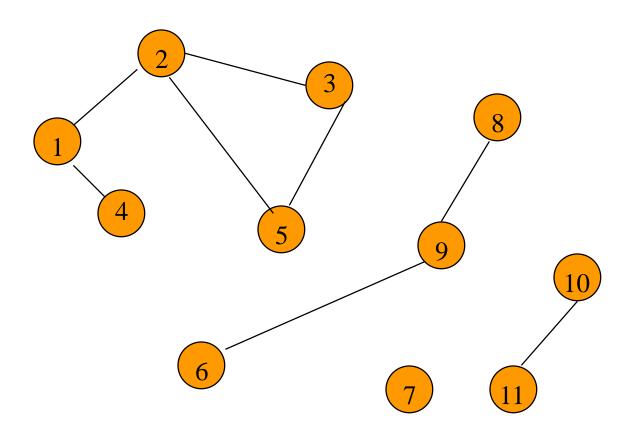
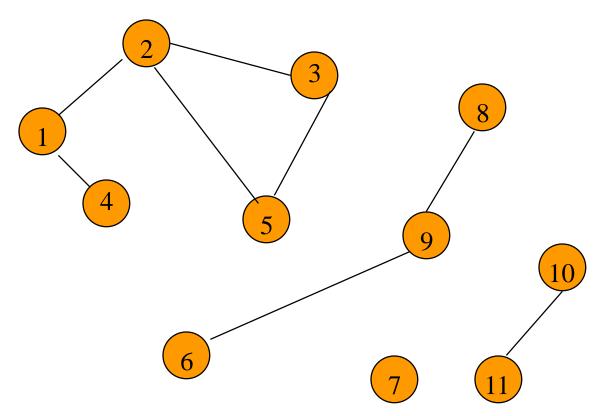
Data structures Fall 2018

• A vertex u is reachable from vertex v iff there is a path from v to u.



 A search method starts at a given vertex v and visits/labels/marks every vertex that is reachable from v.

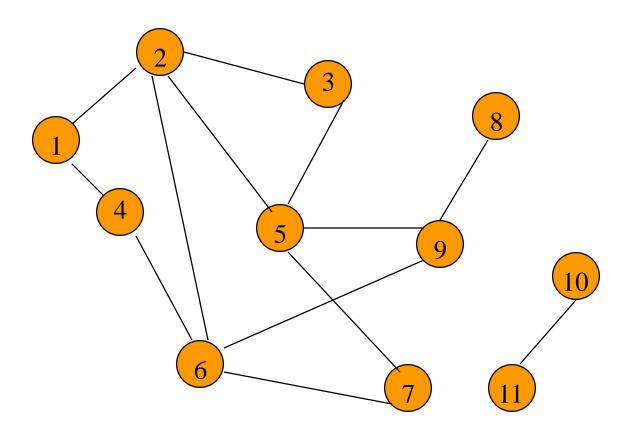


- Many graph problems can be solved using a search method.
 - Path from one vertex to another.
 - Is the graph connected?
 - Find a spanning tree.
 - •
- Commonly used search methods:
 - Breadth-first search.
 - Depth-first search.

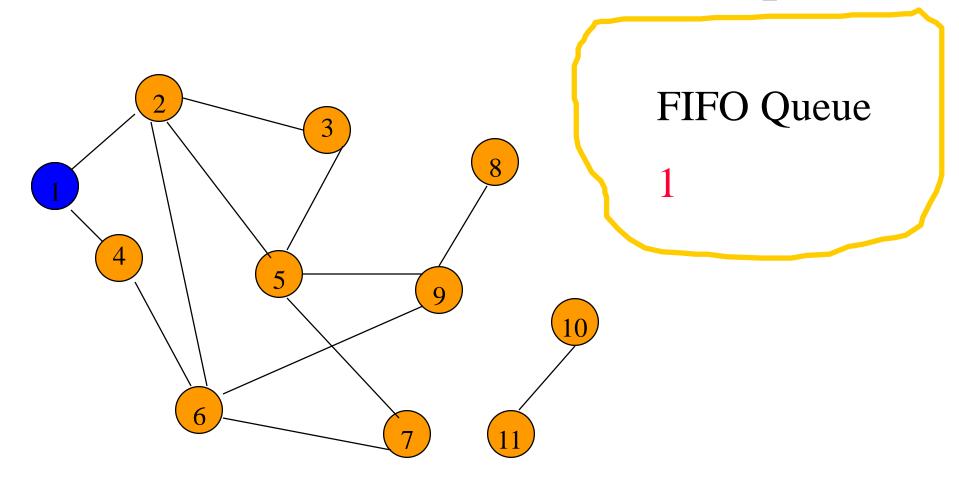
Breadth-First Search

Visit start vertex and put into a FIFO queue.

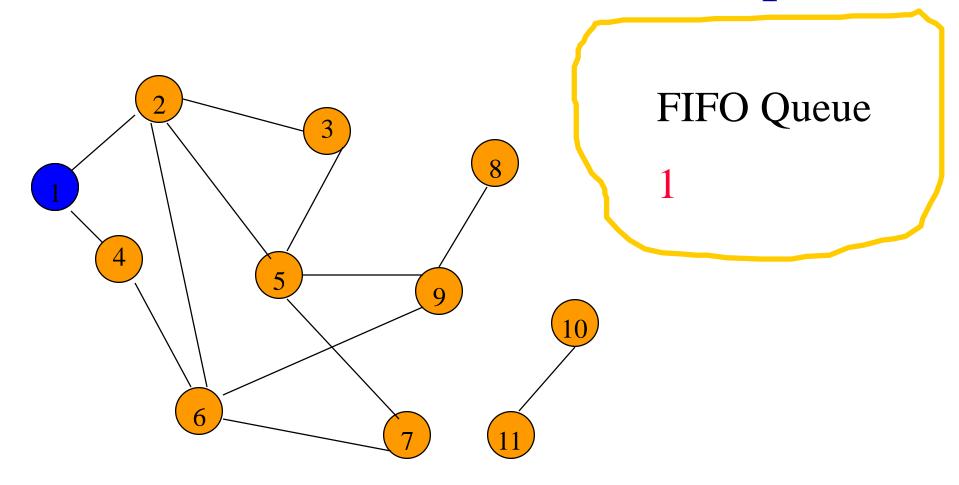
• Repeatedly remove a vertex from the queue, visit its unvisited adjacent vertices, put newly visited vertices into the queue until the queue is empty.



