## Funcons-beta: Abrupting \*

## The PLanCompS Project

Abrupting.cbs | PLAIN | PRETTY

## **Abruptly terminating**

```
[ Funcon stuck
   Entity abrupted
   Funcon finalise-abrupting
   Funcon abrupt
   Funcon handle-abrupt
   Funcon finally ]

Meta-variables T, T', T" <: values

Funcon stuck : ⇒ empty-type
```

stuck does not have any computation. It is used to represent the result of a transition that causes the computation to terminate abruptly.

```
Entity \underline{\hspace{1cm}}^{\text{abrupted}(\underline{\hspace{1cm}}:values?)}
```

abrupted(V) in a label on a transstion indicates abrupt termination for reason V. abrupted() indicates the absence of abrupt termination.

```
Funcon finalise-abrupting(X : \Rightarrow T) : \Rightarrow T | null-type \rightarrow handle-abrupt(X, null-value)
```

finalise-abrupting (X) handles abrupt termination of X for any reason.

```
Funcon abrupt(\_: values): \Rightarrow empty-type
```

abrupt(V) terminates abruptly for reason V.

Rule 
$$abrupt(V : values) \xrightarrow{abrupted(V)} stuck$$

Funcon handle-abrupt(
$$\_: T' \Rightarrow T, \_: T'' \Rightarrow T$$
):  $T' \Rightarrow T$ 

<sup>\*</sup>Suggestions for improvement: plancomps@gmail.com.
Reports of issues: https://github.com/plancomps/CBS-beta/issues.

handle-abrupt(X, Y) first evaluates X. If X terminates normally with value V, then V is returned and Y is ignored. If X terminates abruptly for reason V, then Y is executed with V as given value.

handle-abrupt(X, Y) is associative, with abrupt(given) as left and right unit. handle-abrupt(X, else(Y, abrupt(given))) ensures propagation of abrupt termination for the given reason if Y fails

Rule 
$$X \xrightarrow{\operatorname{abrupted(\ )}} X'$$
 $X \xrightarrow{\operatorname{handle-abrupt(\ X,\ Y)}} X'$ 

handle-abrupt(\(X,\ Y) \)

 $X \xrightarrow{\operatorname{abrupted(\ V:\ T'')}} X'$ 

Funcon finally( $\_: \Rightarrow T$ ,  $\_: \Rightarrow null-type$ ):  $\Rightarrow T$ 

finally(X, Y) first executes X. If X terminates normally with value V, then Y is executed before terminating normally with value V. If X terminates abruptly for reason V, then Y is executed before terminating abruptly with the same reason V.

Rule 
$$\frac{X \xrightarrow{\operatorname{abrupted}()} X'}{\operatorname{finally}(X,Y) \xrightarrow{\operatorname{abrupted}()} \operatorname{finally}(X',Y)}$$

$$X \xrightarrow{\operatorname{abrupted}(V:\operatorname{values})} X'$$

$$\operatorname{finally}(X,Y) \xrightarrow{\operatorname{abrupted}()} \operatorname{sequential}(Y,\operatorname{abrupt}(V))$$
Rule 
$$\operatorname{finally}(V:T,Y) \rightsquigarrow \operatorname{sequential}(Y,V)$$