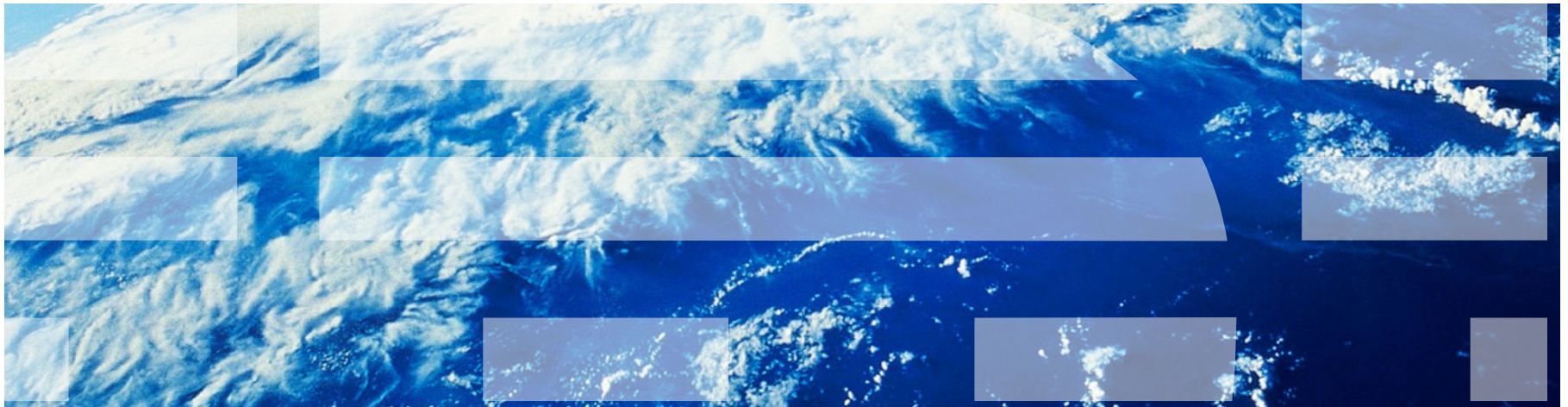


# Modeling and Solving Scheduling Problems with CP Optimizer

Philippe Laborie  
CPLEX Optimization Studio Team



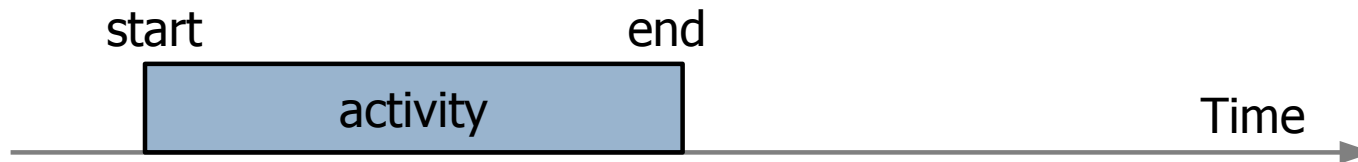
May 28, 2014

# Agenda

- IBM ILOG **CP Optimizer** for **Scheduling**
- Scheduling **concepts** in CP Optimizer
- **Tips** and **Tools**
- **Q&A**

## IBM ILOG CP Optimizer for Scheduling

- Scheduling consist of assigning **starting** and **completion times** to a set of activities while satisfying different types of constraints (resource availability, precedence relationships, ... ) and optimizing some criteria (minimizing tardiness, ...)



- Time is considered as a continuous dimension: domain of possible start/completion times for an activity is potentially **very large**
- Beside start and completion times of activities, other types of decision variables are often involved in real industrial scheduling problems (resource **allocation**, **optional** activities ...)

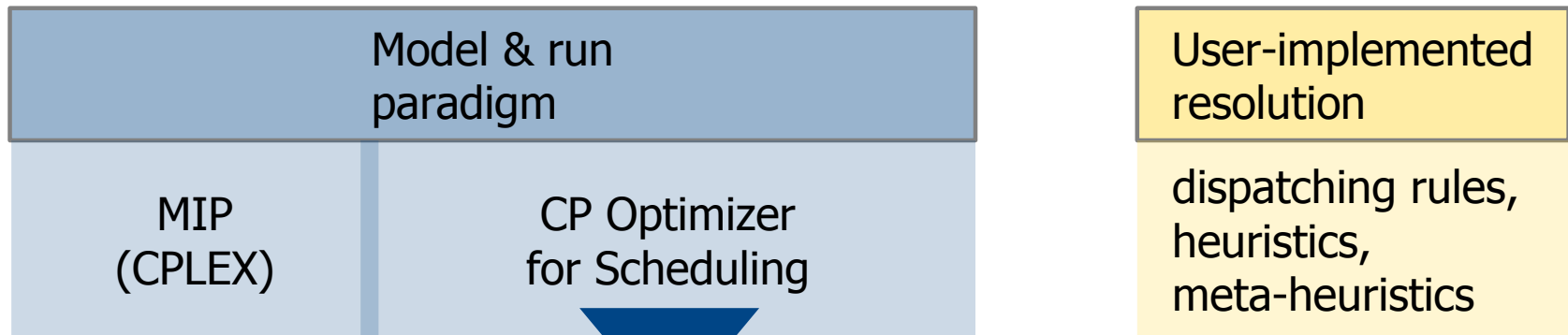
# IBM ILOG CP Optimizer for Scheduling

- Modeling and solving scheduling problems

Model & run paradigm	User-implemented resolution
MIP (CPLEX)	dispatching rules, heuristics, meta-heuristics

# IBM ILOG CP Optimizer for Scheduling

## ■ Modeling and solving scheduling problems



- Dedicated modeling concepts for scheduling
  - Consistent with Optimization paradigm:  
new decision variables, expressions and constraints
  - Available in OPL and APIs (C++, Java, .NET)
- Automatic search
  - Complete
  - Combines tree search, large neighborhood search, meta-heuristics, relaxations, learning, ...

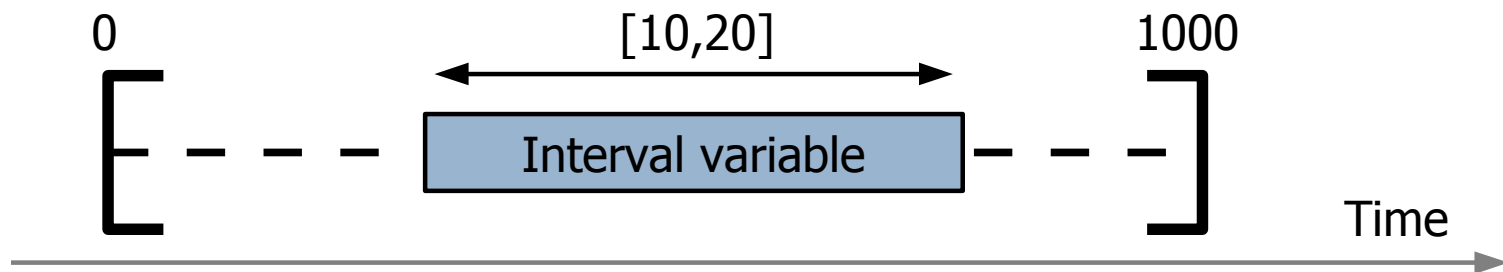
# Agenda

- IBM ILOG CP Optimizer for Scheduling
- Scheduling **concepts** in CP Optimizer
- **Tips** and **Tools**
- **Q&A**

## Concept: **interval variable**

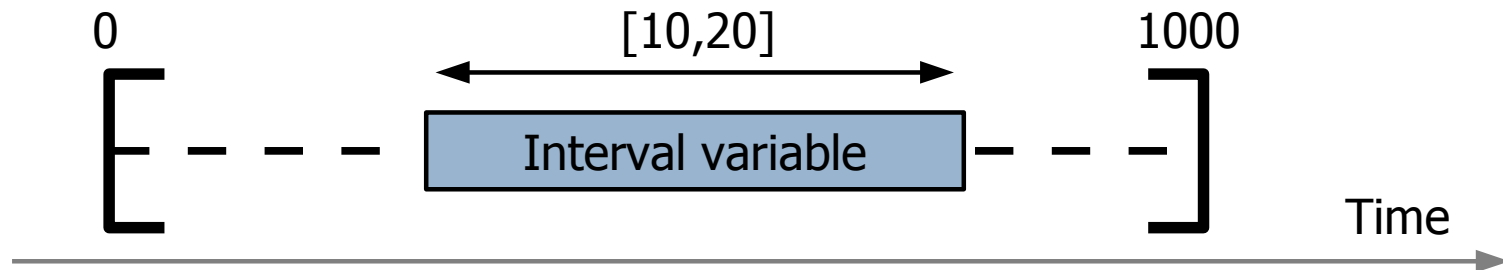
- What for?
  - modeling an interval of time during which a particular property holds (an activity executes, a resource is idle, a tank must be non-empty, ...)
- Example:

**dvar interval x in 0..1000 size in 10..20**



## Concept: interval variable

`dvar interval x in 0..1000 size in 10..20`



- Properties:
  - The **value** of an interval variable is an integer interval [start,end)
  - Domain** of possible values: [0,10), [1,11), [2,12),... [990,1000), [0,11),[1,12),...
  - Domain of interval variables is represented **compactly** in CP Optimizer (a few bounds: `smin`, `smax`, `emin`, `emax`, `szmin`, `szmax`)



## Concepts: **optional interval variable**

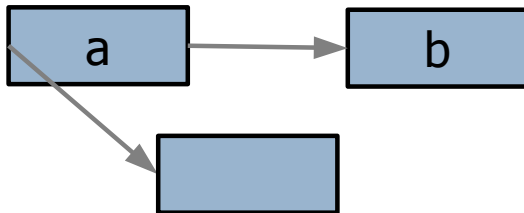
- Interval variables can be defined as being **optional** that is, it is part of the decisions of the problem to decide whether the interval will be **present** or **absent** in the solution
- What for?
  - Modeling optional activities, alternative execution modes for activities, and ... most of the discrete decisions in a schedule
- Example:

```
dvar interval x optional in 0..1000  
              size in 10..20
```

- Properties:
  - An optional interval variable has an additional possible value in its domain (absence value)
  - **Optionality** is a powerful property that you must learn to leverage in your models (more on this later ...)

