# Stack

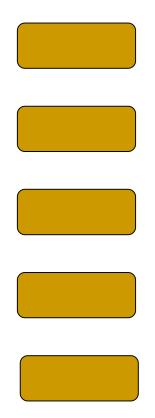
Joy Mukherjee

#### Stack

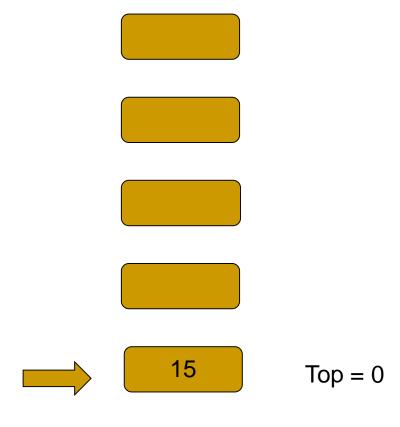
- Stack is a linear data structure
- Insertion and deletion take place on same end
- LIFO(Last In First Out)
  - The last one inserted is the first one deleted
- FILO(First In Last Out)
  - The first one inserted is the last one deleted

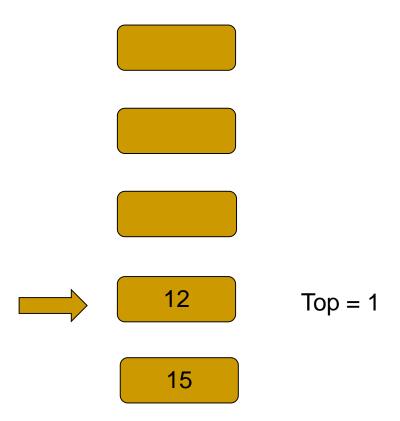
- int A[5], top;
- The maximum size of the stack is 5
- Top points to the topmost element of the stack.

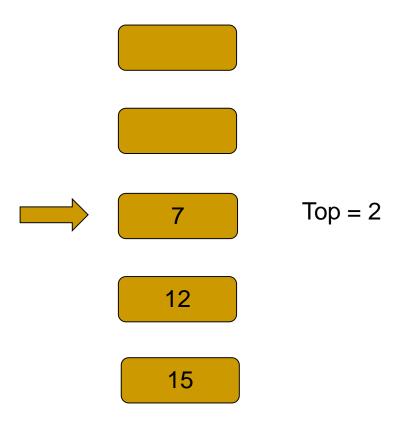
#### **Operations: Initially Stack is Empty**

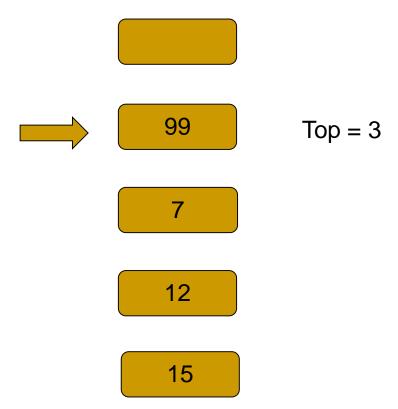


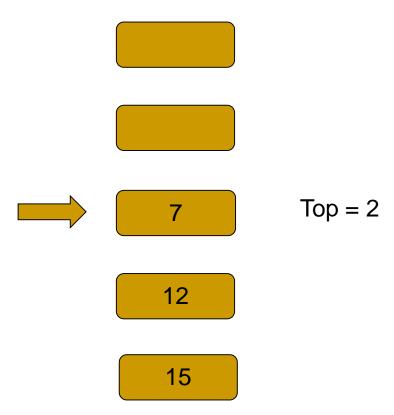


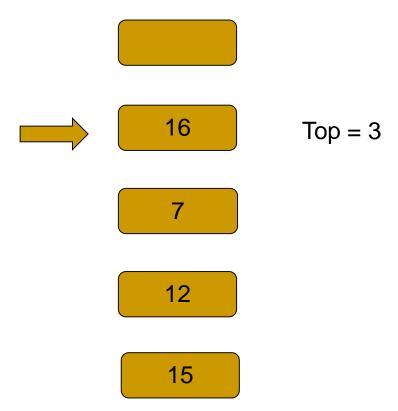




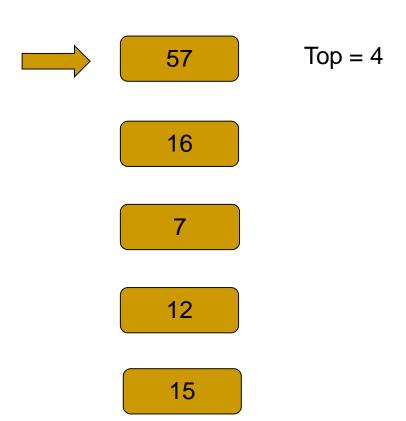




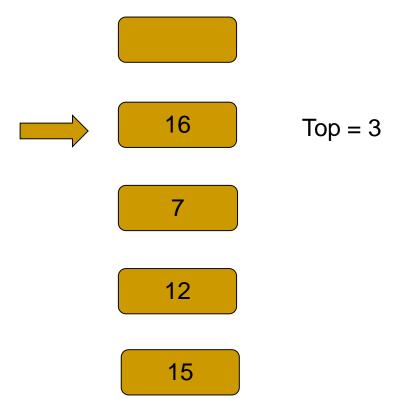


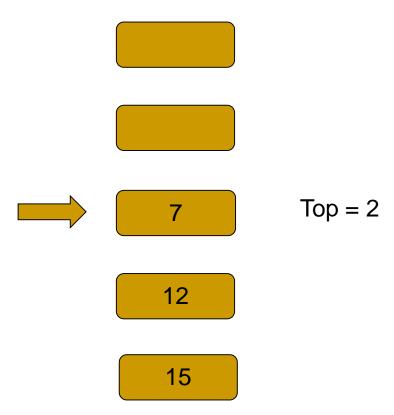


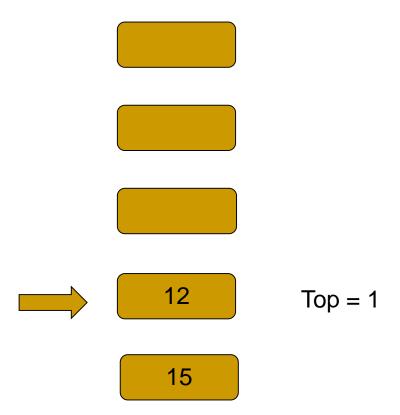
Top = 4

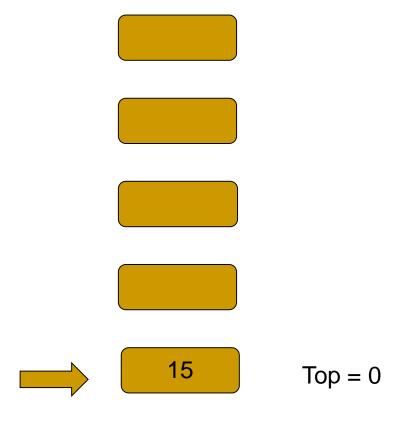


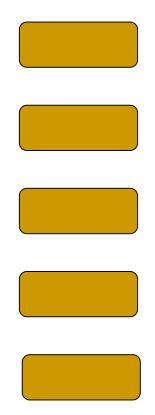




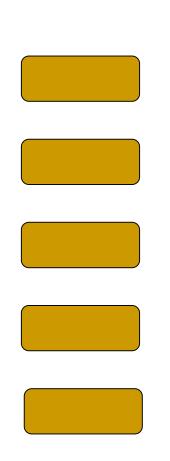


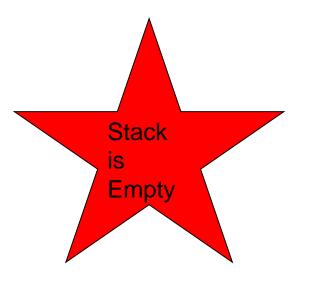














#### Stack

```
#define MAXLEN 5
struct st{
       int A[MAXLEN];
       int top;
}; // User-defined data type struct st is created only
typedef struct st stack; // struct st and stack
represent the same data type
```

