



공지사항

- 기말고사 일정
 - 기말고사1: 6/8 (월) 17:00~18:15
 - 장소: 302-106/107
 - 기말고사2: 6/15(월) 17:00~18:15
 - 장소: 302-105
- 성적
 - 절대평가 기준

과제 (%)	기말고사1 (%)	기말고사2 (%)	출석 (%)	합계 (%)
35	30	30	5	100

Summary of previous lecture

- Deutsch's algorithm
 - Pattern of the algorithm: Initialization → **Superposition** of multiple possibilities → Processing → **Interference** of the multiple outputs → Measurement of the output
- Reversible gate
 - Conversion of digital gates with reversible gates generally incurs the **auxiliary qubits** and **garbage qubits**.
 - Garbage qubit generally gets entangled with other qubits and can be detrimental to the overall circuits.
 - By **un-computing**, the garbage qubits can be un-entangled and the auxiliary qubits can be re-cycled.
- Factoring algorithm (Shor's algorithm)
 - Origin of quantum speed-up
 - Simultaneous calculation of $a^x \pmod{N}$ for x from 0 to $2^{2 \cdot \text{ceil}(\log_2 N)}$
 - Fast Fourier transform



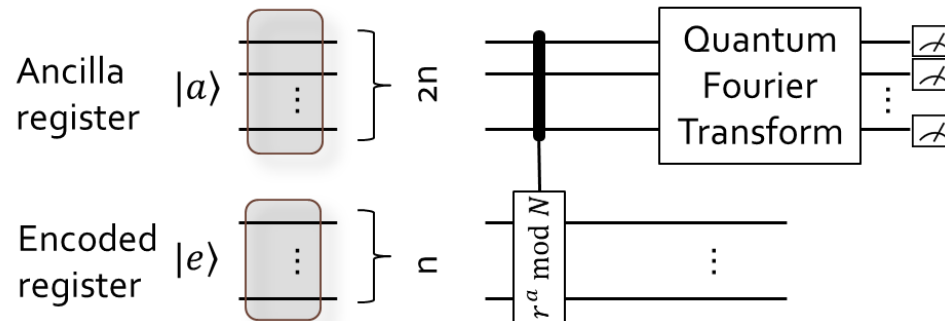
Shor's Algorithm (Factoring algorithm)

- Chapter 5
- Example for factorization of number 15
 - Choose a random number that has the following properties
 - No common divisor with 15 (target of factorization)
 - Smaller than 15 (target of factorization)
 - Ex) $r = 7$
 - Calculate $r^a \pmod{15}$ for all a between 0 and 255
 - Find the period among these values