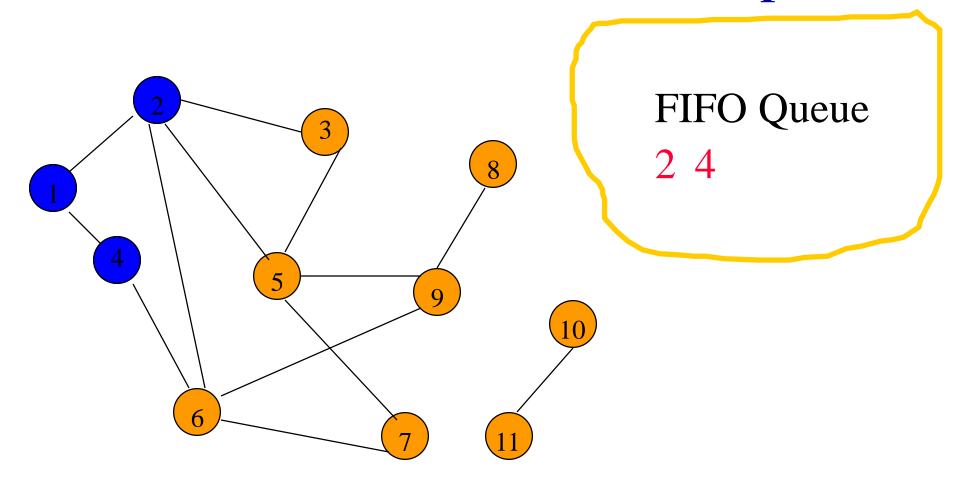
Start search at vertex 1.



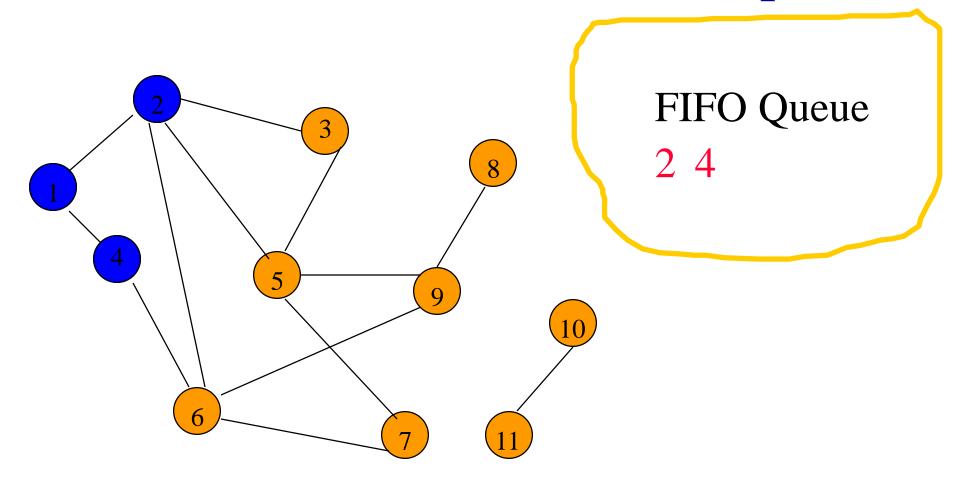
Visit/mark/label start vertex and put in a FIFO queue.



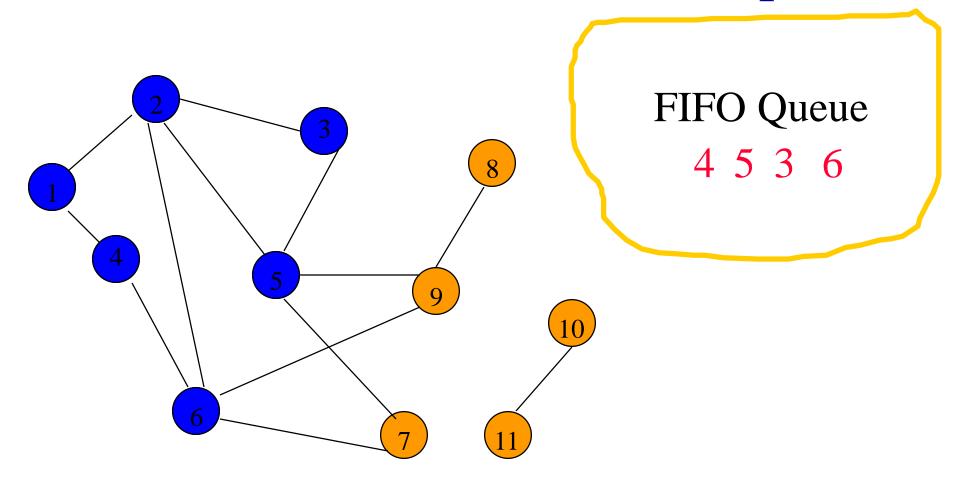
Remove 1 from Q; visit adjacent unvisited vertices; put in Q.



