UMassAmherst

Garbage Collection Without Paging

Matthew Hertz, Yi Feng, Emery Berger

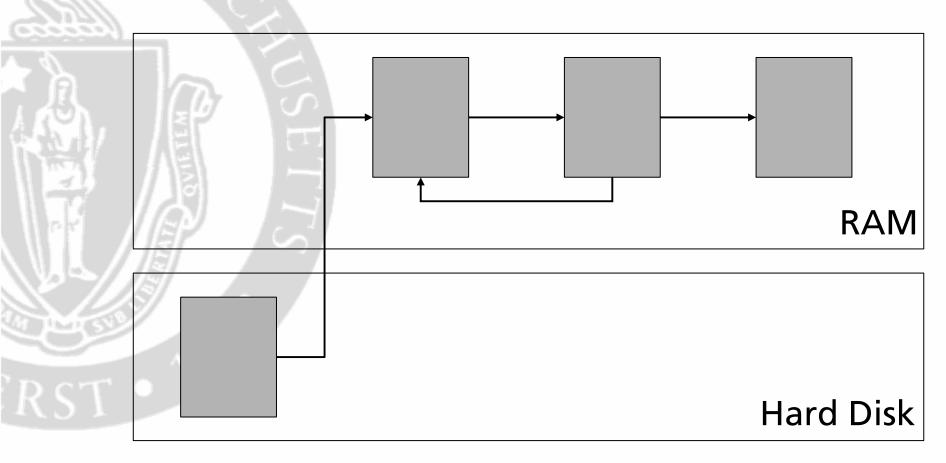
University of Massachusetts Amherst



Garbage Collection Performance

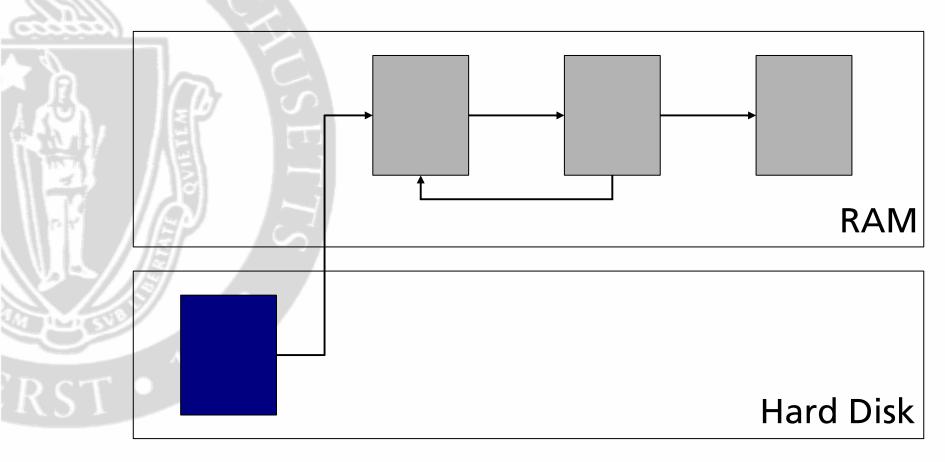
- Garbage collection now performs reasonably well
 - High throughput
 - Low pause times
 - Given large heap and sufficient memory
- But: what happens when there's not enough RAM?





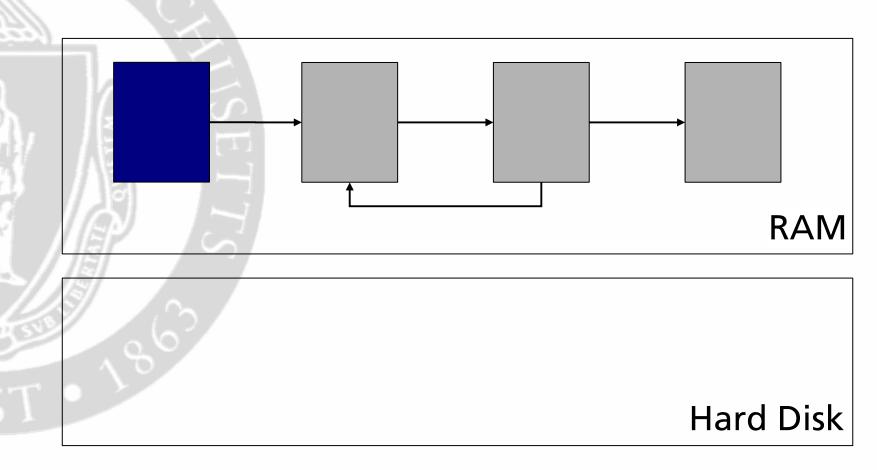
Heap: most pages in RAM, one on disk *GC begins*





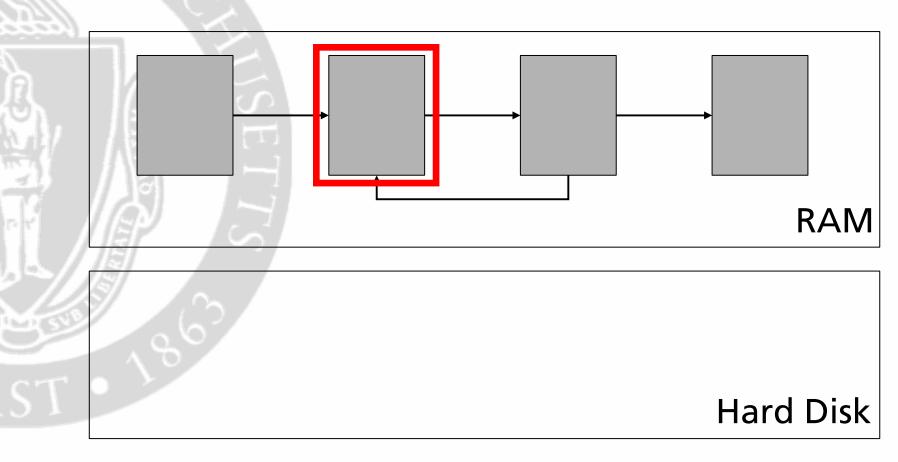
Collector touches an evicted page...





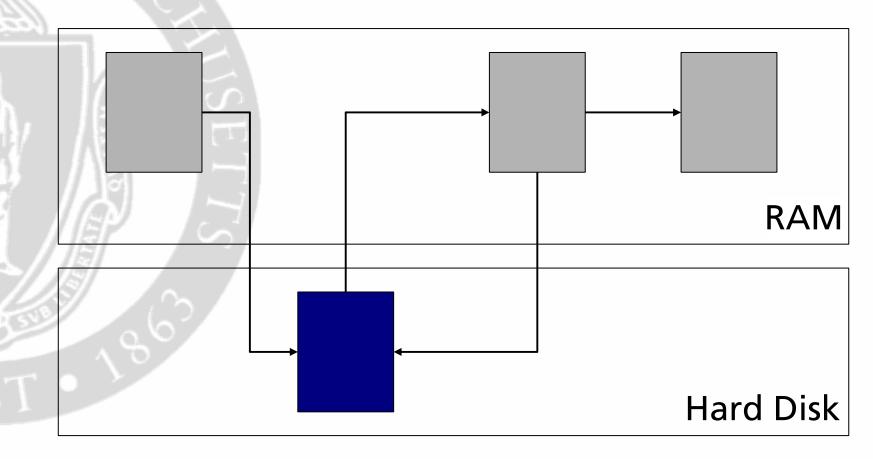
... bringing the page into memory ...





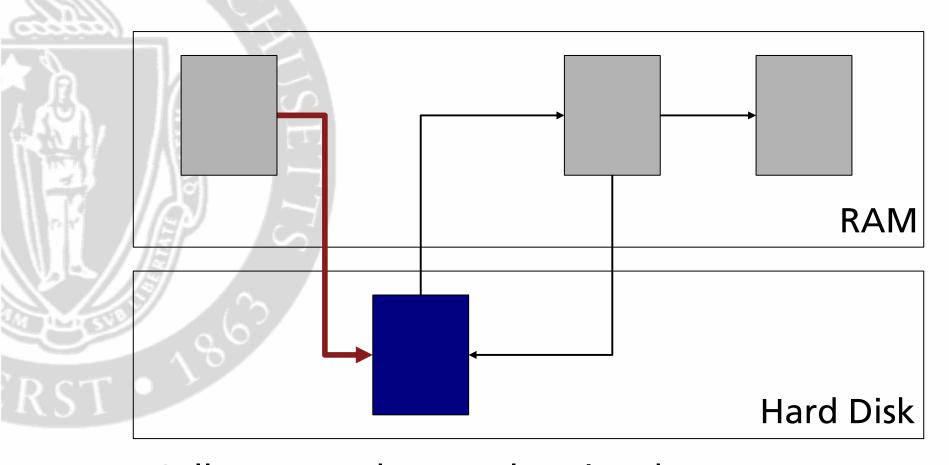
... but triggers another page eviction.





