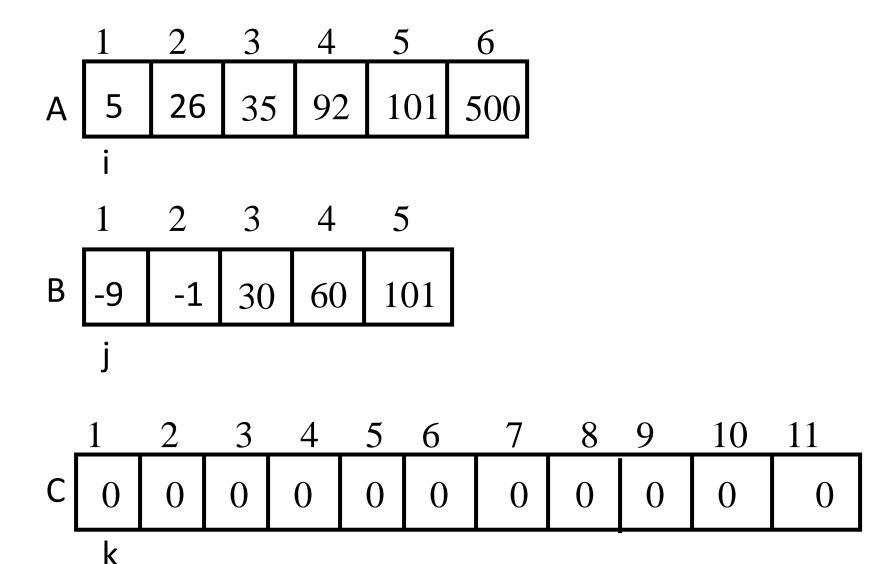
Joy Mukherjee

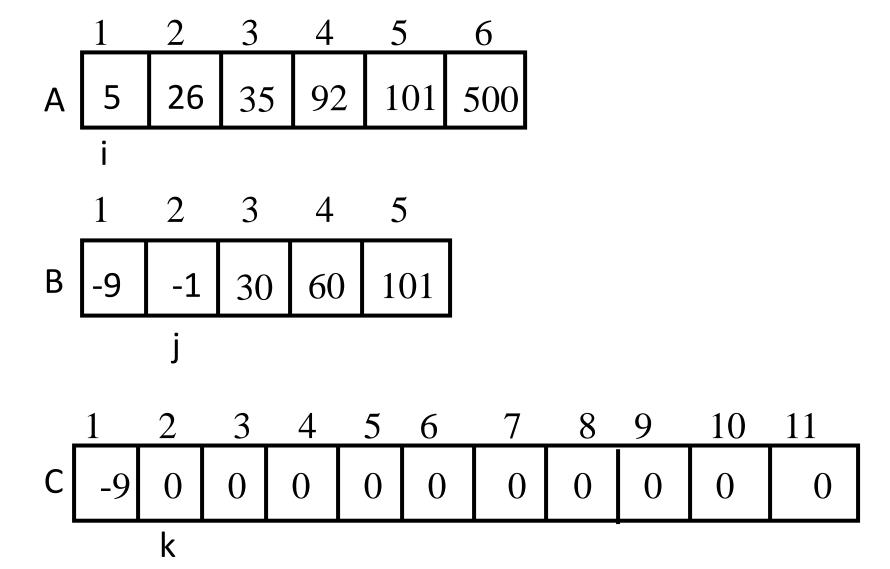
Sorting

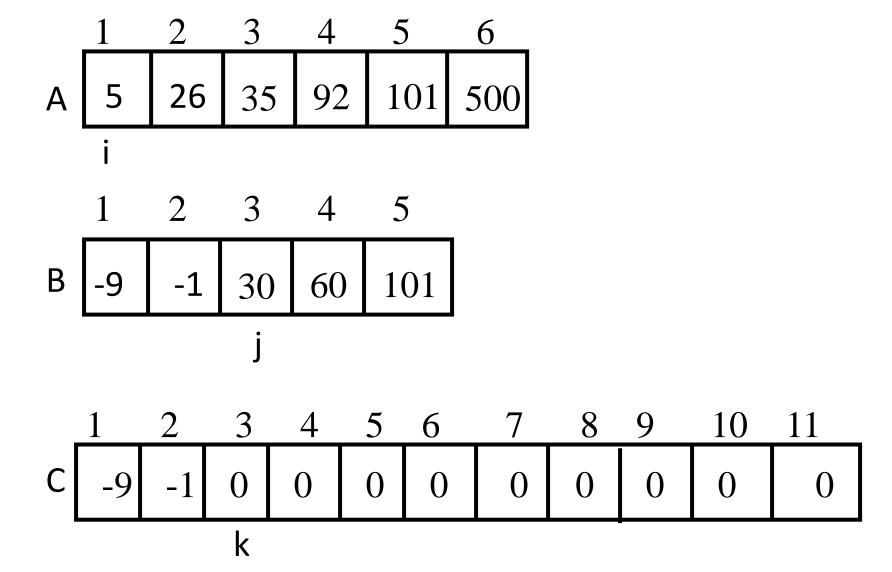
- Given an array A[] of n integers, our objective is to arrange them in non-decreasing order/nonincreasing order.
- Bubble Sort/Selection Sort/Insertion Sort
- Merge Sort/Quick Sort/Heap Sort
- Counting Sort/Radix Sort
- Bucket Sort

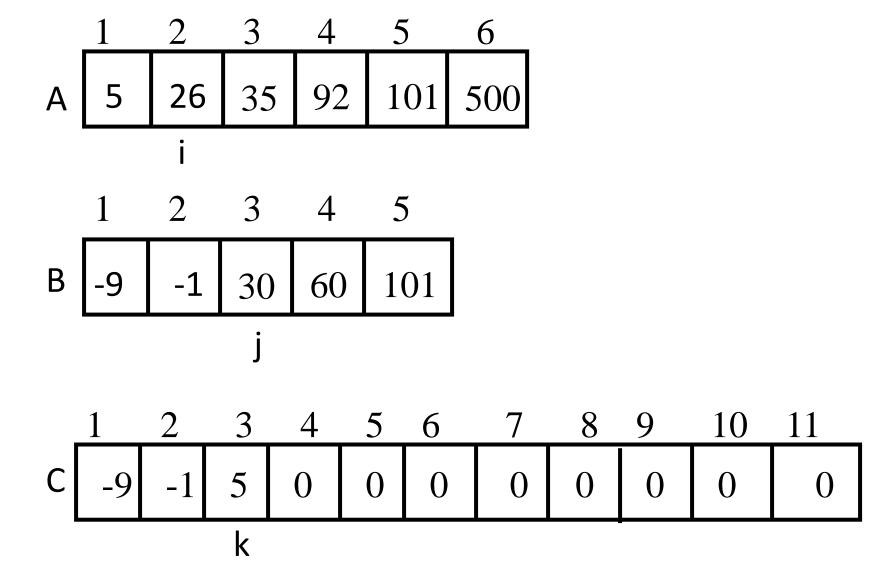
 Given two sorted arrays of length m and n, how do we merge them into a single sorted array of size m+n?

