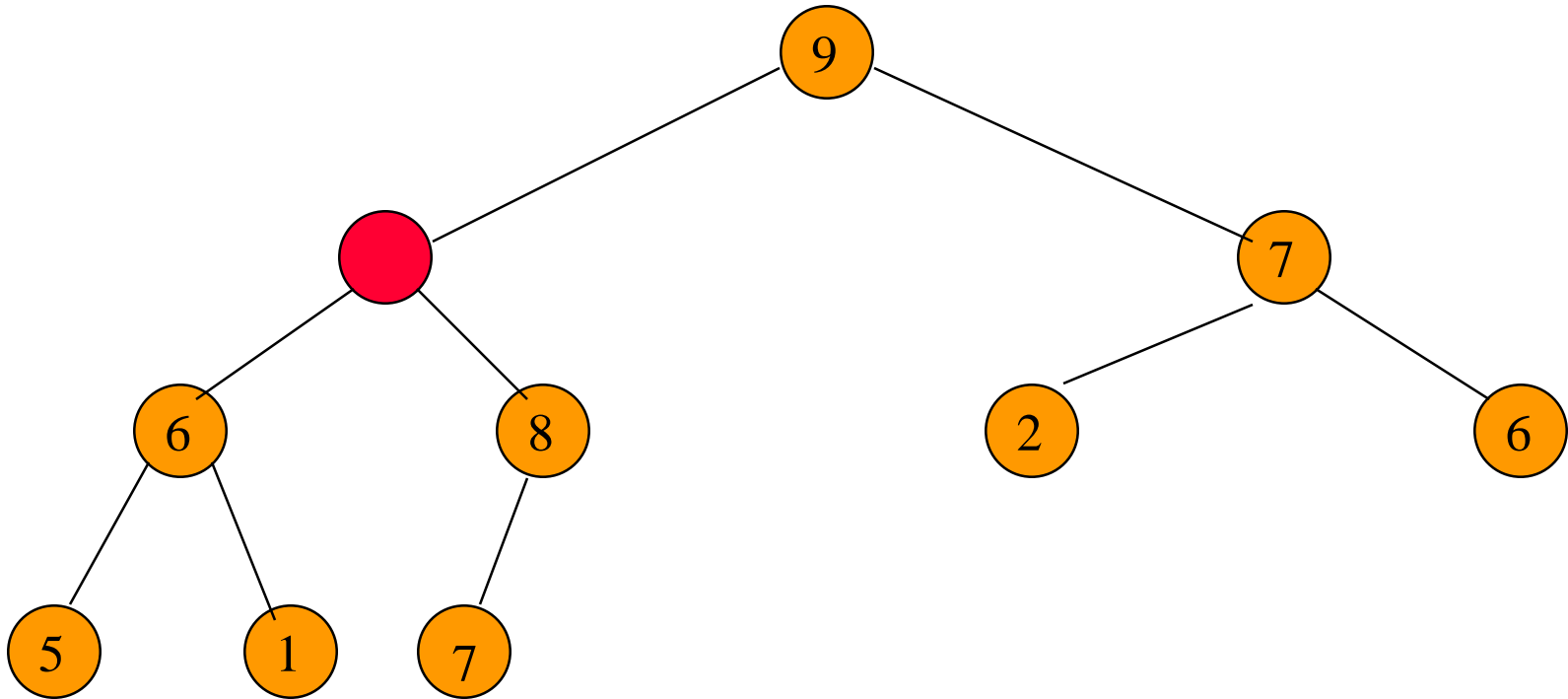
New element is 20.

Putting An Element Into A Max Heap



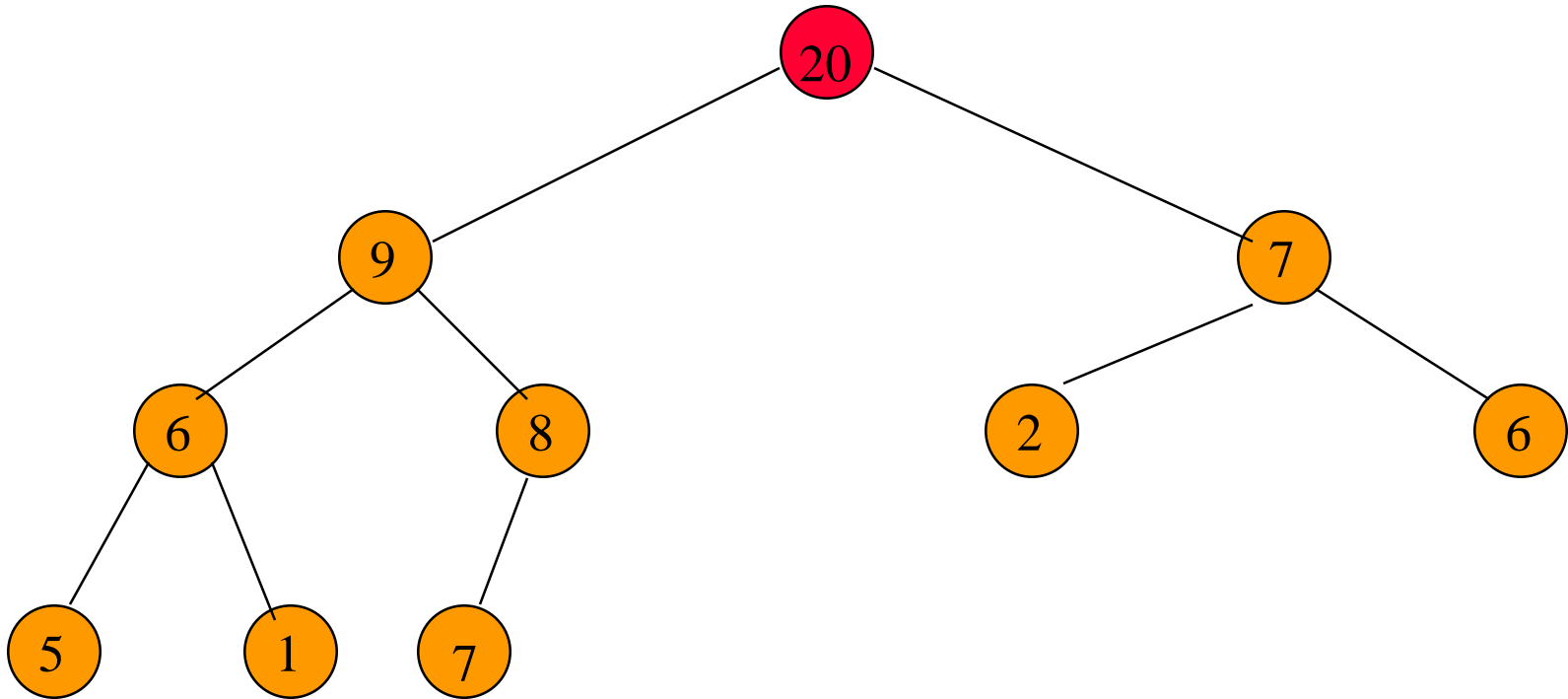
New element is 20.

Putting An Element Into A Max Heap



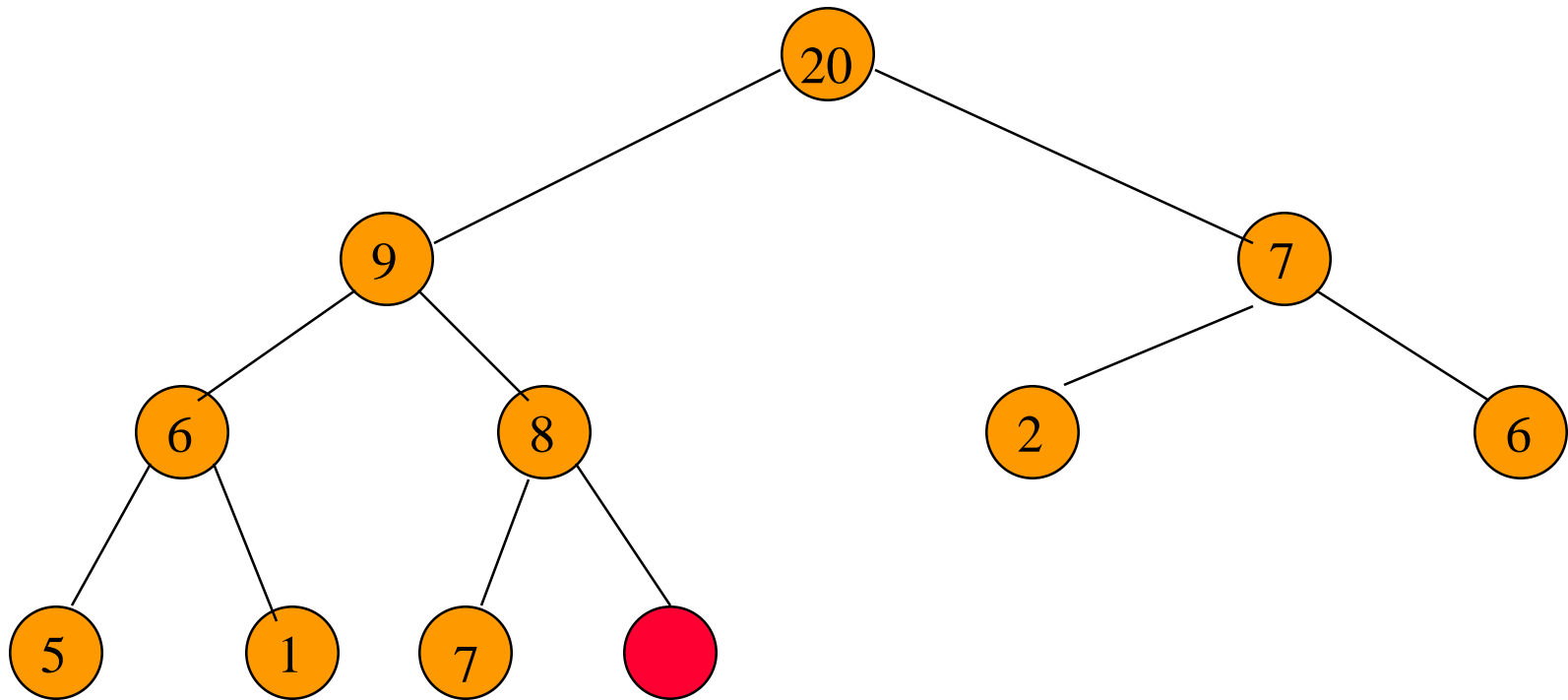
New element is 20.

Putting An Element Into A Max Heap



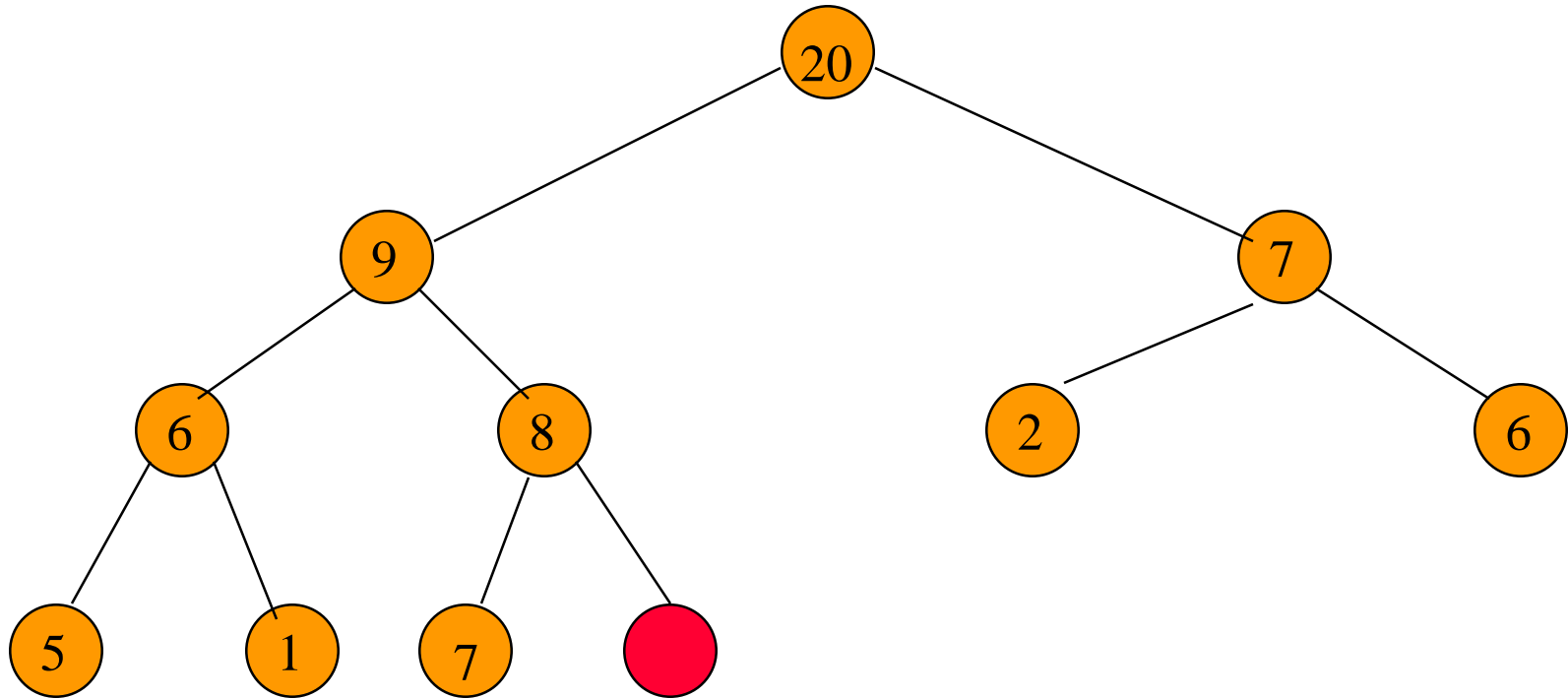
New element is 20.

Putting An Element Into A Max Heap



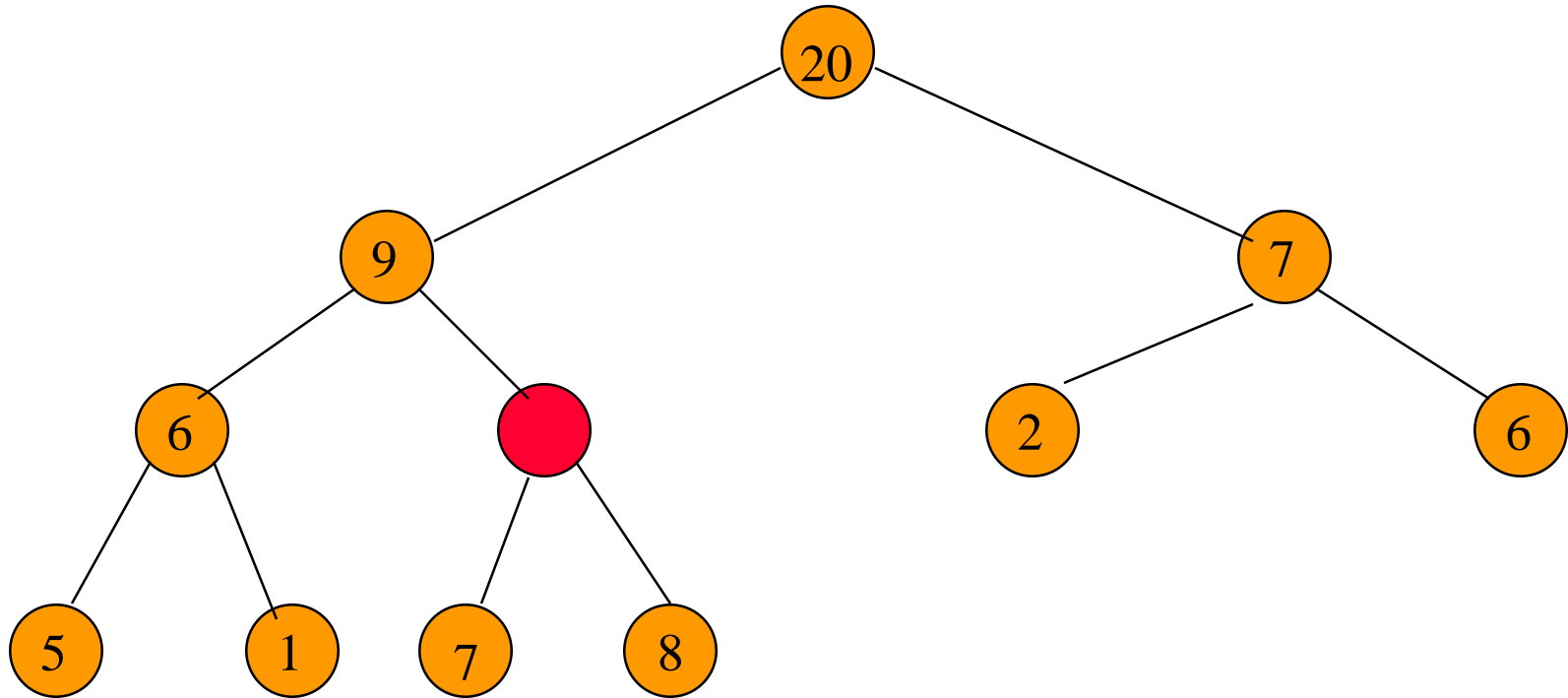
Complete binary tree with **11** nodes.