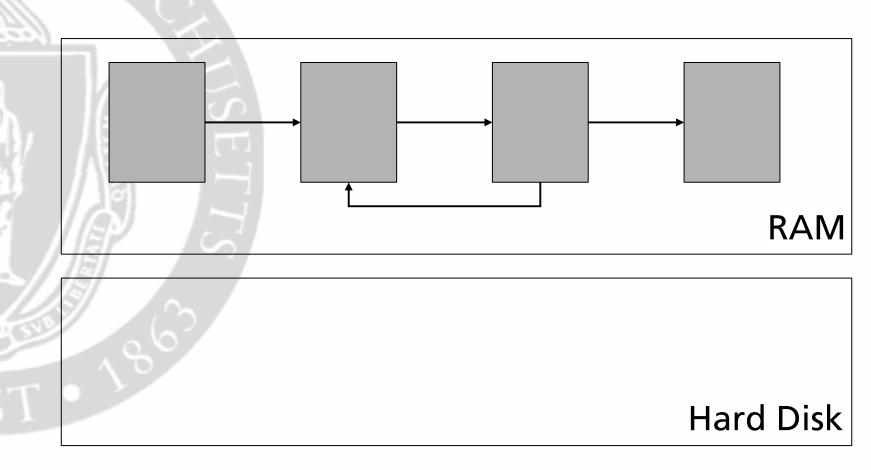
... but triggers another page eviction.





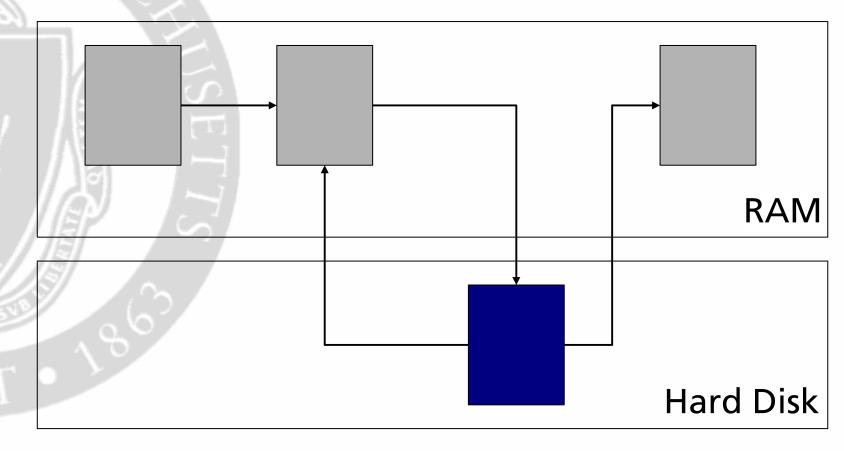
Collector touches newly-evicted page...





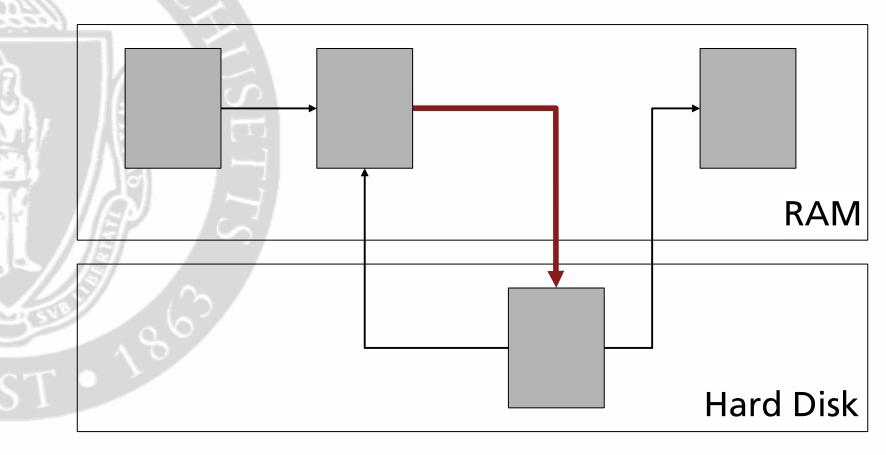
Collector touches newly-evicted page...





... leading to the eviction of yet another page...

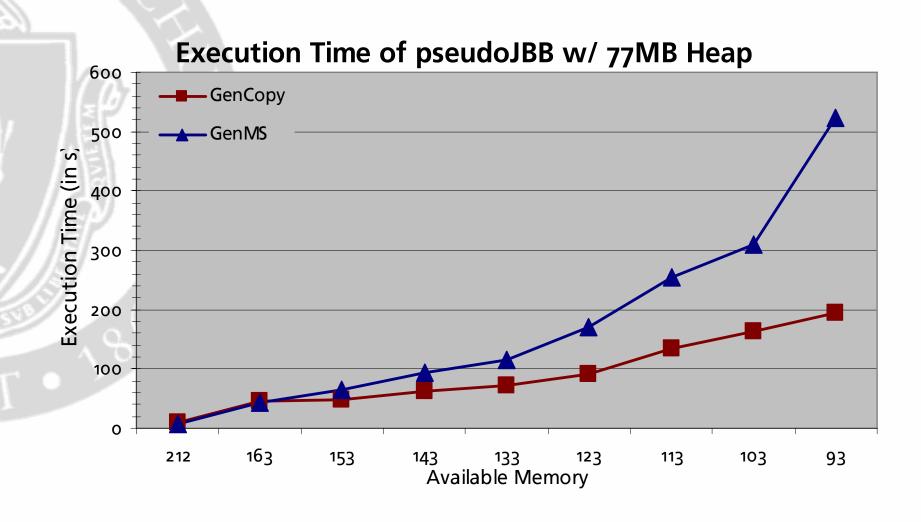


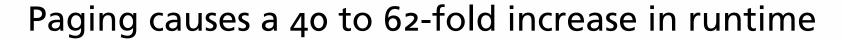


... which eventually triggers more paging.



Program Throughput







Outline

- Motivation
- GC without paging
 - Cooperative garbage collection
 - Bookmarking
- Results

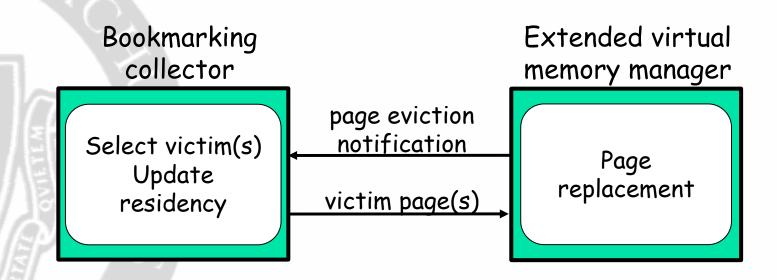


The problem

- Garbage collector: VM-oblivious
 - Examines all reachable objects
 - Lacks knowledge of page residency
 - Treats evicted and resident pages identically
- Virtual memory: GC-oblivious
 - GC & application have different access patterns
 - Likely to evict pages needed by GC

