Funcons-beta: Abrupting

The PLanCompS Project

Funcons-beta/Computations/Abnormal/Abrupting/Abrupting.cbs*

Abruptly terminating

```
[ Funcon stuck
    Entity abrupted
    Funcon finalise-abrupting
    Funcon abrupt
    Funcon handle-abrupt
    Funcon finally ]

Meta-variables T, T', T" <: values

Funcon stuck : ⇒ empty-type
```

stuck does not have any computation. It is used to represent the result of a transition that causes the computation to terminate abruptly.

```
Entity \_ \xrightarrow{\text{abrupted}(\_: values?)} \_
```

 $\operatorname{abrupted}(V)$ in a label on a transition indicates abrupt termination for reason V. $\operatorname{abrupted}()$ indicates the absence of abrupt termination.

```
Funcon finalise-abrupting (X:\Rightarrow T):\Rightarrow T\mid \mathsf{null-type} \leadsto \mathsf{handle-abrupt}(X, \mathsf{null-value})
```

^{*}Suggestions for improvement: plancomps@gmail.com. Issues: https://github.com/plancomps/CBS-beta/issues.

finalise-abrupting (X) handles abrupt termination of X for any reason.

Funcon abrupt(_: values) : ⇒ empty-type

abrupt(V) terminates abruptly for reason V.

Rule
$$abrupt(V : values) \xrightarrow{abrupted(V)} stuck$$

Funcon handle-abrupt(
$$_{-}: T' \Rightarrow T,_{-}: T'' \Rightarrow T$$
): $T' \Rightarrow T$

 $\mathsf{handle\text{-}abrupt}(X,Y)$ first evaluates X. If X terminates normally with value V, then V is returned and Y is ignored. If X terminates abruptly for reason V, then Y is executed with V as given value.

handle-abrupt(X, Y) is associative, with abrupt(given) as left and right unit. handle-abrupt(X, else(Y, abrupt(given))) ensures propagation of abrupt termination for the given reason if Y fails

$$Rule \xrightarrow{X \xrightarrow{abrupted(\)} X'} X'$$

$$handle-abrupt(X,Y) \xrightarrow{abrupted(\)} handle-abrupt(X',Y)$$

$$Rule \xrightarrow{X \xrightarrow{abrupted(V:T'')} X'} X'$$

$$handle-abrupt(X,Y) \xrightarrow{abrupted(\)} give(V,Y)$$

$$Rule \ handle-abrupt(V:T,Y) \leadsto V$$

Funcon finally($_: \Rightarrow T, _: \Rightarrow \text{null-type}$): $\Rightarrow T$

finally (X, Y) first executes X. If X terminates normally with value V, then Y is executed before terminating normally with value V. If X terminates abruptly for reason V, then Y is executed before terminating abruptly with the same reason V.

$$\begin{aligned} &Rule & \xrightarrow{X} & \xrightarrow{abrupted(\)} & X' \\ \hline &finally(X,Y) & \xrightarrow{abrupted(\)} & finally(X',Y) \\ &Rule & \xrightarrow{X} & \xrightarrow{abrupted(\ V: values)} & X' \\ \hline &finally(X,Y) & \xrightarrow{abrupted(\)} & sequential(Y,abrupt(V)) \\ &Rule & finally(V:T,Y) & \Rightarrow sequential(Y,V) \end{aligned}$$

