#### ACM SIGCOMM 2017

# Credit-Scheduled Delay-Bounded Congestion Control for Datacenters

Inho Cho, Keon Jang\*, Dongsu Han





### **Datacenter Network**

### **Small Latency**

< 100 *μs* 



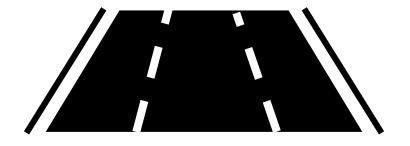
### **Shallow Buffer**

< 30 MB for ToR



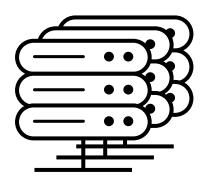
### **High Bandwidth**

10/40 ~ 100 Gbps



### Large Scale

> 10,000 machines



### **Datacenter Network**

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< 100 μs

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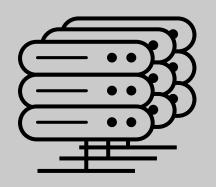
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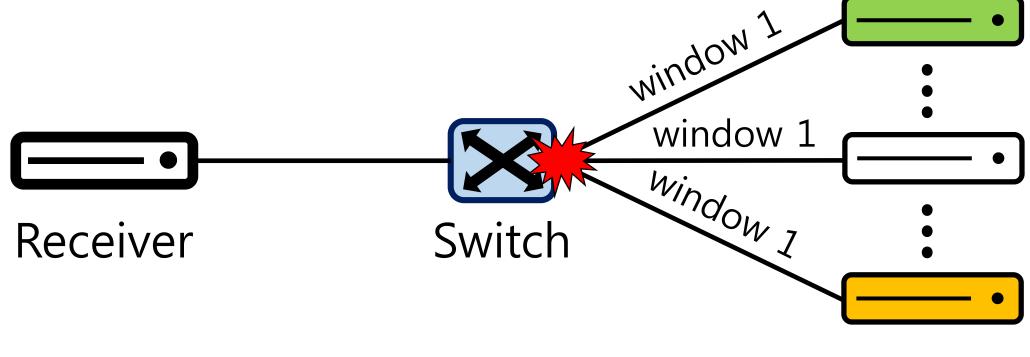
### Large Scale

> 10,000 machines



## Challenge with small BDP

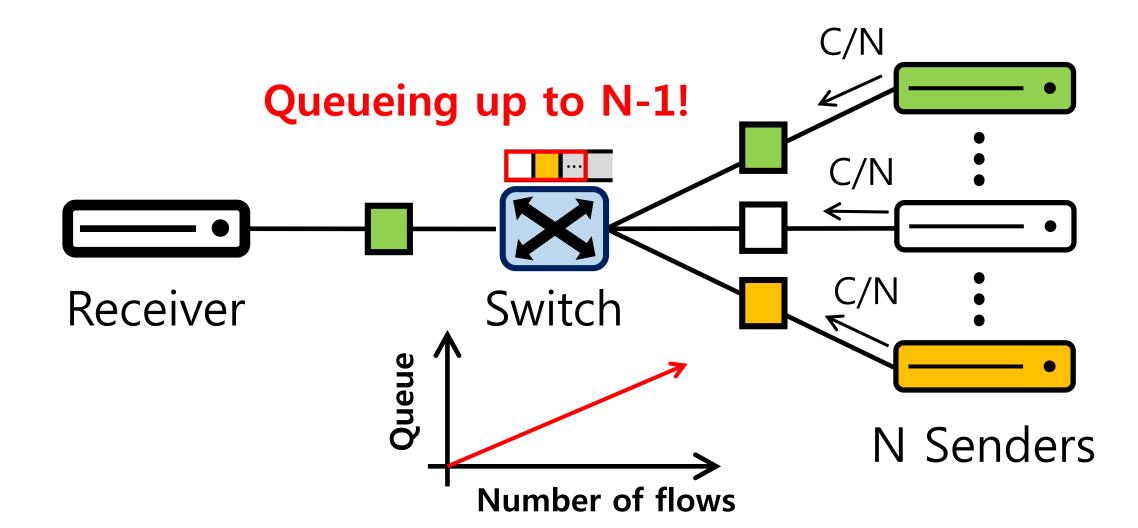
BDP\*( $100\mu s$ , 40Gbps)  $\approx 300$  MTUs



