TCP - Part I

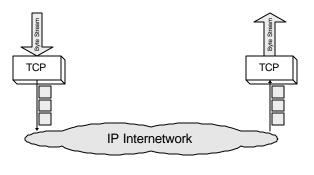
Relates to Lab 5. First module on TCP which covers packet format, data transfer, and connection management.

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Overview

TCP = Transmission Control Protocol

- Connection-oriented protocol
- Provides a reliable unicast end-to-end byte stream over an unreliable internetwork.



Connection-Oriented

- Before any data transfer, TCP establishes a **connection**:
 - One TCP entity is waiting for a connection ("server")
 - The other TCP entity ("client") contacts the server
- The actual procedure for setting up connections is more complex.
- Each connection is CLIENT full duplex

 Request a connection

 Accept a connection

 Data Transer

 Disconnect

 SERVER

 Waiting for connection request

Reliable

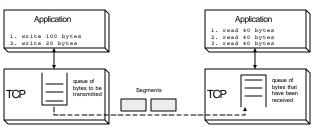
- Byte stream is broken up into chunks which are called segments
 - Receiver sends acknowledgements (ACKs) for segments
 - TCP maintains a timer. If an ACK is not received in time, the segment is retransmitted

•Detecting errors:

- TCP has checksums for header and data. Segments with invalid checksums are discarded
- Each byte that is transmitted has a sequence number

Byte Stream Service

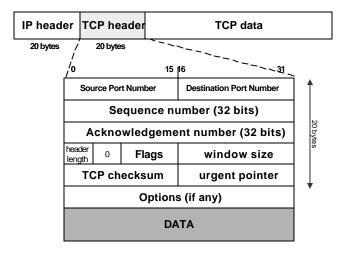
- To the lower layers, TCP handles data in blocks, the segments.
- To the higher layers TCP handles data as a sequence of bytes and does not identify boundaries between bytes
- So: Higher layers do not know about the beginning and end of segments!



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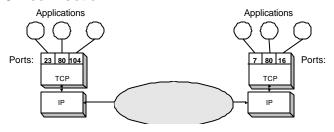
TCP Format

• TCP segments have a 20 byte header with >= 0 bytes of data.



Port Number:

- A port number identifies the endpoint of a connection.
- A pair <IP address, port number> identifies one endpoint of a connection.
- Two pairs <client IP address, server port number> and <server IP address, server port number> identify a TCP connection.



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TCP header fields

• Sequence Number (SeqNo):

- Sequence number is 32 bits long.
- So the range of SeqNo is

$$0 \le \text{SeqNo} \le 2^{32} - 1 \approx 4.3 \text{ Gbyte}$$

- Each sequence number identifies a byte in the byte stream
- Initial Sequence Number (ISN) of a connection is set during connection establishment

Q: What are possible requirements for ISN?

- Acknowledgement Number (AckNo):
 - Acknowledgements are piggybacked, I.e
 a segment from A -> B can contain an
 acknowledgement for a data sent in the B -> A direction
 Q: Why is piggybacking good?
 - A hosts uses the AckNo field to send acknowledgements.
 (If a host sends an AckNo in a segment it sets the "ACK flag")
 - The AckNo contains the next SeqNo that a hosts wants to receive

Example: The acknowledgement for a segment with sequence numbers 0-1500 is AckNo=1501

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TCP header fields

- Acknowledge Number (cont'd)
 - TCP uses the sliding window flow protocol (see CS 457) to regulate the flow of traffic from sender to receiver
 - TCP uses the following variation of sliding window:
 - no NACKs (Negative ACKnowledgement)
 - only cumulative ACKs
- Example:

Assume: Sender sends two segments with "1..1500" and "1501..3000", but receiver only gets the second segment.

In this case, the receiver cannot acknowledge the second packet. It can only send AckNo=1

- Header Length (4bits):
 - Length of header in 32-bit words
 - Note that TCP header has variable length (with minimum 20 bytes)

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TCP header fields

- Flag bits:
 - URG: Urgent pointer is valid
 - If the bit is set, the following bytes contain an urgent message in the range:

SeqNo <= urgent message <= SeqNo+urgent pointer

- ACK: Acknowledgement Number is valid
- PSH: PUSH Flag
 - Notification from sender to the receiver that the receiver should pass all data that it has to the application.
 - Normally set by sender when the sender's buffer is empty

Flag bits:

- RST: Reset the connection
 - The flag causes the receiver to reset the connection
 - Receiver of a RST terminates the connection and indicates higher layer application about the reset
- SYN: Synchronize sequence numbers
 - Sent in the first packet when initiating a connection
- FIN: Sender is finished with sending
 - Used for closing a connection
 - Both sides of a connection must send a FIN

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TCP header fields

Window Size:

- Each side of the connection advertises the window size
- Window size is the maximum number of bytes that a receiver can accept.
- Maximum window size is 2¹⁶-1= 65535 bytes