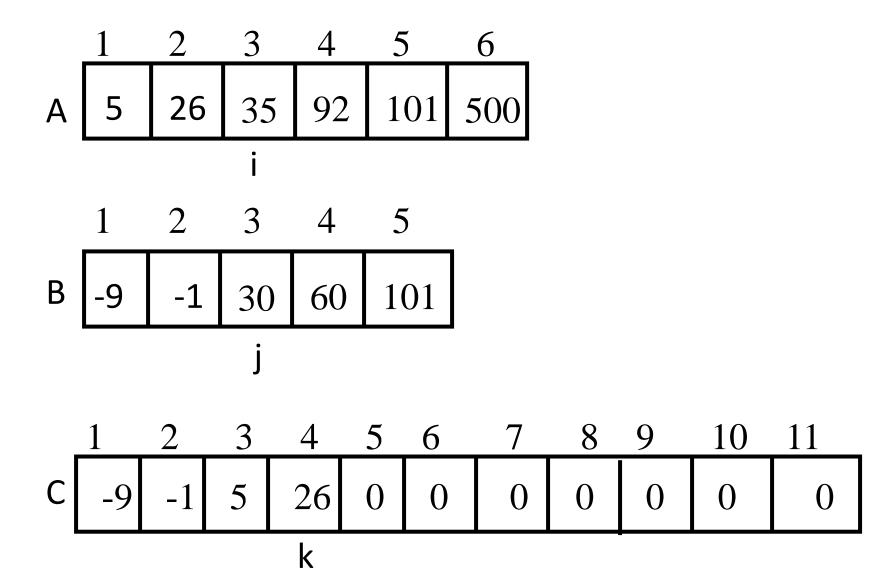
 We use an extra array of size m+n to merge the two sorted arrays.