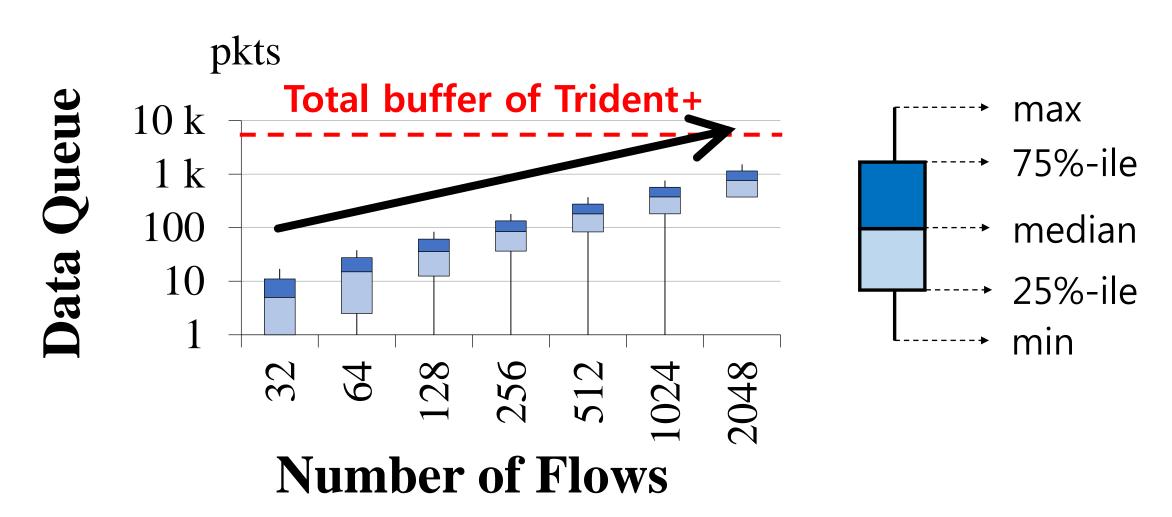
N Senders

\* BDP: Bandwidth-delay Product

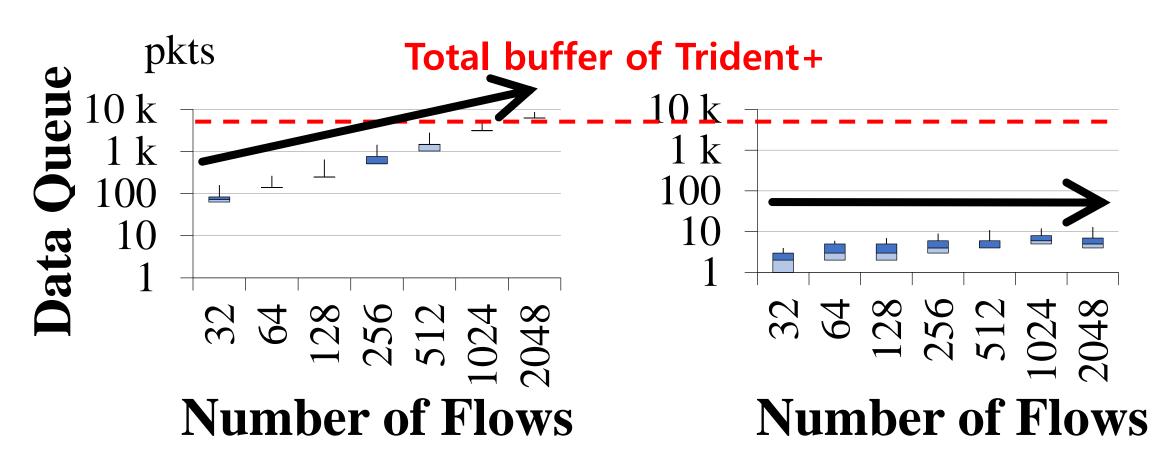
### Rate-based CC + incast traffic



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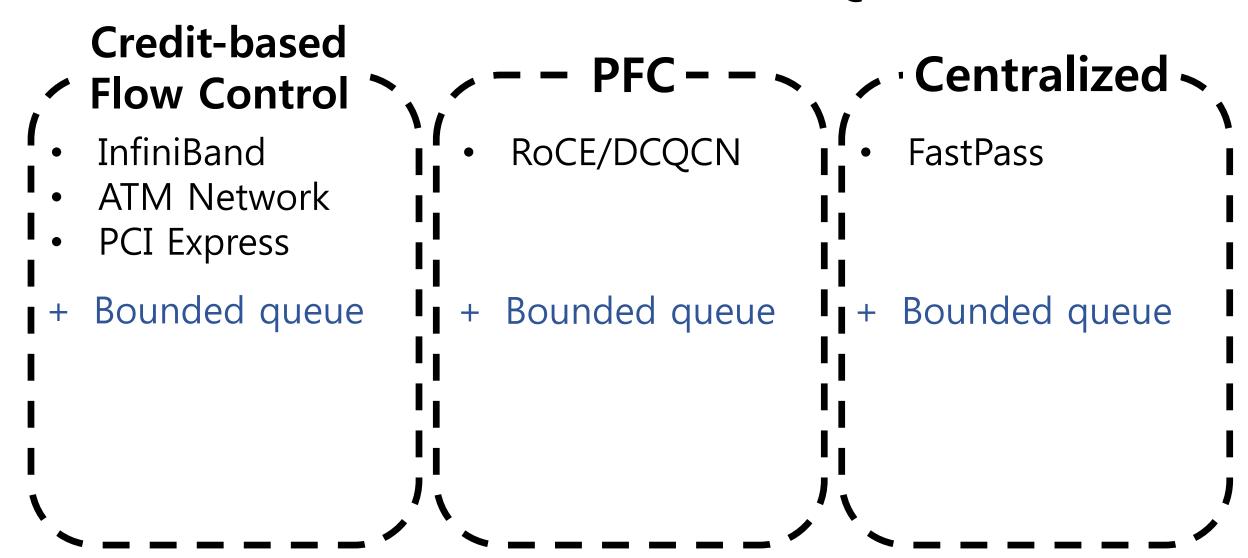
### Rate-based CC vs. credit-based CC



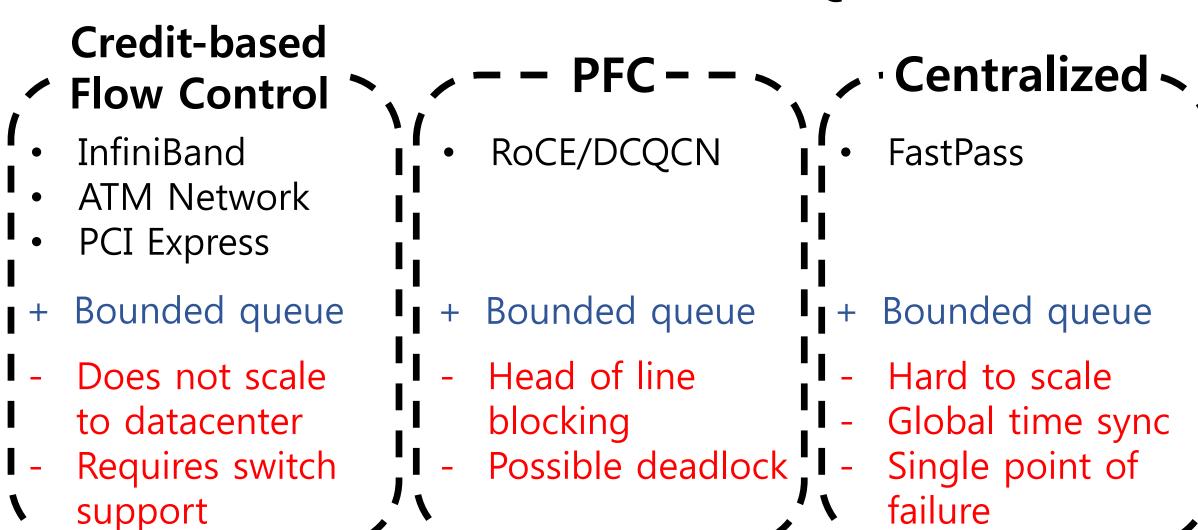
**DCTCP** 

**Credit-based Approach** 

## Prior Work with Bounded Queue



# Prior Work with Bounded Queue



# Prior Work with Bounded Queue

- Credit-based
  Flow Control
  InfiniBand
  RoCE/DCQCN
  FastPass
- How can we get the benefits of credit-based flow control on Ethernet?
- Does not scale
  to datacenter
  Requires switch
  support
  Head of line
  Global time sync
  Single point of failure

## Goal & Our Approach

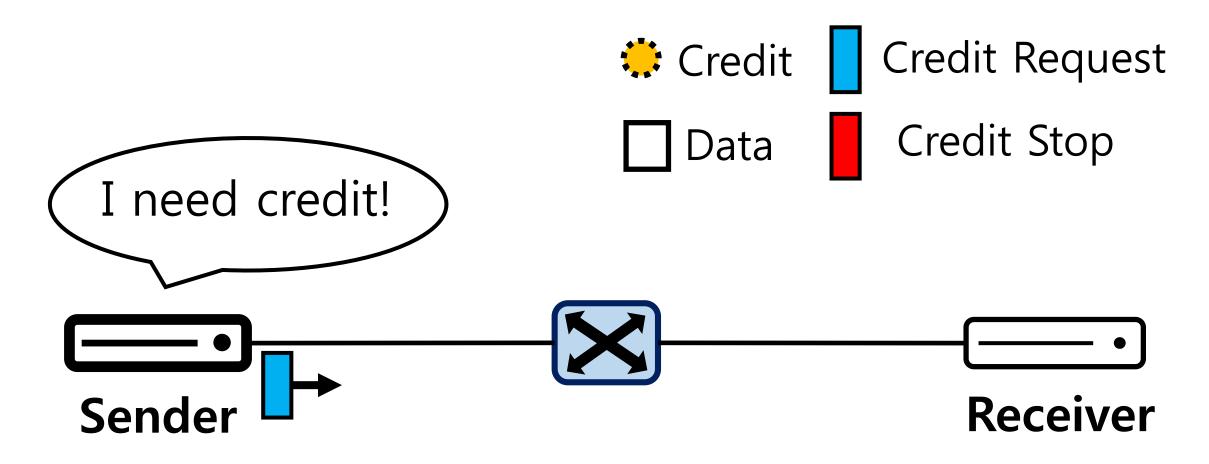
### Goal

To achieve **bounded queue** even with heavy incast using **Ethernet switches**.

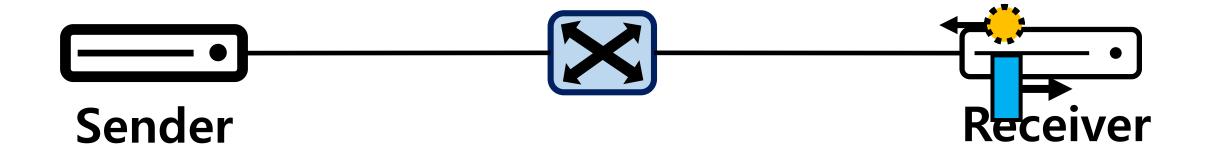
### **ExpressPass**

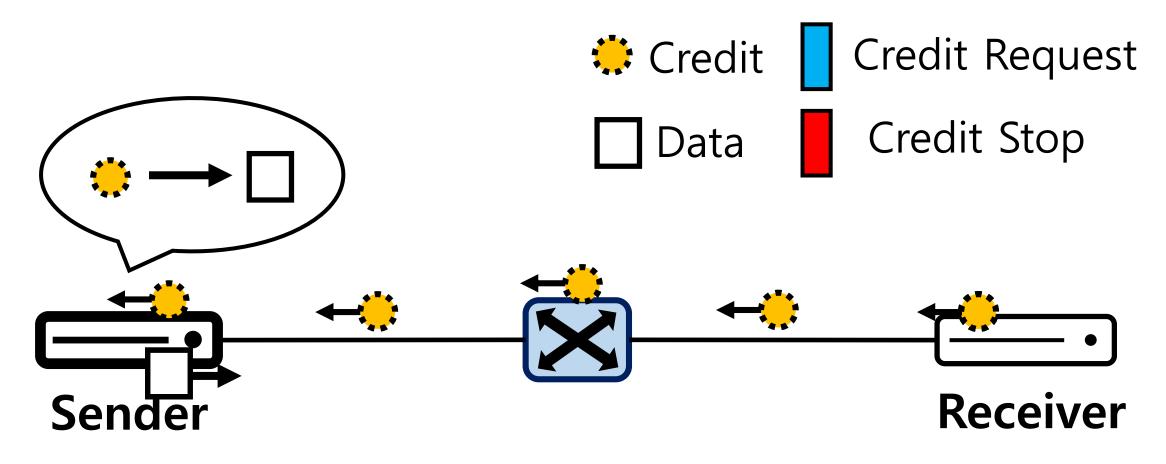
Proactive end-to-end credit-based congestion control using unreliable credits.

