# **UMassAmherst**

# Exterminator: Automatically Correcting Memory Errors with High Probability

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Amherst

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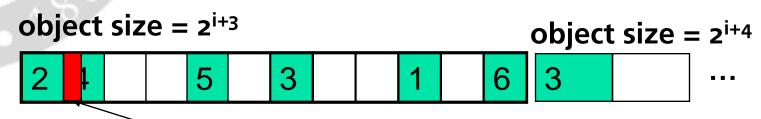
## Problems with Unsafe Languages

- C, C++: pervasive apps, but unsafe
- Numerous opportunities for security vulnerabilities, errors
  - Double/Invalid free
  - Uninitialized reads
  - Dangling pointers
  - Buffer overflows (stack & heap)
- DieHard: eliminates some, **probabilistically** avoids others [PLDI 2006]
  - Exterminator: builds on DieHard



#### **DieHard Overview**

- Use randomization & (optionally)
   replication to reduce risk of memory errors
  - Objects randomly spread across heap
- Different run = different heap
  - Probabilistic memory safety
    - Errors across heaps independent



Run 1: "malignant" overflow

Run 2: "benign" overflow



#### **DieHard Limitations**

#### DieHard:

- Fine for single error
  - But multiple errors eventually swamp probabilistic protection
  - Not great for large overflows
- Tolerates errors
  - But doesn't find them
  - No information for programmer
- Exterminator:

Automatically isolate and fix memory errors



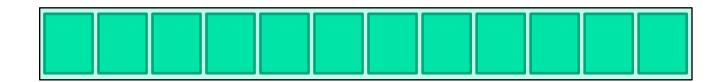
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  - Allocate object too small
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    - Not necessarily contiguous

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char * str = new char[8];
strcpy (str, "goodbye cruel world");
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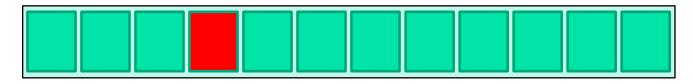
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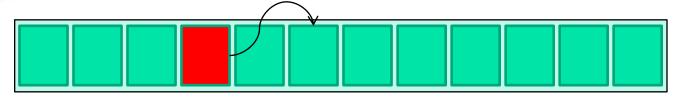
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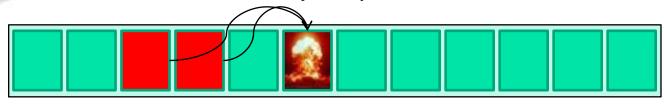


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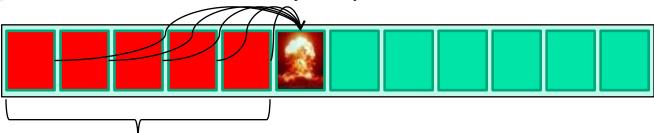




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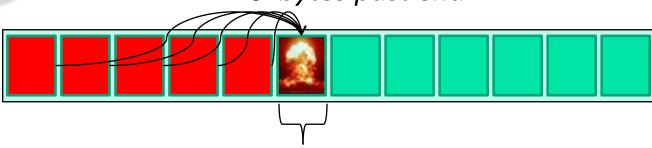


1. Heap provides no useful information

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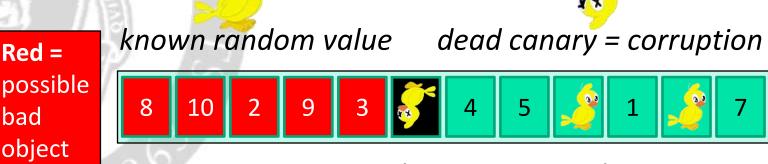
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2. No way to detect corruption

Canaries in freed space detect corruption



Green =
not
bad
object

# = object id (allocation time)



- Canaries in freed space detect corruption
  - Run multiple times with "DieFast" allocator

Red = possible bad object

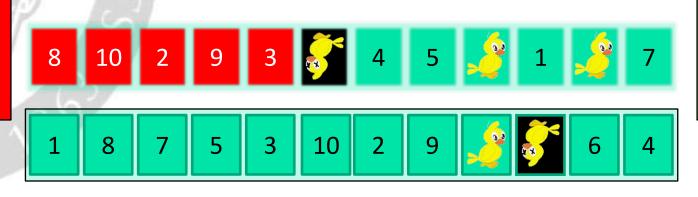


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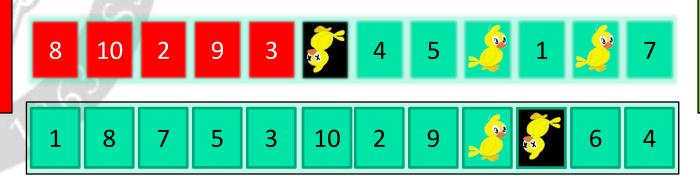