Putting An Element Into A Max Heap



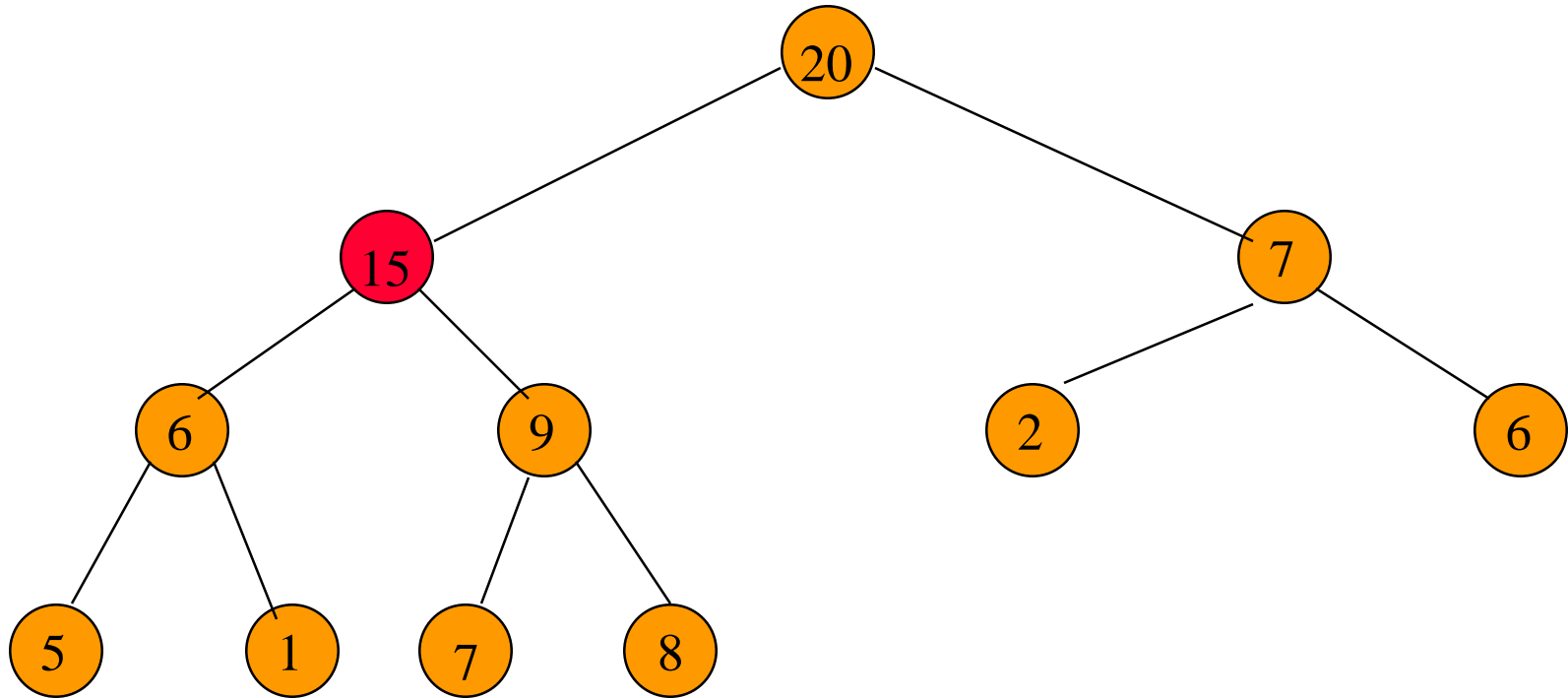
New element is 15.

Putting An Element Into A Max Heap



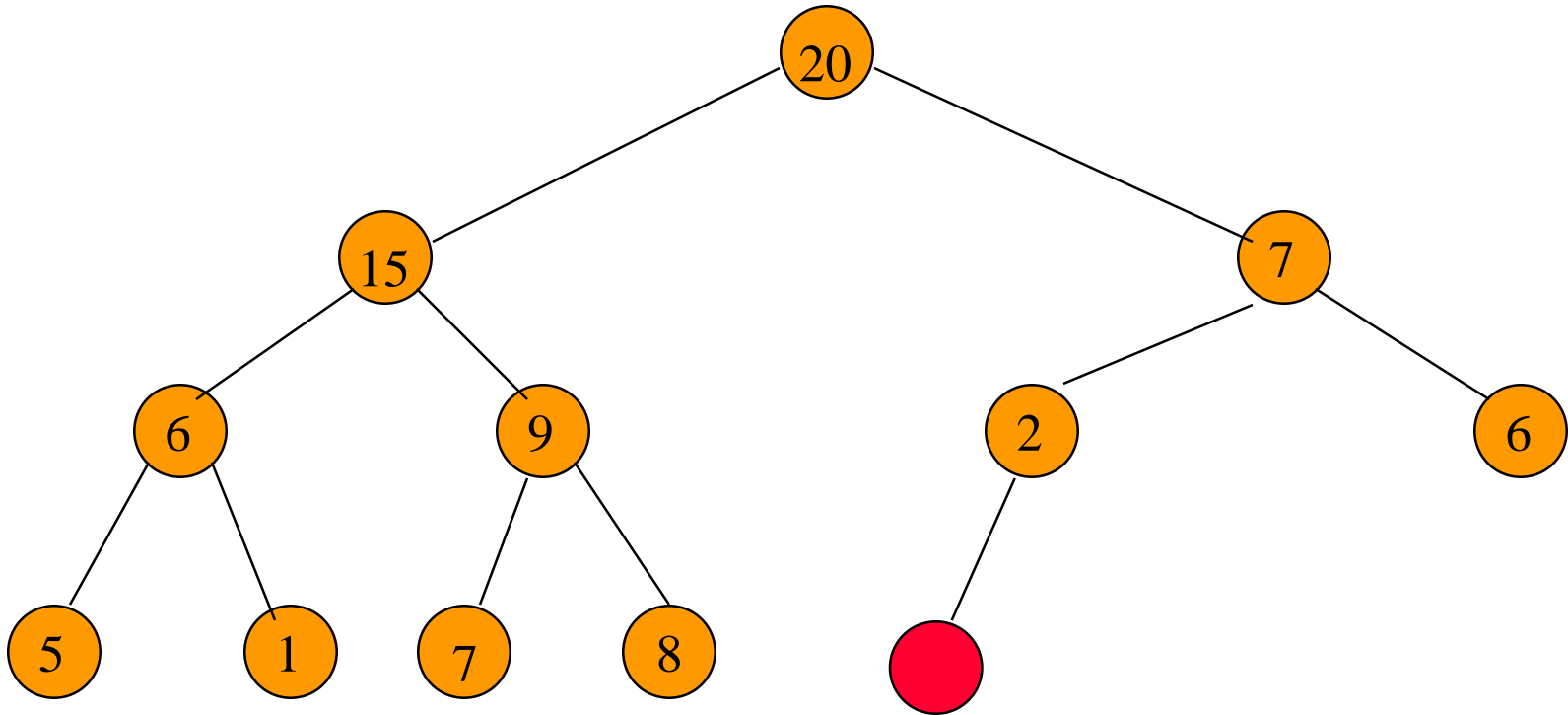
New element is 15.

Putting An Element Into A Max Heap



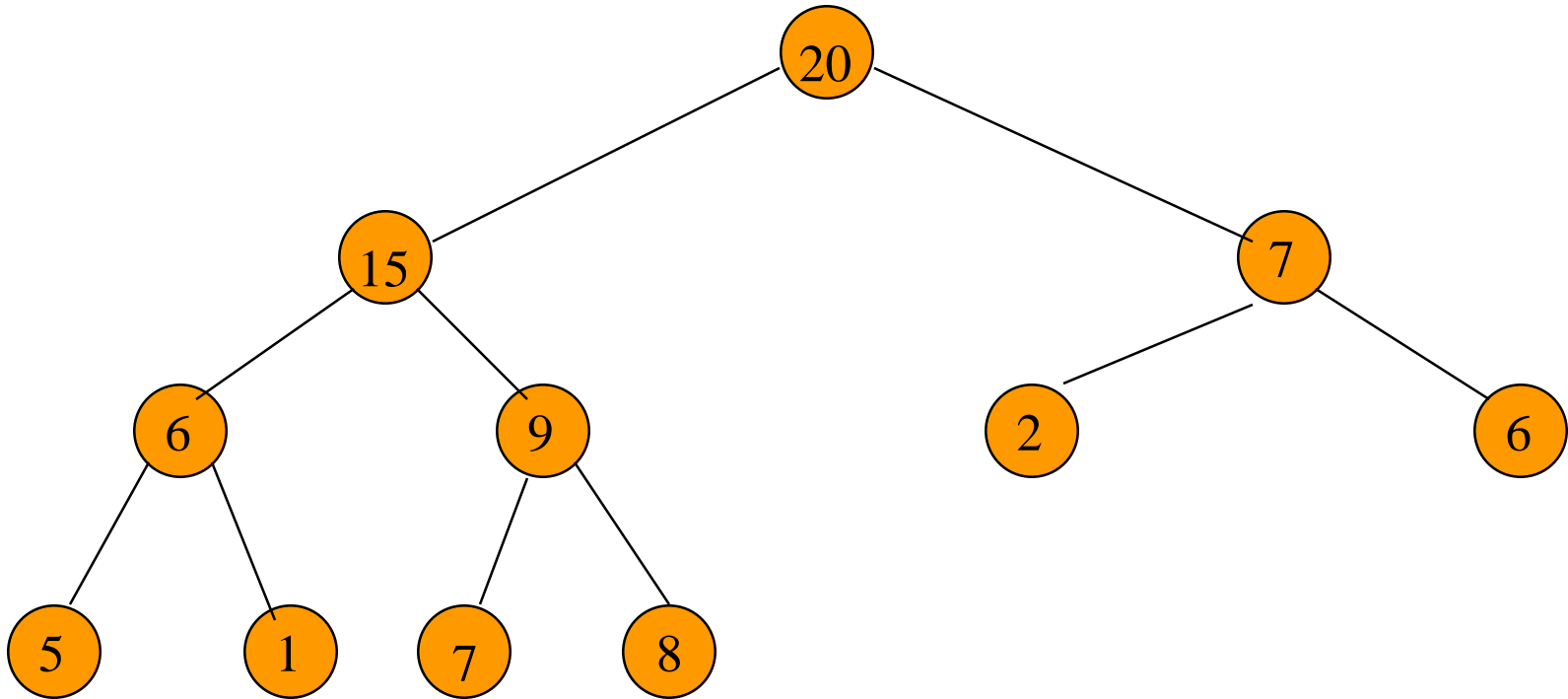
New element is 15.

Complexity Of Put



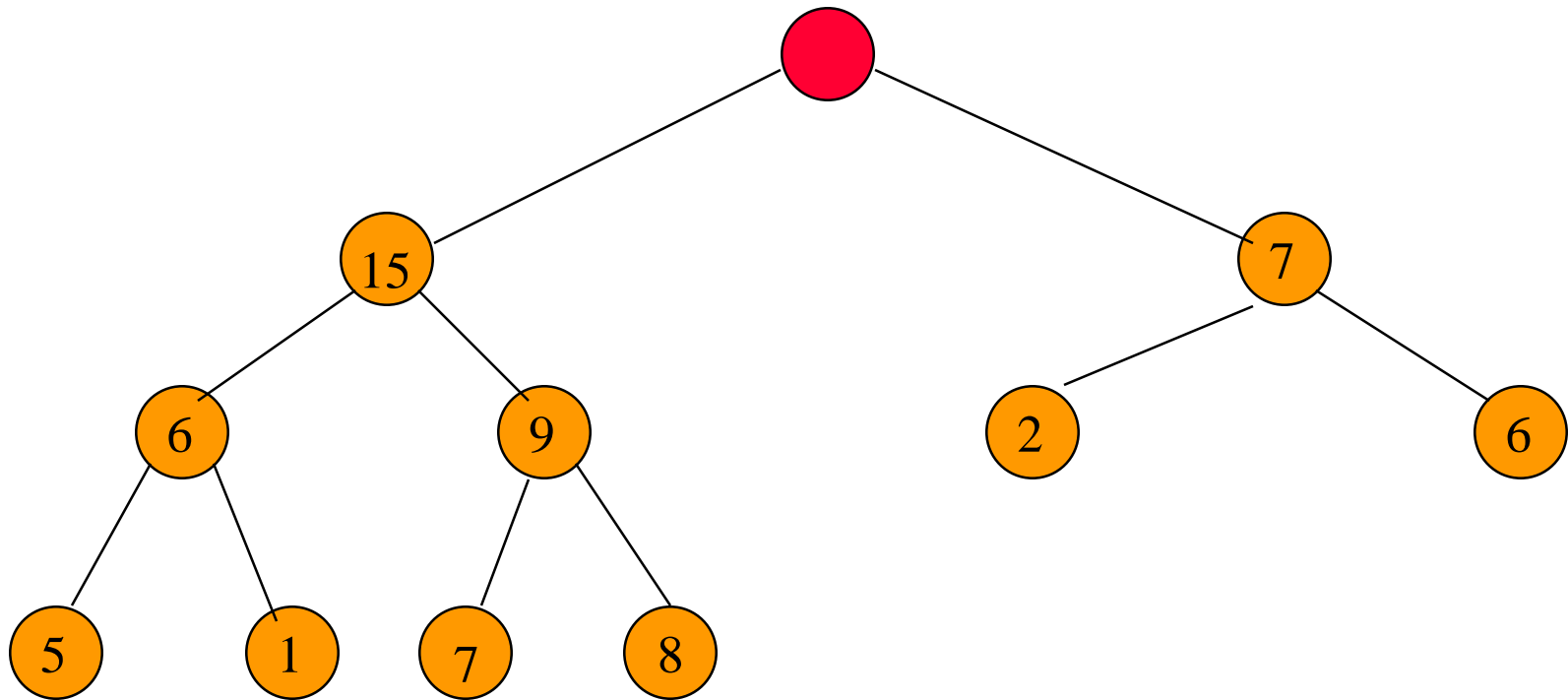
Complexity is $O(\log n)$, where n is heap size.

Removing The Max Element



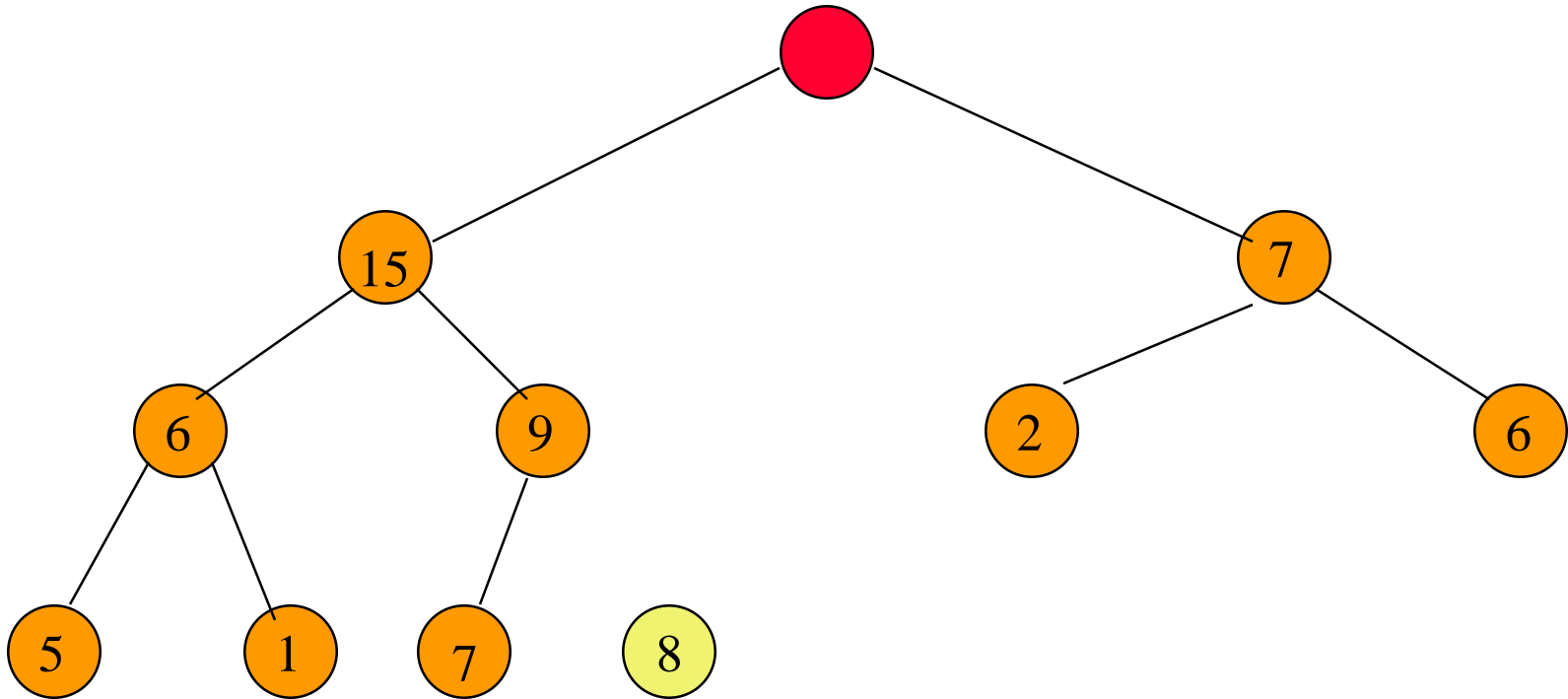
Max element is in the root.

Removing The Max Element



After max element is removed.