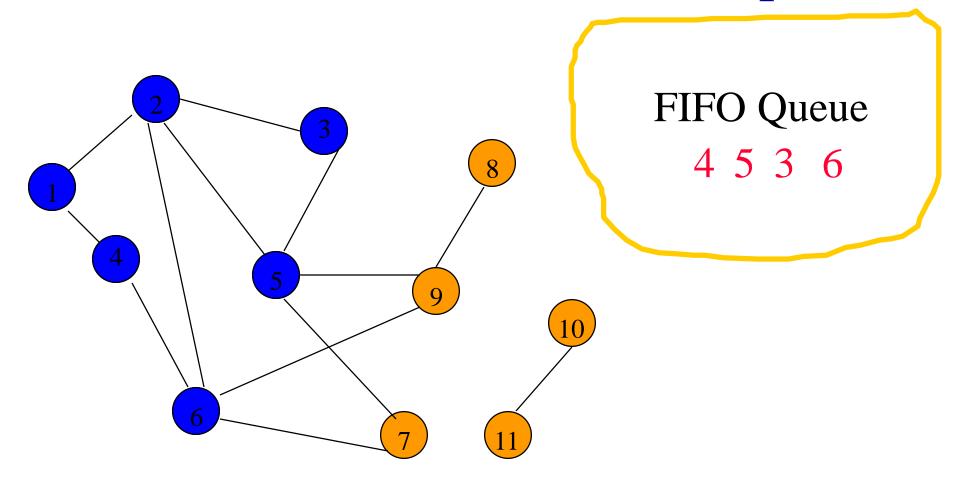
Remove 1 from Q; visit adjacent unvisited vertices; put in Q.



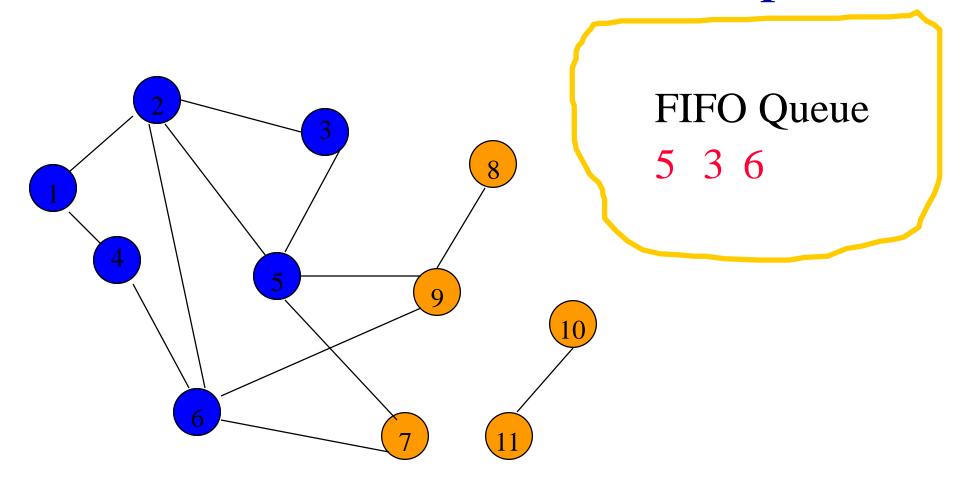
Remove 2 from Q; visit adjacent unvisited vertices; put in Q.



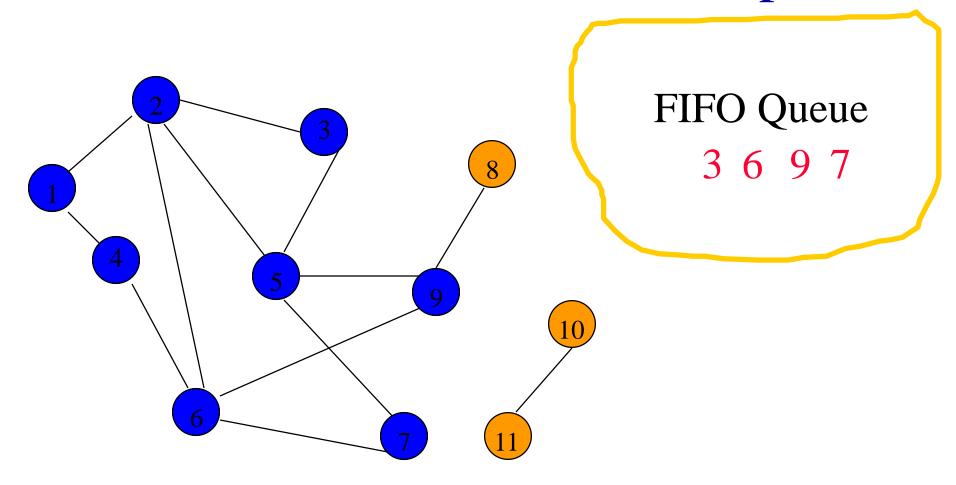
Remove 2 from Q; visit adjacent unvisited vertices; put in Q.



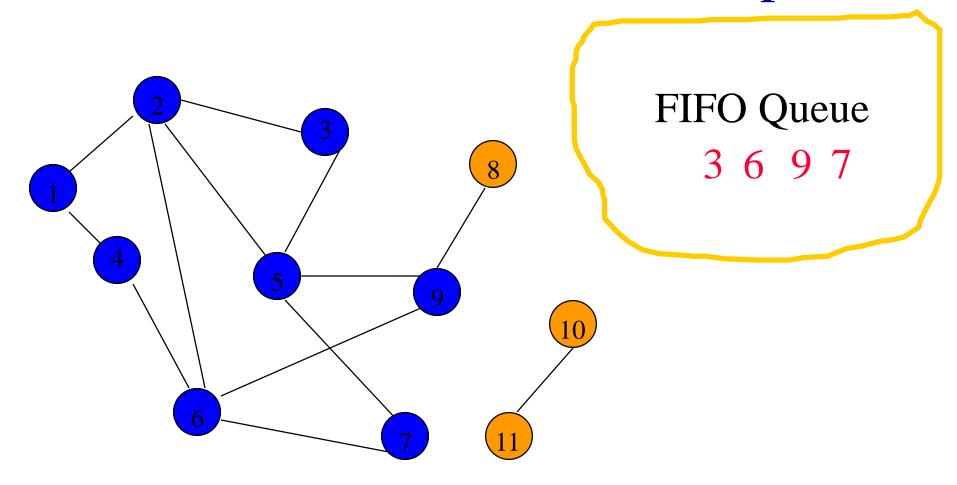
Remove 4 from Q; visit adjacent unvisited vertices; put in Q.



