The STEX3 Manual *

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Abstract

STEX is a collection of LaTeX package that allow to markup documents semantically without leaving the document format, essentially turning LaTeX into a document format for mathematical knowledge management (MKM). STeX augments LaTeX with

- Semantic macros that denote and distinguish between mathematical concepts, operators, etc. independent of their notational presentation,
- A powerful module system that allows for authoring and importing individual
 fragments containing document text and/or semantic macros, independent of

 and without hard coding directory paths relative to the current document,
- A mechanism for exporting ST_EX documents to (modular) XHTML, preserving all the semantic information for semantically informed knowledge management services.

This is the user manual for the STEX package and associated software. It is primarily directed at end-users who want to use STEX to author semantically enriched documents. For the full documentation, see the STEX documentation

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Boxes like this one contain implementation details that are mostly relevant for more advanced use cases, might be useful to know when debugging, or might be good to know to better understand how something works. They can easiyl be skipped on a first read.



 $\stackrel{\longleftarrow}{M} \stackrel{\longrightarrow}{\longrightarrow} \text{Boxes like this one explain how some STEX concept relates to the MMT/OMDoc system, philosophy or language.}$

Chapter 1

What is STEX?

Formal systems for mathematics (such as interactive theorem provers) have the potential to significantly increase both the accessibility of published knowledge, as well as the confidence in its veracity, by rendering the precise semantics of statements machine actionable. This allows for a plurality of added-value services, from semantic search up to verification and automated theorem proving. Unfortunately, their usefulness is hidden behind severe barriers to accessibility; primarily related to their surface languages reminiscent of programming languages and very unlike informal standards of presentation.

STEX minimizes this gap between informal and formal mathematics by integrating formal methods into established and widespread authoring workflows, primarily LATEX, via non-intrusive semantic annotations of arbitrary informal document fragments. That way formal knowledge management services become available for informal documents, accessible via an IDE for authors and via generated *active* documents for readers, while remaining fully compatible with existing authoring workflows and publishing systems.

Additionally, an extensible library of reusable document fragments is being developed, that serve as reference targets for global disambiguation, intermediaries for content exchange between systems and other services.

Every component of the system is designed modularly and extensibly, and thus lay the groundwork for a potential full integration of interactive theorem proving systems into established informal document authoring workflows.

The general STEX workflow combines functionalities provided by several pieces of software:

- $\bullet\,$ The STEX package to use semantic annotations in IATEX documents,
- RusTeX to convert tex sources to (semantically enriched) xhtml,
- The MMT software, that extracts semantic information from the thus generated xhtml and provides semantically informed added value services.

Chapter 2

Quickstart

2.1 Setup

2.1.1 The STEX IDE

TODO: VSCode Plugin

2.1.2 Manual Setup

Foregoing on the STFX IDE, we will need several pieces of software; namely:

- The STEX-Package available here.
 STEX is also available on CTAN and in TeXLive.
- To make sure that STEX too knows where to find its archives, we need to set a global system variable MATHHUB, that points to your local MathHub-directory (see section 3.2).
- The Mmt System available here¹. We recommend following the setup routine documented here.
 - Following the setup routine (Step 3) will entail designating a MathHub-directory on your local file system, where the MMT system will look for STEX/MMT content archives.
- STEX Archives If we only care about IATEX and generating pdfs, we do not technically need MMT at all; however, we still need the MATHHUB system variable to be set. Furthermore, MMT can make downloading content archives we might want to use significantly easier, since it makes sure that all dependencies of (often highly interrelated) STEX archives are cloned as well.
 - Once set up, we can run mmt in a shell and download an archive along with all of its dependencies like this: lmh install <name-of-repository>, or a whole group of archives; for example, lmh install smglom will download all smglom archives.
- RusTeX The Mmt system will also set up RusTeX for you, which is used to generate (semantically annotated) xhtml from tex sources. In lieu of using Mmt, you can also download and use RusTeX directly here.

EdN:1

¹EdNote: For now, we require the sTeX-branch, requiring manually compiling the MMT sources

2.2 A First STEX Document

Having set everything up, we can write a first STEX document. As an example, we will use the smglom/calculus and smglom/arithmetics archives, which should be present in the designated MathHub-folder, and write a small fragment defining the *geometric series*:

TODO: use some sTeX-archive instead of smglom, use a convergence-notion that includes the limit, mark-up the theorem properly

```
1 \documentclass{article}
  \usepackage{stex,xcolor,stexthm}
4 \begin{document}
 5 \begin{smodule}{GeometricSeries}
       \importmodule[smglom/calculus]{series}
      \importmodule[smglom/arithmetics]{realarith}
 7
 8
9
      \symdef{geometricSeries}[name=geometric-series]{\comp{S}}
10
      \begin{sdefinition} [for=geometricSeries]
11
          The \definame{geometricSeries} is the \symname{?series}
13
          \[\defeq{\geometricSeries}{\definiens{
14
              \displaystyle \inf \{ \sup \{ svar\{n\} \} \} \} 
15
                  \realdivide[frac]{1}{
                      \realpower{2}{\svar{n}}
17
              }}
18
          }}.\]
19
      \end{sdefinition}
20
21
      \begin{sassertion} [name=geometricSeriesConverges, type=theorem]
      The \symname{geometricSeries} \symname{converges} towards $1$.
      \end{sassertion}
24 \end{smodule}
25 \end{document}
```

Compiling this document with pdflatex should yield the output

Definition 0.1. The **geometric series** is the series

$$S := \sum_{n=1}^{\infty} \frac{1}{2^n}.$$

Theorem 0.2. The geometric series converges towards 1.

Feel free to move your cursor over the various highlighted parts of the document – depending on your pdf viewer, this should yield some interesting (but possibly for now cryptic) information.

Remark 2.2.1:

Note that all of the highlighting, tooltips, coloring and the environment headers come from stexthm – by default, the amount of additional packages loaded is kept to a minimum and all the presentations can be customized, see <a href="https://chapter.com/chap

Let's investigate this document in detail now:

```
\begin{smodule}{GeometricSeries}
...
\end{smodule}
```

smodule

First, we open a new *module* called GeometricSeries. This module is assigned a *globally unique* identifier (URI), which (depending on your pdf viewer) should pop up in a tooltip if you hover over the word **geometric series**.

```
\importmodule[smglom/calculus]{series}
\importmodule[smglom/arithmetics]{realarith}
```

\importmodule

Next, we *import* two modules — series in the smglom/calculus-archive, and realarith in the smglom/arithmetics-archive. If we investigate these archives, we find the files series.en.tex and realarith.en.tex (respectively) in their respective source-folders, which contain the statements \begin{smodule}{smodule}{series} and \begin{smodule}{frealarith} (respectively).

The \importmodule-statements make all STEX symbols and associated semantic macros (e.g. \infinitesum, \realdivide, \realpower) in the desired module available. Additionally, they "export" these symbols to all further modules which include the current module – i.e. if in some future module we would put \importmodule {GeometricSeries}, we would also have \infinitesum etc. at our disposal.

\usemodule

If we only want to *use* the content of some module Foo, e.g. in remarks or examples, but none of the symbols in our current module actually *depend* on the content of Foo, we can use \usemodule instead – like \importmodule, this will make the module content available, but will *not* export it to other modules.

```
\symdef{GeometricSeries}[name=geometric-series]{\comp{S}}
```

\symdef

Next, we introduce a new symbol with name geometric-series and assign it the semantic macro \geometricSeries. \symdef also immediately assigns this symbol a notation, namely S.

\comp

The macro \comp marks the S in the notation as a notational component, as opposed to e.g. arguments to \geometricSeries. It is the notational components that get highlighted and associated with the corresponding symbol (i.e. in this case geometricSeries). Since \geometricSeries takes no arguments, we can wrap the whole notation in a \comp.

```
\begin{sdefinition} [for=geometricSeries]
...
\end{sdefinition}
\begin{sassertion} [name=geometricSeriesConverges, type=theorem]
...
\end{sassertion}
```

What follows are two STEX-statements (e.g. definitions, theorems, examples, proofs, ...). These are semantically marked-up variants of the usual environments, which take additional optional arguments (e.g. for=, type=, name=). Since many LATEX templates predefine environments like definition or theorem with different syntax, we use sdefinition, sassertion, sexample etc. instead. You can customize these environments to e.g. simply wrap around some predefined theorem-environment. That way, we can still use sassertion to provide semantic information, while being fully compatible with (and using the document presentation of) predefined environments.

In our case, the stexthm-package patches e.g. \begin{sassertion} [type=theorem] to use a theorem-environment defined (as usual) using amsthm.

