

# Objective landscapes in CP Optimizer

! Ongoing work !

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- Introduction and notations
- Objective landscape definition
- Objective landscape properties
- Objective landscape computation
- Objective landscape exploitation
- Preliminary results

## Objective

**Improve performance** of CP Optimizer by exploiting some information on the relation between decision variables and objective value

Given a current solution  $S$  in the LNS, we would like to **estimate**:

- How much a variable  $x_i$  contributes to the cost
- The impact on the cost to modify the value of  $x_i$  by  $\epsilon$  (reduced cost)
- The ideal values (with respect to the cost) in the current domain of  $x_i$

This is useful to:

- Decide which fragment to relax
- Guide the completion by :
  - Enforcing improvement/non-degradation of the cost for some variables
  - Selecting next variable to fix
  - Fixing variables to values close to their ideal ones

- We need to **automatically** compute how decision variables are related with the objective function
- It must work for **continuous** variables with large domains (typically start/end of interval variables)
- This computation has to be **generic** (work for any objective function)
- This computation has to be **cheap**

~> **Objective landscape**

Comparison with **impacts** (on cost):

- **Impacts** are computed by changing the domain of a variable  $x_i$  and measuring how it impacts the cost function
- **Objective landscapes** work the other way round: the cost function is changed and it measures the impact on the domain of variables  $x_i$

Comparison with **impacts** (on cost):

- **Impacts** are mostly designed for discrete variables with reasonably small domains
- **Objective landscapes** are mostly designed for continuous domains; their complexity is not affected by domain size

Comparison with **impacts** (on cost):

- **Impacts** are continuously updated during the search and make a kind of *average* assumption about the value of other variables
- **Objective landscapes** are computed once and for all at the beginning of the search and make a kind of *optimistic* assumption about the value of other variables



Comparison with **temporal linear relaxation** :

- **Temporal linear relaxation**
  - Does not give the contribution of an individual variable to the cost
  - Does not give reduced costs for all variables
- **Objective landscapes** do (in a certain sense)

Comparison with **temporal linear relaxation** :

- **Temporal linear relaxation** heavily exploit the constraints of the problem (their linearization) and is computed at each LNS move
- **Objective landscapes** focus on the cost function and are computed once and for all at the beginning of the search

Comparison with **temporal linear relaxation** :

- **Temporal linear relaxation** need to convexify the objective terms
- **Objective landscapes** make less assumption about convexity of objective terms and is more related to the notion of quasiconvexity (see later)

Comparison with **temporal linear relaxation** :

- **Temporal linear relaxation** needs building and solving an LP
- **Objective landscapes** are much lighter to build

# Introduction: notations

We suppose an optimization problem:

$$\begin{array}{ll}\text{minimize} & z = f(x) \\ \text{subject to} & c(x, y)\end{array}$$

Variables  $x = [x_1, \dots, x_n]$ ,  $y = [y_1, \dots, y_m]$  are **decision variables** in the general sense, that is they can represent:

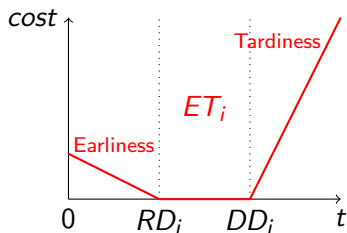
- Actual integer variables (IloIntVar) of the model
- Attributes of an interval variable (presence, start, end, size, length)
- Contribution of interval variables to cumul functions (IloHeightAtXXX)

The objective function  $z$  functionally depends on a set of **objective decision variables**  $x$

# Introduction: example

Flow-shop scheduling problem with earliness/tardiness cost:

$$\begin{aligned} &\text{minimize } z = \sum_{i=1}^n \text{endEval}(o_{i,m}, ET_i) \\ &\text{subject to } \forall i \in [1, n], \forall j \in [1, m-1], \text{endBeforeStart}(o_{i,j}, o_{i,j+1}) \\ &\quad \forall j \in [1, m], \text{noOverlap}(\{o_{i,j}\}_{i \in [1, n]}) \end{aligned}$$



