Funcons-beta: Binding

The PLanCompS Project

`Funcons-beta/Computations/Normal/Binding/Binding.cbs`*

Binding

[*Type* environments
Alias envs
Datatype identifiers
Alias ids
Funcon identifier-tagged
Alias id-tagged
Funcon fresh-identifier
Entity environment
Alias env
Funcon initialise-binding
Funcon bind-value
Alias bind
Funcon unbind
Funcon bound-directly
Funcon bound-value
Alias bound
Funcon closed
Funcon scope
Funcon accumulate
Funcon collateral
Funcon bind-recursively
Funcon recursive]

Meta-variables T <: values

*Suggestions for improvement: `plancomps@gmail.com`.
Issues: `https://github.com/plancomps/CBS-beta/issues`.

Environments

Type environments $\rightsquigarrow$ maps(identifiers, values?)

Alias envs = environments

An environment represents bindings of identifiers to values. Mapping an identifier to ( ) represents that its binding is hidden.

Circularity in environments (due to recursive bindings) is represented using bindings to cut-points called links. Funcons are provided for making declarations recursive and for referring to bound values without explicit mention of links, so their existence can generally be ignored.

Datatype identifiers ::= {_ : strings}
| identifier-tagged(_ : identifiers, _ : values)

Alias ids = identifiers

Alias id-tagged = identifier-tagged

An identifier is either a string of characters, or an identifier tagged with some value (e.g., with the identifier of a namespace).

Funcon fresh-identifier : $\Rightarrow$ identifiers

fresh-identifier computes an identifier distinct from all previously computed identifiers.

Rule fresh-identifier $\rightsquigarrow$ identifier-tagged("`generated`", fresh-atom)

Current bindings

Entity environment(_ : environments) $\vdash$ _ $\longrightarrow$ _

Alias env = environment

The environment entity allows a computation to refer to the current bindings of identifiers to values.

Funcon initialise-binding(X : $\Rightarrow$ T) : $\Rightarrow$ T
$\rightsquigarrow$ initialise-linking(initialise-generating(closed(X)))

initialise-binding(X) ensures that X does not depend on non-local bindings. It also ensures that the linking entity (used to represent potentially cyclic bindings) and the generating entity (for creating fresh identifiers) are initialised.

Funcon bind-value(I : identifiers, V : values) : $\Rightarrow$ environments
$\rightsquigarrow \{I \mapsto V\}$

Alias bind = bind-value

bind-value(I, X) computes the environment that binds only I to the value computed by X.

Funcon unbind(I : identifiers) : $\Rightarrow$ environments
$\rightsquigarrow \{I \mapsto ( )\}$

unbind(I) computes the environment that hides the binding of I.

Funcon bound-directly(_ : identifiers) : $\Rightarrow$ values

bound-directly(I) returns the value to which I is currently bound, if any, and otherwise fails.

bound-directly(I) does *not* follow links. It is used only in connection with recursively-bound values when references are not encapsulated in abstractions.

Rule $$\frac{\text{lookup}(\rho, I) \rightsquigarrow (V : \text{values})}{\text{environment}(\rho) \vdash \text{bound-directly}(I : \text{identifiers}) \longrightarrow V}$$

Rule $$\frac{\text{lookup}(\rho, I) \rightsquigarrow ( )}{\text{environment}(\rho) \vdash \text{bound-directly}(I : \text{identifiers}) \longrightarrow \text{fail}}$$

Funcon bound-value(I : identifiers) : $\Rightarrow$ values
$\rightsquigarrow$ follow-if-link(bound-directly(I))

Alias bound = bound-value

bound-value(I) inspects the value to which I is currently bound, if any, and otherwise fails. If the value is a link, bound-value(I) returns the value obtained by following the link, if any, and otherwise fails. If the inspected value is not a link, bound-value(I) returns it.

bound-value(I) is used for references to non-recursive bindings and to recursively-bound values when references are encapsulated in abstractions.

Scope

Funcon $\mathsf{closed}(X :\Rightarrow T) :\Rightarrow T$

$\mathsf{closed}(X)$ ensures that X does not depend on non-local bindings.

$$\textit{Rule}\ \frac{\mathsf{environment}(\mathsf{map}(\ )) \vdash X \longrightarrow X'}{\mathsf{environment}(\_) \vdash \mathsf{closed}(X) \longrightarrow \mathsf{closed}(X')}$$

$$\textit{Rule}\ \mathsf{closed}(V : T) \rightsquigarrow V$$

Funcon $\mathsf{scope}(\_ : \mathsf{environments}, \_ :\Rightarrow T) :\Rightarrow T$

$\mathsf{scope}(D, X)$ executes D with the current bindings, to compute an environment ρ representing local bindings. It then executes X to compute the result, with the current bindings extended by ρ, which may shadow or hide previous bindings.