### **Stack Operations**

```
int isEmpty (stack S)
stack init ()
                                     if(S.top == -1)
  stack S;
                                         return 1;
  S.top = -1;
                                      return 0;
  return S;
               int isFull (stack S)
                 if (S.top == MAXLEN - 1)
                    return 1;
                 return 0;
```

### **Stack Operations**

```
stack push (stack S, int k)
                                    stack pop (stack S)
  if (isFull(S) == 1) {
                                      if (isEmpty(S) == 1) {
     printf("Stack is Full");
                                            printf("Stack is Empty");
     return S;
                                            return S;
  S.top = S.top + 1;
                                       S.top = S.top - 1;
  S.A[S.top] = k;
                                       return S;
  return S;
```

# Stack: Applications

Infix to Postfix Conversion
Postfix Evaluation
Share Span Problem

### Assignment - 1

- Subject Line: A1\_ROLLNO
- File Name: A1\_ROLLNO.c/cpp
- Email: pds2016autumn@gmail.com
- Submission: 08-August-2021 11:59 P.M.
- Write a C program that takes a mathematical expression as input (infix), and output its value.
- Input: 4+(8-3)\*2-(3+4)
- Output: 7
- Input: 40/8 30\*2
- Output: -55

- Write a C program that takes a mathematical expression as input (infix), and output its value.
- Input: 4+(8-3)\*2-(3+4)
- Output: 7
- Input: 40/8 30\*2
- Output: -55
- Convert Infix to Postfix (Stack)
- Evaluate the Postfix (Stack)

#### Example

- Infix: A B + C (Operator is present in between operands)
- Prefix: +-ABC
- Postfix: AB-C+
- Infix: A B + C \* D(Operator is present in between operands)
- Prefix: +-AB\*CD
- Postfix: AB-CD\*+

- 4+(8-3)\*2-(3+4)
- 4 + 8 3 -\* 2-(3+4)
- 4+83-2\*-(3+4)
- 4 + 8 3 2 \* 3 4+
- 483-2\*+-34+
- 483-2\*+34+-
- Input: Infix
- Output: Postfix
- Assumption:
  - Each operator is a binary operator
  - Associativity of an operator is from left to right
    - Example: A+B+C-D = (((A+B)+C)-D)

#### **Postfix Evaluation**

- 483-2\*+ 34+-
- **7**
- Input: Infix
- Output: Postfix
- Assumption:
  - Each operator is a binary operator
  - Associativity of an operator is from left to right
    - Example: A+B+C-D = (((A+B)+C)-D)

Any mathematical expression is written as an infix expression

It consists of tokens: (, ), operators, and operands

Operator: + - \* /

Operand: Variable, or Value

Infix expression: An operator exists between two operands

$$A + B$$
,  $A - 5 / C$ ,  $(A - 5) / C$ 

Postfix expression: Operands followed by an operator

$$AB+, A5C/-, A5-C/$$

### Infix to postfix conversion: Algorithm

For each token from an infix expression from left to right:

- If (token = operand)
  - Push it to postfix expression
- □ If (token = '(')
  - Push ( to the stack
- □ If (token = ')')
  - Repeatedly pop a token y from stack and push y to postfix expression until '(' is encountered in stack.
  - Pop '(' from stack.
  - If stack is already empty before finding a '('
    - Expression is invalid

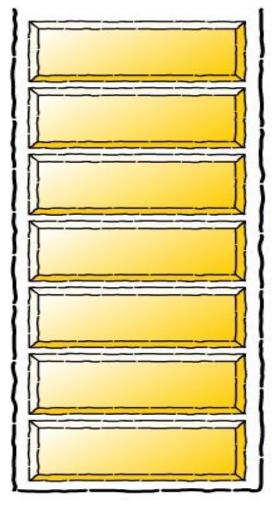
### Infix to postfix conversion: Algorithm

- If (token x = operator)
  - While(the following three conditions are true)
    - The stack is not empty
    - 2. The token y is not a '('.
    - 3. The token y is an operator && Precedence(y) >= Precedence (x),
      - Push the token y into postfix expression
      - Pop the token y from stack
  - Push the token x into stack.

### Infix to postfix conversion: Algorithm

- After the loop has processed all the tokens in the infix expression, use another loop to repeatedly do the following as long as the stack is not empty
  - Pop token y from the stack.
  - Push the token y into postfix expression.