## Concept: **precedence constraint**

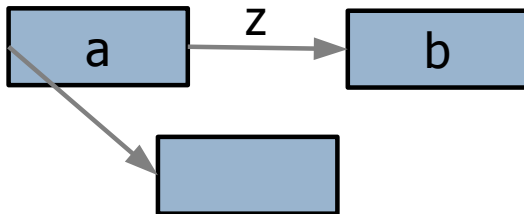
- What for ?
  - Modeling temporal constraints between interval variables
- Example:



```
endBeforeStart(a,b)  
startBeforeStart(a,b)  
startBeforeEnd(a,b)  
endBeforeEnd(a,b)  
endAtStart(a,b)  
startAtStart(a,b)  
startAtEnd(a,b)  
endAtEnd(a,b)
```

## Concept: **precedence constraint**

- What for ?
  - Modeling temporal constraints between interval variables
  - Modeling constant or variable minimal delays
- Example:



```
endBeforeStart(a,b,z)  
startBeforeStart(a,b,z)  
startBeforeEnd(a,b,z)  
endBeforeEnd(a,b,z)  
endAtStart(a,b,z)  
startAtStart(a,b,z)  
startAtEnd(a,b,z)  
endAtEnd(a,b,z)
```

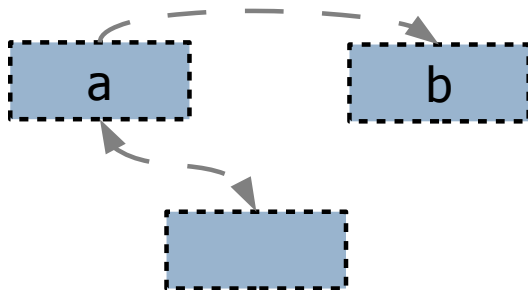
## Concept: **precedence constraint**

- Properties

- Semantic of the constraints handles optionality (as for all constraints in CP Optimizer). Example of endBeforeStart:  
 $\text{present}(a) \text{ AND } \text{present}(b) \Rightarrow \text{end}(a)+z \leq \text{start}(b)$
- All precedence constraints are aggregated in a temporal network and handled by dedicated graph algorithms (fast global propagation, negative cycle detection, ...)

## Concept: **presence constraint**

- What for:
  - Expressing dependency constraints between execution of optional activities or between allocated resources ...
- Examples:



`presenceOf (a) => presenceOf (b)`  
`presenceOf (a) == presenceOf (b)`  
`presenceOf (a) => !presenceOf (b)`  
`!presenceOf (a) => presenceOf (b)`

## Concept: **presence constraint**

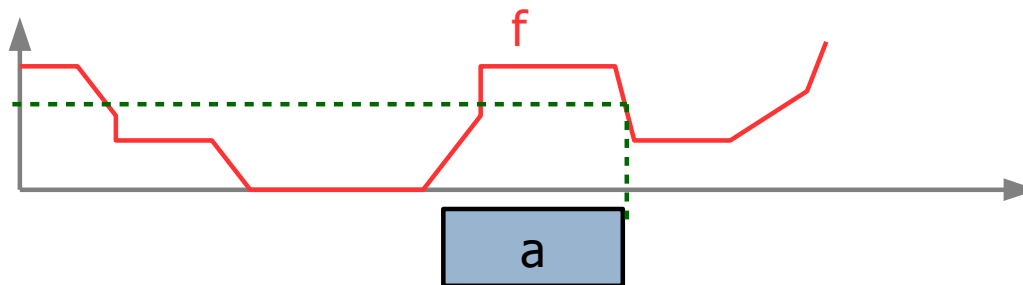
- Properties:
  - All binary constraints between presence status are aggregated in an implication graph
  - CP Optimizer maintains hypothetical bounds on interval variables (e.g. start min would the interval be present)
  - CP Optimizer exploits the implication graph to perform conditional reasoning between related interval variables

## Concept: **expressions on interval variables**

- What for?
  - Evaluating a characteristic of an interval variable (start, end, size) for use in the objective function or in a constraint
- Example:
  - **endOf(a, ABSVAL)** takes value ABSVAL if interval  $a$  is absent otherwise it takes value  $e$  if  $a$  is present and  $a=[s,e]$
  - Typical makespan expression: **max(i in 1..n) endOf(a[i])**

## Concept: **expressions on interval variables**

- What for?
  - Evaluating a piecewise linear function at a particular end-point of an interval variable (earliness-tardiness cost, temporal preference)
- Example:
  - `endEval(f,a, ABSVAL)` takes value ABSVAL if  $a$  is absent otherwise it takes value  $f(e)$  if  $a$  is present and  $a=[s,e]$



- Property:
  - $f$  can be any piecewise linear function (non-continuous, non-convex, ...)

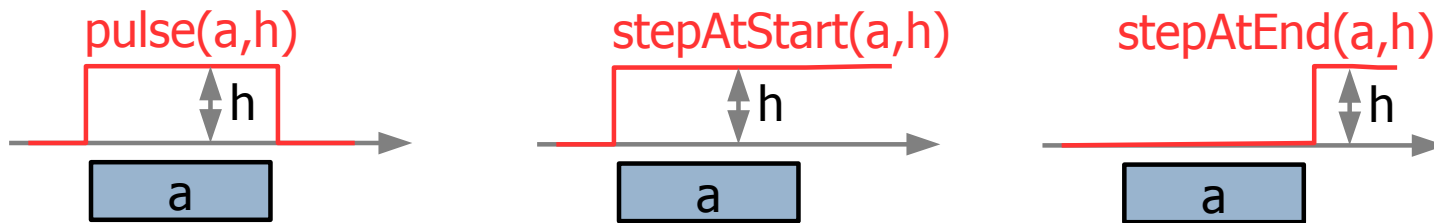


## Concept: **cumul function**

- What for?
  - Modeling and constraining cumulative quantities over time (cumulative use of a discrete resource, level of an inventory with producing and consuming tasks)
  - Restricting the number of intervals that overlap a given date  $t$
  - Forcing a minimal number of intervals to overlap a given date  $t$
  - Complex synchronization constraints between interval variables: e.g. between activities producing/consuming in a tank and the set of time-intervals during which the tank is non-empty

## Concept: **cumul function**

- Cumul functions are built from atomic functions

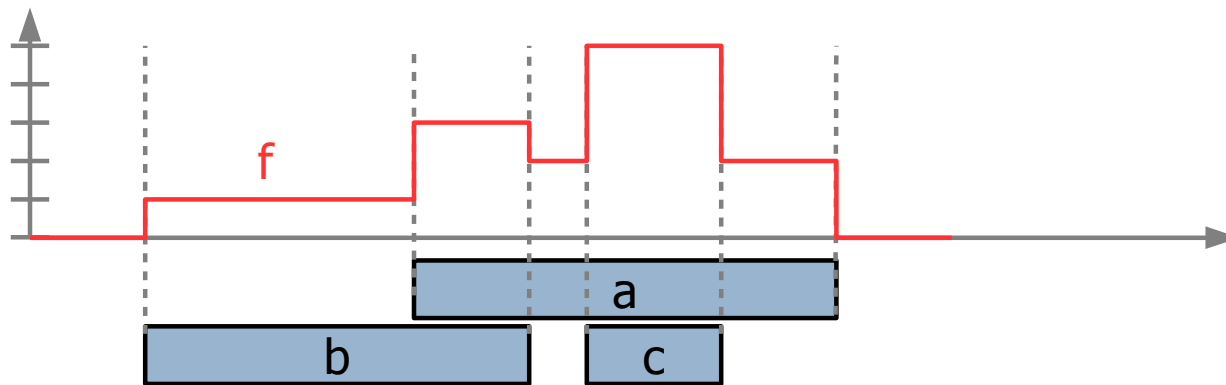


- A cumul function is the sum of atomic functions or their opposite

## Concept: **cumul function**

- Examples:

- cumulFunction  $f = \text{pulse}(a,2) + \text{pulse}(b,1) + \text{pulse}(c,3)$



- Constraint  $f \leq C$  (global capacity limit)
  - Constraint  $\text{alwaysIn}(f, t_0, t_1, C_{\min}, C_{\max})$  (capacity profile)
  - Constraint  $\text{alwaysIn}(f, x, C_{\min}, C_{\max})$  (condition holds over an interval variable  $x$ )