The \definame{geometricSeries} is the \symname{?series}

\symname

The \symname-command prints the name of a symbol, highlights it (based on customizable settings) and associates the text printed with the corresponding symbol. If you hover over the word series in the pdf output, you should see a tooltip showing the full URI of the symbol used.

\symref

The \symname-command is a special case of the more general \symref-command, which allows customizing the precise text associated with a symbol.

\definame \definiendum

The sdefinition-environment provides two additional macros, \definame and \definiendum which behave similar to \symname and \symref, but explicitly mark the symbols as being defined in this environment, to allow for special highlighting.

```
\[\defeq{\geometricSeries}{\definiens{
   \infinitesum{\svar{n}}{1}{
      \realdivide[frac]{1}{
      \realpower{2}{\svar{n}}
   }}
}}.\]
```

The next snippet – set in a math environment – uses several semantic macros imported from (or recursively via) series and realarithmetics, such as \defeq , \infinitesum , etc. In math mode, using a semantic macro inserts its (default) definition. A semantic macro can have several notations – in that case, we can explicitly choose a specific notation by providing its identifier as an optional argument; e.g. $\realdivide[frac]{a}{b}$ will use the explicit notation named $\frac{frac}{frac}$ of the semantic macro \realdivide , which yields $\frac{a}{b}$ instead of $\frac{a}{b}$.

\svar

The \sqrt{n} command marks up the n as a variable with name n and notation n.

\definiens

The **sdefinition**-environment additionally provides the \definiens-command, which allows for explicitly marking up its argument as the *definiens* of the symbol currently being defined.

2.2.1 OMDoc/xhtml Conversion

So, if we run pdflatex on our document, then STEX yields pretty colors and tooltips¹. But STEX becomes a lot more powerful if we additionally convert our document to xhtml.

TODO VSCode Plugin

Using $R_{US}T_{E}X$, we can convert the document to xhtml using the command rustex -i /path/to/file.tex -o /path/to/outfile.xhtml. Investigating the resulting file, we notice additional semantic information resulting from our usage of semantic macros, \symmetric Elow is the (abbreviated) snippet inside our \definiens block:

```
<mrow resource="" property="stex:definiens">
<mrow resource="...?series?infinitesum" property="stex:OMBIND">
  <munderover displaystyle="true">
   <mo resource="...?series?infinitesum" property="stex:comp">\Sigma</mo>
    <mrow resource="1" property="stex:arg">
     <mi resource="var://n" property="stex:OMV">n</mi>
    <mo resource="...?series?infinitesum" property="stex:comp">=</mo>
    <mi resource="2" property="stex:arg">1</mi>
   </mrow>
   <mi resource="...?series?infinitesum" property="stex:comp"></mi>
  </munderover>
  <mrow resource="3" property="stex:arg">
<mfrac resource="...?realarith?division#frac#" property="stex:OMA">
    <mi resource="1" property="stex:arg">1</mi>
<mrow resource="2" property="stex:arg">
<msup resource="...realarith?exponentiation" property="stex:OMA">
      <mi resource="1" property="stex:arg">2</mi>
<mrow resource="2" property="stex:arg"></mi>
       <mi resource="var://n" property="stex:OMV">n</mi>
      </mrow>
     </msup>
    </mrow>
   </mfrac>
  </mrow>
 </mrow>
</mrow>
```

...containing all the semantic information. The MMT system can extract from this the following OPENMATH snippet:

```
<OMBIND>
<OMID name="...?series?infinitesum"/>
<OMV name="n"/>
<OMLIT name="1"/>
<OMA>

<OMS name="...?realarith?division"/>
<OMLIT name="1"/>
<OMA>

<OMS name="...realarith?exponentiation"/>
<OMLIT name="2"/>
<OMLIT name="2"/>
<OMV name="n"/>
</OMA>
</OMA>
</OMBIND>
```

¹...and hyperlinks for symbols, and indices, and allows reusing document fragments modularly, and...

...giving us the full semantics of the snippet, allowing for a plurality of knowledge management services – in particular when serving the xhtml.

Remark 2.2.2:

Note that the html when opened in a browser will look slightly different than the pdf when it comes to highlighting semantic content – that is because naturally html allows for much more powerful features than pdf does. Consequently, the html is intended to be served by a system like MMT, which can pick up on the semantic information and offer much more powerful highlighting, linking and similar features, and being customizable by readers rather than being prescribed by an author.

Additionally, not all browsers (most notably Chrome) support MATHML natively, and might require additional external JavaScript libraries such as MathJax to render mathematical formulas properly.

Chapter 3

Creating STeX Content

We can use STEX by simply including the package with \usepackage{stex}, or - primarily for individual fragments to be included in other documents - by using the STEX document class with \documentclass{stex} which combines the standalone document class with the stex package.

Both the stex package and document class offer the following options:

lang $(\langle language \rangle *)$ Languages to load with the babel package.

mathhub ($\langle directory \rangle$) MathHub folder to search for repositories – this is not necessary if the MATHHUB system variable is set.

sms $(\langle boolean \rangle)$ use persisted mode (not yet implemented).

image $(\langle boolean \rangle)$ passed on to tikzinput.

debug $(\langle log\text{-}prefix\rangle *)$ Logs debugging information with the given prefixes to the terminal, or all if all is given. Largely irrelevant for the majority of users.

3.1 How Knowledge is Organized in STEX

STFX content is organized on multiple levels:

- STEX archives (see section 3.2) contain individual .tex-files.
- These may contain STFX modules, introduced via \begin{smodule}{ModuleName}.
- Modules contain STEX symbol declarations, introduced via \symdecl{symbolname}, \symdef{symbolname} and some other constructions. Most symbols have a notation that can be used via a semantic macro \symbolname generated by symbol declarations.
- STeX expressions finally are built up from usages of semantic macros.



- STEX archives are simultaneously MMT archives, and the same directory structure is consequently used.
- STEX modules correspond to OMDoc/MMT theories. \importmodules (and



similar constructions) induce MMT includes and other theory morphisms, thus giving rise to a theory graph in the OMDoc sense.

- Symbol declarations induce OMDoc/MMT constants, with optional (formal) type and definiens components.
- Finally, STEX expressions are converted to OMDoc/MMT terms, which use the syntax of OPENMATH.

3.2 STEX Archives

3.2.1 The Local MathHub-Directory

\userodule, \importmodule, \inputref etc. allow for including content modularly without having to specify absolute paths, which would differ between users and machines. Instead, STEX uses archives that determine the global namespaces for symbols and statements and make it possible for STEX to find content referenced via such URIs.

All STEX archives need to exist in the local MathHub-directory. STEX knows where this folder is via one of three means:

- 1. If the STEX package is loaded with the option mathhub=/path/to/mathhub, then STEX will consider /path/to/mathhub as the local MathHub-directory.
- 2. If the mathhub package option is *not* set, but the macro \mathhub exists when the STEX-package is loaded, then this macro is assumed to point to the local MathHub-directory; i.e. \def\mathhub{/path/to/mathhub}\usepackage{stex} will set the MathHub-directory as path/to/mathhub.
- 3. Otherwise, STEX will attempt to retrieve the system variable MATHHUB, assuming it will point to the local MathHub-directory. Since this variant needs setting up only once and is machine-specific (rather than defined in tex code), it is compatible with collaborating and sharing tex content, and hence recommended.

3.2.2 The Structure of ST_FX Archives

An STEX archive group/name needs to be stored in the directory /path/to/mathhub/group/name; e.g. assuming your local MathHub-directory is set as /user/foo/MathHub, then in order for the smglom/calculus-archive to be found by the STEX system, it needs to be in /user/foo/MathHub/smglom/calculus.

Each such archive needs two subdirectories:

- /source this is where all your tex files go.
- /META-INF a directory containing a single file MANIFEST.MF, the content of which we will consider shortly

An additional lib-directory is optional, and is where STEX will look for files included via \\libinput.

Additionally a *group* of archives <code>group/name</code> may have an additional archive <code>group/meta-inf</code>. If this <code>meta-inf-archive</code> has a <code>/lib-subdirectory</code>, it too will be searched by <code>\libinput</code> from all tex files in any archive in the <code>group/*-group</code>.

We recommend this additional directory structure in the source-folder of an STEX archive:

- /source/mod/ individual STEX modules, containing symbol declarations, notations, and \begin{sparagraph} [type=symdoc,for=...] environments for "encyclopedic" symbol documentations
- /source/def/ definitions
- /source/ex/ examples
- /source/thm/ theorems, lemmata and proofs; preferably proofs in separate files to allow for multiple proofs for the same statement
- /source/snip/ individual text snippets such as remarks, explanations etc.
- /source/frag/ individual document fragments, ideally only \inputrefing snippets, definitions, examples etc. in some desirable order
- /source/tikz/ tikz images, as individual .tex-files
- /source/pic/ image files.

