



Lambda Calculus Part II



Overview

- Computing with Lambda
- How to represent
 - Booleans, Naturals
- Combinators
 - I
 - K
 - S
- Recursion
 - Y



Lambda Calculus

- As a programming language
 - We must be able to express
 - Constants
 - Booleans
 - Naturals
 - Elementary operations
 - Successor
 - Addition
 - Multiplication
 - We must be able to deal with recursion



Booleans

- How to represent true & false.
 - Hint
 - Think about them as functions
 - What are boolean used for ?
- Branching



Naturals

- What we need
 - A function to represent ZERO
 - A function to get the successor of a natural n
 - A function to get the predecessor of a natural n
 - A function to test for ZERO
- Church numerals
 - Zero: $\lambda s. \lambda z. z$
 - One: $\lambda s. \lambda z. s \ z$
 - Two: $\lambda s. \lambda z. s \ (s \ z)$
 - Three: $\lambda s. \lambda z. s \ (s \ (s \ z))$



Successor

- How to define successor ?

- Natural $n = \lambda s. \lambda z. s \underbrace{(s (s \dots (s z)) \dots)}$

n of them

- Natural $(n+1) = \lambda s. \lambda z. s \underbrace{(s (s (s \dots (s z)) \dots))}$

$n+1$ of them

- Successor



IsZero

- Testing for Zero
 - Input
 - A Natural: n
 - Output
 - True if n is zero
 - False in all other cases otherwise
 - Idea
 - Make an induction on the church numeral
 - For zero $\lambda s.\lambda z.z$ return true. How ?
 - For others: return false, always.



Pairing & Lists

- How to define a Pair ?
 - A function
 - Takes two inputs
 - Returns something that denotes a pair (a function!)
 - Objective
 - We should be able to ask a pair for
 - Its first element
 - Its second element
 - Suggestions ?



Predecessor

- How to define $n-1$ from n ?
 - Question
 - Can we eliminate applications within a numeral directly ?



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 - Form a function that computes $\langle n, n+1 \rangle$ from $\langle n-1, n \rangle$
 - Apply the function n times starting with $\langle ?, \text{zero} \rangle$
 - What is the final pair ?



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```
Pair    =  $\lambda x. \lambda y. \lambda z. z \ x \ y$ 
First   =  $\lambda p. p \ \text{true}$ 
Second  =  $\lambda p. p \ \text{false}$ 
np      =  $\lambda p. \text{Pair} \ (\text{Second } p) \ (\text{succ} \ (\text{Second } p))$ 
pred    =  $\lambda n. \text{First} \ (n \ np \ (\text{Pair} \ \text{zero} \ \text{zero}))$ 
```



Addition

- Objective
 - Compute $m+n$ from naturals m, n
- Idea ?



Multiplication

- Objective
 - Compute $m * n$ from naturals m, n
- Idea ?



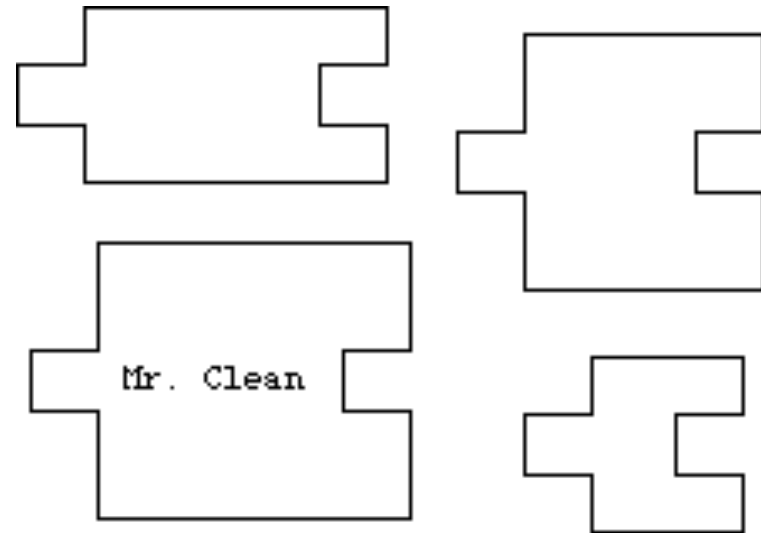
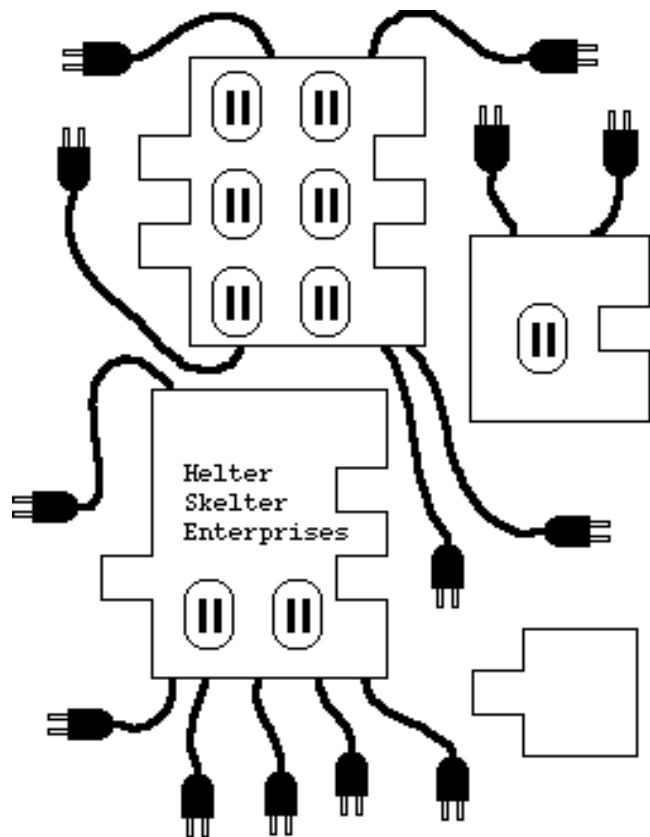
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Combinators

- What are these things?





Combinator I

- **Purpose**
 - Provide the Identity transformation
 - Take a lambda expression as input
 - Produce the lambda expression unchanged
 - Definition

$$I = \lambda x. x$$



Combinator K

- **Purpose**
 - Create a constant *function*
 - Take as input an expression
 - Produce as output a function that returns the expression
 - Definition

$$K = \lambda x . \lambda y . x$$



Combinator S

- **Purpose**
 - Provide a general form of composition.
 - Definition

$$S = \lambda x . \lambda y . \lambda z . x \ z \ (y \ z)$$



Expressiveness

- What does SKK simplify to ?

I	=	$\lambda x. x$
K	=	$\lambda x. \lambda y. x$
S	=	$\lambda x. \lambda y. \lambda z. x \ z \ (y \ z)$



Recursion

- Is achieved with one more combinator
 - Y
- Topic of next lecture is exclusively recursion