## Concept: **cumul function**

- Properties:
  - Complexity of cumul functions is **independent of the time scale**
  - CP Optimizer is able to reason globally over cumul functions

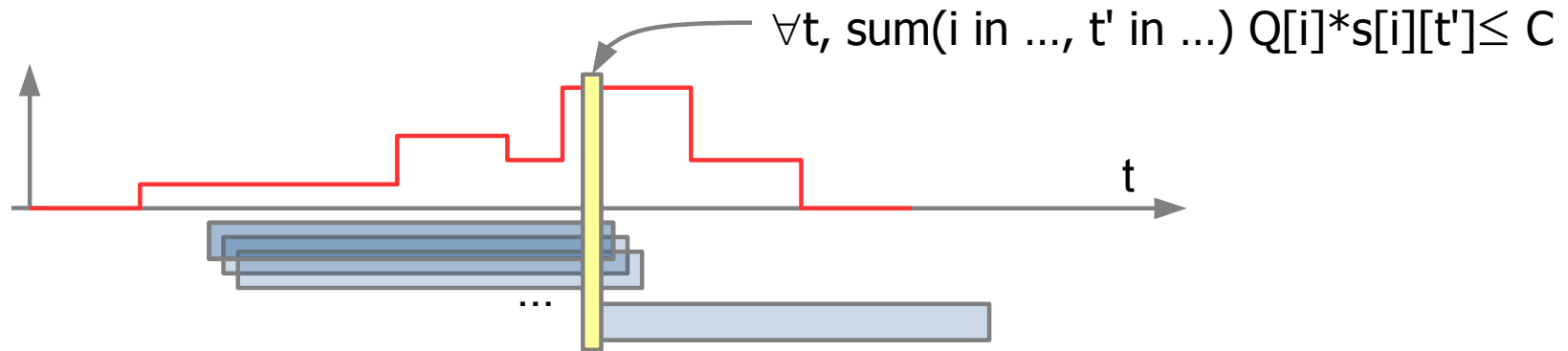
## Concept: **cumul function**

### ■ Properties:

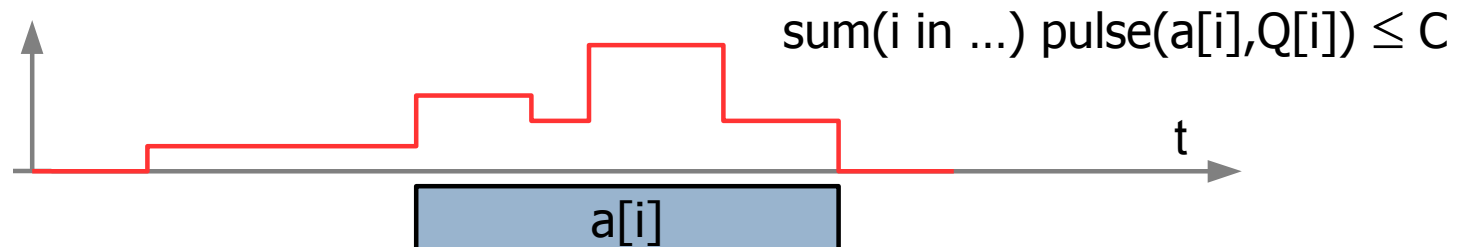
#### –Compare:

##### • Time-indexed MIP model:

$s[i][t] \in \{0,1\}$ :  $s[i][t]=1$  iff  $a[i]$  starts at date  $t$

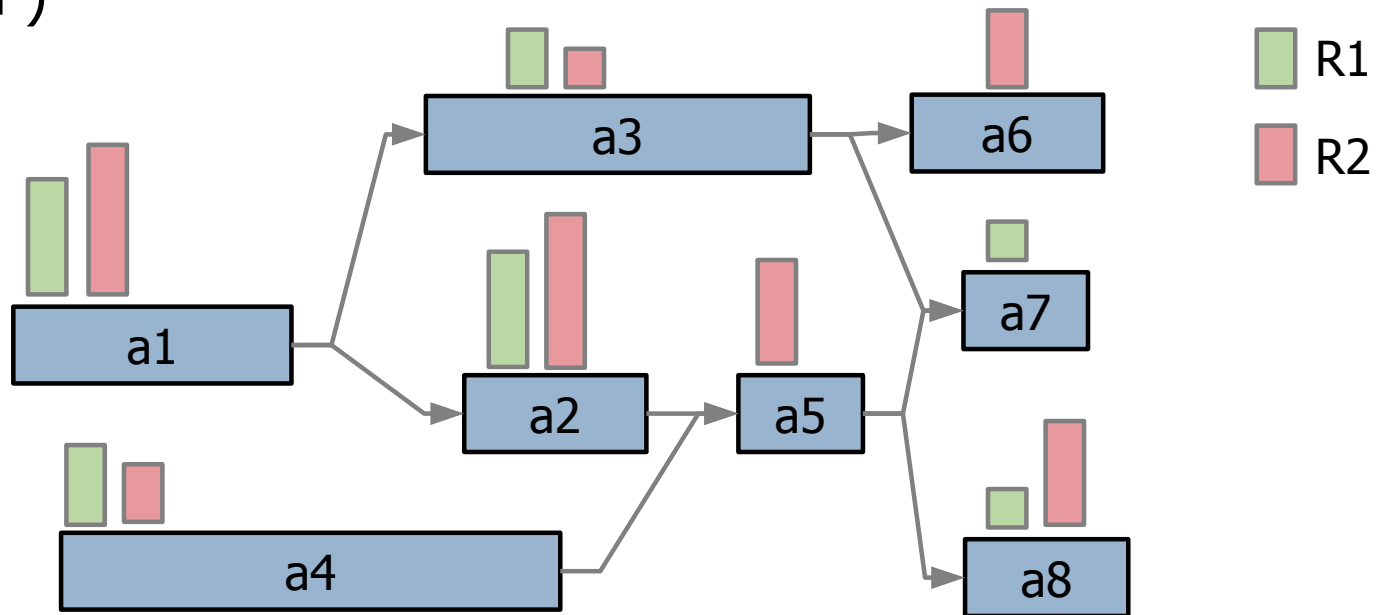


##### • CP Optimizer model:



## Concept: **cumul function**

- Example: Resource-Constrained Project Scheduling Problem (RCPSP)



- Minimization of project makespan

## Concept: **cumul function**

- CP Optimizer model for RCPSP:

```
dvar interval a[i in Tasks] size i.pt;

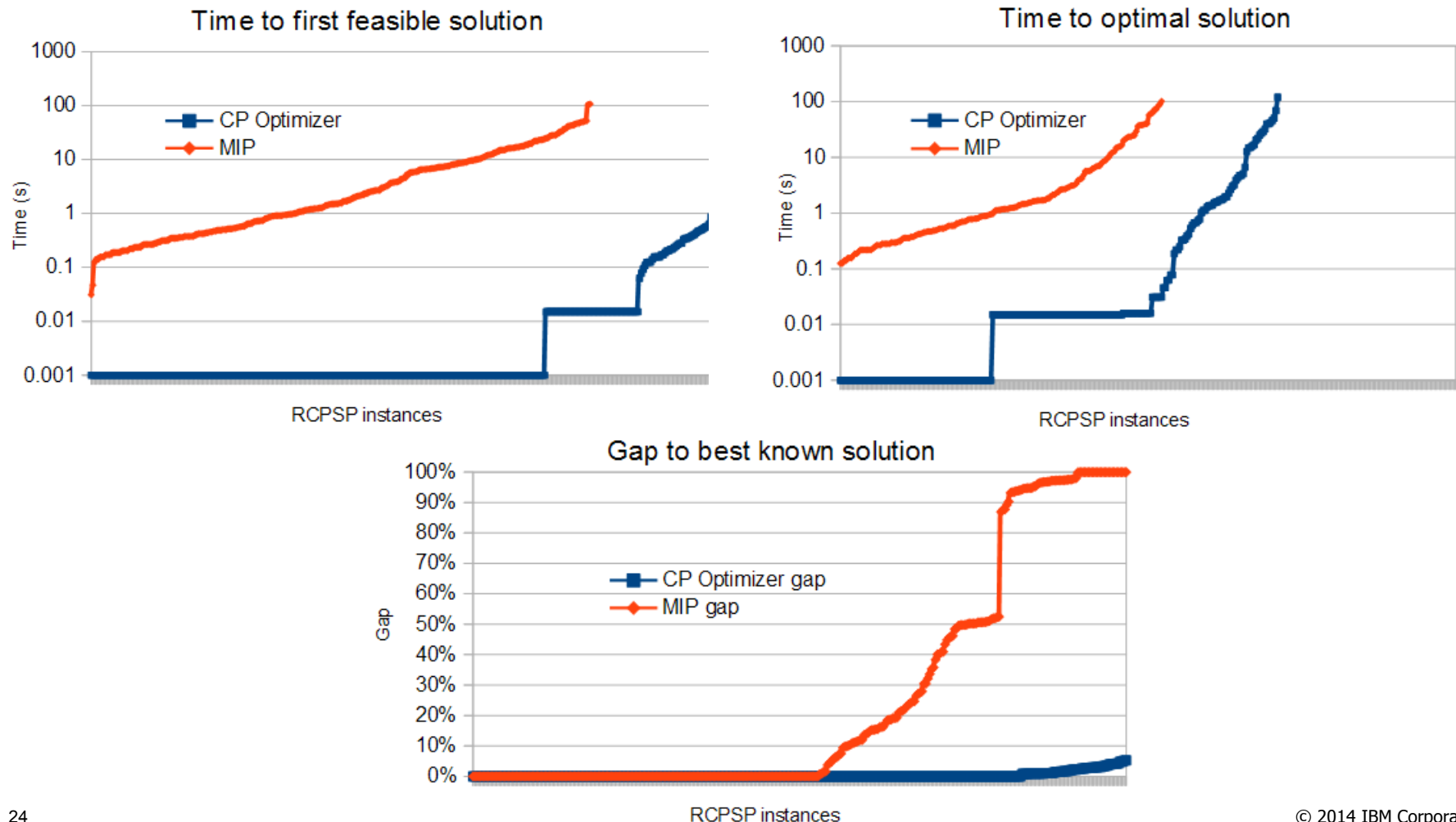
cumulFunction usage[r in Resources] =
    sum (i in Tasks: i.qty[r]>0) pulse(a[i], i.qty[r]);

minimize max(i in Tasks) endOf(a[i]);
subject to {
    forall (r in Resources)
        usage[r] <= Capacity[r];
    forall (i in Tasks, j in i.succs)
        endBeforeStart(a[i], a[<j>]);
}
```