3.2.3 MANIFEST.MF-Files

The MANIFEST.MF in the META-INF-directory consists of key-value-pairs, instructing STEX (and associated software) of various properties of an archive. For example, the MANIFEST.MF of the smglom/calculus-archive looks like this:

Many of these are in fact ignored by STEX, but some are important:

id: The name of the archive, including its group (e.g. smglom/calculus),

source-base or

ns: The namespace from which all symbol and module URIs in this repository are formed, see (TODO),

narration-base: The namespace from which all document URIs in this repository are formed, see (TODO),

url-base: The URL that is formed as a basis for external references, see (TODO),

dependencies: All archives that this archive depends on. STEX ignores this field, but MMT can pick up on them to resolve dependencies, e.g. for lmh install.

3.2.4 Using Files in STEX Archives Directly

Several macros provided by ST_EX allow for directly including files in repositories. These are:

\mhinput

\mhinput[Some/Archive] {some/file} directly inputs the file some/file in the source-folder of Some/Archive.

\inputref

\inputref[Some/Archive]{some/file} behaves like \mhinput, but wraps the input in a \begingroup ... \endgroup. When converting to xhtml, the file is not input at all, and instead an html-annotation is inserted that references the file.

In the majority of cases \inputref is likely to be preferred over \mhinput.

\ifinput

Both \mhinput and \inputref set \iffinput to "true" during input. This allows for selectively including e.g. bibliographies only if the current file is not being currently included in a larger document.

\addmhbibresource

\addmhbibresource [Some/Archive] {some/file} searches for a file like \mhinput does, but calls \addbibresource to the result and looks for the file in the archive root directory directly, rather than the source directory.

\libinput

\libinput{some/file} searches for a file some/file in

- the lib-directory of the current archive, and
- the lib-directory of a meta-inf-archive in (any of) the archive groups containing the current archive

and include all found files in reverse order; e.g. \libinput{preamble} in a .tex-file in smglom/calculus will first input .../smglom/meta-inf/lib/preamble.tex and then ../smglom/calculus/lib/preamble.tex.

Will throw an error if *no* candidate for some/file is found.

\libusepackage

\libusepackage[package-options]{some/file} searches for a file some/file.sty in the same way that \libinput does, but will call \usepackage[package-options]{path/to/some/file} instead of \input.

Will throw an error if not exactly one candidate for some/file is found.

Remark 3.2.1:

Then the preamble.tex files can take care of loading the generally required packages, setting presentation customizations etc. (per archive or archive group or both), and postamble.tex can e.g. print the bibliography, index etc.

3.3 Module, Symbol and Notation Declarations

3.3.1 The smodule-Environment

smodule A new module is declared using the basic syntax

```
\begin{smodule} [options] {ModuleName}...\end{smodule}.
```

A module is required to declare any new formal content such as symbols or notations (but not variables, which may be introduced anywhere).

The smodule-environment takes several optional arguments, all of which are optional:

```
title (\langle token \ list \rangle) to display in customizations.
```

type $(\langle string \rangle *)$ for use in customizations.

deprecate $(\langle module \rangle)$ if set, will throw a warning when loaded, urging to use $\langle module \rangle$ instead.

id $(\langle string \rangle)$ for cross-referencing.

ns $(\langle \mathit{URI} \rangle)$ the namespace to use. Should not be used, unless you know precisely what you're doing. If not explicitly set, is computed using $\text{stex_modules_current_namespace:}$.

lang (\language \rangle) if not set, computed from the current file name (e.g. foo.en.tex).

sig (\language\rangle) if the current file is a translation of a file with the same base name but a
different language suffix, setting sig=<lang> will preload the module from that language
file. This helps ensuring that the (formal) content of both modules is (almost) identical
across languages and avoids duplication.

creators ($\langle string \rangle *$) names of the creators.

contributors ($\langle string \rangle *$) names of contributors.

srccite $(\langle string \rangle)$ a source citation for the content of this module.

```
 \begin{array}{l} \overset{\longleftarrow}{\longrightarrow} \text{ An SIEX module corresponds to an MMT/OMDoc } \textit{theory.} & \text{As such it} \\ -\mathbb{M} & \text{gets assigned a module URI } \textit{(universal resource identifier)} & \text{of the form} \\ & \overset{\longleftarrow}{\longrightarrow} \text{ ``namespace'} & \text{``module-name'}. \\ \end{array}
```

By default, opening a module will produce no output whatsoever, e.g.:

Example 1

Input:

```
1 \begin{smodule}[title={This is Some Module}]{SomeModule}
2  Hello World
3 \end{smodule}
```

Output:

Hello World

\stexpatchmodule

We can customize this behavior either for all modules or only for modules with a specific type using the command \stexpatchmodule[optional-type]{begin-code}{end-code}. Some optional parameters are then available in \smodule*-macros, specifically \smoduletitle, \smoduletype and \smoduleid.

For example:

Example 2

```
Input:
```

```
1 \stexpatchmodule[display]
2 {\textbf{Module (\smoduletitle)}\par}
3 {\par\noindent\textbf{End of Module (\smoduletitle)}}
4
5 \begin{smodule}[type=display,title={Some New Module}]{SomeModule2}
6 Hello World
7 \end{smodule}
```

Output:

```
Module (Some New Module)

Hello World

End of Module (Some New Module)
```

3.3.2 Declaring New Symbols and Notations

Inside an smodule environment, we can declare new STEX symbols.

\symdecl

The most basic command for doing so is using \symdecl{symbolname}. This introduces a new symbol with name symbolname, arity 0 and semantic macro \symbolname.

The starred variant \symdecl*{symbolname} will declare a symbol, but not introduce a semantic macro. If we don't want to supply a notation (for example to introduce concepts like "abelian", which is not something that has a notation), the starred variant is likely to be what we want.

Without a semantic macro or a notation, the only meaningful way to reference a symbol is via \symref,\symname etc.

```
Example 3
Input:

1 \symdecl*{foo}
2 Given a \symname{foo}, we can...

Output:
```

Given a foo, we can...

Obviously, most semantic macros should take actual *arguments*, implying that the symbol we introduce is an *operator* or *function*. We can let \symdecl know the *arity* (i.e. number of arguments) of a symbol like this:

```
Example 4
```

```
Input:
```

```
1 \symdecl{binarysymbol}[args=2] 2 \symref{binarysymbol}{this} is a symbol taking two arguments.
```

Output:

this is a symbol taking two arguments.

.

\notation

In that case, we probably want to supply a notation as well, in which case we can finally actually use the semantic macro in math mode. We can do so using the \notation command, like this:

Example 5 Input: 1 \notation{binarysymbol}{\text{First: }#1\text{; Second: }#2} 2 \$\binarysymbol{a}{b}\$ Output: First: a; Second: b . -M Applications of semantic macros, such as \binarysymbol{a}{b} are translated to -M MMT/OMDoc as OMA-terms with head <OMS name="...?binarysymbol"/>.

\comp

Unfortunately, we have no highlighting whatsoever now. That is because we need to tell STEX explicitly which parts of the notation are *notation components* which *should* be highlighted. We can do so with the \comp command.

→T→ Semantic macros with no arguments correspond to OMS directly.

We can introduce a new notation highlight for \binarysymbol that fixes this flaw, which we can subsequently use with \binarysymbol[highlight]:

```
Example 6
Input:
```

```
1 \notation{binarysymbol}[highlight]
2 {\comp{\text{First: }}#1\comp{\text{; Second: }}#2}
3 $\binarysymbol[highlight]{a}{b}$
```

Output:

```
First: a; Second: b
```



Ideally, \comp would not be necessary: Everything in a notation that is not an argument should be a notation component. Unfortunately, it is computationally expensive to determine where an argument begins and ends, and the argument markers #n may themselves be nested in other macro applications or TeX groups, making it ultimately almost impossible to determine them automatically while also remaining compatible with arbitrary highlighting customizations (such as tooltips, hyperlinks, colors) that users might employ, and that are ultimately invoked by \comp.

Note that it is required that

- 1. the argument markers #n never occur inside a \comp, and
- 2. no semantic arguments may ever occur inside a notation.

Both criteria are not just required for technical reasons, but conceptionally meaningful:

The underlying principle is that the arguments to a semantic macro represent arguments to the mathematical operation represented by a symbol. For example, a semantic macro $\addition\{a\}\{b\}$ taking two arguments would represent the actual addition of (mathematical objects) a and b. It should therefore be impossible for a or b to be part of a notation component of \addition .



Similarly, a semantic macro can not conceptually be part of the notation of \addition, since a semantic macro represents a distinct mathematical concept with its own semantics, whereas notations are syntactic representations of the very symbol to which the notation belongs.