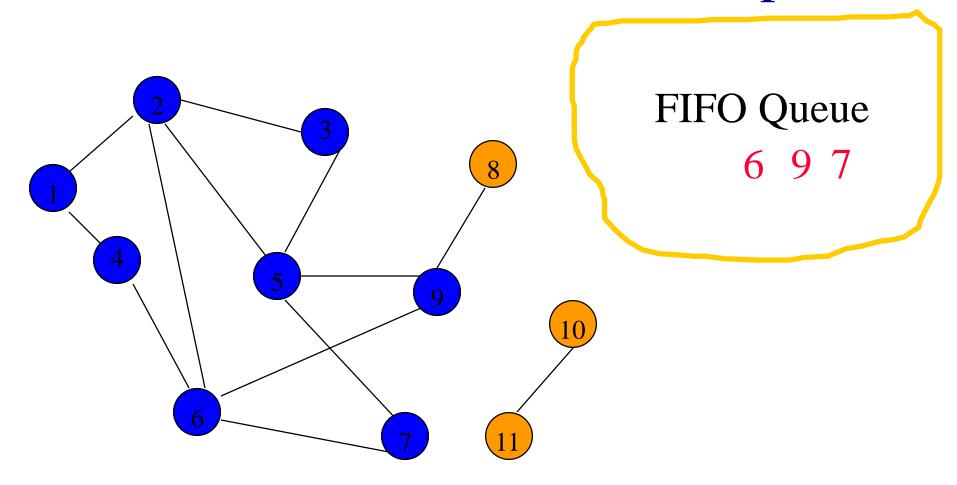
Remove 5 from Q; visit adjacent unvisited vertices; put in Q.



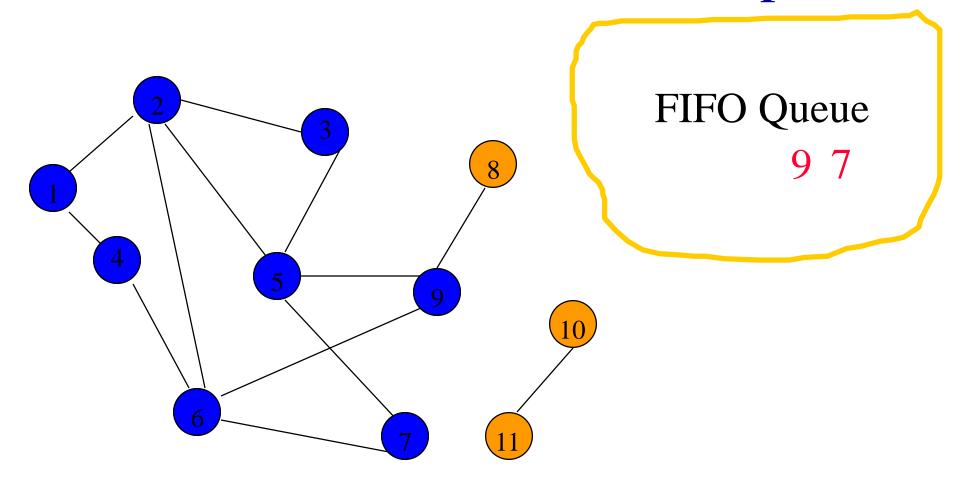
Remove 5 from Q; visit adjacent unvisited vertices; put in Q.



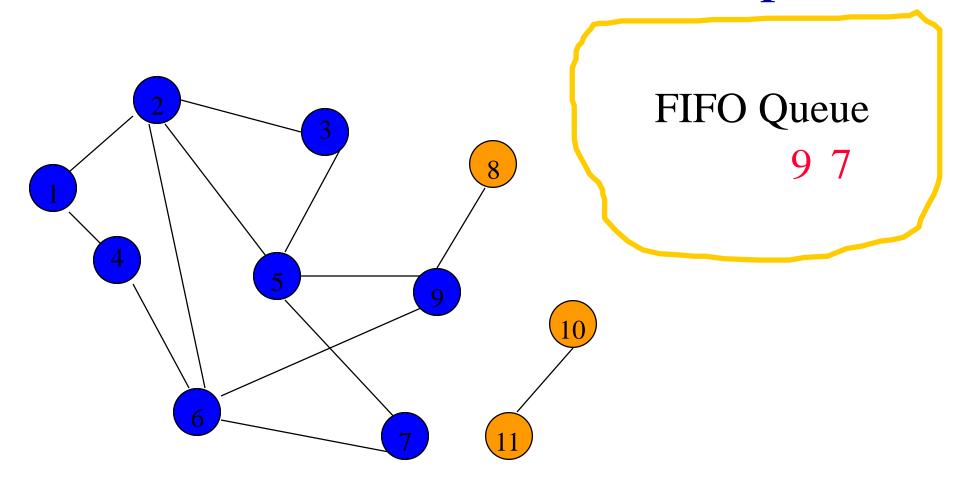
Remove 3 from Q; visit adjacent unvisited vertices; put in Q.