- Comparison of this CP Optimizer model vs a time-indexed MIP model on a set of 300 classical small RCPSP instances (30-120 tasks) + 40 larger ones (900 tasks), time-limit: 2mn, 4 threads

## Concept: **cumul function**

### ■ Comparison of CP Optimizer and MIP performance on RCPSP

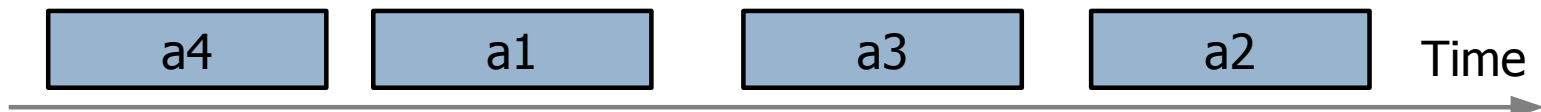




## Concept: **sequence variable**

- What for?
  - modeling sequences of events
  - constraining the transitions between consecutive events (setup times/costs,...)
- Example:  

```
dvar sequence s in all(i in Tasks) a[i]  
noOverlap(s)
```
- Properties:
  - A value for the sequence variable is a total order on the present interval variables. E.g.  $a_4 \rightarrow a_1 \rightarrow a_3 \rightarrow a_2$
  - A constraint `noOverlap` states that intervals of the sequence do not overlap and follows the total order of the sequence

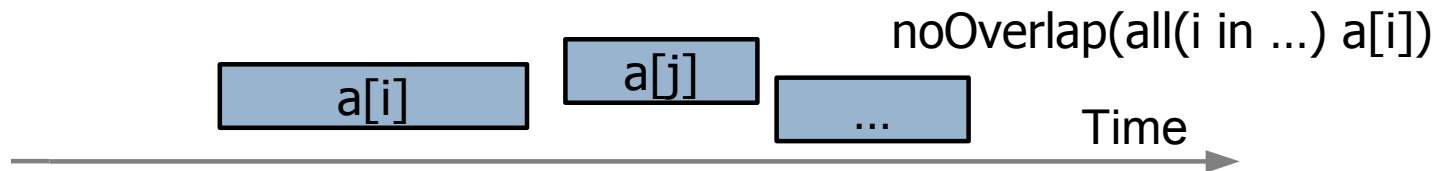


## Concept: **sequence variable**

- Properties:
  - Complexity is **independent of the time scale**
  - CP Optimizer is able to reason **globally** over a sequence variable
  - **Avoid quadratic models** over each pair (i,j) of intervals in the sequence
- Compare:
  - Quadratic disjunctive MIP formulation with big-Ms:
 
$$b[i][j] \in \{0,1\}: b[i][j]=1 \text{ iff } a[i] \text{ before } a[j]$$

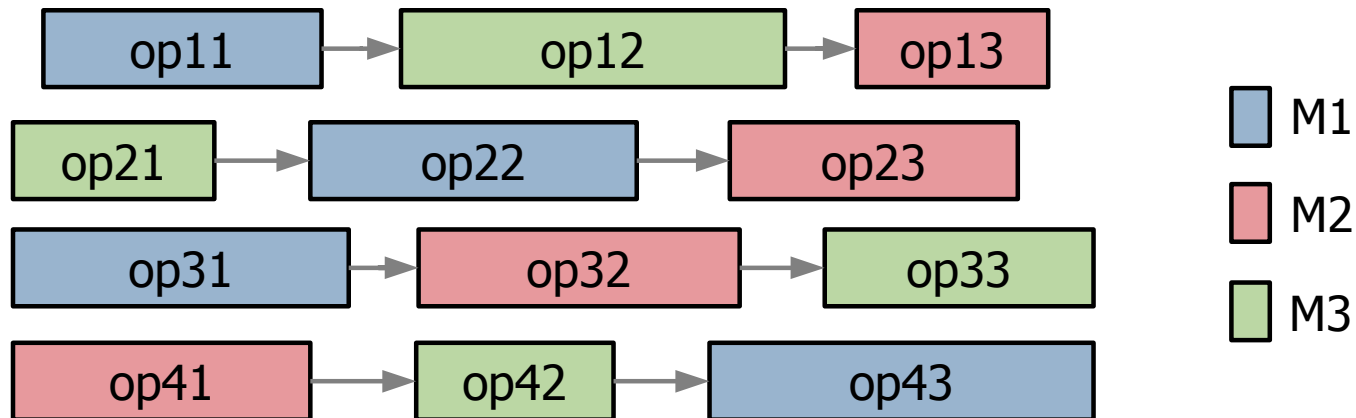
$$\text{end}[i] \leq \text{start}[j] + M \cdot (1 - b[i][j])$$

$$\text{end}[j] \leq \text{start}[i] + M \cdot b[i][j]$$
  - CP Optimizer model:



## Concept: **sequence variable**

- Example: Job-shop Scheduling Problem



- Minimization of makespan

## Concept: **sequence variable**

- CP Optimizer model for Job-shop:

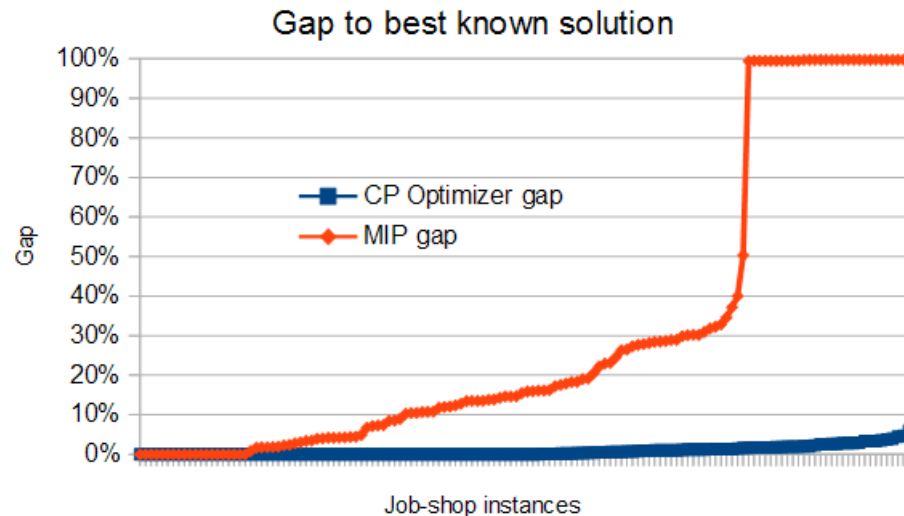
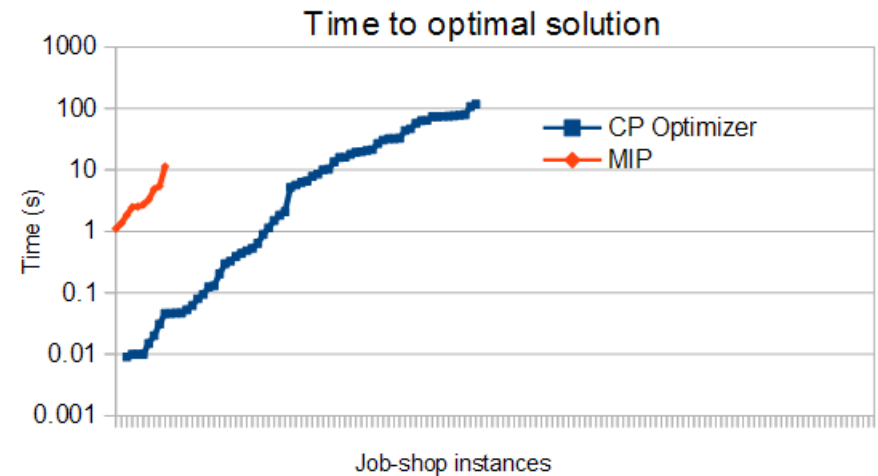
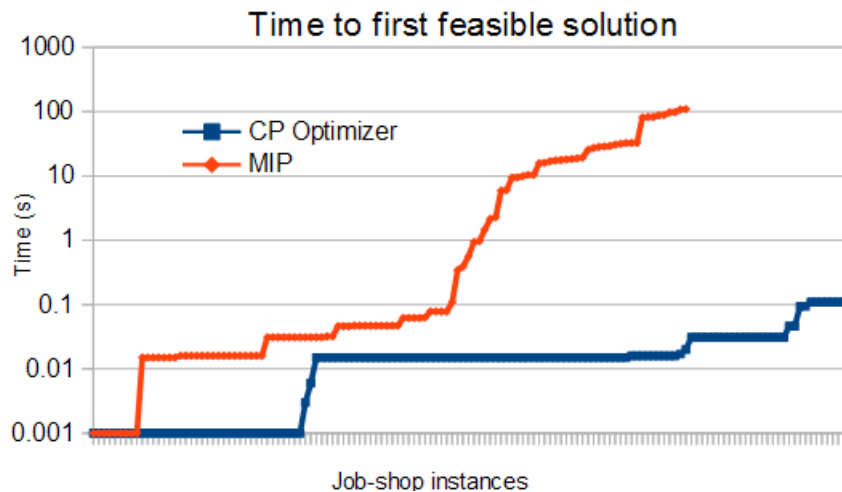
```
dvar interval op[j in Jobs][p in Pos] size Ops[j][p].pt;
dvar sequence mchs[m in Mchs] in
  all(j in Jobs, p in Pos : Ops[j][p].mch == m) op[j][p];

minimize max(j in Jobs) endOf(op[j][nbPos-1]);
subject to {
  forall (m in Mchs)
    noOverlap(mchs[m]);
  forall (j in Jobs, p in 1..nbPos-1)
    endBeforeStart(op[j][p-1], op[j][p]);
}
```