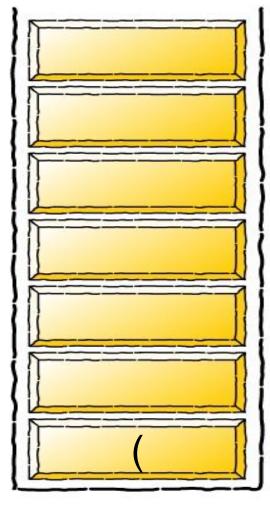


infixVect

$$(a+b-c)*d-(e+f)$$

postfixVect

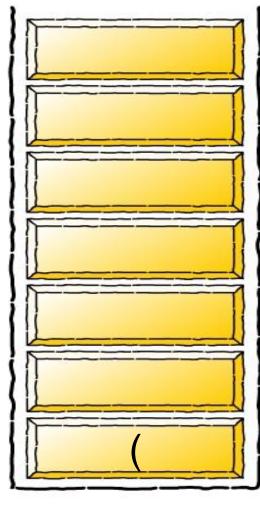




infixVect

postfixVect



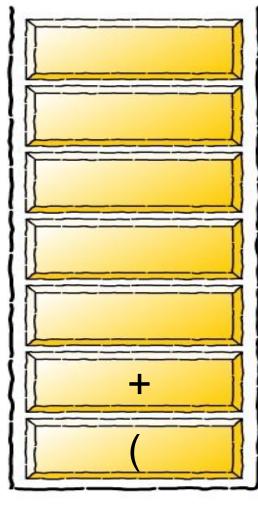


infixVect

postfixVect

a



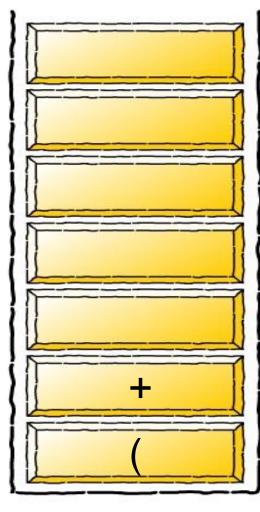


infixVect

postfixVect

a



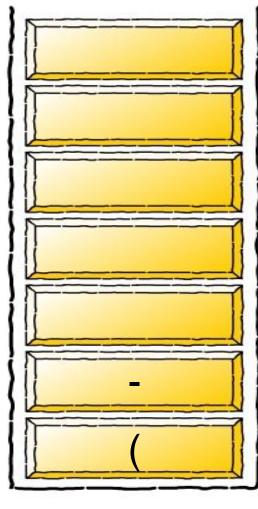


infixVect

postfixVect

a b



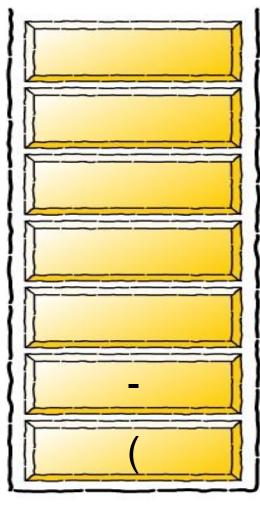


infixVect

postfixVect

a b +

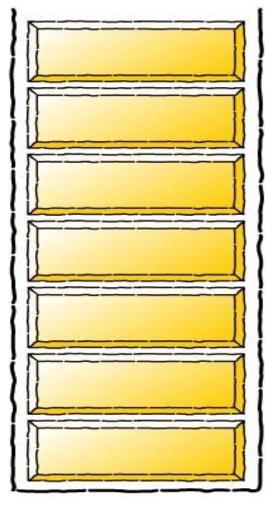




infixVect

$$ab+c$$

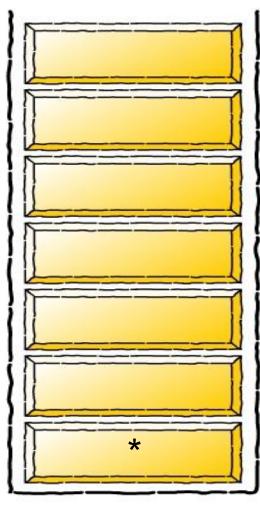
#### stack



infixVect

$$ab+c-$$



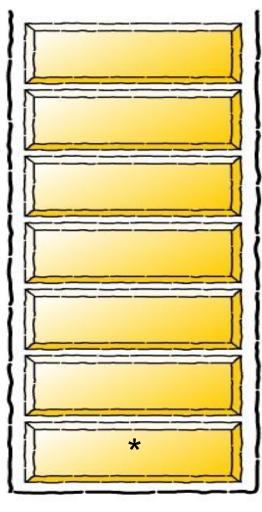


infixVect

$$d-(e+f)$$

$$ab+c-$$

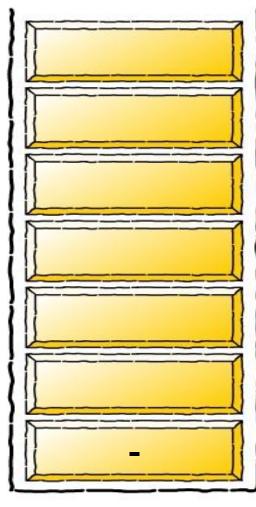




infixVect

$$ab+c-d$$

#### stack

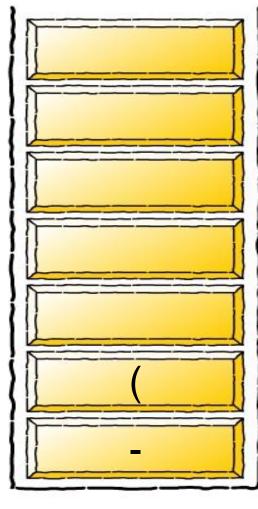


infixVect

$$(e+f)$$

$$ab+c-d*$$



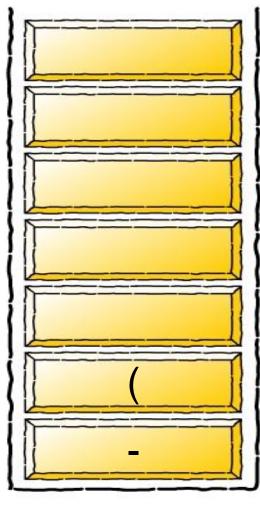


infixVect

$$e + f$$
)

$$ab+c-d*$$

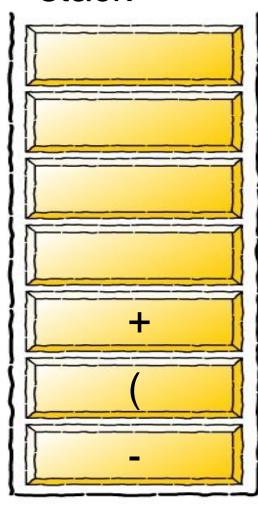




infixVect

$$ab+c-d*e$$



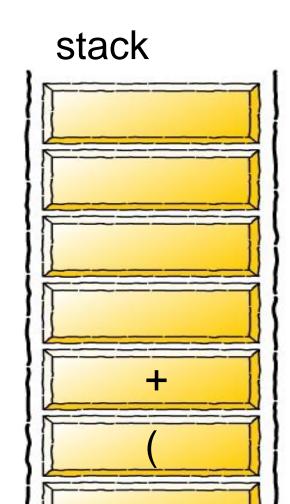


infixVect

f )

postfixVect

ab+c-d\*e



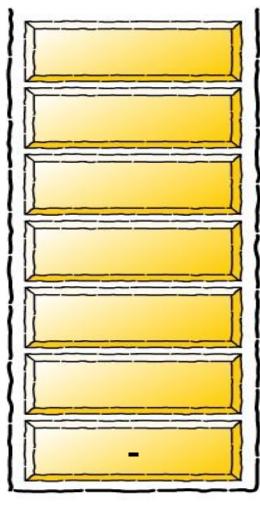
infixVect

)

postfixVect

ab+c-d\*ef



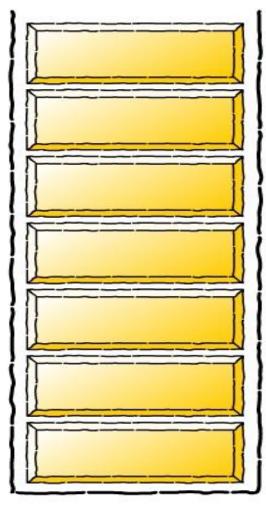


infixVect

postfixVect

ab+c-d\*ef+

#### stack



infixVect

postfixVect

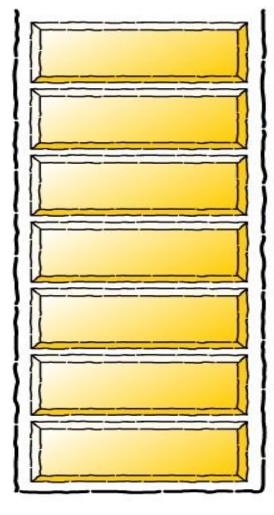
ab + c - d \* ef + -

#### infixVect

$$(a+b-c)*d-(e+f)$$

$$ab + c - d * ef + -$$

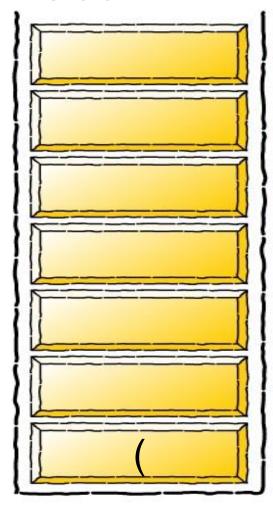




infixVect

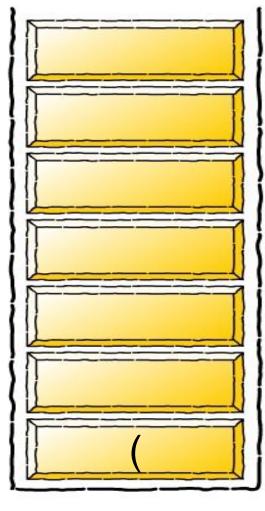
$$(a+b*c)$$





infixVect

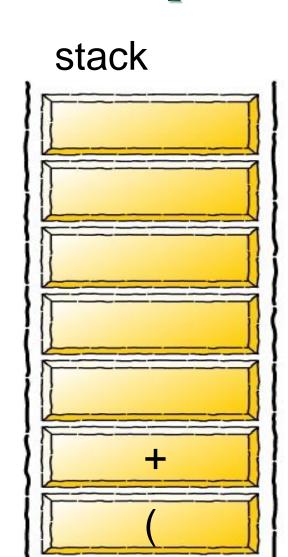




infixVect

postfixVect

a

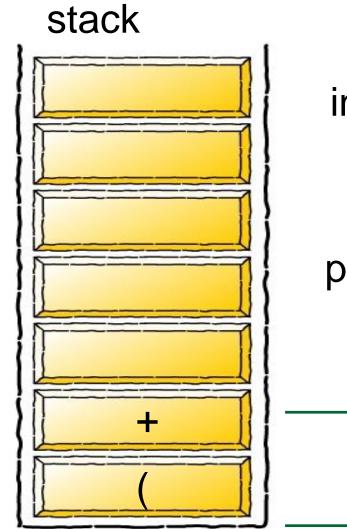


infixVect

b \* c )