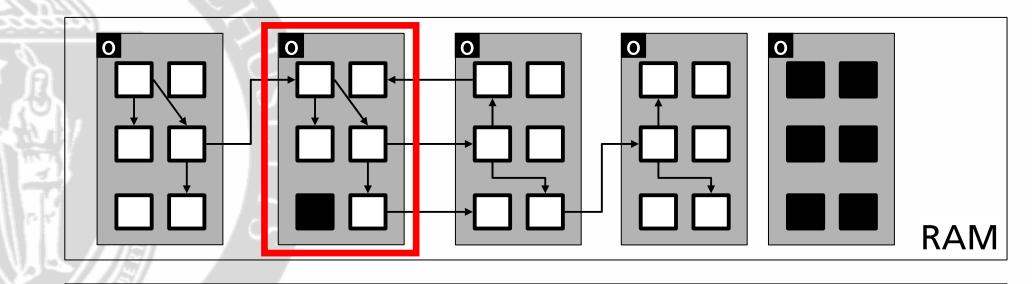


Cooperative Garbage Collection



- In response to notifications, BC:
 - Adjusts heap size to fit in main memory
 - Prevents eviction of important pages
 - Avoids touching non-resident pages

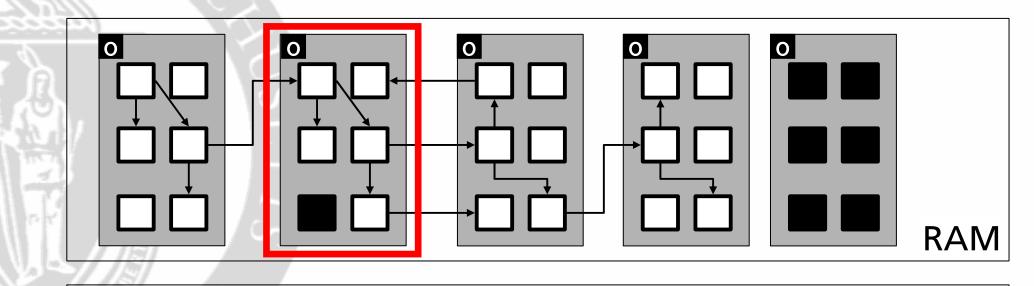




Hard Disk

When notified, avoid a pending eviction...

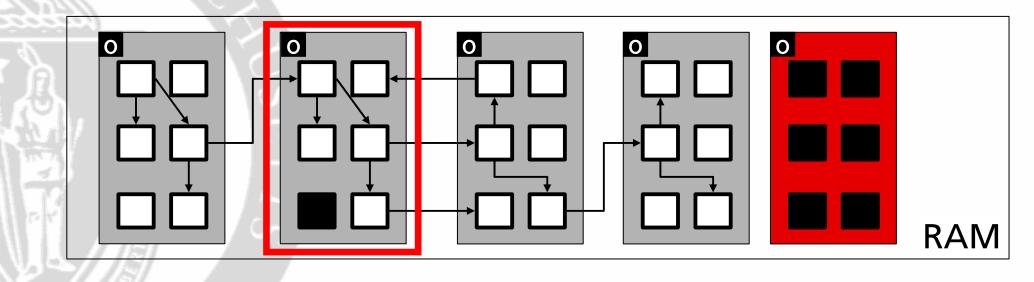




Hard Disk

...find a page BC knows to be empty...

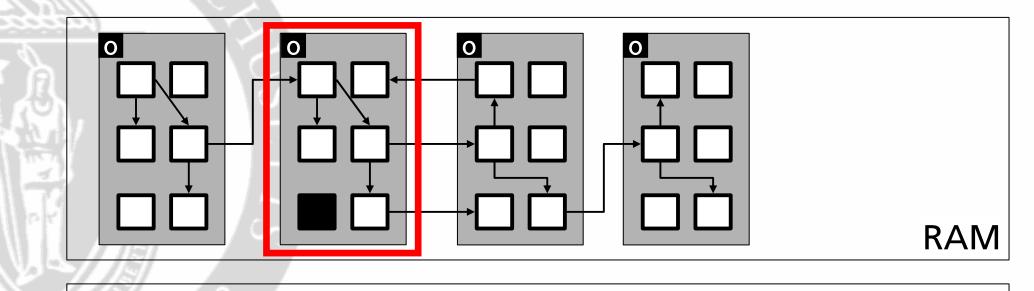




Hard Disk

... and discard it...





Hard Disk

... eliminating the need to evict a page.



Limit of Heap Sizing

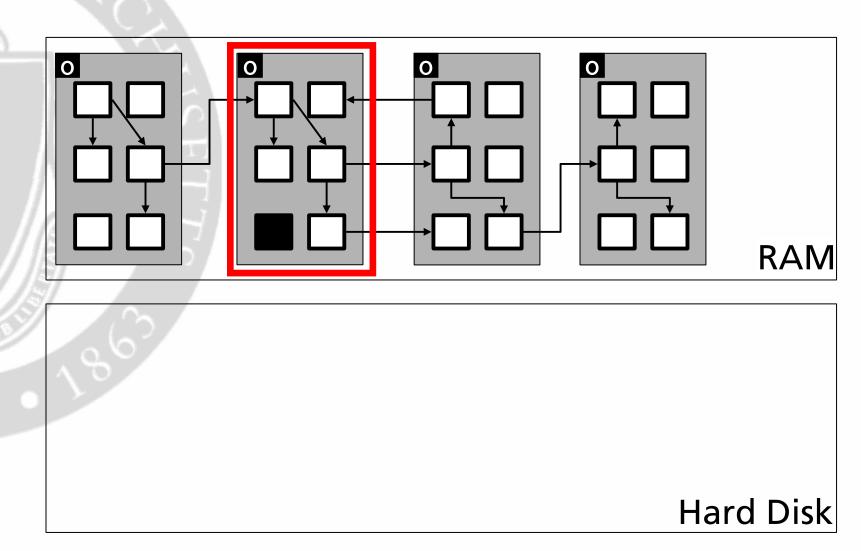
- Could collect, compact, compress, etc.
- Eventually:
 - Will run out of pages to discard
 - Going to have to evict non-empty pages
 - Result: Paging
- Can we avoid this?



Bookmarks

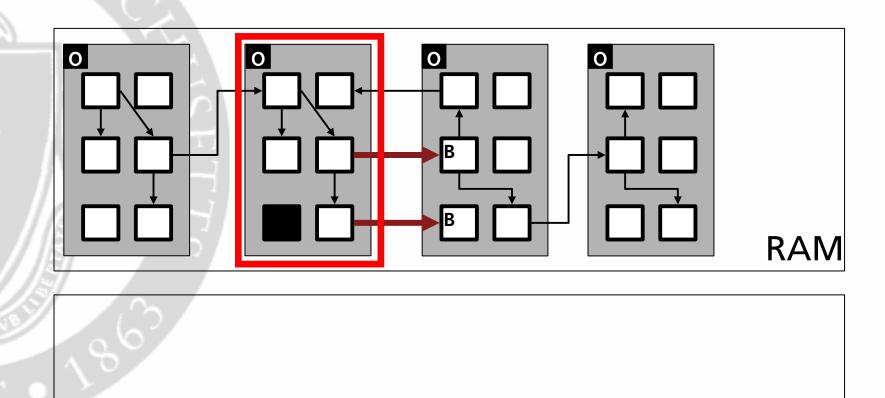
- We introduce bookmarks:
 - Summaries of connectivity info on evicted pages
 - References from objects on the page
 - These summaries enable GC w/o paging





Process page before eviction...

