

A Concise Horn Theory for RCC8

Matthias Westphal and Julien Hué¹

Abstract. RCC8 is a well-known constraint language for expressing and reasoning about spatial knowledge. We state a simple and concise Horn theory for RCC8 analogous to the ORD-Horn theory for temporal reasoning. This theory allows for expressing RCC8 and retains tractability of the well-known Horn reduct of RCC8. Further, it is much more adequate for practical purposes in the area of logic programming and surpasses previous attempts.

1 BACKGROUND

We assume the reader is familiar with standard FOL concepts, notation in the CSP literature, and syntactic interpretations [2, 6]. A constraint language Γ is simply a relational FO structure. The *Constraint Satisfaction Problem* for Γ , $\text{CSP}(\Gamma)$, is the decision problem if some given primitive positive FO sentence formulated over the signature of Γ is satisfiable in Γ . An algorithm *solves* $\text{CSP}(\Gamma)$ if it solves the decision problem.

The spatial reasoning formalism RCC8 [3] describes binary relations between abstract spatial regions. There are 8 relation primitives: equal EQ , disconnected DC , externally connected EC , partially overlaps PO , non-tangential proper part $NTPP$, tangential proper part TPP , and the converses $NTPPI$, $TPPI$. For example, the formalism allows to formally state that “Central Park is a (non-tangential proper) part of Manhattan”. These eight relations are denoted in the signature σ_{RCC8A} of atomic RCC8 relation symbols. Further, RCC8 contains all disjunctions of these relation symbols forming the signature σ_{RCC8} of 256 relations. We denote by Γ the constraint language RCC8 on signature σ_{RCC8} constructed as an ω -categorical structure as in [2].

There are two important reducts of Γ : the σ_{RCC8A} -reduct Γ_{RCC8A} , and the $\sigma_{\hat{\mathcal{H}}_8}$ -reduct $\Gamma_{\hat{\mathcal{H}}_8}$. Symbols of $\sigma_{\hat{\mathcal{H}}_8}$ denote the RCC8 Horn relations. Renz and Nebel [4] gave a transformation of any instance of $\text{CSP}(\Gamma_{\hat{\mathcal{H}}_8})$ with n FO variables into an equivalent propositional Horn formula with $O(n^4)$ literals. Well-known results (e.g. [4]) are: $\sigma_{\text{RCC8A}} \subsetneq \sigma_{\hat{\mathcal{H}}_8} \subsetneq \sigma_{\text{RCC8}}$, $|\sigma_{\hat{\mathcal{H}}_8}| = 148$, and strong 3-consistency solves $\text{CSP}(\Gamma_{\hat{\mathcal{H}}_8})$. The reduct $\Gamma_{\hat{\mathcal{H}}_8}$ is the largest maximal tractable subclass of RCC8 that includes all atomic relations. Further, $\text{CSP}(\Gamma_{\hat{\mathcal{H}}_8}) \in \text{P}$ while $\text{CSP}(\Gamma)$ is NP-complete.

We present a syntactic interpretation of RCC8 and a concise Horn theory for solving $\text{CSP}(\Gamma_{\hat{\mathcal{H}}_8})$ in polynomial time. The theory is weaker than strong 3-consistency. The Herbrand expansion yields propositional Horn formulas with $O(n^2)$ literals and $O(n^3)$ clauses.

2 A CONCISE HORN THEORY

Bennett discussed a number of different representations for logical reasoning about spatial relations [1]. Among them is the RCC7 rep-

resentation which consists of relations representable as equations in interior algebra. These 7 relations are: equality EQ , disconnectedness DC , discreteness DR , parthood P , non-tangential parthood NTP , and the converses P^{-1} , NTP^{-1} . Bennett found 115 distinct relations to be definable using purely conjunctive formulas if the complements of the relations are included.

In the following we introduce a set of relations closely related to RCC7 and its use for defining further relations. Here, EQ is not included as a primitive because it can be defined using parthood. Further, we drop converse symbols and use the complements of DC and DR : connectedness C and overlap O . We refer to this set as RCC4. Its signature can be thought of as a shorthand notation of specific RCC8 relation symbols, i.e., we treat σ_{RCC4} as a proper subset of σ_{RCC8} .

Definition 1 (RCC4) In the binary signature $\sigma_{\text{RCC4}} := \{C, O, NTP, P\}$ NTP denotes the relation symbol $NTPP$ in σ_{RCC8} , C the relation symbol representing $\neg DC(x, y)$ over Γ , O the relation symbol representing $\neg(DC(x, y) \vee EC(x, y))$, and P the symbol representing $TPP(x, y) \vee NTPP(x, y) \vee EQ(x, y)$. Further, Γ_{RCC4} denotes the σ_{RCC4} -reduct of Γ .