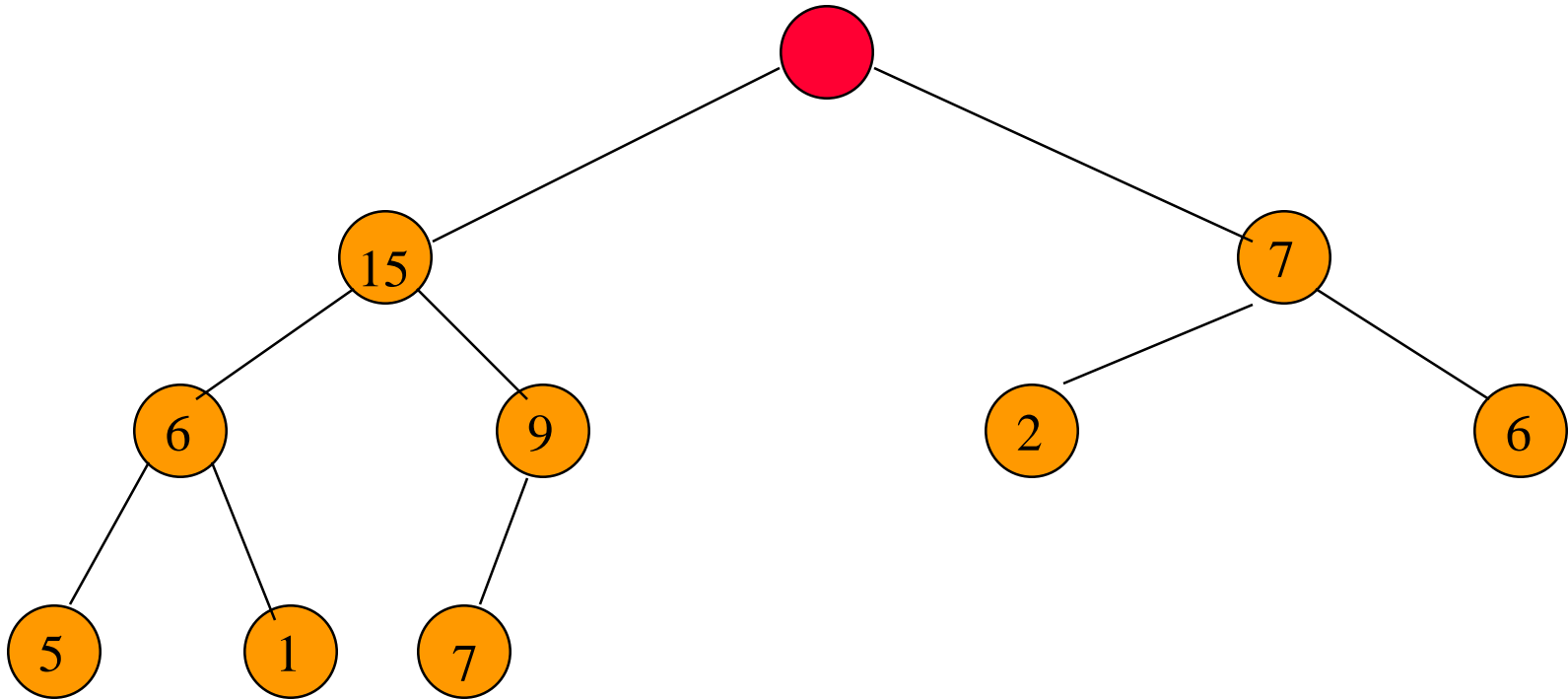
Removing The Max Element



Heap with 10 nodes.

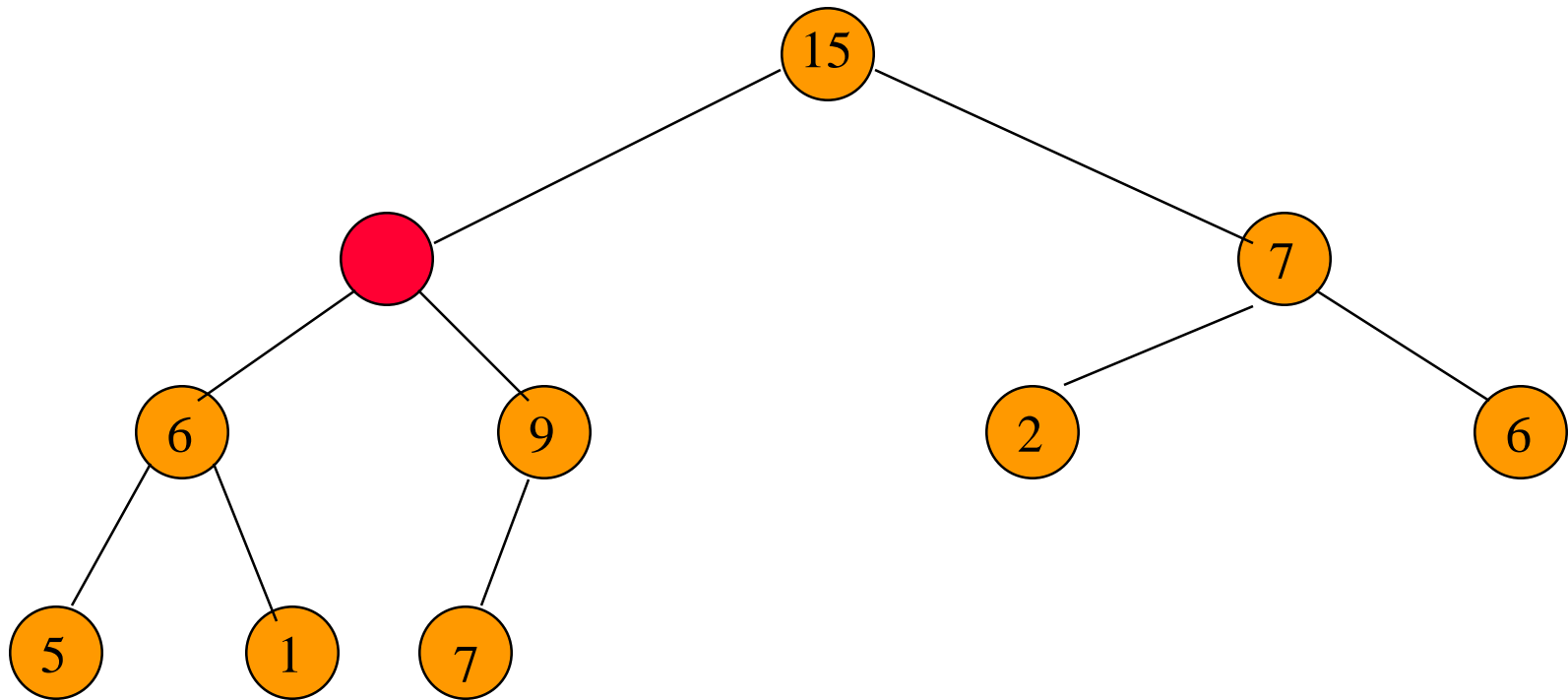
Reinsert 8 into the heap.

Removing The Max Element



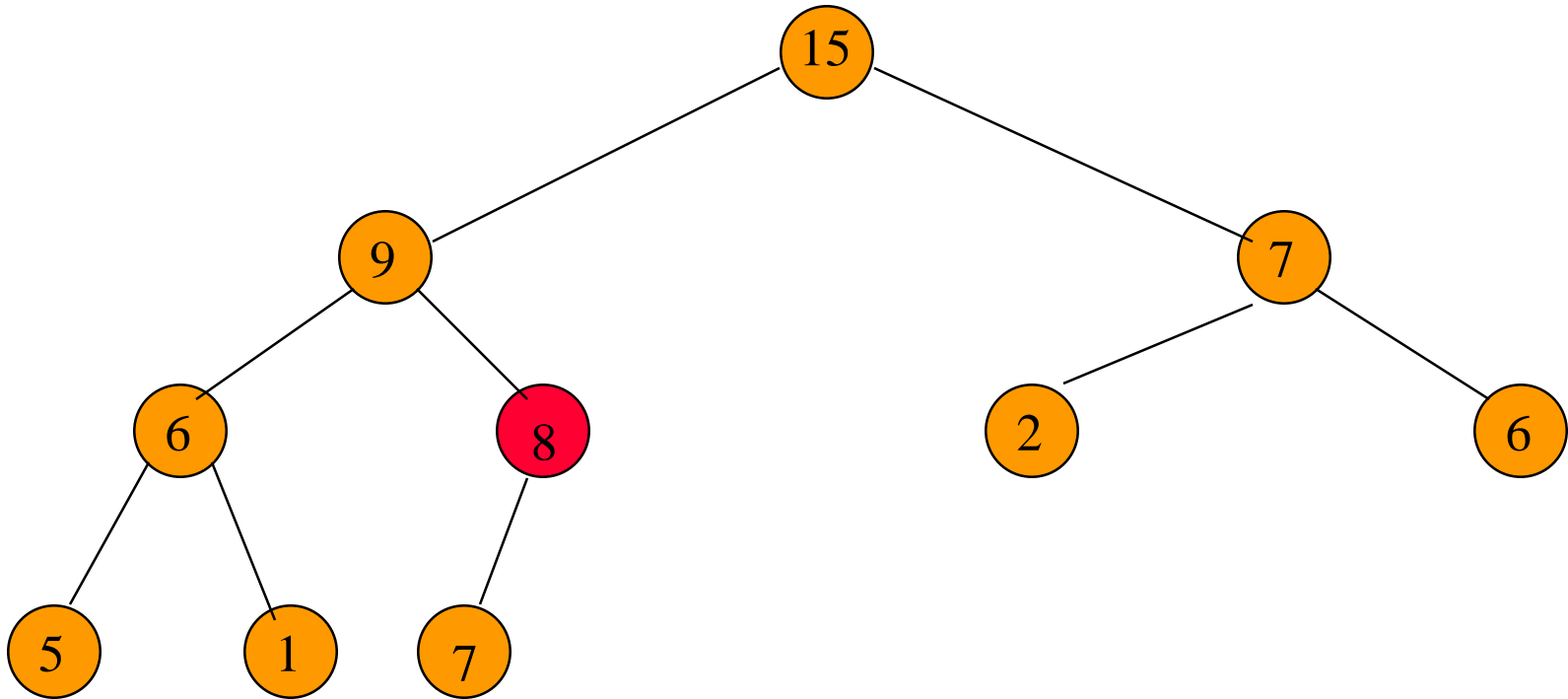
Reinsert **8** into the heap.

Removing The Max Element



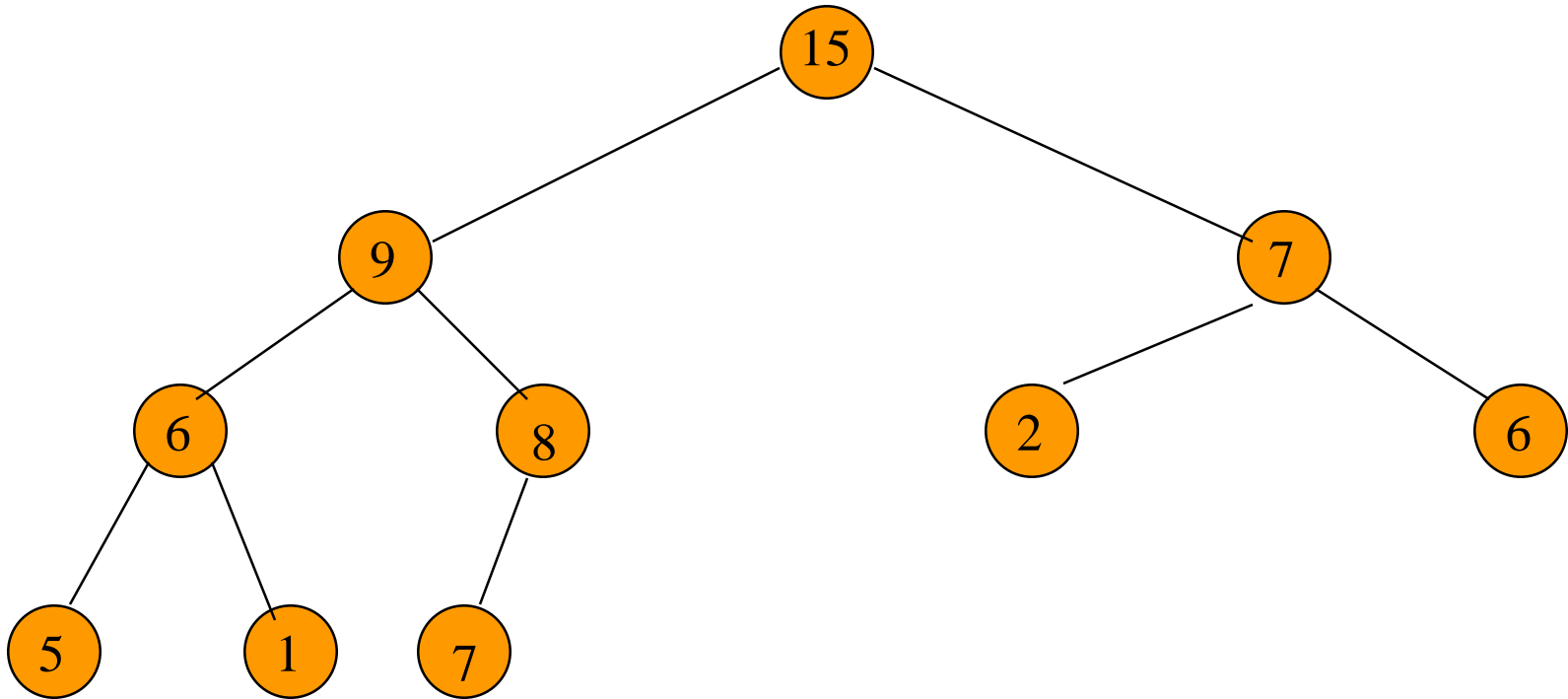
Reinsert **8** into the heap.

Removing The Max Element



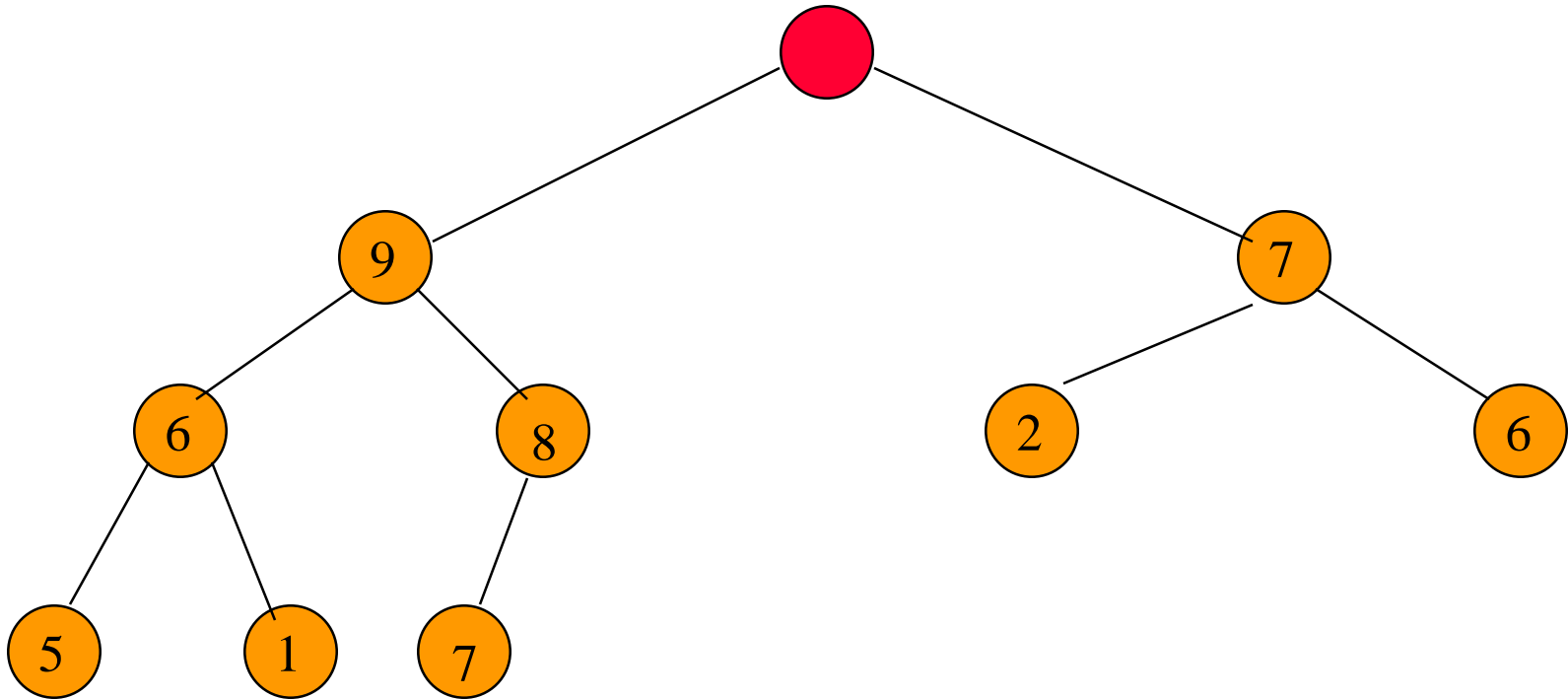
Reinsert **8** into the heap.

