

# UNIT - 5

## "GRAPH THEORY"

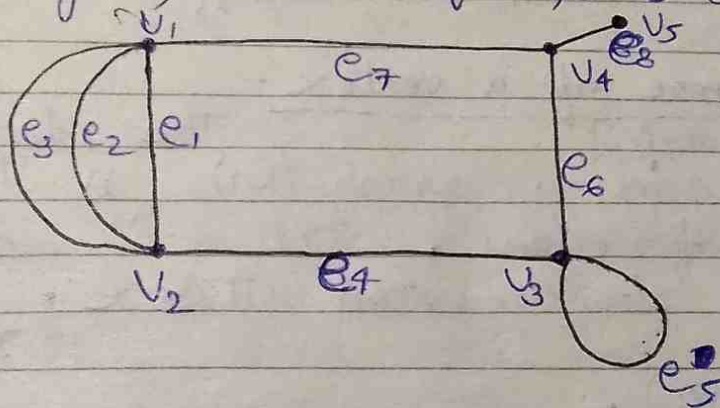
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### ★ Graph -


A Graph  $G = (V, E)$  consist of a set of object  $V = \{V_1, V_2, V_3, \dots\}$  whose elements are called vertices and a another ~~set~~ set  $E = \{e_1, e_2, e_3, \dots\}$  whose elements are called edges, and set of  $V, E$  is called a graph.

### Example -

Sari definati  
mai example  
dikhna hai



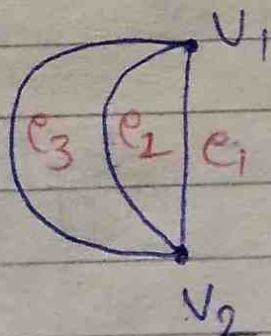
② Self loop - (Self loop) An edge is said to be a self loop, if its both end vertices are same.

example - 

### ③ Parallel edge -

If there are two or more than two having same pair of vertices, then such edges are called parallel edge.

example -



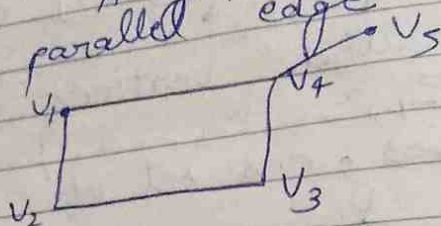


graph

④ Simple graph -

A graph that has neither self loop nor parallel edge are called simple graph.

example -



⑤ Degree of a vertex -

The degree of a vertex  $\deg(V_i)$  in a graph  $(G)$  is equal to the no. of edges, which are incident on  $V_i$  with self loop counted twice. It is denoted by  $\deg(V_i)$ .

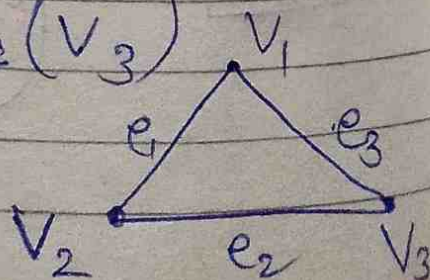
ex  $\Rightarrow \deg(V_1) = \deg(V_2) = \deg(V_3) = 4$

$\deg(V_4) = 3, \deg(V_5) = 1$

⑥ Regular graph -

A graph  $(G)$  in which all vertices are of equal degree is called regular graph.

Ex.  $\Rightarrow \deg(V_1) = \deg(V_2) = \deg(V_3)$

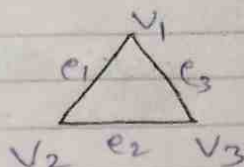




## ⑦ Isolated vertex -

A vertex of degree zero is called isolated vertex, or an end vertex.

ex.  $\Rightarrow$

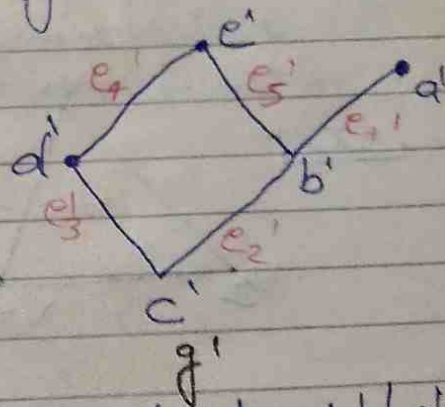
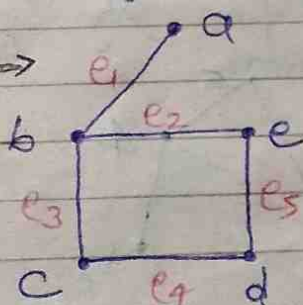


