

Languages-beta: OC-L-07-Expressions

The P_LanCompS Project

Languages-beta/OC-L/OC-L-07-Expressions/OC-L-07-Expressions.cbs*

Language "OCaml Light"

*Suggestions for improvement: plancomps@gmail.com.
Issues: <https://github.com/plancomps/CBS-beta/issues>.

7 Expressions

```
Syntax  $E : \text{expr} ::=$  value-path  
| constant  
| ( expr )  
| begin expr end  
| ( expr : typexpr )  
| expr comma-expr+  
| expr :: expr  
| [ expr semic-expr* ]  
| [ expr semic-expr* ; ]  
| [ | expr semic-expr* | ]  
| [ | expr semic-expr* ; | ]  
| { field = expr semic-field-expr* }  
| { field = expr semic-field-expr* ; }  
| { expr with field = expr semic-field-expr* }  
| { expr with field = expr semic-field-expr* ; }  
| expr argument+  
| prefix-symbol expr  
| - expr  
| -. expr  
| expr infix-op-1 expr  
| expr infix-op-2 expr  
| expr infix-op-3 expr  
| expr infix-op-4 expr  
| expr infix-op-5 expr  
| expr infix-op-6 expr  
| expr infix-op-7 expr  
| expr infix-op-8 expr  
| expr . field  
| expr .( expr )  
| expr .( expr ) <- expr  
| if expr then expr (else expr)?  
| while expr do expr done  
| for value-name = expr (to | downto) expr do expr done  
| expr ; expr  
| match expr with pattern-matching  
| function pattern-matching  
| fun pattern+ -> expr  
| try expr with pattern-matching  
| let-definition in expr  
| assert expr
```

$A : \text{argument} ::=$ expr

$PM : \text{pattern-matching} ::=$ pattern -> expr pattern-expr*

Rule $\llbracket (E) \rrbracket : \text{expr} =$
 $\llbracket E \rrbracket$
Rule $\llbracket \text{begin } E \text{ end} \rrbracket : \text{expr} =$
 $\llbracket E \rrbracket$
Rule $\llbracket (E : T) \rrbracket : \text{expr} =$
 $\llbracket E \rrbracket$
Rule $\llbracket E_1 E_2 A A^* \rrbracket : \text{expr} =$
 $\llbracket ((E_1 E_2)) A A^* \rrbracket$
Rule $\llbracket PS E \rrbracket : \text{expr} =$
 $\llbracket ((PS)) E \rrbracket$
Rule $\llbracket - E \rrbracket : \text{expr} =$
 $\llbracket ((\sim -)) E \rrbracket$
Rule $\llbracket -. E \rrbracket : \text{expr} =$
 $\llbracket ((\sim -.)) E \rrbracket$
Rule $\llbracket E_1 IO-1 E_2 \rrbracket : \text{expr} =$
 $\llbracket ((IO-1)) E_1 E_2 \rrbracket$
Rule $\llbracket E_1 IO-2 E_2 \rrbracket : \text{expr} =$
 $\llbracket ((IO-2)) E_1 E_2 \rrbracket$
Rule $\llbracket E_1 IO-3 E_2 \rrbracket : \text{expr} =$
 $\llbracket ((IO-3)) E_1 E_2 \rrbracket$
Rule $\llbracket E_1 IO-4 E_2 \rrbracket : \text{expr} =$
 $\llbracket ((IO-4)) E_1 E_2 \rrbracket$
Rule $\llbracket E_1 IO-5 E_2 \rrbracket : \text{expr} =$
 $\llbracket ((IO-5)) E_1 E_2 \rrbracket$
Rule $\llbracket E_1 \& E_2 \rrbracket : \text{expr} =$
 $\llbracket E_1 \&\& E_2 \rrbracket$
Rule $\llbracket E_1 \text{ or } E_2 \rrbracket : \text{expr} =$
 $\llbracket E_1 || E_2 \rrbracket$
Rule $\llbracket E_1 IO-8 E_2 \rrbracket : \text{expr} =$
 $\llbracket ((IO-8)) E_1 E_2 \rrbracket$
Rule $\llbracket E_1 .(E_2) \rrbracket : \text{expr} =$
 $\llbracket \text{array_get } E_1 E_2 \rrbracket$
Rule $\llbracket E_1 .(E_2) <- E_3 \rrbracket : \text{expr} =$
 $\llbracket \text{array_set } E_1 E_2 E_3 \rrbracket$
Rule $\llbracket \text{if } E_1 \text{ then } E_2 \rrbracket : \text{expr} =$
 $\llbracket \text{if } E_1 \text{ then } E_2 \text{ else } (()) \rrbracket$
Rule $\llbracket \text{fun } P \rightarrow E \rrbracket : \text{expr} =$
 $\llbracket \text{function } P \rightarrow E \rrbracket$
Rule $\llbracket \text{fun } P P^+ \rightarrow E \rrbracket : \text{expr} = 3$
 $\llbracket \text{fun } P \rightarrow (\text{fun } P^+ \rightarrow E) \rrbracket$
Rule $\llbracket [E SE^* ;] \rrbracket : \text{expr} =$
 $\llbracket [E SE^*] \rrbracket$
Rule $\llbracket [| E SE^* ; |] \rrbracket : \text{expr} =$
 $\llbracket [| E SE^* |] \rrbracket$
Rule $\llbracket \{ F = E SFE^* ; \} \rrbracket : \text{expr} =$