The structure Γ_{RCC4} admits a FO interpretation of Γ . We first give the syntactic interpretation of the atomic RCC8 relation symbols (virtually the same map as provided in Table 5.4 of [1]).

Definition 2 (π_{RCC4}) The syntactic interpretation π_{RCC4} of atomic RCC8 relation symbols in the RCC4 symbols is defined as follows:

$$\begin{aligned} \pi_{\text{RCC4}}(DC)(x, y) &:= \neg C(x, y) \\ \pi_{\text{RCC4}}(EC)(x, y) &:= C(x, y) \wedge \neg O(x, y) \\ \pi_{\text{RCC4}}(PO)(x, y) &:= O(x, y) \wedge \neg P(x, y) \wedge \neg P(y, x) \\ \pi_{\text{RCC4}}(EQ)(x, y) &:= P(x, y) \wedge P(y, x) \\ \pi_{\text{RCC4}}(TPP)(x, y) &:= P(x, y) \wedge \neg P(y, x) \wedge \neg NTP(x, y) \\ \pi_{\text{RCC4}}(NTPP)(x, y) &:= NTP(x, y) \\ \pi_{\text{RCC4}}(TPPI)(x, y) &:= \pi_{\text{RCC4}}(TPP)(y, x) \\ \pi_{\text{RCC4}}(NTPPI)(x, y) &:= \pi_{\text{RCC4}}(NTPP)(y, x) \end{aligned}$$

It is easy to verify that this syntactic interpretation provides an interpretation of Γ_{RCC8A} in Γ_{RCC4} . The map can be easily extended to all RCC8 relations, thus Γ has an interpretation in Γ_{RCC4} . Note, although RCC4 is very close to RCC7, the formula $\neg EQ(x, y)$ is written as a disjunction over σ_{RCC4} – contrary to RCC7 where EQ is a primitive. From the RCC subsumption lattice [3] it follows:

Proposition 1 Let ψ be a formula $L_O(x, y) \wedge L_C(x, y) \wedge L_P(x, y) \wedge L_{NTP}(x, y) \wedge L_P(y, x) \wedge L_{NTP}(y, x)$ in which $L_R(x, y) = R(x, y)$ or $L_R(x, y) = \neg R(x, y)$. Then ψ is either unsatisfiable in Γ_{RCC4} or equivalent to an atomic RCC8 relation.

For convenience we consider the expansion of Γ_{RCC4} to $\Gamma_{\text{RCC4}+}$ by adding a symbol \bar{R} for the complement of each relation symbol R

¹ Albert-Ludwigs-Universität Freiburg, Department of Computer Science, Germany, email: westpham@informatik.uni-freiburg.de, hue@informatik.uni-freiburg.de

in σ_{RCC4} (Γ_{RCC4+} is still a reduct of Γ). Then, π_{RCC4} can be seen as primitive positive. In the following, we write π_{RCC4+} to denote the syntactic interpretation of each RCC8 relation symbol over Γ_{RCC4+} that uses CNF defining formulas consisting of prime implicates.

Computer analysis shows the following properties.

Proposition 2 *There are 42 RCC8 relations which have a primitive positive definition in Γ_{RCC4+} . All 148 RCC8 Horn relations have a Horn definition in Γ_{RCC4} .*

In the following we give a Horn theory over Γ_{RCC4} written as the Datalog program Π_{RCC4+}^{RCC4+} with symbols in σ_{RCC4+} and a special symbol false denoting \perp . The program contains the necessary subsumption rules and rules on 3 variables that derive from the axiomatization of the atomic composition table of RCC8 given by Bodirsky and Wöfl [2] using π_{RCC4+} and simplifications. It solves $\text{CSP}(\Gamma_{RCC4+})$ by Proposition 1 and the results of [2].

