

# A Concise Horn Theory for RCC8

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**Abstract.** RCC8 is a well-known constraint language for expressing and reasoning about spatial knowledge. We state a simple and concise Horn theory for RCC8 analogous to the ORD-Horn theory for temporal reasoning. This theory allows for expressing RCC8 and retains tractability of the well-known Horn reduct of RCC8. Further, it is much more adequate for practical purposes in the area of logic programming and surpasses previous attempts.

## 1 BACKGROUND

We assume the reader is familiar with standard FOL concepts, notation in the CSP literature, and syntactic interpretations [2, 6]. A constraint language  $\Gamma$  is simply a relational FO structure. The *Constraint Satisfaction Problem for  $\Gamma$* ,  $CSP(\Gamma)$ , is the decision problem if some given primitive positive FO sentence formulated over the signature of  $\Gamma$  is satisfiable in  $\Gamma$ . An algorithm *solves*  $CSP(\Gamma)$  if it solves the decision problem.

The spatial reasoning formalism RCC8 [3] describes binary relations between abstract spatial regions. There are 8 relation primitives: equal  $EQ$ , disconnected  $DC$ , externally connected  $EC$ , partially overlaps  $PO$ , non-tangential proper part  $NTPP$ , tangential proper part  $TPP$ , and the converses  $NTPPI$ ,  $TPPI$ . For example, the formalism allows to formally state that “Central Park is a (non-tangential proper) part of Manhattan”. These eight relations are denoted in the signature  $\sigma_{RCC8A}$  of atomic RCC8 relation symbols. Further, RCC8 contains all disjunctions of these relation symbols forming the signature  $\sigma_{RCC8}$  of 256 relations. We denote by  $\Gamma$  the constraint language RCC8 on signature  $\sigma_{RCC8}$  constructed as an  $\omega$ -categorical structure as in [2].

There are two important reducts of  $\Gamma$ : the  $\sigma_{RCC8A}$ -reduct  $\Gamma_{RCC8A}$ , and the  $\sigma_{\widehat{\mathcal{H}}_8}$ -reduct  $\Gamma_{\widehat{\mathcal{H}}_8}$ . Symbols of  $\sigma_{\widehat{\mathcal{H}}_8}$  denote the RCC8 Horn relations. Renz and Nebel [4] gave a transformation of any instance of  $CSP(\Gamma_{\widehat{\mathcal{H}}_8})$  with  $n$  FO variables into an equivalent propositional Horn formula with  $O(n^4)$  literals. Well-known results (e.g. [4]) are:  $\sigma_{RCC8A} \subsetneq \sigma_{\widehat{\mathcal{H}}_8} \subsetneq \sigma_{RCC8}$ ,  $|\sigma_{\widehat{\mathcal{H}}_8}| = 148$ , and strong 3-consistency solves  $CSP(\Gamma_{\widehat{\mathcal{H}}_8})$ . The reduct  $\Gamma_{\widehat{\mathcal{H}}_8}$  is the largest maximal tractable subclass of RCC8 that includes all atomic relations. Further,  $CSP(\Gamma_{\widehat{\mathcal{H}}_8}) \in P$  while  $CSP(\Gamma)$  is NP-complete.

We present a syntactic interpretation of RCC8 and a concise Horn theory for solving  $CSP(\Gamma_{\widehat{\mathcal{H}}_8})$  in polynomial time. The theory is weaker than strong 3-consistency. The Herbrand expansion yields propositional Horn formulas with  $O(n^2)$  literals and  $O(n^3)$  clauses.

## 2 A CONCISE HORN THEORY

Bennett discussed a number of different representations for logical reasoning about spatial relations [1]. Among them is the RCC7 rep-

resentation which consists of relations representable as equations in interior algebra. These 7 relations are: equality  $EQ$ , disconnectedness  $DC$ , discreteness  $DR$ , parthood  $P$ , non-tangential parthood  $NTP$ , and the converses  $P^{-1}$ ,  $NTP^{-1}$ . Bennett found 115 distinct relations to be definable using purely conjunctive formulas if the complements of the relations are included.

In the following we introduce a set of relations closely related to RCC7 and its use for defining further relations. Here,  $EQ$  is not included as a primitive because it can be defined using parthood. Further, we drop converse symbols and use the complements of  $DC$  and  $DR$ : connectedness  $C$  and overlap  $O$ . We refer to this set as RCC4. Its signature can be thought of as a shorthand notation of specific RCC8 relation symbols, i.e., we treat  $\sigma_{RCC4}$  as a proper subset of  $\sigma_{RCC8}$ .

**Definition 1 (RCC4)** In the binary signature  $\sigma_{RCC4} := \{C, O, NTP, P\}$   $NTP$  denotes the relation symbol  $NTPP$  in  $\sigma_{RCC8}$ ,  $C$  the relation symbol representing  $\neg DC(x, y)$  over  $\Gamma$ ,  $O$  the relation symbol representing  $\neg (DC(x, y) \vee EC(x, y))$ , and  $P$  the symbol representing  $TPP(x, y) \vee NTPP(x, y) \vee EQ(x, y)$ . Further,  $\Gamma_{RCC4}$  denotes the  $\sigma_{RCC4}$ -reduct of  $\Gamma$ .