$\mathsf{closed}(\mathsf{scope}(\rho, X))$ ensures that X can reference only the bindings provided by ρ.

$$\textit{Rule}\ \frac{\mathsf{environment}(\mathsf{map\text{-}override}(\rho_1, \rho_0)) \vdash X \longrightarrow X'}{\mathsf{environment}(\rho_0) \vdash \mathsf{scope}(\rho_1 : \mathsf{environments}, X) \longrightarrow \mathsf{scope}(\rho_1, X')}$$

$$\textit{Rule}\ \mathsf{scope}(\_ : \mathsf{environments}, V : T) \rightsquigarrow V$$

Funcon $\mathsf{accumulate}(\_ : (\Rightarrow \mathsf{environments})^*) :\Rightarrow \mathsf{environments}$

$\mathsf{accumulate}(D_1, D_2)$ executes D_1 with the current bindings, to compute an environment ρ_1 representing some local bindings. It then executes D_2 to compute an environment ρ_2 representing further local bindings, with the current bindings extended by ρ_1, which may shadow or hide previous current bindings. The result is ρ_1 extended by ρ_2, which may shadow or hide the bindings of ρ_1.

$\mathsf{accumulate}(\_, \_)$ is associative, with $\mathsf{map}(\ )$ as unit, and extends to any number of arguments.

$$\textit{Rule}\ \frac{D_1 \longrightarrow D_1'}{\mathsf{accumulate}(D_1, D_2) \longrightarrow \mathsf{accumulate}(D_1', D_2)}$$

$$\textit{Rule}\ \mathsf{accumulate}(\rho_1 : \mathsf{environments}, D_2) \rightsquigarrow \mathsf{scope}(\rho_1, \mathsf{map\text{-}override}(D_2, \rho_1))$$

$$\textit{Rule}\ \mathsf{accumulate}(\ ) \rightsquigarrow \mathsf{map}(\ )$$

$$\textit{Rule}\ \mathsf{accumulate}(D_1) \rightsquigarrow D_1$$

$$\textit{Rule}\ \mathsf{accumulate}(D_1, D_2, D^+) \rightsquigarrow \mathsf{accumulate}(D_1, \mathsf{accumulate}(D_2, D^+))$$

Funcon collateral(ρ^* : environments *) : $\Rightarrow$ environments
$\rightsquigarrow$ checked map-unite(ρ^*)

collateral($D_1, \cdots$) pre-evaluates its arguments with the current bindings, and unites the resulting maps, which fails if the domains are not pairwise disjoint.

collateral(D_1, D_2) is associative and commutative with map( ) as unit, and extends to any number of arguments.

Recurse

Funcon bind-recursively(I : identifiers, E : $\Rightarrow$ values) : $\Rightarrow$ environments
$\rightsquigarrow$ recursive({I},
bind-value(I,
E))

bind-recursively(I, E) binds I to a link that refers to the value of E, representing a recursive binding of I to the value of E. Since bound-value(I) follows links, it should not be executed during the evaluation of E.

Funcon recursive(SI : sets(identifiers), D : $\Rightarrow$ environments) : $\Rightarrow$ environments
$\rightsquigarrow$ re-close(bind-to-forward-links(SI),
D)

recursive(SI, D) executes D with potential recursion on the bindings of the identifiers in the set SI (which need not be the same as the set of identifiers bound by D).

Auxiliary Funcon re-close(M : maps(identifiers, links), D : $\Rightarrow$ environments) : $\Rightarrow$ environments
$\rightsquigarrow$ accumulate(scope(M,
D),
sequential(set-forward-links(M),
map( )))

re-close(M, D) first executes D in the scope M, which maps identifiers to freshly allocated links. This computes an environment ρ where the bound values may contain links, or implicit references to links in abstraction values. It then sets the link for each identifier in the domain of M to refer to its bound value in ρ, and returns ρ as the result.

Auxiliary Funcon bind-to-forward-links(*SI* : sets(identifiers)) : ⇒ maps(identifiers, links)
⇝ map-unite(interleave-map(bind-value(given,
fresh-link(values)),
set-elements(*SI*)))

bind-to-forward-links(*SI*) binds each identifier in the set *SI* to a freshly allocated link.

Auxiliary Funcon set-forward-links(*M* : maps(identifiers, links)) : ⇒ null-type
⇝ effect(interleave-map(set-link(map-lookup(*M*,
given),
bound-value(given)),
set-elements(map-domain(*M*))))

For each identifier *I* in the domain of *M*, set-forward-links(*M*) sets the link to which *I* is mapped by *M* to the current bound value of *I*.