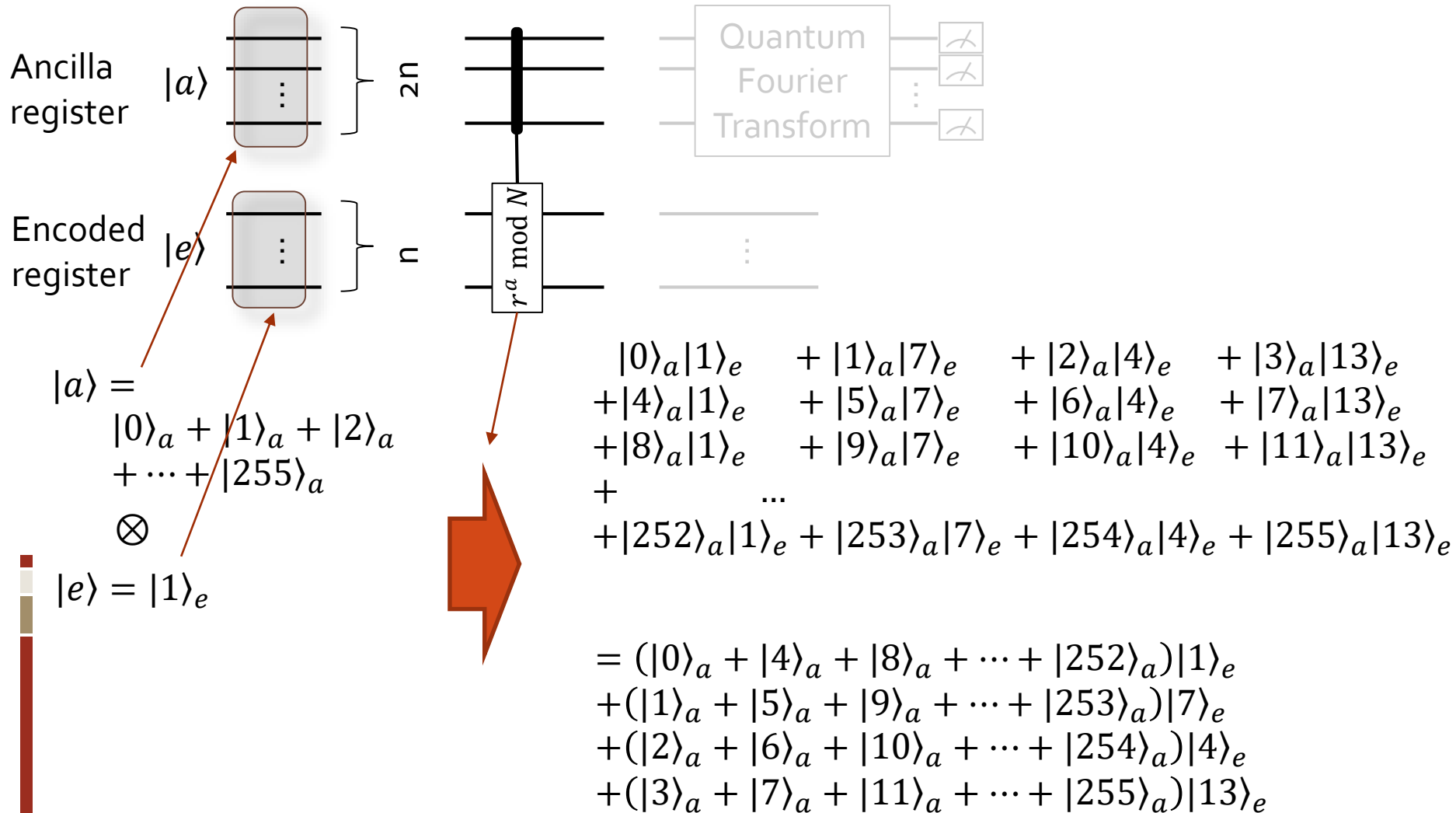
Ex)

7^0	7^1	7^2	7^3	7^4	7^5	7^6	7^7	7^8	7^9	7^{10}	7^{11}	7^{12}	...
1	7	4	13	1	7	4	13	1	7	4	13	1	...

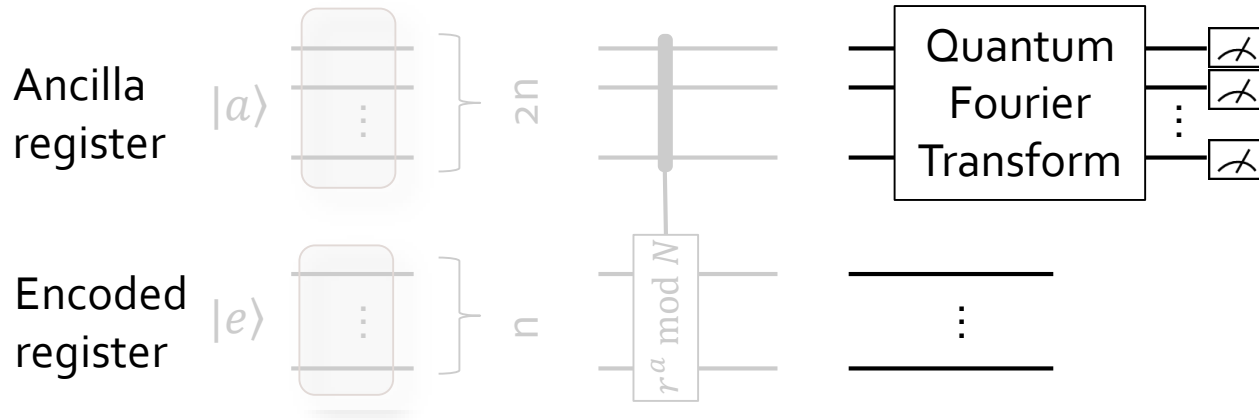
- $7^4 = 1 \pmod{15} \Rightarrow 7^4 - 1 = (7^2 - 1)(7^2 + 1) = N * 15$
- $\gcd(7^2 - 1, 15) = 3, \gcd(7^2 + 1, 15) = 5$