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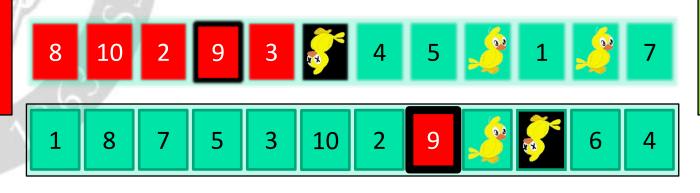


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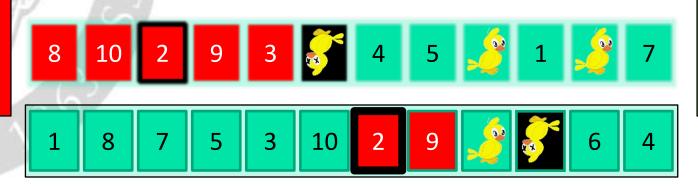


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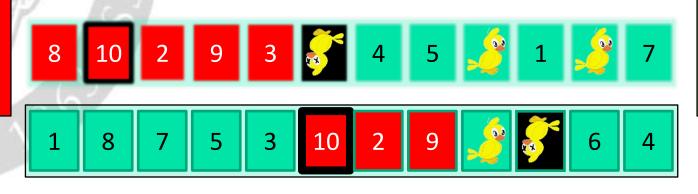


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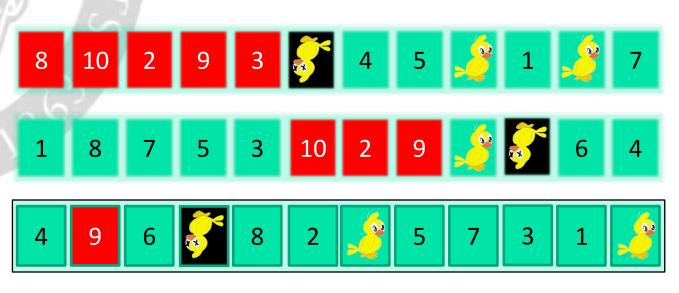
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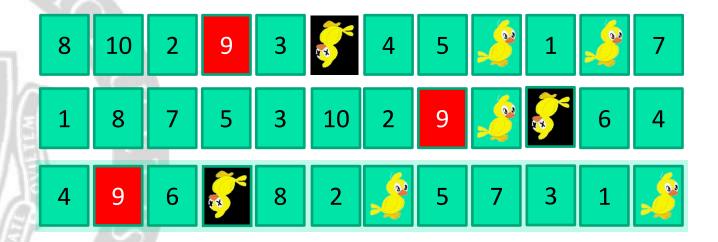
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⇒ object 9 overflowed, with high probability



# **Buffer Overflow Analysis**



$$E({\rm false\ positives}) = \frac{1}{(H-1)^{k-2}}$$
 H = # heap objects K = # iterations

- Example: H = 1,000,000 objects 3 iterations  $\approx \frac{1}{1,000,000}$  false positives
- Iterations exponentially increase precision



## **Isolating Dangling Pointers**

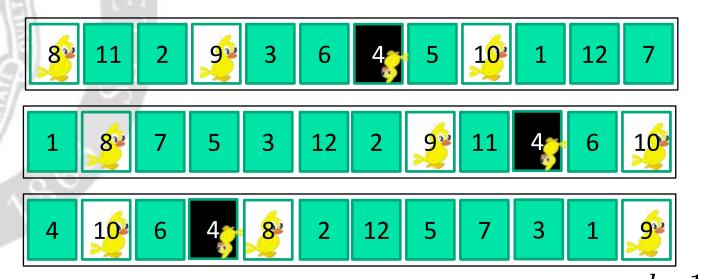
- Dangling pointer error:
  - Live object freed too soon
  - Overwritten by some other object

```
int * v = new int[4];
...
delete [] v; // oops
...
char * str = new char[16];
strcpy (str, "die, pointer");
v[3] = 12;
... use of v[0]
```



## **Isolating Dangling Pointers**

- Unlike buffer overflow:
  - dangling pointer ⇒ same corruption in all



$$P(identical overflow) \leq \left(\frac{1}{H}\right)$$

• 
$$k = 3 \Rightarrow false \ negatives \approx \frac{1}{1,000,000}$$



# **Correcting Allocator**

- Generate runtime patches to correct errors
  - Track object call sites in allocator
- Prevent overflows: pad overflowed objects

$$malloc(8) \Rightarrow malloc(8 + \delta)$$





$$\Rightarrow$$

1

Prevent dangling pointers: defer frees

free(ptr) 
$$\Rightarrow$$
 delay  $\delta$  mallocs; free(ptr)



