Network Security (NetSec)



Summer 2015

Chapter 01: Fundamentals

Module 04: Reconnaissance



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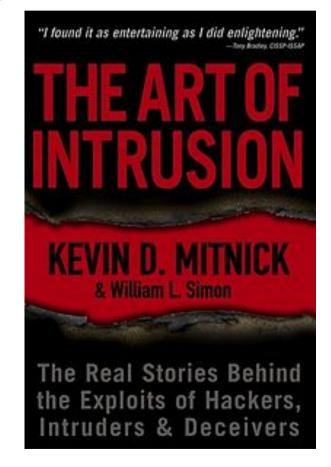
Learning Objectives



Some practical insights into reconnaissance (to prepare attacks on networks)

- Know basic means on how to identify and map networking infrastructure
- Identify which pieces of information are public and do hence not provide any stronghold/defense (in a sense of security by obscurity)
- Develop awareness of the attacker's moves; know the preparatory steps potential attackers are going to perform

To probe further







Take an attacker's perspective before thinking about defenses!

Here: Reconnaissance





Overview of this Module



- (1) "Casing the joint"
- (2) The Robin Hood hacker
- (3) Reconnaissance with public sources of information
- (4) Reconnaissance with active scanning/probing (in part in Appendix)
- (5) Recommended readings

Chapter 01, Module 04





"Casing the Joint" - Definitions



"Casing the Joint" – informal for reconnoitering before a robbery

Random House Dictionary: re-con-nais-sance

- 1.the act of reconnoitering.
- 2.Military. a search made for useful military information in the field, esp. by examining the ground.
- 3.Surveying, Civil Engineering. a general examination or survey of a region, usually followed by a detailed survey.
- 4.Geology. an examination or survey of the general geological characteristics of a region.

Wordnet: reconnaissance

• the act of reconnoitring (especially to gain information about an enemy or potential enemy); "an exchange of fire occurred on a reconnaissance mission"











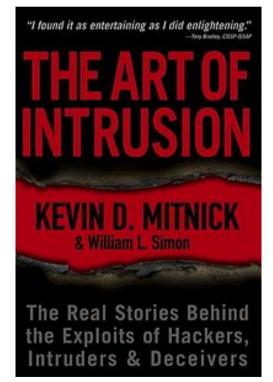
Who of you has heard about the "New York Times Hack"? Who of you has heard about the "Robin Hood Hacker" behind it?

Here comes part of the story of Adrian Lamo ...

- It gives some insights on what creative minds will do to hack a network
- ... and might teach you something about trust in humans.

"[Hacking] has always been for me less about technology and more about religion"

- Adrian Lamo



Source: Ch.5 from the above book





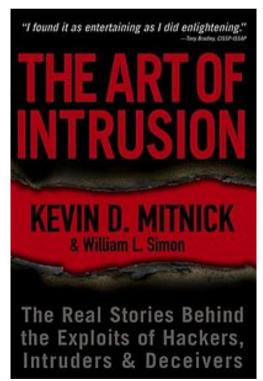


Some background

- The young Adrian Lamo is called Robin Hood hacker, since he does not hide his identity and he informs the owners of the hacked networks about the problems – he hacks out of curiosity, but for good
- Adrian hacked into Microsoft, Yahoo, MCI WorldCom, Excite@Home, Cingular, and other telcos

Hacking Excite@home

- Curiosity ... a network managing all cable customers in the country surely is well protected
- From an open network in a students lounge, he accessed the webpage, finds a misconfigured proxy that gives access to one internal webserver



Source: Ch.5 from the above book







Hacking Excite@home contd.

- He posts message on this internal server on having problems to log in
- The response contains a URL pointing to the system for managing IT problems, this URL gives clues about other (similar) systems in the network that handle support related questions
- There is no authentication for the support systems they are designed for "internal" access (anyone calling the system must have access since this is privileged info – security by obscurity)
- Netcraft.com to find out OS of server systems
- Network operation center offers helpdesk system ... with plenty of details on customer account data, password, on handling trouble tickets, script for generating authentication cookies (allowing helpdesk members to login as the customer who is calling)
- He calls a guy with a particular strange ticket that has never been resolved by Excite@home support and is about stolen credit card details







Hacking Excite@home contd.

- Open proxy stops working
- He performs reverse DNS lookups and probes IP addresses, ends up finding dialup00.corp.home.net
- Dialup-users run Win98, have open network shares, etc.
- Changes startup script so remote users run software for Adrian
- Installs remote desktop management software (virus checkers do not alarm, since this is legitimate software)
- Learns what internal systems the dialup users access using netstat
- On the discovered servers, he finds
 - Database of > 3 million subscribers with details such as OS the users use, cable modem serial numbers, names, email addresses
 - Adrian nows a lead network engineer at Excite@home, tells him all details (meeting at 4:30 AM in the morning)
 - Finally wants to see proxy server in data-center ... is asked: "how would you secure this server", cuts ethernet wire with pocket knife ... engineer says "that's good enough", attaches note "do not re-attach"







Hacking NY-Times

- Curiosity: have already access to washington post, wouldn't t be nice to be in the NYT network
- Website hosted with company, i.e. outside the network
- Using whois (arin.net) to get some basic server info
- Port scan of addresses belonging to NYT
- Various open proxies that allowed him in, but not of direct use
- Email of all NYT staffers from external website (that agrred on being reachable via the Internet) ... sends mail and gets back answer
- IP addresses in mail headers as a starting point ... manual scan of IP addresses, but only internal we servers without much info
- Finds old NYT Intranet site, decommissioned, but with link to production system ... with material teaching staff on how to operate the system
- Search engine on this machine that allows free-form SQL queries







Hacking NY-Times (contd.)