Analysis of Factorization Process I

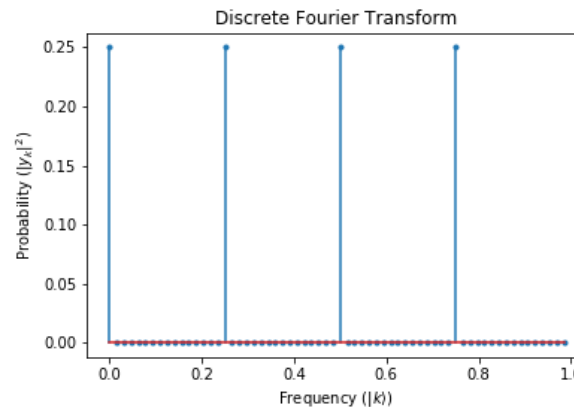
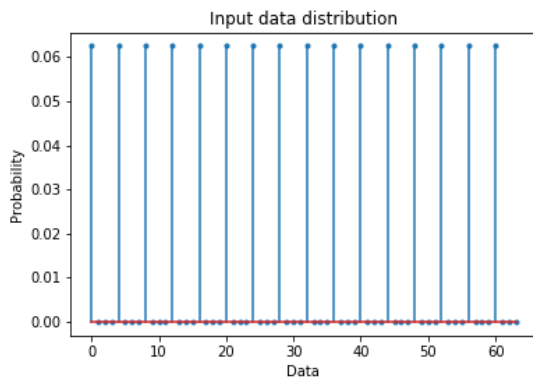


Analysis of Factorization Process II



$$\begin{aligned}
 &= (|0\rangle_a + |4\rangle_a + |8\rangle_a + \dots + |252\rangle_a)|1\rangle_e \\
 &+ (|1\rangle_a + |5\rangle_a + |9\rangle_a + \dots + |253\rangle_a)|7\rangle_e \\
 &+ (|2\rangle_a + |6\rangle_a + |10\rangle_a + \dots + |254\rangle_a)|4\rangle_e \\
 &+ (|3\rangle_a + |7\rangle_a + |11\rangle_a + \dots + |255\rangle_a)|13\rangle_e
 \end{aligned}$$

$$\begin{aligned}
 &= (|k=0\rangle + |k=64\rangle + |k=128\rangle + |k=192\rangle)_y |1\rangle_e \\
 &+ (|k=0\rangle + e^{i3\pi/2}|k=64\rangle + e^{i\pi}|k=128\rangle + e^{i\pi/2}|k=192\rangle)_y |7\rangle_e \\
 &+ (|k=0\rangle + e^{i\pi}|k=64\rangle + e^{i2\pi}|k=128\rangle + e^{i\pi}|k=192\rangle)_y |4\rangle_e \\
 &+ (|k=0\rangle + e^{i\pi/2}|k=64\rangle + e^{i\pi}|k=128\rangle + e^{i3\pi/2}|k=192\rangle)_y |13\rangle_e
 \end{aligned}$$



The above plots are generated using 64 inputs instead of 256 for readability

If $|k = 192\rangle$ quantum state is measured, the corresponding frequency is $192/256 = 3/4$. From this value, we can learn that there exist a high probability that the period is 4.