Semantics $\text{evaluate}[_ : \text{expr}] : \Rightarrow \text{implemented-values}$

Rule $\text{evaluate}[VP] =$

$\text{bound}(\text{value-name}[VP])$

Rule $\text{evaluate}[CNST] =$

$\text{value}[CNST]$

Rule $\text{evaluate}[(E : T)] =$

$\text{evaluate}[E]$

Rule $\text{evaluate}[E_1, E_2 CE^*] =$

$\text{tuple}(\text{evaluate-comma-sequence}[E_1, E_2 CE^*])$

Rule $\text{evaluate}[E_1 :: E_2] =$

$\text{cons}(\text{evaluate}[E_1],$
 $\text{evaluate}[E_2])$

Rule $\text{evaluate}[[E SE^*]] =$

$[\text{evaluate-semicolon-sequence}[E SE^*]]$

Rule $\text{evaluate}[[|E SE^*|]] =$

$\text{vector}(\text{left-to-right-map}(\text{allocate-initialised-variable}(\text{implemented-values},$
 $\text{given}),$
 $\text{evaluate-semicolon-sequence}[E SE^*]))$

Rule $\text{evaluate}[[| |]] =$

$\text{vector}()$

Rule $\text{evaluate}[\{F = E SFE^*\}] =$

$\text{record}(\text{collateral}(\text{evaluate-field-sequence}[F = E SFE^*]))$

Rule $\text{evaluate}[\{E_1 \text{ with } F = E_2 SFE^*\}] =$

$\text{record}(\text{map-override}(\text{evaluate-field-sequence}[F = E_2 SFE^*],$
 $\text{checked record-map}(\text{evaluate}[E_1])))$

Rule $\text{evaluate}[CSTR E] =$

$\text{variant}(\text{constr-name}[CSTR],$
 $\text{evaluate}[E])$

Otherwise $\text{evaluate}[E_1 E_2] =$

$\text{apply}(\text{evaluate}[E_1],$
 $\text{evaluate}[E_2])$

Rule $\text{evaluate}[E . F] =$

$\text{record-select}(\text{evaluate}[E],$
 $\text{field-name}[F])$

Rule $\text{evaluate}[E_1 \ \&\& \ E_2] =$

$\text{if-true-else}(\text{evaluate}[E_1],$
 $\text{evaluate}[E_2],$
 $\text{false})$

Rule $\text{evaluate}[E_1 \ || \ E_2] =$

$\text{if-true-else}(\text{evaluate}[E_1],$
 $\text{true},$
 $\text{evaluate}[E_2])$

Rule $\text{evaluate}[\text{if } E_1 \text{ then } E_2 \text{ else } E_3] =$

$\text{if-true-else}(\text{evaluate}[E_1],$
 $\text{evaluate}[E_2],$

Expression sequences and maps

Semantics `evaluate-comma-sequence` $\llbracket _ : (\text{expr comma-expr}^*) \rrbracket : (\Rightarrow \text{implemented-values})^+$

Rule `evaluate-comma-sequence` $\llbracket E_1 , E_2 CE^* \rrbracket =$

`evaluate` $\llbracket E_1 \rrbracket$,

`evaluate-comma-sequence` $\llbracket E_2 CE^* \rrbracket$

Rule `evaluate-comma-sequence` $\llbracket E \rrbracket =$

`evaluate` $\llbracket E \rrbracket$

Semantics `evaluate-semicolon-sequence` $\llbracket _ : (\text{expr semic-expr}^*) \rrbracket : (\Rightarrow \text{implemented-values})^+$