- System used Lotus Nodes, full database access here meant all information of all employees, newsstand owners, etc ... down to social security number, how much they earn, any complaints about service of NYT, etc.
- Yet Adrian does not know, which OS the system was running ... "I don't analyze networks this way", "its about the people and how they configure the network. Most people are very predictable"
- Internal search engine in fact indexed the entire site, SQL query tool on server ... by probing he finds database names, etc. containing
 - Complete username-password list at NYT
 - List of all persons held on terrorist charge in US
 - List of all op-ed (opposite the editorial page) writers, containing info of celebs such as robert redford
 - Adds his own number and contact detail as "computer hacking/ security and communicatoins intelligence"







Hacking NY-Times (contd.)

- **-** ...
- Discovers that he can use LexisNexis database system with access to legal and news information ... is claimed to have performed 3000 searches ... is undiscovered for some 3 months
- He decides to tell a friend (reporter) with details under agreement to not publish before he informs the NYT and told them how to fix problem
- NYT has no interest to learn this, has to be pressed to learn about the vulnerability ... fixes it in the following 48 hours ... call the FBI, after they learn that Adrian used LexisNexis ... Adrian hides for five days, then "surrenders" at StarBucks
- Is sued with damages in the 300.000 USD range (calculated using the price of pay as you go searches with 12USD each, which would mean 270 queries per day for 3 months) ...
- 65.000 USD charges + 6 months home confinement
- Btw. Is not fluent in any programming language ...





Robin Hood turning Evil?



And now for something completely different

- The very same Adrian Lamo is also known as the person who turned in Bradley Manning with the report that Manning had leaked thousands of sensitive U.S. government documents, thus starting the wikileaks scandal ... apparently, Manning trusted Lamo after having chatted with him extensively ...
- If you are interested to learn more, the Internet is full of Information on this case ... good starting points are wired.com (with a pro Lamo coverage) or for opposing views wikileaks.org and salon.com (condemning wired.com) ...









Some Basic Reconnaissance



Reconnaissance based on Publicly Available Information



To be reachable, need to register domain names and to allow for resolving them (to allow for mapping "names" into "addresses")

Here we are taking a closer look at

- Reconnaissance with whois
- Reconnaissance with DNS

A few words about a Registrar:

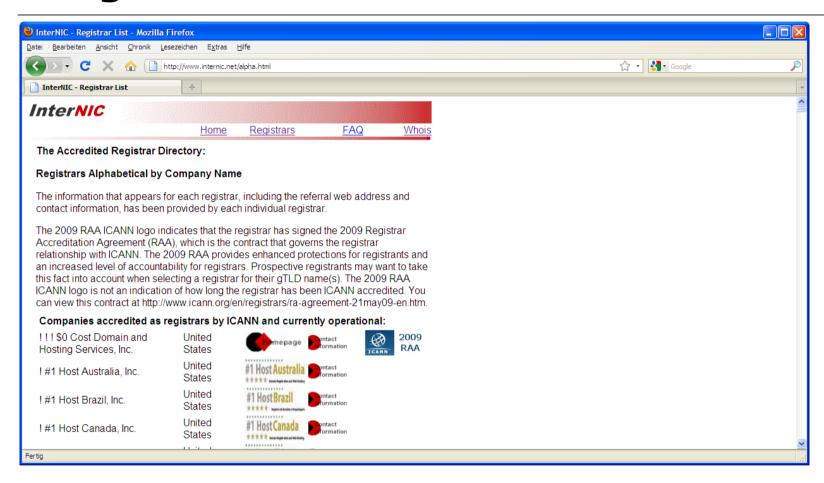
- Organization where you register a domain name
- Verifies uniqueness of name
- Enters domain name into various databases: whois & DNS
- List of registrars can be obtained from internic.net





Querying Databases: Target "tu-darmstadt.de"

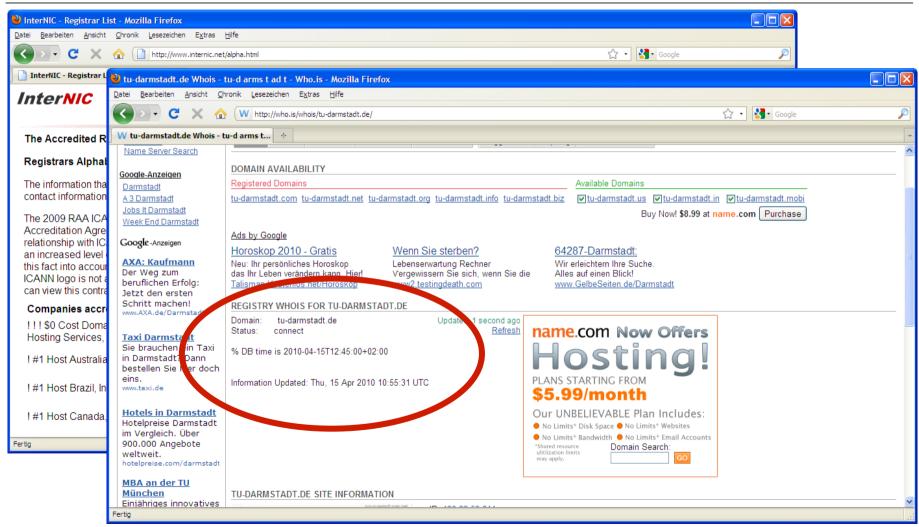






Querying Databases: Target "tu-darmstadt.de"







Whois Databases



Two steps

First find target's registrar, then whois target at registrar

Input: domain name or company name

Output: registrar, whois server, dns server, names of people (administrator, billing contact), phone numbers, E-mail addresses

Some useful whois sites:

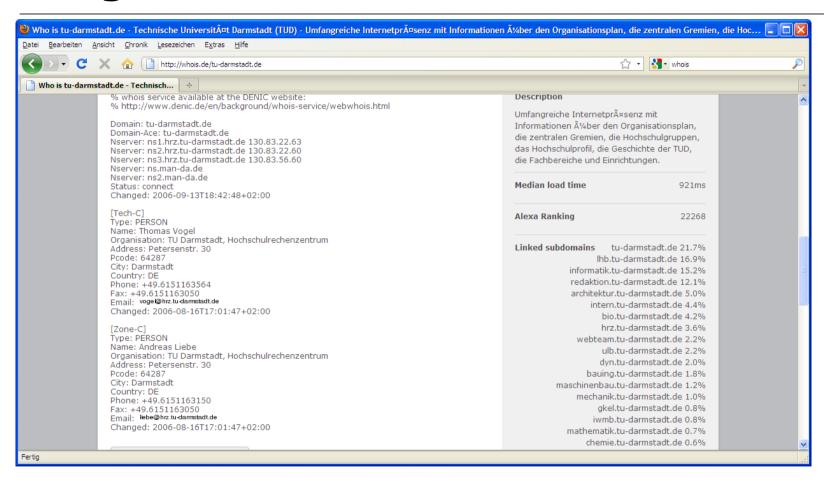
- www.internic.net
 - For com, net and org top-level domains
- whois.net, who.is, www.networksolutions.com/whois
 - For country-code top-level domains or arbitrary domains, e.g., jp, fr





Querying Databases: Target "tu-darmstadt.de"

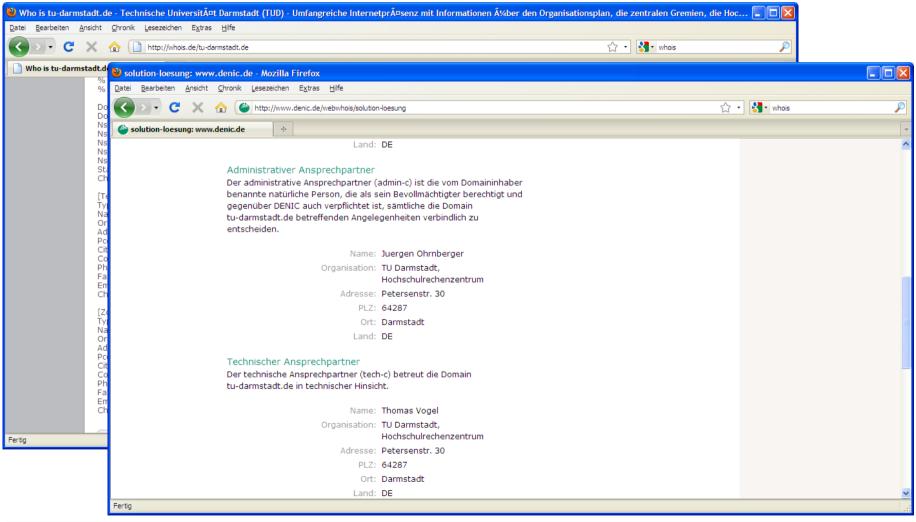






Querying Databases: Target "tu-darmstadt.de"







Reconnaissance: IP Ranges



IANA: Internet Assigned Numbers

Authority

ARIN: American Registry for Internet Numbers

Maintains whois database that includes IP address ranges in US

RIPE: Europe

APNIC: Asia

[Source: http://xkcd.com/195/]

THIS CHART SHOWS THE IP ADDRESS SPACE ON A PLANE USING A FRACTAL MAPPING WHICH PRESERVES GROUPING -- ANY CONSECUTIVE STRING OF IPS WILL TRANSLATE TO A SINGLE COMPACT, CONTIGUOUS REGION ON THE MAP. EACH OF THE 256 NUMBERED BLOCKS REPRESENTS ONE /8 SUBNET (CONTAINING ALL IPS THAT START WITH THAT NUMBER). THE UPPER LEFT SECTION SHOWS THE BLOCKS SOLD DIRECTLY TO CORPORATIONS AND GOVERNMENTS IN THE 1990'S BEFORE THE RIRS TOOK OVER ALLOCATION.

14 15 16 19 -> 13 12 17 18





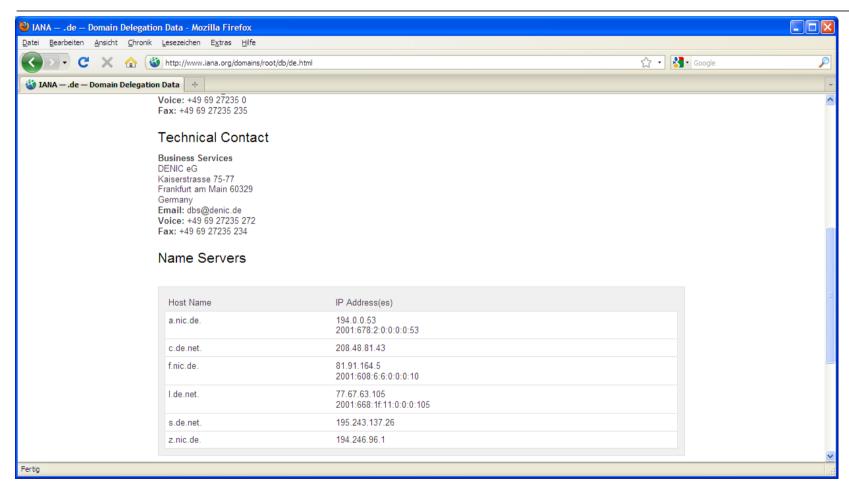
THE IPV4 SPACE, 2006 HP DEC FORD CSC PON-XEROX BELL APPLE MIT EUROPE ASIA-PACIFIC CARIBBEAN VARIOUS http://www.iana.org/ AMERICA assignments/ ipv4-address-space/ ipv4-address-space.xml