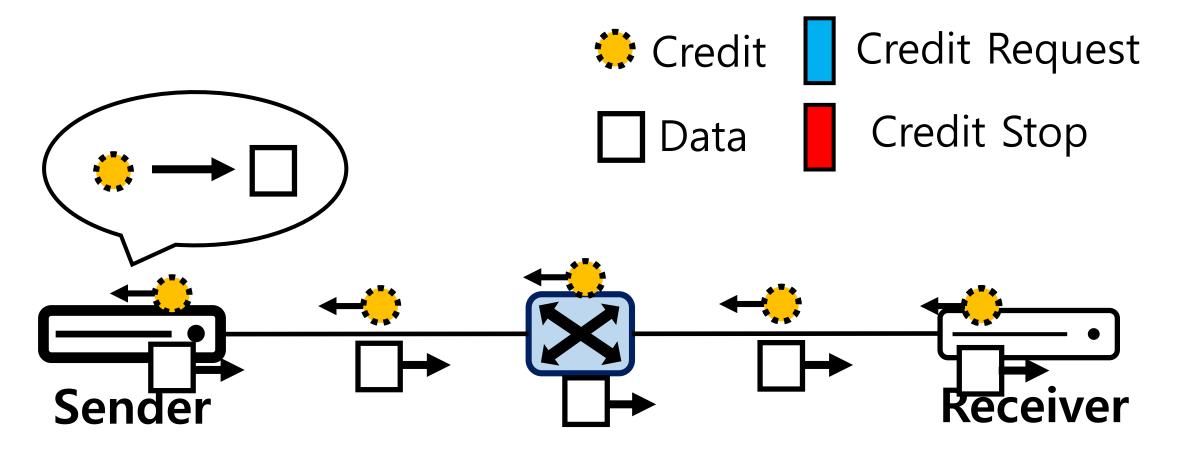
# **ExpressPass**End host behavior

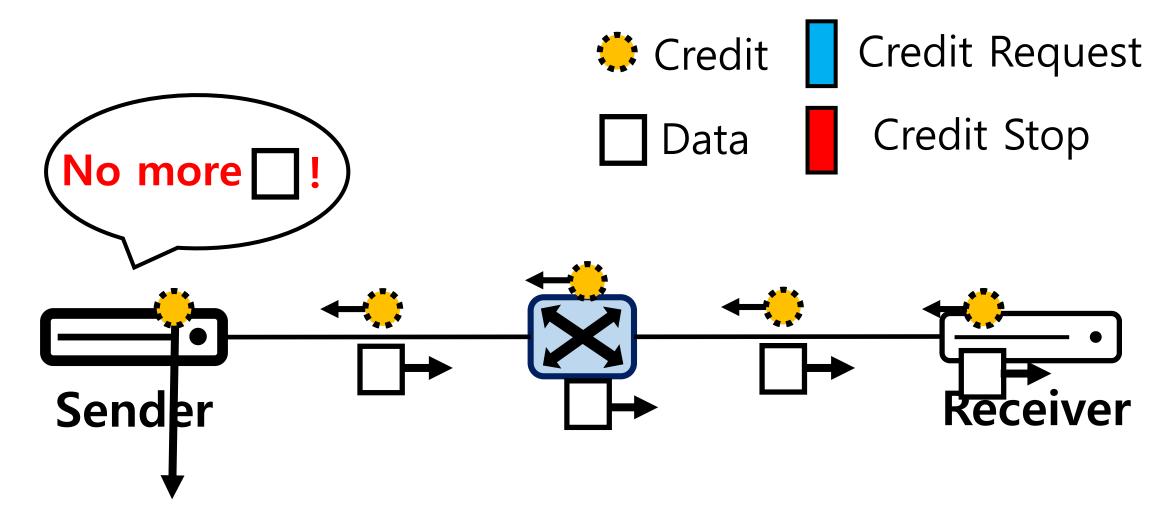


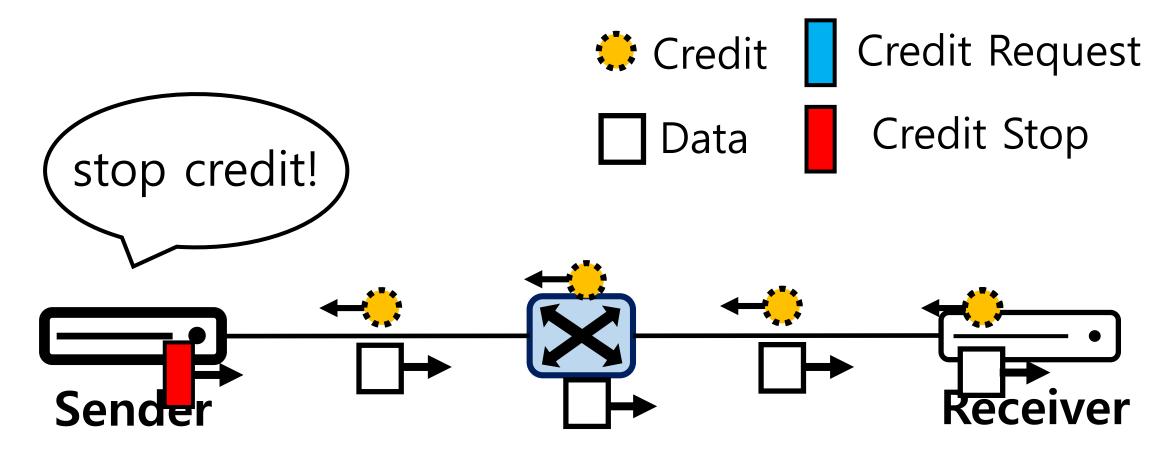
# **ExpressPass**End host behavior

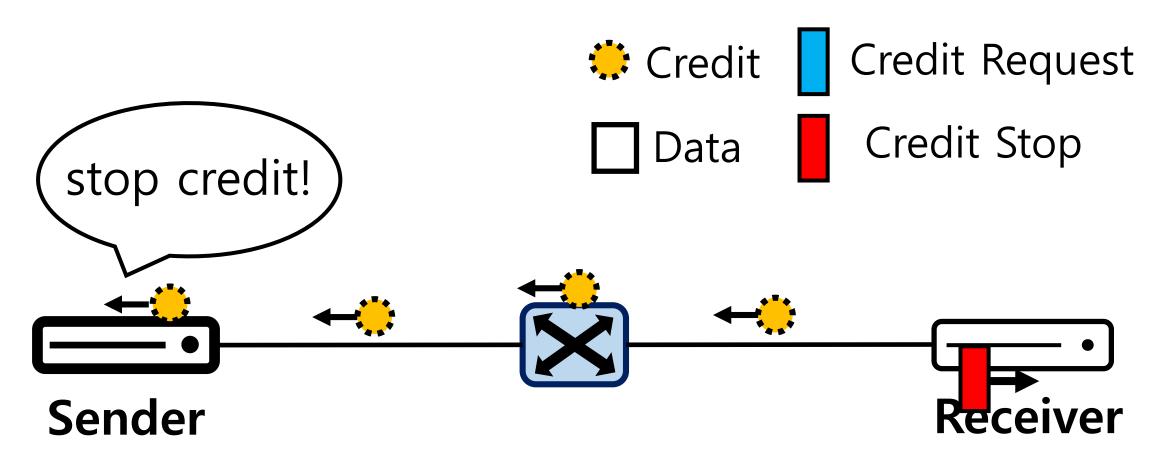




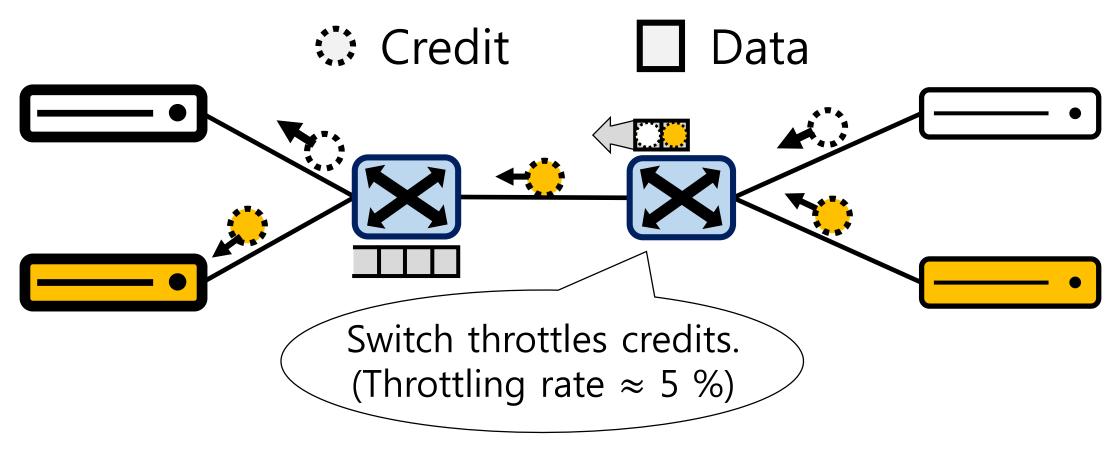








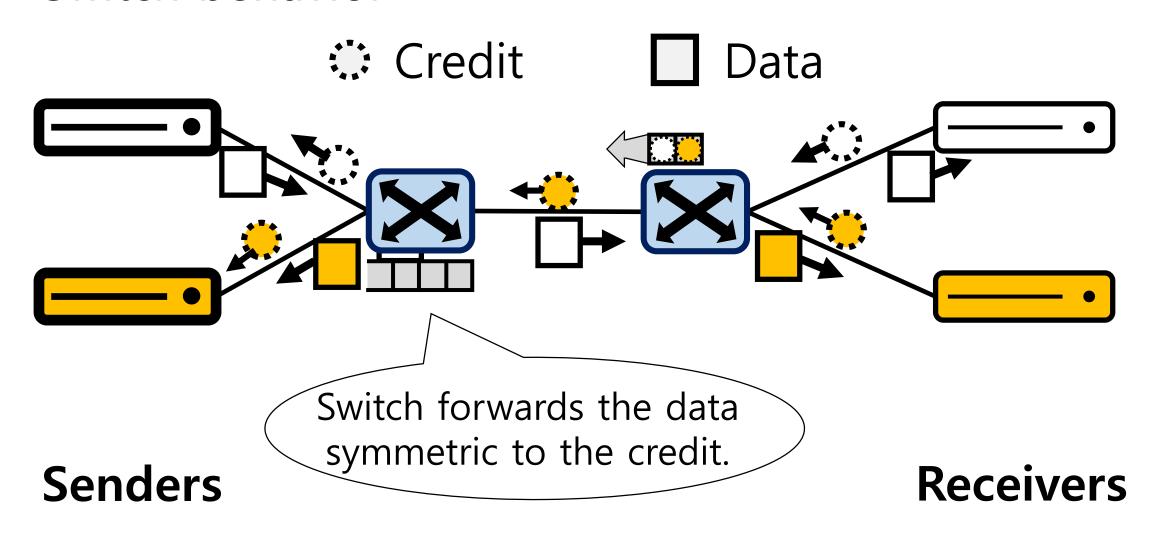
Switch behavior



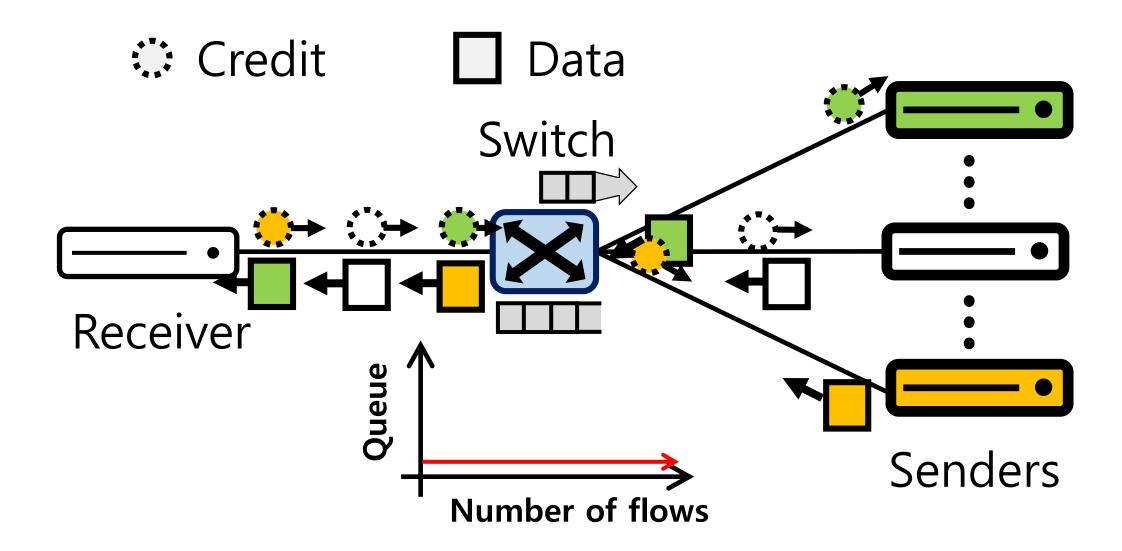
**Senders** 

Receivers

Switch behavior



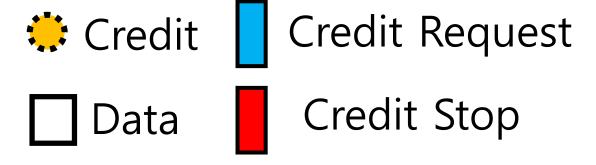
### Credit-scheduled data transmission

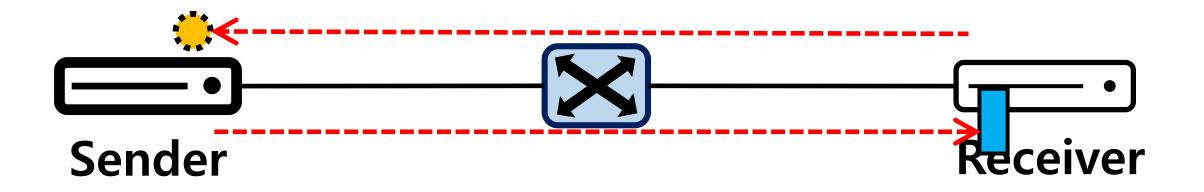