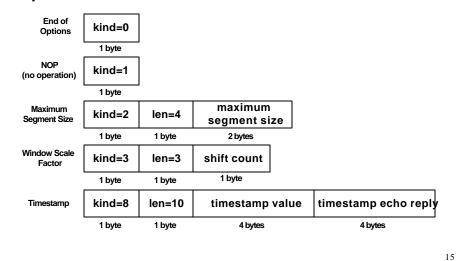
TCP Checksum:

 TCP checksum covers over both TCP header and TCP data (also covers some parts of the IP header)

• Urgent Pointer:

- Only valid if **URG** flag is set

• Options:



TCP header fields

- Options:
 - NOP is used to pad TCP header to multiples of 4 bytes
 - Maximum Segment Size
 - Window Scale Options
 - » Increases the TCP window from 16 to 32 bits, I.e., the window size is interpreted differently
 - Q: What is the different interpretation?
 - » This option can only be used in the SYN segment (first segment) during connection establishment time
 - Timestamp Option
 - » Can be used for roundtrip measurements

Connection Management in TCP

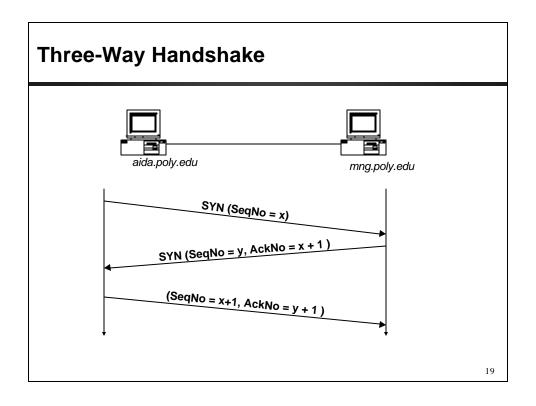
- Opening a TCP Connection
- Closing a TCP Connection
- Special Scenarios
- State Diagram

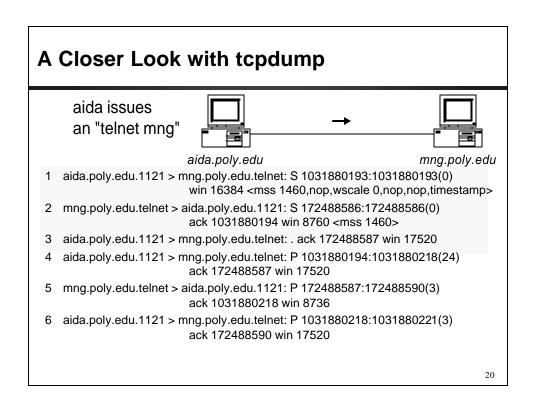
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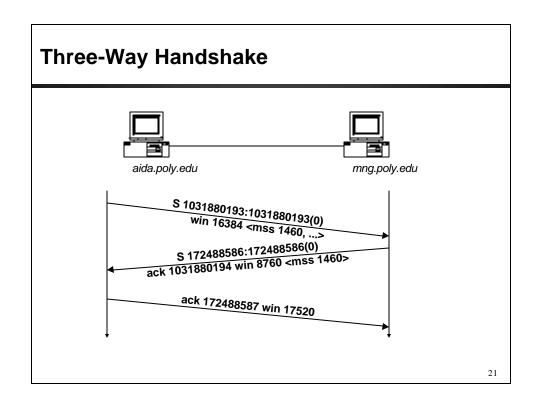
TCP Connection Establishment

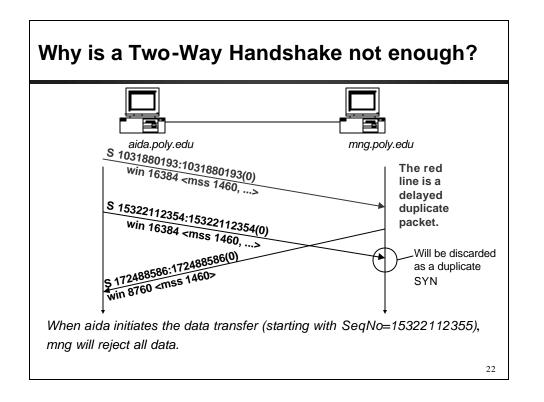
- TCP uses a three-way handshake to open a connection:
 - (1) ACTIVE OPEN: Client sends a segment with
 - SYN bit set *
 - port number of client
 - initial sequence number (ISN) of client
 - (2) PASSIVE OPEN: Server responds with a segment with
 - SYN bit set *
 - initial sequence number of server
 - ACK for ISN of client
 - (3) Client acknowledges by sending a segment with:
 - ACK ISN of server

(* counts as one byte)









TCP Connection Termination

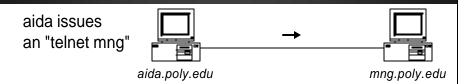
- Each end of the data flow must be shut down independently ("half-close")
- If one end is done it sends a FIN segment. This means that no more data will be sent
- Four steps involved:
 - (1) X sends a FIN to Y (active close)
 - (2) Y ACKs the FIN,