MAP OF THE INTERNET



IANA Databases Query at RIPE.NET

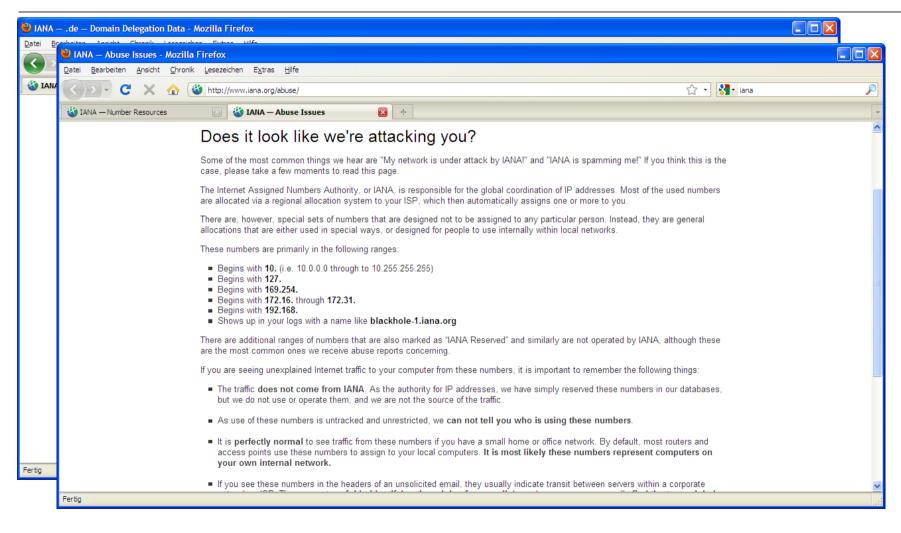






IANA Databases Query at RIPE.NET

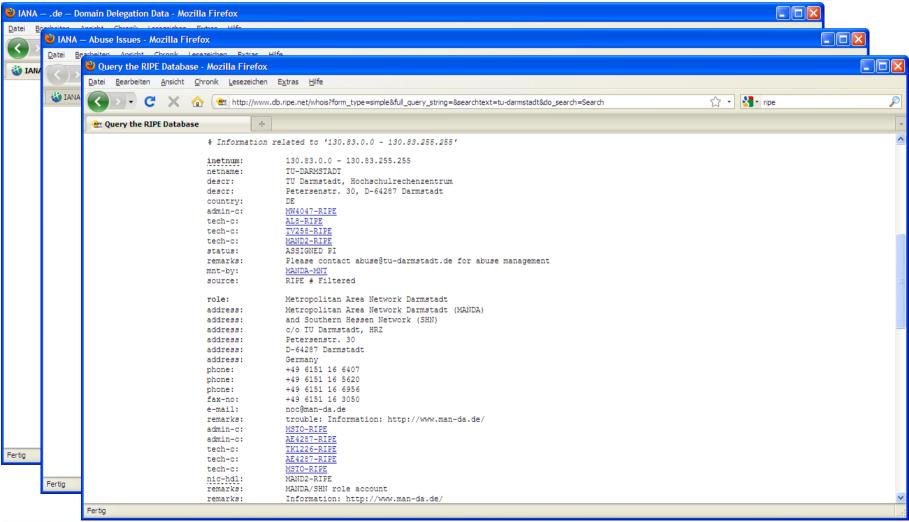






IANA Databases Query at RIPE.NET







Why are Whois Databases Publicly Available?



Troubleshooting is one reason, security another:

- If you're under attack, can analyze source address of packets
- Can use whois database to obtain info about the domain from where the attack is coming
- Can inform admin that their systems are source of an attack



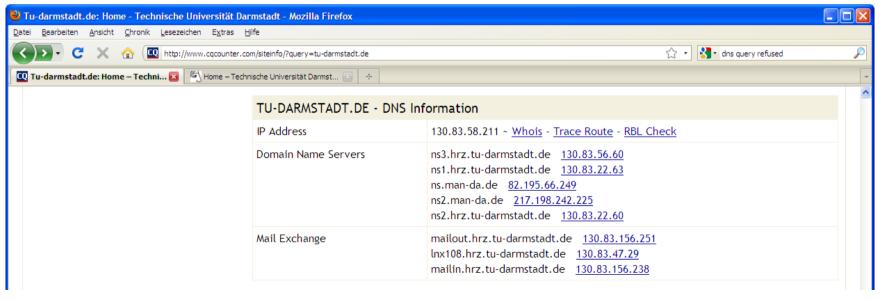


Reconnaissance: DNS Database



Let's quickly review DNS

- Distributed database implemented in hierarchy of many DNS servers
 Authoritative name server
- For a given domain (e.g., tu-darmstadt.de), provides server name to IP address mappings for servers (web, email, ftp, etc) in domain
 Primary and secondary name server for reliability

