

4190.204

Data Structures

Fall 2018

Instructor

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Course web site

- On etl.snu.ac.kr

Contains

Schedule, lecture slides, assignments/homework, and other course related information

Text books

- **Data Structures, Algorithms, and Applications in Java**, Second Edition
by **Sartaj Sahni**, Silicon Press, 2005.

or

- **Data Abstraction and Problem Solving with Java**, Third Edition
by **Janet Prichard and Frank M. Carrano**, Pearson, 2011.

Prerequisite

- Java programming



Evaluation

- Assignments – 30%
- Midterm – 30%
- Final exam – 30%
- **Surprise** quizzes – 5%
- Lab Participation – 5%
- Assignments – programming problems + written exercises

Goals

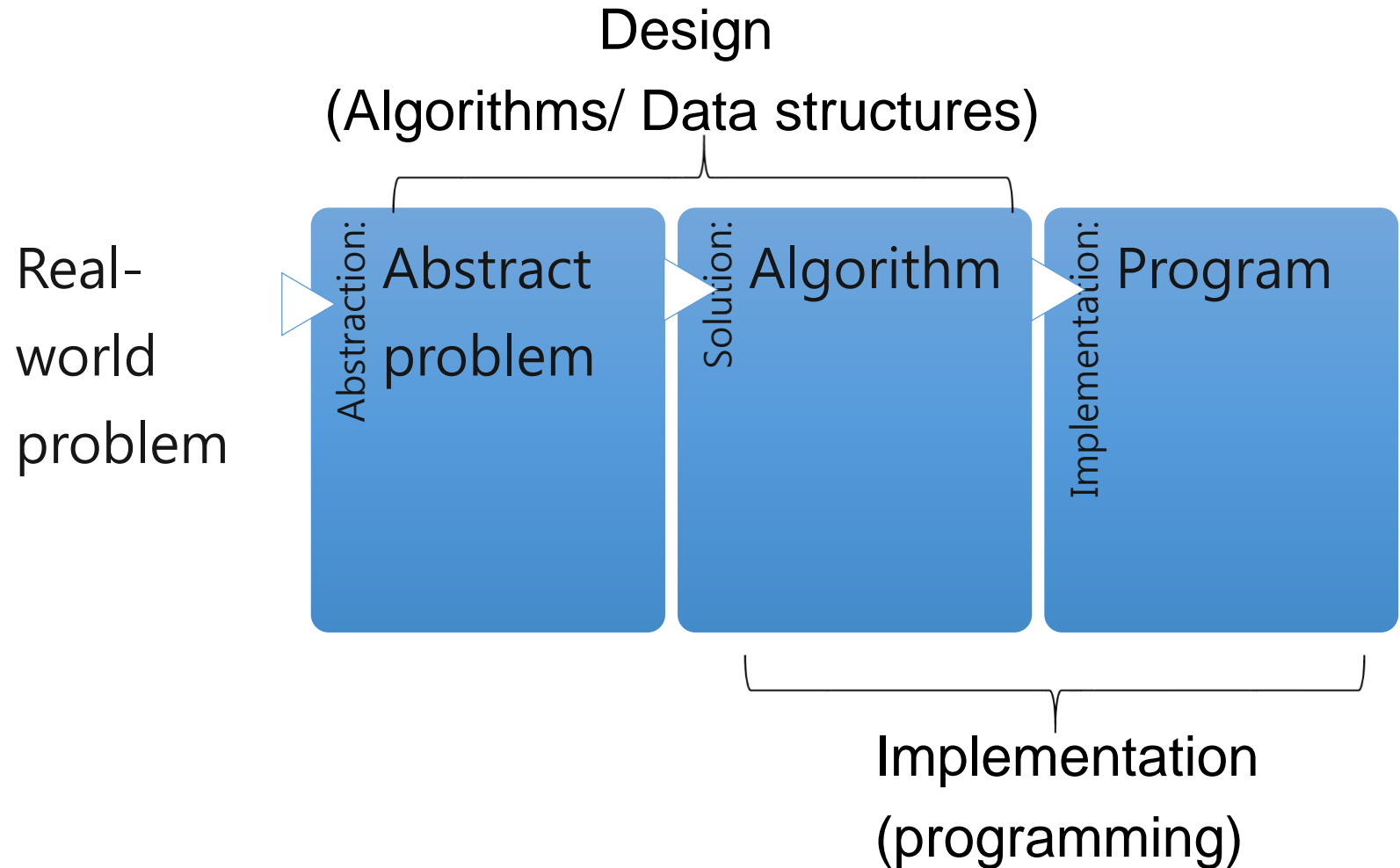
- Understanding how data structures impact the performance of programs
- Solving computational problems using some basic data structures like lists, stacks, trees, graphs etc.
- Designing and writing large programs using object-oriented design principles