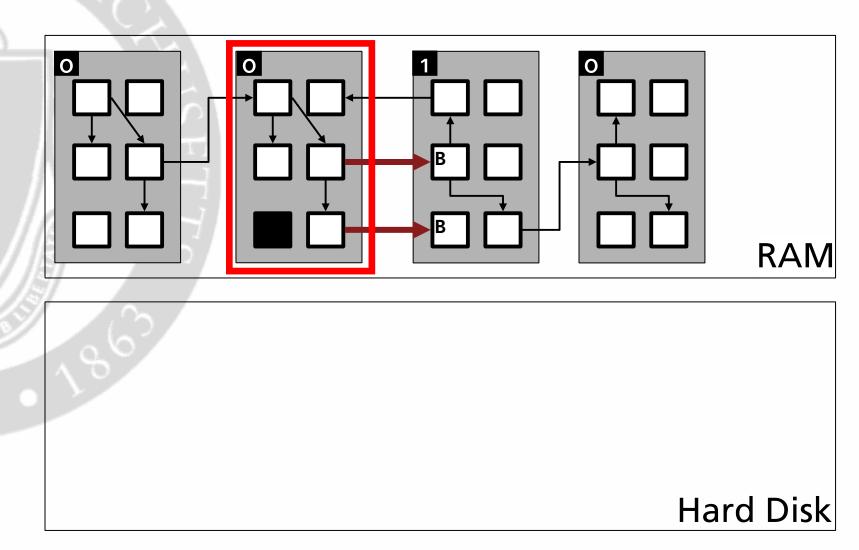




Hard Disk

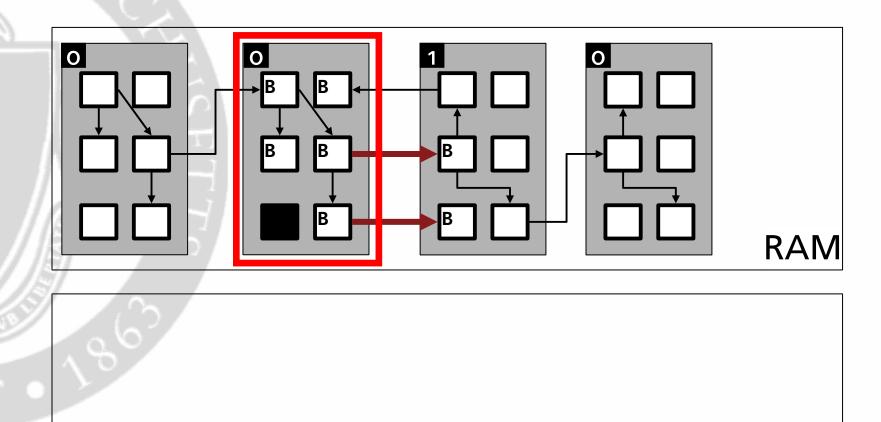
... by following pointers & bookmark-ing targets...





... increment the referring page count...

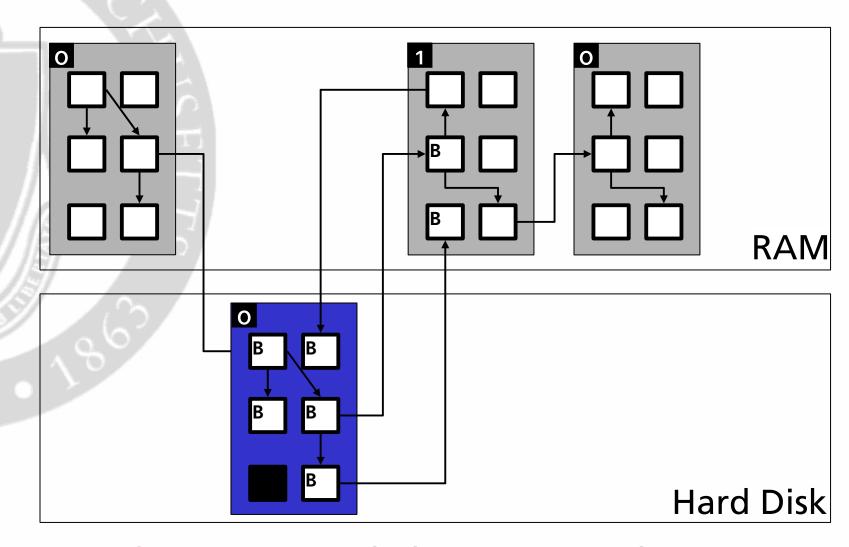




Hard Disk

... conservatively bookmark objects on the page...





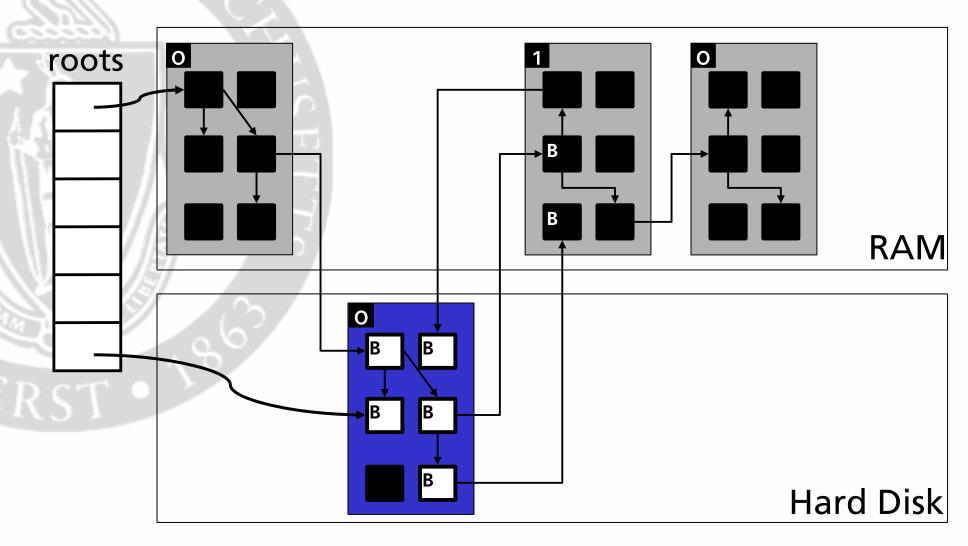
... then tell extended VM to evict the page.



Bookmarking Details

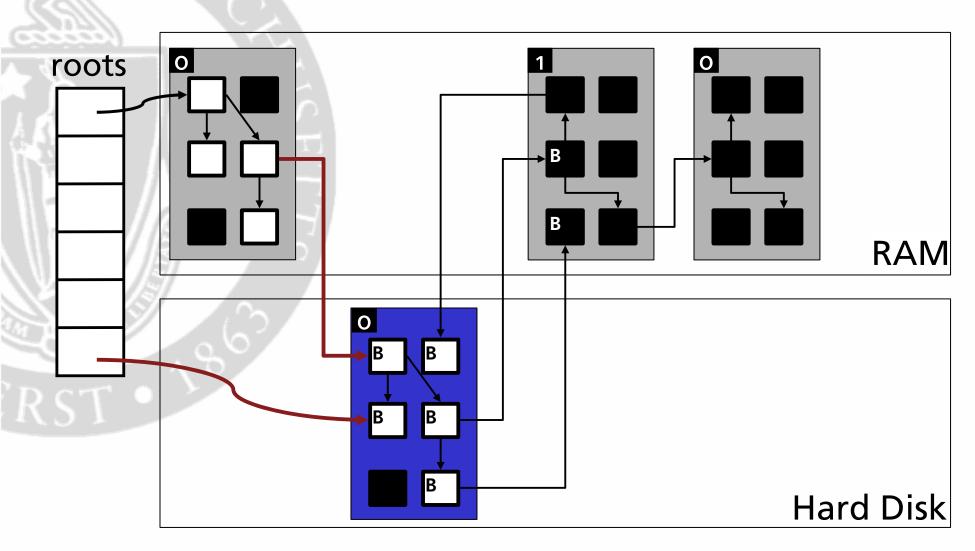
- Cheap summary of connectivity
 - One bit per object: free
 - One word per page: referring page count
 - Bookmarks cleared when count = zero
- Use bookmarks as secondary roots during garbage collection





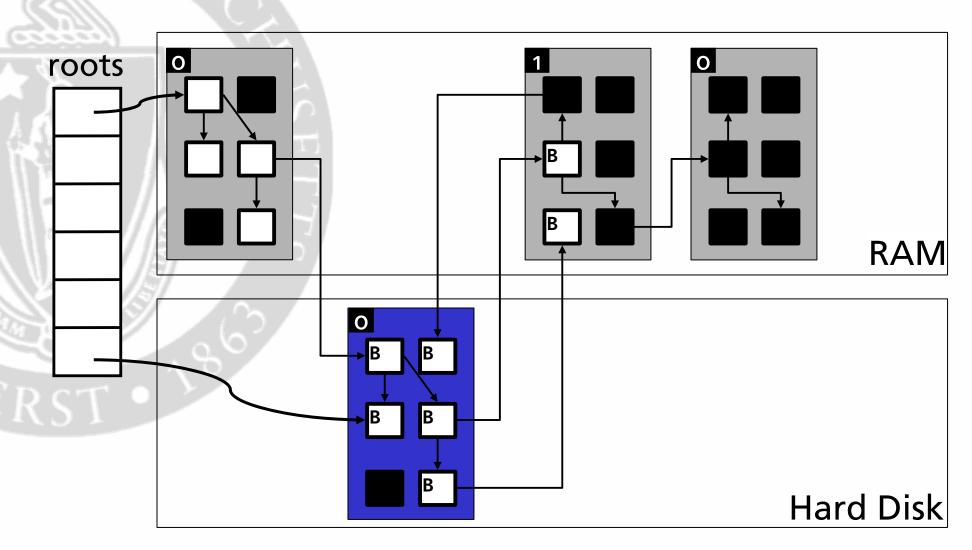
Process objects as usual, but...





