Scanned by CamScanner



	MACER'S M
3)	How does a DSC mode & where can I purchase DSC?
- Ans	A unique digital Angeopoint (hash) of a document
	is executed using any hashing algorithm.
	Than hash is encoypted using signess polvate key
	- Encrybed hash & signer's public kay are appended
	are combined into a digital signature, which is
	then appended to the document.
	- DSC Applicants can disectly approach Costifying
	Anthosities wand issue a DSC using Andhay
	ekyc authentication. A licensed CA issues
	Digital Signature.
	- Some CAS are Mational Informaties Contor (NIC)
	IDRBT CA, SafeScoypt CA services, E-MUDHRA, ato
	(n) code solutions CA, TCS, CDAC, NSDL Capticoon.
4)	Mat are different classes of DSCs?
Bors	Different classes of DSCs are-
	class 2: Here, the identity of a person is
	resified against a bousted, pre-verified database.
	o vocación por viciniza accomina
	Class 3 - This is 12 191 A 1 1
	class 3: This is the highest level where the
	person needs to present himself or horself
	1 Front of a Regular Authority (CA)
	Prove his/hor identily.

3) what is difference between digitally signing &	_
encoypting an email?	
is Encoypting an email probects the privacy of the	_
message by convexting from plain text to	_
cipher beat. Only the recipient who has private	
key that matches the public key you used to	
encoypt the message can deligher the message	,
- By digitally signing a message, you apply your	
migne digital mast bo ble message, The digital)
signature includes your coobificate & public key.	
This information possibles bother poures by the	
recipient that you signed the contents of a	2
message & not an imposter, & that the contain	t
have not been alberted in boarsit	
6) mad actually hoppens when I digitally	_
sign any Exampaction?	
- Transforming a process to make it entirely	
electronic requires sensing 2 giving legal	
value to the transactions in rational	
00 0 - 0000 3000 G - 1110 G - 1	
Hence secure & legally binding brompactions	_
Hence secure & legally binding bromsactions	<u> </u>
Dave necessary in mobility when transactions	<u>—</u>
is going on between two parties.	_
Digitally signing any boarsactions generates	_
Digitally signing any boarsaction openerates a legally-binding pood proof of bransaction,	_
Digitally signing any boardaction generates	_
Digitally signing any boardaction openerates a legally-binding pood proof of bransaction,	_

+ FAQ's :-1. what is encoyption & decoyption of data? ANT The process of connecting the plain bext message be cipher beat message is known as encryption. leb, P is plainteat, E is ancoyptron & C is appropried. If we perform F on P we get C. P E(P) The process of restoring the plain beat from the cipher beat message is known as decryption or process of converting ciphertest message into plaintest message is called as decoyption let, C se aphabent, Dis decryption & Pisplaintext. IP we perform Donc it result in original P 1-e P= D(C), 2. How can Cassas Cipher be cracked? 19- These are 2 cases-(i) It abtacker does not know encryption bechnique then using bechniques such as Frequency analysis can be used to a crack the ciphor & find bretey. (1) If abbacker knows that caesar cipher is used then busboute Porce abback can be used to find the key as there are only 25 shift values possible in case of casas apper bechingue.

ATT.			
100		p substi	bublen boxes
0)	and is the in	mpostance of substi	x) in DGS!
3)	(C 1 + 401) A	# 6	30/10000
Ade	Clost is a bos	acc which	pextooms substitutes
Hes	key algorithe	ne like DES acres ontion 48-bib input	pextorms substitution.
	In S-box substit	ention 48-610	nded sight plain heat
Action Lawy to second Style Action			
	M		
		CIENT EST	ound the outsid
	+1 10 AV	at on the race	7000
	Som each St	ox immediately after	eers as many
	officers on box	sible. The 32-bits a	se permied by
	replacement of e	ach 416 with another	a bit in specified
	bable called a	is P-box permutati	on,
1		E	
		ne between stream	
	Book for comparision	Stream Cipher	Stream Hocks Cipher
	Basic	Converts the plain	Converbs the beat by
		text by baking its	taleng one byte of
-		block at a time	plain beat at a time
	Complexity	Simple design	
3	No of buts used	64-bits or niose	Complex Compasatively
	Confusion & Diffusion	uses both	8 bits
- 0)	Agorian moder	FCR(Flaction	Uses Confusionouly
And the Party of t	Used	ECB(Blacksonic Code Books	1 CFB (Cipher Feedback)
	leversibility_	CBC (Cipher Block Chaining	9) OFB(Ontput Foodball)
	0	encoupled bed	we XOR 60 encoypo
(6)		is hard	112-1 - be easily
(1)	eg Edanigle		As plant
		Eistef Cipher	Cioled
			Ne gram Cipheo