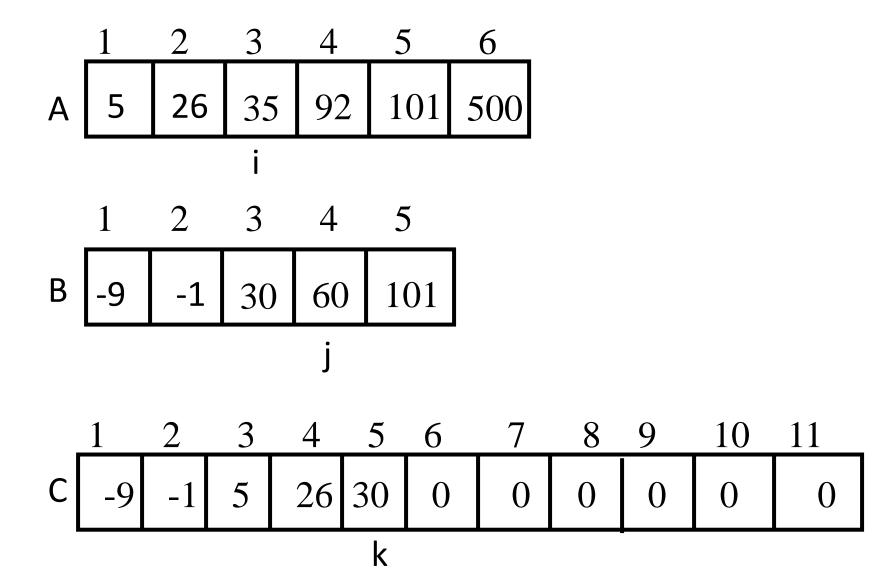


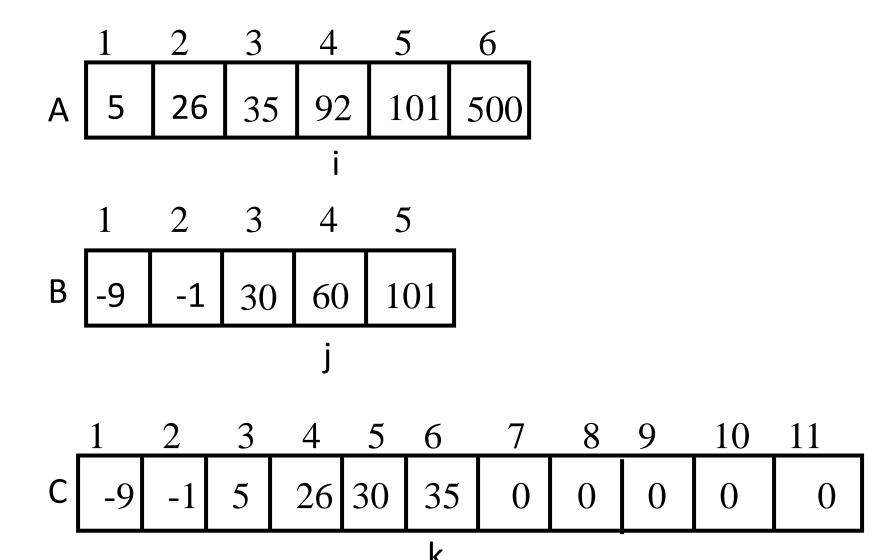


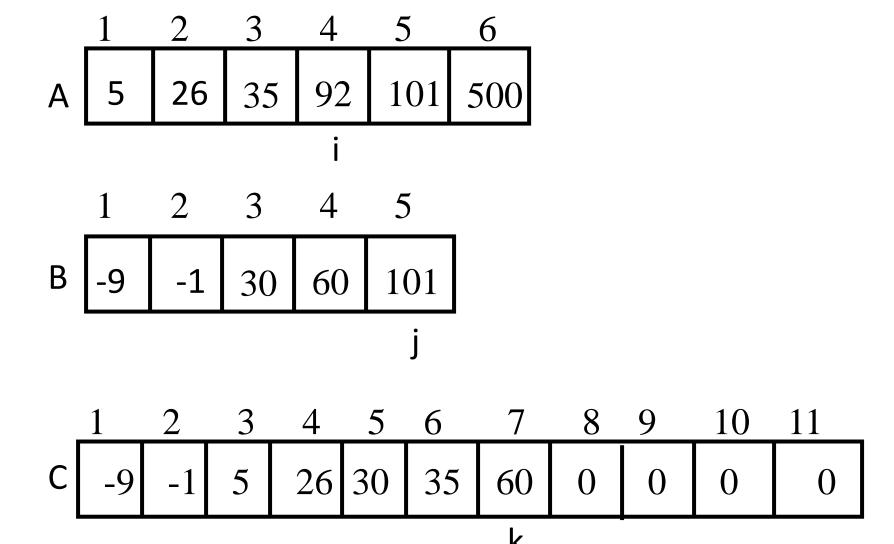


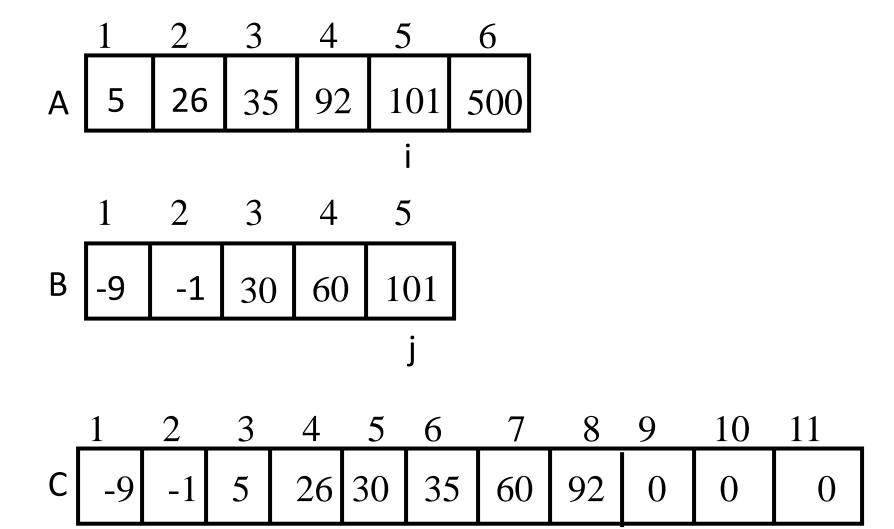




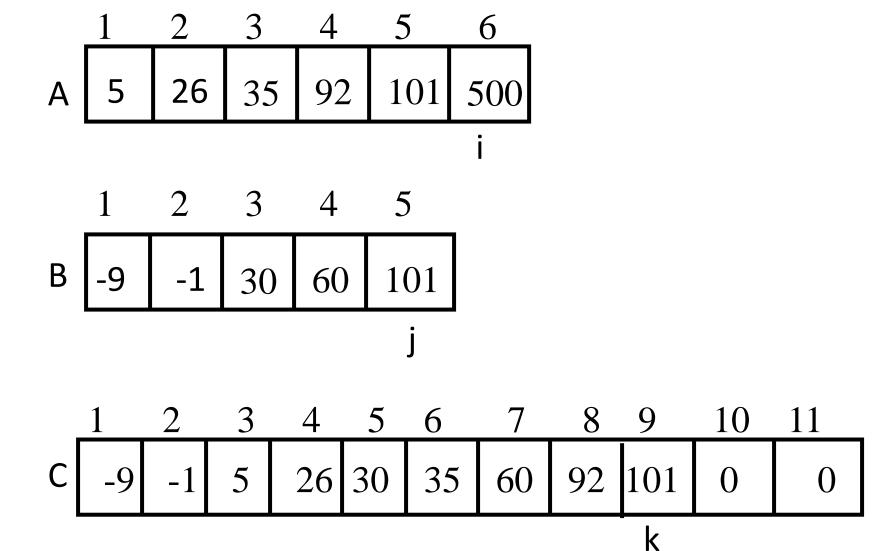


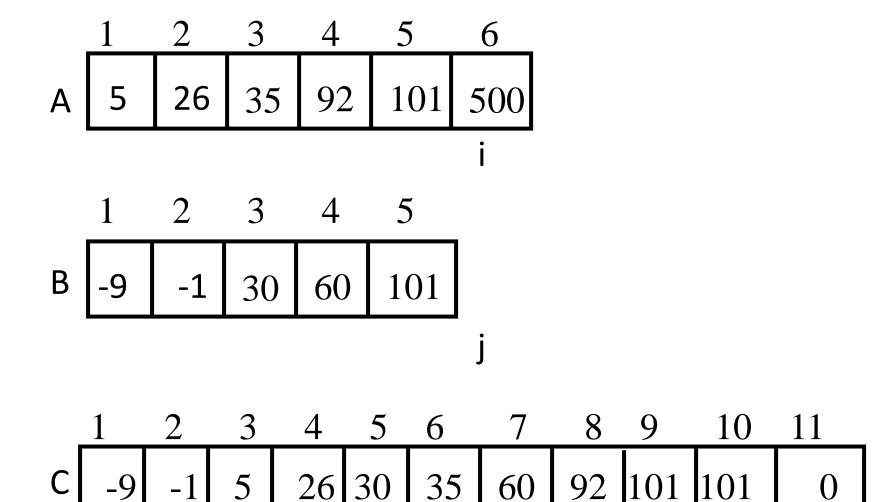




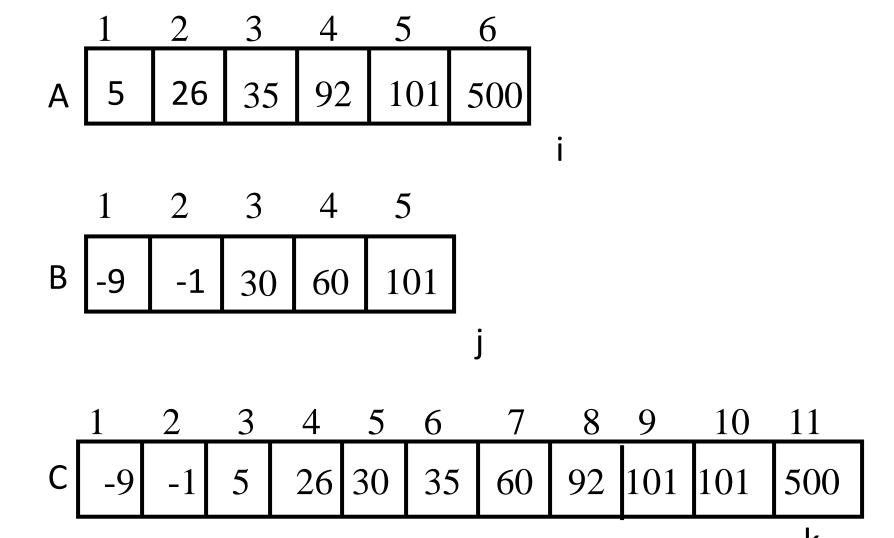


k



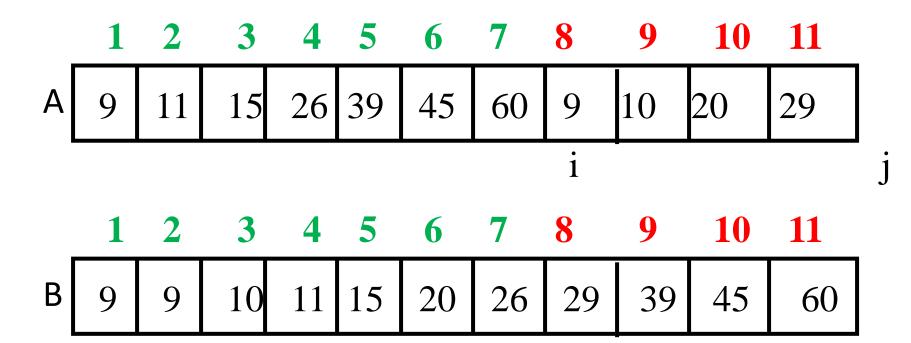


k



Merge

- A [low..high]
- A[low..mid] and A[mid+1..high]