If you want an argument to a semantic macro to be a purely syntactic parameter, then you are likely somewhat confused with respect to the distinction between the precise syntax and semantics of the symbol you are trying to declare (which happens quite often even to experienced STEX users), and might want to give those another thought - quite likely, the macro you aim to implement does not actually represent a semantically maningful mathematical concept, and you will want to use \def and similar native LATEX macro definitions rather than semantic macros.

\symdef

In the vast majority of cases where a symbol declaration should come with a semantic macro, we will want to supply a notation immediately. For that reason, the \symdef command combines the functionality of both \symdecl and \notation with the optional arguments of both:

Example 7

Input:

```
1 \symdef{newbinarysymbol}[hl,args=2]
2 {\comp{\text{1.: }}#1\comp{\text{; 2.: }}#2}
3 $\newbinarysymbol{a}{b}$
```

Output:

```
1.: a; 2.: b
```

We just declared a new symbol newbinarysymbol with args=2 and immediately provided it with a notation with identifier hl. Since hl is the *first* (and so far, only) notation supplied for newbinarysymbol, using \newbinarysymbol without optional argument defaults to this notation.

\setnotation

The first notation provided will stay the default notation unless explicitly changed — this is enabled by the \setnotation command: \setnotation{symbolname} {notation-id} sets the default notation of \symbolname to notation-id, i.e. henceforth, \symbolname behaves like \symbolname[notation-id] from now on.

Often, a default notation is set right after the corresponding notation is introduced – the starred version \notation* for that reason introduces a new notation and immediately sets it to be the new default notation. So expressed differently, the first \notation for a symbol behaves exactly like \notation*, and \notation*{foo}[bar]{...} behaves exactly like \notation{foo}{bar}.

Operator Notations

Once we have a semantic macro with arguments, such as \newbinarysymbol, the semantic macro represents the application of the symbol to a list of arguments. What if we want to refer to the operator itself, though?

We can do so by supplying the \notation (or \symdef) with an operator notation, indicated with the optional argument op=. We can then invoke the operator notation using \symbolname! [notation-identifier]. Since operator notations never take arguments, we do not need to use \comp in it, the whole notation is wrapped in a \comp automatically:

Example 8 Input: 1 \notation{newbinarysymbol}[ab, 2 op={\text{a:}\cdot\text{; b:}\cdot}] 3 {\comp{\text{a:}}#1\comp{\text{; b:}}#2} 4 \symname{newbinarysymbol} is also occasionally written 5 \$\newbinarysymbol![ab]\$ Output: newbinarysymbol is also occasionally written a: ·; b:

3.3.3 Argument Types

The notations so far used *simple* arguments which we call i-type arguments. Declaring a new symbol with \symdecl{foo}[args=3] is equivalent to writing \symdecl{foo}[args=iii], indicating that the semantic macro takes three i-type arguments. However, there are three more argument types which we will investigate now, namely b-type, a-type and B-type arguments.

b-Type Arguments

A b-type argument represents a variable that is bound by the symbol in its application, making the symbol a binding operator. Typical examples of binding operators are e.g. sums \sum , products \prod , integrals \int , quantifiers like \forall and \exists , that λ -operator, etc.

```
\buildrel M \buildrel b-type arguments behave exactly like i-type arguments within TeX, but applications of binding operators, i.e. symbols with b-type arguments, are translated to \buildrel T \buildrel DMBIND-terms in OMDOC/MMT, rather than OMA.
```

Fo example, we can implement a summation operator binding an index variable and taking lower and upper index bounds and the expression to sum over like this:

```
Example 9
```

Input:

```
1 \symdef{summation}[args=biii]
2 {\mathop{\comp{\sum}}_{#1\comp{=}#2}^{#3}#4}
3 $\summation{\svar{x}}{1}{\svar{n}}{\svar{x}}^2$
```

Output:

```
\sum_{x=1}^{n} x^2
```

where the variable x is now bound by the \summation-symbol in the expression.

a-Type Arguments

a-type arguments represent a *flexary argument sequence*, i.e. a sequence of arguments of arbitrary length. Formally, operators that take arbitrarily many arguments don't "exist", but in informal mathematics, they are ubiquitous. a-type arguments allow us to write e.g. \addition{a,b,c,d,e} rather than having to write something like \addition{a}{\addition{b}{\addition{b}}}!

\notation (and consequently \symdef, too) take one additional argument for each a-type argument that indicates how to "accumulate" a comma-separated sequence of arguments. This is best demonstrated on an example.

Let's say we want an operator representing quantification over an ascending chain of elements in some set, i.e. $\ascendingchain\{S\}\{a,b,c,d,e\}\{t\}$ should yield $\forall a <_S b <_S c <_S d <_S e.t$. The "base"-notation for this operator is simply

 ${\comp{\forall} #2\comp{.},}#3}$, where #2 represents the full notation fragment *accumulated* from {a,b,c,d,e}.

The additional argument to \notation (or \symdef) takes the same arguments as the base notation and two additional arguments ##1 and ##2 representing successive pairs in the a-type argument, and accumulates them into #2, i.e. to produce $a <_S b <_S c <_S d <_S e$, we do {##1 \comp{<}_{#1} ##2}:

Example 10

Input:

```
1 \symdef{ascendingchain}[args=iai]
2 {\comp{\forall} #2\comp{.\,}#3}
3 {##1 \comp{<}_{#1} ##2}
4
5 Tadaa: $\ascendingchain{S}{a,b,c,d,e}{t}$</pre>
```

Output:

```
Tadaa: \forall a <_S b <_S c <_S d <_S e. t
```

If this seems overkill, keep in mind that you will rarely need the single-hash arguments #1,#2 etc. in the a-notation-argument. For a much more representative and simpler example, we can introduce flexary addition via:

Example 11

bind a single variable etc.

```
Input:

1 \symdef{addition}[args=a]{#1}{##1 \comp{+} ##2}
2 3 Tadaa: $\addition{a,b,c,d,e}$

Output:

Tadaa: a+b+c+d+e
```

The assoc-key We mentioned earlier that "formally", flexary arguments don't really "exist". Indeed, formally, addition is usually defined as a binary operation, quantifiers

Consequently, we can tell STEX (or, rather, MMT/OMDOC) how to "resolve" flexary arguments by providing \symdecl or \symdef with an optional assoc-argument, as in \symdecl{addition}[args=a,assoc=bin]. The possible values for the assoc-key are:

bin: A binary, assoiative argument, e.g. as in \addition

binl: A binary, left-associative argument, e.g. $a^{b^{c^d}}$, which stands for $((a^b)^c)^d$

binr: A binary, right-associative argument, e.g. as in $A \to B \to C \to D$, which stands for $A \to (B \to (C \to D))$

pre: Successively prefixed, e.g. as in $\forall x, y, z. P$, which stands for $\forall x. \forall y. \forall z. P$

conj: Conjunctive, e.g. as in a=b=c=d or $a,b,c,d\in A$, which stand for $a=d\wedge b=d\wedge c=d$ and $a\in A\wedge b\in A\wedge c\in A\wedge d\in A$, respectively

pwconj: Pairwise conjunctive, e.g. as in $a \neq b \neq c \neq d$, which stands for $a \neq b \land a \neq c \land a \neq d \land b \neq c \land b \neq d \land c \neq d$

B-Type Arguments

Finally, B-type arguments simply combine the functionality of both a and b - i.e. they represent an arbitrarily long sequence of variables to be bound, e.g. for implementing quantifiers:

Example 12

```
Input:

1 \symdef{quantforall}[args=Bi]
2 {\comp{\forall}#1\comp{.}#2}
3 {##1\comp,##2}
4
5 $\quantforall{\svar{x},\svar{y},\svar{z}}{P}$
```

Output:

```
\forall x,y,z.P
```

3.3.4 Type and Definiens Components

\symdecl and \symdef take two more optional arguments. TEX largely ignores them (except for special situations we will talk about later), but MMT can pick up on them for additional services. These are the type and def keys, which expect expressions in math-mode (ideally using semantic macros, of course!)

```
The type and def keys correspond to the type and definiens components of CM OMDOC/MMT constants.

M Correspondingly, the name "type" should be taken with a grain of salt, since OMDOC/MMT—being foundation-independent—does not a priori implement a fixed typing system.
```

The type-key allows us to provide additional information (given the necessary STEX symbols), e.g. for addition on natural numbers:

Example 13

Input:

```
1 \symdef{Nat}[type=\set]{\comp{\mathbb N}}
2 \symdef{addition}[
3     type=\funtype{\Nat,\Nat}{\Nat},
4     op=+,
5     args=a
6 ]{#1}{##1 \comp+ ##2}
7
8 \symname{addition} is an operation $\funtype{\Nat,\Nat}{\Nat}$
```