postfixVect

a

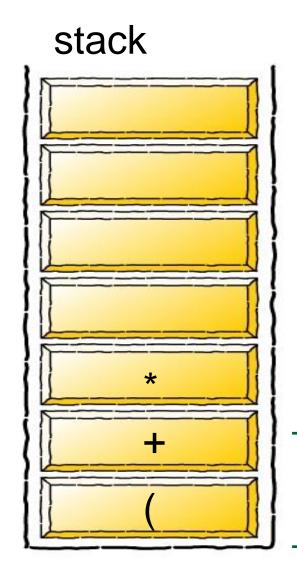


infixVect

\* C )

postfixVect

a b

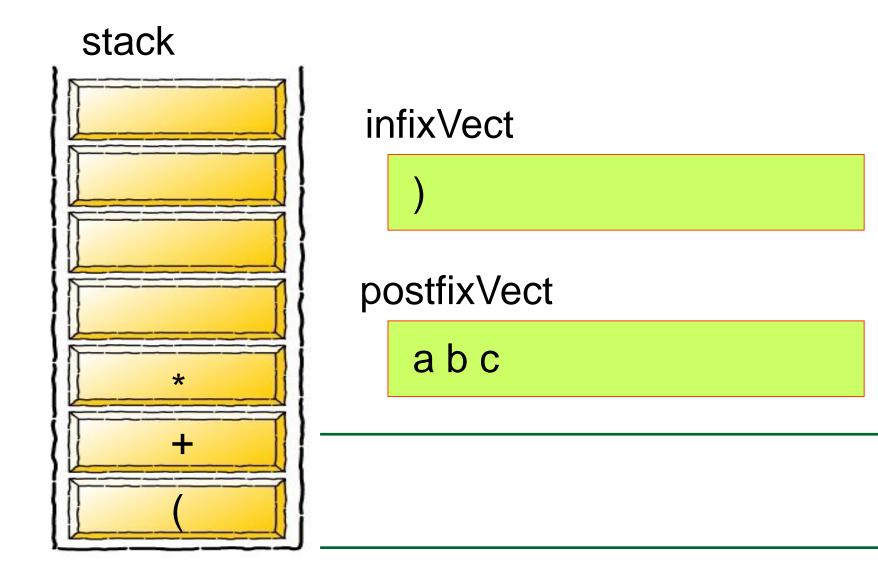


infixVect

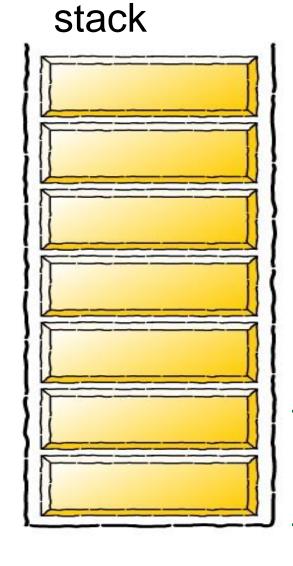
**c**)

postfixVect

a b







infixVect

postfixVect

a b c \* +

```
For each token in the postfix from left to right

if (token = Operand)

Push Operand to the stack

if (token = Operator)

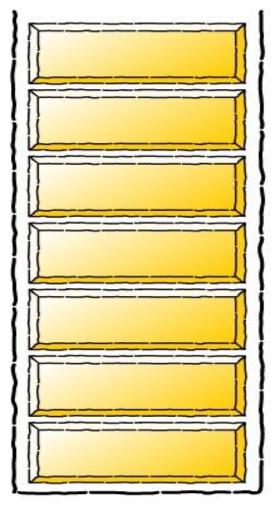
Second = Pop from the stack

First = Pop from the stack

Push First Operator Second to the stack

Return Stack[top]
```