Removing The Max Element



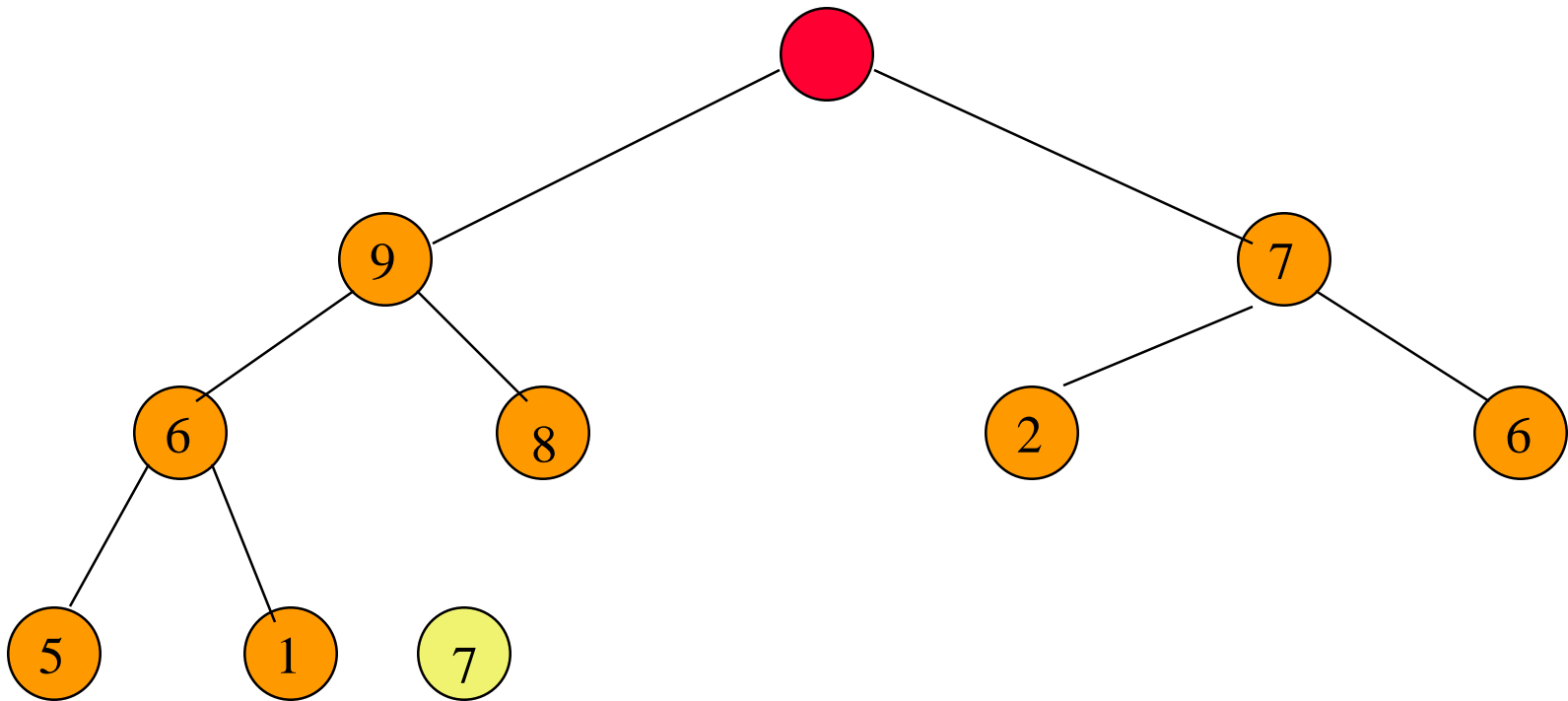
Max element is 15.

Removing The Max Element



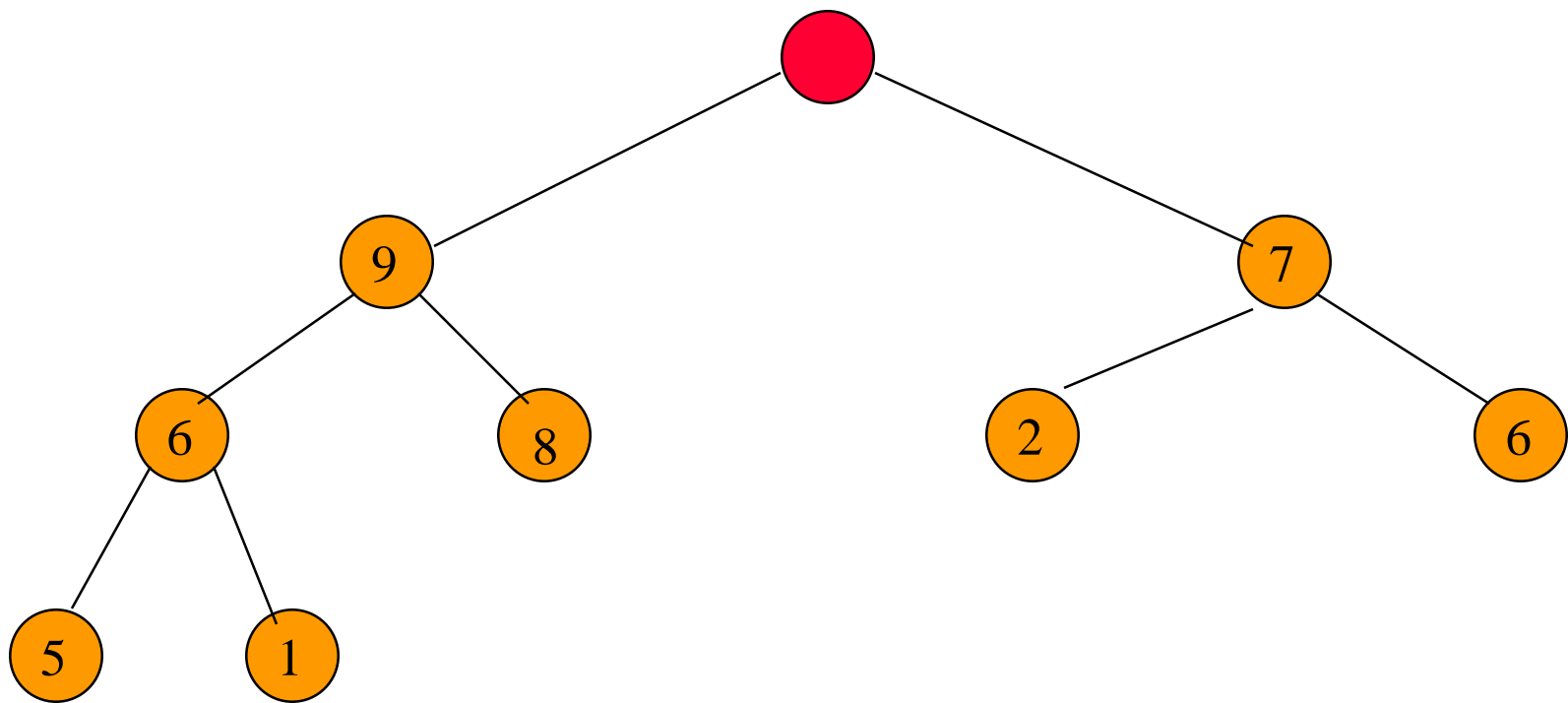
After max element is removed.

Removing The Max Element



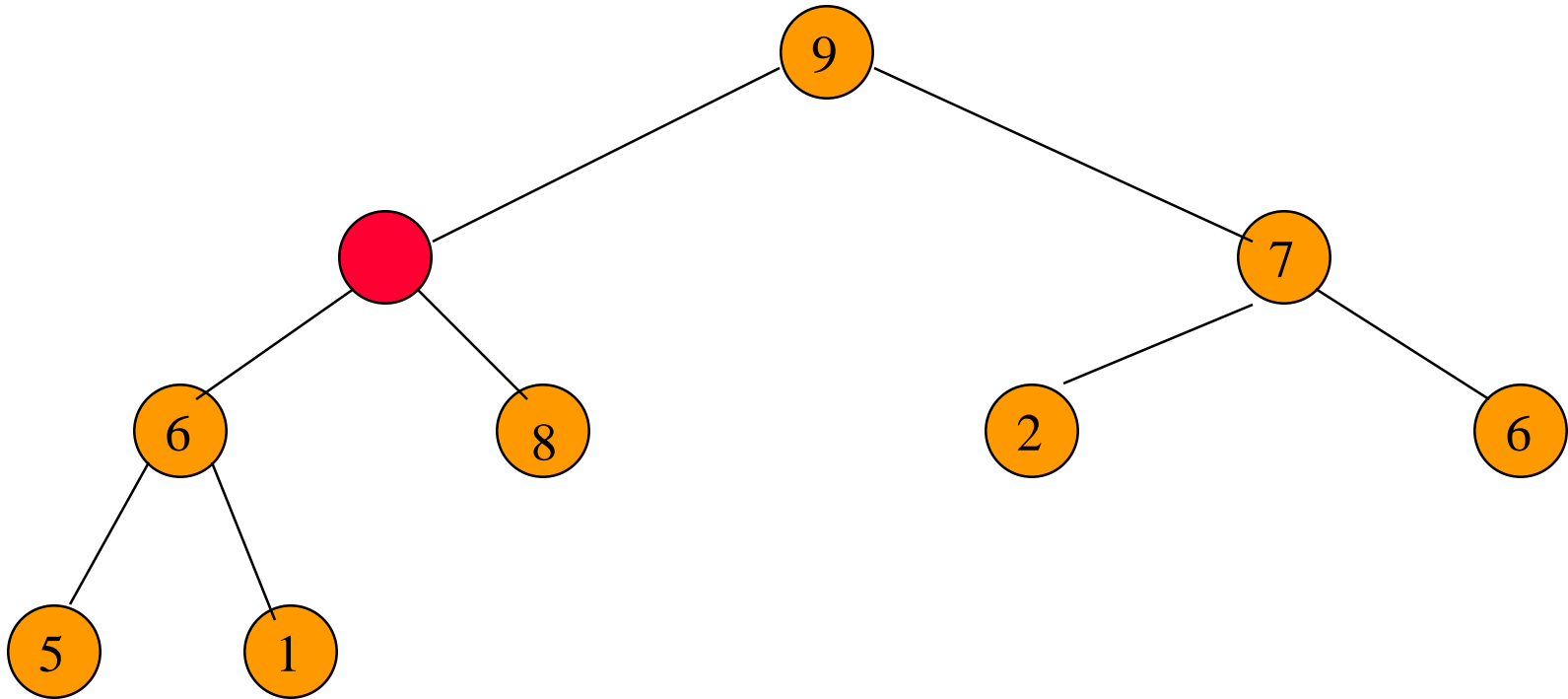
Heap with 9 nodes.

Removing The Max Element



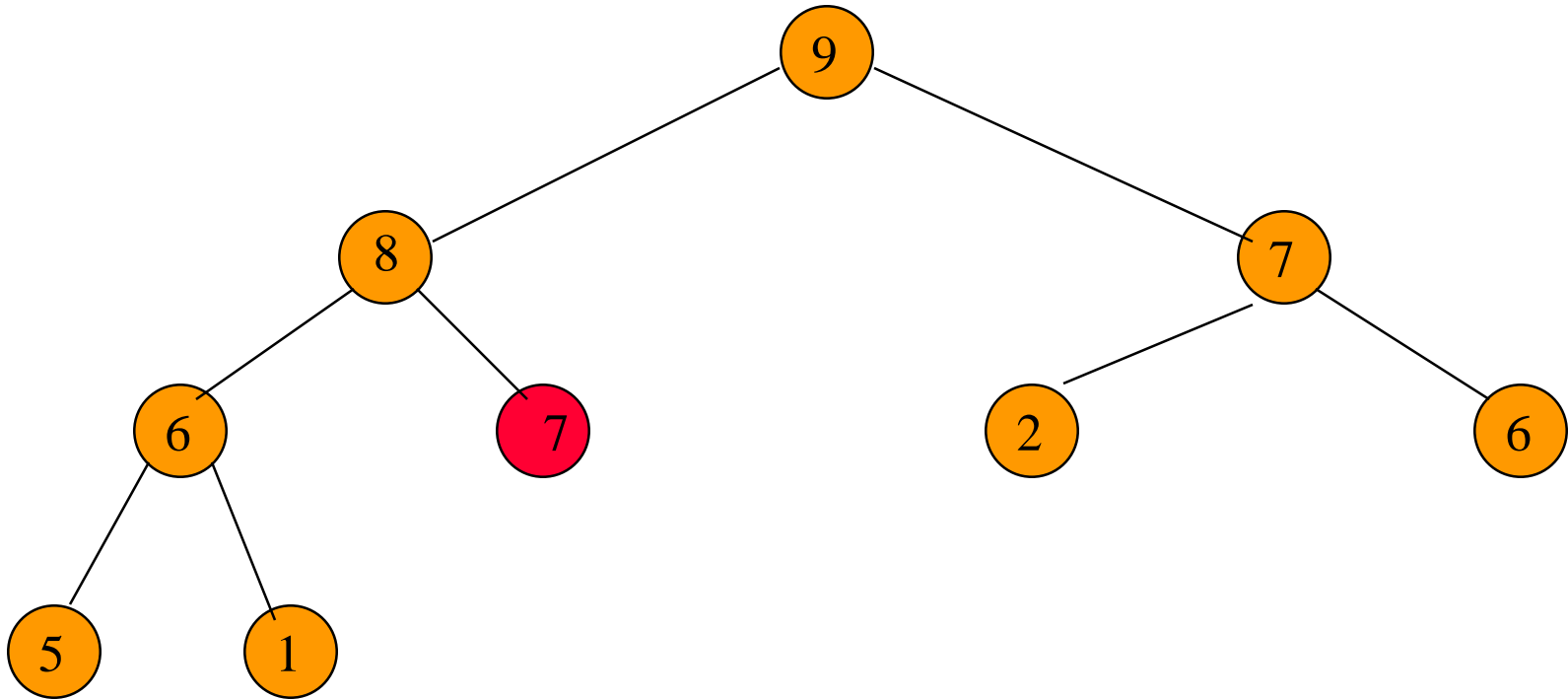
Reinsert **7**.

Removing The Max Element



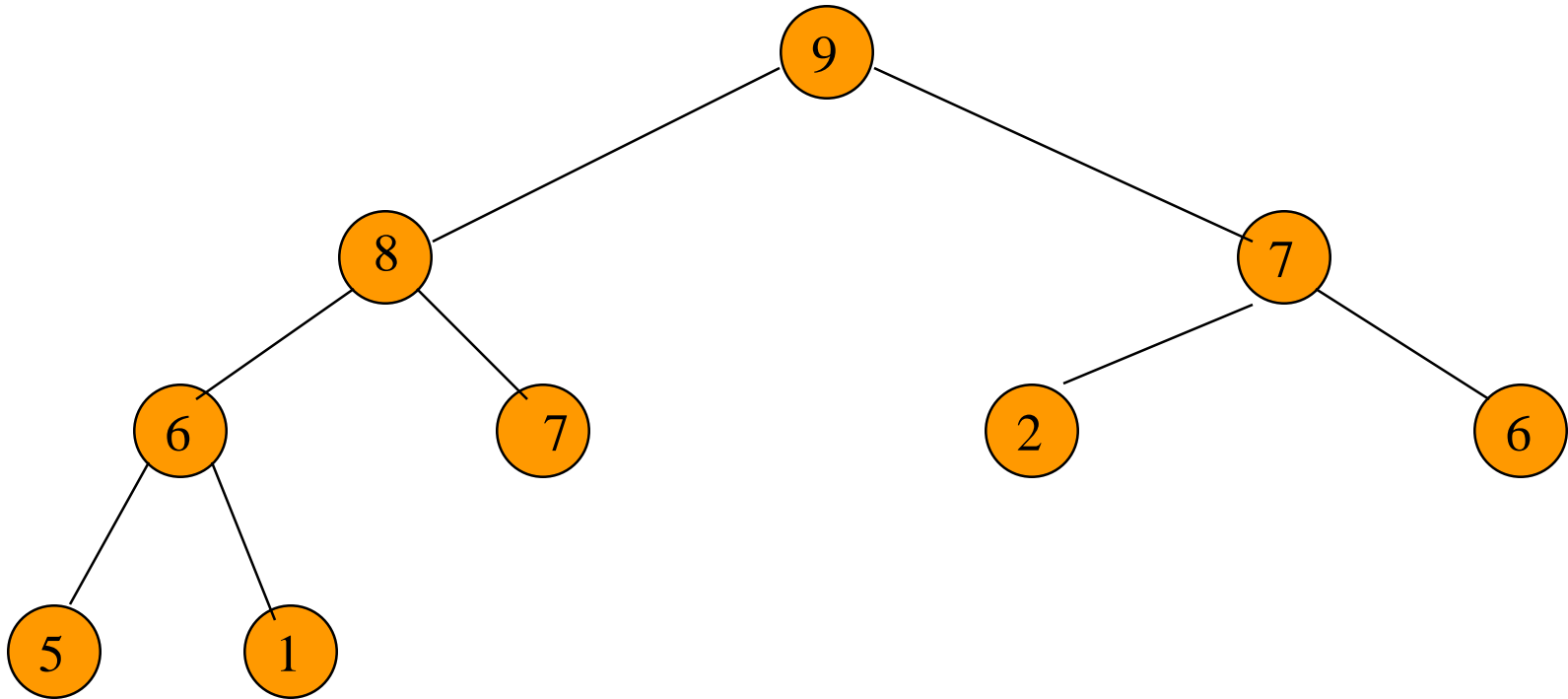
Reinsert **7**.