Output:

```
addition is an operation \mathbb{N}{\times}\mathbb{N}{\to}\mathbb{N}
```

The def-key allows for declaring symbols as abbreviations:

Example 14

Input:

```
\symdef{successor}[
     type=\funtype{\Nat}{\Nat},
     def = \{ x } { \addition { \xxx{x}, 1} },
     op=\mathtt{succ},
args=1
\frac{4}{5}
6 ]{\comp{\mathtt{succ(}#1\comp{)}}}
8 The \symname{successor} operation \int \int \int {\mathbb R} {\mathbb R}^{n} dt
9 is defined as \int x{x}}{\addition{xxx{x},1}}
```

Output:

```
The successor operation \mathbb{N} \rightarrow \mathbb{N} is defined as x \mapsto x+1
```

3.3.5 Precedences and Automated Bracketing

Having done \addition, the obvious next thing to implement is \multiplication. This is in theory straight-forward:

Example 15

Input:

```
\symdef{multiplication}[
   type=\funtype{\Nat,\Nat}{\Nat},
3
   op=\cdot,
   args=a
5 ]{#1}{##1 \comp\cdot ##2}
```

Output:

```
multiplication is an operation \mathbb{N} \times \mathbb{N} \to \mathbb{N}
```

However, if we combine \addition and \multiplication, we notice a problem:

Example 16

Input:

```
1 \alpha, \
```

Output:

```
a+b\cdot c+d\cdot e
```

'We all know that \cdot binds stronger than +, so the output $a+b\cdot c+d\cdot e$ does not actually reflect the term we wrote. We can of course insert parentheses manually

but we can also do better by supplying precedences and have ST_EX insert parentheses automatically.

For that purpose, \notation (and hence \symdef) take an optional argument prec=<opprec>;<argprec1>x...x<argprec n>.

We will investigate the precise meaning of <opprec> and the <argprec>s shortly – in the vast majority of cases, it is prefectly sufficient to think of prec= taking a single number and having that be *the* precedence of the notation, where lower precedences (somewhat counterintuitively) bind stronger than higher precedences. So fixing our notations for \addition and \multiplication, we get:

Example 18

```
Input:

1 \notation{multiplication}[
2    op=\cdot,
3    prec=50
4 ]{#1}{##1 \comp\cdot ##2}
5 \notation{addition}[
6    op=+,
7    prec=100
8 ]{#1}{##1 \comp+ ##2}
9
10 $\addition{a, \multiplication{b, \addition{c, \multiplication{d,e}}}}$
```

Output:

```
a + b \cdot (c + d \cdot e)
```

'Note that the precise numbers used for precedences are pretty arbitrary - what matters is which precedences are higher than which other precedences when used in conjunction.

\infprec \neginfprec

It is occasionally useful to have "infinitely" high or low precedences to enforce or forbid automated bracketing entirely – for those purposes, \infprec and \neginfprec exist (which are implemented as the maximal and minimal integer values accordingly).



More precisely, each notation takes

1. One operator precedence and

2. one argument precedence for each argument.

By default, all precedences are 0, unless the symbol takes no argument, in which case the operator precedence is \neginfprec (negative infinity). If we only provide a single number, this is taken as both the operator precedence and all argument precedences.

STEX decides whether to insert parentheses by comparing operator precedences to a downward precedence p_d with initial value \infprec. When encountering a semantic macro, STEX takes the operator precedence p_{op} of the notation used and checks whether $p_{op} > p_d$. If so, STEX insert parentheses.

When STEX steps into an argument of a semantic macro, it sets p_d to the respective argument precedence of the notation used.

In the example above:

- 1. STEX starts out with $p_d = \$
- STEX encounters \addition with p_{op} = 100. Since 100 ≯\infprec, it inserts no parentheses.



- 4. Next, STEX encounters \multiplication{b,...}, whose notation has $p_{op} = 50$.
- 5. We compare to the current downward precedence p_d set by \addition, arriving at $p_{op} = 50 > 100 = p_d$, so SIEX again inserts no parentheses.
- 6. Since the notation of \multiplication has no explicitly set argument precedences, STEX uses the operator precedence for all arguments of \multiplication, hence sets $p_d = p_{op} = 50$ and recurses.
- 7. Next, gT_EX encounters the inner $\addition\{c,...\}$ whose notation has $p_{op} = 100$.
- 8. We compare to the current downward precedence p_d set by \multiplication, arriving at $p_{op} = 100 > 50 = p_d$ which finally prompts STEX to insert parentheses, and we proceed as before.

3.3.6 Variables

All symbol and notation declarations require a module with which they are associated, hence the commands \symdecl, \notation, \symdef etc. are disabled outside of smodule-environments.

Variables are different – variables are allowed everywhere, are not exported when the current module (if one exists) is imported (via \importmodule or \usemodule) and (also unlike symbol declarations) "disappear" at the end of the current TeX group.

\svar

So far, we have always used variables using n , which marks-up n as a variable with name n. More generally, code marks-up the arbitrary code as representing a variable with name foo.

Of course, this makes it difficult to reuse variables, or introduce "functional" variables with arities > 0, or provide them with a type or definiens.

\vardef

For that, we can use the \vardef command. Its syntax is largely the same as that of \symdef, but unlike symbols, variables have only one notation (TODO: so far?), hence there is only \vardef and no \vardecl.

Example 19

```
Input:
```

```
1  \vardef{varf}[
2    name=f,
3    type=\funtype{\Nat}{\Nat},
4    op=f,
5    args=1,
6    prec=0;\neginfprec
7  ]{\comp{f}#1}
8  \vardef{varn}[name=n,type=\Nat]{\comp{n}}
9  \vardef{varx}[name=x,type=\Nat]{\comp{x}}
10
11  Given a function $\varf!:\funtype{\Nat}{\Nat}$,
12  by $\addition{\varf!,\varn}$ we mean the function
13  $\fun{\varx}{\varf{\addition}\varx,\varn}}$
```

Output:

```
Given a function f: \mathbb{N} \to \mathbb{N}, by f+n we mean the function x \mapsto f(x+n)
```

'(of course, "lifting" addition in the way described in the previous example is an operation that deserves its own symbol rather than abusing \addition, but... well.)

TODO: bind=forall/exists

3.3.7 Variable Sequences

Variable sequences occur quite frequently in informal mathematics, hence they deserve special support. Variable sequences behave like variables in that they disappear at the end of the current TEX group and are not exported from modules, but their declaration is quite different.

\varseq

A variable sequence is introduced via the command \varseq, which takes the usual optional arguments name and type. It then takes a starting index, an end index and a notation for the individual elements of the sequence parametric in an index.

This is best shown by example:

Example 20

Input:

```
1 \vardef{varn}[name=n,type=\Nat]{\comp{n}}
2 \varseq{seqa}[name=a,type=\Nat]{1}{\varn}{\comp{a}_{#1}}
3
4 The $i$th index of $\seqa!$ is $\seqa{i}$.
```

Output:

```
The ith index of a_1, \ldots, a_n is a_i.
```

.

Note that the syntax \searrow now automatically generates a presentation based on the starting and ending index.

TODO: more notations for invoking sequences.

Notably, variable sequences are nicely compatible with ${\tt a}\textsc{-type}$ arguments, so we can do the following:

Example 21

```
Input:
```

```
1 \alpha
```

Output:

```
a_1 + \ldots + a_n
```

.