What the course is about

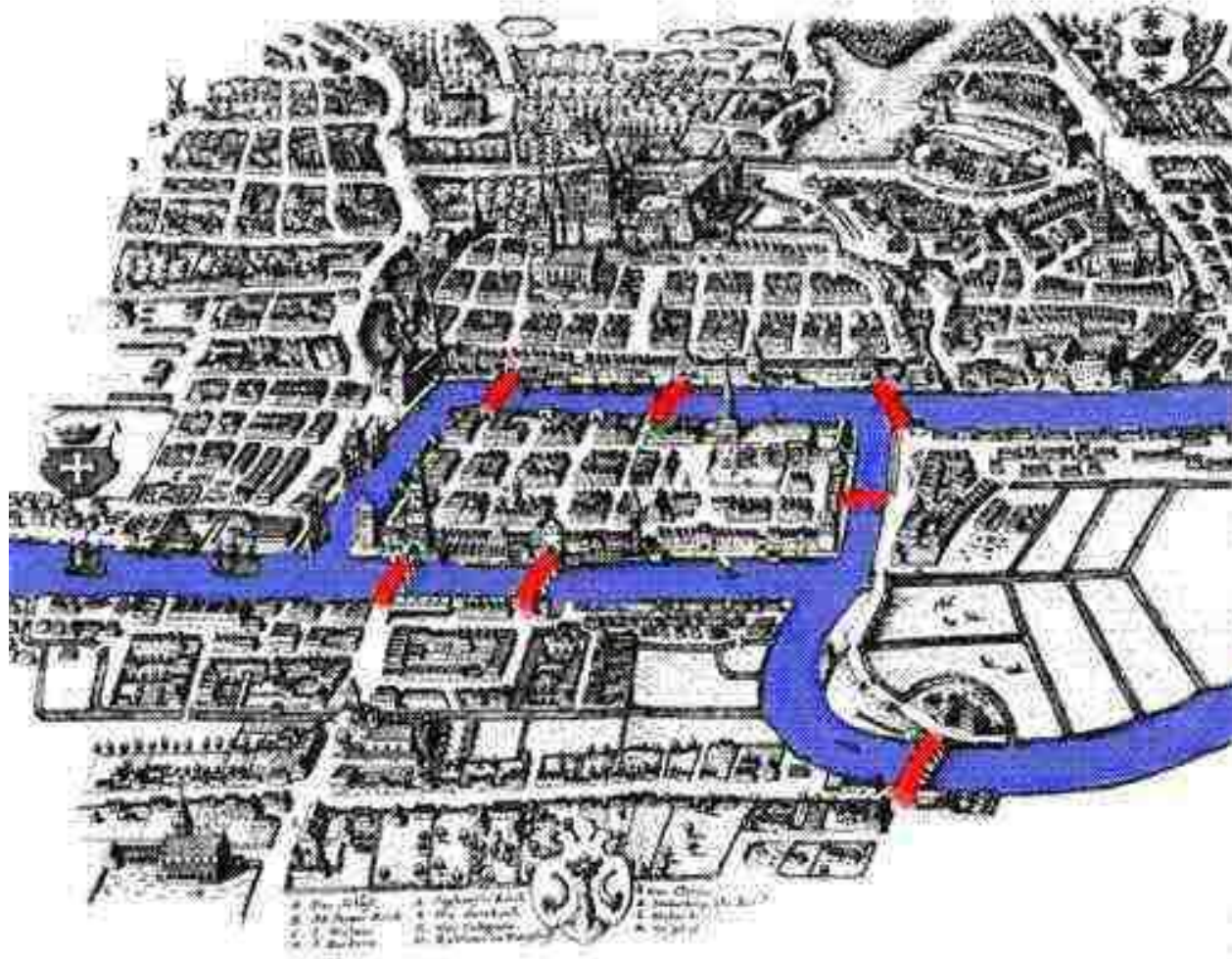
- Data structures is concerned with the representation and manipulation of data.
- All programs represent data in some way and manipulate the data.
- Data is represented in a `data structure' and manipulated through an `algorithm'.
- The study of data structures and algorithms is fundamental to Computer Science.