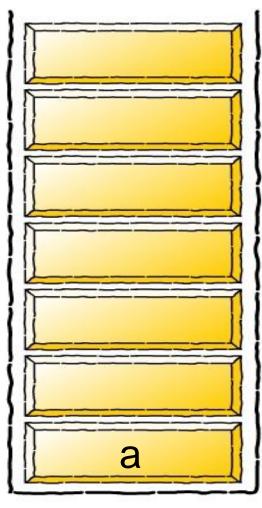
#### stack



$$(a+b-c)*d-(e+f)$$

$$ab + c - d * ef + -$$

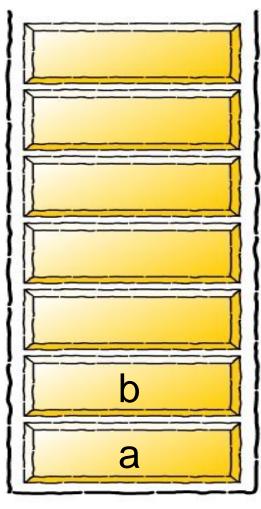
#### stack



$$(a+b-c)*d-(e+f)$$

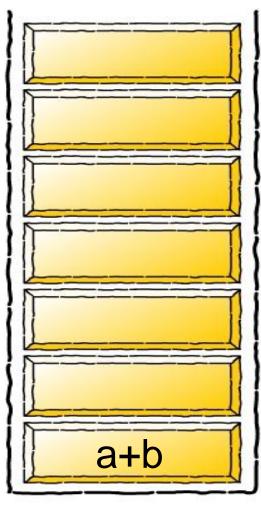
$$b + c - d * e f + -$$

#### stack



$$(a+b-c)*d-(e+f)$$

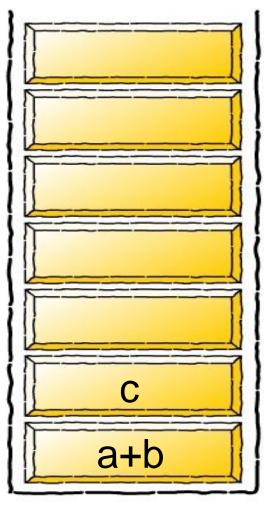
#### stack



$$(a+b-c)*d-(e+f)$$

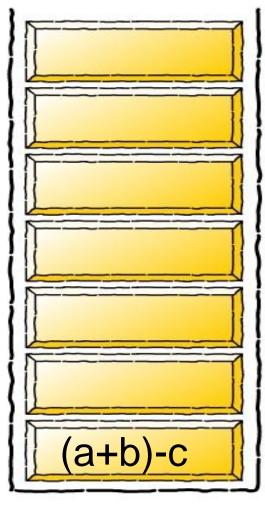
$$c - d * e f + -$$

#### stack



$$(a+b-c)*d-(e+f)$$

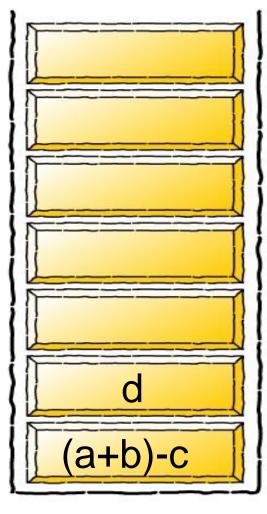




$$(a+b-c)*d-(e+f)$$

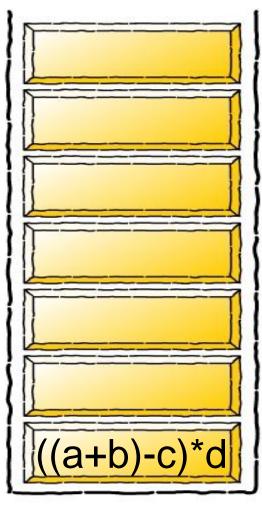
$$d * e f + -$$

#### stack



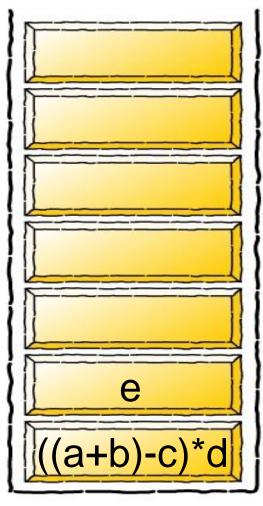
$$(a+b-c)*d-(e+f)$$

#### stack



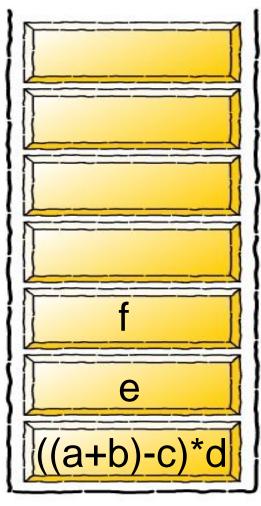
$$(a+b-c)*d-(e+f)$$





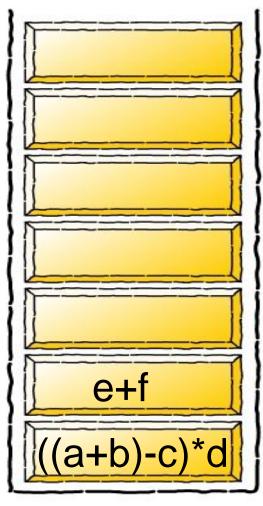
$$(a+b-c)*d-(e+f)$$





$$(a+b-c)*d-(e+f)$$



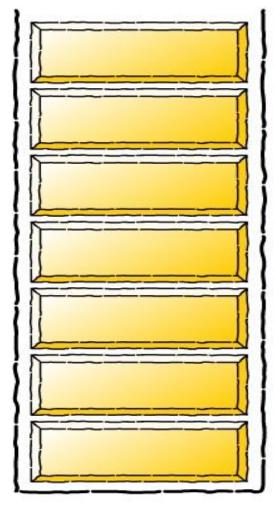


$$(a+b-c)*d-(e+f)$$

postfixVect

\_

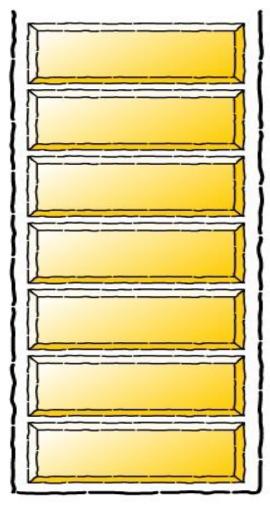
#### stack



$$(a+b-c)*d-(e+f)$$

$$(((a+b)-c)*d)-(e+f)$$

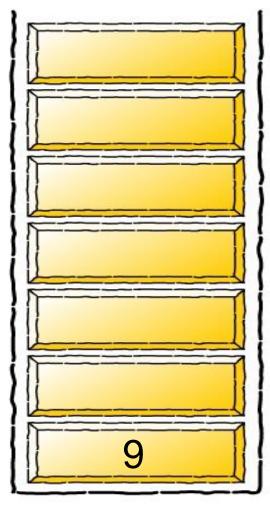
#### stack



$$(9+4-6)*2-(3+6)=5$$

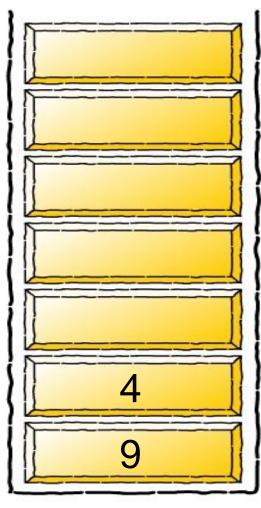
$$94 + 6 - 2 * 36 + -$$

#### stack



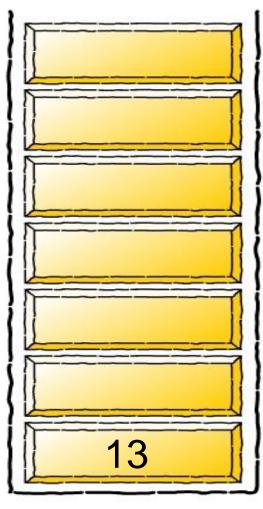
$$(9+4-6)*2-(3+6)=5$$

#### stack



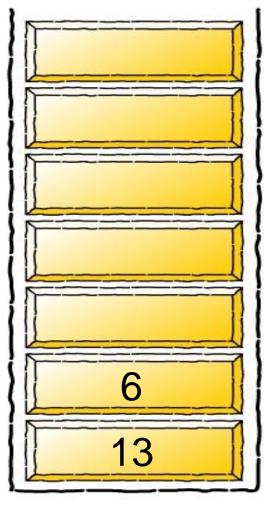
$$(9+4-6)*2-(3+6)=5$$

#### stack



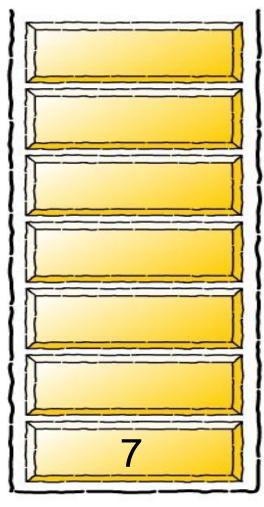
$$(9+4-6)*2-(3+6)=5$$





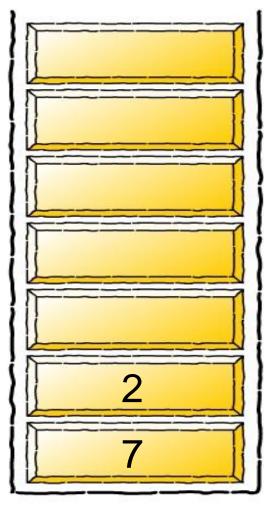
$$(9+4-6)*2-(3+6)=5$$

#### stack



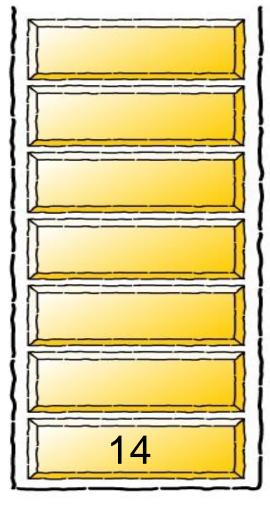
$$(9+4-6)*2-(3+6)=5$$

#### stack



$$(9+4-6)*2-(3+6)=5$$

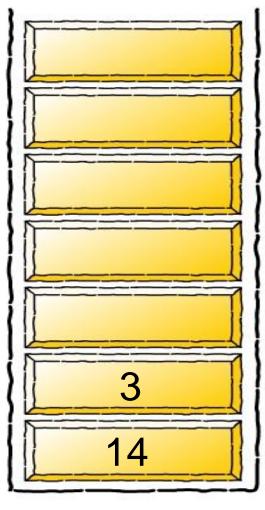




$$(9+4-6)*2-(3+6)=5$$

$$36 + -$$

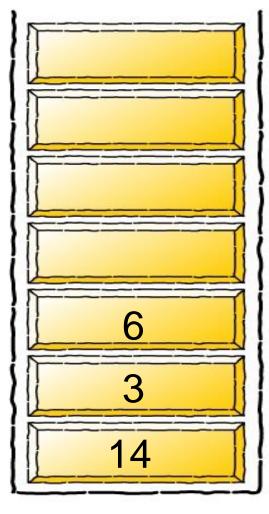




$$(9+4-6)*2-(3+6)=5$$