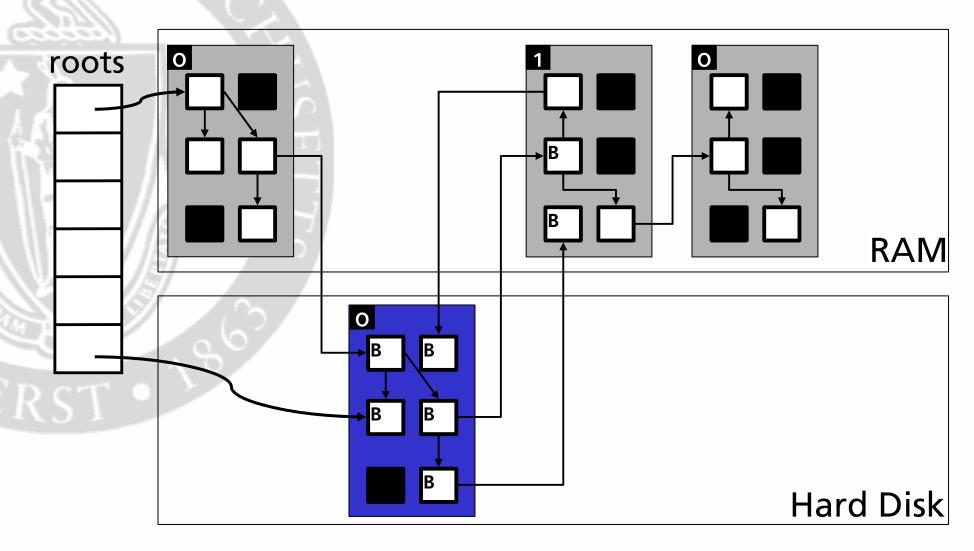
... ignore any references to evicted pages.





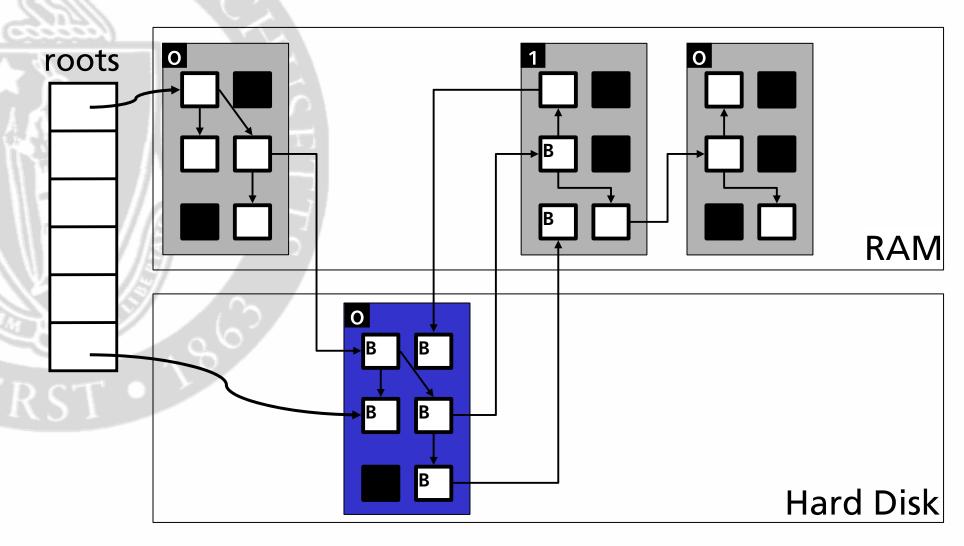
Use bookmarks to recreate evicted references...





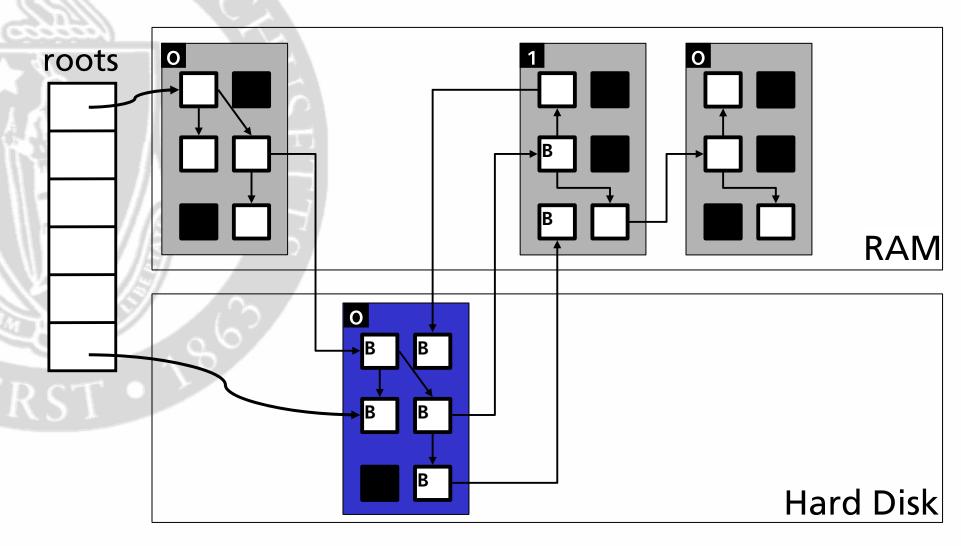
... and continue collection.





Result: Garbage collection without paging!





Note: can waste space on evicted pages.

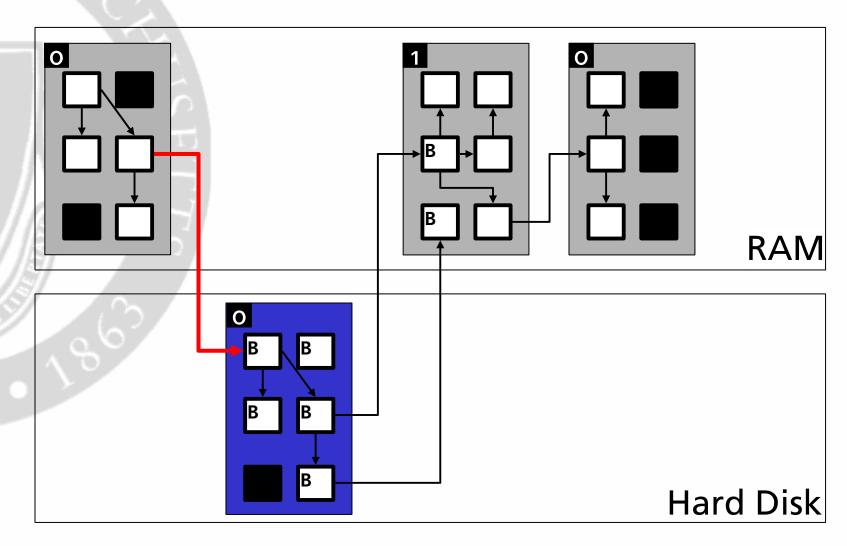


Bookmarking Incompleteness

- Space waste not just on evicted pages
- Collection with bookmarks is necessarily incomplete
 - Not guaranteed to reclaim all memory



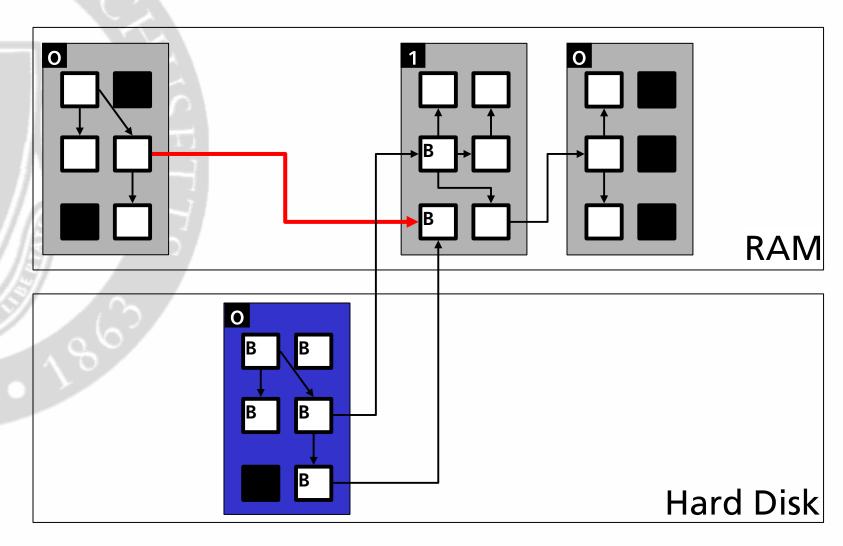
Bookmarking Incompleteness



When a reference to an evicted object changes...