Rule `evaluate-semicolon-sequence` $\llbracket E_1 ; E_2 SE^* \rrbracket =$

`evaluate` $\llbracket E_1 \rrbracket$,

`evaluate-semicolon-sequence` $\llbracket E_2 SE^* \rrbracket$

Rule `evaluate-semicolon-sequence` $\llbracket E \rrbracket =$

`evaluate` $\llbracket E \rrbracket$

Semantics `evaluate-field-sequence` $\llbracket _ : (\text{field} = \text{expr semic-field-expr}^*) \rrbracket : (\Rightarrow \text{envs})^+$

Rule `evaluate-field-sequence` $\llbracket F_1 = E_1 ; F_2 = E_2 SFE^* \rrbracket =$

$\{\text{field-name} \llbracket F_1 \rrbracket \mapsto \text{evaluate} \llbracket E_1 \rrbracket\}$,

`evaluate-field-sequence` $\llbracket F_2 = E_2 SFE^* \rrbracket$

Rule `evaluate-field-sequence` $\llbracket F = E \rrbracket =$

$\{\text{field-name} \llbracket F \rrbracket \mapsto \text{evaluate} \llbracket E \rrbracket\}$

Matching

Semantics `match` $\llbracket _ : \text{pattern-matching} \rrbracket : (\text{implemented-values} \Rightarrow \text{implemented-values})^+$

Rule `match` $\llbracket P_1 \rightarrow E_1 \mid P_2 \rightarrow E_2 PE^* \rrbracket =$

`match` $\llbracket P_1 \rightarrow E_1 \rrbracket$,

`match` $\llbracket P_2 \rightarrow E_2 PE^* \rrbracket$

Rule `match` $\llbracket P \rightarrow E \rrbracket =$

`case-match`(`evaluate-pattern` $\llbracket P \rrbracket$,

`evaluate` $\llbracket E \rrbracket$)

Value definitions

Semantics `define-values` $\llbracket _ : \text{let-definition} \rrbracket : \Rightarrow \text{environments}$

Rule `define-values` $\llbracket \text{let } LB \text{ } ALB^* \rrbracket =$

`define-values-nonrec` $\llbracket LB \text{ } ALB^* \rrbracket$

Rule `define-values` $\llbracket \text{let rec } LB \text{ } ALB^* \rrbracket =$

`recursive`(`set`(`bound-ids-sequence` $\llbracket LB \text{ } ALB^* \rrbracket$),

`define-values-nonrec` $\llbracket LB \text{ } ALB^* \rrbracket$)

Semantics `define-values-nonrec` $\llbracket _ : (\text{let-binding and-let-binding}^*) \rrbracket : \Rightarrow \text{environments}$

Rule `define-values-nonrec` $\llbracket LB_1 \text{ and } LB_2 \text{ } ALB^* \rrbracket =$

`collateral`(`define-values-nonrec` $\llbracket LB_1 \rrbracket$,

`define-values-nonrec` $\llbracket LB_2 \text{ } ALB^* \rrbracket$)

Rule `define-values-nonrec` $\llbracket P = E \rrbracket =$

`else`(`match`(`evaluate` $\llbracket E \rrbracket$,

`evaluate-pattern` $\llbracket P \rrbracket$),

`throw`(`ocaml-light-match-failure`))

Semantics `bound-ids-sequence` $\llbracket _ : (\text{let-binding and-let-binding}^*) \rrbracket : \text{ids}^+$

Rule `bound-ids-sequence` $\llbracket LB \rrbracket =$

`bound-id` $\llbracket LB \rrbracket$

Rule `bound-ids-sequence` $\llbracket LB_1 \text{ and } LB_2 \text{ } ALB^* \rrbracket =$

`bound-id` $\llbracket LB_1 \rrbracket$,

`bound-ids-sequence` $\llbracket LB_2 \text{ } ALB^* \rrbracket$

Semantics `bound-id` $\llbracket _ : \text{let-binding} \rrbracket : \text{ids}$

Rule `bound-id` $\llbracket VN = E \rrbracket =$

`value-name` $\llbracket VN \rrbracket$

Otherwise `bound-id` $\llbracket LB \rrbracket =$

`fail`