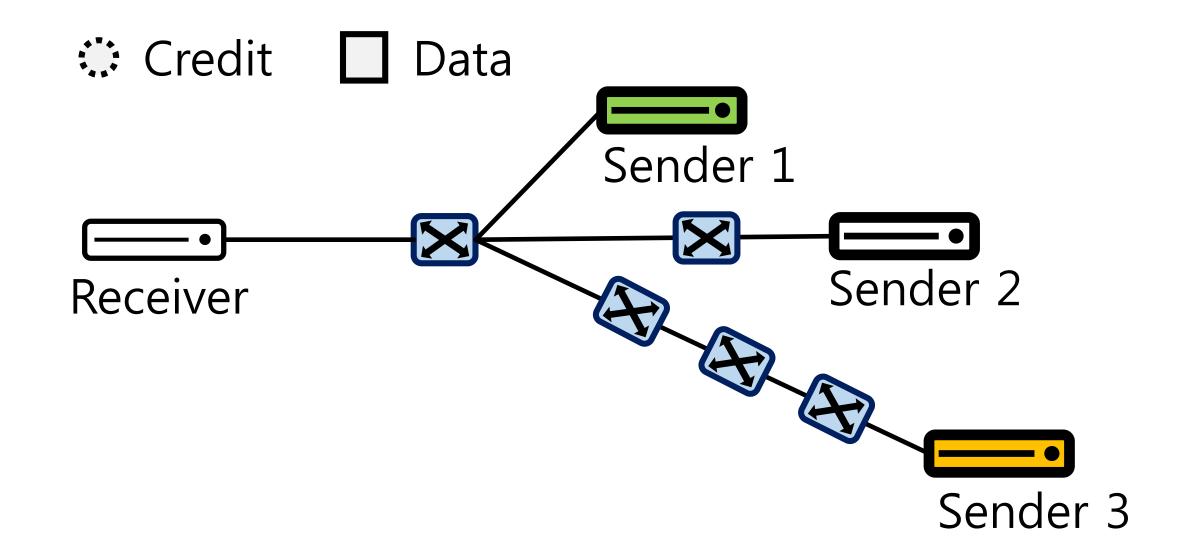
# Challenges

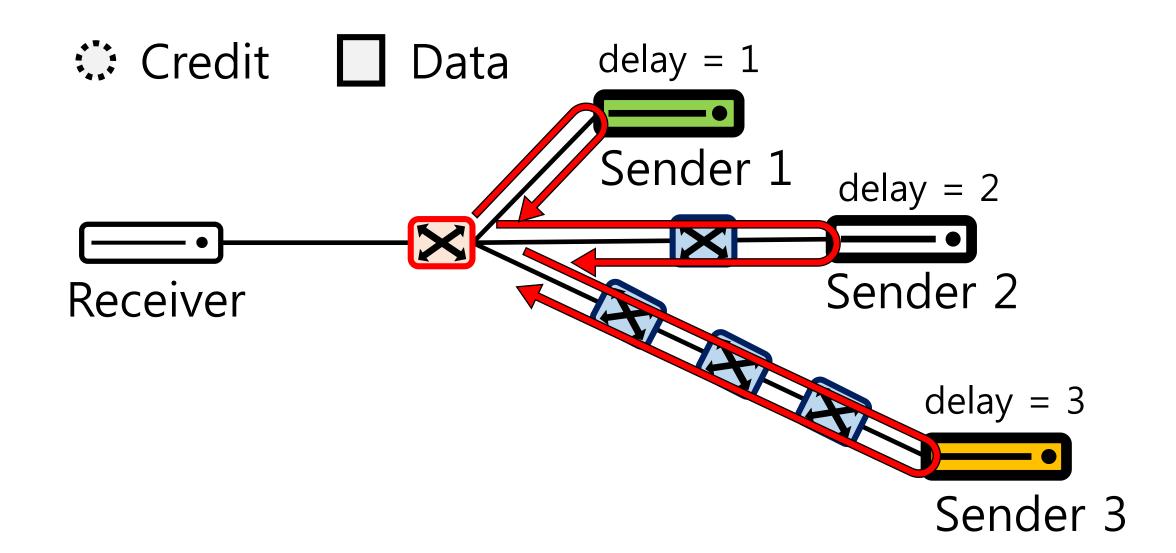
	Challenges	Techniques to address		
٢	Signaling overhead	Piggybacking to handshake packets		
$\left\{ \right.$	Non-zero queueing	Bounded queue		
	Credit waste	Credit feedback control		
	Fair drop on switch	Jitter, variable-sized credits		
	Path symmetry	Deterministic ECMP, packet level loa balancing		
	Multiple traffic classes	Prioritizing credits rather than data		