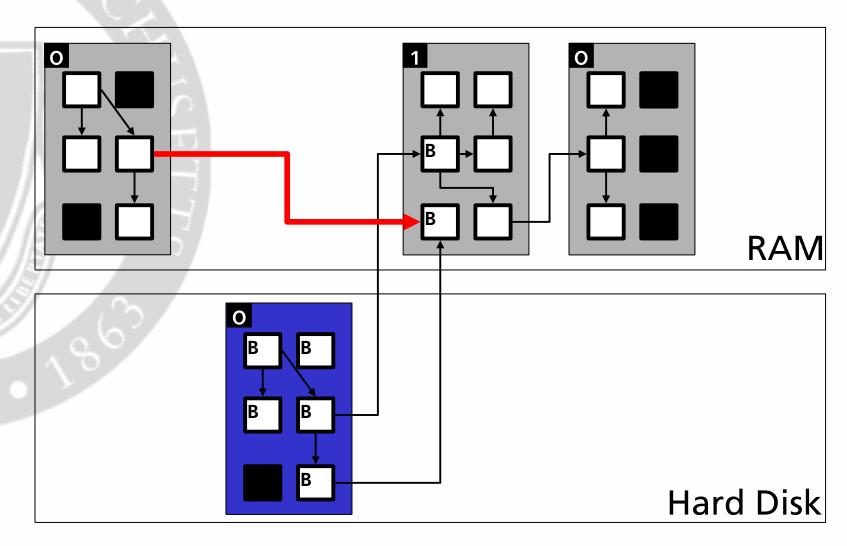


Bookmarking Incompleteness



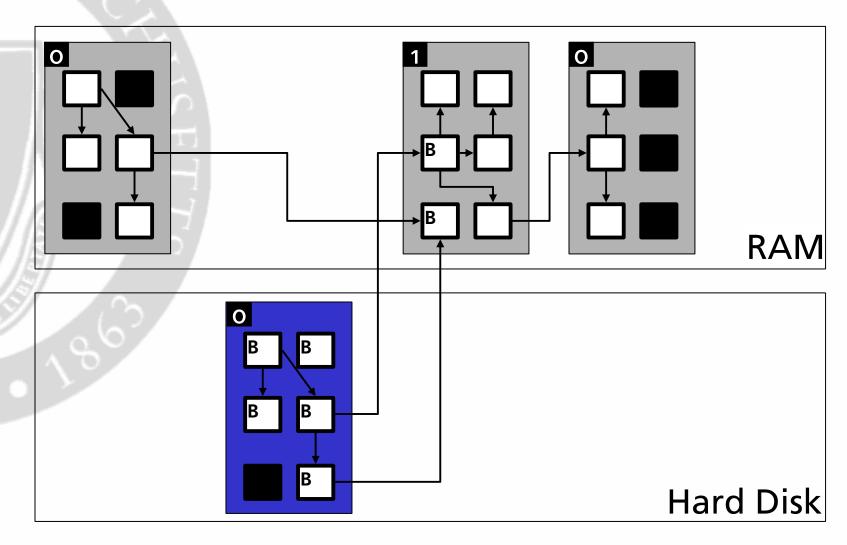
When a reference to an evicted object changes...





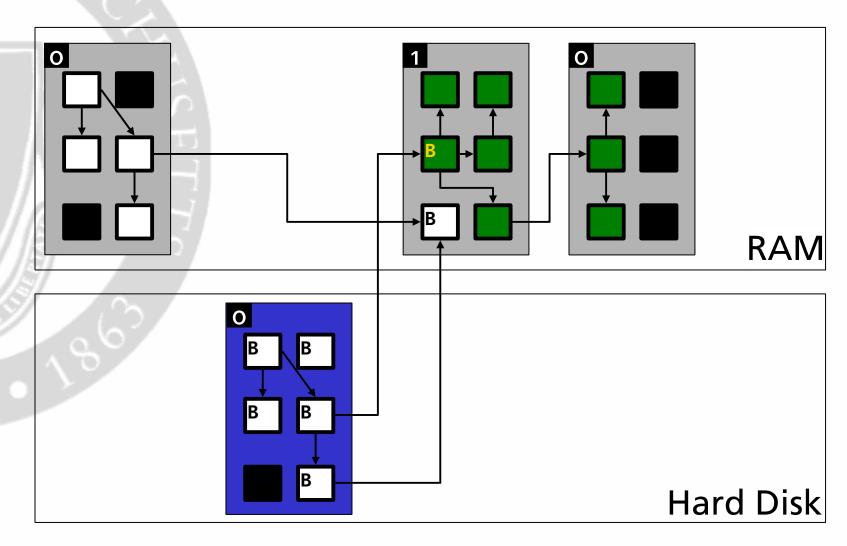
...it can make evicted objects unreachable.





But bookmarks cannot be removed...





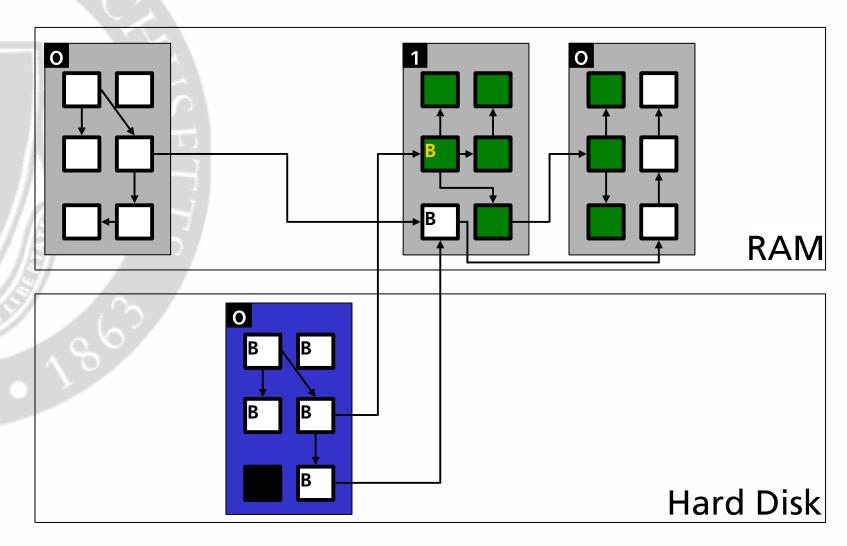
...retaining unreachable heap objects.



- Worst-case: completeness requires duplicating evicted pages
 - See paper for more info

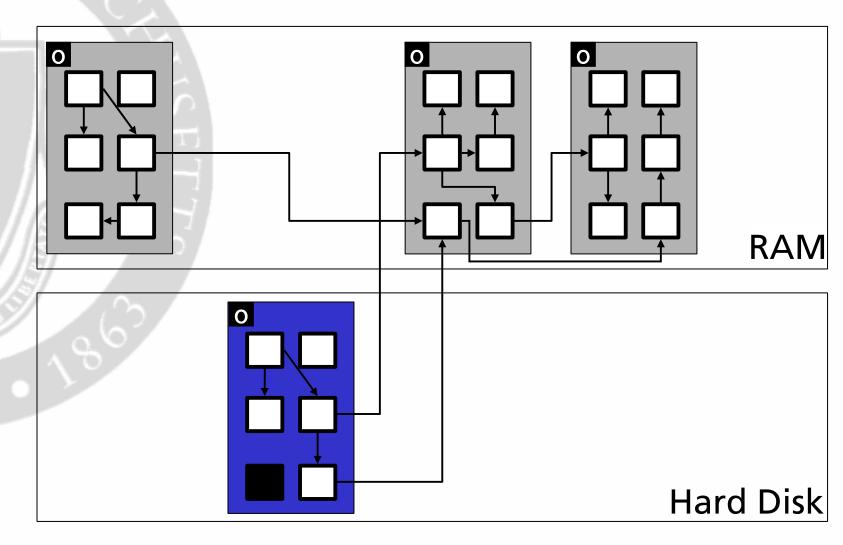
How can we preserve completeness?