#### **Exterminator Architecture**

- Three main pieces:
  - DieHard-based allocator (DieFast)
    - Reveals bugs
  - Error isolator
    - Finds bugs across multiple heaps w.h.p.
  - Correcting allocator
    - Fixes bugs
- Multiple modes suitable for testing (debugging) or deployment



#### **Exterminator Modes**

#### Iterative

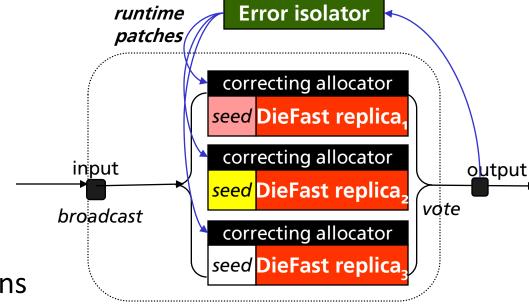
- Run multiple times
- Same inputs
- Debugging

#### Replicated

- Run simultaneously
- Deployable w/limitations
- Can fix errors on-the-fly

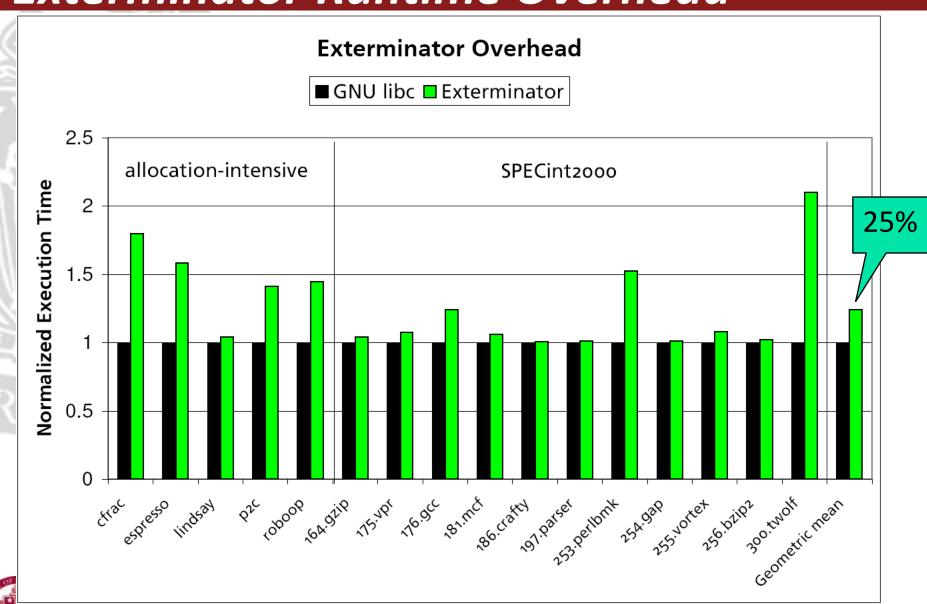
#### Cumulative

- Different inputs, nondeterminism
- Deployable; see paper for details





#### **Exterminator Runtime Overhead**



#### **Empirical Results: Real Faults**

- Squid heap overflow
  - Crashes glibc 2.8.0 and BDW collector
  - lacksquare 3 iterations to fix  $\Rightarrow$  6 byte pad
    - Prevents overflow for all subsequent executions



#### **Empirical Results: Real Faults**

#### Mozilla 1.7.3 buffer overflow

#### Debug scenario:

repeated load of PoC: 23 runs to fix overflow







#### Deployed scenario:

different browsing sessions: 34 runs to fix





#### **Exterminator Conclusion**

- Exterminator: automatic error correction w.h.p.
  - Randomization ⇒ bugs have different effects
  - Statistical analysis combines information from multiple runs to isolate error
  - Correcting allocator eliminates bugs at runtime

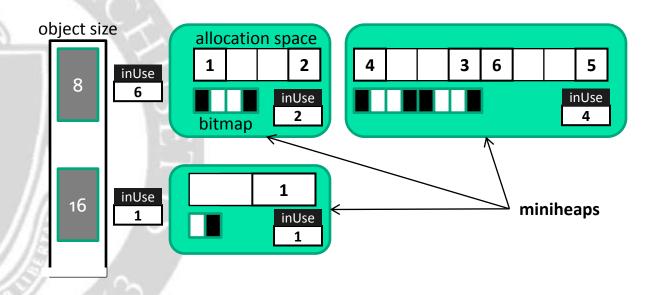
http://www.cs.umass.edu/~gnovark/







# DieHard, heap layout



- Bitmap-based, **segregated** size classes
  - Bit represents one object of given size
    - i.e., one bit =  $2^{i+3}$  bytes, etc.
- malloc(): randomly probe bitmap for free space
- free(): just reset bit

#### **Exterminator Extensions**