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Merge

```
void merge(int a[], int low, int mid,
                                             a[low] to a[mid] = A
 int high)
                                             a[mid+1] to a[high] = B
                                             b[0] to b[high-low] = C
    int i = low, j = mid+1, k = 0;
    int b[high-low+1];
    while(i <= mid && j <= high) { // Both
    array are not exhausted
      if(a[i] \leq a[i])
            \{b[k] = a[i]; i++; k++;\}
      else
            \{b[k] = a[j]; j++; k++;\}
```

Merge

```
while(j != high+1) // a[mid+1] to a[high] is not exhausted
  b[k++] = a[j++];
while(i != mid+1) // a[low] to a[mid] is not exhausted
  b[k++] = a[i++];

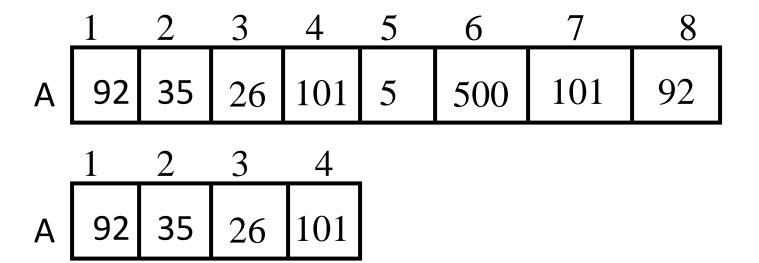
for(i = low, j = 0; i <= high; i++, j++)
  a[i] = b[j]; // Copy sorted b[0 .. High-low] to a[low to high]</pre>
```

- It is a recursive function.
 - ➤ Divide the array into almost equal two halves
 - ➤ Recursively sort left half A[low..mid]
 - ➤ Recursively sort right half A[mid+1..high]
 - ➤ Merge the two halves to get the array sorted A[low..high]

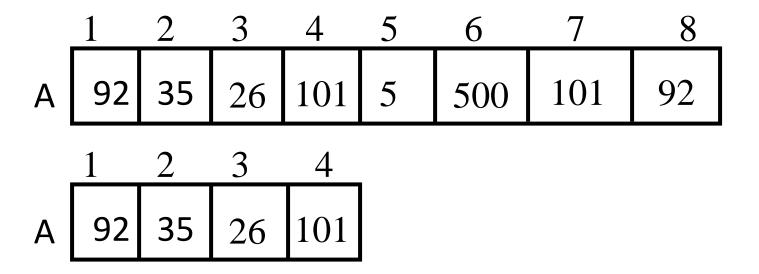
```
void mergesort(int a[], int low, int high)
 if(low < high) {
     int mid = (high+low)/2;
     mergesort(a, low, mid); // Sort the left half
     mergesort(a, mid+1, high); // Sort the right half
     merge(a, low, mid, high); // Merge the sorted halves
From main(): mergesort(a, 0, n-1);
```

	1	2	3	4	5	6	7	8
Α	92	35	26	101	5	500	101	92

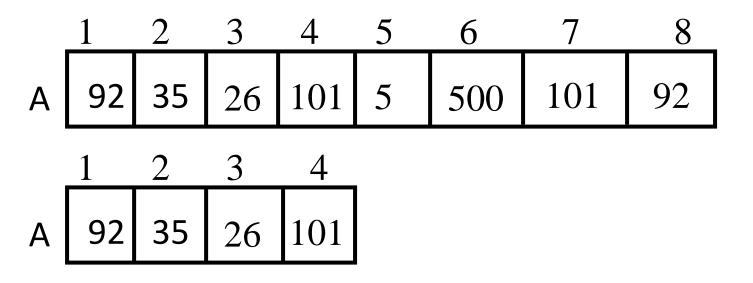
Merge Sort: Divide (Level 1)



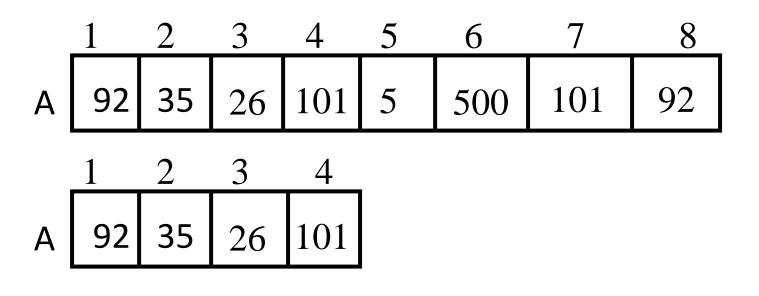
Merge Sort: Divide (Level 2)



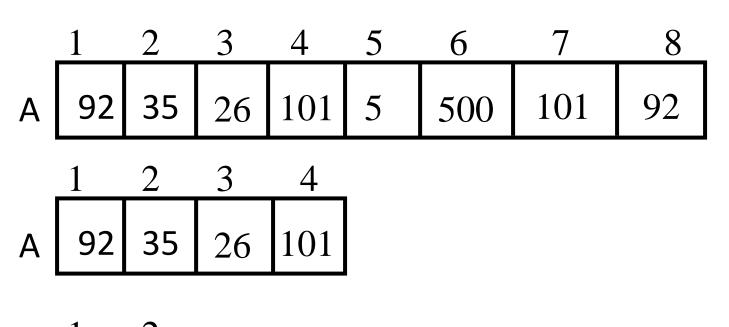
Merge Sort : Divide (Level 3) & Conquer (Level 4)