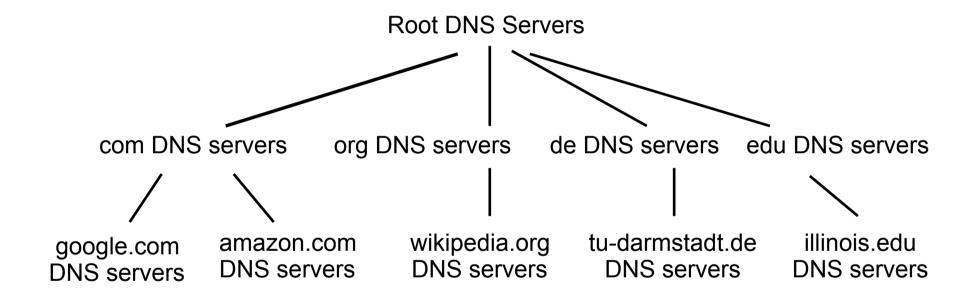






DNS Hierarchy

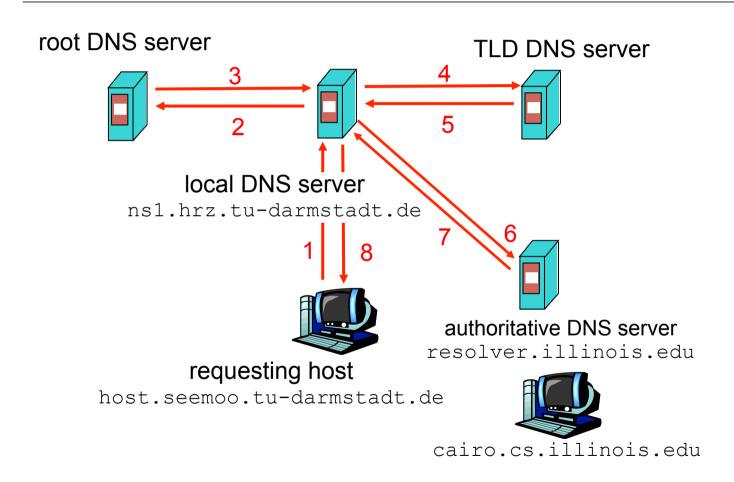






DNS: Queries







DNS Records



DNS: distributed db storing resource records (RR)

RR format: (name, value, type, ttl)

Which RRs do you know?

Type=NS

- name is domain (e.g. foo.com)
- value is IP address of authoritative name server for this domain

Type=A

- name is hostname
- value is IP address

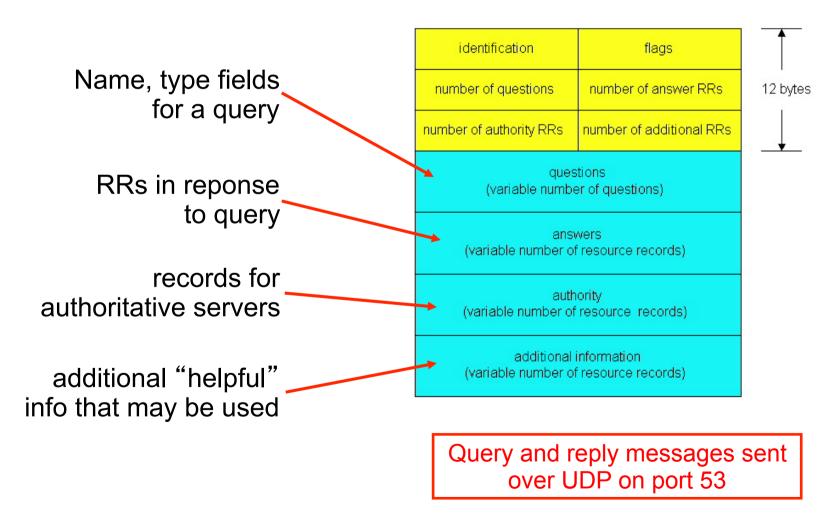
Type=MX

value is name of mailserver associated with name



DNS Protocol, Messages







DNS Contd.



DNS: Caching and Updating Records

- Once (any) DNS server learns mapping, it caches mapping
 - Cache entries timeout (disappear) after some time
 - Cache poisoning is well-known attack

Interrogating DNS Servers

- Attacker first gets primary or secondary authoritative server for target organization using whois
- Attacker can then query the DNS by sending DNS query messages.
- Tools (often available in Unix and Windows machines; also available at web sites):
 - nslookup
 - host
 - dig (domain information groper)