Removing The Max Element



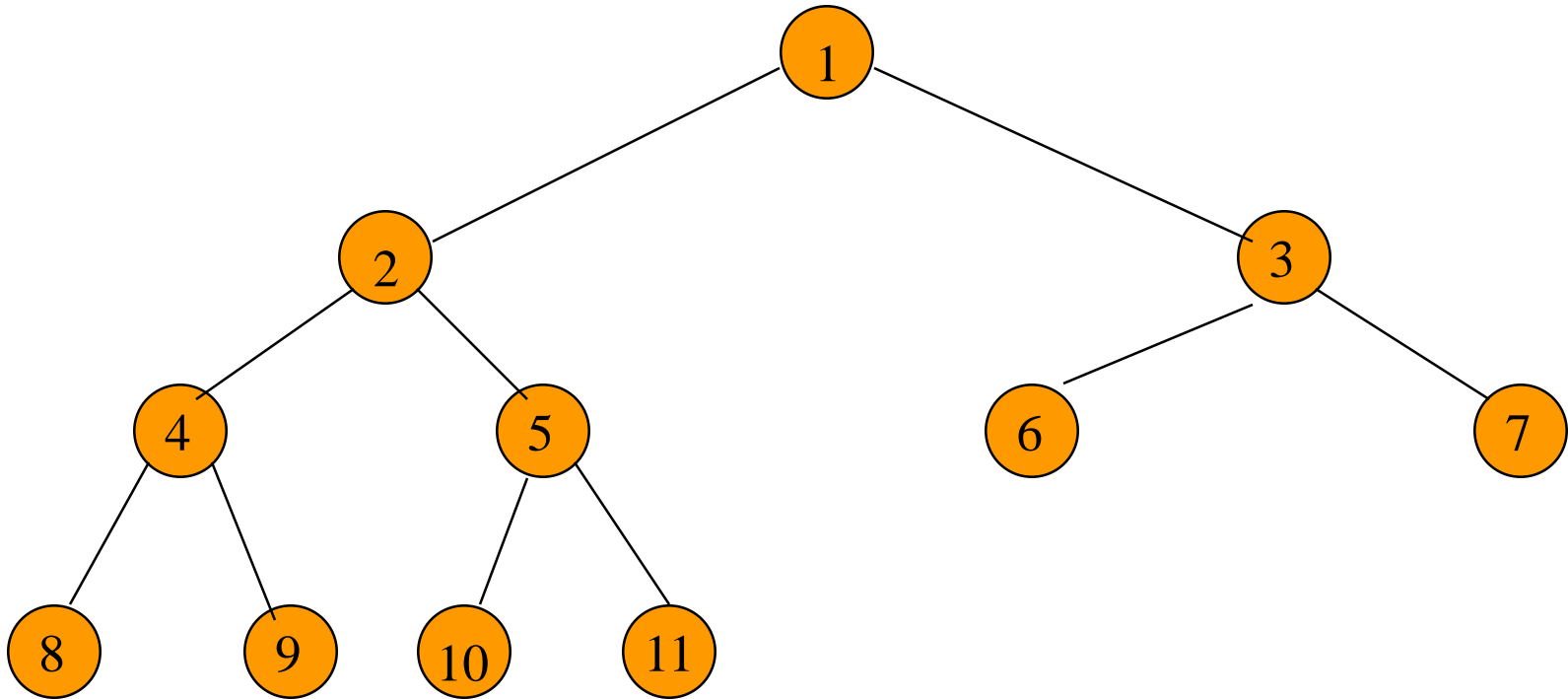
Reinsert **7**.

Complexity Of Remove Max Element



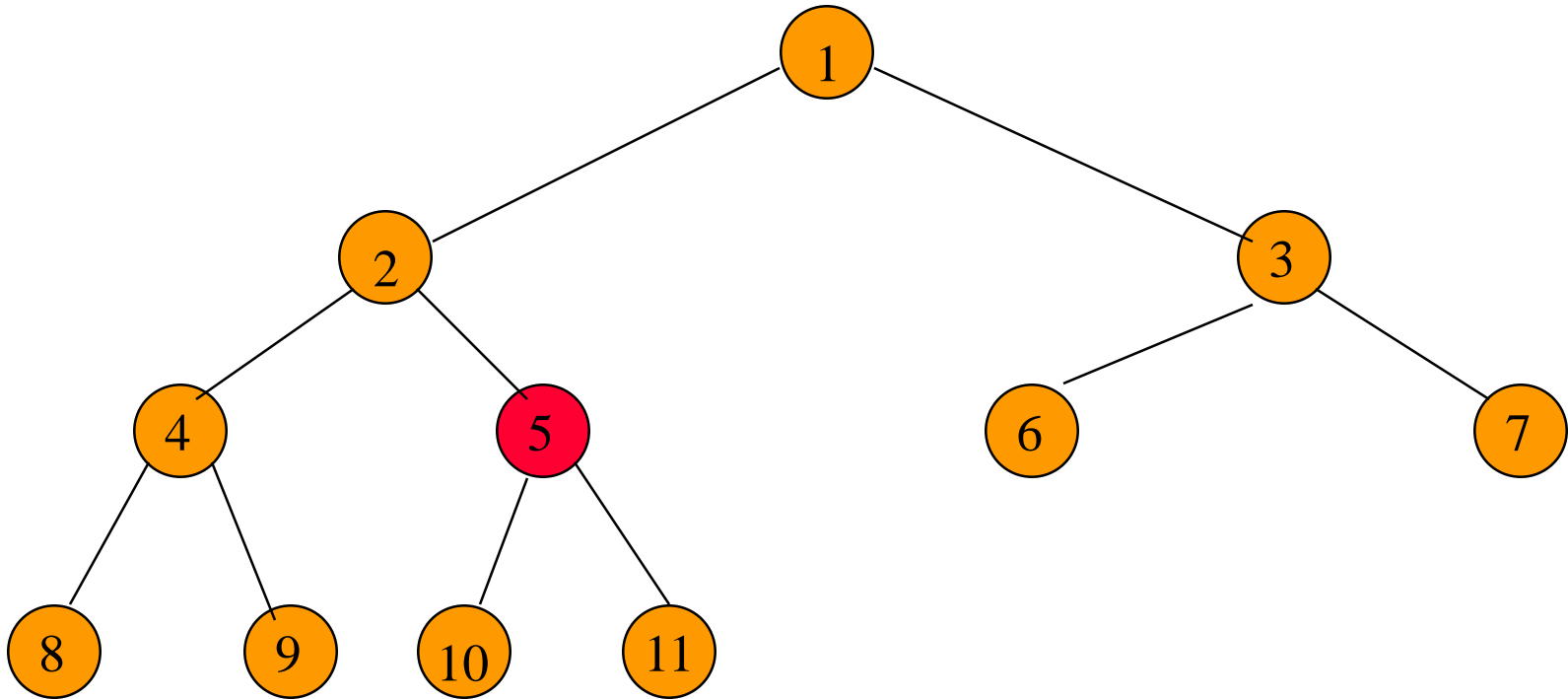
Complexity is $O(\log n)$.

Initializing A Max Heap



input array = $[-, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11]$

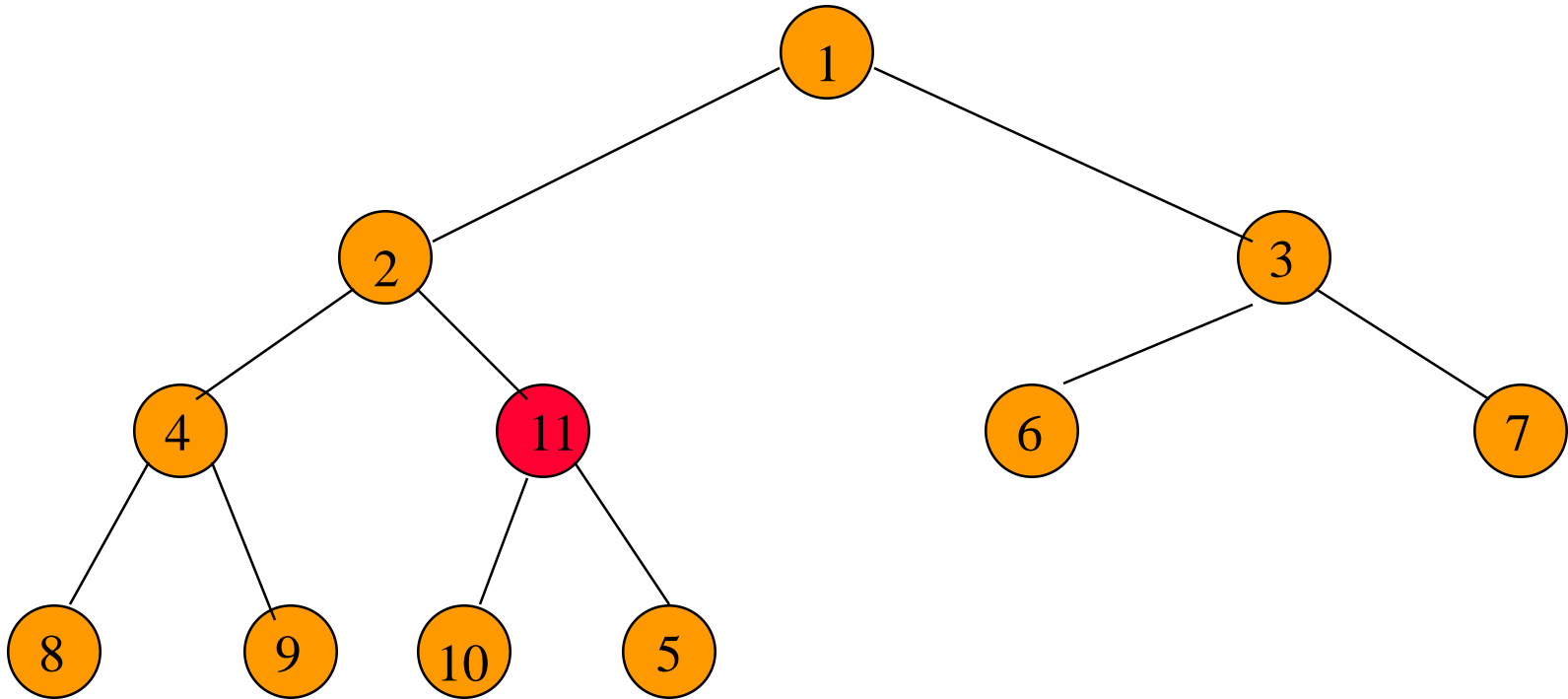
Initializing A Max Heap



Start at rightmost array position that has a child.

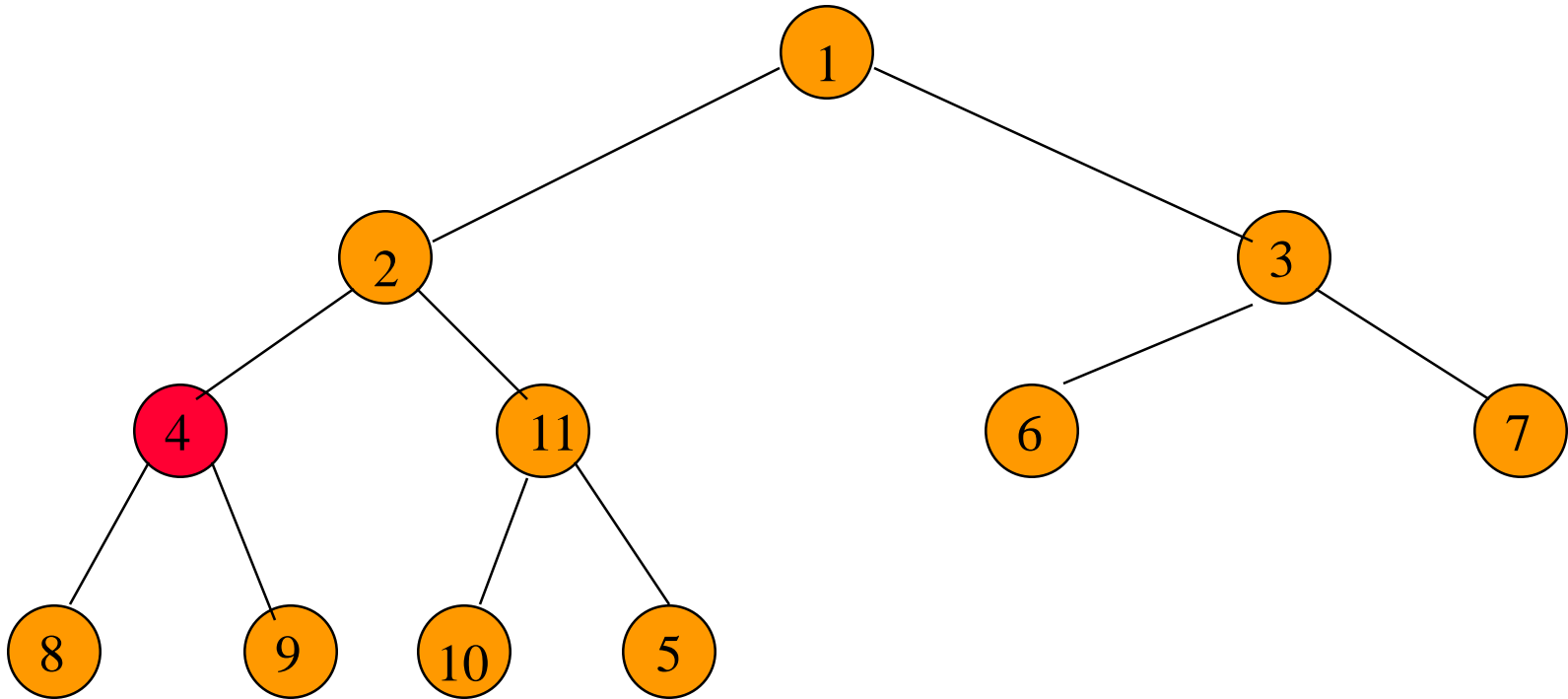
Index is $n/2$.

Initializing A Max Heap

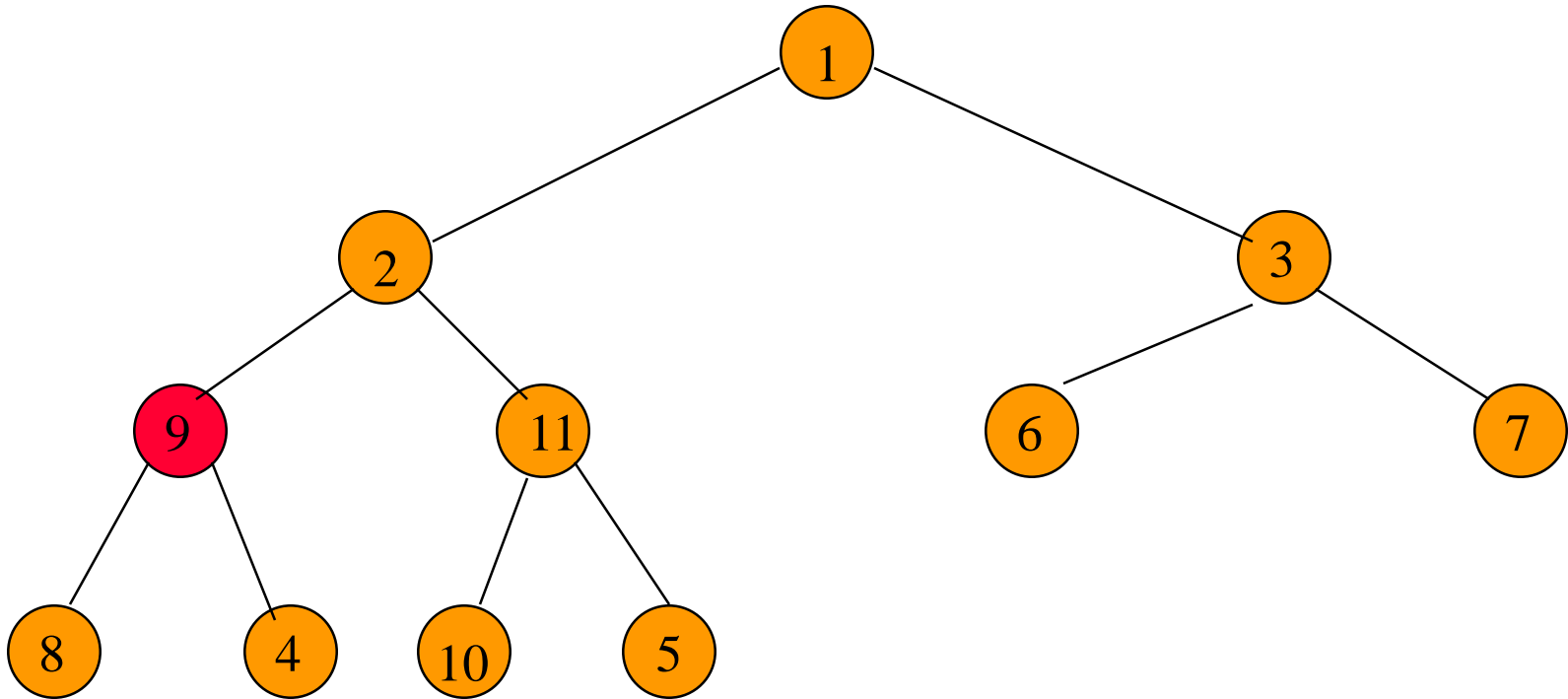


Move to next lower array position.