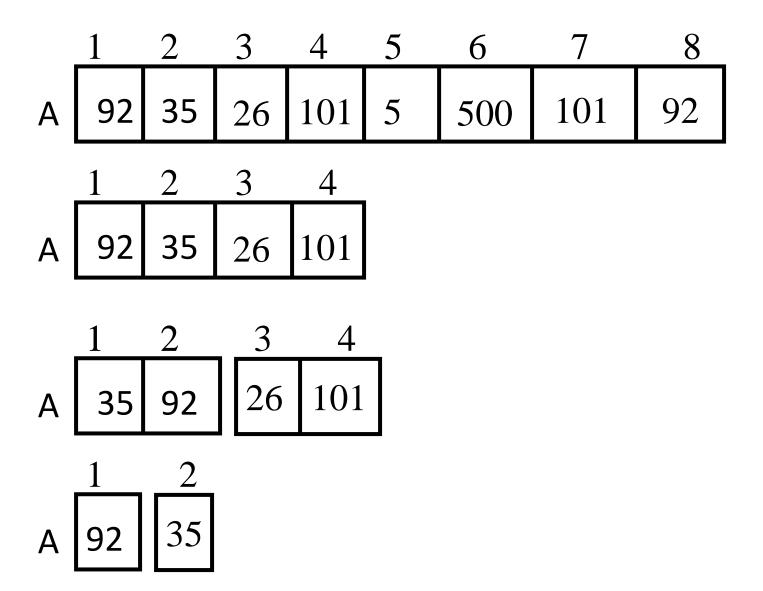
Merge Sort: Conquer (Level 4)



Merge Sort: Conquer (Level 3)

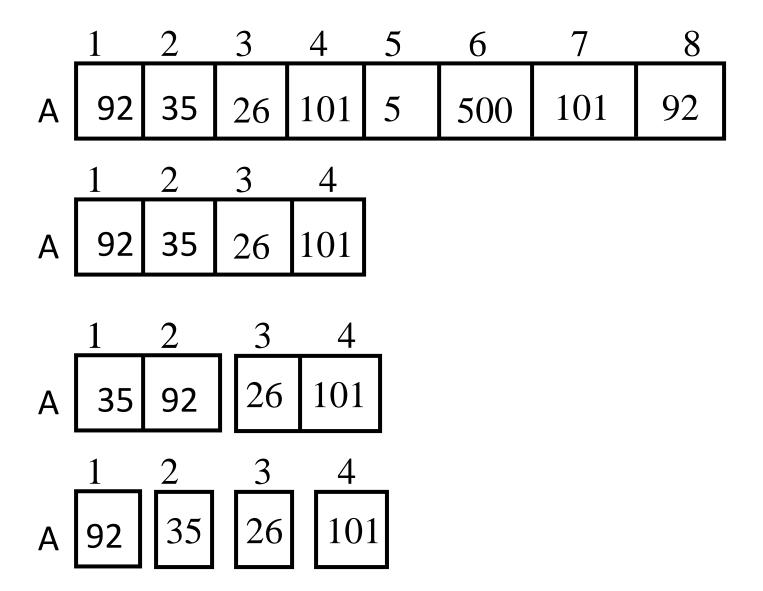


Merge Sort : Divide (Level 2)

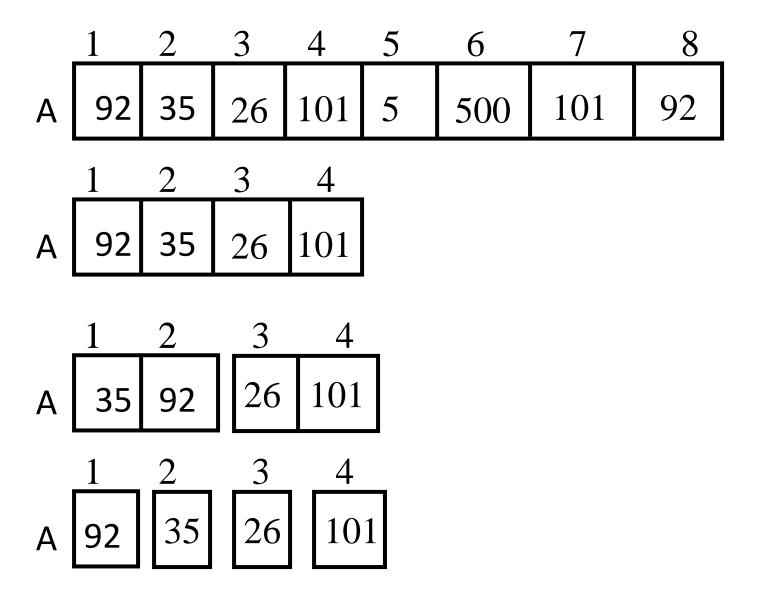


Merge Sort : Divide (Level 3) & Conquer (Level 4)

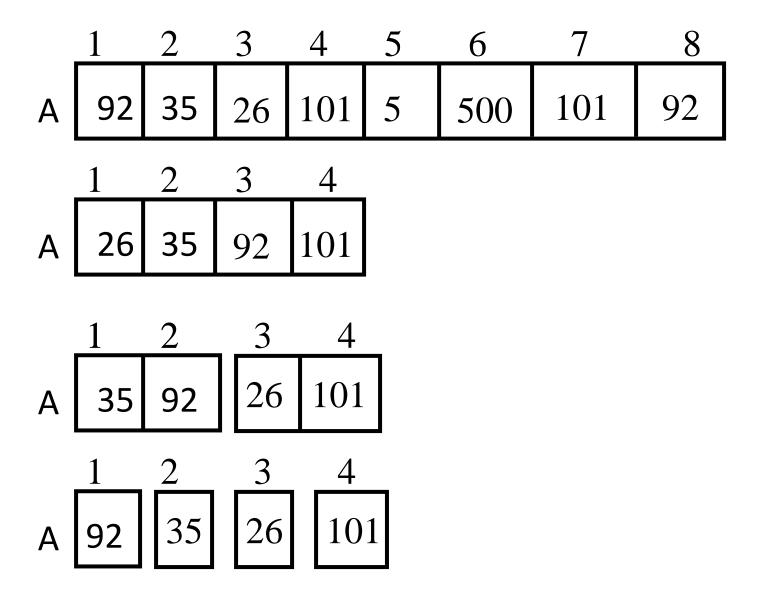
Merge Sort: Conquer (Level 4)