Objective decision variables are the end values of  $o_{i,m}$

# Introduction: assumptions

- Focus on objective variables **without holes** in the domain:
  - Interval variables attributes (start, end, size, length, presence, heights)
  - Boolean variables
  - Integer variables with a “continuous” semantics
  - ~> Domain of a variable is a range  $[L, U]$
- Single objective function (no multi-objective for now)

# Landscape definition: notations

Let  $x_i$  be a decision variable and  $S$  a feasible solution

We denote:

- $D_i$  the initial domain of variable  $x_i$
- $D = D_1 \times \dots \times D_n$
- $X_i^L$  a lower bound on the value of variable  $x_i$
- $X_i^U$  an upper bound on the value of variable  $x_i$
- $X_i^S$  the value of variable  $x_i$  in solution  $S$
- $Z^L$  a lower bound on the objective function
- $Z^U$  an upper bound on the objective function



# Landscape definition: notations

Let  $x_i$  be an objective variable, function  $f_i^* : D_i \rightarrow \mathbb{R}$  denotes the optimal objective value one can obtain when  $x_i = v$ . That is:

$$f_i^*(v) = \min_{x \in D \text{ s.t. } x_i = v} f(x)$$

## Idea

The idea of objective landscapes is to build, for each decision variable  $x_i$ , a function that approximates  $f_i^*$

# Landscape definition: notations

We suppose the existence of a given propagation algorithm that is able to propagate constraints like  $f(x) \leq z$  in order to reduce the domain of variables  $x_i$

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Let  $x_i$  be an objective variable, we define:

- $Z^L$  the smallest value of  $z \in \mathbb{R}$  such that the propagation of objective cut  $f(x) \leq z$  does not fail
- $X_i^L(z)$  for  $z \geq Z^L$  as the lower bound on variable  $x_i$  obtained after the propagation of an objective cut  $f(x) \leq z$
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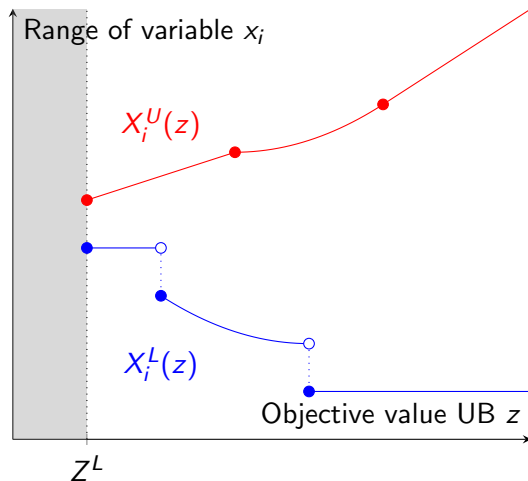
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Clearly, by monotonicity of the propagation,  $X_i^L(z)$  (resp.  $X_i^U(z)$ ) is a non-increasing (resp. non-decreasing) function of  $z$  and  $X_i^L(z) \leq X_i^U(z)$

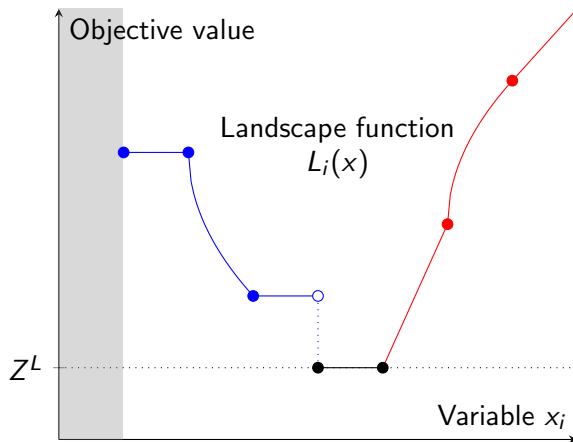
# Landscape definition



## Objective landscape (informal definition)

The **objective landscape function** of an objective variable  $x_i$  is a function  $L_i$  whose graph is the  $90^\circ$  rotate of the union of the graphs of the two functions  $X_i^L(z)$  and  $X_i^U(z)$ .