Sequences can be multidimensional using the args-key, in which case the notation's arity increases and starting and ending indices have to be provided as a comma-separated list:

Example 22

Input:

```
1 \vardef{varm}[name=m,type=\Nat]{\comp{m}}
2 \varseq{seqa}[
3     name=a,
4     args=2,
5     type=\Nat,
6 ]{1,1}{\varm,\varm}{\comp{a}_{#1}^{#2}}
7
8 $\seqa!$ and $\addition{\seqa}$
```

Output:

```
a_1^1, \dots, a_n^m and a_1^1 + \dots + a_n^m
```

We can also explicitly provide a "middle" segment to be used, like such:

Example 23

```
Input:
```

```
1 \varseq{seqa}[
2    name=a,
3    type=\Nat,
4    args=2,
5    mid={\comp{a}_{\varn}^1,\comp{a}_1^2,\ellipses,\comp{a}_{1}^{\varm}}}
6 ]{1,1}{\varn,\varm}{\comp{a}_{\text{#1}}^{\text{#2}}}
7    8 $\seqa!$ and $\addition{\seqa}$
```

Output:

```
a_1^1, \dots, a_n^1, a_1^2, \dots, a_1^m, \dots, a_n^m and a_1^1 + \dots + a_n^1 + a_1^2 + \dots + a_n^m + \dots + a_n^m
```

3.4 Module Inheritance and Structures

3.4.1 Multilinguality and Translations

If we load the STEX document class or package with the option lang=<lang>, STEX will load the appropriate babel language for you – e.g. lang=de will load the babel language ngerman. Additionally, it makes STEX aware of the current document being set in (in this example) german. This matters for reasons other than mere babel-purposes, though:

Every module is assigned a language. If no STEX package option is set that allows for inferring a language, STEX will check whether the current file name ends in e.g. .en.tex (or .de.tex or .fr.tex, or...) and set the language accordingly. Alternatively, a language can be explicitly assigned via \begin{smodule}[lang=<language>]{Foo}.

```
Technically, each smodule-environment induces two OMDoc/MMT theories: \begin{smodule}[lang=<lang>]{Foo} generates a theory some/namespace?Foo that only contains the "formal" part of the module – i.e. exactly the content—M→ that is exported when using \importmodule.

TAMAGINATION Additionally, MMT generates a language theory some/namespace/Foo?<lang> that includes some/namespace?Foo and contains all the other document content – variable declarations, includes for each \usenbodule, etc.
```

Notably, the language suffix in a filename is ignored for \usemodule, \importmodule and in generating/computing URIs for modules. This however allows for providing translations for modules between languages without needing to duplicate content:

If a module Foo exists in e.g. english in a file Foo.en.tex, we can provide a file Foo.de.tex right next to it, and write \begin{smodule}[sig=en]{Foo}. The sig-key then signifies, that the "signature" of the module is contained in the english version of the module, which is immediately imported from there, just like \importmodule would.

Additionally to translating the informal content of a module file to different languages, it also allows for customizing notations between languages. For example, the least common multiple of two numbers is often denoted as $\mathtt{lcm}(a,b)$ in english, but is called kleinstes gemeinsames Vielfaches in german and consequently denoted as $\mathtt{kgV}(a,b)$ there.

We can therefore imagine a german version of an lcm-module looking something like this:

```
1 \begin{smodule}[sig=en]{lcm}
2 \notation*{lcm}[de]{\comp{\mathtt{kgV}}(#1,#2)}
3
4 Das \symref{lcm}{kleinste gemeinsame Vielfache}
5 $\lcm{a,b}$ von zwei Zahlen $a,b$ ist...
6 \end{smodule}
```

If we now do \importmodule{lcm} (or \usemodule{lcm}) within a german document, it will also load the content of the german translation, including the de-notation for \lcm.

3.4.2 Simple Inheritance and Namespaces

\importmodule \usemodule

\importmodule[Some/Archive]{path?ModuleName} is only allowed within an smodule-environment and makes the symbols declared therein available. Additionally the content of ModuleName will be exported if the current module is imported somewhere else via \importmodule.

\usemodule behaves the same way, but without exporting the content of the used module.

It is worth going into some detail how exactly \importmodule and \usemodule resolve their arguments to find the desired module – which is closely related to the namespace generated for a module, that is used to generate its URI.

Ideally, STEX would use arbitrary URIs for modules, with no forced relationships between the logical namespace of a module and the physical location of the file declaring the module – like MMT does things.

Unfortunately, TEX only provides very restricted access to the file system, so we are forced to generate namespaces systematically in such a way that they reflect the physical location of the associated files, so that STEX can resolve them accordingly. Largely, users need not concern themselves with namespaces at all, but for completenesses sake, we describe how they are constructed:



- If \begin{smodule}{Foo} occurs in a file /path/to/file/Foo[.\lang\].tex which does not belong to an archive, the namespace is file://path/to/file.
- If the same statement occurs in a file /path/to/file/bar[.\(\lang\rang\rangle\)].tex, the namespace is file://path/to/file/bar.

In other words: outside of archives, the namespace corresponds to the file URI with the filename dropped iff it is equal to the module name, and ignoring the (optional) language suffix.

If the current file is in an archive, the procedure is the same except that the initial segment of the file path up to the archive's **source**-folder is replaced by the archive's namespace URI.

Conversely, here is how name spaces/URIs and file paths are computed in import statements, examplary $\verb|\info| import module:$

• \importmodule{Foo} outside of an archive refers to module Foo in the current namespace. Consequently, Foo must have been declared earlier in the same document or, if not, in a file Foo[.\langle langle].tex in the same directory.



- The same statement within an archive refers to either the module Foo declared earlier in the same document, or otherwise to the module Foo in the archive's top-level namespace. In the latter case, is has to be declared in a file Foo[. $\langle lang \rangle$].tex directly in the archive's source-folder.
- Similarly, in \importmodule(some/path?Foo) the path some/path refers to either the sub-directory and relative namespace path of the current directory and namespace outside of an archive, or relative to the current archive's top-level namespace and source-folder, respectively.

The module Foo must either be declared in the

file $\langle top\text{-}directory \rangle$ /some/path/Foo[. $\langle lang \rangle$].tex, or in $\langle top\text{-}directory \rangle$ /some/path[. $\langle lang \rangle$].tex (which are checked in that order).



- Similarly, \importmodule[Some/Archive] {some/path?Foo} is resolved like
 the previous cases, but relative to the archive Some/Archive in the mathhubdirectory.
- Finally, \importmodule{full://uri?Foo} naturally refers to the module Foo in the namespace full://uri. Since the file this module is declared in can not be determined directly from the URI, the module must be in memory already, e.g. by being referenced earlier in the same document. Since this is less compatible with a modular development, using full URIs directly is strongly discouraged, unless the module is delared in the current file directly.

\STEXexport

\importmodule and \usemodule import all symbols, notations, semantic macros and (recursively) \importmodules. If you want to additionally export e.g. convenience macros and other code from a module, you can use the command \STEXexport{<code>} in your module. Then <code> is executed (both immediately and) every time the current module is opened via \importmodule or \usemodule.



Note, that \newcommand defines macros globally and throws an error if the macro already exists, potentially leading to low-level IATEX errors if we put a \newcommand in an \STEXexport and the <code> is executed more than once in a document — which can happen easily.

A safer alternative is to use macro definition principles, that are safe to use even if the macro being defined already exists, and ideally are local to the current T_EX group, such as \def or \let .

3.4.3 The mathstructure Environment

A common occurrence in mathematics is bundling several interrelated "declarations" together into *structures*. For example:

- A monoid is a structure $\langle M, \circ, e \rangle$ with $\circ : M \times M \to M$ and $e \in M$ such that...
- A topological space is a structure (X,\mathcal{T}) where X is a set and \mathcal{T} is a topology on X
- A partial order is a structure $\langle S, \leq \rangle$ where \leq is a binary relation on S such that...

This phenomenon is important and common enough to warrant special support, in particular because it requires being able to *instantiate* such structures (or, ratherer, structure *signatures*) in order to talk about (concrete or variable) *particular* monoids, topological spaces, partial orders etc.

mathstructure

The mathstructure environment allows us to do exactly that. It behaves exactly like the smodule environment, but is itself only allowed inside an smodule environment, and allows for instantiation later on.

How this works is again best demonstrated by example:

Example 24

```
Input:
    \begin{mathstructure} { monoid}
 2
      \symdef{universe}[type=\set]{\comp{U}}}
 3
      \symdef{op}[
 4
          args=2,
          type=\funtype{\universe,\universe}{\universe},
 6
          op=\circ
      ]{#1 \comp{\circ} #2}
 8
      \symdef{unit}[type=\universe]{\comp{e}}
 9 \end{mathstructure}
10
11 A \symname{monoid} is...
```

Output:

```
A monoid is...
```

'Note that the \symname{monoid} is appropriately highlighted and (depending on your pdf viewer) shows a URI on hovering – implying that the mathstructure environment has generated a *symbol* monoid for us. It has not generated a semantic macro though, since we can not use the monoid-symbol *directly*. Instead, we can instantiate it, for example for integers:

Example 25

```
Input:

1 \symdef{Int}[type=\set]{\comp{\mathbb Z}}
2 \symdef{addition}[
3     type=\funtype{\Int,\Int}{\Int},
4     args=2,
5     op=+
6 ]{##1 \comp{+} ##2}
7 \symdef{zero}[type=\Int]{\comp{0}}
8
9 $\mathstruct{\Int,\addition!,\zero}$ is a \symname{monoid}.
```

Output:

```
\langle \mathbb{Z}, +, 0 \rangle is a monoid.
```

So far, we have not actually instantiated monoid, but now that we have all the symbols to do so, we can:

Example 26 Input:

```
1 \instantiate{intmonoid}{
2    universe = Int ,
3    op = addition ,
4    unit = zero
5 }{monoid}{\mathbb{Z}_{+,0}}
6
7    $\intmonoid{\universe}$, $\intmonoid{\unit}$ and $\intmonoid{\unitmonoid}$.
8
9 Also: $\intmonoid!$
```

Output:

```
\mathbb{Z}, 0 and a+b.
Also: \mathbb{Z}_{+,0}
```

\instantiate

So summarizing: \instantiate takes four arguments: The (macro-)name of the instance, a key-value pair assigning declarations in the corresponding mathstructure to symbols currently in scope, the name of the mathstructure to instantiate, and lastly a notation for the instance itself.