$\rightarrow V_4$  is isolated vertex

## ⑧ Iso-morphic graph -

Two graphs are called iso-morphic, if their graph ~~theoretic~~ properties are same that is they have same no. of edge, they have equal no. of vertices with a given degree, they have same no. of vertices.

ex.  $\Rightarrow$



The vertices  $a, b, c, d, e$  correspond to  $a', b', c', d', e'$  respectively and the edges  $e_1, e_2, e_3, e_4, e_5$  correspond to  $e_1', e_2', e_3', e_4', e_5'$  respectively.

A graph  $G'$  is said to be sub-graph of graph  $G$ . If all the vertices and all edges of  $G'$  are in  $G$ .



Each edge of  $(g')$  has the same end vertices in  $(g)$  as in  $g$ .

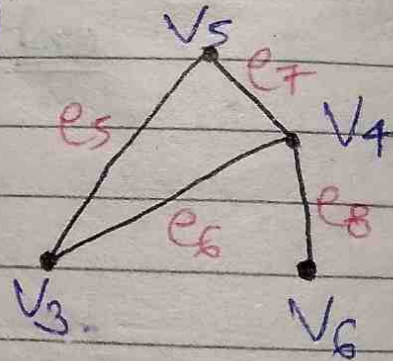
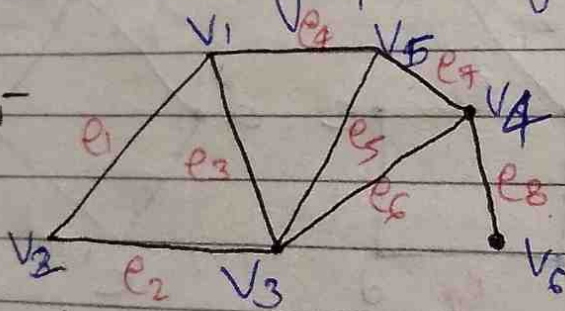
GATE  $\rightarrow$  MCQs  
#Note

From the definition of sub-graph

- ① Every graph is its own sub-graph.
- ② A single vertex in  $(g)$  is a sub-graph of  $(g)$ .
- ③ A single edge in  $(g)$  ~~is~~ together with its end vertices is a sub-graph of  $(g)$ .

- ④ A sub-graph of a sub-graph of  $(g)$  is also a sub-graph of  $(g)$ .

ex -



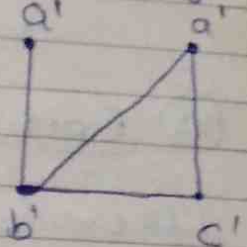
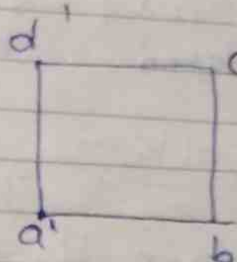
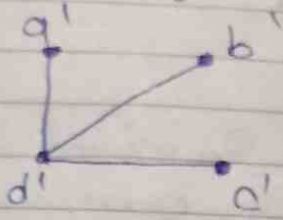
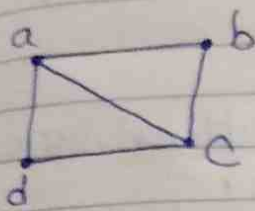
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vertices

## Spanning sub-graph

Let  $G' = (V', E')$  be a subgraph of a graph  $G = (V, E)$ . If  $V' = V$ , then  $G'$  is said to be spanning sub-graph of  $G$ .



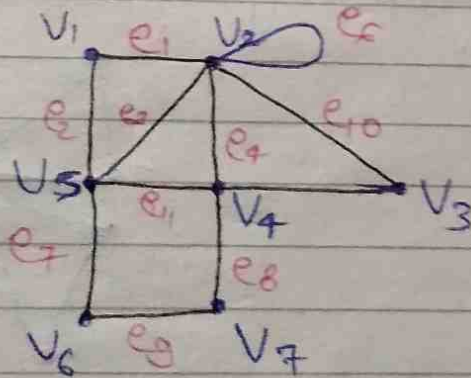
(g).

its

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## Walk

A walk in a graph  $(g)$  is defined as a finite alternating sequence of vertices & edges which begins and ends, with vertices.