# Landscape definition



## Objective landscape (formal definition)

The **objective landscape function** of an objective variable  $x_i$  is a function  $L_i : D_i \rightarrow [Z^L, +\infty)$  defined as follow:

- For  $v \in [X_i^L(Z^L), X_i^U(Z^L)] : L_i(v) = Z^L$
- For  $v < X_i^L(Z^L) : L_i(v) = \min \{z \mid X_i^L(z) \leq v\}$
- For  $v > X_i^U(Z^L) : L_i(v) = \min \{z \mid X_i^U(z) \geq v\}$

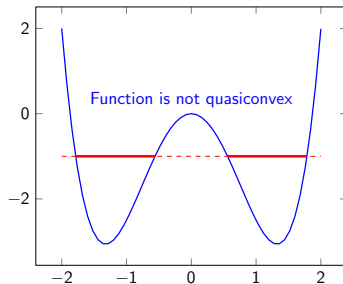
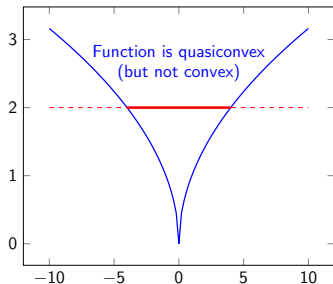


# Landscape properties: quasiconvex functions

## Definition

A function  $f : S \rightarrow \mathbb{R}$  defined on a convex subset  $S$  of a real vector space is **quasiconvex** if for all  $x, y \in S$  and  $\lambda \in [0, 1]$  we have  $f(\lambda x + (1 - \lambda)y) \leq \max \{f(x), f(y)\}$

Informally, along any stretch of the curve the highest point is one of the endpoints. Example of one-variable functions:



## Property 1 (Quasiconvexity of landscape functions)

For any objective variable  $x_i$ ,  $L_i$  is quasiconvex

Proof is a direct consequence of the fact  $X_i^L(z)$  (resp.  $X_i^U(z)$ ) is a non-increasing (resp. non-decreasing) function of  $z$  and  $X_i^L(z) \leq X_i^U(z)$



## Property 2 (Landscape functions as lower bounds)

For any objective variable  $x_i$ ,  $L_i \leq f_i^*$

Proof is a direct consequence of the soundness of propagation □

## Definition (Exact landscape function)

A landscape function  $L_i$  for an objective variable  $x_i$  is said to be **exact** if and only if  $L_i = f_i^*$

# Landscape properties: a particular (but common) case

## Assumption

Let's suppose:

- 1 For all objective variable  $x_i$ , function  $f_i^*$  is quasiconvex
- 2 Propagation of  $f(x) \leq z$  performs bound-consistency on the  $x_i$ 's

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This is for instance the case if:

- $f$  is the sum of individual terms:  $f(x_1, \dots, x_n) = \sum_i f_i(x_i)$  and
- Each function  $f_i$  is quasiconvex

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It also hold more generally for any combination of sum, min, max, product (of non-negative terms) of quasiconvex individual terms  $f_i(x_i)$ , as far as  $x_i$  appears only once in the expression, for instance:

$$x^2\sqrt{y} + \max(u^3, \min(v, w))$$

# Landscape properties: a particular but common case

## Property 3

If function  $f_i^*$  is quasiconvex and if the propagation of  $f(x) \leq z$  performs bound-consistency on  $x_i$ 's then  $L_i$  is exact ( $L_i = f_i^*$ )