It then generates a semantic macro that takes as argument the name of a declaration in the instantiated **mathstructure** and resolves it to the corresponding instance of that particular declaration.

\instantiate and mathstructure make use of the *Theories-as-Types* paradigm:

mathstructure{<name>} does in fact simply create a nested theory with name

<name>-structure. The constant <name> is defined as Mod(<name>-structure)
a dependent record type with manifest fields, the fields of which are generated from

(and correspond to) the constants in <name>-structure.

\instantiate appropriately generates a constant whose definions is a record torm.

\instantiate appropriately generates a constant whose definiens is a record term of type Mod(<name>-structure), with the fields assigned appropriately based on the key-value-list.

Notably, \instantiate throws an error if not *every* declaration in the instantiated mathstructure is being assigned.

You might consequently ask what the usefulness of mathstructure even is.

\varinstantiate

The answer is that we can also instantiate a mathstructure with a *variable*. The syntax of \varianstantiate is equivalent to that of \instantiate, but all of the key-value-pairs are optional, and if not explicitly assigned (to a symbol *or* a variable declared with \vardef) inherit their notation from the one in the mathstructure environment.

This allows us to do things like:

Example 27 Input:

Eco 1/12 .— (23,5,6) & monord on 22...

We will return to this example later, when we also know how to handle the axioms of a monoid.

3.4.4 The copymodule Environment

TODO: explain

Given modules:

```
Example 28
```

```
Input:

1 \begin{smodule}{magma}
2 \symdef{universe}{\comp{\mathcal U}}
3 \symdef{operation}[args=2,op=\circ]{#1 \comp\circ #2}
4 \end{smodule}
5 \begin{smodule}{monoid}
6 \importmodule{magma}
7 \symdef{unit}{\comp e}
8 \end{smodule}
9 \begin{smodule}{group}
10 \importmodule{monoid}
11 \symdef{inverse}[args=1]{{#1}^{\comp{-1}}}
12 \end{smodule}
```

Output:

.

We can form a module for *rings* by "cloning" an instance of <code>group</code> (for addition) and <code>monoid</code> (for multiplication), respectively, and "glueing them together" to ensure they share the same universe:

Example 29

Input:

```
\begin{smodule}{ring}
      \begin{copymodule} { group } { addition }
3
          \renamedecl[name=universe] {universe} {runiverse}
4
          \renamedecl[name=plus]{operation}{rplus}
5
          \renamedecl[name=zero]{unit}{rzero}
6
          \renamedecl[name=uminus]{inverse}{ruminus}
 7
      \end{copymodule}
8
      \notation*{rplus}[plus,op=+,prec=60]{#1 \comp+ #2}
9
          \notation*{rzero}[zero]{\comp0}
10
          \notation*{ruminus}[uminus,op=-]{\comp- #1}
          \begin{copymodule}{monoid}{multiplication}
11
          \assign{universe}{\runiverse}
12
13
          \renamedecl[name=times] {operation} {rtimes}
14
          \renamedecl[name=one]{unit}{rone}
15
      \end{copymodule}
16
      \notation*{rtimes}[cdot,op=\cdot,prec=50]{#1 \comp\cdot #2}
17
          \notation*{rone}[one] {\comp1}
18
          Test: $\rtimes a{\rplus c{\rtimes de}}$
19 \end{smodule}
```

Output:

```
Test: a \cdot (c + d \cdot e)
```

TODO: explain donotclone

3.4.5The interpretmodule Environment

TODO: explain

```
Example 30
```

```
Input:
   \begin{smodule}{int}
2
      \symdef{Integers}{\comp{\mathbb Z}}
3
      \symdef{plus}[args=2,op=+]{#1 \comp+ #2}
      \symdef{zero}{\comp0}
 4
 5
6
      \symdef{uminus}[args=1,op=-]{\comp-#1}
      \begin{interpretmodule}{group}{intisgroup}
          \assign{universe}{\Integers}
9
          \assign{operation}{\plus!}
10
          \assign{unit}{\zero}
          \assign{inverse}{\uminus!}
      \end{interpretmodule}
13 \end{smodule}
```

Output:

3.5 Primitive Symbols (The ST_EX Metatheory)

TODO: metatheory documentation

Chapter 4

Using STEX Symbols

Given a symbol declaration \symdecl{symbolname}, we obtain a semantic macro \symbolname. We can use this semantic macro in math mode to use its notation(s), and we can use \symbolname! in math mode to use its operator notation(s). What else can we do?

4.1 \symmet and its variants

\symref \symname

We have already seen \symname and \symref, the latter being the more general.

\symref{<symbolname>}{<code>} marks-up <code> as referencing <symbolname>. Since quite often, the <code> should be (a variant of) the name of the symbol anyway, we also have \symname{<symbolname>}.

Note that \symname uses the *name* of a symbol, not its macroname. More precisely, \symname will insert the name of the symbol with "-" replaced by spaces. If a symbol does not have an explicit name = given, the two are equal – but for \symname it often makes sense to make the two explicitly distinct. For example:

Example 31

```
Input:

1 \symdef{Nat}[
2    name=natural-number,
3    type=\set
4 ]{\comp{\mathbb{N}}}
5
6 A \symname{Nat} is...
Output:
```

A natural number is...

\symname takes two additional optional arguments, pre= and post= that get prepended or appended respectively to the symbol name.

 \S ymname

Additionally, \Symname behaves exactly like \symname, but will capitalize the first letter of the name:

Example 32

Input:

1 \Symname[post=s]{Nat} are...

Output:

Natural numbers are...

This is as good a place as any other to explain how STEX resolves a string symbolname to an actual symbol.

If \symbolname is a semantic macro, then STEX has no trouble resolving symbolname to the full URI of the symbol that is being invoked.

However, especially in \symname (or if a symbol was introduced using \symdec1* without generating a semantic macro), we might prefer to use the *name* of a symbol directly for readability — e.g. we would want to write A \symname{natural-number} is... rather than A \symname{Nat} is... STEX attempts to handle this case thusly:



If string does not correspond to a semantic macro $\sl TEX$ checks all symbols currently in scope until it finds one, whose full URI ends with string. This allows for disambiguating more precisely, e.g. by saying $\sl TEX$ where several additions are in scope.

However, this also means that if we have symbols foo and e.g. miraculous-foo, then STEX might resolve \symname{foo} to miraculous-foo if it finds this symbol first. It is therefore a good idea to prefix symbol names with a ?, thus ensuring that STEX will find the symbol ...?foo rather than ...?miraculous-foo.

4.2 Marking Up Text and On-the-Fly Notations

We can also use semantic macros outside of text mode though, which allows us to annotate arbitrary text fragments.

Let us assume again, that we have $\symdef{addition}[args=2]{#1 \comp+ #2}$. Then we can do

Example 33

Input:

```
1 \addition{\comp{The sum of} <math display="inline">\arg{s\argn}} \and \arg{s\argm}} \argnet{2 is...}
```

Output:

The sum of n and m is...

'...which marks up the text fragment as representing an application of the addition-symbol to two argument n and m.

\arg

In text mode, every semantic macro takes exactly one argument, namely the text-fragment to be annotated. The \arg command is only valid within the argument to a semantic macro and marks up the *individual arguments* for the symbol.

We can also use semantic macros in text mode to invoke an operator itself instead of its application, with the usual syntax using !:

```
Example 34
Input:

1 \addition! {Addition} is...

Output:

Addition is...
```

In deed, \symbolname! {<code>} is exactly equivalent to \symref {symbolname} {<code>} (the latter is in fact implemented in terms of the former).

\arg also allows us to switch the order of arguments around and "hide" arguments: For example, \arg[3]{<code>} signifies that <code> represents the *third* argument to the current operator, and \arg*[i]{<code>} signifies that <code> represents the *i*th argument, but it should not produce any output (it is exported in the xhtml however, so that MMT and other systems can pick up on it)

Example 35

```
Input:

1 \addition{\comp{adding}
2 \arg[2]{$\svar{k}$}
3 \arg*{$\addition{\svar{n}}{\svar{m}}$}} yields...

Output:

adding k yields...
```

Note that since the second \arg has no explicit argument number, it automatically represents the first not-yet-given argument – i.e. in this case the first one.