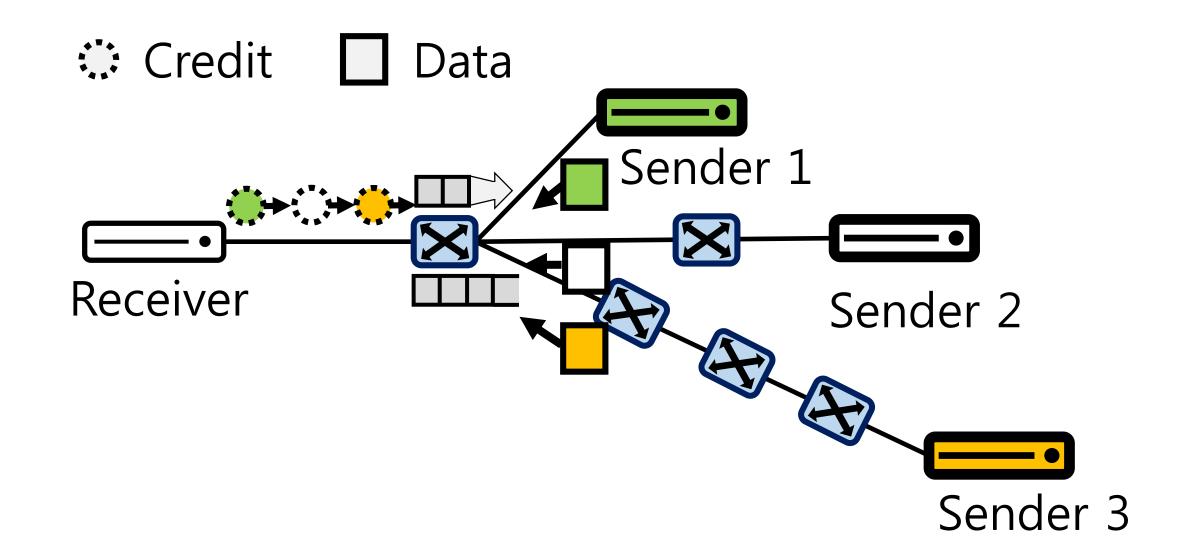
# Signaling Overhead

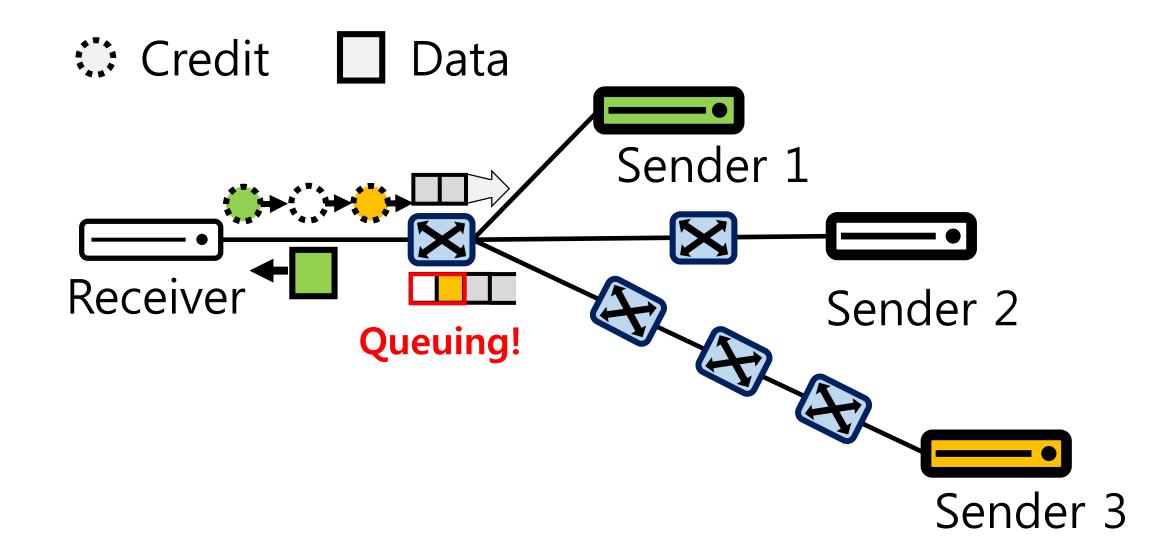




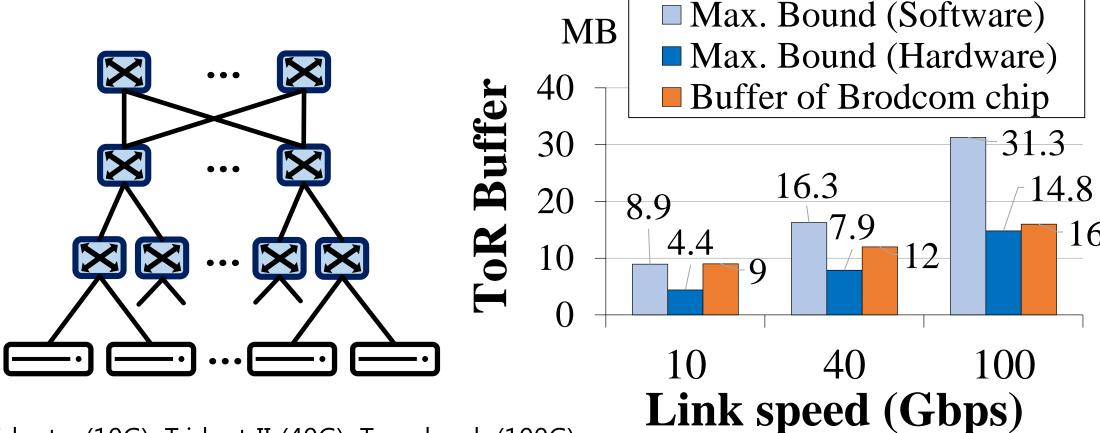






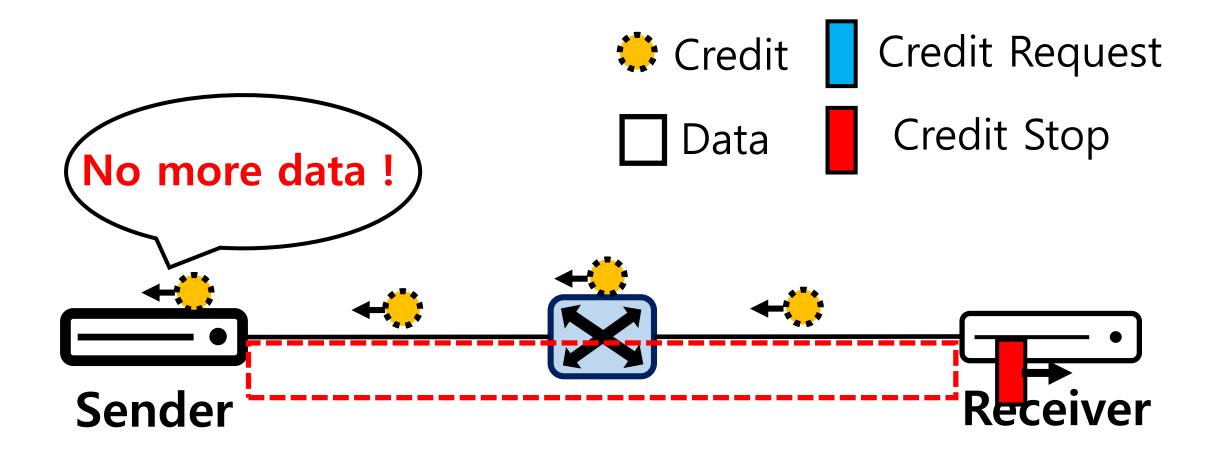


 $\max(buffer) = C * \{\max(delay) - \min(delay)\}$ 



<sup>\*</sup> Trident+ (10G), Trident II (40G), Tomahawk (100G)

### **Credit Waste**



### **Credit Feedback Control**

#### **Proactive Congestion Control**

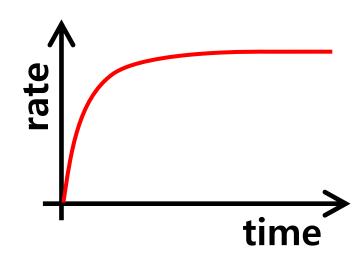
Prevents the congestion <u>before</u> actual congestion happens using credits.