Proof is not very complex. By property 2, we know  $L_i \leq f_i^*$ . The proof that  $f_i^* \leq L_i$  exploits bound-consistency (and quasiconvexity of  $f_i^*$  where bounding functions  $X_i^L(z)$  and  $X_i^U(z)$  are discontinuous). The formal proof is available in a separate document □



## Theorem (Characterization of exact landscapes)

A landscape function  $L_i$  is exact if and only if function  $f_i^*$  is quasiconvex and the propagation of  $f(x) \leq z$  performs bound-consistency on  $x_i$ 's

Proof is easy by combining properties 1 and 3 and the fact that if propagation does not perform bound-consistency, one can easily exhibit cases where for a given  $v$ ,  $L_i(v) > f_i^*(v)$  □

# Landscape computation

Given an upper-bound  $Z^U$  on the objective (e.g. from an existing initial solution) and a lower bound  $Z^L$  (e.g. from root propagation), landscape functions can be computed by discretizing the range of possible objective values  $[Z^L, Z^U]$  as a discrete set  $\{z_1 = Z^L, z_2, \dots, z_j, \dots, z_{m-1}, z_m = Z^U\}$  and storing the min/max bounds of each decision variables  $x_i$  for each  $z_j$

## Basic algorithm for landscape computation

```
for  $j$  in  $m..1$  do  
   $\text{post}(f(x) \leq z_j)$   
  for  $i$  in  $1..n$  do  
     $X_i^L(z_j) \leftarrow$  current lower bound of  $x_i$   
     $X_i^U(z_j) \leftarrow$  current upper bound of  $x_i$ 
```

## Questions:

- 1 Beside objective expression, should we also propagate constraints ?
- 2 How to best discretize objective values ?
- 3 How to interpolate landscapes ?

Question: should we also propagate constraints ?

Pros:

- It is a way to (modestly) handle constraints in the landscapes
- The algorithm to compute landscape can work directly on the complete problem as extracted in the engine
- The landscapes can easily be extended so that they are also computed for **non-objective** decision variables, we just need to record the bounds for these variables too

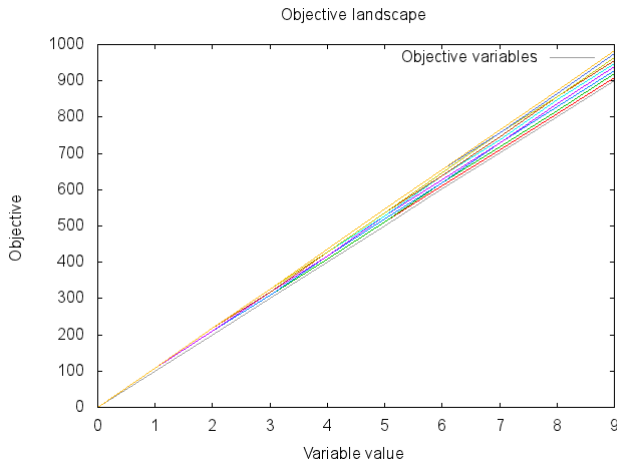
Question: should we also propagate constraints ?

## Cons:

- It introduces dependencies between the objective variables and messes-up the theoretical framework about exactness of landscapes
- More importantly: it hides some important information

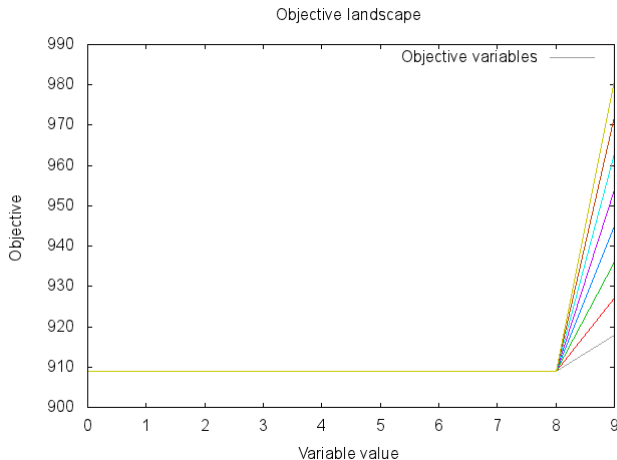
# Landscape computation

Example:  $i \in [0, 10), x_i \in [0, 10)$ : minimize  $\sum_i (100 + i)x_i$  subject to allDiff $\{x_i\}$



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Question: How to best discretize objective values ?