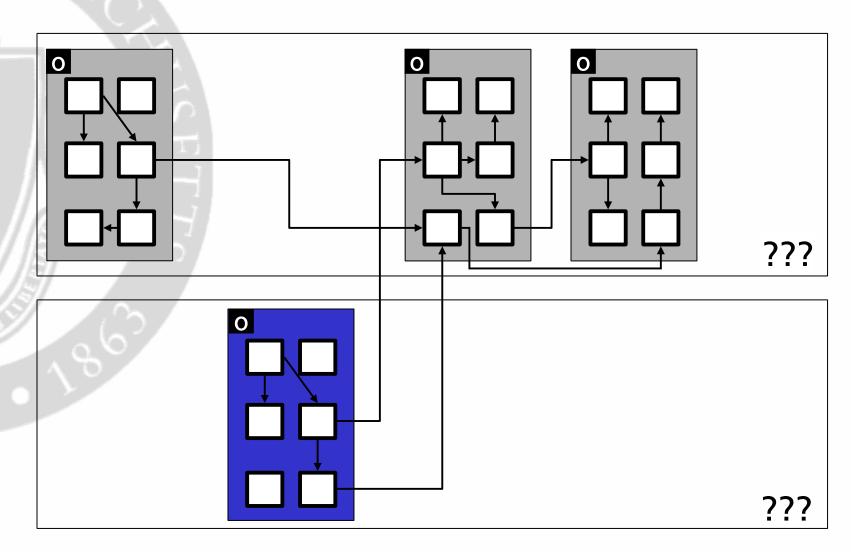
If the heap becomes full...





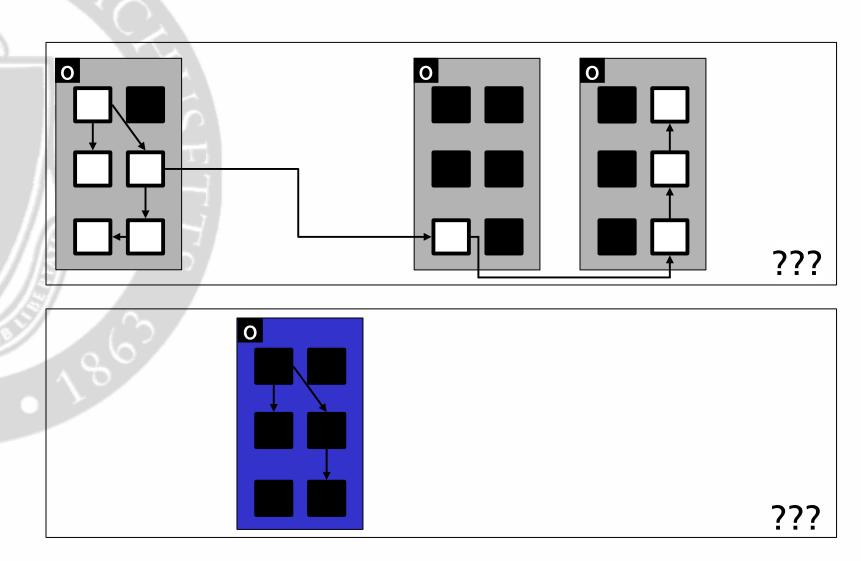
...BC removes all bookmarks...





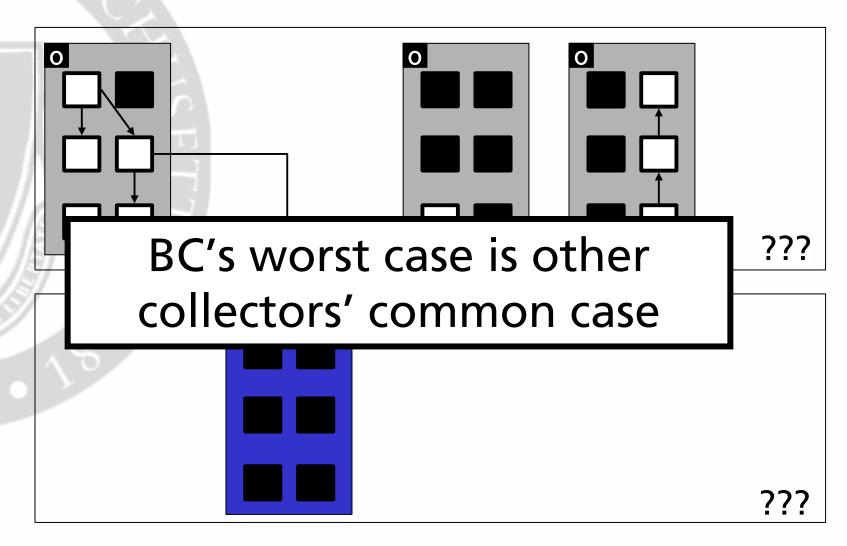
...and performs a VM-oblivious collection...





...reclaiming all unreachable objects.





...reclaiming all unreachable objects.



BC Performance Optimizations

- Uses generational design similar to GenMS
 - Yields good performance when not paging
- Compacts heap to reduce pressure
 - Prevents single object from tying down a page



Outline

- Motivation
- GC without paging
 - Cooperative garbage collection
 - Bookmarking
- Results

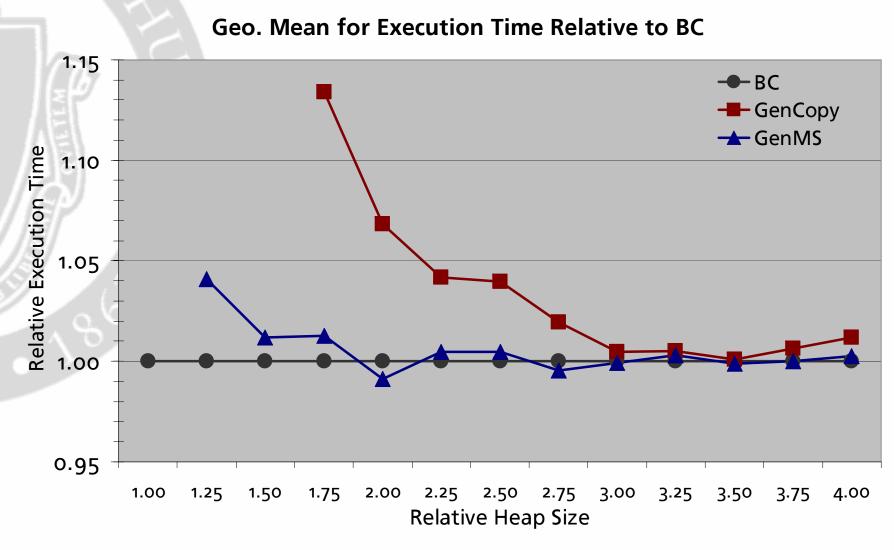


Experimental Methodology

- Extended Linux kernel 2.4.20
 - Eviction notification
 - vm_relinquish()
 - Added only 600 LOC
- Jikes RVM 2.3.2 & MMTk
 - Compare BC to MarkSweep, SemiSpace, CopyMS, GenCopy, GenMS