Discrete Fourier Transform

- N -dimensional vector space composed of $[x_0, x_1, \dots, x_{N-1}]$

- One option for the basis is $[0, \dots, 0, 1, 0, \dots, 0]$

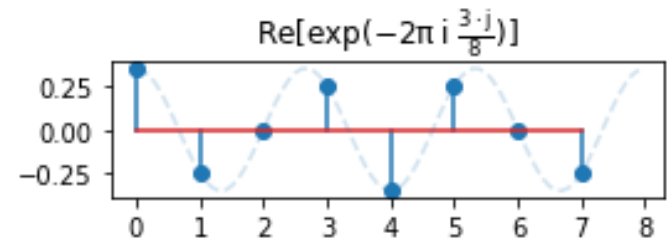
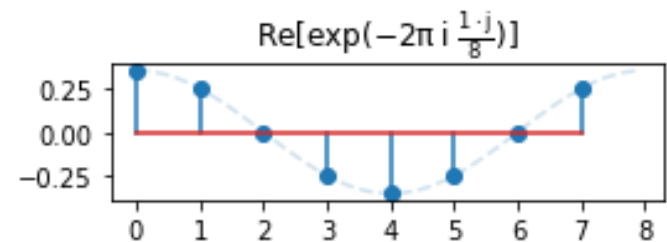
- Orthonormal basis: $|k\rangle \leftrightarrow \frac{1}{\sqrt{N}} \left[e^{-2\pi i \frac{k \cdot 0}{N}}, e^{-2\pi i \frac{k \cdot 1}{N}}, \dots, e^{-2\pi i \frac{k \cdot (N-1)}{N}} \right]$

for $0 \leq k < N$

- For example, if $N = 8$

- $k = 1 \rightarrow \frac{1}{\sqrt{8}} \left[e^{-2\pi i \frac{1 \cdot 0}{8}}, e^{-2\pi i \frac{1 \cdot 1}{8}}, \dots, e^{-2\pi i \frac{1 \cdot 7}{8}} \right]$

- $k = 3 \rightarrow \frac{1}{\sqrt{8}} \left[e^{-2\pi i \frac{3 \cdot 0}{8}}, e^{-2\pi i \frac{3 \cdot 1}{8}}, \dots, e^{-2\pi i \frac{3 \cdot 7}{8}} \right]$



- Inner product between $[x_0, x_1, \dots, x_{N-1}]$ and $[y_0, y_1, \dots, y_{N-1}]$ is defined as $\sum_{j=0}^{N-1} x_j^* y_j$.

Discrete Fourier Transform

- Orthonormality between $|k\rangle$ and $|k'\rangle$

- $$\langle k|k'\rangle = \sum_{j=0}^{N-1} \frac{1}{\sqrt{N}} e^{2\pi i \frac{k \cdot j}{N}} \frac{1}{\sqrt{N}} e^{-2\pi i \frac{k' \cdot j}{N}} = \frac{1}{N} \sum_{j=0}^{N-1} e^{2\pi i \frac{(k-k') \cdot j}{N}} = \frac{1}{N} \sum_{j=0}^{N-1} \alpha^j$$

where $\alpha = e^{2\pi i \frac{(k-k')}{N}}$

- When $k = k'$: $\alpha = 1 \rightarrow \langle k|k'\rangle = \frac{1}{N} \sum_{j=0}^{N-1} \alpha^j = \frac{1}{N} \sum_{j=0}^{N-1} 1 = 1$

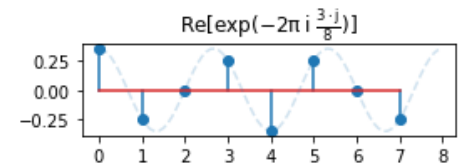
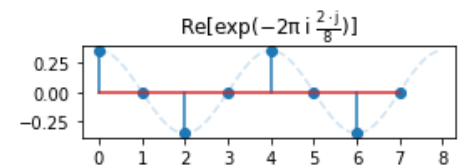
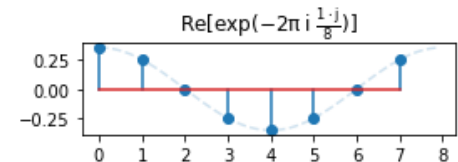
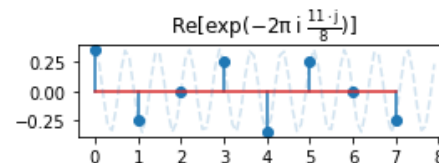
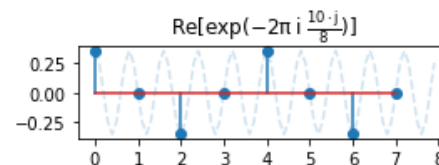
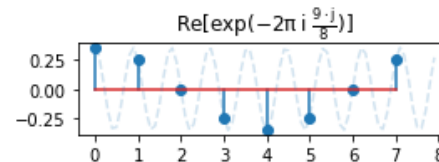
- When $k \neq k'$: $\alpha \neq 1 \rightarrow \langle k|k'\rangle = \frac{1}{N} \sum_{j=0}^{N-1} \alpha^j = \frac{1}{N} \frac{\alpha^N - 1}{\alpha - 1}$,