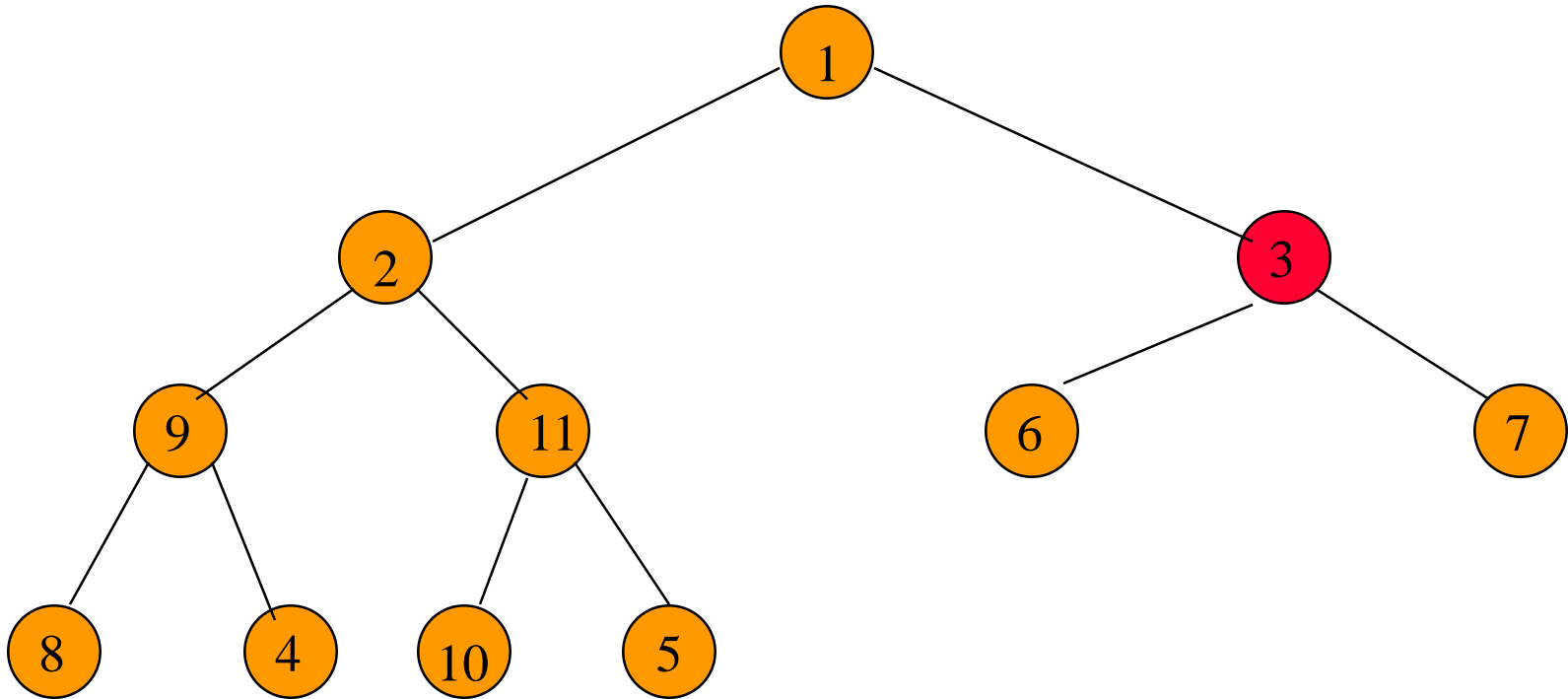
Initializing A Max Heap



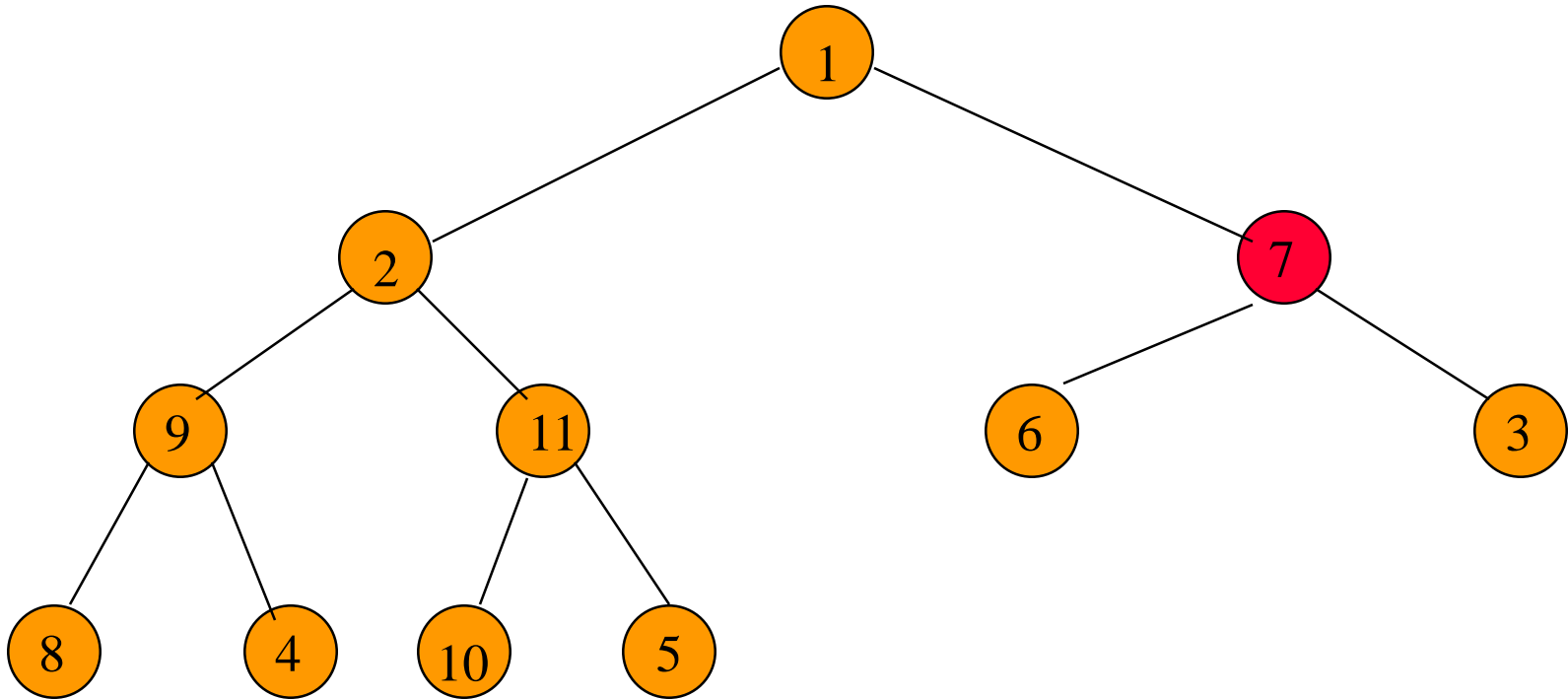
Initializing A Max Heap



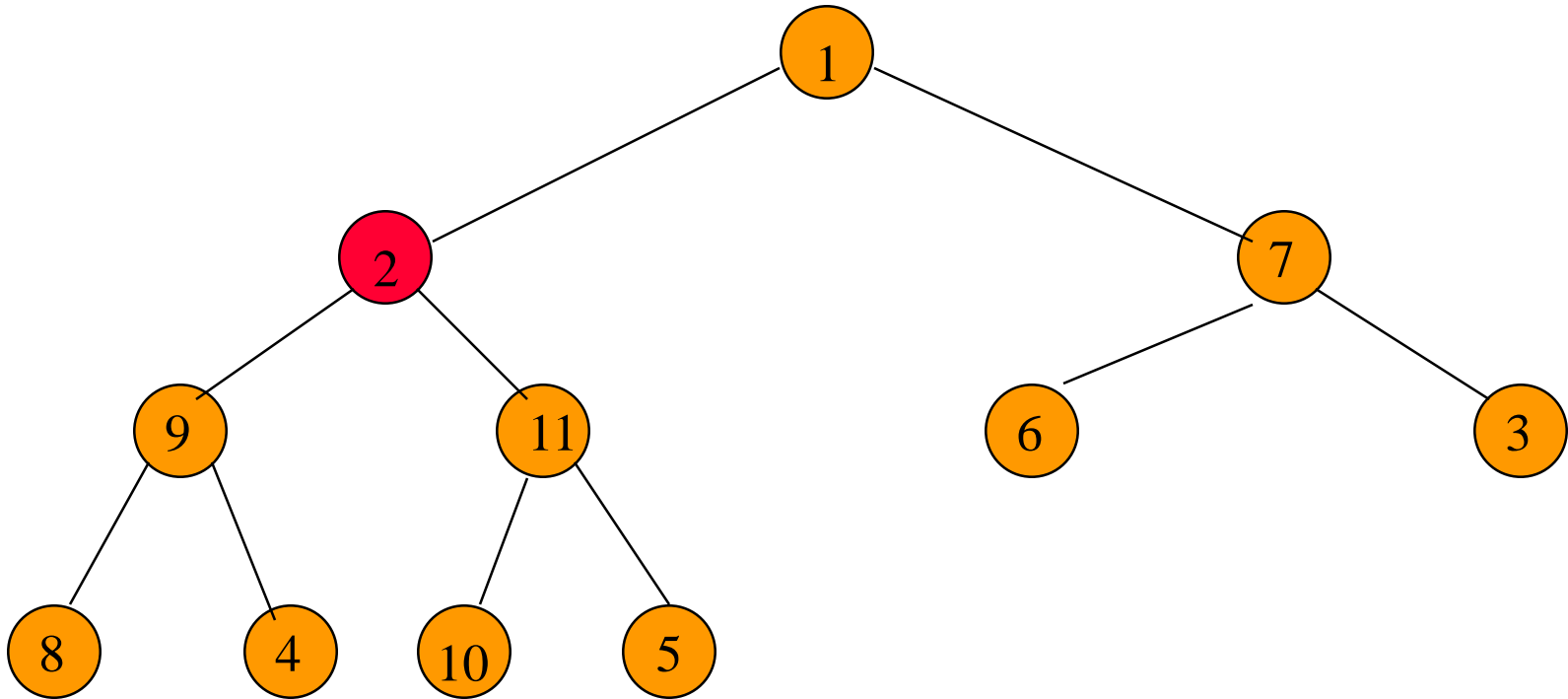
Initializing A Max Heap



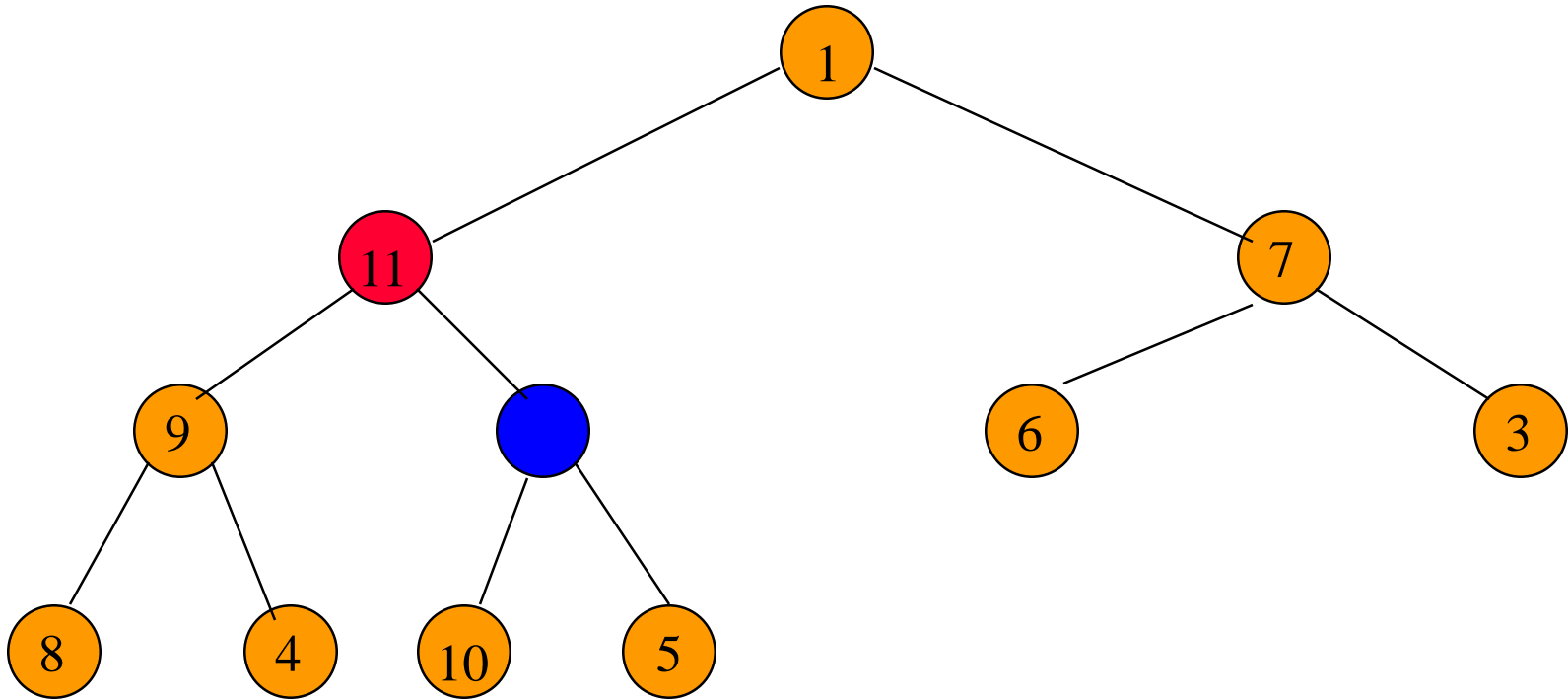
Initializing A Max Heap



Initializing A Max Heap

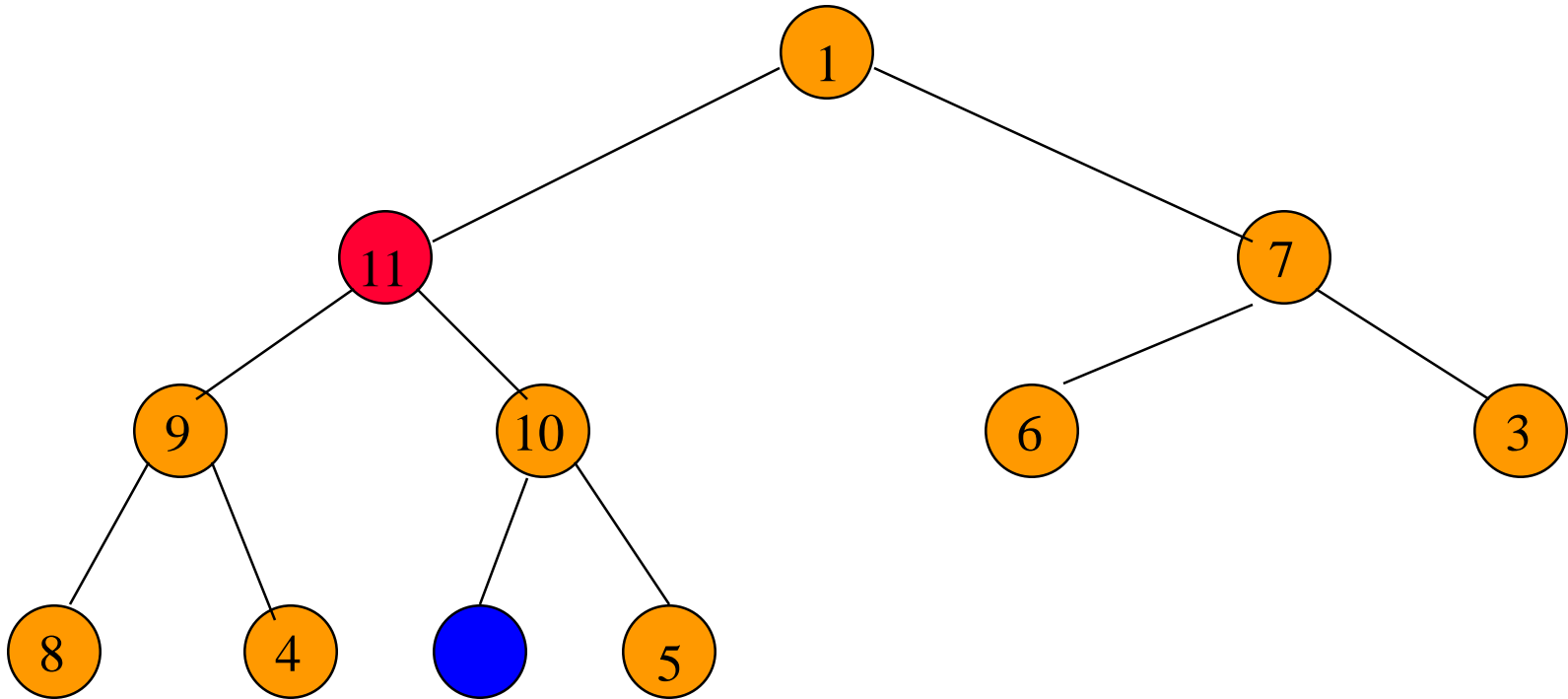


Initializing A Max Heap



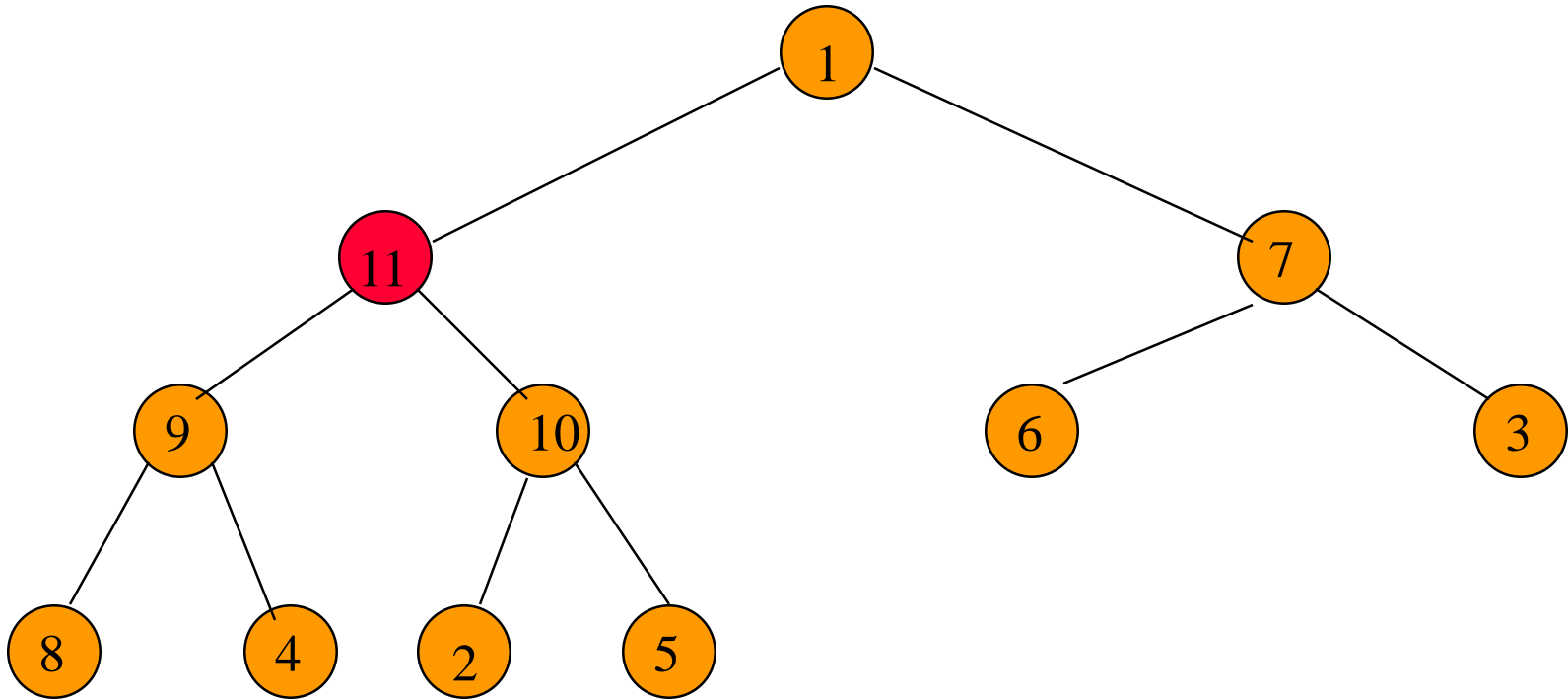
Find a home for 2.