#### Cheap credit drop

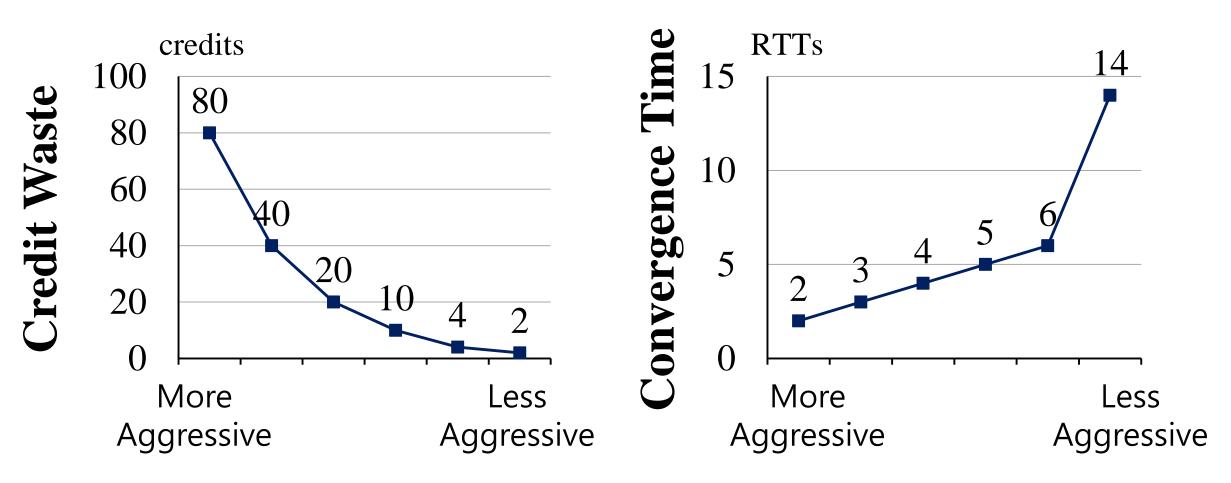
We can increase rate aggressively.

Bandwidth probing is cheap.

Convergence can be faster.



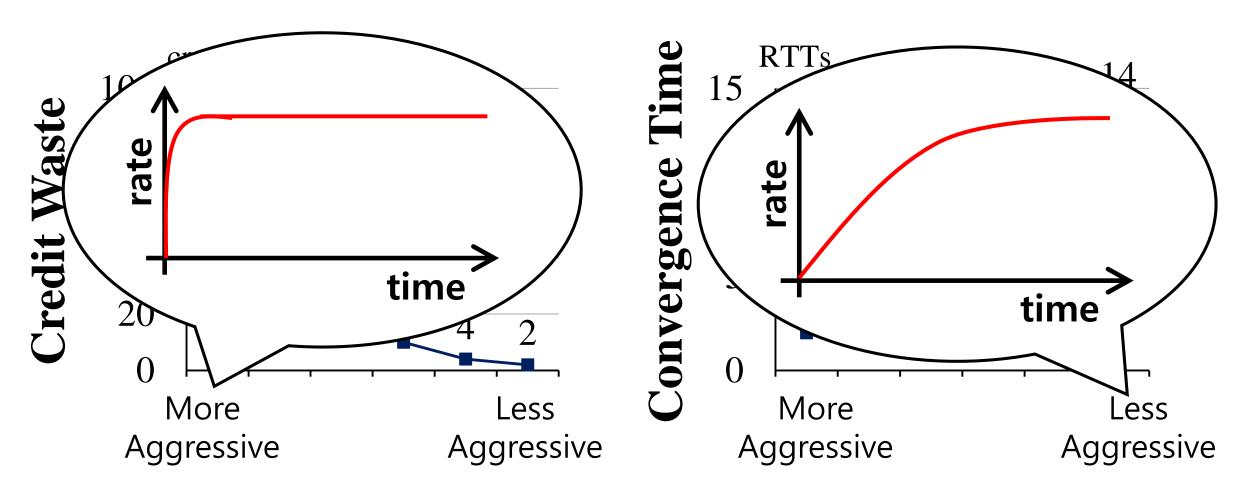
# Credit Waste & Convergence Time



Level of Aggressiveness

Level of Aggressiveness

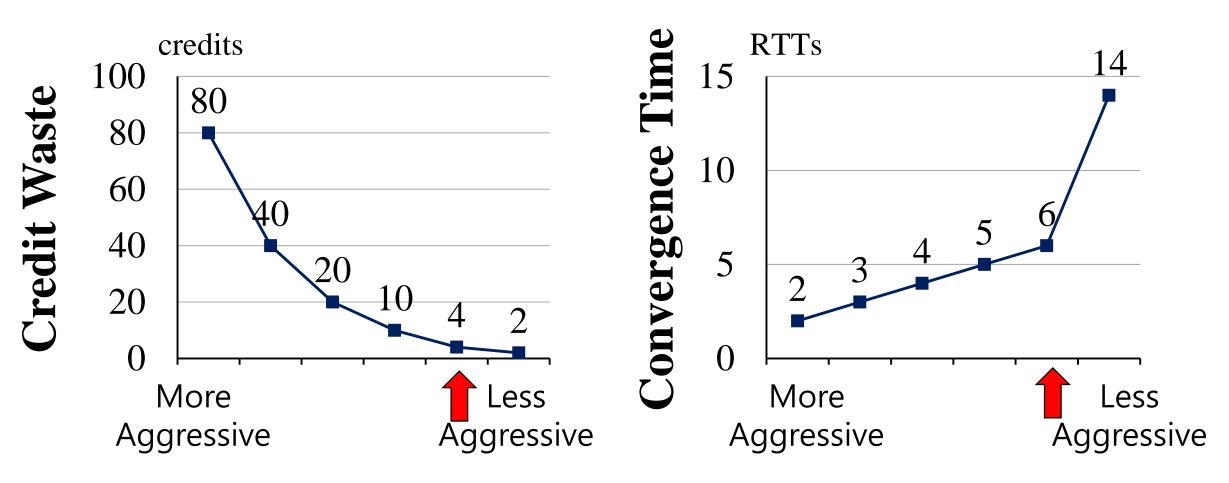
# Credit Waste & Convergence Time



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# Credit Waste & Convergence Time



Level of Aggressiveness

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# **Evaluation Setup**

### **Testbed setup**

- Dumbbell topology
- Implementation on SoftNIC
- 12 hosts (Xeon E3/E5) connected to single ToR (Quanta T3048)
- Each host has 10Gbps x 1port

### **NS-2 Simulation Setup**

- Fat-tree topology
- 192 hosts / 32 ToR / 16 aggr. / 8 core switches
- Each host has 10Gbps x 1port

### **Evaluation**

- (1) Does ExpressPass provides low & bounded queueing with realistic workloads?
- (2) Is the convergence fast and stable?
- (3) How low & bounded queuing and fast & stable convergence translate into the flow completion time?