- Comparison of this CP Optimizer model vs a disjunctive MIP formulation on a set of 140 classical Job-shop instances (50-2000 tasks), time-limit: 2mn, 4 threads

## Concept: **sequence variable**

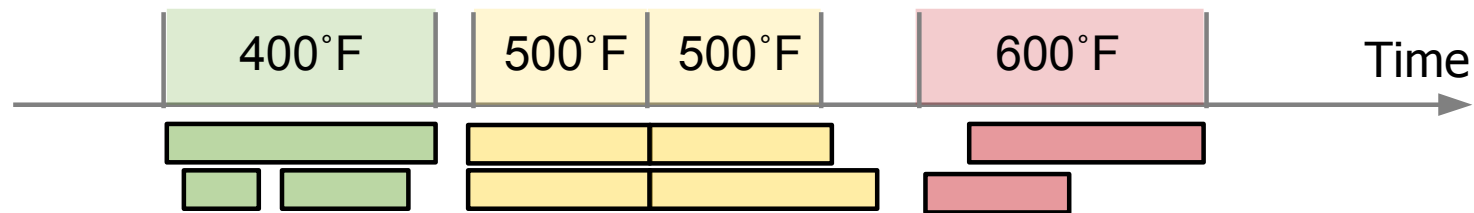
### ■ Comparison of CP Optimizer and MIP performance on Job-Shop



## Concept: state function

- What for?
  - Modeling evolution of a state variable over time
  - Modeling incompatibility relations between activities overlapping a given date  $t$
  - Synchronizing start/end value of compatible tasks (batching)

- Example

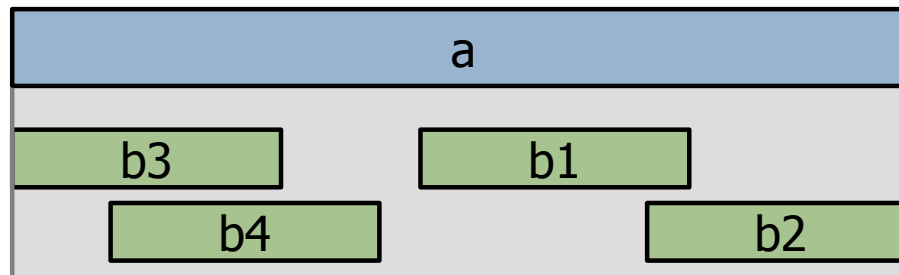


- Properties
  - Complexity is **independent of the time scale**
  - CP Optimizer is able to reason **globally** over a state function
  - **Avoid quadratic models** over each pair  $(i,j)$  of incompatible intervals on the state function

## Concept: **span constraint**

- What for?
  - Modeling task → sub-task decomposition
  - Modeling immobilization time of a resource
- Example

**$\text{span}(a, [b1, \dots, bn])$**

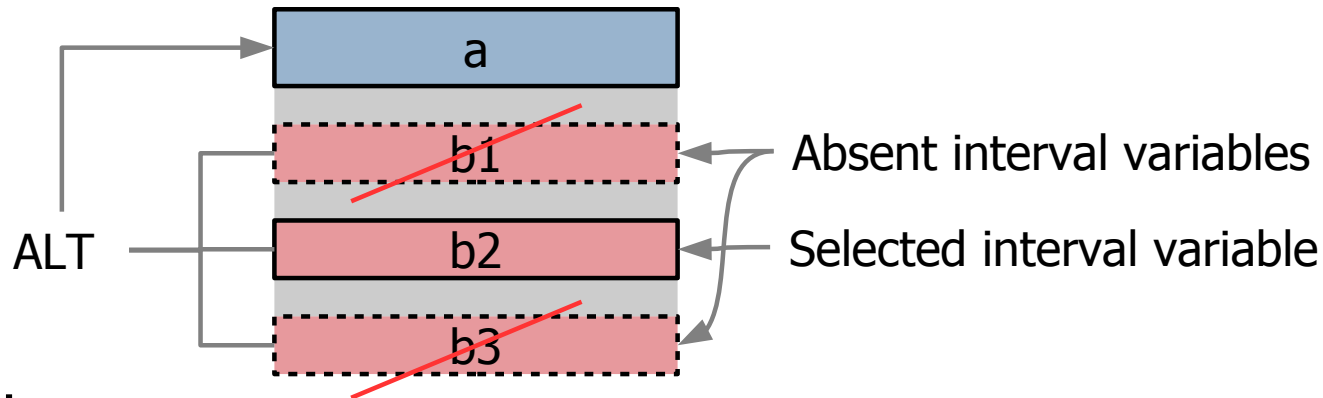


- Properties
  - The constraint of course handles optional interval variables

## Concept: **alternative constraint**

- What for?
  - Modeling alternative resource/modes/recipes
  - In general modeling a discrete selection in the schedule
- Example

**alternative(a, [b1, ..., bn])**

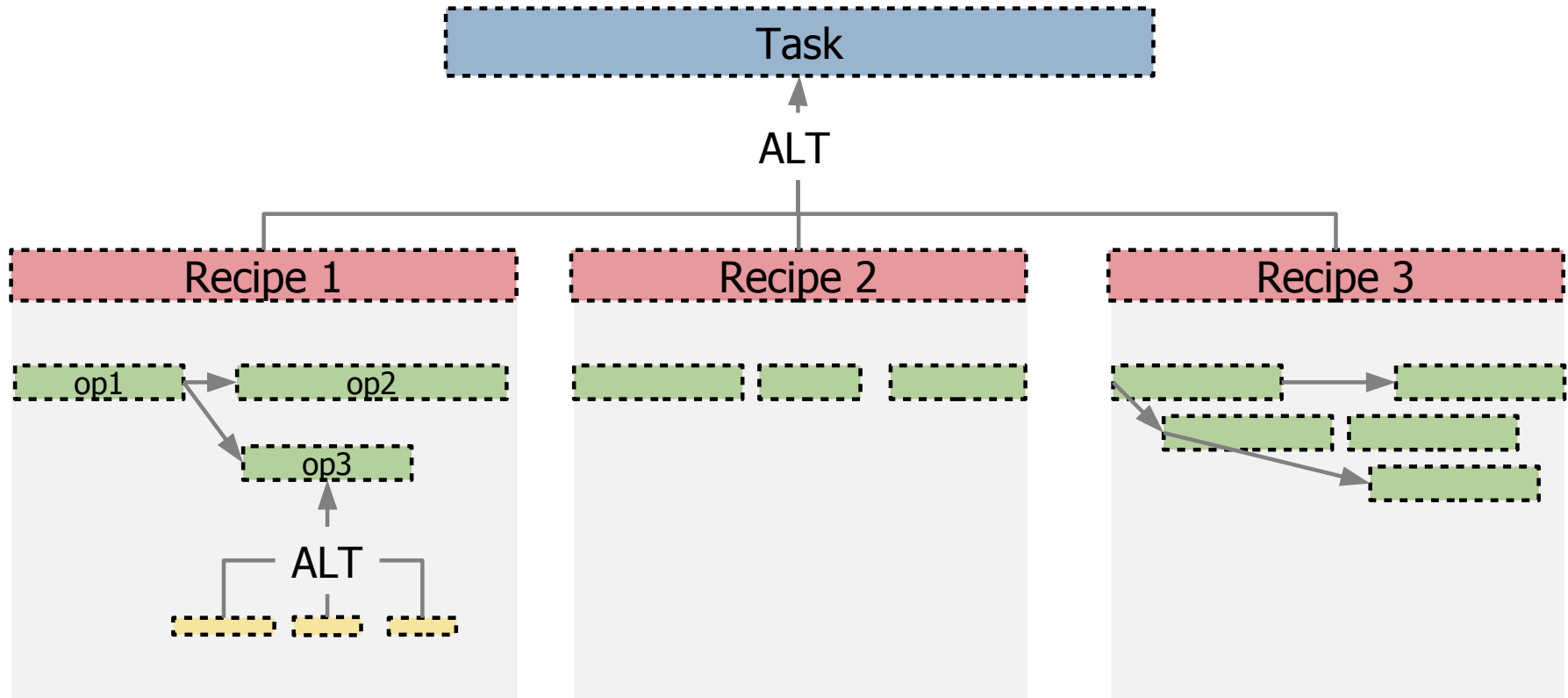


- Properties
  - Conditional reasoning allows a strong pruning on the alternative master interval variable *a*
  - Master interval variable *a* can of course be optional