Ideally, we would like to select objective values

$\{z_1 = Z^L, z_2, \dots, z_j, \dots, z_{m-1}, z_m = Z^U\}$  that result in a good sampling of the typical values of  $x_i$  met in the solutions  $(X_i^S)$  as the landscape function will be accessed at these values

Note:

- The basic algorithm for landscape computation can of course be adapted to select  $z_i$  on the fly given the variable bounds computed with the previous  $z_j, j < i$
- We could even think of collecting the typical values of variables in solutions and then have a step that refines the landscape function around those values



Question: How to interpolate landscapes ?

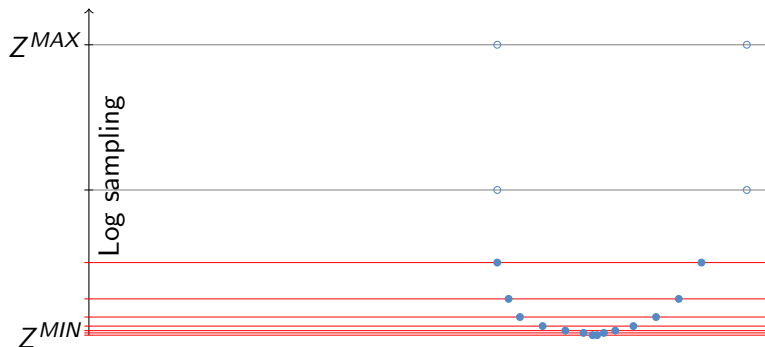
- Interpolation with lower bounding step functions preserves the lower bounding property of landscapes function
- Linear interpolation may be fine too

Prototype implementation:

- Landscapes are computed at the end of presolve, in a transformer that converts only the objective expression (plus some functional constraints like:  $z - x_1 - \dots - x_n == K$ )
- Let  $Z^{MIN}$  and  $Z^{MAX}$  denote the objective range (in particular after propagation-based transformer)

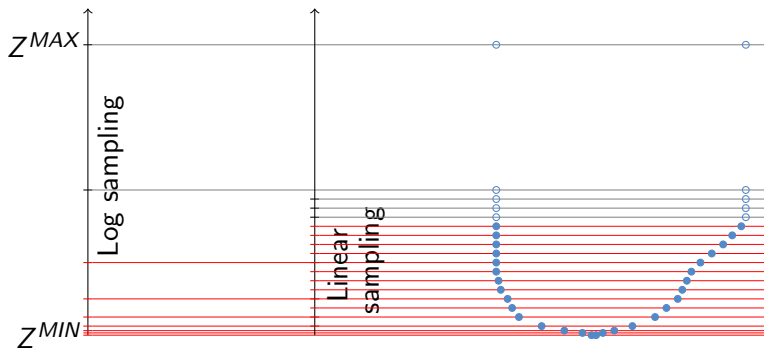
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Prototypical implementation:



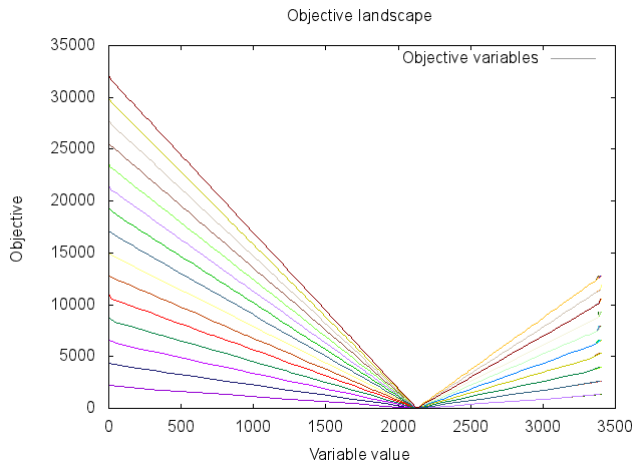
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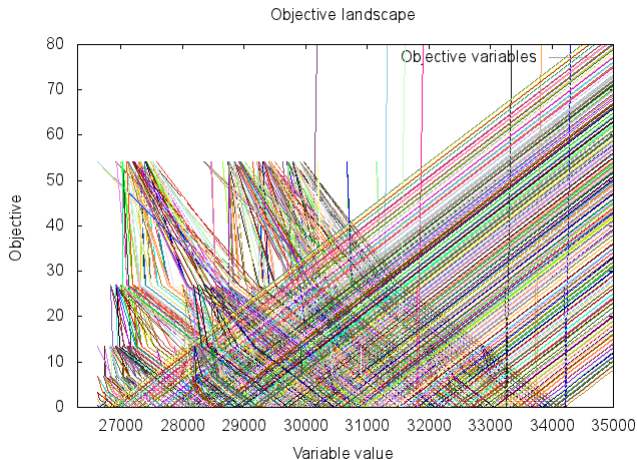
# Landscape computation

Example: an instance of the CommonDueDate family



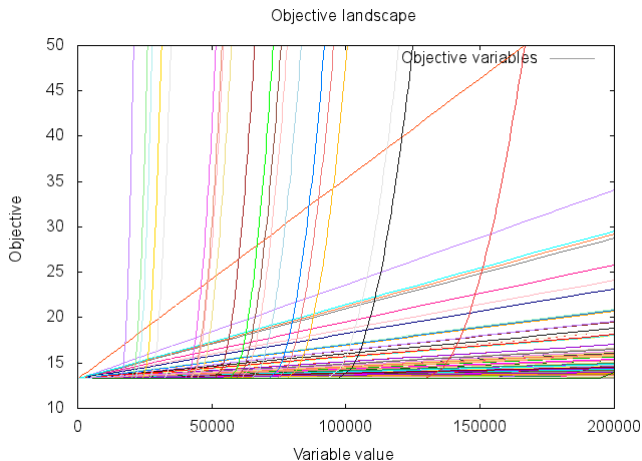
# Landscape computation

Example: all 1816 landscape functions for FLETC instance



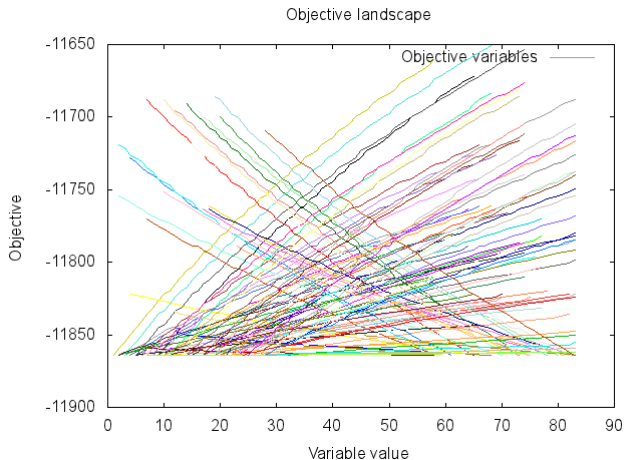
# Landscape computation

Example: all 601 landscape functions for an ST Micro instance



# Landscape computation

Example: RCPSP with discounted cash flow





# Landscape exploitation

- ① Preventing cost degradation for some variables in LNS move
- ② Defining new types of neighborhoods
- ③ Selecting variables and values in completion goals
- ④ Integration in BOLIDE meta-heuristics
- ⑤ ... any other idea welcomed!