### **Realistic Workloads**

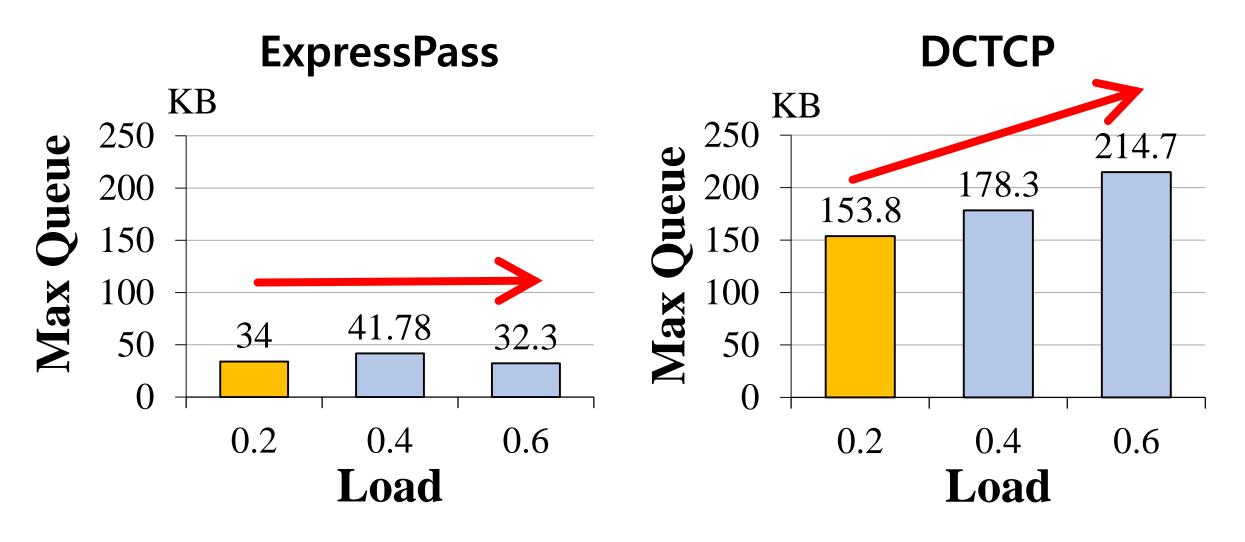
	Data Mining	Web Search	Cache Follower	Web Server
0 – 10KB (S)	78%	49%	50%	63%
10 - 100KB (M)	5%	3%	3%	18%
100KB-1MB (L)	8%	18%	18%	19%
1MB- (XL)	9%	20%	29%	-
Average flow size	7.41MB	1.6MB	701KB	64KB

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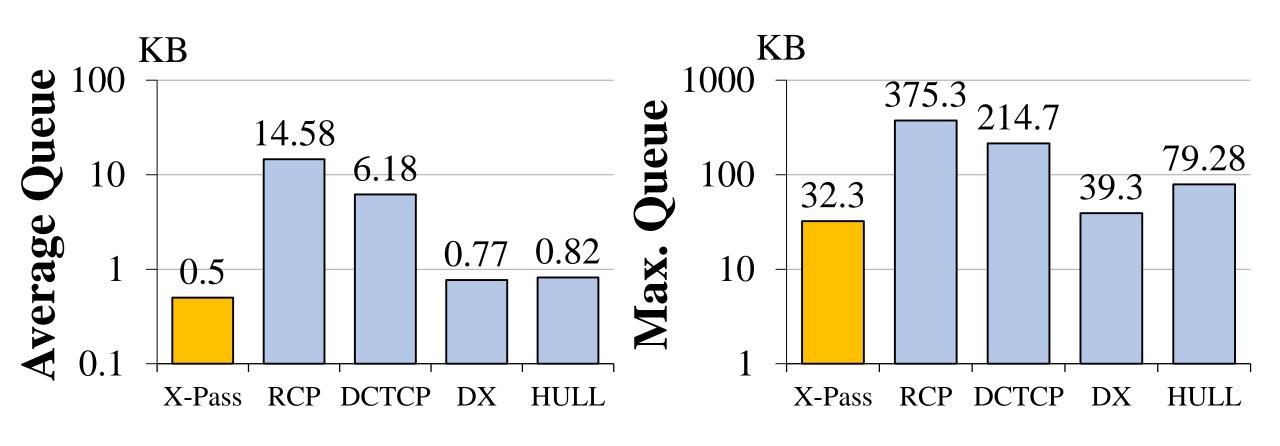
#### **Bounded Queue**

cache follower workload / load 0.2 - 0.4 / 0KB ~ (All Size)