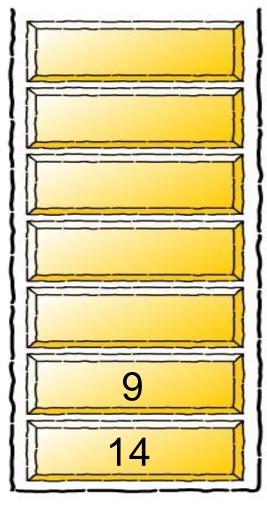
$$6 + -$$

#### stack



$$(9+4-6)*2-(3+6)=5$$



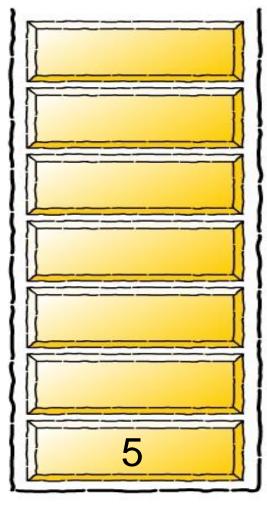


$$(9+4-6)*2-(3+6)=5$$

postfixVect

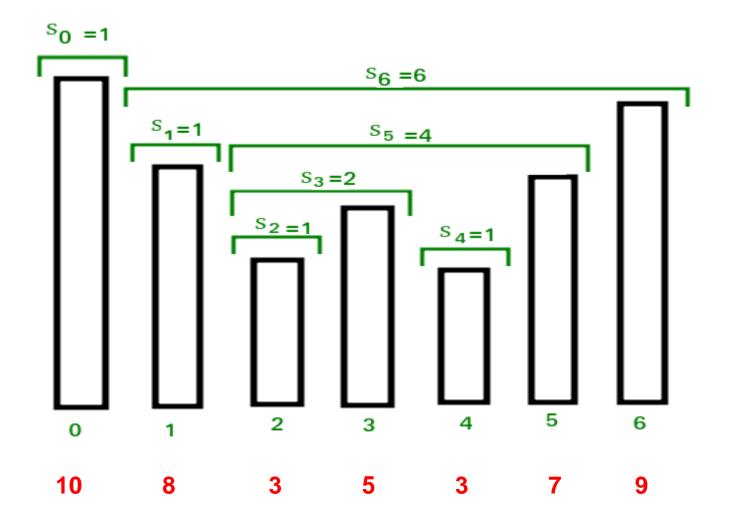
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$$(9+4-6)*2-(3+6)=5$$

- Input: A sequence of n prices for a share for n consecutive days
- Output: Calculate span of share's price for all n days.
- Share Span for Day i
  - Span[i]: Maximum no of consecutive days on or before i-th day such that price(i) >= price(all the previous consecutive days)



- Input: A sequence of n prices for a share for n consecutive days
- Output: Calculate span of share's price for all n days.
- Example:

$$Price[7] = \{10, 9, 6, 7, 6, 7, 8\}$$

**Output:** Span[7] =  $\{1, 1, 1, 2, 1, 4, 5\}$ 

■ Input: n = 7 and

$$Price[7] = \{5, 9, 6, 7, 6, 7, 10\}$$

**Output:** Span[7] =  $\{1, 2, 1, 2, 1, 4, 7\}$ 

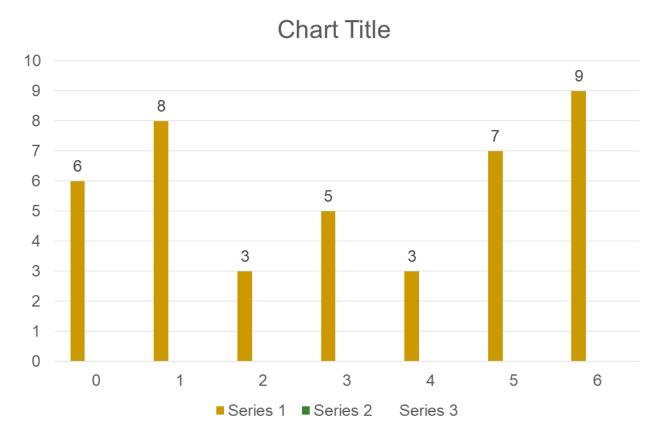
#### **Naïve Solution**

```
Span[0] = 1;
For(i = 1; i < n; i++) {
       Span[i] = 1;
       For(i = i - 1; i >= 0; j--)
               if(A[i] \le A[i])
                       Span[i] = Span[i] + 1;
Time complexity = O(n(n-1)/2) = O(n^2)
Can we do better? Yes (using stack)
```

#### **Linear Time Solution**

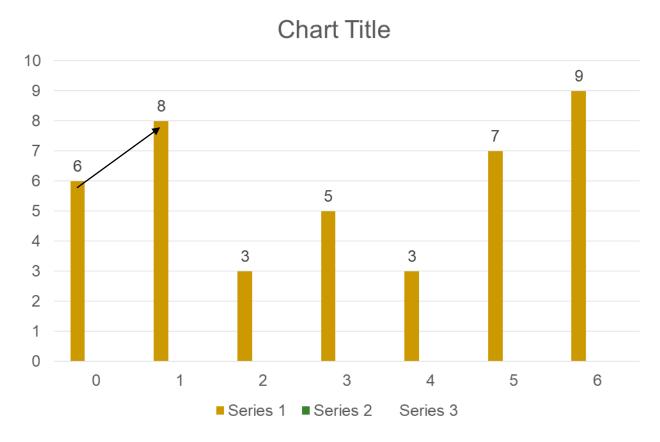
```
stack.push(0);
span[0] = 1;
for (i = 1; i < n; i++) {
   while (stack is not empty and price[top of stack] <= price[i])
       stack.pop();
   If (stack becomes empty) then it implies price[i] is greater than all
   elements on left of it, i.e., price[0], price[1], ..price[i-1].
      span[i] = i + 1;
   Else price[i] is less than the element at top of stack
      span[i] = i - stack.top();
    stack.push(i);
```