In a walk, no edge appear more than one however, a ~~vertex~~ vertex may appear more than once.

→  $V_1 e_1 V_2 e_6 V_2 e_{10} V_3 e_5 V_4 e_8 V_7$   
walk →



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(11) Closed Walk -

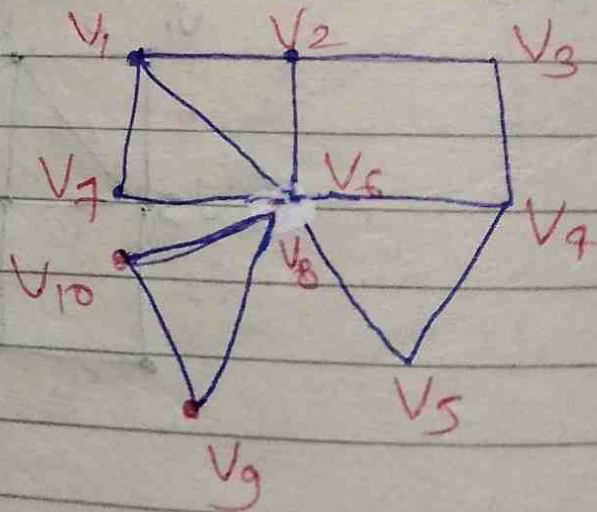
If terminal vertices are same then walk is called closed walk. A walk  $V_5 e_1 V_4 e_5 V_2 e_6 V_2 e_1 V_1 e_2 V_5$  is a closed walk.

(12) Open walk -

When terminal vertices are different, then it is an open walk.  $V_5 e_1 V_4 e_5 V_7$  is an open walk.

(13) Connected Graph -

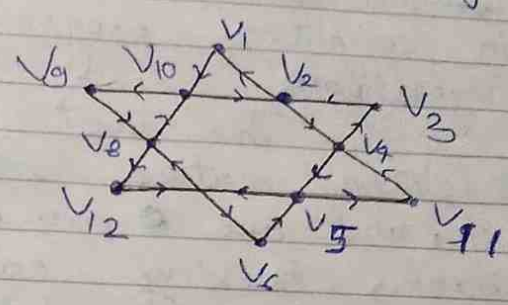
A graph  $(G)$  is called connected, if there is at least one path b/w every pair of vertices. Otherwise ~~graph~~ it is disconnected.





★ Euler graph -

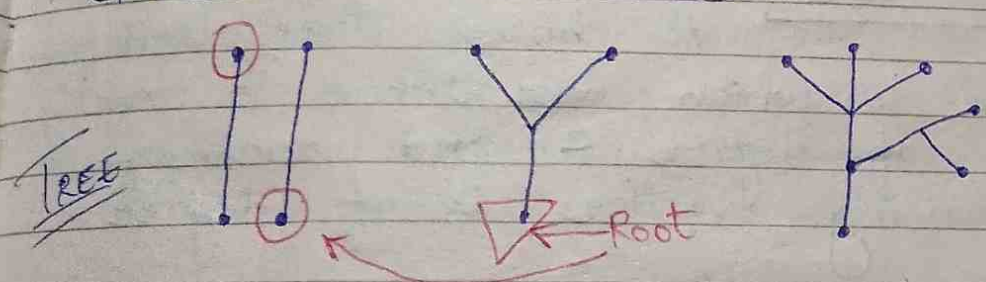
A closed walk in a graph which include all the edges of the graph is euler graph.



IMP PYQs

★ Tree -

A connected graph having no circuit is called tree.



• Some Types of Trees -

① Rooted tree -

A rooted tree is a tree with distinguish vertices called the root.

Denoted by  $\bigcirc$   $\&$   $\nabla$



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(2) Decision tree -

A decision tree are labelled rooted tree which occur specially in application programming & computer algorithm.

The root represent a starting point later vertices represent later decision point, and one proceed downward to the tree, choosing edge at each step occur to observed data.