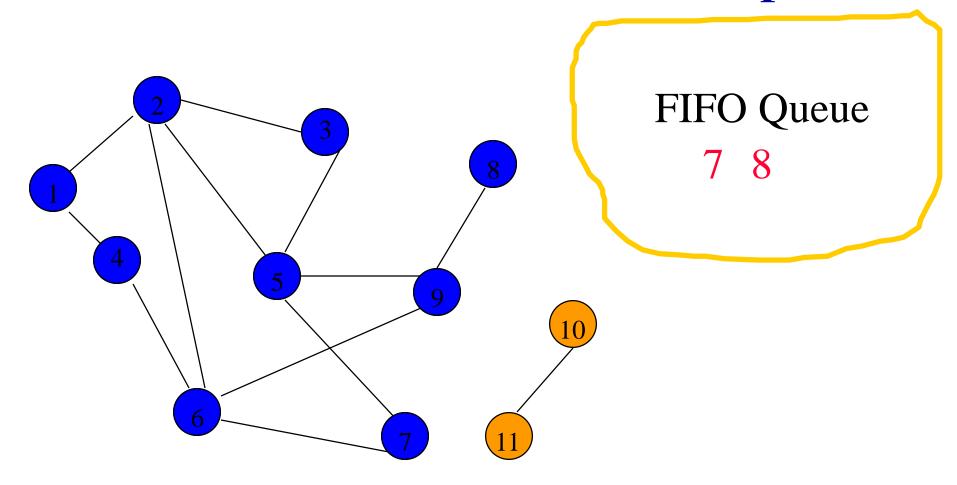
Remove 6 from Q; visit adjacent unvisited vertices; put in Q.



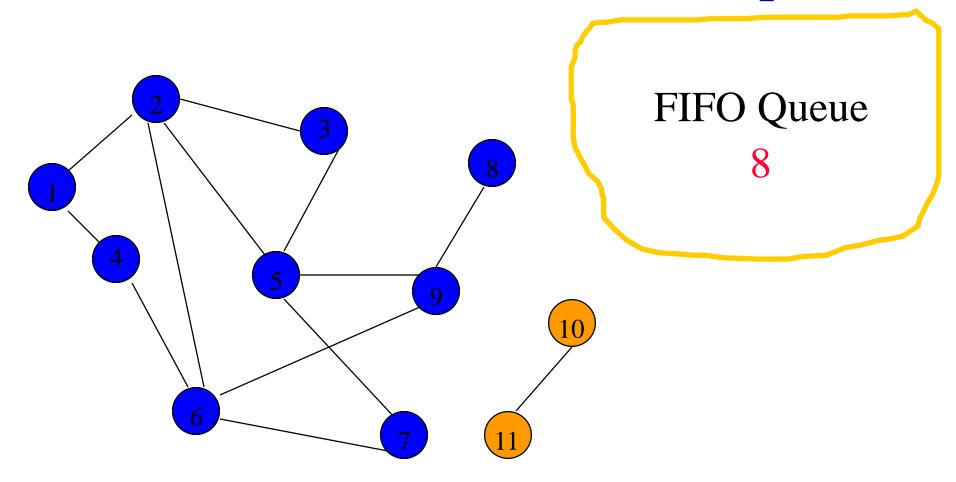
Remove 6 from Q; visit adjacent unvisited vertices; put in Q.



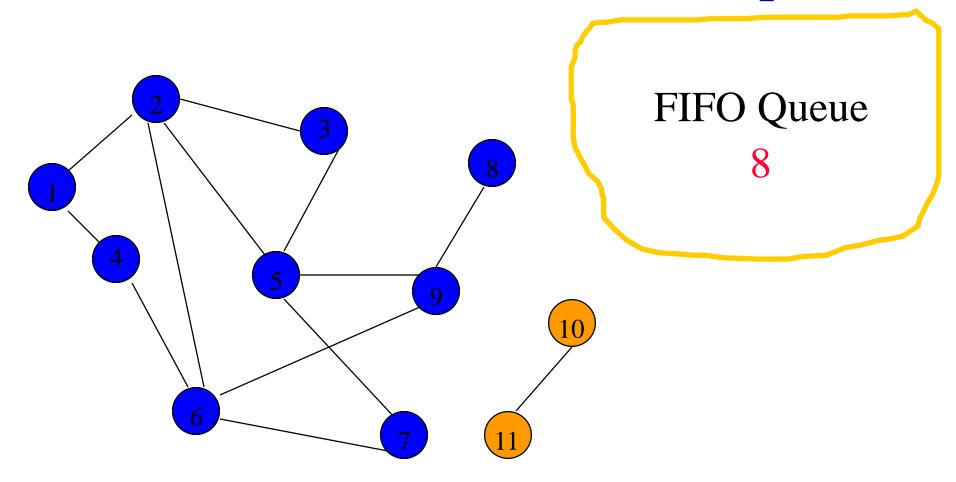
Remove 9 from Q; visit adjacent unvisited vertices; put in Q.



