

Func: print the last element in a LL.

What is the big-oh?  $O(n)$  - head pointer only

$O(1)$  - head & tail pointers

## Code for LL - DRAW A LOT OF PICTURES

Build a Singly-linked-list class to hold ints.

Node class



Node 2 pieces  $\rightarrow$  data  
 $\rightarrow$  pointer to the next item in the list.  
 reference

```
class Node {
```

```
    int data;
```

```
    Node next;
```

```
}
```

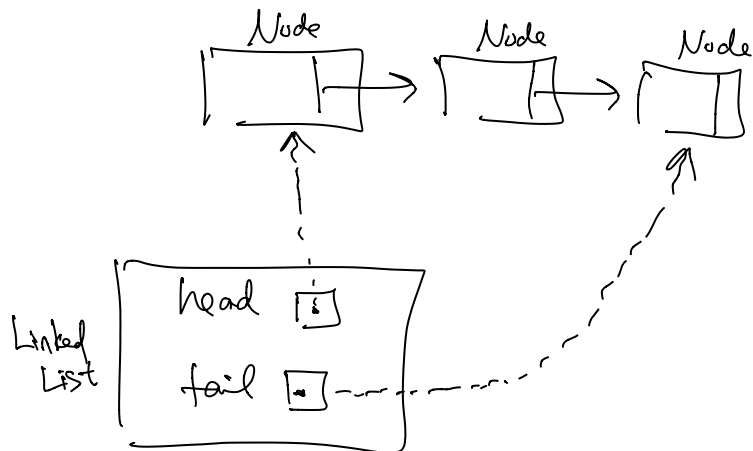
$\leftarrow$  this is not a Node "inside" a Node, it is a reference to the next Node.

```
class LinkedList {
```

```
    Node head;
```

```
    (Node tail;)
```

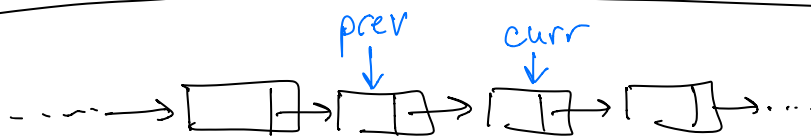
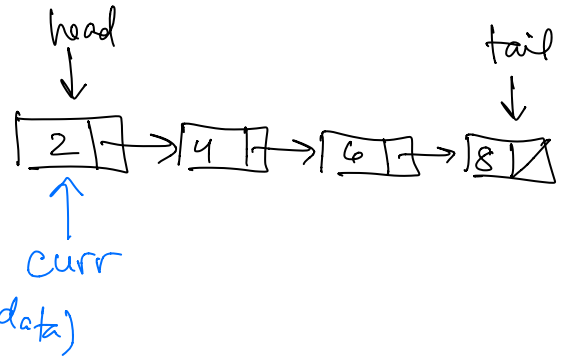
```
}
```



Traversals : going through the LL one element at a time. } 1-pointer traversal

Forward direction (left to right)

```
for(int x=0; x < arr.size(); x++)  
{  
    // Node curr = (wherever you want to start traversing) [normally this is the head]  
    while (stopping condition is not true (normally curr != null))  
    {  
        // processing step with curr.data  
        // example: System.out.println(curr.data)  
        curr = curr.next;  
    }  
}
```



2-pointer traversal

Node curr = head (or somewhere)

Node prev = null

while (curr != null)

{  
 // process curr.data

...

prev = curr;

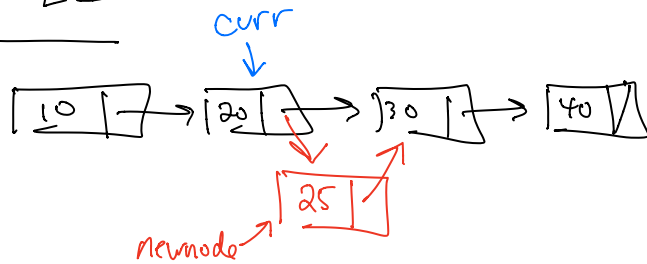
curr = curr.next;