(at this time: Y can still send data to X)

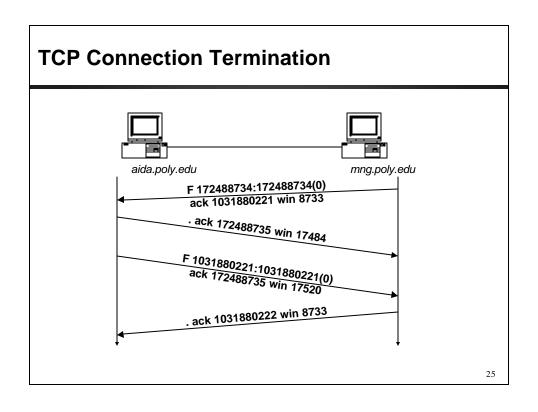
- (3) and Y sends a FIN to X (passive close)
- (4) X ACKs the FIN.

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Connection termination with tcpdump

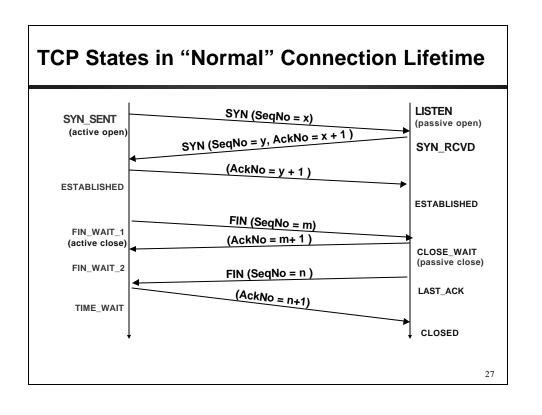


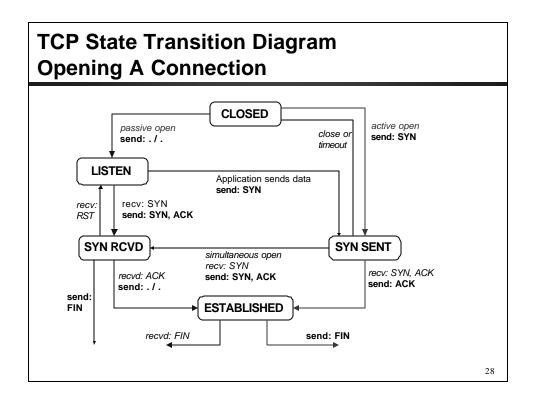
- 1 mng.poly.edu.telnet > aida.poly.edu.1121: F 172488734:172488734(0) ack 1031880221 win 8733
- 2 aida.poly.edu.1121 > mng.poly.edu.telnet: . ack 172488735 win 17484
- 3 aida.poly.edu.1121 > mng.poly.edu.telnet: F 1031880221:1031880221(0) ack 172488735 win 17520
- 4 mng.poly.edu.telnet > aida.poly.edu.1121: . ack 1031880222 win 8733



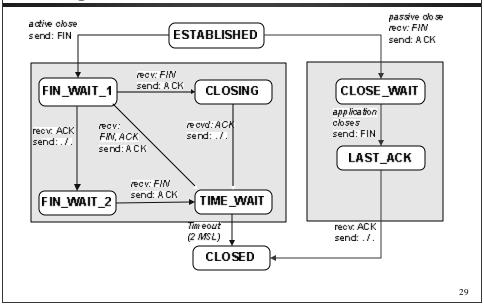
TCP States

State	Description
CLOSED	No connection is active or pending
LISTEN	The server is waiting for an incoming call
SYN RCVD	A connection request has arrived; wait for Ack
SYN SENT	The client has started to open a connection
ESTABLISHED	Normal data transfer state
FIN WAIT 1	Client has said it is finished
FIN WAIT 2	Server has agreed to release
TIMED WAIT	Wait for pending packets ("2MSL wait state")
CLOSING	Both Sides have tried to close simultanesously
CLOSE WAIT	Server has initiated a release
LAST ACK	Wait for pending packets





TCP State Transition Diagram Closing A Connection



2MSL Wait State

2MSL Wait State = TIME WAIT

 When TCP does an active close, and sends the final ACK, the connection must stay in in the TIME_WAIT state for twice the maximum segment lifetime.

2MSL= 2 * Maximum Segment Lifetime

- Why?
 TCP is given a chance to resent the final ACK. (Server will timeout after sending the FIN segment and resend the FIN)
- The MSL is set to 2 minutes or 1 minute or 30 seconds.

Resetting Connections

- Resetting connections is done by setting the RST flag
- When is the RST flag set?
 - Connection request arrives and no server process is waiting on the destination port
 - Abort (Terminate) a connection
 Causes the receiver to throw away buffered data. Receiver does not acknowledge the RST segment