The same syntax can be used in math mode, too, which allows us to spontaneously introduce new notations on the fly. We can activate it using the starred variants of semantic macros:

Example 36

```
Input:
```

```
1 Given $\addition{\svar{n}}{\svar{m}}$, then
2 $\addition*{
3  \arg*{\addition{\svar{n}}}{\svar{m}}}
4  \comp{+}
5  \arg{\svar{k}}
6 }$ yields...
```

Output:

```
Given n+m, then +k yields...
```

4.3 Referencing Symbols and Statements

TODO: references documentation

Chapter 5

STEX Statements

5.1 Definitions, Theorems, Examples, Paragraphs

As mentioned earlier, we can semantically mark-up *statements* such as definitions, theorems, lemmata, examples, etc.

The corresponding environments for that are:

- sdefinition for definitions,
- sassertion for assertions, i.e. propositions that are declared to be *true*, such as theorems, lemmata, axioms,
- sexample for examples, and
- sparagraph for other semantic paragraphs, such as comments, remarks, conjectures, etc.

The *presentation* of these environments can be customized to use e.g. predefined theorem-environments, see chapter 6 for details.

All of these environments take optional arguments in the form of key=value-pairs. Common to all of them are the keys id= (for cross-referencing, see section 4.3), type= for customization (see chapter 6) and additional information (e.g. definition principles, "difficulty" etc), title=, and for=.

The for= key expects a comma-separated list of existing symbols, allowing for e.g. things like

Example 37

```
Input:

1 \begin{sexample}[
2   id=additionandmultiplication.ex,
3   for={addition,multiplication},
4   type={trivial,boring},
5   title={An Example}
6 ]
7   $\addition{2,3}$ is $5$, $\multiplication{2,3}$ is $6$.
8 \end{sexample}
```

Output:

Example 5.1.1 (An Example). 2+3 is 5, $2\cdot 3$ is 6.

\definiendum \definame \definiens \Definame

sdefinition (and sparagraph with type=symdoc) introduce three new macros: definiendum behaves like symref (and definame/Definame like symname/Symname, respectively), but highlights the referenced symbol as being defined in the current definition.

\definiens [<optional symbolname>] {<code>} marks up <code> as being the explicit definiens of <optional symbolname> (in case for= has multiple symbols).

The special type=symdoc for sparagraph is intended to be used for "informal definitions", or encyclopedia-style descriptions for symbols.

The MMT-system can use those (in lieu of an actual sdefinition in scope) to present to users, e.g. when hovering over symbols.

All four environments also take an optional parameter name = – if this one is given a value, the environment will generate a *symbol* by that name (but with no semantic macro). Not only does this allow for \symref et al, it allows us to resume our earlier example for monoids much more nicely:

Example 38

```
Input:
    \begin{mathstructure} { monoid}
 2
       \symdef{universe}[type=\set]{\comp{U}}}
 3
       \symdef{op}[
 4
           args=2,
 5
           type=\funtype{\universe, \universe}{\universe},
 6
           op=\circ
 7
       ]{#1 \comp{\circ} #2}
 8
       \symdef{unit}[type=\universe]{\comp{e}}
10
       \begin{sparagraph} [type=symdoc,for=monoid]
           A \definame{monoid} is a structure
11
12
           $\mathstruct{\universe,\op!,\unit}$
13
           where $\op!:\funtype{\universe}{\universe}$ and
14
           $\inset{\unit}{\universe}$ such that
15
\frac{16}{17}
           \begin{sassertion} [name=associative,
               type=axiom,
18
               title=Associativity]
19
               $\op!$ is associative
20
           \end{sassertion}
\frac{21}{22}
           \begin{sassertion} [name=isunit,
               type=axiom,
23
               title=Unit]
24
               \displaystyle {\displaystyle \{ \op{\svar}(x)}{\svar}(x)} 
25
               for all $\inset{\svar{x}}{\universe}$
26
           \end{sassertion}
27
       \end{sparagraph}
28 \end{mathstructure}
30 An example for a \symname{monoid} is...
```

Output:

```
A monoid is a structure \langle U, \circ, e \rangle where \circ : U \to U and e \in U such that 
Axiom 5.1.2 (Associativity). \circ is associative 
Axiom 5.1.3 (Unit). x \circ e = x for all x \in U 
An example for a monoid is...
```

Now the mathstructure monoid contains two additional symbols, namely the axioms for associativity and that e is a unit. Note that both symbols do not represent the mere propositions that e.g. \circ is associative, but the assertion that it is actually true that \circ is associative.

If we now want to instantiate monoid (unless with a variable, of course), we also need to assign associative and neutral to analogous assertions. So the earlier example

```
1 \instantiate{intmonoid}{
2    universe = Int ,
3    op = addition ,
4    unit = zero
5 }{monoid}{\mathbb{Z}_{+,0}}
```

...will not work anymore. We now need to give assertions that addition is associative and that zero is a unit with respect to addition.²

5.2 Proofs

TODO

 $^{^2}$ Of course, STeX can not check that the assertions are the "correct" ones – but if the assertions (both in monoid as well as those for addition and zero) are properly marked up, MMT can. TODO: should

Chapter 6

Highlighting and Presentation Customizations

The environments starting with s (i.e. smodule, sassertion, sexample, sdefinition, sparagraph and sproof) by default produce no additional output whatsoever (except for the environment content of course). Instead, the document that uses them (whether directly or e.g. via inputref) can decide how these environments are supposed to look like.

The stexthm defines some default customizations that can be used, but of course many existing LATEX templates come with their own definition, theorem and similar environments that authors are supposed (or even required) to use. Their concrete syntax however is usually not compatible with all the additional arguments that STEX allows for semantic information.

Therefore we introduced the separate environments **sdefinition** etc. instead of using **definition** directly, and allow authors to specify how these environments should be styled via the commands **stexpatch***.

\stexpatchmodule \stexpatchdefinition \stexpatchassertion \stexpatchexample \stexpatchparagraph \stexpatchproof All of these commands take one optional and two proper arguments, i.e. \stexpatch*[<type>]{<begin-code>}{end-code}.

After SIEX reads and processes the optional arguments for these environments, (some of) their values are stored in the macros \s*<field> (i.e. sexampleid, \sassertionname, etc.). It then checks for all the values <type> in the type=-list, whether an \stexpatch*[<type>] for the current environment has been called. If it finds one, it uses that patches <begin-code> and <end-code> to mark up the current environment. If no patch for (any of) the type(s) is found, it checks whether and \stexpatch* was called without optional argument.

For example, if we want to use a predefined theorem environment for sassertions with type=theorem, we can do

1 \stexpatchassertion[theorem] {\begin{theorem}} {\end{theorem}}

...or, rather, since e.g. theorem-environments defined using amsthm take an optional title as argument, we can do:

- 1 \stexpatchassertion[theorem]
- 2 {\ifx\sassertiontitle\@empty
- 3 \begin{theorem}

```
4  \else
5  \begin{theorem}[\sassertiontitle]
6  \fi}
7  {\end{theorem}}
```

Or, if we want all sdefinitions to use a predefined definition-environment, we can do

```
1 \stexpatchdefinition
2 {\ifx\sdefinitiontitle\@empty
3 \begin{definition}
4 \else
5 \begin{definition}[\sdefinitiontitle]
6 \fi}
7 {\end{definition}}
```

\compemph
\varemph
\symrefemph
\defemph

Apart from the environments, we can control how STEX highlights variables, notation components, \symrefs and \definiendums, respectively.

To do so, we simply redefine these four macros. For example, to highlight notation components (i.e. everything in a \comp) in blue, as in this document, we can do \def\compemph#1{\textcolor{blue}{#1}}. By default, \compemph et al do nothing.

\compemph@uri \varemph@uri \symrefemph@uri \defemph@uri For each of the four macros, there exists an additional macro that takes the full URI of the relevant symbol currently being highlighted as a second argument. That allows us to e.g. use pdf tooltips and links. For example, this document uses

```
1 \protected\def\symrefemph@uri#1#2{
2 \pdftooltip{
3 \srefsymuri{#2}{\symrefemph{#1}}}
4 }{
5 URI:~\detokenize{#2}
6 }
7 }
```

By default, $\compemph@uri$ is simply defined as $\compemph{#1}$ (analogously for the other three commands).

Chapter 7

Additional Packages

TODO: tikzinput documentation

7.1 Modular Document Structuring

TODO: document-structure documentation

7.2 Slides and Course Notes

TODO: notesslides documentation

7.3 Homework, Problems and Exams

TODO: problem documentation
TODO: hwexam documentation