(3) Binary Tree -

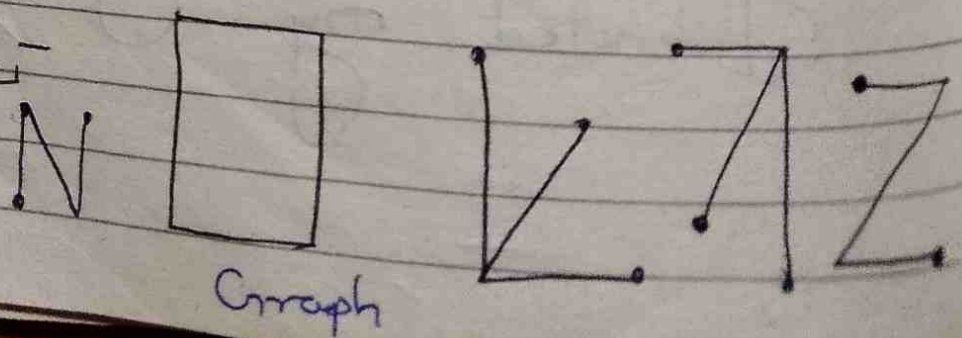
A binary tree is defined as a tree in which ~~var~~ there is exactly one vertex of degree 2 and each of the remaining vertices is of degree 1 or 3.

(4) Spanning Tree -

If  $g = (V, E)$  is any connected graph, a spanning tree in  $g$  is a sub-graph  $T = (V, E')$  which is a tree.

Example -

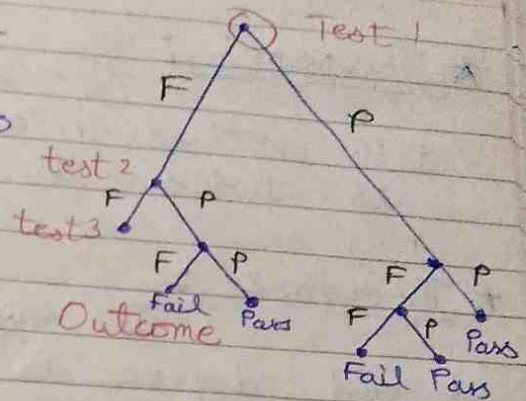
Spanning Tree





### ⑤ Binary decision tree,-

A student must pass two of three test to pass a course. The binary decision reappears in adjoining figure with F denoting Fail, P denoting pass.



If a student passes both of the first two test or fail both there is no decision, needed from test 3. We have ~~emphasized~~ emphasized the path to the outcome for a student who passes the first test fail the second and passes the third: the student passes the course.

In a binary tree, each decision has only two possible outcomes that is yes (OR) No, True (OR) False, 0 (OR) 1).



6 Marks Que.

★ Theorem 1 :

The sum of the degrees of all vertices in a graph is equal to twice the no. of edges.

★ Theorem 2 :

In any graph, the no. of vertices of odd degree is always even.

★ Theorem 5 :

The maximum no. of edges in a simple graph with  $n$  vertices is  $\frac{n(n-1)}{2}$ .



★ Properties of double integration —

① If  $K$  is constant func.<sup>n</sup> then, double integration are  $\iint_R K f(x, y) dx dy = K \iint_R f(x, y) dx dy$

② Linear property of double integral

$$\iint_R \{K_1 f_1(x, y) + K_2 f_2(x, y)\} dx dy$$

$$= K_1 \iint_R f_1(x, y) dx dy + K_2 \iint_R f_2(x, y) dx dy$$

③ If the region  $(R)$  is par into  $R_1, R_2$

$$\iint_R f(x, y) dx dy = \iint_{R_1} f(x, y) dx dy + \iint_{R_2} f(x, y) dx dy$$

④ The double integral  $\int_a^b \int_{f_1(x)}^{f_2(x)} dx dy$  defines

the area enclosed by the  $y = f_1(x), y = f_2(x)$  & the ordinates  $x = a, x = b$ .



## (DNF) & (CNF)

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★ Involution law, -  $(a')' = a$

★ Demorgan's law, -  $(a+b)' = a'b'$   
 $(a.b)' = a' + b'$

★ In the Boolean Algebra, prove the following.