}

~~prev = prev.next??~~

## INSERTING INTO A LL

inserting after a node



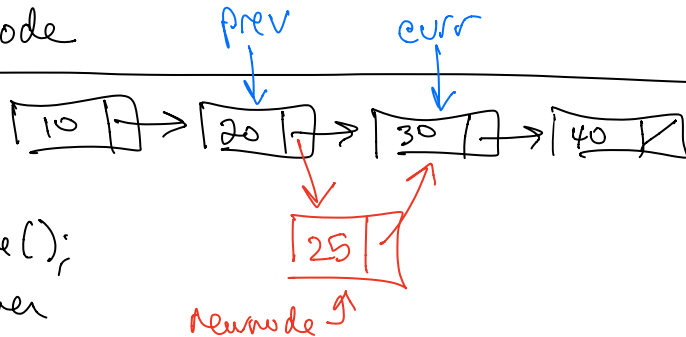
Node newnode = new Node();

newnode.data = whatever

newnode.next = curr.next

curr.next = newnode

inserting before a node



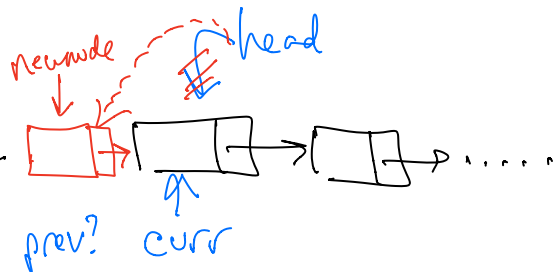
Node newnode = new Node();

newnode.data = whatever

newnode.next = curr

prev.next = newnode

insert @ the beginning of a LL

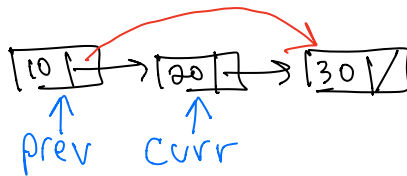


newnode.next = curr/head;

→ head = newnode;

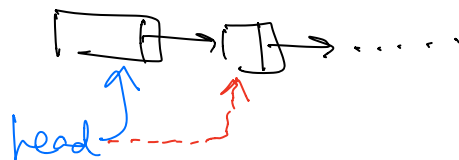
## Deletion

prev.next = curr.next



Deleting the head of a LL

head = head.next;



## Big-oh of insert/delete

Most of the time, <sup>\*</sup> insert/delete are  $O(1)$

\*  $\rightarrow$  Takes time to traverse the list.