The structure  $\Gamma_{RCC4}$  admits a FO interpretation of  $\Gamma$ . We first give the syntactic interpretation of the atomic RCC8 relation symbols (virtually the same map as provided in Table 5.4 of [1]).

**Definition 2 ( $\pi_{RCC4}$ )** The syntactic interpretation  $\pi_{RCC4}$  of atomic RCC8 relation symbols in the RCC4 symbols is defined as follows:

$$\begin{aligned}\pi_{RCC4}(DC)(x, y) &:= \neg C(x, y) \\ \pi_{RCC4}(EC)(x, y) &:= C(x, y) \wedge \neg O(x, y) \\ \pi_{RCC4}(PO)(x, y) &:= O(x, y) \wedge \neg P(x, y) \wedge \neg P(y, x) \\ \pi_{RCC4}(EQ)(x, y) &:= P(x, y) \wedge P(y, x) \\ \pi_{RCC4}(TPP)(x, y) &:= P(x, y) \wedge \neg P(y, x) \wedge \neg NTP(x, y) \\ \pi_{RCC4}(NTPP)(x, y) &:= NTP(x, y) \\ \pi_{RCC4}(TPPI)(x, y) &:= \pi_{RCC4}(TPP)(y, x) \\ \pi_{RCC4}(NTPPI)(x, y) &:= \pi_{RCC4}(NTPP)(y, x)\end{aligned}$$

It is easy to verify that this syntactic interpretation provides an interpretation of  $\Gamma_{RCC8A}$  in  $\Gamma_{RCC4}$ . The map can be easily extended to all RCC8 relations, thus  $\Gamma$  has an interpretation in  $\Gamma_{RCC4}$ . Note, although RCC4 is very close to RCC7, the formula  $\neg EQ(x, y)$  is written as a disjunction over  $\sigma_{RCC4}$  – contrary to RCC7 where  $EQ$  is a primitive. From the RCC subsumption lattice [3] it follows:

**Proposition 1** Let  $\psi$  be a formula  $L_O(x, y) \wedge L_C(x, y) \wedge L_P(x, y) \wedge L_{NTP}(x, y) \wedge L_P(y, x) \wedge L_{NTP}(y, x)$  in which  $L_R(x, y) = R(x, y)$  or  $L_R(x, y) = \neg R(x, y)$ . Then  $\psi$  is either unsatisfiable in  $\Gamma_{RCC4}$  or equivalent to an atomic RCC8 relation.

For convenience we consider the expansion of  $\Gamma_{RCC4}$  to  $\Gamma_{RCC4+}$  by adding a symbol  $\overline{R}$  for the complement of each relation symbol  $R$

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in  $\sigma_{\text{RCC}4}$  ( $\Gamma_{\text{RCC}4+}$  is still a reduct of  $\Gamma$ ). Then,  $\pi_{\text{RCC}4}$  can be seen as primitive positive. In the following, we write  $\pi_{\text{RCC}4+}$  to denote the syntactic interpretation of each RCC8 relation symbol over  $\Gamma_{\text{RCC}4+}$  that uses CNF defining formulas consisting of prime implicants.

Computer analysis shows the following properties.

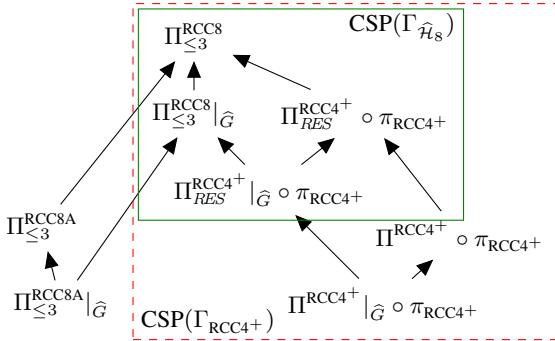
**Proposition 2** *There are 42 RCC8 relations which have a primitive positive definition in  $\Gamma_{\text{RCC}4+}$ . All 148 RCC8 Horn relations have a Horn definition in  $\Gamma_{\text{RCC}4}$ .*

In the following we give a Horn theory over  $\Gamma_{\text{RCC}4}$  written as the Datalog program  $\Pi^{\text{RCC}4+}$  with symbols in  $\sigma_{\text{RCC}4+}$  and a special symbol `false` denoting  $\perp$ . The program contains the necessary subsumption rules and rules on 3 variables that derive from the axiomatization of the atomic composition table of RCC8 given by Bodirsky and Wölfl [2] using  $\pi_{\text{RCC}4+}$  and simplifications. It solves  $\text{CSP}(\Gamma_{\text{RCC}4+})$  by Proposition 1 and the results of [2].