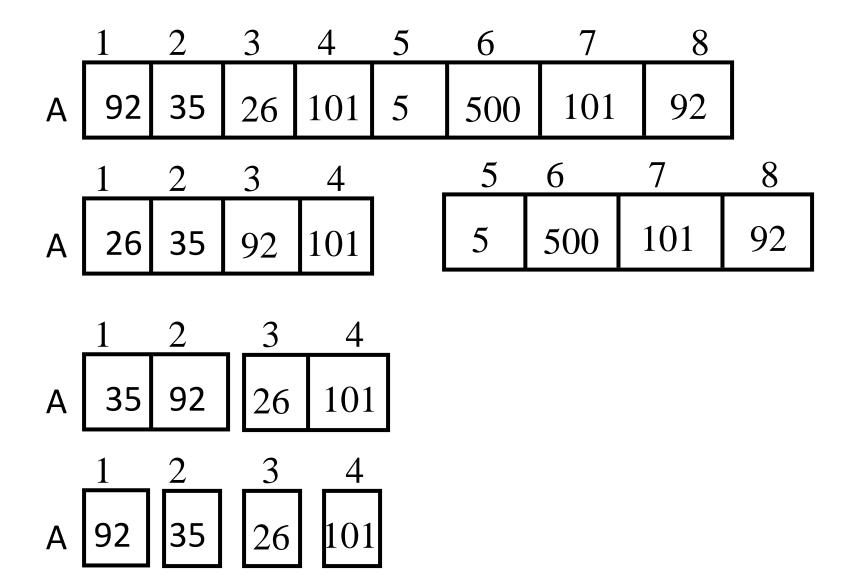
Merge Sort: Conquer (Level 3)



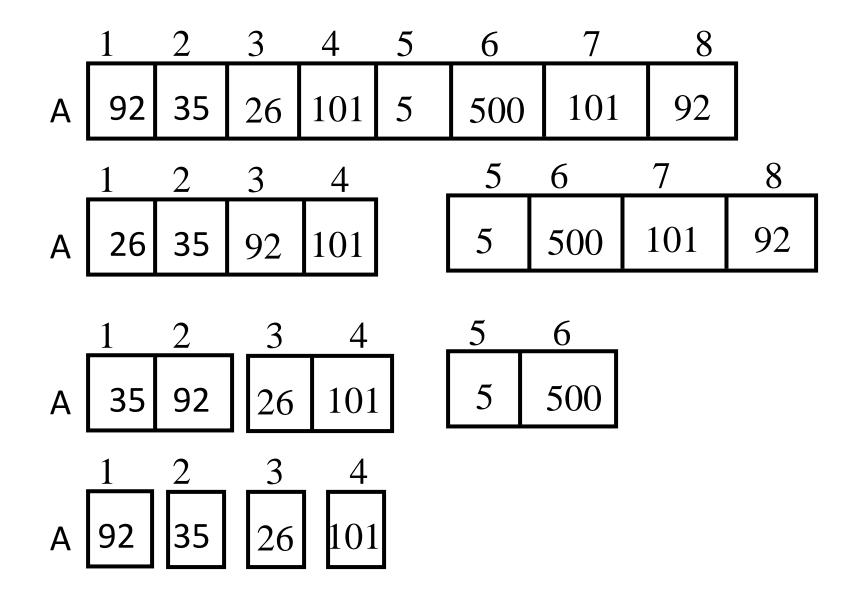
Merge Sort: Conquer (Level 2)



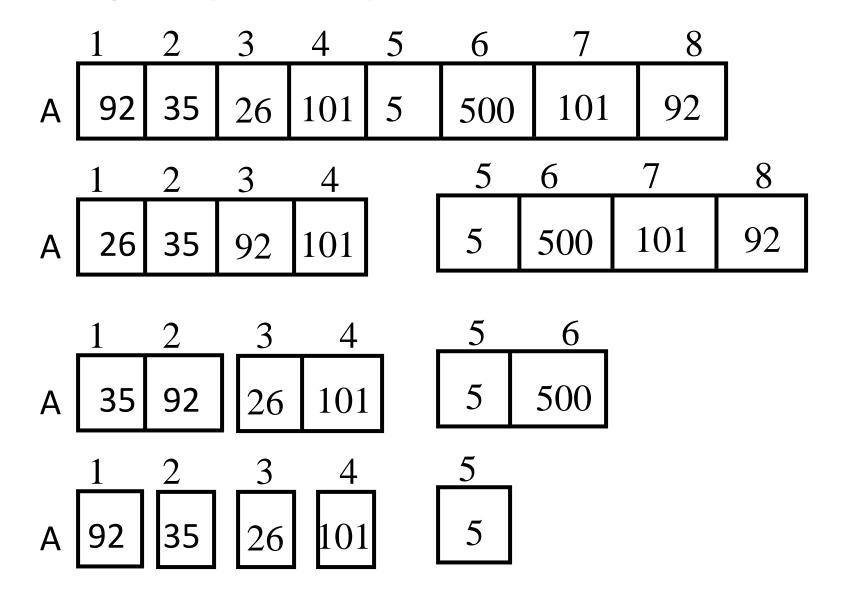
Merge Sort : Divide (Level 1)



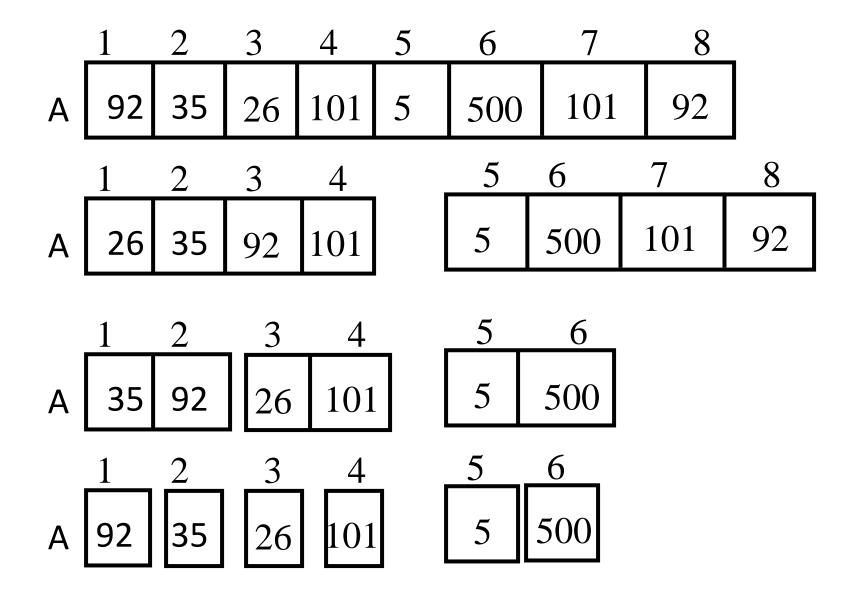
Merge Sort: Divide (Level 2)