but because $\alpha^N = \left(e^{2\pi i \frac{(k-k')}{N}}\right)^N = 1$, $\langle k|k'\rangle = 0$

- What about $k > N$?

- It corresponds to the case for $0 \leq k' < N$ due

$$\begin{aligned} & \text{to } e^{-2\pi i \frac{(k'+N) \cdot j}{N}} = \\ & e^{-2\pi i \frac{k' \cdot j}{N}} e^{-2\pi i \frac{N \cdot j}{N}} = \\ & e^{-2\pi i \frac{k' \cdot j}{N}} \end{aligned}$$



Discrete Fourier Transform

- Any arbitrary sequence of number $[x_0, x_1, \dots, x_{N-1}] = |x\rangle$ can be decomposed into the weighted sum of basis

$$\frac{1}{\sqrt{N}} \left[e^{-2\pi i \frac{k \cdot 0}{N}}, e^{-2\pi i \frac{k \cdot 1}{N}}, \dots, e^{-2\pi i \frac{k \cdot (N-1)}{N}} \right] = |k\rangle$$

- $|x\rangle = \sum_{k=0}^{N-1} c_k |k\rangle$

- Example for $N = 4$

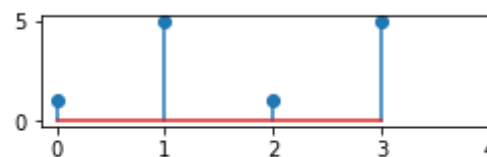
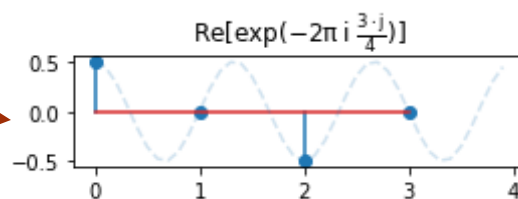
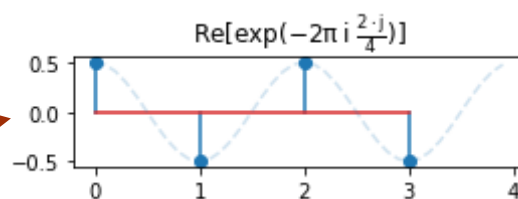
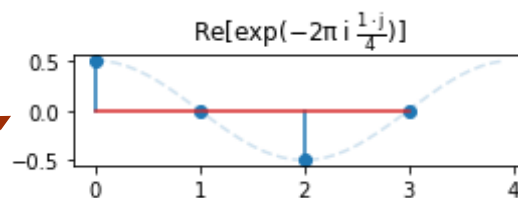
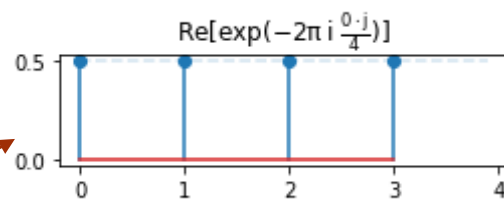
- $|k = 0\rangle = |0\rangle = \frac{1}{2} [1, 1, 1, 1]$

