

Funcons-beta: Patterns

The P_{LAN}CompS Project

Funcons-beta/Values/Abstraction/Patterns/Patterns.cbs*

Patterns

```
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```

General patterns are simple patterns or structured patterns. Matching a pattern to a value either computes an environment or fails.

Simple patterns are constructed from abstractions whose bodies depend on a given value, and whose executions either compute environments or fail.

Structured patterns are composite values whose components may include simple patterns as well as other values.

Matching a structured value to a structured pattern is similar to assigning a structured value to a structured variable, with simple pattern components matching component values analogously to simple variable components assigned component values.

Note that patterns match only values, not (empty or proper) sequences.

*Suggestions for improvement: plancomps@gmail.com.
Issues: <https://github.com/plancomps/CBS-beta/issues>.

Meta-variables $T, T' <:$ **values**

Simple patterns

Datatype **patterns** ::= **pattern**($_ :$ **abstractions**(**values** \Rightarrow **environments**))

patterns is the type of simple patterns that can match values of a particular type.

pattern(**abstraction**(X)) constructs a pattern with dynamic bindings, and **pattern**(**closure**(X)) computes a pattern with static bindings. However, there is no difference between dynamic and static bindings when the pattern is matched in the same scope where it is constructed.

Funcon **pattern-any** : \Rightarrow **patterns**
 \rightsquigarrow **pattern**(**abstraction**(**map**($_$)))

pattern-any matches any value, computing the empty environment.

Funcon **pattern-bind**($I :$ **identifiers**) : \Rightarrow **patterns**
 \rightsquigarrow **pattern**(**abstraction**(**bind-value**(I ,
 given)))

pattern-bind(I) matches any value, computing the environment binding I to that value.

Funcon **pattern-type**(T) : \Rightarrow **patterns**
 \rightsquigarrow **pattern**(**abstraction**(**if-true-else**(**is-in-type**(given ,
 T),
map($_$),
fail)))

pattern-type(T) matches any value of type T , computing the empty environment.

Funcon **pattern-else**($_ :$ **values**, $_ :$ **values**) : \Rightarrow **patterns**
Rule **pattern-else**($P_1 :$ **values**, $P_2 :$ **values**) \rightsquigarrow **pattern**(**abstraction**(**else**(**match**(given , P_1), **match**(given , P_2)))

`pattern-else`(P_1, P_2) matches all values matched by P_1 or by P_2 . If a value matches P_1 , that match gives the computed environment; if a value does not match P_1 but matches P_2 , that match gives the computed environment; otherwise the match fails.

Funcon `pattern-unite`($_ : \text{values}, _ : \text{values}$) : \Rightarrow `patterns`

Rule `pattern-unite`($P_1 : \text{values}, P_2 : \text{values}$) \rightsquigarrow `pattern`(`abstraction`(`collateral`(`match`(`given`, P_1), `match`(`given`, P_2))

`pattern-unite`(P_1, P_2) matches all values matched by both P_1 and P_2 , then uniting the computed environments, which fails if the domains of the environments overlap.

Pattern matching

Funcon `match`($_ : \text{values}, _ : \text{values}$) : \Rightarrow `environments`

`match`(V, P) takes a (potentially structured) value V and a (potentially structured) pattern P . Provided that the structure and all components of P exactly match the structure and corresponding components of V , the environments computed by the simple pattern matches are united.