### I = 0 Push(0)



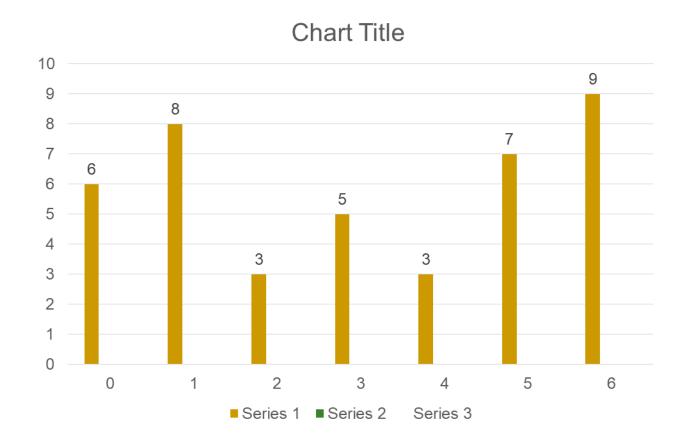
Stack 0

### I = 1



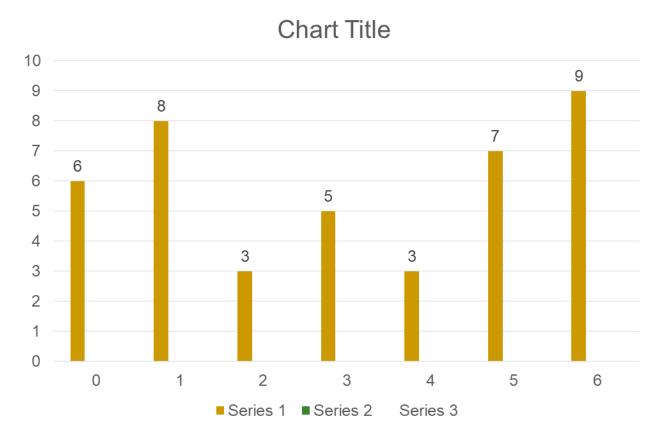
Stack 0

# Pop



#### Stack

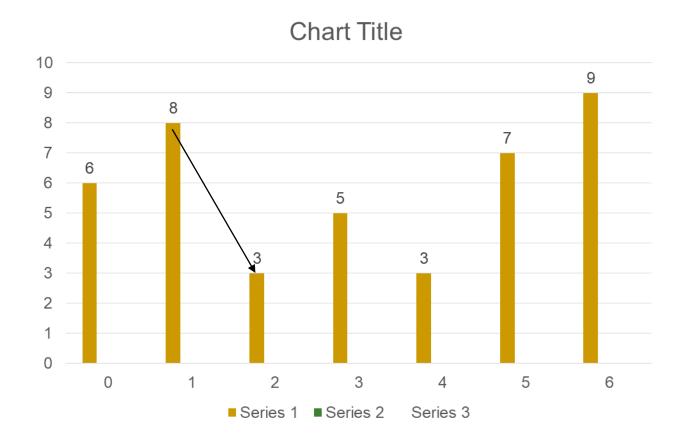
# **Push(1)**



Stack 1

Span 1 2

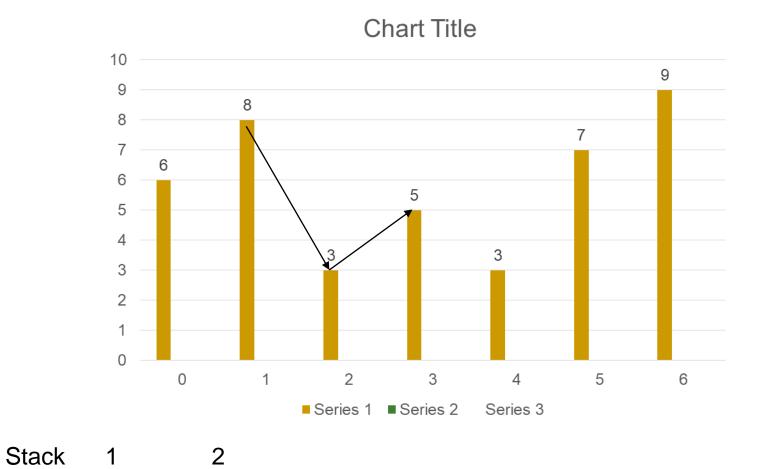
# I = 2 Push(2)



Span 1 2 1

Stack

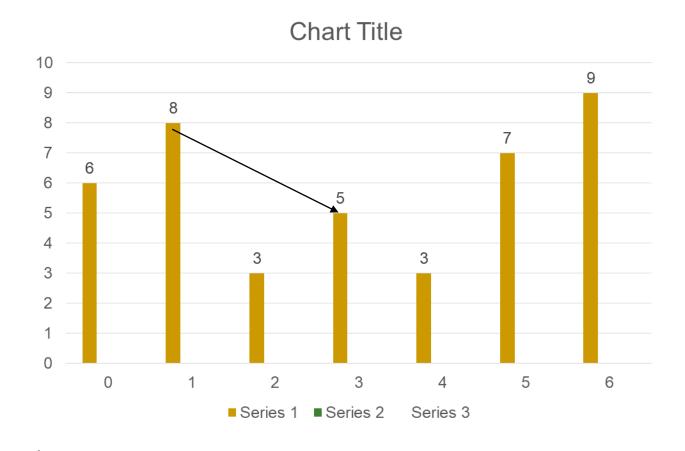
### I = 3



Span 1 2 1

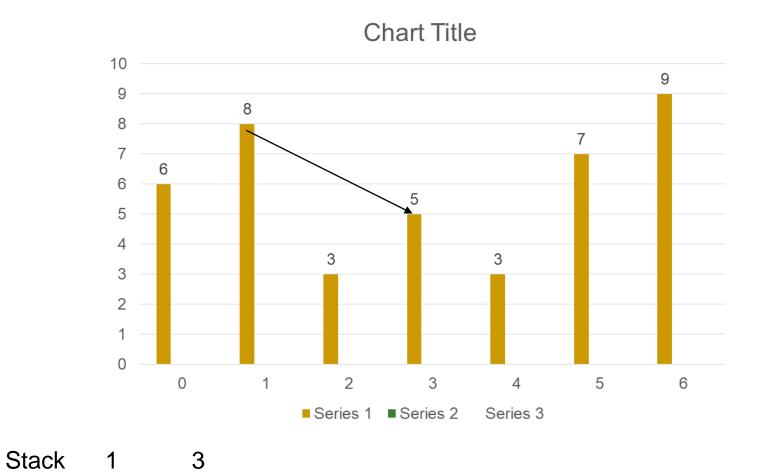
# Pop

Stack



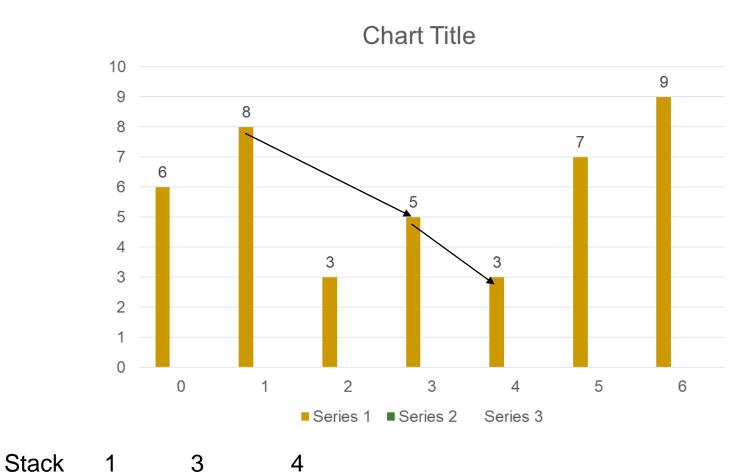
Span 1 2 1

# **Push(3)**

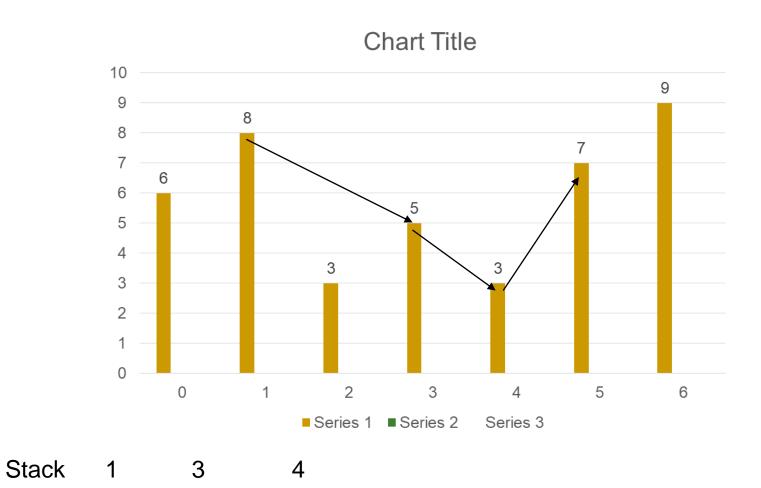


Span 1 2 1 2

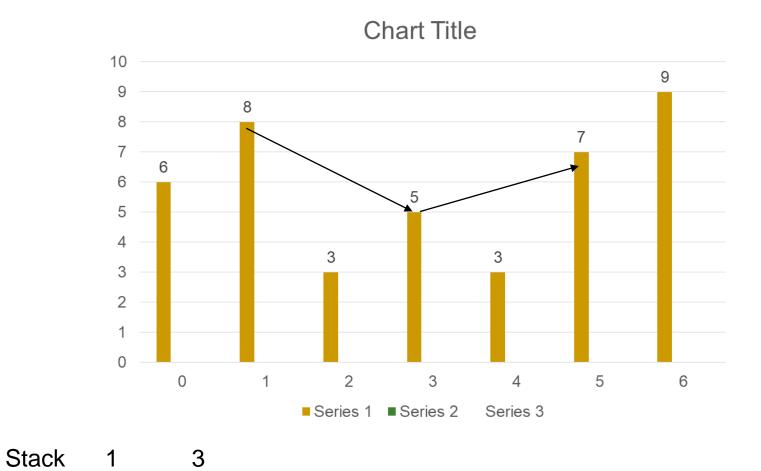
# I = 4 Push(4)