- $|k = 1\rangle = |1\rangle = \frac{1}{2} [1, -i, -1, i]$

- $|k = 2\rangle = |2\rangle = \frac{1}{2} [1, -1, 1, -1]$

- $|k = 3\rangle = |3\rangle = \frac{1}{2} [1, i, -1, -i]$

- $|x\rangle = [1, 5, 1, 5] = 6|0\rangle - 4|2\rangle$



Derivation of Quantum Fourier Transform I

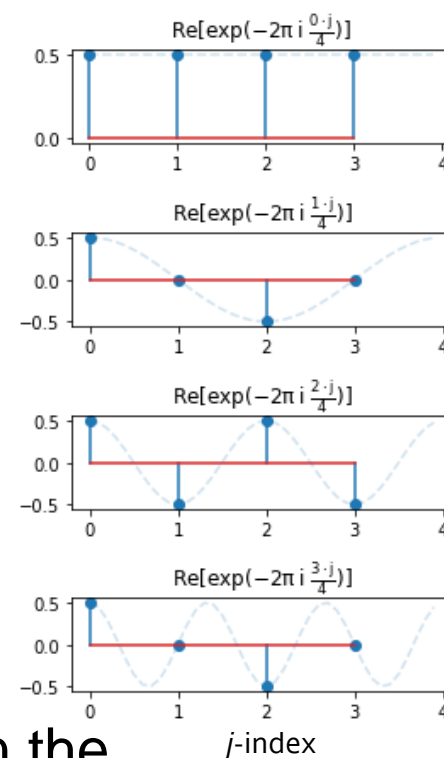
- Discrete Fourier transform finds the period embedded in the given random sequence x_0, \dots, x_{N-1}
- Example: assume that we are given a sequence x_0, x_1, x_2, x_3 composed of 4 numbers. The following calculations allow us to find y_0, y_1, y_2, y_3 that is the relative importance of the signal with the corresponding period.

$$y_0 = \frac{1}{\sqrt{4}} \left(x_0 e^{2\pi i \frac{0 \cdot 0}{4}} + x_1 e^{2\pi i \frac{0 \cdot 1}{4}} + x_2 e^{2\pi i \frac{0 \cdot 2}{4}} + x_3 e^{2\pi i \frac{0 \cdot 3}{4}} \right)$$

$$y_1 = \frac{1}{\sqrt{4}} \left(x_0 e^{2\pi i \frac{1 \cdot 0}{4}} + x_1 e^{2\pi i \frac{1 \cdot 1}{4}} + x_2 e^{2\pi i \frac{1 \cdot 2}{4}} + x_3 e^{2\pi i \frac{1 \cdot 3}{4}} \right)$$

$$y_2 = \frac{1}{\sqrt{4}} \left(x_0 e^{2\pi i \frac{2 \cdot 0}{4}} + x_1 e^{2\pi i \frac{2 \cdot 1}{4}} + x_2 e^{2\pi i \frac{2 \cdot 2}{4}} + x_3 e^{2\pi i \frac{2 \cdot 3}{4}} \right)$$

$$y_3 = \frac{1}{\sqrt{4}} \left(x_0 e^{2\pi i \frac{3 \cdot 0}{4}} + x_1 e^{2\pi i \frac{3 \cdot 1}{4}} + x_2 e^{2\pi i \frac{3 \cdot 2}{4}} + x_3 e^{2\pi i \frac{3 \cdot 3}{4}} \right)$$



- Goal of quantum Fourier transform (QFT): when the input quantum state is $x_0|0\rangle + x_1|1\rangle + x_2|2\rangle + x_3|3\rangle$, the outcome of QFT should be $y_0|0\rangle + y_1|1\rangle + y_2|2\rangle + y_3|3\rangle$.