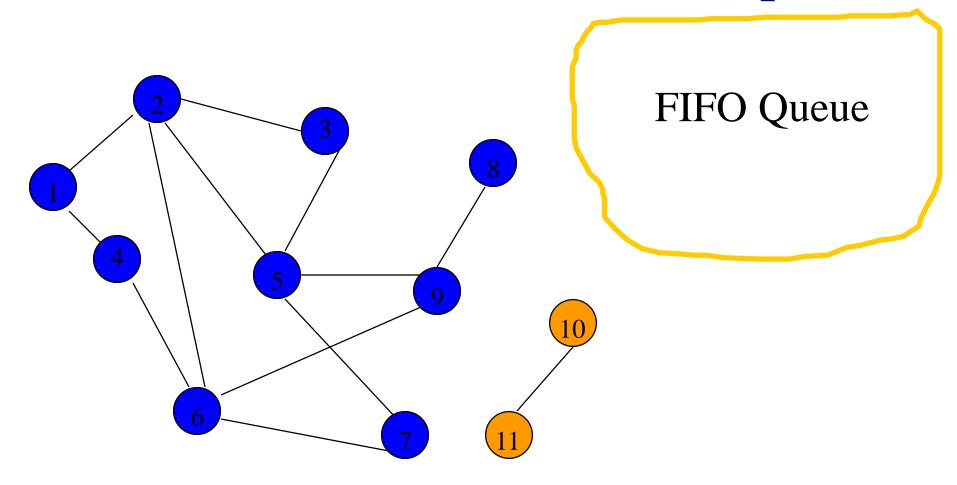
Remove 9 from Q; visit adjacent unvisited vertices; put in Q.



Remove 7 from Q; visit adjacent unvisited vertices; put in Q.



Remove 8 from Q; visit adjacent unvisited vertices; put in Q.



Queue is empty. Search terminates.

Breadth-First Search Property

• All vertices reachable from the start vertex (including the start vertex) are visited.

Time Complexity



- Each visited vertex is put on (and removed from) the queue exactly once.
- When a vertex is removed from the queue, we examine its adjacent vertices.
 - O(n) if adjacency matrix used
 - O(vertex degree) if adjacency lists used
- Total time, when adjacency matrix is used:
 - O(mn), where m is number of vertices in the component that is searched.

Time Complexity



When adjacency lists are used:

O(n + sum of degrees of the vertices in the component)

= O(n + number of edges in the component)

Path From Vertex v To Vertex u

- Start a breadth-first search at vertex v.
- Terminate when vertex u is visited or when
 Q becomes empty (whichever occurs first).
- Time
 - $O(n^2)$ when adjacency matrix used
 - O(n+e) when adjacency lists used (e is number of edges)

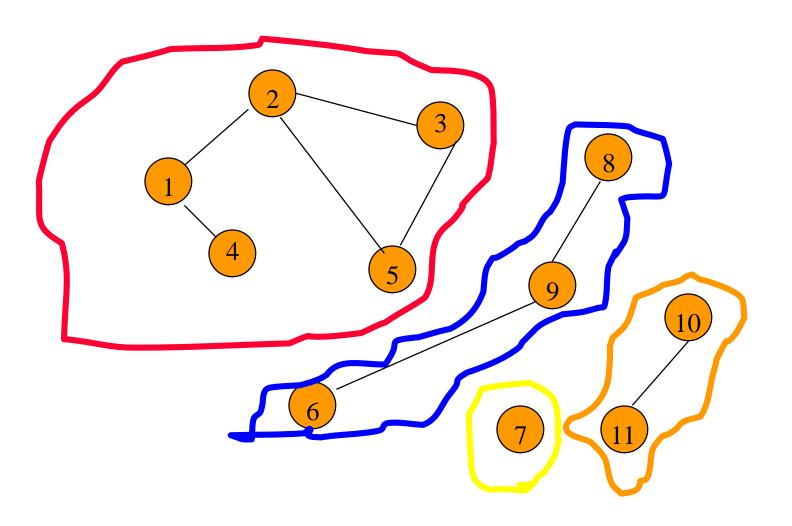
Is The Graph Connected?

- Start a breadth-first search at any vertex of the graph.
- Graph is connected iff all n vertices get visited.
- Time
 - $O(n^2)$ when adjacency matrix used
 - O(n+e) when adjacency lists used (e is number of edges)

Connected Components

- Start a breadth-first search at any as yet unvisited vertex of the graph.
- Newly visited vertices (plus edges between them) define a component.
- Repeat until all vertices are visited.

Connected Components

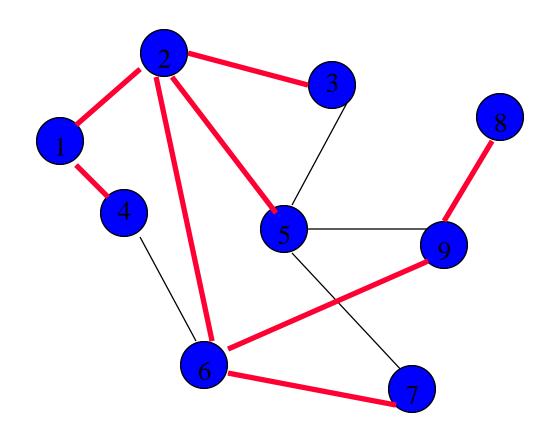


Time Complexity



- $O(n^2)$ when adjacency matrix used
- O(n+e) when adjacency lists used (e is number of edges)

Spanning Tree



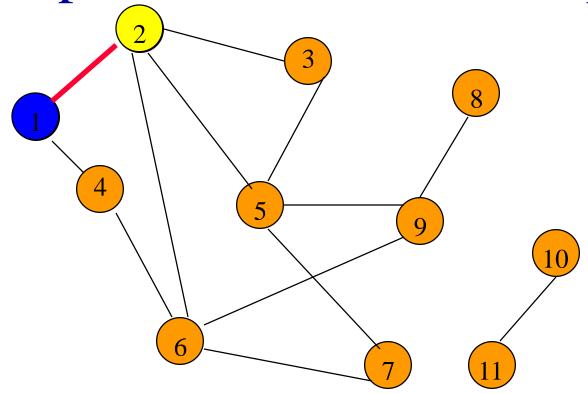
Breadth-first search from vertex 1.
Breadth-first spanning tree.