nslookup



Available in most unix & Windows **Machines**

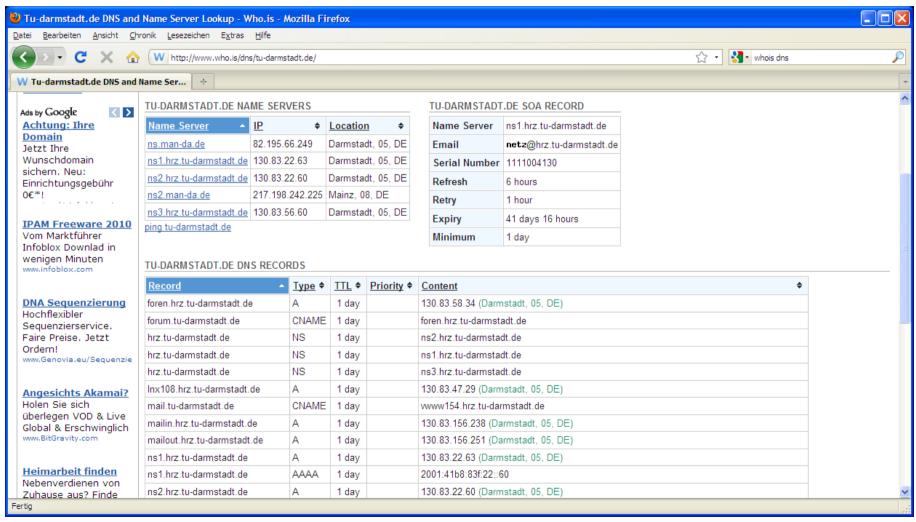
- Get DNS server IP address from whois
- set type=any "get all" ...
- however, today best practice is to disable

```
_ 🗆 ×
                                                             Command Prompt - nslookup
                                                                                                                                                                                                                              _ 🗆 ×
                                                             C:\Documents and Set
                                                                                                 C:\WINDOWS\system32\cmd.exe - nslookup
                                                            Default Server: ns1
Address: 207.172.3.
                                                                                                 C:\Dokumente und Einstellungen\Lecture>nslookup
                                                                                                Standardserver: arcor.easybox
Address: 192.168.2.1
                                                                ls dialpad
                                                            Insl.dns.rcn.netl > server 130.83.22.63

*** Can't list domai Standardserver: ns1.hrz.tu-darmstadt.de
> server 192.5.6.34 Address: 130.83.22.63
                                                          www.tu-darmstadt.de canonical name = vwcms-live01.hrz.tu-darmstadt.de
tu-darmstadt.de nameserver = ns3.hrz.tu-darmstadt.de
                                                                                               tu-darmstadt.de nameserver = ns.man-da.de
tu-darmstadt.de nameserver = ns2.man-da.de
tu-darmstadt.de nameserver = ns2.hrz.tu-darmstadt.de
tu-darmstadt.de nameserver = ns1.hrz.tu-darmstadt.de
ns.man-da.de internet address = 82.195.66.249
ns.man-da.de internet address = 2001:41b8:0:1::53
ns1.hrz.tu-darmstadt.de internet address = 130.83.22.63
ns1.hrz.tu-darmstadt.de AAAA IPv6 address = 2001:41b8:83f:22::60
ns2.hrz.tu-darmstadt.de internet address = 130.83.22.60
ns2.hrz.tu-darmstadt.de AAAA IPv6 address = 2001:41b8:83f:22::63
ns2.man-da.de internet address = 217.198.242.225
ns3.hrz.tu-darmstadt.de internet address = 130.83.56.60
ns3.hrz.tu-darmstadt.de AAAA IPv6 address = 2001:41b8:83f:56::60
> vwcms-live01.hrz.tu-darmstadt.de
                                                                                                tu-darmstadt.de nameserver = ns.man-da.de
                                                              dialpad.com.
                                                              dialpad.com.
                                                              dialpad.com.
dialpad.com.
dialpad.com.
                                                              televoip
                                                              wram
                                                               kosova
                                                              vietnam
                                                              asianet
                                                                                                > vwcms-liveO1.hrz.tu-darmstadt.de
Server: ns1.hrz.tu-darmstadt.de
Address: 130.83.22.63
                                                               reseller
                                                               now
                                                                                                                                                                  internet address = 130.83.58.211
HINFO CPU = PC OS = Linux
MX preference = 10, mail exchanger = vwc
                                                                                                vwcms-live01.hrz.tu-darmstadt.de
                                                              www.cn
                                                                                                vwcms-live01.hrz.tu-darmstadt.de
vwcms-live01.hrz.tu-darmstadt.de
                                                              espano1
                                                              www.espanol
                                                                                                ms-live01.hrz.tu-darmstadt.de
vwcms-live01.hrz.tu-darmstadt.de
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                                                              www.pk
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                                                                                                hrz.tu-darmstadt.de
                                                                                                                                        nameserver = ns1.hrz.tu-darmstadt.de
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                                                                                                                                        nameserver = ns2.hrz.tu-darmstadt.de