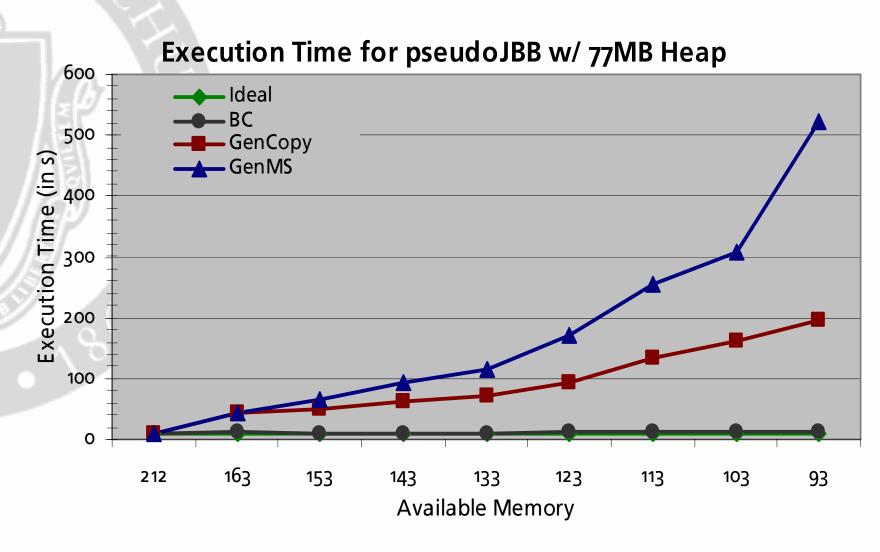
Throughput w/o Memory Pressure





BC runtime comparable to GenMS

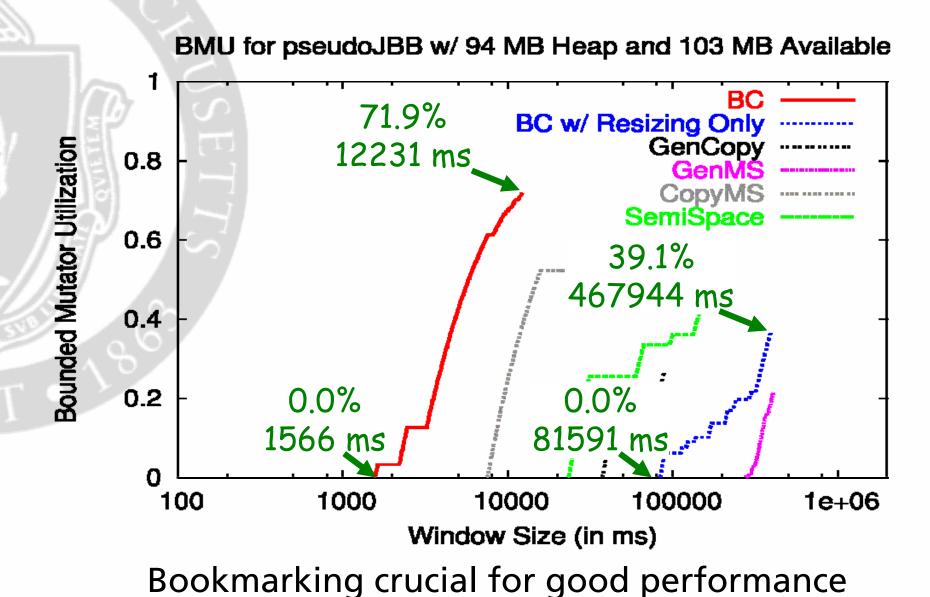
Throughput while Paging







BMU Curve w/ Memory Pressure





Summary of Results

- When not paging:
 - BC as fast as GenMS
- When paging:
 - vs. GenMS (fastest when not paging)
 - 41x faster, avg pause up to 218x smaller
 - vs. CopyMS (next fastest when paging)
 - 5x faster, avg pause up to 45x smaller

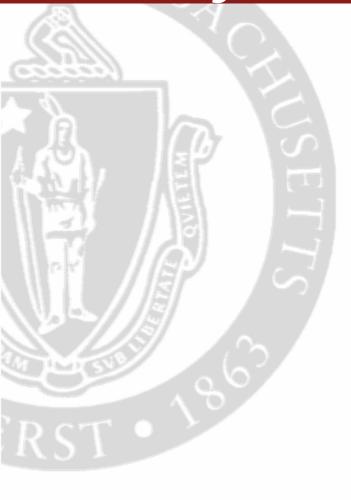


Conclusion

Bookmarking collector available at: http://www.cs.umass.edu/~hertz



Thank you





Other Pointer Summarizations

- Pointer Compression
 - Save space by compressing pointers
 - Tracing pointers becomes expensive
- Remembered Sets
 - Add evicted pointers to buffers
 - Requires space upon pages eviction
- Per Object Referring Page Counts
 - Increases BC overheads



Related Work

- VM-Sensitive Garbage Collection
 - Linearizing LISP lists [Bobrow/Murphy]
 - Does not know or use which heap pages are memory resident
 - Independent heap regions [Bishop]
 - Cannot avoid page faults when region contains evicted pages
 - Ephemeral garbage collection [Moon]
 - Must access evicted pages when they contain pointers into generations being collected



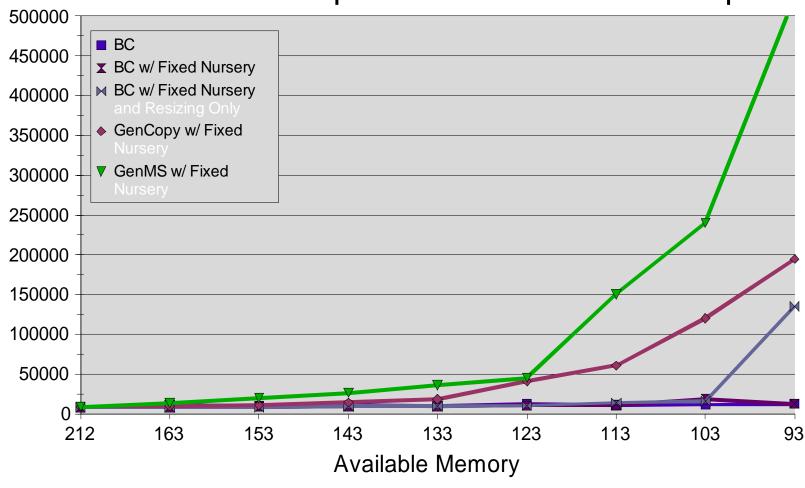
Related Work

- VM-Cooperative Garbage Collection
 - Discarding empty pages [Cooper, et al.]
 - No mechanism for evicting non-empty pages
 - Shrinking heap size [Alonso/Appel]
 - Responds to memory pressure changes only after a collection
 - Automatic heap sizing [Yang, et al.]
 - Orthogonal approach that determines proper heap size



Throughput w/ Memory Pressure



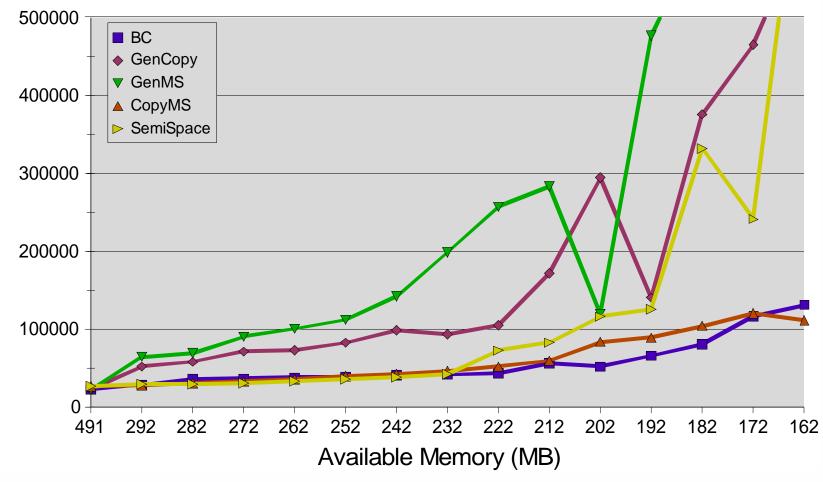


BC outperforms fixed nursery collectors



Multiprogramming Performance

Elapsed Time for 2 runs of pseudoJBB, 77MB Heap



BC performs well in many environments



Experimental Methodology

Benchmarks: SPECjvm98, ipsixql, jython, and pseudoJBB

- "Opt-and-reset" methodology
 - First iteration optimizes all code
 - Record results from second run