## Concept: **span/alternative constraint**

- What for?
  - Modeling Work Breakdown Structure
- Example



# Agenda

- **IBM ILOG CP Optimizer** for **Scheduling**
- Scheduling **concepts** in CP Optimizer
- **Tips** and **Tools**
- **Q&A**

## Tips and Tools

- A measure of complexity of a model:
  - **Number of variables**
  - Number of groups of variables with different semantics
  - Number of variable types (integer variables / interval variables)
- Take advantage of the expressiveness of the CP Optimizer modeling language to compact the model by decreasing the above indicators
- In particular if in a scheduling model, you are creating 1 decision variable for each time point  $t$  you are usually in bad shape
- Exploit the scheduling concepts to reason on time-lines (sequence variables, cumul functions, state functions) and avoid an explicit representation of time

## Tips and Tools

- A measure of complexity of a model is the **number of constraints**
- If the number of constraints grows more than linearly with the number of variables or the size of variables domains, the model will probably not scale well
- Furthermore, if a set of many small constraints can be reformulated more compactly, this often leads to stronger inference in the engine
- Example: when you need to model activities (and more generally intervals of time) that -under some conditions- cannot overlap, think of using sequence variables, noOverlap constraints and/or state functions

## Tips and Tools

- Be careful when using **composite constraints**
- In scheduling models, exploit optionality of interval variables
- Examples :



- $\text{presenceOf}(a) * \text{endOf}(a) + (1 - \text{presenceOf}(a)) * K$
- $\text{presenceOf}(a) \Rightarrow (10 \leq \text{startOf}(a))$



- USE:  $\text{endOf}(a, K)$
- USE:  $10 \leq \text{startOf}(a, 10)$

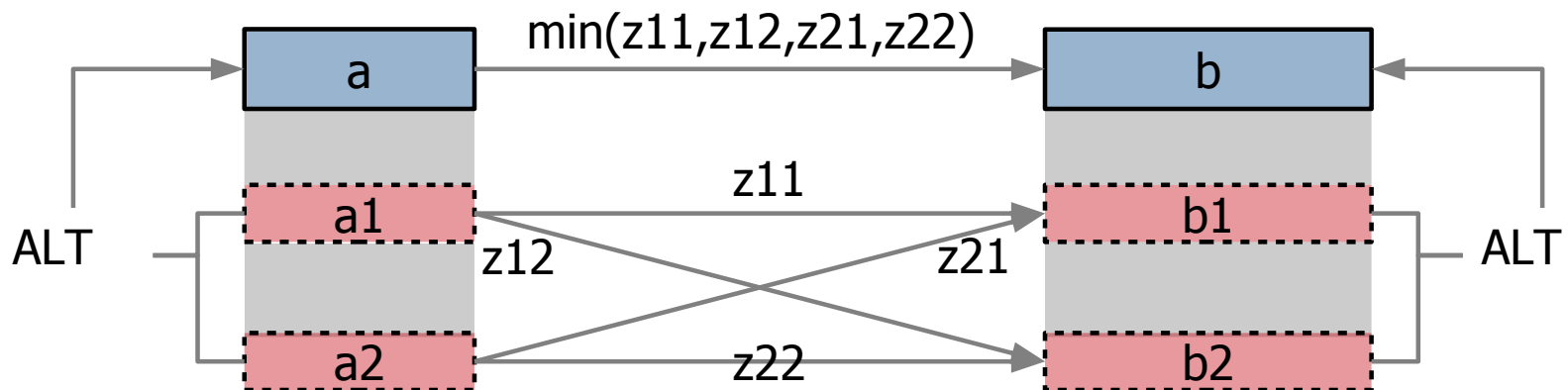
- Remarks:
  - **Exception:** Binary logical constraints on presence status (like  $\text{presenceOf}(a) \Rightarrow \text{presenceOf}(b)$ ) are handled in a special (and efficient) way in the engine
  - Except for constraint “presenceOf”, all constraints on scheduling constructs (interval and sequence variables, cumul and state functions) cannot be used in composite constraints

## Tips and Tools

- Surrogate constraints **may** help to reduce the search space
- A surrogate constraint is a constraint that does not change the set of solutions (but may be useful for the search by allowing a better pruning of the search space)
- Some examples :
  - Redundant cumul function for alternative resources

## Tips and Tools

- Surrogate constraints **may** help to reduce the search space
- A surrogate constraint is a constraint that does not change the set of solutions (but may be useful for the search by allowing a better pruning of the search space)
- Some examples :
  - Redundant cumul function for alternative resources
  - Redundant precedence constraints



## Tips and Tools

- Some classical patterns in scheduling models
  - Chain of optional interval variables
  - Optional intervals in paths of precedence constraints
  - Multilevel alternatives
  - Temporal alternatives
- *See bonus material at the end of the slide deck*



## Tips and Tools

- Questions:
  - Why is my model infeasible ?
  - Search does not find any feasible solution, what should I do ?
  - Search does not converge fast enough, what should I do ?
  - Does my model improvement really help?

# Tips and Tools

- Why is my model infeasible ?
  - Conflict refiner → Fix/relax model

The screenshot displays the IBM ILOG CPLEX Studio interface. The left pane shows the model file `sched_jobshop.mod` with the following code:

```

19 dvar interval op[j in Jobs][p in Pos] size Ops[j][p].pt;
20 dvar sequence mchs[m in Mchs] in
21   all(j in Jobs, p in Pos : Ops[j][p].mch == m) op[j][p];
22 dexpr int makespan = max(j in Jobs) endOf(op[j][nbPos-1]);
23 minimize makespan;
24 subject to {
25   makespan <= 300;
26   forall (m in Mchs)
27     machine: noOverlap(mchs[m]);
28   forall (j in Jobs, p in 1..nbPos-1)
29     precedence: endBeforeStart(op[j][p-1], op[j][p]);
30 }
  
```

The right pane shows the `Engine log` with the following output:

```

*          48          6
*          49          6
*          50          6
*          51          6
! Conflict refining terminated
! -----
! Conflict status      : Terminated normally, conflict found
! Conflict size       : 6 constraints
! Number of iterations : 51
! Total memory usage  : 1.1 MB
! Conflict computation time : 0.04s
! -----
  
```

The bottom pane shows the `Conflicts` tab with the following table:

Line	In conflict	Element (6)
28	Yes	precedence[1][4]
28	Yes	precedence[1][5]
28	Yes	precedence[1][6]
28	Yes	precedence[1][7]
28	Yes	precedence[1][8]
28	Yes	precedence[1][9]

# Tips and Tools

- The engine log is an essential tool for understanding how the engine behaves on a given model

```

! -----
! Minimization problem - 110 variables, 100 constraints
! LogPeriod          = 100,000
! Workers            = 2
! Initial process time : 0.00s (0.00s extraction + 0.00s propagation)
!   . Log search space : 664.4 (before), 664.4 (after)
!   . Memory usage     : 1.1 MB (before), 1.2 MB (after)
! Using parallel search with 2 workers.
! -----
!
!      Best Branches  Non-fixed   W      Branch decision
*      1,062         739 0.01s      1      -
...
*      936          136k 3.69s      1      -
      936          300k      85      2      F      on op#0#9
! Time = 4.05s, Explored branches = 450,928, Memory usage = 7.7 MB
!      Best Branches  Non-fixed   W      Branch decision
*      930          166k 4.55s      1      -
      930          200k      110     1      F      on op#5#1
! -----
! Search terminated normally, 17 solutions found.
! Best objective      : 930 (optimal - effective tol. is 0)
! Number of branches  : 641,352
! Number of fails     : 266,640
! Total memory usage   : 6.9 MB (6.4 MB CP Optimizer + 0.5 MB Concert)
! Time spent in solve  : 5.99s (5.99s engine + 0.00s extraction)
! Search speed (br. / s) : 107,070.5
! -----

```

Problem  
characteristics

## Tips and Tools

- The engine log is an essential tool for understanding how the engine behaves on a given model

```

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! Minimization problem - 110 variables, 100 constraints
! LogPeriod          = 100,000
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! -----

```

Modified  
parameter values

# Tips and Tools

- The engine log is an essential tool for understanding how the engine behaves on a given model

```

! -----
! Minimization problem - 110 variables, 100 constraints
! LogPeriod           = 100,000
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!   . Log search space : 664.4 (before), 664.4 (after)
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! -----
!
!      Best Branches  Non-fixed   W      Branch decision
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...
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      936          300k      85      2      F      on op#0#9
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! Search speed (br. / s) : 107,070.5
! -----