                                                                                                                                        nameserver = ns3.hrz.tu-darmstadt.de
                                                                                                hrz.tu-darmstadt.de
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nslookup via web







Reconnaissance via Google



Google sees lots of content that is not properly protected

- Sample query:
 - "intitle: Cisco Systems, Inc. VPN 3000 Concentrator"
 - Online devices that might be open to anyone, if not properly configured
 - "allinurl:auth_user_file.txt"
 - Search Operators
 - site: search for references to the specified site
 - link: find sites containing search term as a link
 - cache: display the cached version of pages found
 - intitle: find sites containing terms in the title of a page
 - inurl: find sites containing terms in the URL of a page
 - filetype: search specific document type
 - Look for webcams, printer, etc. maybe outdated security patches
- Source:

http://www.au.af.mil/au/awc/awcgate/nist/ljutic_fissea_22mar05.pdf





Reconnaissance Intermediate Summary



So far: Obtaining information from public databases:

- whois databases
 - Tool: web sites
- DNS database & server OS for web servers
 - Tool: nslookup, web sites (netcraft.com), Google

Defense

Keep to a minimum what you put in the public database: only what is necessary

Whats up next: active scans

- Network mapping
- Port scanning (e.g. using nmap)
- Sniffing (e.g. using wireshark)

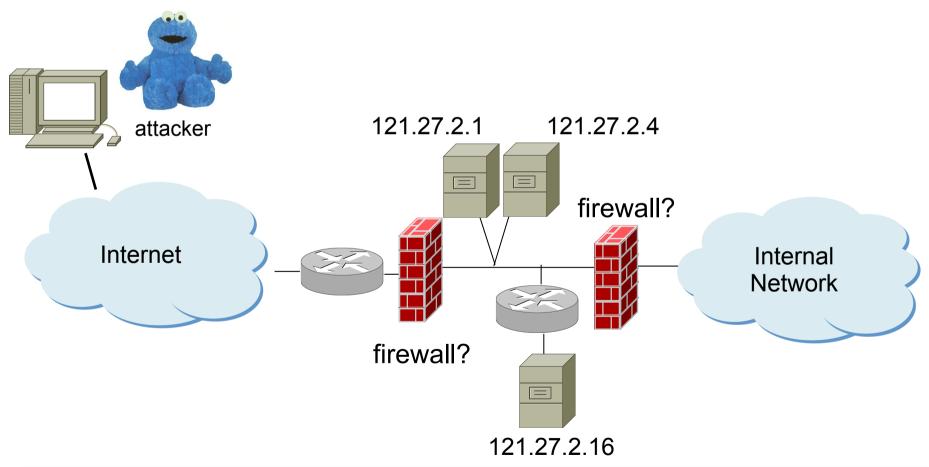




Network Mapping Tools



Goal: Learn about a remote network





Network Mapping



Attacker uses ping sweeps to determine live hosts

Attacker uses port scans to determine live services

Attacker often uses traceroute/pathping to determine path to each host discovered during ping sweep

- Traceroute: How it works
 - Source sends UDP packets to target
 - Each to an unlikely port
 - 3 packets with the same TTL, then increments TTL
 - When router decrements TTL to 0, sends back to source ICMP packet
 - type 11, code 0, TTL expired
 - When target receives packet, sends back to source ICMP packet
 - type 3, code 0, destination port unreachable
- Overlay results from traceroute to create an approximate network diagram





Ping Sweep



Ping

- Recall ICMP messages are directly encapsulated in IP datagrams (protocol 1)
- To ping a host:
 - Send ICMP Echo Request (ICMP type 8)
 - Host responds with ICMP Echo Reply (type 0)
- So let's ping the entire IP address range
 - Use automated tool for this ping sweep
- If firewall blocks ping packets:
 - Try sweeping with TCP SYN packets to port 80
 - Or try sending UDP packets to possible ports





Port Scanning



Now that we have a map with some hosts, let's find out what ports are open on a target host

■ 65,535 TCP ports; 65,535 UDP ports

Web server: TCP port 80

DNS server: UDP port 53

Mail server: TCP port 25

Port scanning tools can scan:

List of ports

Range of ports

• All possible TCP and UDP ports

Attacker may scan a limited set of ports, to avoid detection





Traceroute/Pathping



traceroute: gaia.cs.umass.edu to www.eurecom.fr

Three delay measements from gaia.cs.umass.edu to cs-gw.cs.umass.edu 1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms 2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms 3 cht-vbns.gw.umasš.edu (128.119.3.130) 6 ms 5 ms 5 ms 4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms 5 jn1-so7-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms 6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms 7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms trans-oceanic 8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms link 9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms 10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms 11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms 12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms 13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms 14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms 15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms 16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms means no reponse (probe lost, router not replying) 19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms



Nmap (1) – see also Appendix



Extremely popular

- usually run over linux
- rich feature set, exploiting raw sockets
- need root to use all features

Ping sweeping

- over any range of IP addresses
- with ICMP, SYN, ACK
- OS determination

Port scanning

- Over any range of ports
- Almost any type of TCP, UDP packet

Source IP address spoofing

Decoy scanning

Packet fragmentation

Timing Options

nmap MAN page details use and is excellent source of information





Nmap (2)



Input

nmap [Scan Type] [Options] <target hosts>

 Default for port scanning: ports 1-1024 plus ports listed in nmap service file

Output

- open ports: syn/ack returned; port is open
- unfiltered ports: RST returned: port is closed but not blocked by firewall
- filtered ports: nothing returned; port is blocked by firewall

See Appendix for further examples





Network Mapping Defenses