$P(x, x)$	$\overline{NTP}(x, x)$
$C(x, y) \leftarrow O(x, y)$	$O(x, y) \leftarrow P(x, y)$
$P(x, y) \leftarrow NTP(x, y)$	$\text{false} \leftarrow P(y, x), NTP(x, y)$
$O(x, y) \leftarrow O(y, x)$	$C(x, y) \leftarrow C(y, x)$
$\text{false} \leftarrow C(x, y), \overline{C}(x, y)$	$\text{false} \leftarrow NTP(x, y), \overline{NTP}(x, y)$
$\text{false} \leftarrow O(x, y), \overline{O}(x, y)$	$\text{false} \leftarrow P(x, y), \overline{P}(x, y)$
$P(x, y) \leftarrow P(x, z), P(z, y)$	$C(x, y) \leftarrow C(x, z), P(z, y)$
$O(x, y) \leftarrow C(x, z), NTP(z, y)$	$O(x, y) \leftarrow O(x, z), P(z, y)$
$NTP(x, y) \leftarrow P(x, z), NTP(z, y)$	$NTP(x, y) \leftarrow NTP(x, z), P(z, y)$

### 3 PROPERTIES AND EVALUATION

Analogous to the study of temporal formalisms in [6], we can analyze the relation of this new program to commonly used ones: strong 3-consistency ( $\Pi_{\leq 3}^{\text{RCC}8}$ ), the atomic composition table ( $\Pi_{\leq 3}^{\text{RCC}8A}$ ), and the augmented  $\Pi^{\text{RCC}4+}$  that emulates positive unit resolution on Horn clauses in  $\pi_{\text{RCC}4+}$  ( $\Pi_{\text{RES}}^{\text{RCC}4+}$ ). Additionally, each of them can be weakened by restricted inference to arcs of an instance that are part of some fixed chordal graph of an instance (denote by  $|_{\widehat{G}}$ ). Finally, we can define a strict partial order on the programs, where  $A \prec B$  if program  $A$  rules out more partial solutions than program  $B$ . For details, references, and discussion of these concepts see [6].



**Figure 1.** Lattice of propagation strength of Datalog programs given interpretations for RCC8. Frames indicate CSPs solved by the programs.

The lattice for  $\prec$  is given in Figure 1 (where  $A \leftarrow B$  if  $A \prec B$ ) which also states which CSPs of reducts are solved by the programs.

In order to encode instances to propositional CNF it has been common practice to consider the *support encoding* of  $\Pi_{\leq 3}^{\text{RCC}8A}|_{\widehat{G}}$  (see [6] for references). However, this encoding is quite weak.

**Proposition 3** *Unit propagation on the support encoding of  $\Pi_{\leq 3}^{\text{RCC}8A}$  does not solve  $\text{CSP}(\Gamma_{\text{RCC}8A})$ , while unit propagation on the Herbrand expansion of  $\Pi_{\text{RES}}^{\text{RCC}4+}|_{\widehat{G}} \circ \pi_{\text{RCC}4+}$  solves  $\text{CSP}(\Gamma_{\widehat{H}_8})$ .*

If **(a)** the support encoding of  $\Pi_{\leq 3}^{\text{RCC}8A}|_{\widehat{G}}$  is used SAT solvers have to rely on search to refute instances, but when **(b)** the Herbrand expansion of  $\Pi_{\text{RES}}^{\text{RCC}4+}|_{\widehat{G}} \circ \pi_{\text{RCC}4+}$  is used search is only necessary on non-Horn relations of instances in  $\text{CSP}(\Gamma)$ . Further, both encodings result in propositional CNF with  $O(n^2)$  literals and  $O(n^3)$  clauses for instances with  $n$  FO variables, but interpretation and program of the new concise theory have less relation symbols and rules.

For a brief evaluation we consider 1000 random instances of  $\text{CSP}(\Gamma)$  constructed from the  $H(150, 15.5, 4.0)$  model with hard relations  $\mathcal{NP}_8$  [5]. These instances have 150 FO variables and a constraint network degree of 15.5. The test setup is the same as in [6].<sup>2</sup>

SAT Encoding	Runtime (seconds) of the Glucose-2.2 SAT solver			
	Avg.	25-PCTL	50-PCTL	99-PCTL
<b>(a)</b>	440.90	227.05	349.48	545.62
<b>(b)</b>	45.87	8.92	18.69	103.75
				441.32

This is a factor 4 to 25 improvement. Further, the average size of an instance in DIMACS drops from **(a)** 360 MiB to **(b)** 160 MiB.

### 4 CONCLUSION

We presented a concise Horn theory for RCC8 that allows for expressing the entire RCC8 language and retains the tractability of the well-known Horn class of relations. The theory is of broad interest for incorporating RCC8 in logic programming, e.g. SAT and ASP. In particular, the theory allows for an improved  $O(n^3)$  propositional representation that is superior to previous work.

For future work it would be interesting to consider propagators based on this theory in Constraint Programming. In particular, a potential graph based algorithm could form a fast propagator.

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<sup>2</sup> The tool to compute the encoding is available; see this publication’s entry at <http://www.informatik.uni-freiburg.de/~ki/publications/>