Spanning Tree

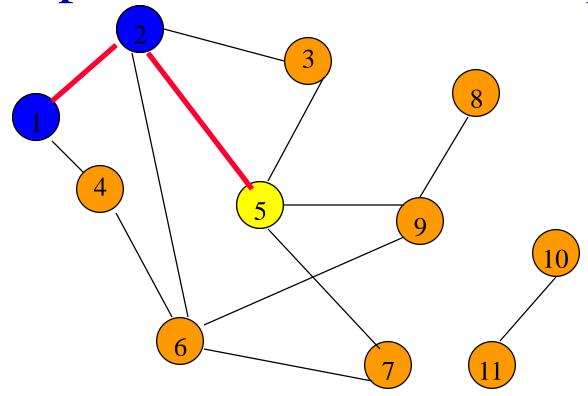
- Start a breadth-first search at any vertex of the graph.
- If graph is connected, the n-1 edges used to get to unvisited vertices define a spanning tree (breadth-first spanning tree).
- Time
 - $O(n^2)$ when adjacency matrix used
 - O(n+e) when adjacency lists used (e is number of edges)

Depth-First Search

```
depthFirstSearch(v)
 Label vertex v as reached.
 for (each unreached vertex u
                      adjacenct from v)
   depthFirstSearch(u);
```

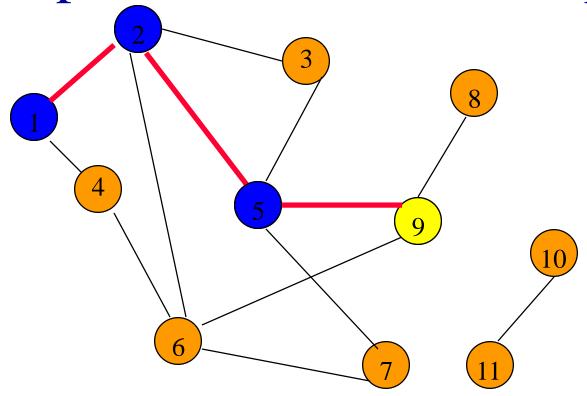


- Start search at vertex 1.
- Label vertex 1 and do a depth first search from either 2 or 4.
- Suppose that vertex 2 is selected.



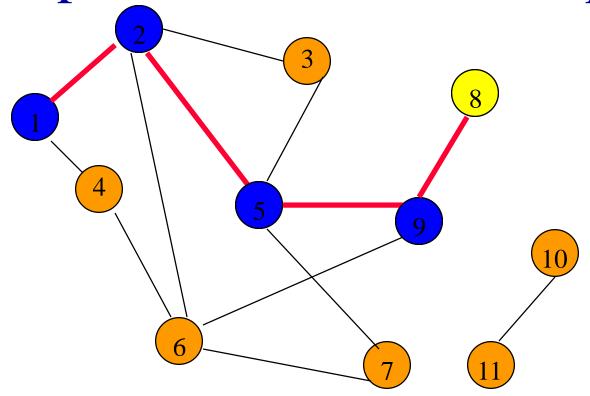
Label vertex 2 and do a depth first search from either 3, 5, or 6.

Suppose that vertex 5 is selected.



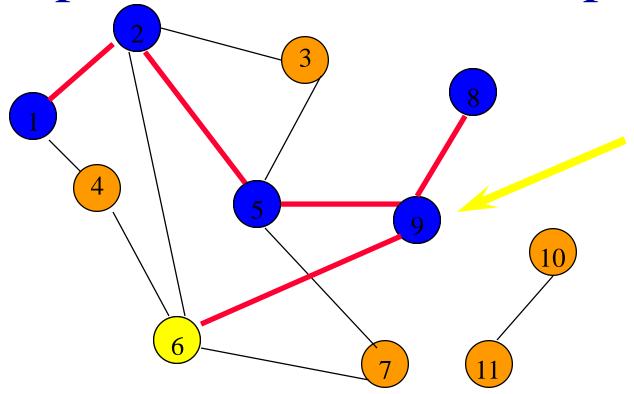
Label vertex 5 and do a depth first search from either 3, 7, or 9.

Suppose that vertex 9 is selected.



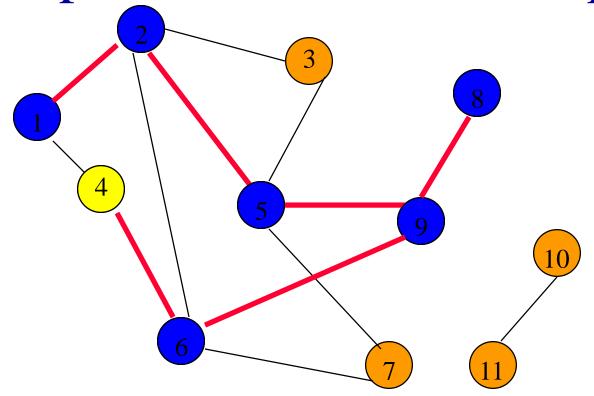
Label vertex 9 and do a depth first search from either 6 or 8.

Suppose that vertex 8 is selected.



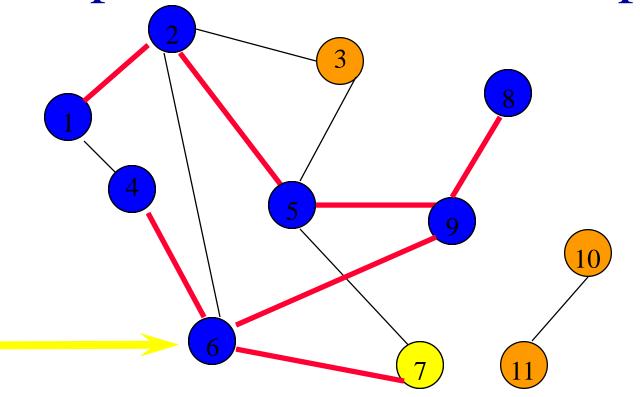
Label vertex 8 and return to vertex 9.

