

Unstable-Languages-beta: SIMPLE-THR-3-Statements *

The P_{LAN}CompS Project

SIMPLE-THR-3-Statements.cbs | PLAIN | PRETTY

Language "SIMPLE-THR"

3 Statements

Syntax $Block : \text{block} ::= \{ \text{stmts?} \}$
 $Stmts : \text{stmts} ::= \text{stmt stmts?}$
 $Stmt : \text{stmt} ::= \text{imp-stmt} \mid \text{vars-decl}$
 $ImpStmt : \text{imp-stmt} ::= \text{block}$
| $\text{exp} \text{ ;}$
| $\text{if} \text{ ' (exp)' block ('else' block)?}$
| $\text{while} \text{ ' (exp)' block}$
| $\text{for} \text{ ' (stmt exp ; exp)' block}$
| $\text{print} \text{ ' (exps)' ;}$
| $\text{return} \text{ exp? ;}$
| $\text{try} \text{ block 'catch' ' (id)' block}$
| $\text{throw} \text{ exp ;}$
| $\text{join} \text{ exp ;}$
| $\text{acquire} \text{ exp ;}$
| $\text{release} \text{ exp ;}$
| $\text{rendezvous} \text{ exp ;}$

Rule $\llbracket \text{'if' ' (Exp)' Block} \rrbracket : \text{stmt} =$
 $\llbracket \text{'if' ' (Exp)' Block 'else' '{ '}' } \rrbracket$

Rule $\llbracket \text{'for' ' (Stmt Exp_1 ; Exp_2)'}$
 $\text{'{ ' Stmts '}' } \rrbracket : \text{stmt} =$
 $\llbracket \text{'{ ' Stmt}$
 $\text{'while' ' (Exp_1)'}$
 $\text{'{ ' { ' Stmts '}' Exp_2 ; ' '}'}$
 $\text{'}' } \rrbracket$

*Suggestions for improvement: plancomps@gmail.com.
Reports of issues: <https://github.com/plancomps/CBS-beta/issues>.

Semantics $\text{exec}[_ : \text{stmts}] : \Rightarrow \text{null-type}$

Rule $\text{exec}[\text{'{' '}'}] = \text{null}$

Rule $\text{exec}[\text{'{' Stmt '}'}] = \text{exec}[\text{Stmt}]$

Rule $\text{exec}[\text{ImpStmt Stmt}] =$
 $\text{sequential}(\text{exec}[\text{ImpStmt}], \text{exec}[\text{Stmt}])$

Rule $\text{exec}[\text{VarsDecl Stmt}] =$
 $\text{scope}(\text{declare}[\text{VarsDecl}], \text{exec}[\text{Stmt}])$

Rule $\text{exec}[\text{VarsDecl}] = \text{effect}(\text{declare}[\text{VarsDecl}])$

Rule $\text{exec}[\text{Exp ';' }] = \text{effect}(\text{rval}[\text{Exp}])$

Rule $\text{exec}[\text{'if' '(' Exp ')' Block₁ 'else' Block₂}] =$
 $\text{if-else}(\text{rval}[\text{Exp}], \text{exec}[\text{Block₁}], \text{exec}[\text{Block₂}])$

Rule $\text{exec}[\text{'while' '(' Exp ')' Block}] = \text{while}(\text{rval}[\text{Exp}], \text{exec}[\text{Block}])$

Rule $\text{exec}[\text{'print' '(' Exps ')' ';' }] = \text{print}(\text{rvals}[\text{Exps}])$

Rule $\text{exec}[\text{'return' Exp ';' }] = \text{return}(\text{rval}[\text{Exp}])$

Rule $\text{exec}[\text{'return' ';' }] = \text{return}(\text{null})$

Rule $\text{exec}[\text{'try' Block₁ 'catch' '(' Id ')' Block₂}] =$
 $\text{handle-thrown}(\text{exec}[\text{Block₁}],$
 $\text{scope}(\text{bind}(\text{id}[\text{Id}], \text{allocate-initialised-variable}(\text{values}, \text{given})),$
 $\text{exec}[\text{Block₂}]))$

Rule $\text{exec}[\text{'throw' Exp ';' }] = \text{throw}(\text{rval}[\text{Exp}])$

SIMPLE uses natural numbers to identify threads; the use of $\text{lookup-index}(_)$ below converts a natural number to the associated thread-id.

Rule $\text{exec}[\text{'join' Exp ';' }] =$
 $\text{thread-join } \text{lookup-index}(\text{rval}[\text{Exp}])$

The use of $\text{memo-value}(V, SY)$ below associates V with a lock. When a thread requests a lock already held by another thread, the requesting thread is suspended until the request is granted. The use of $\text{postpone}(_)$ below automatically releases held locks when the current thread terminates.

Rule $\text{exec}[\text{'acquire' Exp ';' }] =$
 $\text{give}(\text{memo-value}(\text{rval}[\text{Exp}], \text{reentrant-lock-create}),$
 $\text{sequential}(\text{postpone}$
 $\text{if-true-else}(\text{is-exclusive-lock-holder given},$
 $\text{reentrant-lock-release given},$
 $\text{null-value}),$
 $\text{reentrant-lock-sync-else-wait given}))$

The use of $\text{memo-value-recall}(V)$ below gives the lock associated with V .

Rule $\text{exec}[\text{'release' Exp ';' }] =$
 $\text{reentrant-lock-exit } \text{memo-value-recall } \text{rval}[\text{Exp}]$

The use of $\text{memo-value}(V, SY)$ below associates V with a rendezvous. When a thread requests a rendezvous on a particular value, and there is no previous uncompleted request for a rendezvous on the same value, the requesting thread is suspended until the request is granted.

Rule `exec`[['rendezvous' *Exp* ';']] =
rendezvous-sync-else-wait(
memo-value("rendezvous", rendezvous-create(2)),
`rval`[[*Exp*]])