### I = 5

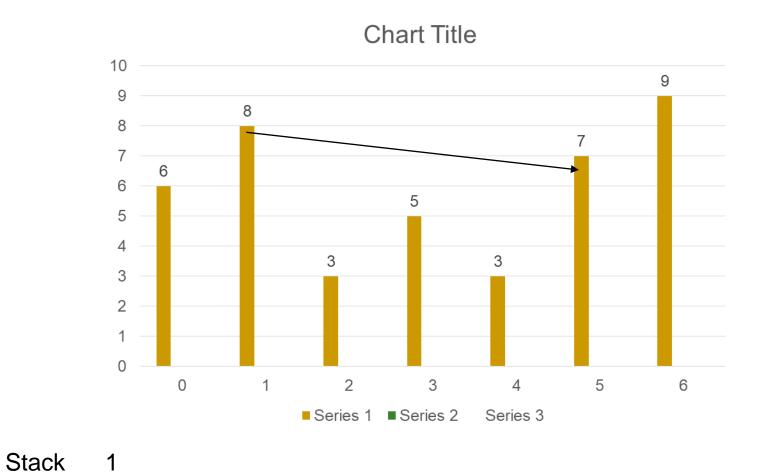


## Pop



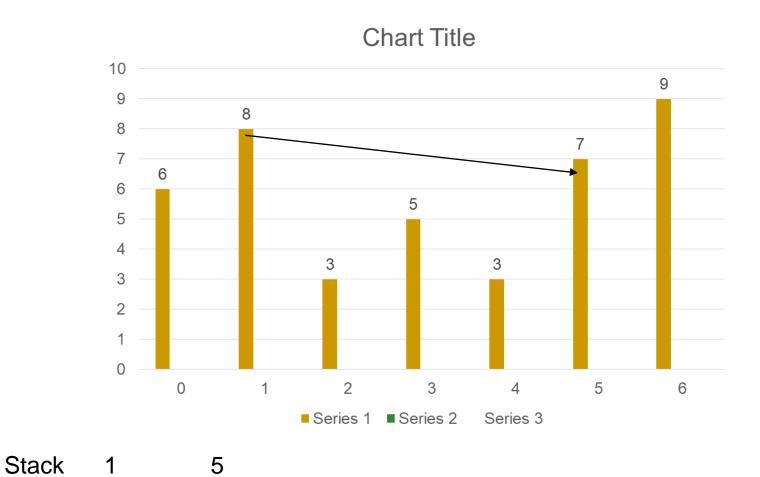
Span 1 2 1 2 1

## Pop



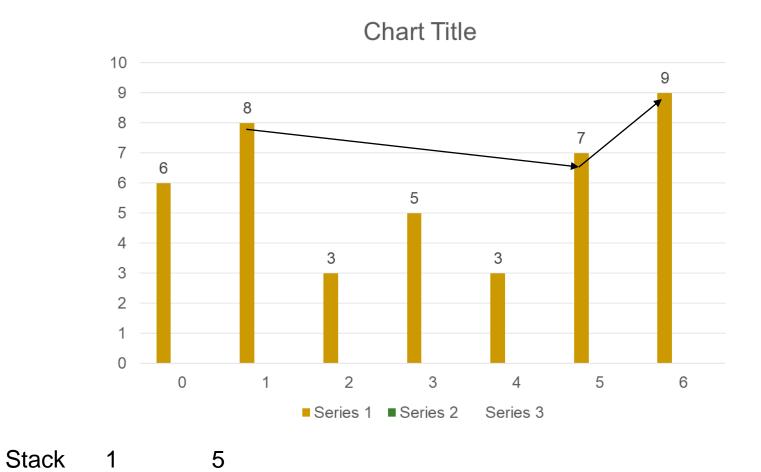
Span 1 2 1 2 1

### **Push (5)**



Span 1 2 1 4

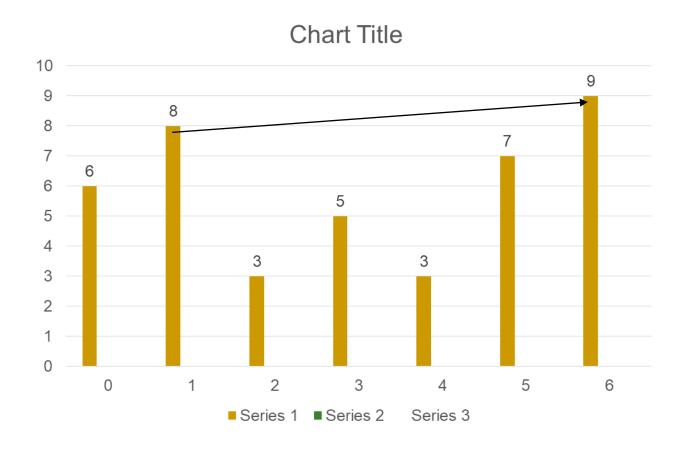
### I = 6



Span 1 2 1 2 1 4

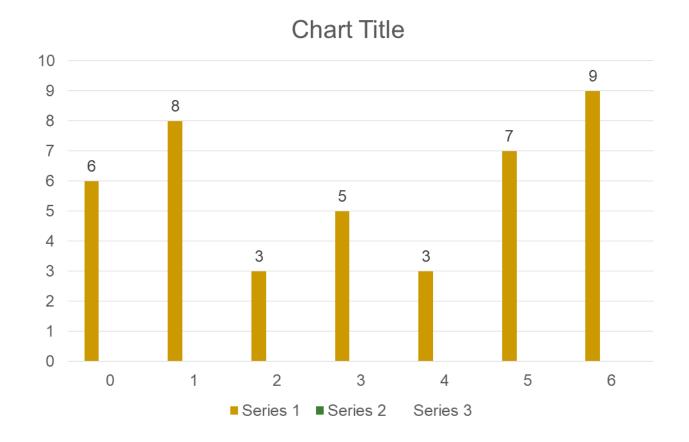
## Pop

Stack



Span 1 2 1 4

## Pop

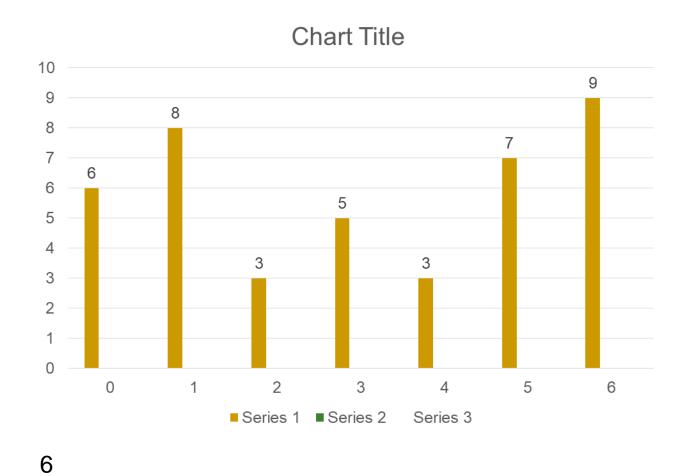


Stack

Span 1 2 1 2 1 4

### **Push (6)**

Stack



Span 1 2 1 2 1 4 7

#### **Linear Time Solution**

```
stack.push(0);
span[0] = 1;
for (i = 1; i < n; i++) {
   while( Stack is not empty && Price[Stack[top]] <= Price[i] )</pre>
      Stack.Pop()
   If ( stack is empty )
      span[i] = i + 1
   If( Price[i] < Price[Stack[top]] )</pre>
      span[i] = i - Stack[top]
   Stack.Push(i)
                                   Each day is pushed/popped at
                                   most once in the stack.
                                   Time complexity = O(2n) = O(n)
```