(i)  $a' + ab = a' + b$

(ii)  $a.b \Rightarrow a.b' = 0 \quad \forall a, b \in B$

(iii)  $a.b + b.c + c.a = (a+b)(b+c)(c+a)$

★ Boolean func<sup>n</sup>, -

An expression obtained by the application of binary operations '+' & '.' and a unary operation (') on the finite number of elements of boolean-algebra  $(B, +, ')$  is called Boolean func<sup>n</sup> @ OR polynomial.

★ Minimal Boolean function, -

A minimal boolean func<sup>n</sup> is (n) variable  $x_1, x_2, x_3, \dots, x_n$  is product of (n) letters.

★ Disjunctive normal form (DNF), -

A boolean polynomial which can be written @ sum of the minimal boolean function called disjunctive normal form (or conomital form)

$$f(x, y) = x.y + x'y + x.y' + x'y'$$



2 mai comm

★ Conjunctive Normal form - A boolean polynomial is called a CNF (or dual canonical form) if it is the product of distinct factors where each factor is the sum of variables  $x_1, x_2, x_3, \dots$  and their products  $x_1', x_2', x_3', \dots$  and in each factor variables or their complements do not occur more than once

$$f(x_1, x_2) = (x_1 + x_2)(x_1' + x_2)(x_1 + x_2')(x_1' + x_2')$$

Ex. - Convert  $x + x'y$  to DNF.

$$x(1) + x' \cdot y$$

$$x(y + y') + x'y$$

$$xy + xy' + x'y \quad \text{Ans.}$$

► Disjunctive (DNF)

Ex. - Write the following func.<sup>n</sup> into conjunctive normal form, in which maximum number of variable are used.

$$(i) f(x, y, z) = x \cdot y' + xz + xy$$



Que. Test for Consistency  $AX = B$  Solve :-

$$= \begin{aligned} 2x - 3y + 7z &= 5 \\ 3x + y - 3z &= 13 \\ 2x + 19y - 47z &= 22 \end{aligned}$$

$$C = [A:B]$$

Rank of A = Rank of C  
then consistent

★ If Rank = no. of unknown  $(x, y, z)$   
then unique sol<sup>n</sup>  $\Rightarrow$

★ Rank < no. of unknown  $\Rightarrow$  Infinite sol<sup>n</sup>

$$AX = B$$

$$\begin{bmatrix} 2 & -3 & 7 \\ 3 & 1 & -3 \\ 2 & 19 & -47 \end{bmatrix} \begin{bmatrix} x \\ y \\ z \end{bmatrix} = \begin{bmatrix} 5 \\ 13 \\ 22 \end{bmatrix}$$

Now, Augmented Matrix  
 $C = [A:B]$

$$C = \begin{bmatrix} 2 & -3 & 7 & \vdots & 5 \\ 3 & 1 & -3 & \vdots & 13 \\ 2 & 19 & -47 & \vdots & 22 \end{bmatrix}_{3 \times 4}$$

For Finding Rank  
Convert in  
upper triangular  
Matrix

(-) lower triangular matrix

$$R_2 = R_2 - R_1 \quad \text{(OR)} \quad 2R_2 - 3R_1$$

$$R_3 = R_3 - R_1$$



$$C = \begin{bmatrix} 2 & -3 & 7 & : & 5 \\ 0 & 11 & -27 & : & 11 \\ 0 & 22 & -54 & : & 17 \end{bmatrix}$$

$$R_3 = R_3 - 2R_2$$

$$C = \begin{bmatrix} 2 & -3 & 7 & : & 5 \\ 0 & 11 & -27 & : & 11 \\ 0 & 0 & 0 & : & -5 \end{bmatrix}$$

← Now it is upper triangular matrix

\* Rank of a matrix = No. of non-zero row is called rank of a matrix.

Now, In C matrix

No. of non-zero row = 3

Rank of Matrix = 3

$$A = \begin{bmatrix} 2 & -3 & 7 \\ 0 & 11 & -27 \\ 0 & 0 & 0 \end{bmatrix}$$

Upper triang. matrix

Rank of a matrix

= No. of Non-Zero Row

Rank = 2

of matrix A

Rank of a matrix  
find kme ka  
Yeh rule  
hai Matrix  
Upper  
triangular  
matrix  
hon chahye!!

[∴ Rank of C ≠ Rank of A]

[∴ The given system is not consistent  
and there exist No Sol.]