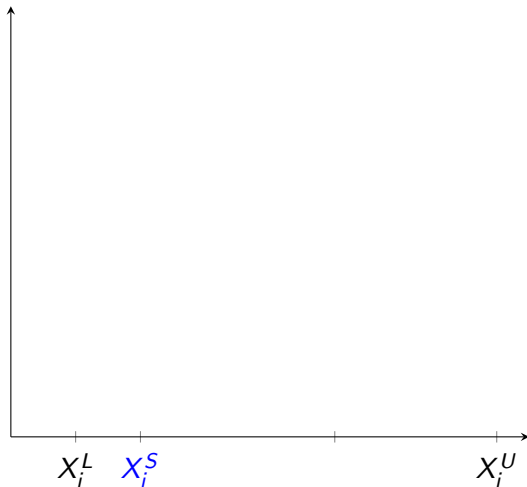
At the root node of an LNS move, after restoration of the LNS fragment, for each objective variable  $x_i$ , we know:

- $X_i^S$ : the value of  $x_i$  in the incumbent solution
- $X_i^U$ : the current upper-bound on  $x_i$
- $X_i^L$ : the current lower-bound on  $x_i$
- $L_i$ : the landscape function of  $x_i$
- $F_i$ : whether or not  $x_i$  belongs to the LNS fragment

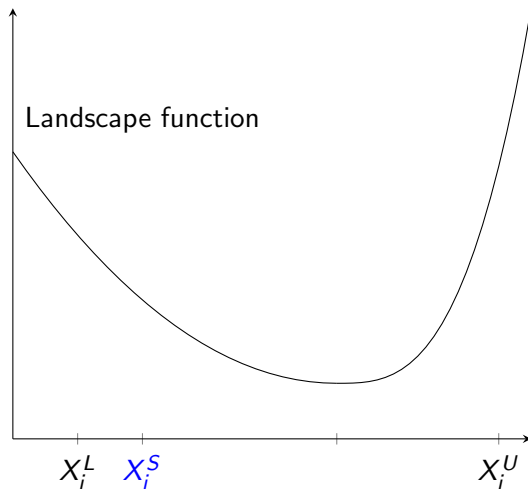
Very simple idea. At a given LNS move:

- Use or not the strategy below depending on a Boolean learner
- If strategy is used, select a certain ratio of objective variables (learned, currently possible values are: 5%, 10%, 20%, 50%, 75%, 100%), according to some **criterion**, on which to enforce non-degradation
- Non-degradation constraint: for selected variables  $x_i$ , restrict their domain to the range of values  $v$  such that  $L_i(v) \leq L_i(X_i^S)$

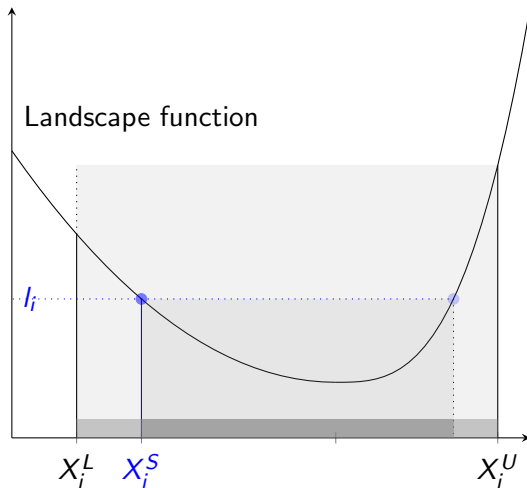
# Landscape exploitation: Preventing cost degradation