5)_	what is the most slewer way to design the
	forgot persurad feature?
4	By using a seconcery phone number or email id.
	If user clicks Forgot password then take
	in out of usex email id fixst.
	Then ark for the recovery mobile no eman id.
	IF recovery details mater them sing are
	(One bime password) bo secovery mobile no
	email id.
	IF OTP is resirred then only ask uses to
	auted new password & verify it and hence,
	change the pass wood successfully.
	anegre propo one server pay.
	ane production of

-			MAEER	's MIT
*	FAG's:	1 * 2		
)	unat semsity prop function exhibits!	objes sho	uld a grad band	
	function exhibits?		Jusa Mag	^
J-	The hash value is data being hoshed.	Fully of.	bemains 1 1 12	
1	- All input data mu	it be us	ed maising	
	distribution of data	scores a	cxirs a Air sot	
	possible values.		Cours only the	ot
	Different hosh valu	e for som	e shrina	
	The state of the s	1	30011	
2)	anat is MAC9 Differen	entrate MA	1C & MO.	
	MAC sometimes know	en as a be	ug is short pie	
	of information used	60 auth	entrate a messa	98
	- 2 Bagic steps: Builde	ing a bag s	verifying a bar	9.
	V 10 2	O	0 0	<u> </u>
	Basis of Comparision	MD	MAC	
	Integrity	Yes	Yes	
_	Aubhortication	No	Yes	
	Non-repudiation	No	No	-
_	kind of koys	none	symmetric keys	
	,			
_	A host of the message	e, if appen	ded to the message	2
	itself, only probects a	gainst accide	ental changes to t	ne !
-	message, as an abbae	ker who me	1 we it inshead	
	can simply calculate	a now passe	aires integriby.	
-	of the original one. So A MAC probects orgain	et mosserge fe	regard by anyone	who
1	doesn't know the sec	get key.		
	asegue know one se	0		



(= 1)					WAEER'S
3)	Compare MO5 &	SHA1			
	Basis of Congasisn		SHA	1 1	105
	Size I message in 1	1	160		128
	Coyptanalysis attack	-	Vulnera	ble Mob	rubnerable
(M)	Speed	·	Slow	,	Fost
(iv)	Number of Steps		8.0	<u> </u>	64
	Buffergize in bits		160		128
					·
4)	Compase vasions	vescro	m of SM	Α.	
	Basis of Compasision		SHA-1	SHA-256/224	SHA-512/384
	Size of message in bits	160	160	256/224	512/384
	Size of internal	160	160	286	SIZ
	state in bits				i i i
(M)	Size of block in bibs	512	512	512	1024
(iv)	length of messenge	264-1	2-1	2 -1	2 -1
(v)	Size of word	32	32	32	64
(vi)	No. of steps	80	80	64	80
(IIV)	Collision	Yes	23 attack	Not yet	None yet
				The organization	100 900
				•	
				*	
					7
				1	
					21 ⁴ 2.5 1 4

9.	W/ALLIY O WIT
*	FAG's:
1)	Enlist varous algoribhus used for primality
	testing.
14-	various algorithms used for primality besting are-
	(i) Miller-Rabin Primality Test
(TiD AKS primality best,
(III) Fermat primality best.
	N) Fro benus primality test.
(v) Pocklington Primality Test.
	vi) Filiptic curve Poimality Test-
2)	what are application of primality bests?
	(i) Gryptography schemes such as RSA algorithm.
	(11) Compubational number theory.
	(III) sproxmation science
	(v) Computer science.
	,
_3)	Demonstrate the Miller Rabin test for any large
	number & show the steps.
	To debernine if N=221 is prime
	$ n-1 = 220 = 2^2 \times 55$
	:. S=2 & d=55
	Select a random number such prat. 1LaKn-1
	(b) $a = 174$ $a^{2.d} \mod n = 174 \mod 221 = 47 \neq 1$, $n-1$ $a^{2.d} \mod n = 174 \mod 221 = 220 = n-1$
	2'-d mod n = 174 mod 221 = 220 = N-1
	Since 220 = n-1 mod n, elling 221 is prime or
	174 is a stong strong falsifier for 221,



MAEER's N

Consider a=137
a2°d mod n = 13755 mod 221 = 188 +1 N-1
a2'd moder = 137 10 mod 221 = 205 + 13, N-1
137 is a witness for compositeness of 221
2 174 was in fact a strong falsifier.