Initializing A Max Heap



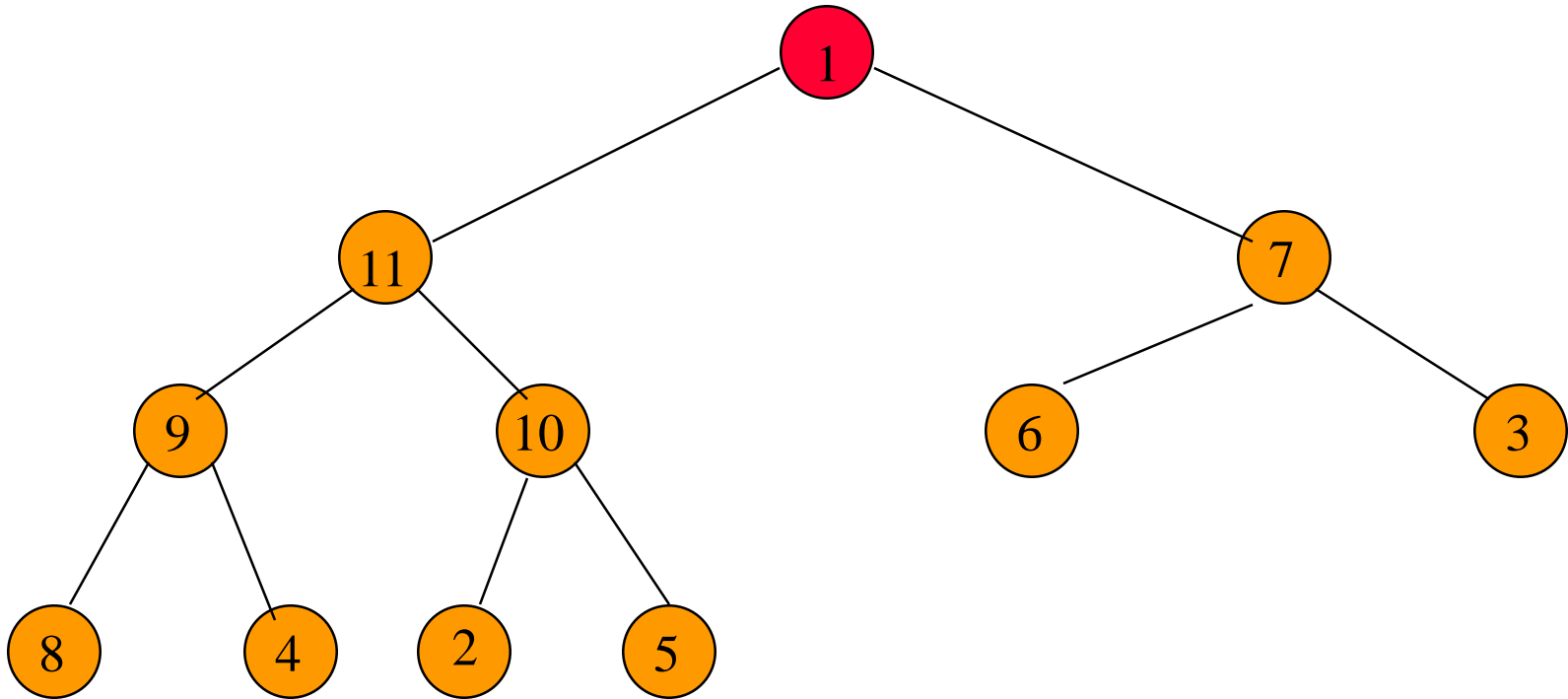
Find a home for 2.

Initializing A Max Heap



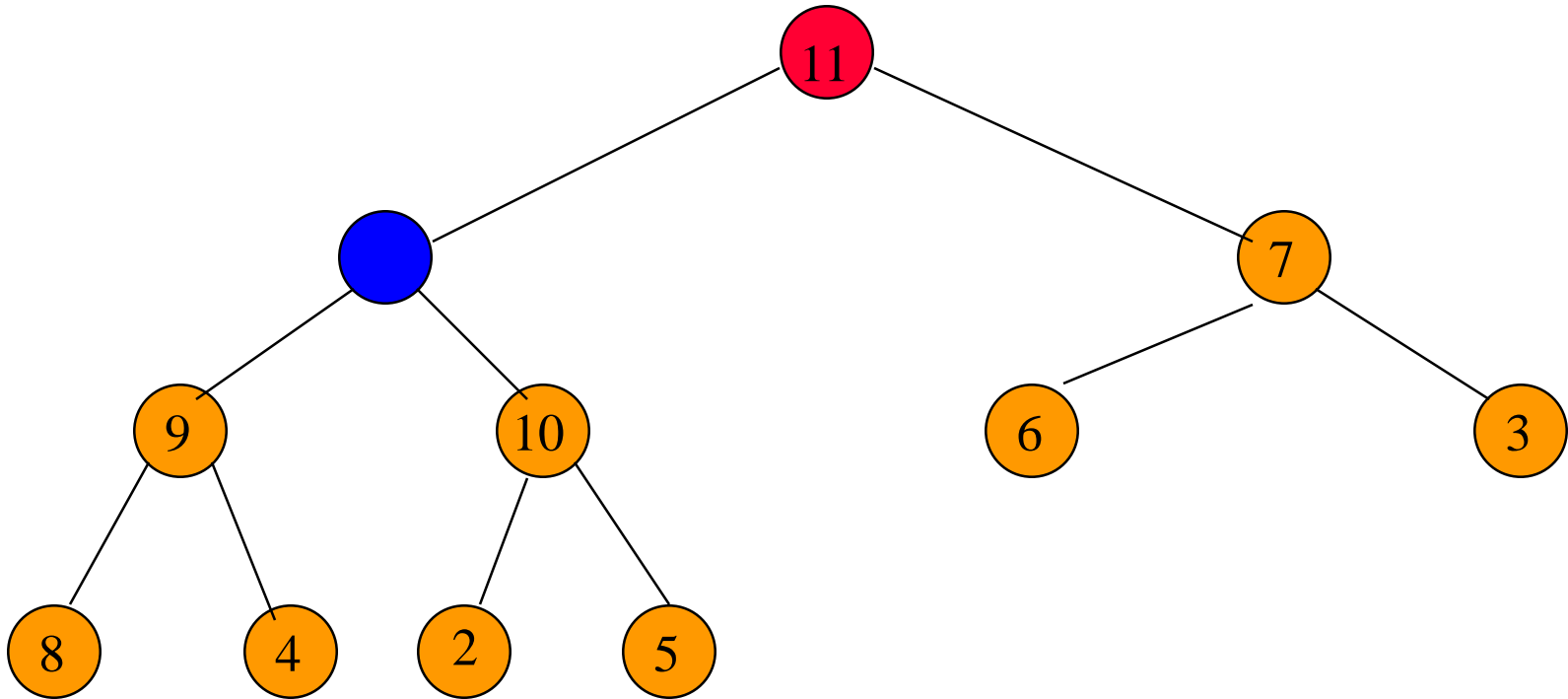
Done, move to next lower array position.

Initializing A Max Heap



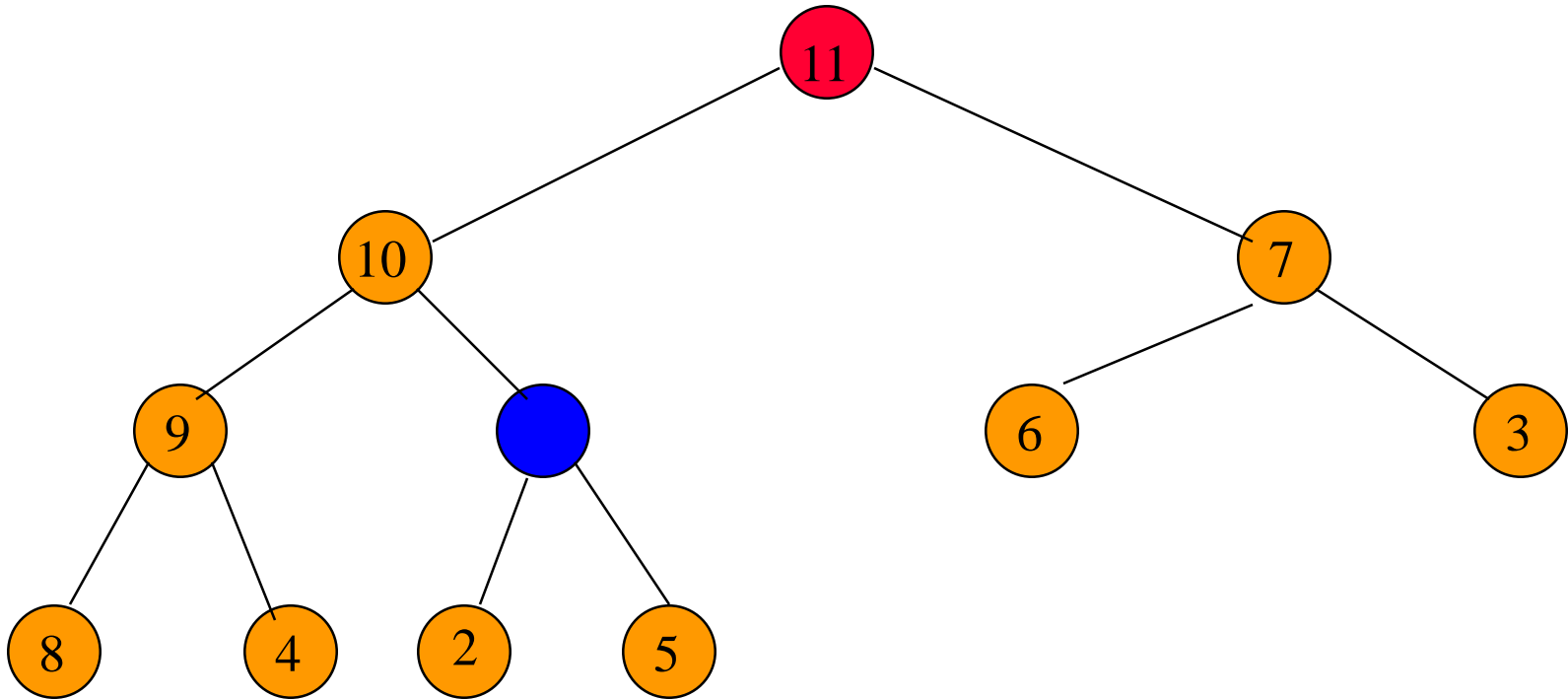
Find home for 1.

Initializing A Max Heap



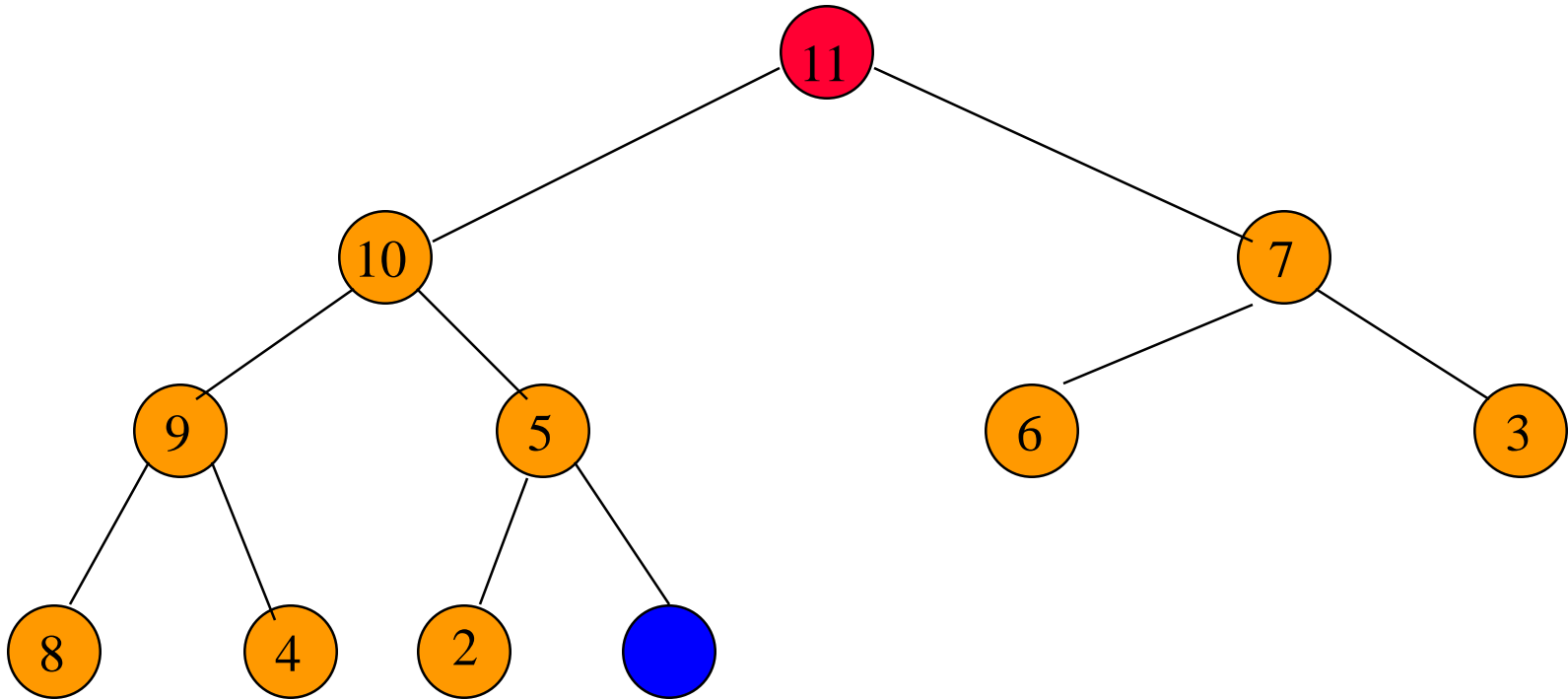
Find home for 1.