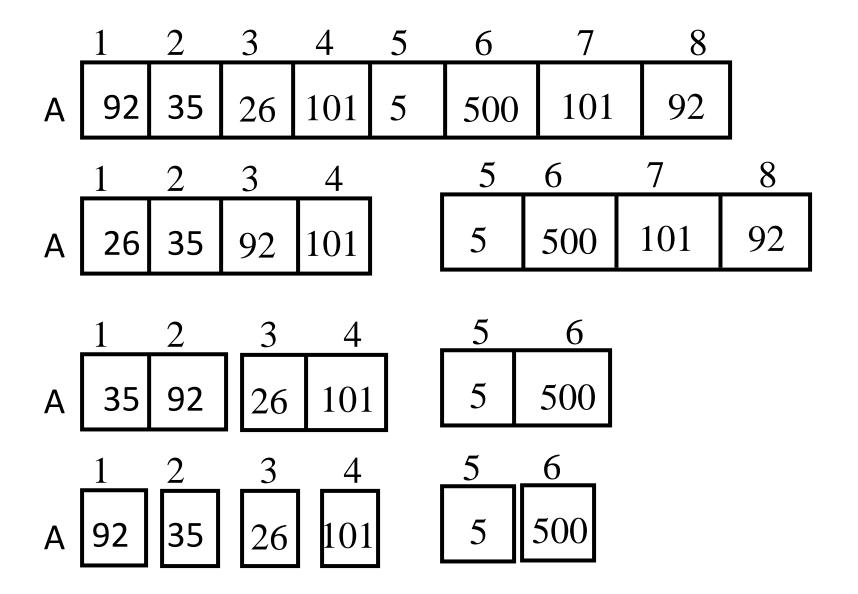
Merge Sort : Divide (Level 3) & Conquer (Level 4)



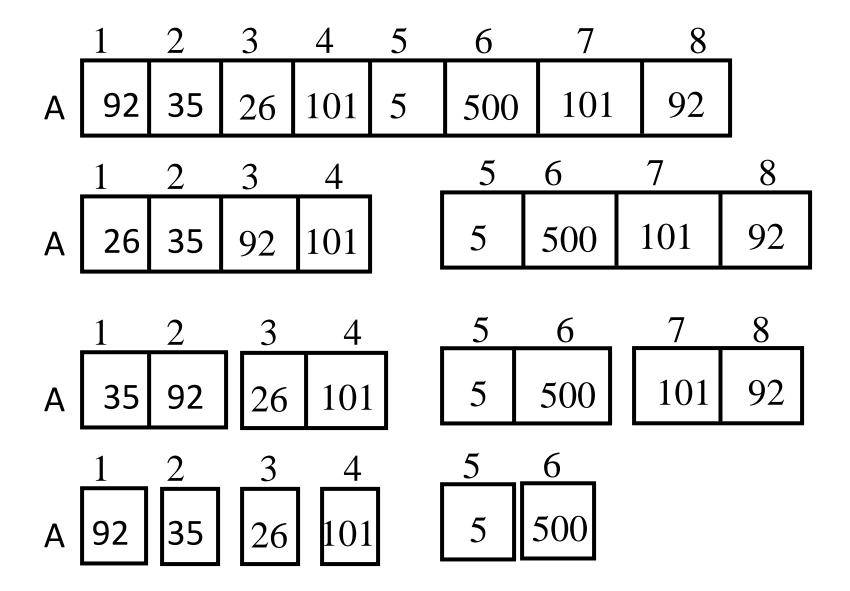
Merge Sort : Conquer (Level 4)



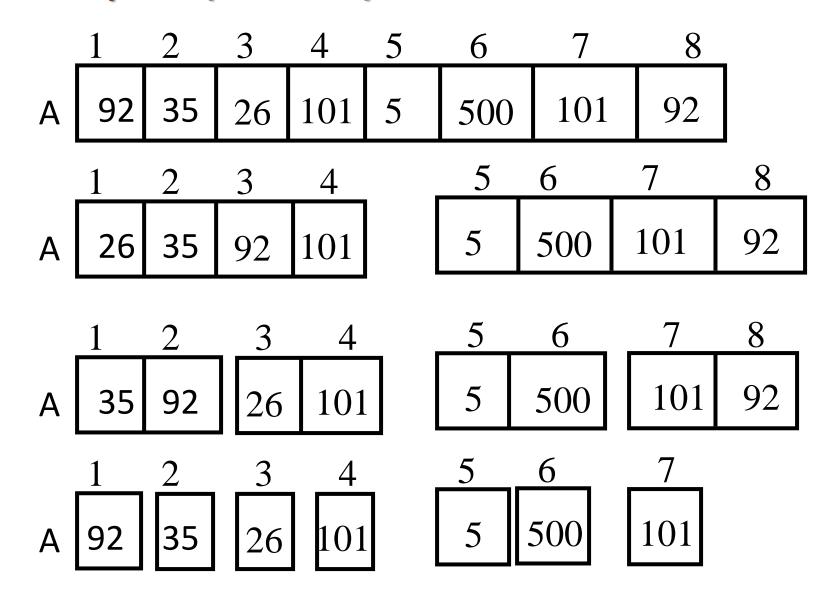
Merge Sort : Conquer (Level 3)



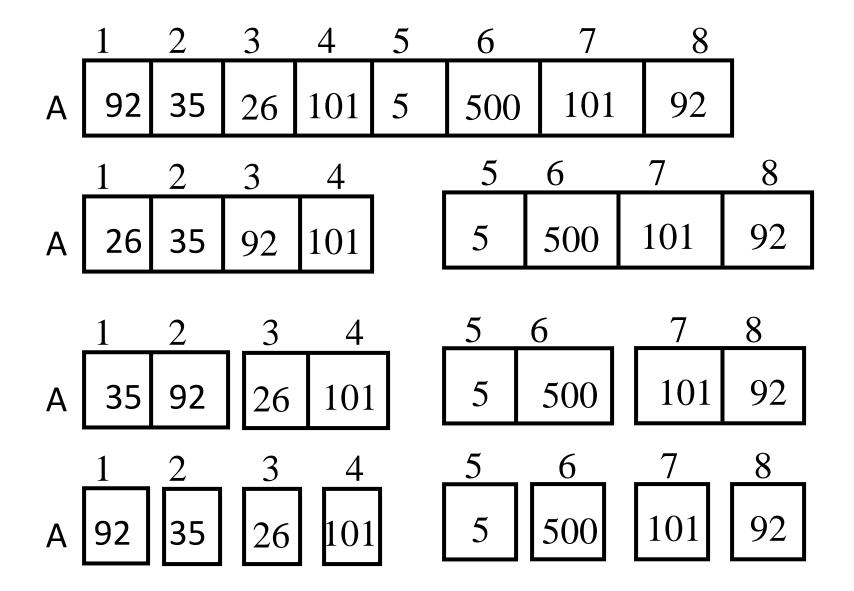
Merge Sort: Divide (Level 2)