Filter using firewalls and packet-filtering capabilities of routers

- Block incoming ICMP packets, except to the hosts that you want to be pingable
- Filter Time Exceeded ICMP messages leaving your network

Close all unused ports

Scan your own systems to verify that unneeded ports are closed Intrusion Detection Systems

But be aware: makes network troubleshooting also harder



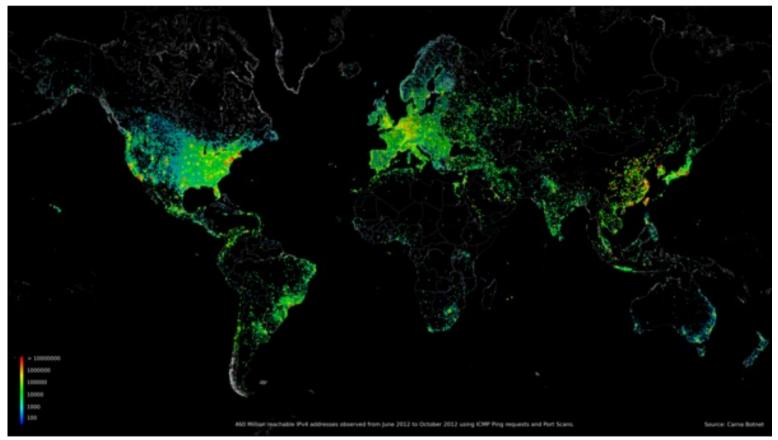


A Cool Project to Map the Internet



Have you heard of the Carna botnet?

Recommended: http://internetcensus2012.bitbucket.org/paper.html







Summary





Acks & Recommended Reading



Selected slides of this chapter courtesy of

Keith Ross

Recommended reading

- Manuals of popular unix tools
 - ping
 - traceroute
 - Nmap

Additional reading

- There are various online resources on "ethical hacking"
- Always a good starting point is the ccc (chaos computer club) ... but do not expect to find "handbooks for hacking" at their site





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Contact







Appendix: Nmap Examples



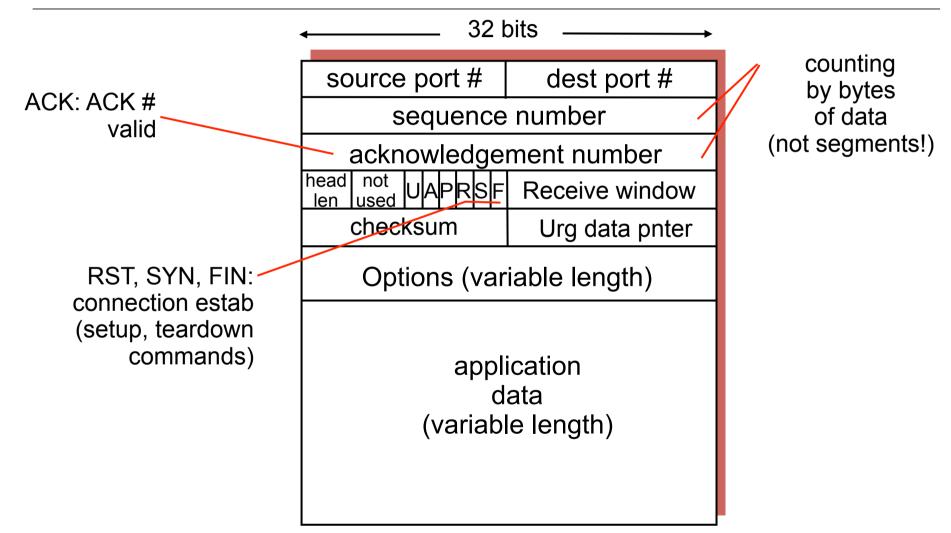
Some details on TCP
Some nmap examples

Please note: Appendices are in general not relevant for the exam



Excursus: TCP Segment Structure







Excursus: TCP seq. #'s and ACKs

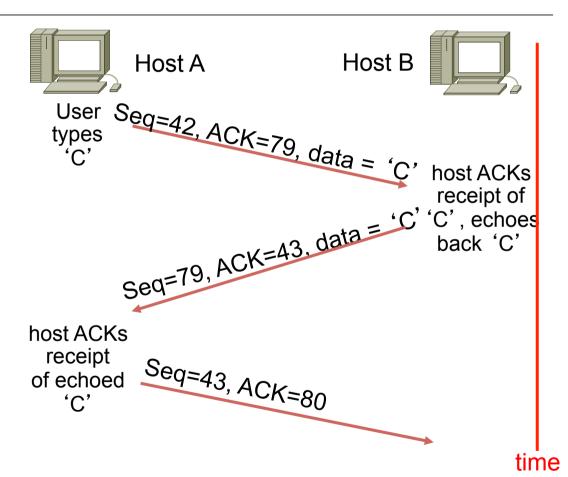


Seq. #'s:

byte stream "number" of first byte in segment's data

ACKs:

seq # of next byte expected from other side



simple telnet scenario





Excursus: TCP Connection Establishment



Three way handshake:

Step 1: client host sends TCP SYN segment to server

- SYN=1, ACK=0
- specifies initial seq #
- no data

Step 2: server host receives SYN, replies with SYN-ACK segment

- SYN=1, ACK=1
- server host allocates buffers
- specifies server initial seq. #

Step 3: client receives SYN-ACK, replies with ACK segment, which may contain data

■ SYN=0, ACK=1





TCP: Reset packet



If machine receives a TCP packet it is not expecting, it responds with TCP packet with RST bit set.

• For example when no process is listening on destination port For UDP, machine returns ICMP "port unreachable" instead



Nmap (3): ping sweep



Sends ICMP echo request (ping) to 256 addresses Can change options so that pings with SYNs, ACKs...

$$-sP = ping$$



Nmap (4): polite port scan



nmap -sT -v target.com

Attempts to complete 3-way handshake with each target port Sends SYN, waits for SYNACK, sends ACK, then sends FIN to close connection

If target port is closed, no SYNACK returned

Instead RST packet is typically returned

TCP connect scans are easy to detect

- Target (e.g. Web server) may log completed connections
- Gives away attacker's IP address





Nmap (5): TCP SYN port scan



nmap -sS -v target.com

Stealthier than polite scan Send SYN, receive SYNACK, send RST

Send RST segment to avoid an accidental DoS attack

Stealthier: hosts do not record connection

But routers with logging enabled will record the SYN packet

Faster: don't need to send FIN packet





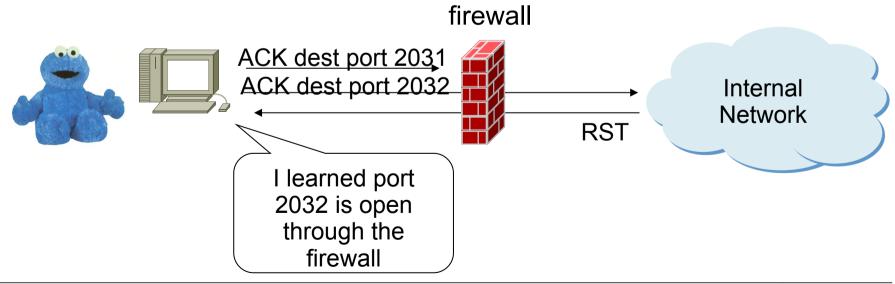
Nmap (6): TCP ACK scans



Many filters (in firewalls and routers) only let internal systems hosts initiate TCP connections

 Drop packets for which ACK=0 (ie SYN packet): no sessions initiated externally

To learn what ports are open through firewall, try an ACK scan (segments with ACK=1)





Nmap (7): UDP port scans



UDP doesn't have SYN, ACK, RST packets

nmap simply sends UDP packet to target port

- ICMP Port Unreachable: interpret port closed
- Nothing comes back: interpret port open
 - False positives common





Nmap (8): Obscure Source



Attacker can enter list of decoy source IP addresses into Nmap For each packet it sends, Nmap also sends packets from decoy source IP addresses

For 4 decoy sources, send five packets

Attacker's actual address must appear in at least one packet, to get a result

If there are 30 decoys, victim network will have to investigate 31 different sources!





Nmap (9): TCP Stack Fingerprinting



In addition to determining open ports, attacker wants to know OS on targeted machine:

- exploit machine's known vulnerabilities
- sophisticated hacker may set up lab environment similar to target network

TCP implementations in different OSes respond differently to (illegal) combinations of TCP flag bits





Nmap (10): Fingerprinting



Nmap sends

- SYN to open port
- NULL to open port (no flag bits set)
- SYN/FIN/URG/PSH to open port
- SYN to closed port
- ACK to closed port
- FIN/PSH/URG to closed port
- UDP to closed port

Nmap includes a database of OS fingerprints for hundreds of platforms

See nmap.org for further details





Nmap (11): examples



nmap -v target.com

Scans all TCP default ports on target.com; verbose mode

 First pings addresses in target network to find hosts that are up. Then scans default ports at these hosts; stealth mode (doesn't complete the connections); tries to determine OS running on each scanned host

■ Sends an Xmas tree scan to the first half of each of the 255 possible subnets in the 198.116/16. Testing whether the systems run ssh, DNS, pop3, or imap

• finds all web servers on machines with IP addresses ending in .2.3, . 2.4, or .2.5