Initializing A Max Heap



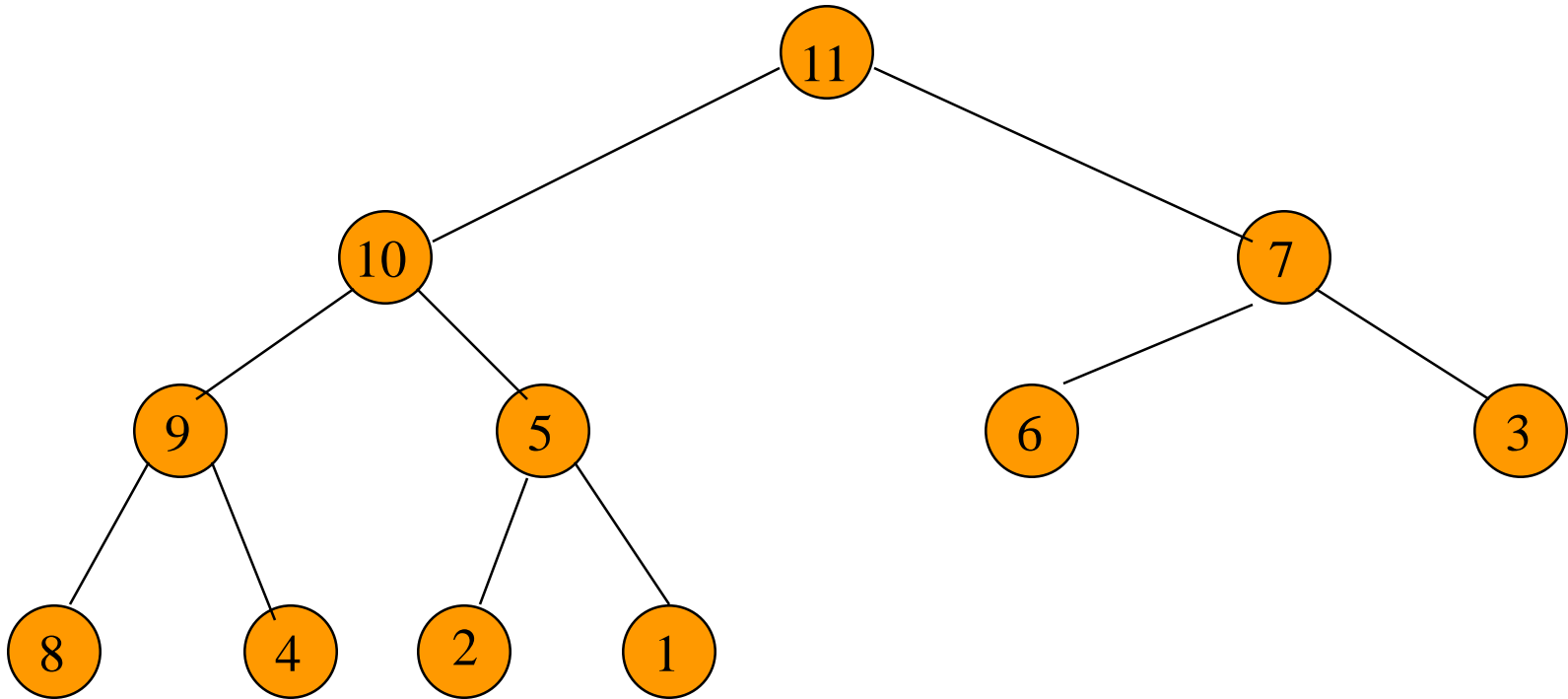
Find home for 1.

Initializing A Max Heap



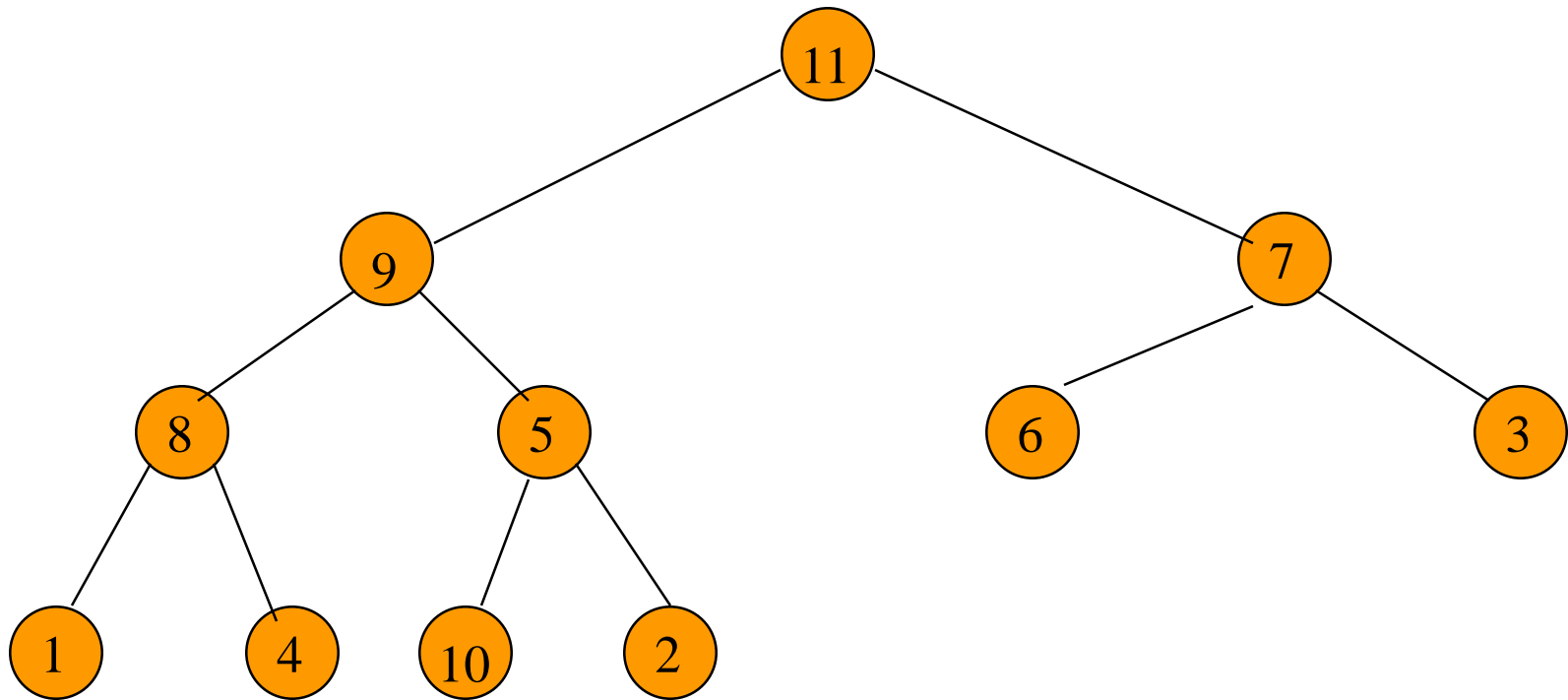
Find home for 1.

Initializing A Max Heap



Done.

Time Complexity



Height of heap = h .

Number of subtrees with root at level j is $\leq 2^{j-1}$.

Time for each subtree is $O(h-j+1)$.

Complexity



Time for level j subtrees is $\leq 2^{j-1}(h-j+1) = t(j)$.

Total time is $t(1) + t(2) + \dots + t(h-1) = O(n)$.

Huffman codes

Fixed length codes

- Want to represent data as a sequence of 0's and 1's

For example: BACADAEAFABBAAGA

- A-000 B-001 C-010 D-011 E-100 F-101 G-110 H-111

001 000 010 000 011 000 100 000 101 000 001 001 000 000
000 110 000 111

- This is a fixed length code.
- Can we make the sequence of 0's and 1's shorter ?

Variable length codes

- A 0 B 100 C 1010 D 1011

E 1100 F 1101 G 1110 H 1111

100010100101101100011010100100000111001111

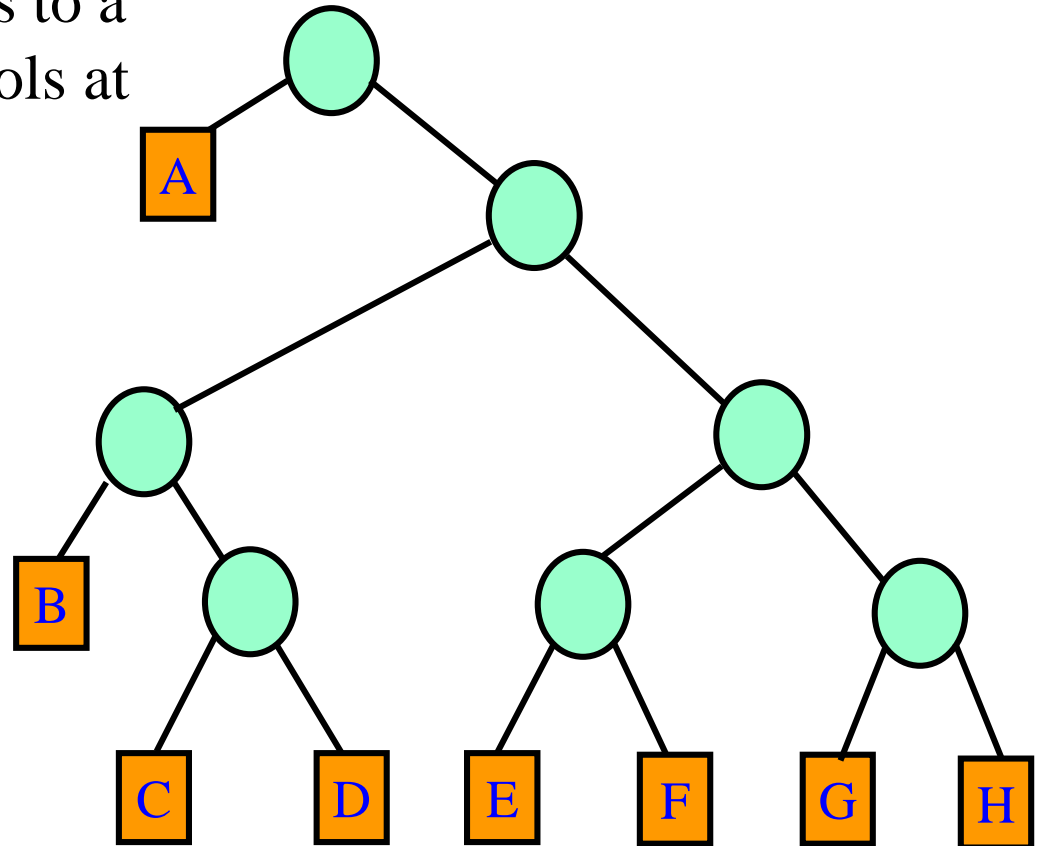
- This is a variable length code.
- How do we decode ?
- Use prefix codes: No codeword is a prefix of the other

Representing prefix codes

A 0 B 100 C 1010 D 1011

E 1100 F 1101 G 1110 H 1111

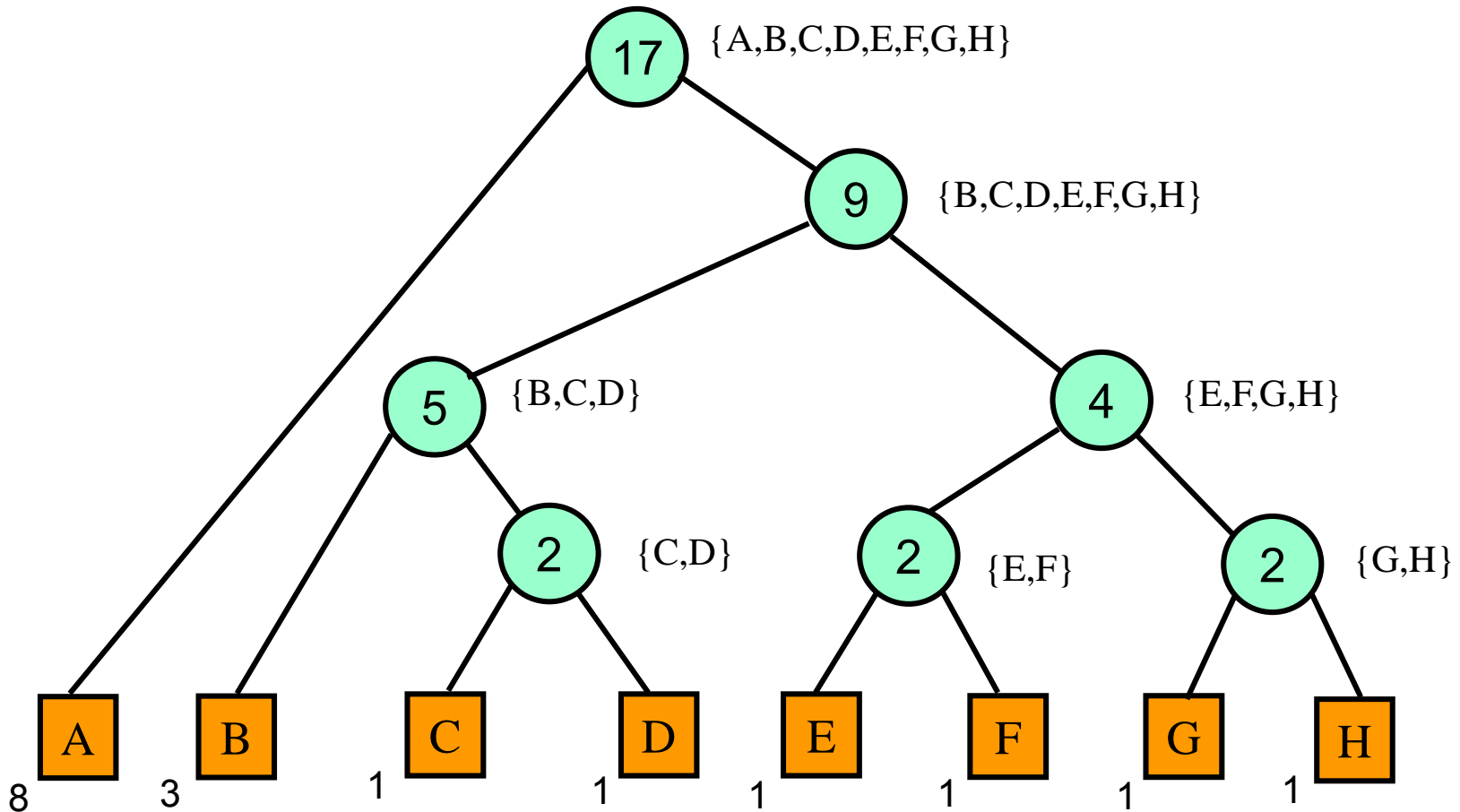
- A prefix code corresponds to a binary tree with the symbols at the **leaves** and vice versa.



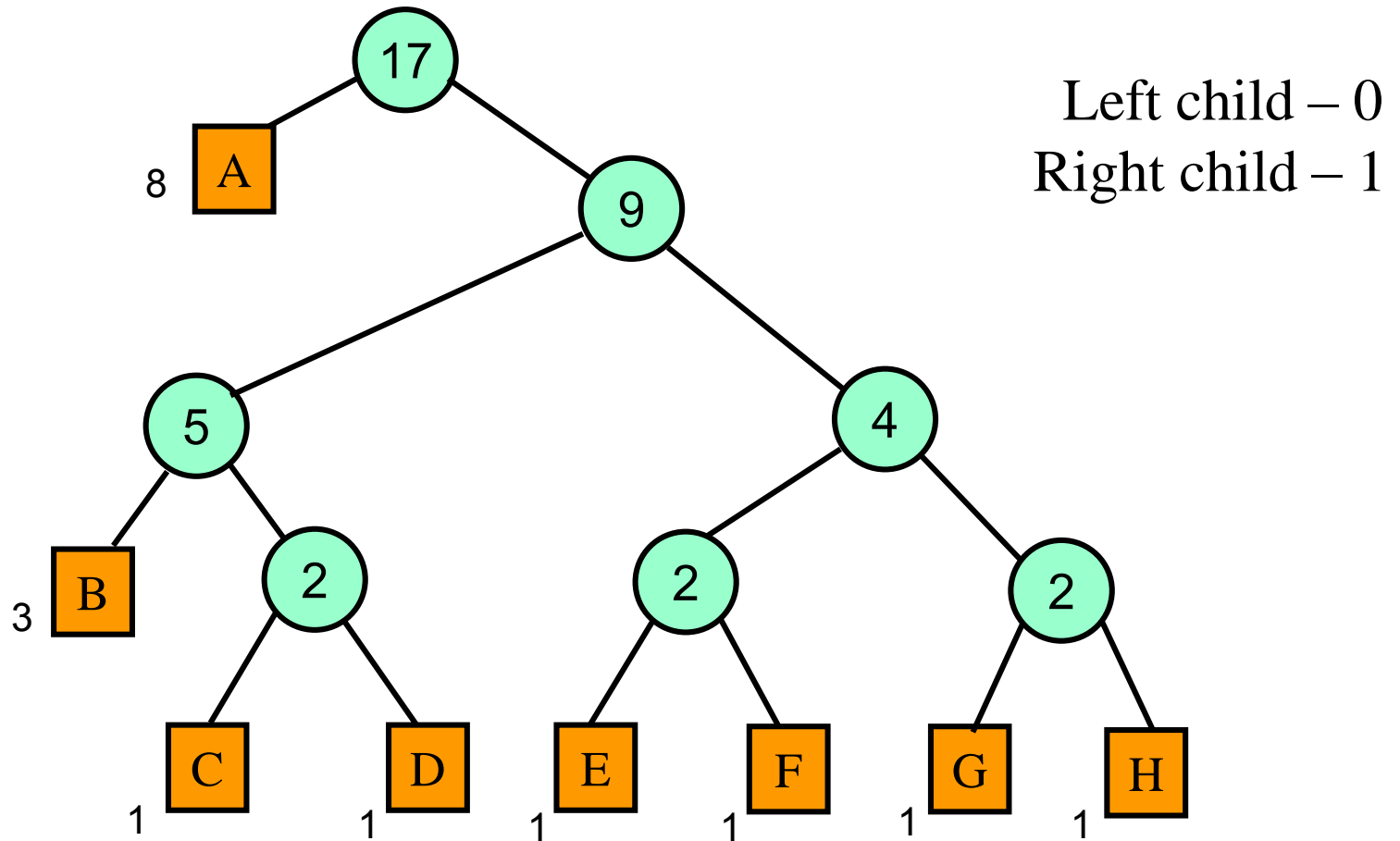
Construction Algorithm

- Compute the character frequencies
- Insert the frequencies into a min heap
- Repeat
 - Delete the two minimum elements, f_1 and f_2 from the heap
 - insert f_1+f_2 into the heap
- Complexity: $O(n \log n)$.

Construction of Huffman tree



Representation



$\text{code}(D) = 1011$; $\text{code}(F) = 1101$