• SLL w/ only a head pointer

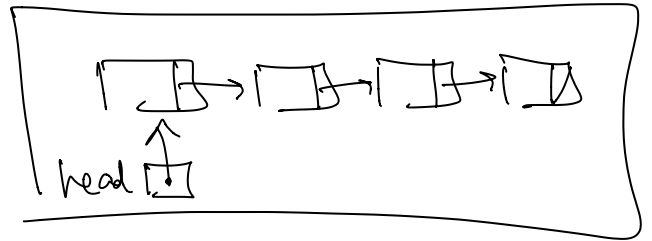
insert @ beginning =  $O(1)$

insert @ end =  $O(n)$

delete @ beginning =  $O(1)$

delete @ end =  $O(n)$

in general: insert/delete =  $O(1) + O(\text{traversal})$



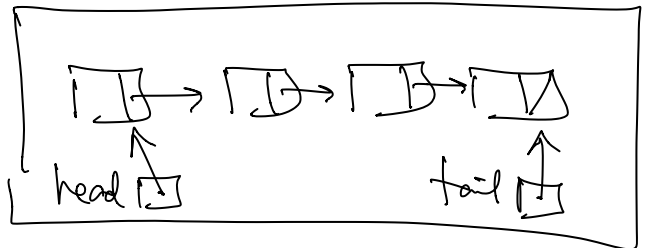
• SLL w/ both head + tail pointers

insert @ beg -  $O(1)$

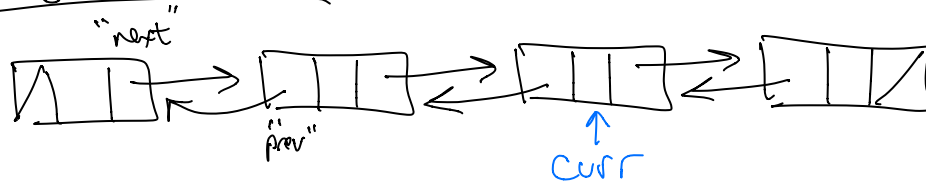
" @ end -  $O(1)$

delete @ beg -  $O(1)$

delete @ end -  $O(n)$

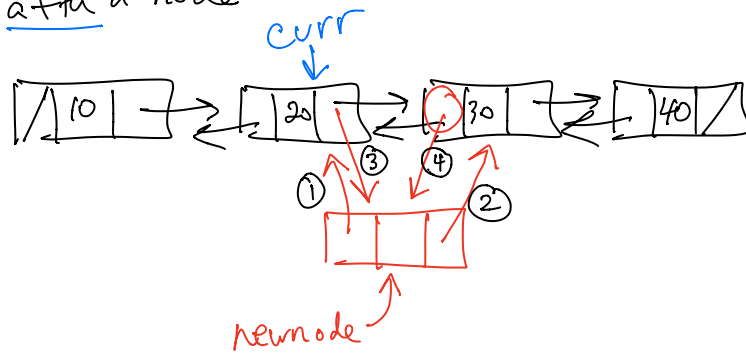


## Doubly-linked-list



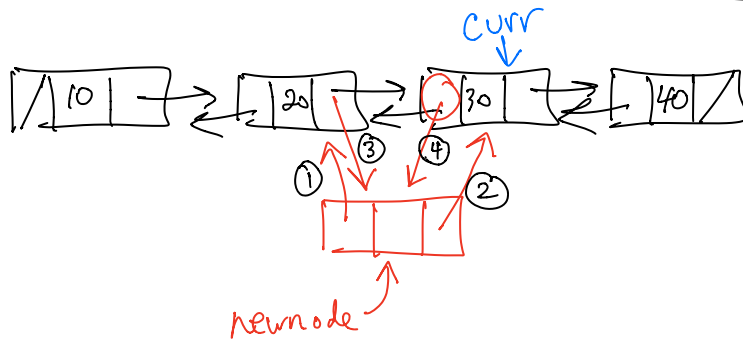
```
class Node {  
    int data;  
    Node next;  
    Node prev;  
}
```

Inserting after a node



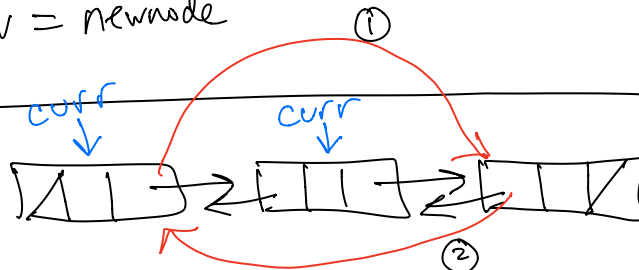
- ①  $\text{newnode.prev} = \text{curr}$
- ②  $\text{newnode.next} = \text{curr.next}$
- ③  $\text{curr.next} = \text{newnode}$
- ④  $\text{curr.next.prev} = \text{newnode}$

Inserting  
before a  
node



- ①  $\text{newnode.prev} = \text{curr.prev}$
- ②  $\text{newnode.next} = \text{curr}$
- ③  $\text{curr.prev.next} = \text{newnode}$
- ④  $\text{curr.prev} = \text{newnode}$

Deletion



- ①  $\text{curr.prev.next} = \text{curr.next}$
- ②  $\text{curr.next.prev} = \text{curr.prev}$

Deletion @ head or tail

Defect that a head of time & do special cases.

if (curr.prev == null) // curr == head

if (curr.next == null) / if (curr == tail)

Big-oh

insert/deleting @ head/tail -  $O(1)$

" " in middle -  $O(1) + O(\text{traversal})$