```

Root node  
information

## Tips and Tools

- The engine log is an essential tool for understanding how the engine behaves on a given model

```

! -----
! Minimization problem - 110 variables, 100 constraints
! LogPeriod           = 100,000
! Workers             = 2
! Initial process time : 0.00s (0.00s extraction + 0.00s propagation)
!   . Log search space : 664.4 (before), 664.4 (after)
!   . Memory usage     : 1.1 MB (before), 1.2 MB (after)
! Using parallel search with 2 workers.
! -----
!
!      Best Branches  Non-fixed   W      Branch decision
*      1,062         739 0.01s      1      -
...
*      936          136k 3.69s      1      -
      936          300k      85      2      F      on op#0#9
! Time = 4.05s, Explored branches = 450,928, Memory usage = 7.7 MB
!      Best Branches  Non-fixed   W      Branch decision
*      930          166k 4.55s      1      -
      930          200k      110     1      F      on op#5#1
! -----
! Search terminated normally, 17 solutions found.
! Best objective       : 930 (optimal - effective tol. is 0)
! Number of branches   : 641,352
! Number of fails      : 266,640
! Total memory usage    : 6.9 MB (6.4 MB CP Optimizer + 0.5 MB Concert)
! Time spent in solve   : 5.99s (5.99s engine + 0.00s extraction)
! Search speed (br. / s) : 107,070.5
! -----

```

New incumbent  
solutions (time, worker)

## Tips and Tools

- The engine log is an essential tool for understanding how the engine behaves on a given model

```

! -----
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! Search speed (br. / s) : 107,070.5
! -----

```

Periodical log  
with fail information,  
number of unfixed  
variables, current decision

## Tips and Tools

- The engine log is an essential tool for understanding how the engine behaves on a given model

```

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! LogPeriod           = 100,000
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! Search speed (br. / s) : 107,070.5
! -----

```

Final information  
with solution status  
and search statistics



## Tips and Tools

- Search does not find any feasible solution, what should I do ?
  - Identify which part of the model is hard:
    - Simplify problem by removing “easy” constraints
    - Search log helps (set LogPeriod=1, look where first fails occur)
  - Strengthen the model:
    - Be careful with composite constraints
    - Add surrogate constraints
    - Break symmetries
  - Increase inference levels in the engine

## Tips and Tools

- Search does not find any feasible solution, what should I do ?
  - Try using search phases to first focus on the difficult decision variables
  - Solve the problem as two sub-problems:
    1. Finding an initial feasible solution (some difficult constraints can be relaxed and put in the objective function\*), then
    2. Optimize starting from this solution. Consider using starting points.
  - Isolate the different sources of complexity in the model by decomposition and sequential resolutions. Consider using starting points

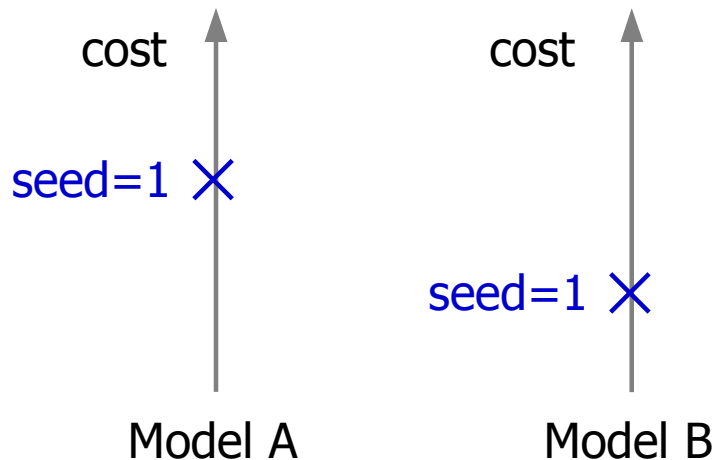
(\*) CP Optimizer's concepts makes it easy to relax part of the model. E.g. make some interval variables *optional* and maximize number of present intervals, relax *deadlines* as *tardiness* costs

## Tips and Tools

- Search does not converge fast enough, what should I do ?
  - Identify and try to **isolate the different sources of complexity** in your model, first focusing on bottleneck/important parts in a “simplified” model
  - Exploit results of previous solve as (1) additional constraints and/or (2) starting point and/or (3) heuristic (search phases)
  - Do not hesitate to use MIP models in some steps, exploit strengths of each technology (MIP/CP)
    - For instance CP is good at deciding activities **start/end times** under resource constraints when required resource quantities and activity durations are fixed or mostly depend on activity start/end dates
    - For problems with very variable **resource quantities** or **activity durations**, it makes sense to consider an initial MIP model that works on these variables (example: lot sizing)

## Tips and Tools

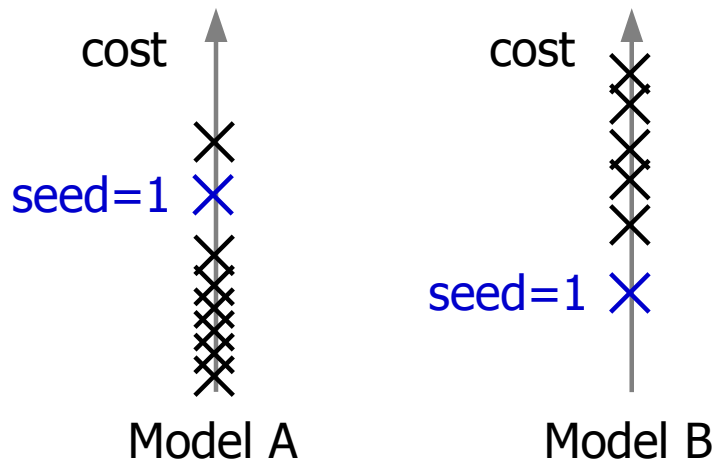
- Does my model improvement really help?
  - CP Optimizer's search relies on a stochastic algorithm initialized by a random seed (parameter IloCP::RandomSeed)
  - For a given model, the solution depends on the random seed
  - Best solution cost at a given time limit:



–Is model B better than model A ?      Yes

## Tips and Tools

- Does my model improvement really help?
  - CP Optimizer's search relies on a stochastic algorithm initialized by a random seed (parameter IloCP::RandomSeed)
  - For a given model, the solution depends on the random seed
  - Best solution cost at a given time limit:



Never evaluate a  
model change on  
a single instance/seed !

**No!**

– Is model B better than model A ?

~~Yes~~

## Conclusion

- For scheduling problems, CP Optimizer extends the **Model&run** paradigm with some (**few**) general concepts (interval & sequence variables, cumul & state functions, ...)
- These concepts can be used to model scheduling problems in a compact way that **avoids enumeration of time**. Typically, the size of a scheduling model in CP Optimizer grows linearly with the size of the data
- When combined together these concepts provide an expressive language capable of handling **complex** and **large** industrial scheduling problems
- The **automatic search** is **complete**, it exploits these concepts using an **efficient** and **robust** combination of techniques from scheduling theory, meta-heuristics, CP, MP

-

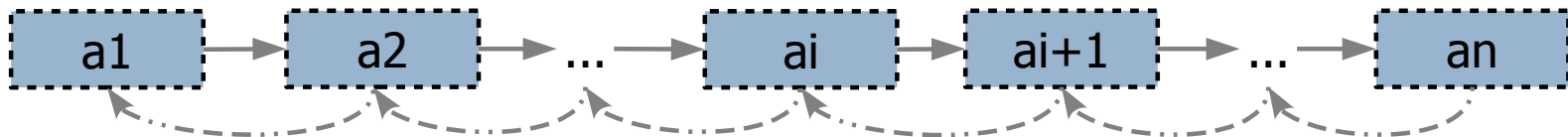
## Bonus material

- List of most useful parameters for search analysis:
  - Workers (=1)
  - LogPeriod (=1)
  - SearchType (=DepthFirst)
  - FailLimit, SolutionLimit (=1)
- List of most useful parameters for performance tuning:
  - TimeLimit
  - TemporalRelaxation (=Off)
  - Workers
  - SearchType (=Restart|MultiPoint)
  - Some inference levels:
    - NoOverlapInferenceLevel
    - CumulFunctionInferenceLevel



## Bonus material: classical patterns in scheduling models

- Chain of optional interval variables
  - Use cases: any phenomenon that is associated with a set of consecutive time intervals whose number is not known a-priori
  - Examples: partially preemptive activities, intervals during which a resource is in use, ...
- Model is a chain of precedence and implication constraints between optional interval variables

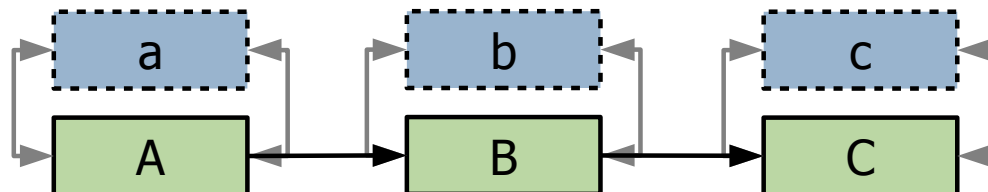


```
endBeforeStart(a[i], a[i+1]);  
presenceOf(a[i+1]) => presenceOf(a[i]);
```