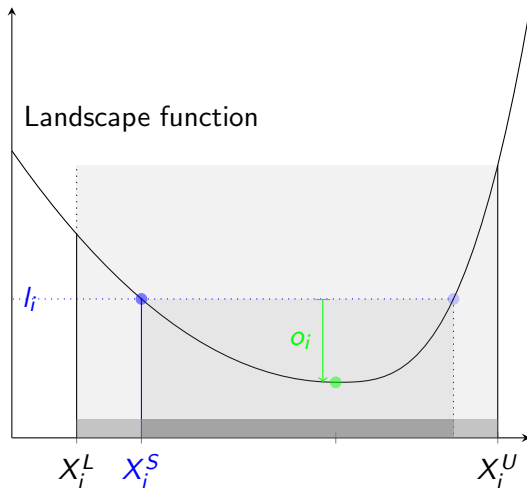
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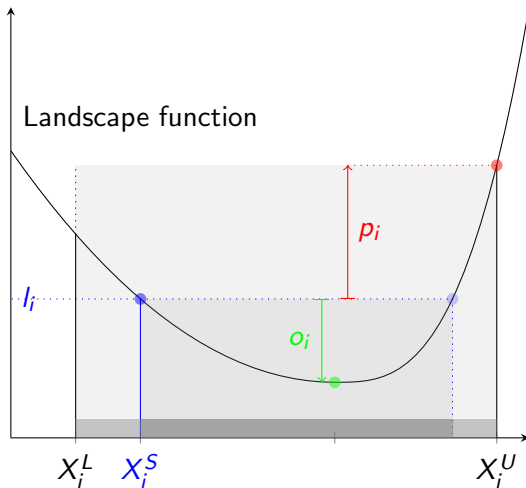
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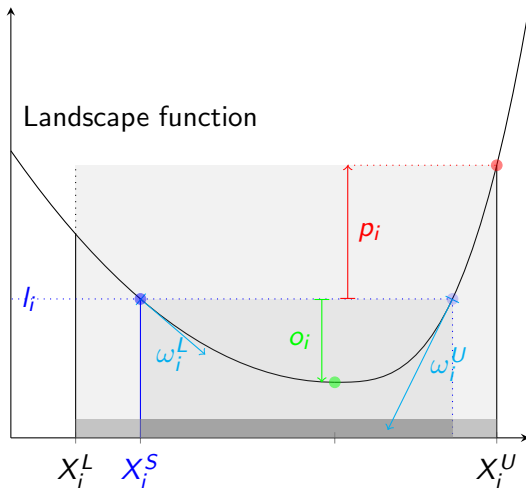


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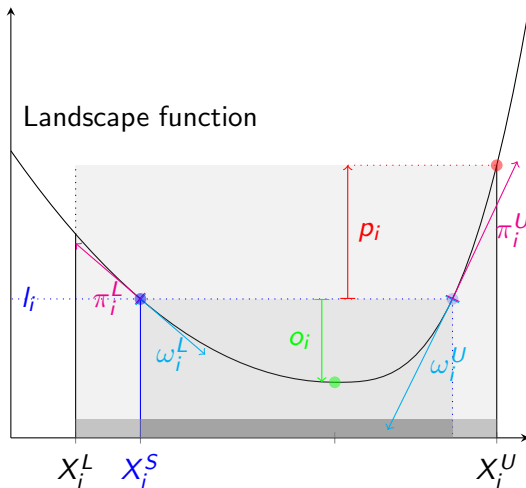




# Landscape exploitation: Preventing cost degradation



# Landscape exploitation: Preventing cost degradation



# Preliminary results

Comparison using the same code (on Head, similar to 12.7) with 1 worker with/without the usage of landscapes. Average over all scheduling tests (3146 instances, 129 families). Average speed-up ratio uses *family average*.

Criterion (argmax)	Speedup
① $l_i$	1.24
② $\max(o_i, p_i)$	1.22
③ $ \pi_i + \omega_i $	1.27
④ random	<b>1.35</b>
⑤ $o_i$	1.24
⑥ $p_i$	1.21

# This is very early work ...

- 1 Preventing cost degradation in LNS moves
- 2 Defining new types of neighborhoods
- 3 Selecting variables and values in completion goals
- 4 Integration in BOLIDE meta-heuristics
- 5 ... any other idea welcomed!

Thank you

Questions?