- Algorithm: the essence of computational procedure – step by step instructions
- Program: an implementation of an algorithm in some programming language
- Data structure: organization of data needed to solve the problem

Problem solving

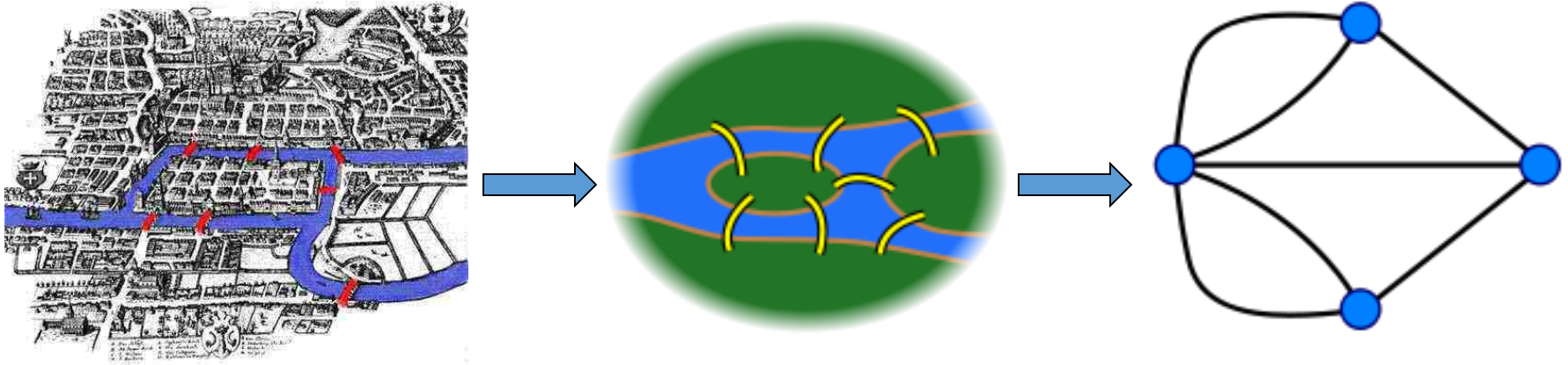


Konigsberg bridge problem



Konigsberg bridge problem

- Abstraction



No Algorithm !

Other (algorithmic) problems

- Show all possible ways of going from place A to place B by crossing at most 4 bridges
- Fastest (shortest) way to get from A to B
- Each problem may require the data to be represented in a different form

A simple problem

- Given a set of processes and their waiting times, return them in the decreasing order of their waiting times.
- Abstract problem: Given a set of numbers, sort them decreasing (or equivalently in increasing) order.

Sorting algorithm

- Let's use Insertion Sort
 - Start with a sequence of size 1
 - Insert the remaining elements, one at a time
 - **Example:** Sort 7, 3, 5, 6, 1
 - Start with 7 and insert 3 => 3, 7
 - Insert 5 => 3, 5, 7
 - Insert 6 => 3, 5, 6, 7
 - Insert 1 => 1, 3, 5, 6, 7

Insertion sort implementation

```
public static void insertionSort(int [] a) {  
    for (int i = 1; i < a.length; i++) {  
        // insert a[i] into a[0:i-1]  
        int t = a[i];  
        int j;  
        for (j = i - 1; j >= 0 && t < a[j]; j--)  
            a[j + 1] = a[j];  
        a[j + 1] = t; }  
}
```


Performance measures/ Complexity

- Space/Memory
- Time
 - Count a particular operation
 - Count number of steps
 - Asymptotic complexity

Comparison Count

```
for (int i = 1; i < a.length; i++) {  
    int t = a[i];  
    int j;  
    for (j = i - 1; j >= 0 && t < a[j]; j--)  
        a[j + 1] = a[j];  
    a[j + 1] = t;  
}
```

Comparison Count

- Pick an instance characteristic, n
 - Eg. $n = a.length$ for insertion sort
- Determine count as a function of this instance characteristic.

The exact number of comparisons depends on the instance ***a***, and varies greatly from instance to instance.

Comparison Count

- Worst-case count = maximum count
- Best-case count = minimum count
- Average count over all instances.

Worst-Case Comparison Count

```
for (j = i - 1; j >= 0 && t < a[j]; j--)  
    a[j + 1] = a[j];
```

$a = [1, 2, 3, 4]$ and $t = 0 \Rightarrow 4$ comparisons

$a = [1, 2, 3, \dots, i]$ and $t = 0 \Rightarrow i$ comparisons

Worst-Case Comparison Count

```
for (int i = 1; i < n; i++)  
    for (j = i - 1; j >= 0 && t < a[j]; j--)  
        a[j + 1] = a[j];
```

total comparisons = $1 + 2 + 3 + \dots + (n-1) =$
 $(n-1)n/2$
(in the worst case)

Insertion sort

- What is the best case?
 - #Comparisons?
- Average case?
 - #Comparisons?

Counting comparisons

- The worst case number of comparisons performed by Insertion Sort can be reduced significantly by using 'Binary Insertion Sort'
- But the overall running time doesn't improve significantly
- Counting only comparisons is not enough

What about counting 'data movement'?

Sorting out sorting

- <http://www.youtube.com/watch?v=YvTW7341kpA>