Derivation of Quantum Fourier Transform II

$$\begin{aligned}
 & y_0|0\rangle + y_1|1\rangle + y_2|2\rangle + y_3|3\rangle \\
 &= \begin{pmatrix} x_0 e^{2\pi i \frac{0 \cdot 0}{4}} |0\rangle \\ + \\ x_1 e^{2\pi i \frac{0 \cdot 1}{4}} |0\rangle \\ + \\ x_2 e^{2\pi i \frac{0 \cdot 2}{4}} |0\rangle \\ + \\ x_3 e^{2\pi i \frac{0 \cdot 3}{4}} |0\rangle \end{pmatrix} + \begin{pmatrix} x_0 e^{2\pi i \frac{1 \cdot 0}{4}} |1\rangle \\ + \\ x_1 e^{2\pi i \frac{1 \cdot 1}{4}} |1\rangle \\ + \\ x_2 e^{2\pi i \frac{1 \cdot 2}{4}} |1\rangle \\ + \\ x_3 e^{2\pi i \frac{1 \cdot 3}{4}} |1\rangle \end{pmatrix} + \begin{pmatrix} x_0 e^{2\pi i \frac{2 \cdot 0}{4}} |2\rangle \\ + \\ x_1 e^{2\pi i \frac{2 \cdot 1}{4}} |2\rangle \\ + \\ x_2 e^{2\pi i \frac{2 \cdot 2}{4}} |2\rangle \\ + \\ x_3 e^{2\pi i \frac{2 \cdot 3}{4}} |2\rangle \end{pmatrix} + \begin{pmatrix} x_0 e^{2\pi i \frac{3 \cdot 0}{4}} |3\rangle \\ + \\ x_1 e^{2\pi i \frac{3 \cdot 1}{4}} |3\rangle \\ + \\ x_2 e^{2\pi i \frac{3 \cdot 2}{4}} |3\rangle \\ + \\ x_3 e^{2\pi i \frac{3 \cdot 3}{4}} |3\rangle \end{pmatrix} / \sqrt{4} \\
 &= x_0 \left(e^{2\pi i \frac{0 \cdot 0}{4}} |0\rangle + e^{2\pi i \frac{1 \cdot 0}{4}} |1\rangle + e^{2\pi i \frac{2 \cdot 0}{4}} |2\rangle + e^{2\pi i \frac{3 \cdot 0}{4}} |3\rangle \right) / \sqrt{4} \\
 &\quad + x_1 \left(e^{2\pi i \frac{0 \cdot 1}{4}} |0\rangle + e^{2\pi i \frac{1 \cdot 1}{4}} |1\rangle + e^{2\pi i \frac{2 \cdot 1}{4}} |2\rangle + e^{2\pi i \frac{3 \cdot 1}{4}} |3\rangle \right) / \sqrt{4} \\
 &\quad + x_2 \left(e^{2\pi i \frac{0 \cdot 2}{4}} |0\rangle + e^{2\pi i \frac{1 \cdot 2}{4}} |1\rangle + e^{2\pi i \frac{2 \cdot 2}{4}} |2\rangle + e^{2\pi i \frac{3 \cdot 2}{4}} |3\rangle \right) / \sqrt{4} \\
 &\quad + x_3 \left(e^{2\pi i \frac{0 \cdot 3}{4}} |0\rangle + e^{2\pi i \frac{1 \cdot 3}{4}} |1\rangle + e^{2\pi i \frac{2 \cdot 3}{4}} |2\rangle + e^{2\pi i \frac{3 \cdot 3}{4}} |3\rangle \right) / \sqrt{4}
 \end{aligned}$$

Note that the initial quantum state is $x_0|0\rangle + x_1|1\rangle + x_2|2\rangle + x_3|3\rangle$.

Then QFT is equivalent to unitary transformation

$$|j\rangle \rightarrow \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} e^{2\pi i \frac{k \cdot j}{N}} |k\rangle.$$

Summary of Quantum Fourier Transform

- Discrete Fourier transform (DFT)
 - Input data for DFT: x_0, \dots, x_{N-1}
 - Output data of DFT: $y_k = \frac{1}{\sqrt{N}} \sum_{j=0}^{N-1} x_j e^{2\pi i j k / N}$
- Quantum Fourier transform (QFT)
 - Input quantum state: each input data is used as the probability amplitude of the corresponding basis $\sum_{j=0}^{N-1} x_j |j\rangle$
 - Output quantum state: has the output of DFT as the probability amplitude of the corresponding basis $\sum_{k=0}^{N-1} y_k |k\rangle$
 - $\sum_{j=0}^{N-1} x_j |j\rangle \rightarrow \sum_{k=0}^{N-1} y_k |k\rangle$
- Implementation of QFT circuit
 - Need a quantum circuit that can transform the basis ket $|0\rangle, \dots, |N-1\rangle$ of the input quantum state in the following way: $|j\rangle \rightarrow \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} e^{2\pi i \frac{k \cdot j}{N}} |k\rangle$
 - Circuit example for QFT where $R_k = \begin{bmatrix} 1 & 0 \\ 0 & e^{2\pi i / 2^k} \end{bmatrix}$