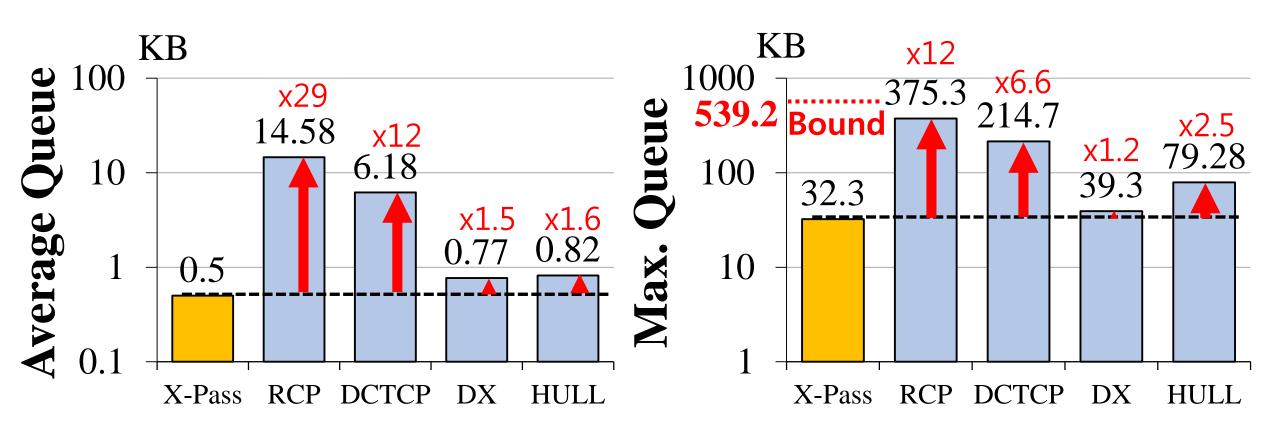
#### Low Average Queue

cache follower workload / load 0.6 / 0KB -

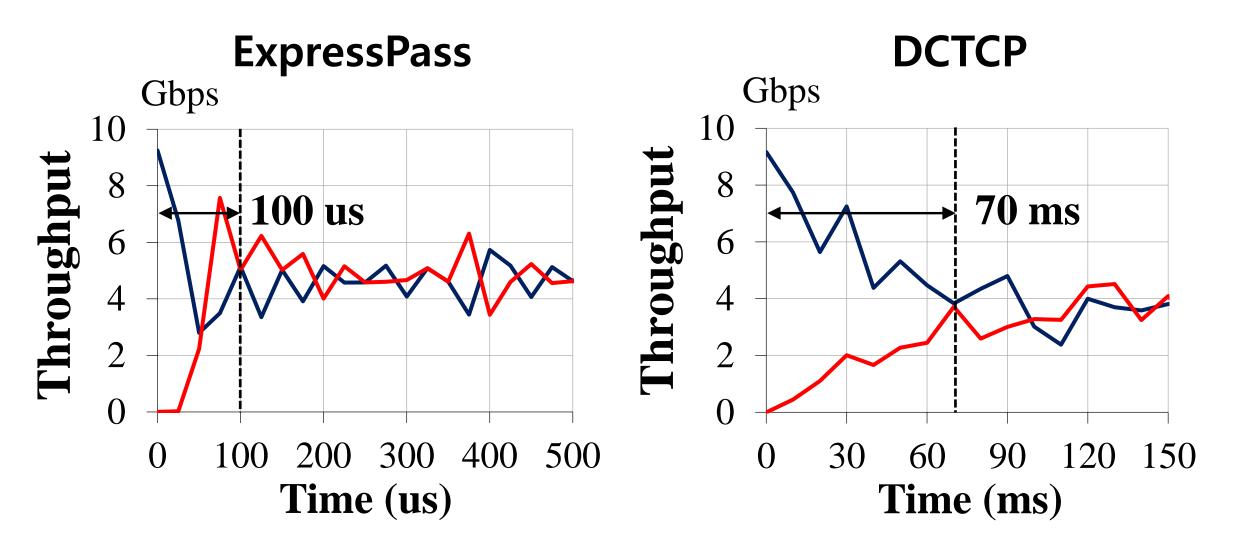


#### Low Average Queue

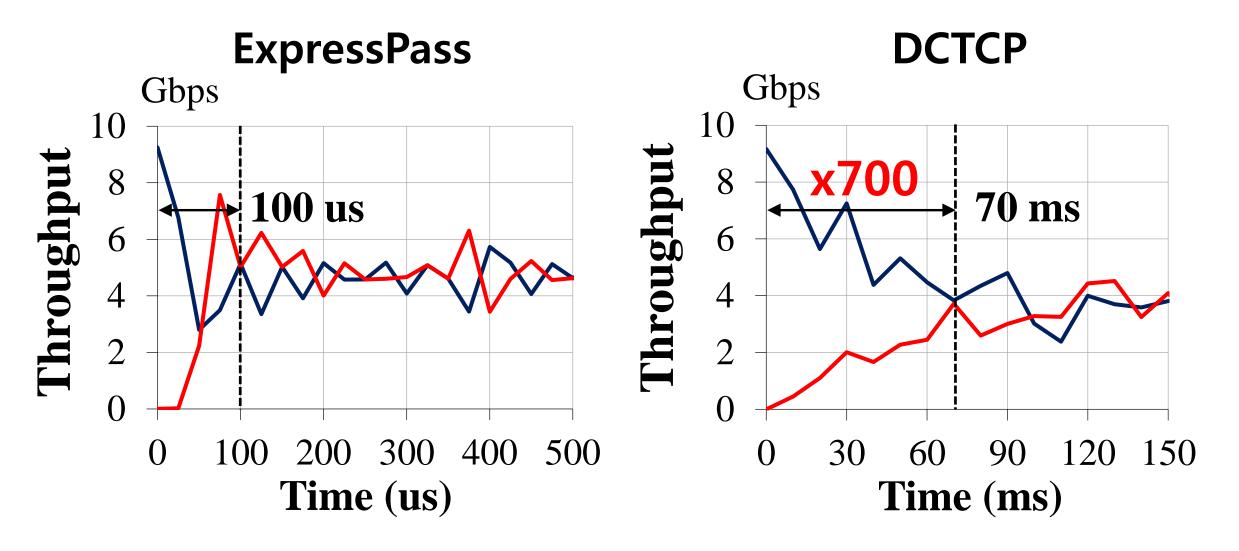
cache follower workload / load 0.6 / 0KB -



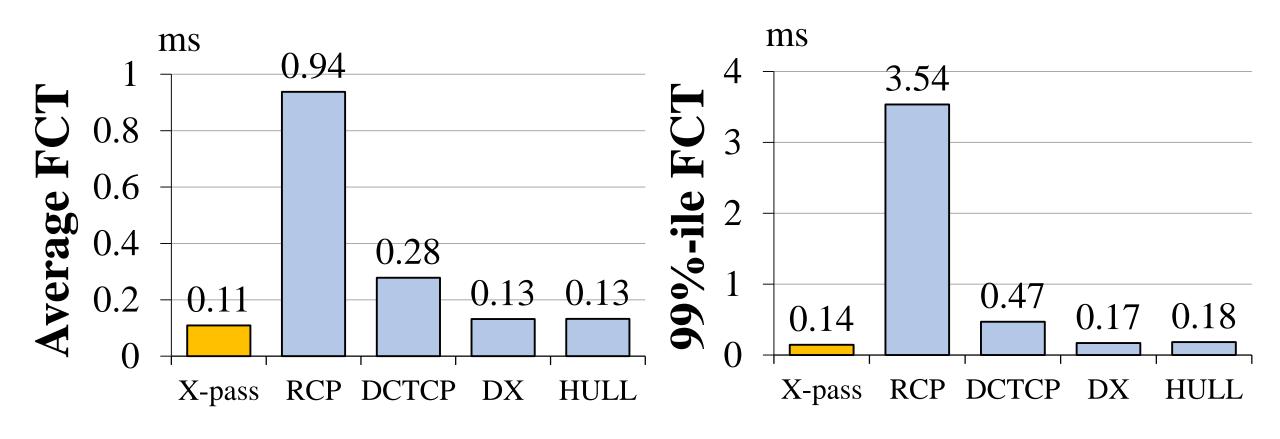
#### Fast & Stable Convergence



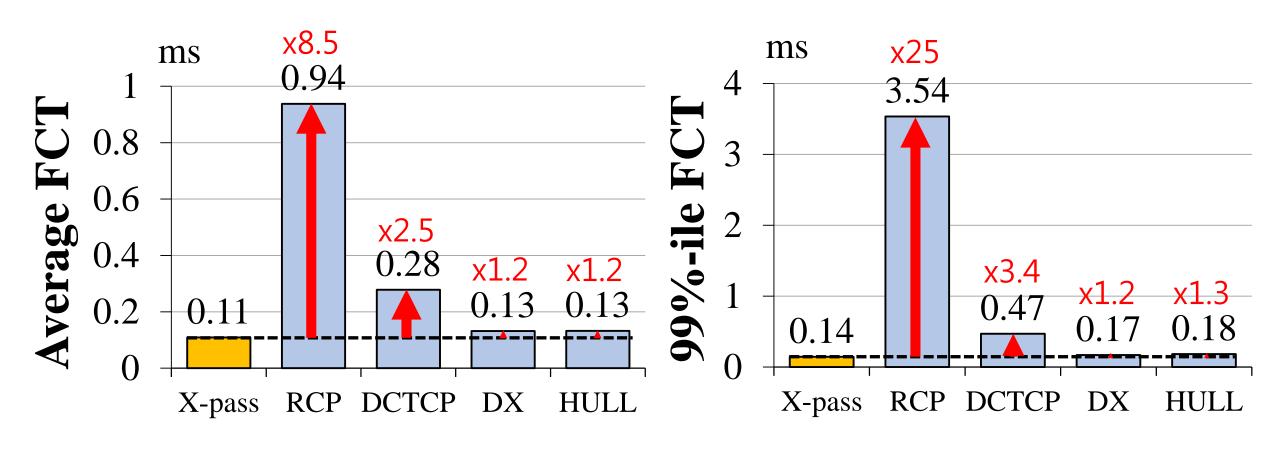
### Fast & Stable Convergence



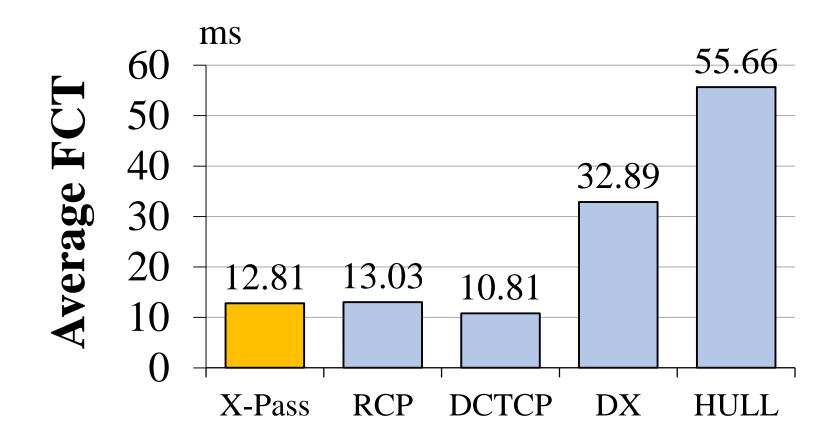
cache follower workload / load 0.6 / 0 - 10KB



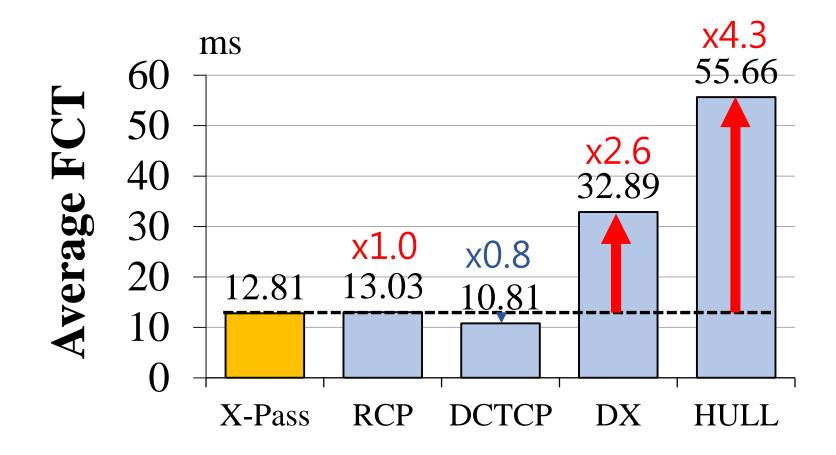
cache follower workload / load 0.6 / 0 - 10KB



cache follower workload / load 0.6 / 1MB -



cache follower workload / load 0.6 / 1MB -



#### Conclusion

- ExpressPass is end-to-end, credit-scheduled, and delay-bounded congestion control for datacenter.
- ExpressPass propose a new **proactive** datacenter congestion control.
- Our evaluation on testbed and ns-2 simulation show that ExpressPass achieves
  - (1) Low & bounded queueing
  - (2) Fast & stable convergence
  - (3) Short flow completion time especially for small flows

# **Thanks**

Happy to answer your questions