$P(x, x)$	$\overline{NTP}(x, x)$
$C(x, y) :- O(x, y)$	$O(x, y) :- P(x, y)$
$P(x, y) :- NTP(x, y)$	$\text{false} :- P(y, x), NTP(x, y)$
$O(x, y) :- O(y, x)$	$C(x, y) :- C(y, x)$
$\text{false} :- C(x, y), \overline{C}(x, y)$	$\text{false} :- NTP(x, y), \overline{NTP}(x, y)$
$\text{false} :- O(x, y), \overline{O}(x, y)$	$\text{false} :- P(x, y), \overline{P}(x, y)$
$P(x, y) :- P(x, z), P(z, y)$	$C(x, y) :- C(x, z), P(z, y)$
$O(x, y) :- C(x, z), NTP(z, y)$	$O(x, y) :- O(x, z), P(z, y)$
$NTP(x, y) :- P(x, z), NTP(z, y)$	$NTP(x, y) :- NTP(x, z), P(z, y)$

3 PROPERTIES AND EVALUATION

Analogous to the study of temporal formalisms in [6], we can analyze the relation of this new program to commonly used ones: strong 3-consistency ($\Pi_{\leq 3}^{RCC8}$), the atomic composition table ($\Pi_{\leq 3}^{RCC8A}$), and the augmented Π_{RCC4+}^{RCC4+} that emulates positive unit resolution on Horn clauses in π_{RCC4+} (Π_{RES}^{RCC4+}). Additionally, each of them can be weakened by restricted inference to arcs of an instance that are part of some fixed chordal graph of an instance (denote by $|\hat{G}$). Finally, we can define a strict partial order on the programs, where $A \prec B$ if program A rules out more partial solutions than program B . For details, references, and discussion of these concepts see [6].

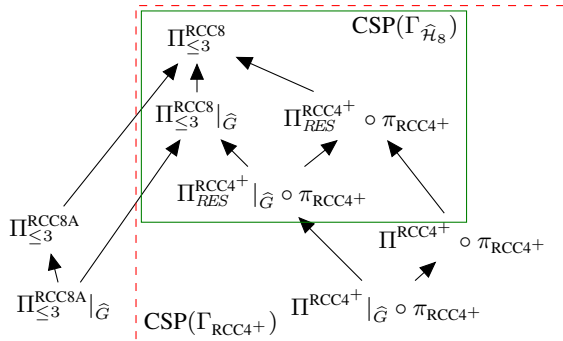


Figure 1. Lattice of propagation strength of Datalog programs given interpretations for RCC8. Frames indicate CSPs solved by the programs.

The lattice for \prec is given in Figure 1 (where $A \leftarrow B$ if $A \prec B$) which also states which CSPs of reducts are solved by the programs.

In order to encode instances to propositional CNF it has been common practice to consider the *support encoding* of $\Pi_{\leq 3}^{RCC8A} | \hat{G}$ (see [6] for references). However, this encoding is quite weak.

Proposition 3 *Unit propagation on the support encoding of $\Pi_{\leq 3}^{RCC8A}$ does not solve $\text{CSP}(\Gamma_{RCC8A})$, while unit propagation on the Herbrand expansion of $\Pi_{RES}^{RCC4+} | \hat{G} \circ \pi_{RCC4+}$ solves $\text{CSP}(\Gamma_{\hat{H}_8})$.*

If (a) the support encoding of $\Pi_{\leq 3}^{RCC8A} | \hat{G}$ is used SAT solvers have to rely on search to refute instances, but when (b) the Herbrand expansion of $\Pi_{RES}^{RCC4+} | \hat{G} \circ \pi_{RCC4+}$ is used search is only necessary on non-Horn relations of instances in $\text{CSP}(\Gamma)$. Further, both encodings result in propositional CNF with $O(n^2)$ literals and $O(n^3)$ clauses for instances with n FO variables, but interpretation and program of the new concise theory have less relation symbols and rules.

For a brief evaluation we consider 1000 random instances of $\text{CSP}(\Gamma)$ constructed from the $H(150, 15.5, 4.0)$ model with hard relations \mathcal{NP}_8 [5]. These instances have 150 FO variables and a constraint network degree of 15.5. The test setup is the same as in [6].²

SAT Encoding	Runtime (seconds) of the Glucose-2.2 SAT solver				
	Avg.	25-PCTL	50-PCTL	90-PCTL	99-PCTL
(a)	440.90	227.05	349.48	545.62	1 966.94
(b)	45.87	8.92	18.69	103.75	441.32

This is a factor 4 to 25 improvement. Further, the average size of an instance in DIMACS drops from (a) 360 MiB to (b) 160 MiB.

4 CONCLUSION

We presented a concise Horn theory for RCC8 that allows for expressing the entire RCC8 language and retains the tractability of the well-known Horn class of relations. The theory is of broad interest for incorporating RCC8 in logic programming, e.g. SAT and ASP. In particular, the theory allows for an improved $O(n^3)$ propositional representation that is superior to previous work.

For future work it would be interesting to consider propagators based on this theory in Constraint Programming. In particular, a potential graph based algorithm could form a fast propagator.

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² The tool to compute the encoding is available; see this publication’s entry at <http://www.informatik.uni-freiburg.de/~ki/publications/>