From vertex 9 do a dfs(6).



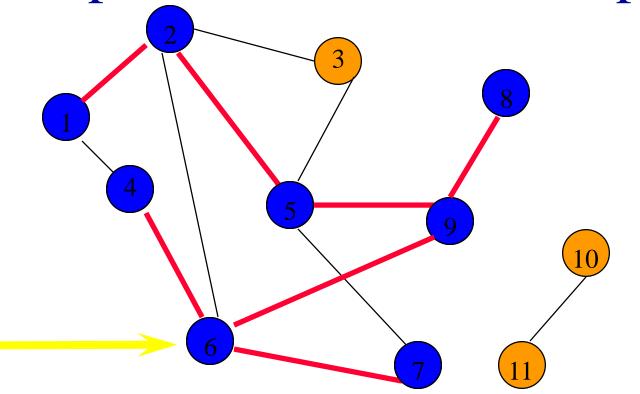
Label vertex 6 and do a depth first search from either 4 or 7.

Suppose that vertex 4 is selected.

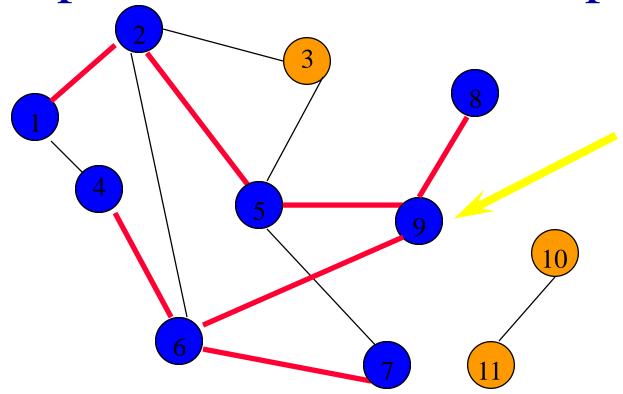


Label vertex 4 and return to 6.

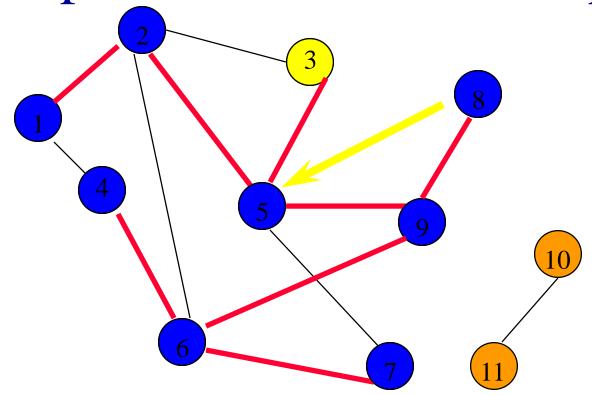
From vertex 6 do a dfs(7).

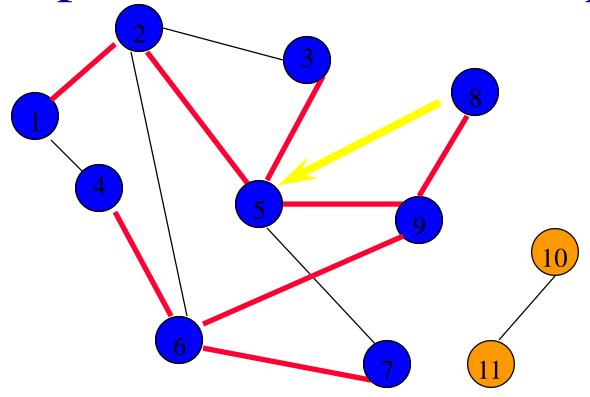


Label vertex 7 and return to 6. Return to 9.



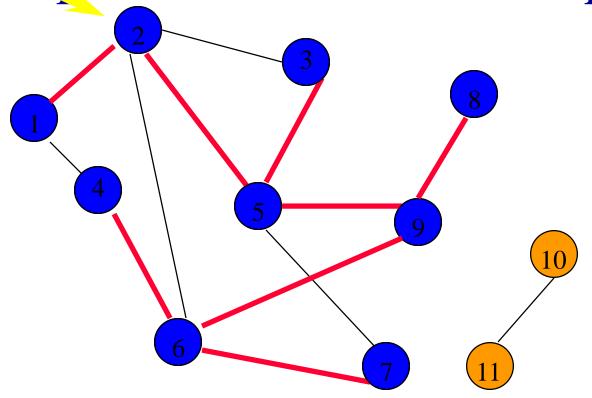
Return to 5.



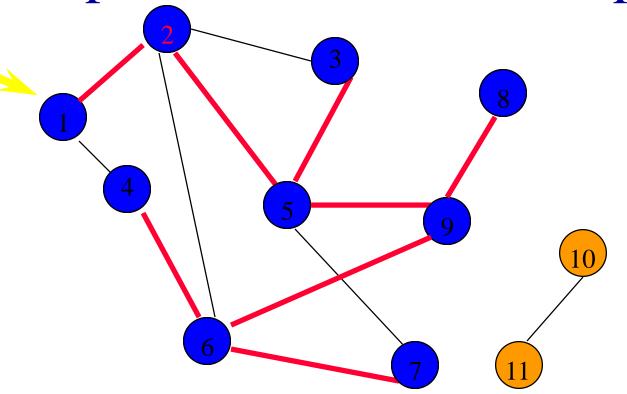


Label 3 and return to 5.

Return to 2.



Return to 1.



Return to invoking method.

Depth-First Search Properties

- Same complexity as BFS.
- Same properties with respect to path finding, connected components, and spanning trees.
- Edges used to reach unlabeled vertices define a depth-first spanning tree when the graph is connected.
- There are problems for which bfs is better than dfs and vice versa.