- This ensures that only the first  $k$  intervals will be present where  $k$  is an implicit decision of the problem

## Bonus material: classical patterns in scheduling models

- Optional intervals in paths of precedence constraints
  - Use case: a precedence graph with optional activities such that path of precedence constraints going through absent activities should still be considered
  - Example:  $a \rightarrow b \rightarrow c$ , precedence  $a \rightarrow c$  should be considered even if  $b$  is absent
  - Direct model using  $\text{endBeforeStart}(a,b)$ ,  $\text{endBeforeStart}(b,c)$  does not work as both constraints are inactive if  $b$  is absent
- Idea 1: Add  $\text{endBeforeStart}(x,y)$  constraints of the transitive closure. Issue: may add a quadratic number of constraints
- Idea 2: Use additional non-optional interval variables  $A, B, C$  of free length synchronized with  $a, b, c$  ( $\text{startAtStart}$ ,  $\text{endAtEnd}$ )

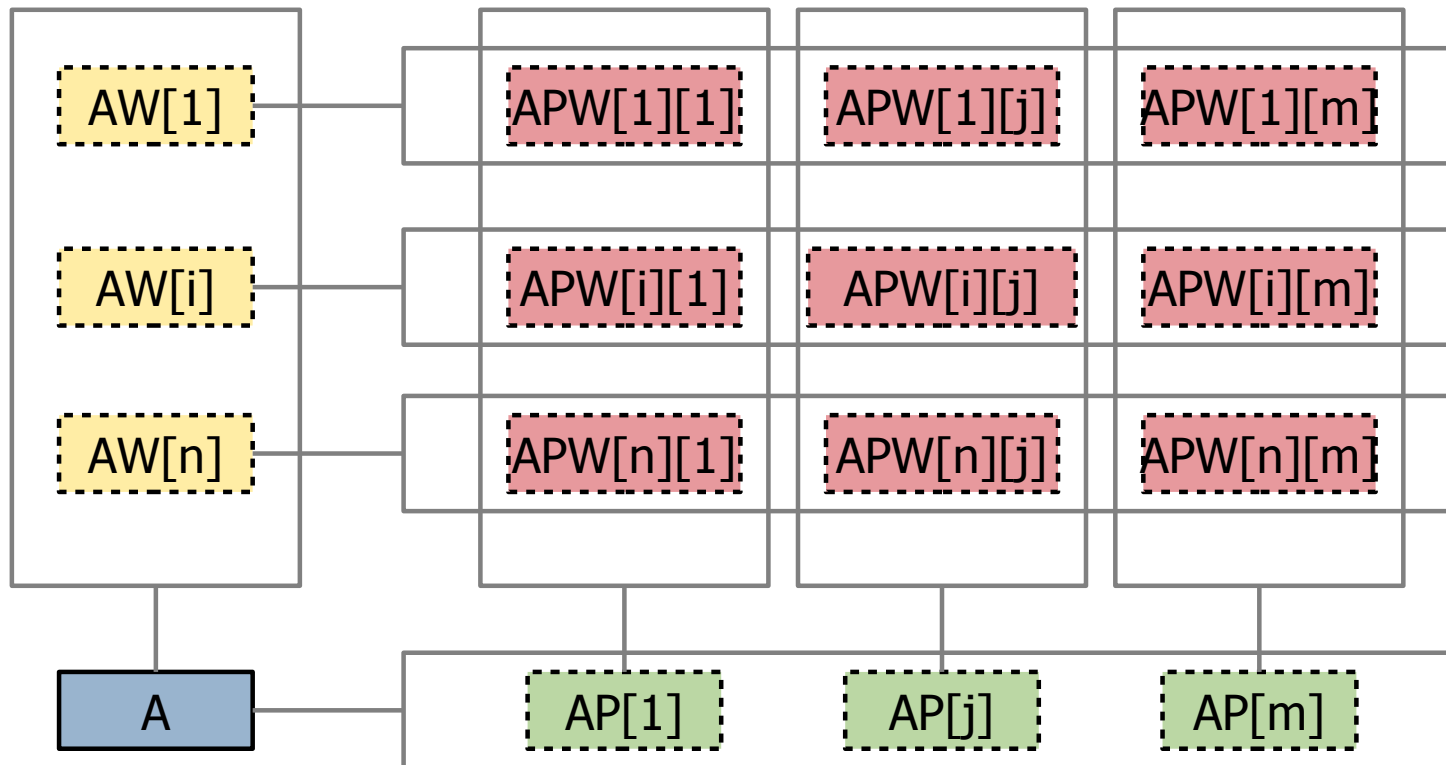


## Bonus material: classical patterns in scheduling models

- Multi-level alternatives
  - Use case: an activity needs to select an element in a Cartesian product
  - Example: an activity A needs 1 worker  $W_i$  among  $n$  to be executed at 1 given position  $P_j$  among  $m$ . Duration depends both on the worker  $W_i$  and the position  $P_j$ . Specific constraints need to be posted on all activities executing at a position  $P_j$  (e.g. space limitation) while other constraints need to be posted on all activities executed by a given worker  $W_i$  (e.g. noOverlap)
- Idea: create one optional interval variable  $AW[i]$  for activity A executing on worker  $W_i$  and one optional interval variable  $AP[j]$  for activity A executing a position  $P_j$

## Bonus material: classical patterns in scheduling models

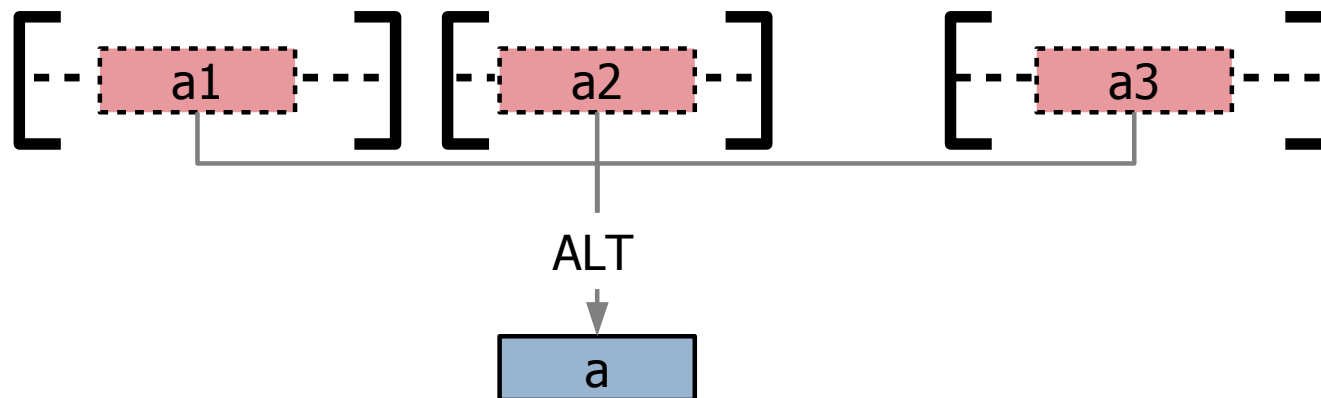
- Multi-level alternatives
- Idea: create one optional interval variable  $AW[i]$  for activity  $A$  executing on worker  $W_i$  and one optional interval variable  $AP[j]$  for activity  $A$  executing at position  $P_j$



## Bonus material: classical patterns in scheduling models

- Temporal alternatives

- Alternative constraints are not restricted to model disjunctions of resources/modes/recipes, they can also be used to represent temporal disjunction
- Use-case: an activity must execute during a disjunctive set of possible time-windows



- Alternative models that directly work on interval variable *a* are possible (use `cumulFunction` or `forbidExtent`). They will be lighter (less variables) but will propagate less.

## Bonus material

### ▪ Technical papers related with scheduling in CP Optimizer

- P. Laborie, J. Rogerie. *Temporal Linear Relaxation in IBM ILOG CP Optimizer*. Proc. MISTA 2013. Ghent, Belgium. Aug. 2013.
- P. Laborie, J. Rogerie, P. Shaw, P. Vilím. *Reasoning with Conditional Time-intervals, Part II: an Algebraical Model for Resources*. Proc. FLAIRS 2009.
- P. Vilím. *Max Energy Filtering Algorithm for Discrete Cumulative Resources*. Proc. CPAIOR 2009. p294-308.
- P. Laborie. *IBM ILOG CP Optimizer for Detailed Scheduling Illustrated on Three Problems*. Proc. CPAIOR 2009. p148-162.
- P. Vilím. *Edge Finding Filtering Algorithm for Discrete Cumulative Resources in  $O(kn \log n)$* . Proc. CP 2009. p802-816.
- P. Laborie, J. Rogerie. *Reasoning with Conditional Time-intervals*. Proc. FLAIRS 2008, p555-560.
- P. Laborie, D. Godard. *Self-Adapting Large Neighborhood Search: Application to single-mode scheduling problems*. Proc. MISTA 2007. Paris, France. Aug. 2007.
- D. Godard, P. Laborie, W. Nuijten. *Randomized Large Neighborhood Search for Cumulative Scheduling*. Proc. ICAPS-05. Monterey, California (USA). June 2005.