Rule `match`($V : \text{values}, \text{pattern}(\text{abstraction}(X))$) \rightsquigarrow `give`(V, X)

Rule
$$\frac{}{\text{match}(\text{datatype-value}(I_1 : \text{identifiers}, V_1^* : \text{values}^*), \text{datatype-value}(I_2 : \text{identifiers}, V_2^* : \text{values}^*)) \rightsquigarrow \text{some-element}(\text{dom}(M_1), \text{dom}(M_2))}$$

Rule
$$\frac{\text{dom}(M_2) == \{ \}}{\text{match}(M_1 : \text{maps}(-, -), M_2 : \text{maps}(-, -)) \rightsquigarrow \text{if-true-else}(\text{is-equal}(\text{dom}(M_1), \{ \}), \text{map}(-, -), \text{fail})}$$

Rule
$$\frac{\text{dom}(M_2) \neq \text{some-element}(\text{dom}(M_1), \text{dom}(M_2))}{\text{match}(M_1 : \text{maps}(-, -), M_2 : \text{maps}(-, -)) \rightsquigarrow \text{if-true-else}(\text{is-in-set}(K, \text{dom}(M_1)), \text{collateral}(\text{match}(\text{map-loop}(M_1, M_2, K), \text{map-loop}(M_1, M_2, K)), \text{fail}))}$$

Rule
$$\frac{P : \sim(\text{datatype-values} \mid \text{maps}(-, -))}{\text{match}(V : \text{values}, P : \text{values}) \rightsquigarrow \text{if-true-else}(\text{is-equal}(V, P), \text{map}(-, -), \text{fail})}$$

Funcon `match-loosely`($_ : \text{values}, _ : \text{values}$) : \Rightarrow `environments`

`match-loosely`(V, P) takes a (potentially structured) value V and a (potentially structured) pattern P . Provided that the structure and all components of P loosely match the structure and corresponding components of V , the environments computed by the simple pattern matches are united.

Rule $\text{match-loosely}(V : \text{values}, \text{pattern}(\text{abstraction}(X))) \rightsquigarrow \text{give}(V, X)$

Rule $\frac{}{\text{match-loosely}(\text{datatype-value}(I_1 : \text{identifiers}, V_1^* : \text{values}^*), \text{datatype-value}(I_2 : \text{identifiers}, V_2^* : \text{values}^*)) \rightsquigarrow \text{match-loosely}(I_1, I_2, V_1^*, V_2^*)}$

Rule $\frac{\text{dom}(M_2) == \{ \}}{\text{match-loosely}(M_1 : \text{maps}(-, -), M_2 : \text{maps}(-, -)) \rightsquigarrow \text{map}(\text{ }, \text{ })}$

Rule $\frac{}{\text{match-loosely}(M_1 : \text{maps}(-, -), M_2 : \text{maps}(-, -)) \rightsquigarrow \text{if-true-else}(\text{is-in-set}(K, \text{dom}(M_1)), \text{collateral}(\text{match-loosely}(M_1, M_2, K, \text{dom}(M_1))), \text{fail})}$

Rule $\frac{P : \sim(\text{datatype-values} \mid \text{maps}(-, -))}{\text{match-loosely}(DV : \text{values}, P : \text{values}) \rightsquigarrow \text{if-true-else}(\text{is-equal}(DV, P), \text{map}(\text{ }, \text{ }), \text{fail})}$

Funcon $\text{case-match}(_ : \text{values}, _ : \Rightarrow T') : \Rightarrow T'$

$\text{case-match}(P, X)$ matches P exactly to the given value. If the match succeeds, the computed bindings have scope X .

Rule $\text{case-match}(P : \text{values}, X) \rightsquigarrow \text{scope}(\text{match}(\text{given}, P), X)$

Funcon $\text{case-match-loosely}(_ : \text{values}, _ : \Rightarrow T') : \Rightarrow T'$

$\text{case-match}(P, X)$ matches P loosely to the given value. If the match succeeds, the computed bindings have scope X .

Rule $\text{case-match-loosely}(P : \text{values}, X) \rightsquigarrow \text{scope}(\text{match-loosely}(\text{given}, P), X)$

Funcon $\text{case-variant-value}(_ : \text{identifiers}) : \Rightarrow \text{values}$

$\text{case-variant-value}(I)$ matches values of variant I , then giving the value contained in the variant.

Rule $\text{case-variant-value}(I : \text{identifiers}) \rightsquigarrow \text{case-match}(\text{variant}(I, \text{pattern-any}), \text{variant-value}(\text{given}))$