

Unstable-Languages-beta: LD-Start *

The PPlanCompS Project

LD-Start.cbs | PLAIN | PRETTY

OUTLINE

1 Lexical constructs

2 Call-by-value lambda-calculus

3 Arithmetic and Boolean expressions

4 References and imperatives

5 Multithreading

6 Programs

Language "LD"

[1 Lexical constructs
2 Call-by-value lambda-calculus
3 Arithmetic and Boolean expressions
4 References and imperatives
5 Multithreading
6 Programs
A Disambiguation]

Lexical syntax:

```
Lexis  X : id ::= ('a'-'z') ('a'-'z' | '0'-'9')*
        N : int ::= ('0'-'9')+
        keyword ::= 'do' | 'else' | 'fork' | 'if'
                  | 'in' | 'join' | 'lambda' | 'let'
                  | 'ref' | 'spawn' | 'then' | 'while'
```

Context-free syntax:

*Suggestions for improvement: plancomps@gmail.com.
Reports of issues: <https://github.com/plancomps/CBS-beta/issues>.

Syntax $E : \text{exp} ::=$

- int
- id
- $\text{'lambda' id '.' exp}$
- exp exp
- $\text{'let' id '=' exp 'in' exp}$
- '(' exp ')
- exp '+' exp
- exp '*' exp
- exp '/' exp
- exp '<=' exp
- exp '&\&' exp
- $\text{'if' exp 'then' exp 'else' exp}$
- 'ref' exp
- exp ':=' exp
- '!' exp
- exp ';' exp
- '(' ')
- $\text{'while' exp 'do' exp}$
- 'spawn' exp
- 'join' exp

Expression evaluation:

Type Id-values

- $\rightsquigarrow \text{functions}(\text{values}, \text{values})$
- integers
- booleans
- variables
- null-type
- thread-ids

Semantics $\text{eval}[\![_ : \text{exp}]\!] : \Rightarrow \text{Id-values}$

1 Lexical constructs

Rule $\text{eval}[\![N]\!] = \text{decimal } "N"$

Rule $\text{eval}[\![X]\!] = \text{bound } "X"$

2 Call-by-value lambda-calculus

Rule $\text{eval}[\![\text{'lambda' } X \text{'.' } E]\!] =$
function closure
scope(
bind("X", given),
 $\text{eval}[\![E]\!]$)

Rule $\text{eval}[\![E_1 E_2]\!] =$
apply($\text{eval}[\![E_1]\!]$, $\text{eval}[\![E_2]\!]$)

Rule $\text{eval}[\![\text{'let' } X \text{'=' } E_1 \text{'in' } E_2]\!] =$
scope(
bind("X", $\text{eval}[\![E_1]\!]$),
 $\text{eval}[\![E_2]\!]$)

Desugaring (alternative to the above rule):

Rule $\llbracket \text{'let' } X \text{'=' } E_1 \text{'in' } E_2 \rrbracket : \text{exp} =$
 $\llbracket \text{'(' } \text{'lambda' } X \text{'.' } E_2 \text{')' ' (' } E_1 \text{')' } \rrbracket$

Rule $\text{eval}[\![\text{'(' } E \text{')' }]\!] = \text{eval}[\![E]\!]$

3 Arithmetic and Boolean expressions

Rule $\text{eval}[\![E_1 \text{'+' } E_2]\!] =$
int-add($\text{eval}[\![E_1]\!]$, $\text{eval}[\![E_2]\!]$)

Rule $\text{eval}[\![E_1 \text{'*'} E_2]\!] =$
int-mul($\text{eval}[\![E_1]\!]$, $\text{eval}[\![E_2]\!]$)

Rule $\text{eval}[\![E_1 \text{'/' } E_2]\!] =$
checked int-div($\text{eval}[\![E_1]\!]$, $\text{eval}[\![E_2]\!]$)

Rule $\text{eval}[\![E_1 \text{'<=' } E_2]\!] =$
is-less-or-equal l-to-r($\text{eval}[\![E_1]\!]$, $\text{eval}[\![E_2]\!]$)

Rule $\text{eval}[\![E_1 \text{'\&\&'} E_2]\!] =$
if-true-else($\text{eval}[\![E_1]\!]$, $\text{eval}[\![E_2]\!]$, false)

Rule $\text{eval}[\![\text{'if' } E_1 \text{'then' } E_2 \text{'else' } E_3]\!] =$
if-true-else($\text{eval}[\![E_1]\!]$, $\text{eval}[\![E_2]\!]$, $\text{eval}[\![E_3]\!]$)

4 References and imperatives

Rule $\text{eval}[\![\text{'ref' } E]\!] =$
allocate-initialised-variable(lid-values, $\text{eval}[\![E]\!]$)

Rule $\text{eval}[\![E_1 \text{' := ' } E_2]\!] =$
assign($\text{eval}[\![E_1]\!]$, $\text{eval}[\![E_2]\!]$)

Rule $\text{eval}[\![\text{'!' } E]\!] = \text{assigned}(\text{eval}[\![E]\!])$

Rule $\text{eval}[\![E_1 \text{' ; ' } E_2]\!] =$
sequential(effect($\text{eval}[\![E_1]\!]$), $\text{eval}[\![E_2]\!]$)

Rule $\text{eval}[\![\text{'(' } \text{')' }]\!] = \text{null-value}$

Rule $\text{eval}[\![\text{'while' } E_1 \text{'do' } E_2]\!] =$
while-true($\text{eval}[\![E_1]\!]$, effect($\text{eval}[\![E_2]\!]$))

5 Multithreading

N.B. The funcons for multithreading have not yet been fully validated, so they are defined in Unstable-Funcons-beta instead of Funcons-beta.

Rule $\text{eval}[\![\text{'spawn'}\ E]\!] =$
 $\text{thread-activate thread-joinable thunk closure eval}[\![E]\!]$
Rule $\text{eval}[\![\text{'join'}\ E]\!] = \text{thread-join}(\text{eval}[\![E]\!])$

6 Programs

Syntax $\text{START} : \text{start} ::= \text{exp}$

Semantics $\text{start}[\![_ : \text{start}]\!] : \Rightarrow \text{values}$

Rule $\text{start}[\![E]\!] =$
 $\text{initialise-binding}$
 $\text{initialise-storing}$
 finalise-failing
 multithread
 $\text{eval}[\![E]\!]$