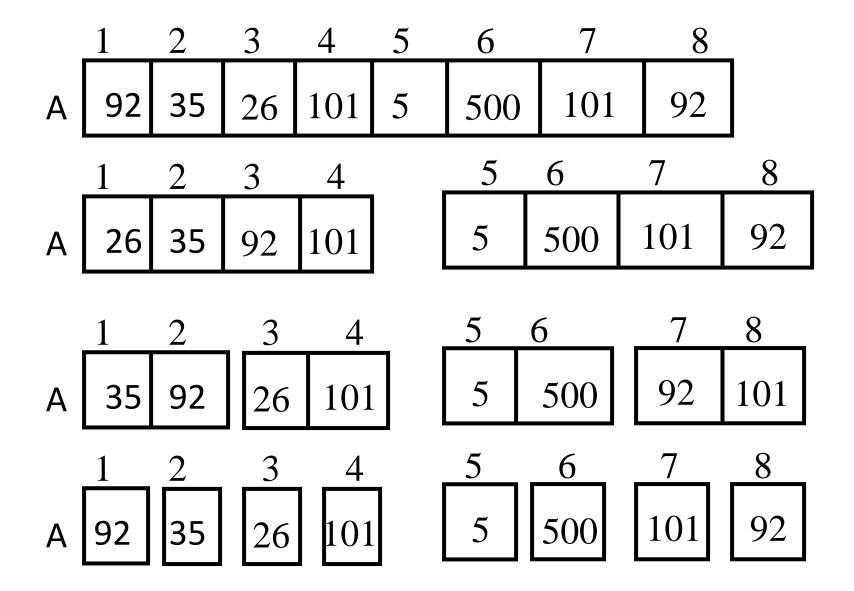
Merge Sort : Divide (Level 3) & Conquer (Level 4)



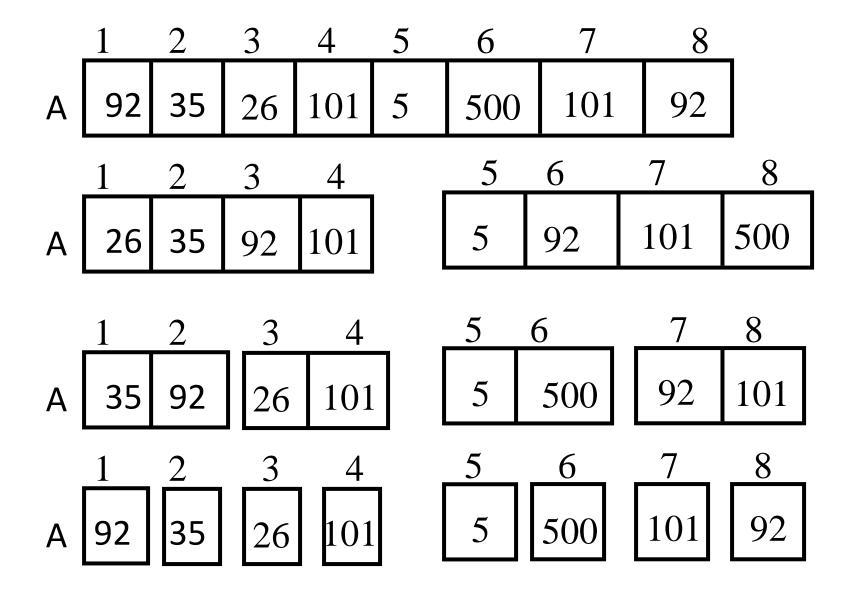
Merge Sort : Conquer (Level 4)



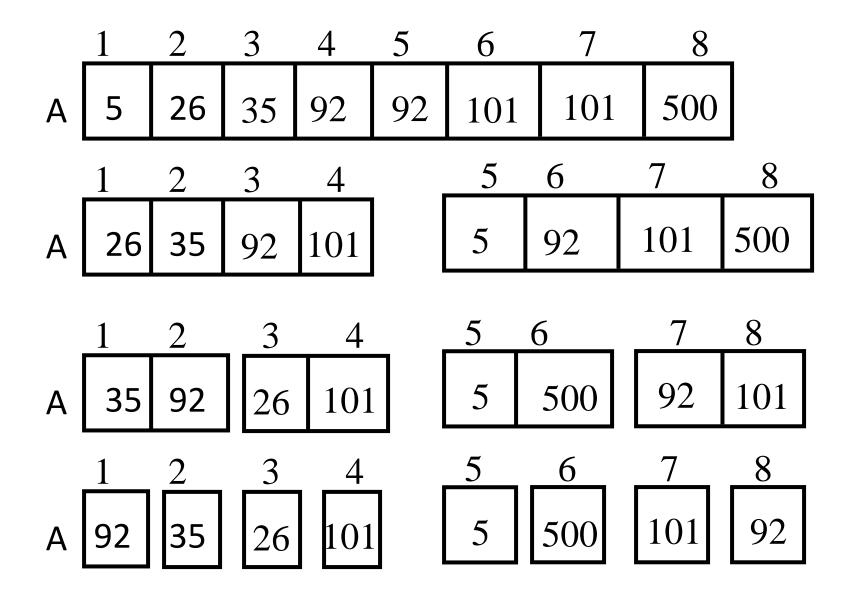
Merge Sort : Conquer (Level 3)



Merge Sort : Conquer (Level 2)



Merge Sort: Conquer (Level 1)



```
void mergesort(int a[], int low, int high)
 if(low < high) {
     int mid = (high+low)/2;
     mergesort(a, low, mid); // Sort the left half
     mergesort(a, mid+1, high); // Sort the right half
     merge(a, low, mid, high); // Merge the sorted halves
```

Time Complexity: Merge Sort

```
T(n) = Time to merge sort array of size n = 2<sup>k</sup> integers
T(1) = 1

T(n) = 2T(n/2) + cn

2T(n/2) = 2<sup>2</sup>T(n/2<sup>2</sup>) + cn

2<sup>2</sup>T(n/2<sup>2</sup>) = 2<sup>3</sup>T(n/2<sup>3</sup>) + cn

......

2<sup>k-1</sup>T(n/2<sup>k-1</sup>) = 2<sup>k</sup>T(n/2<sup>k</sup>) + cn

T(n) = 2<sup>k</sup>T(n/2<sup>k</sup>) + cn. k (Summation of k equations)

T(n) = nT(1) + c.n.log<sub>2</sub>n = n + c.n.log<sub>2</sub>n = O(nlog<sub>2</sub>n)
```

Space Complexity: Merge Sort

- Space needed to solve merge sort
- Do not consider the space for input array
- Consider the extra space needed in terms of input size
- At the bottom level, an array of size 2 is required
- Immediately higher level, an array of size 4 is required
- In the next higher level, an array of size 8 is required
- ...
- In the top level, an array of size $2^k = n$ is required
- Important to note: when array in one level is used/created, the arrays in the bottom levels have already been destroyed.
- Therefore, the space complexity = $O(max(2, 4, 8, ..., 2^k = n)) = O(n)$