MC²: High Performance GC for Memory-Constrained Environments

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Motivation

- Java widely used
  - Safety
  - Portability
- Garbage collector requirements
  - High throughput
  - Short pauses
  - Good memory utilization

Motivation

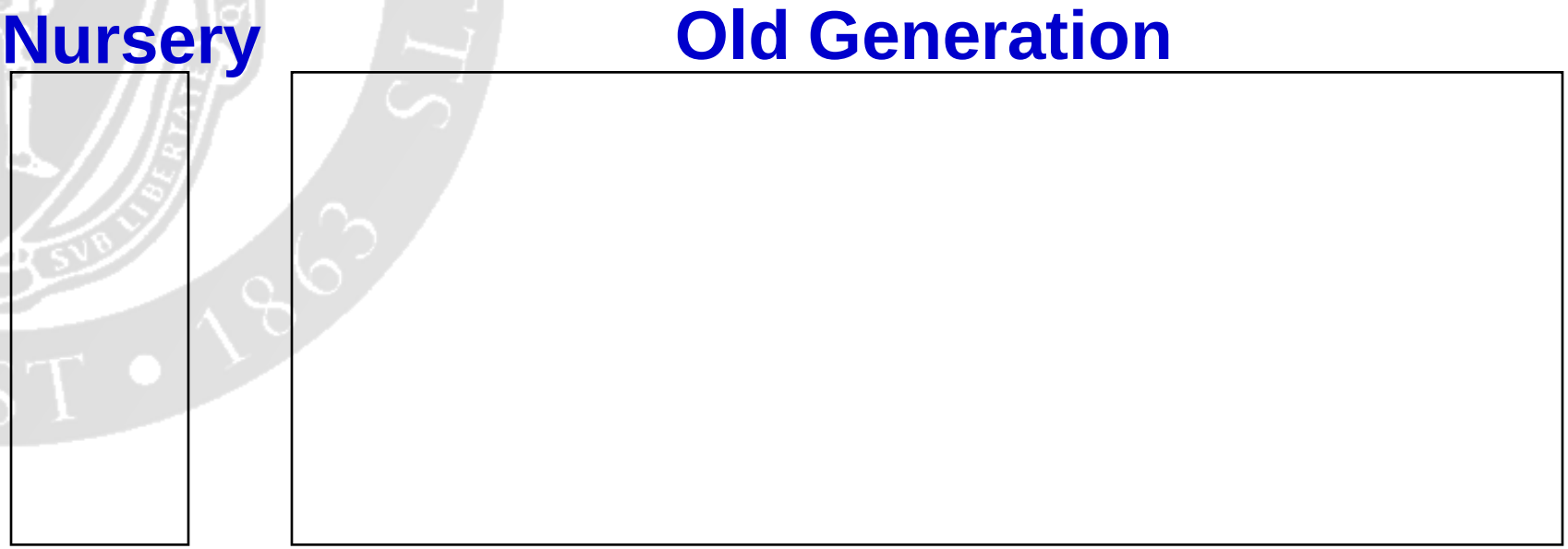
- Handheld Devices
  - Cellular phones, PDAs widely used
  - Constrained memory
- Diverse applications
  - Media players
  - Video games
  - Digital cameras
  - GPS
  - Scaled down desktop apps (e-mail, browser etc.)
- Require high throughput, short response time

Talk Outline

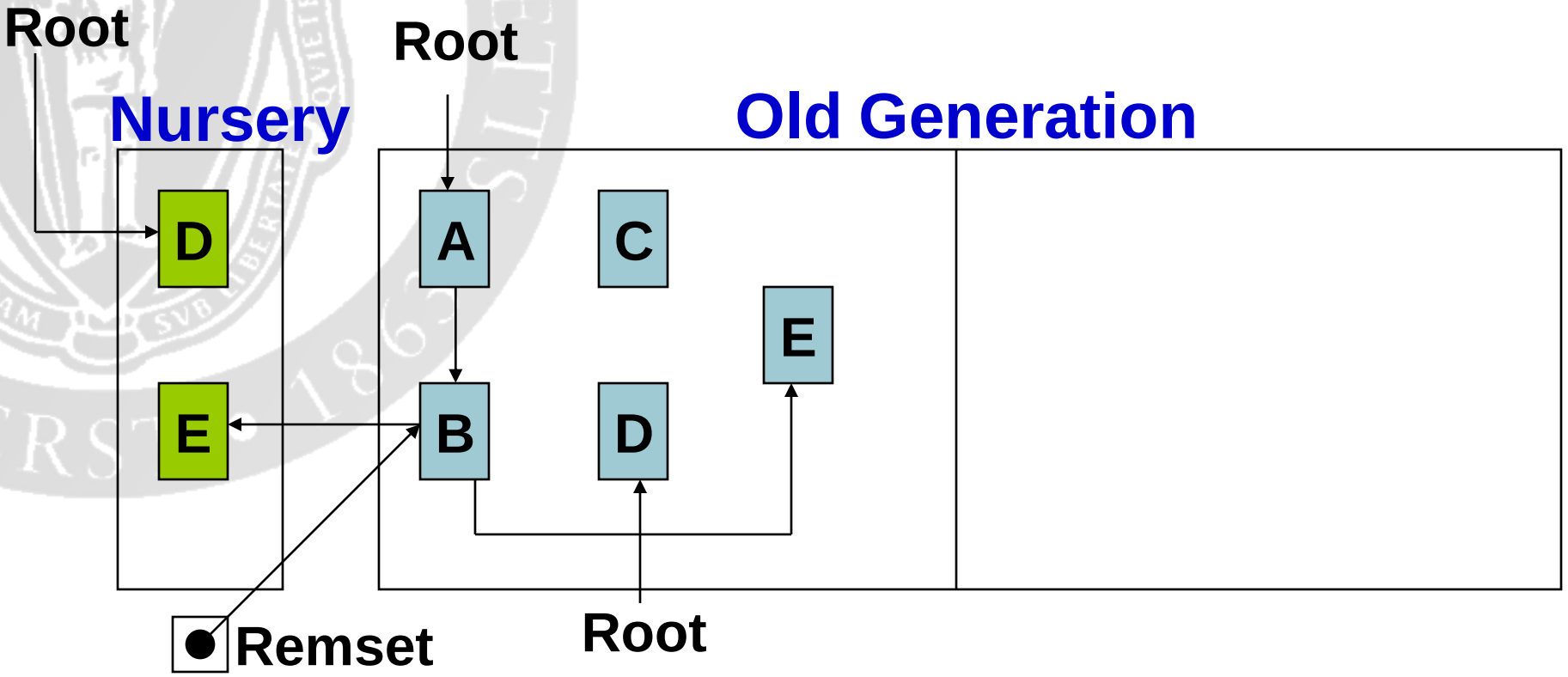
- Generational collection
- MC^2 overview
- Algorithmic Details
- Experimental Results
- Conclusions

Generational Collection

- Divide heap into regions called generations
- Generations segregate objects by age
- Focus GC effort on younger objects



Generational Copying Collection



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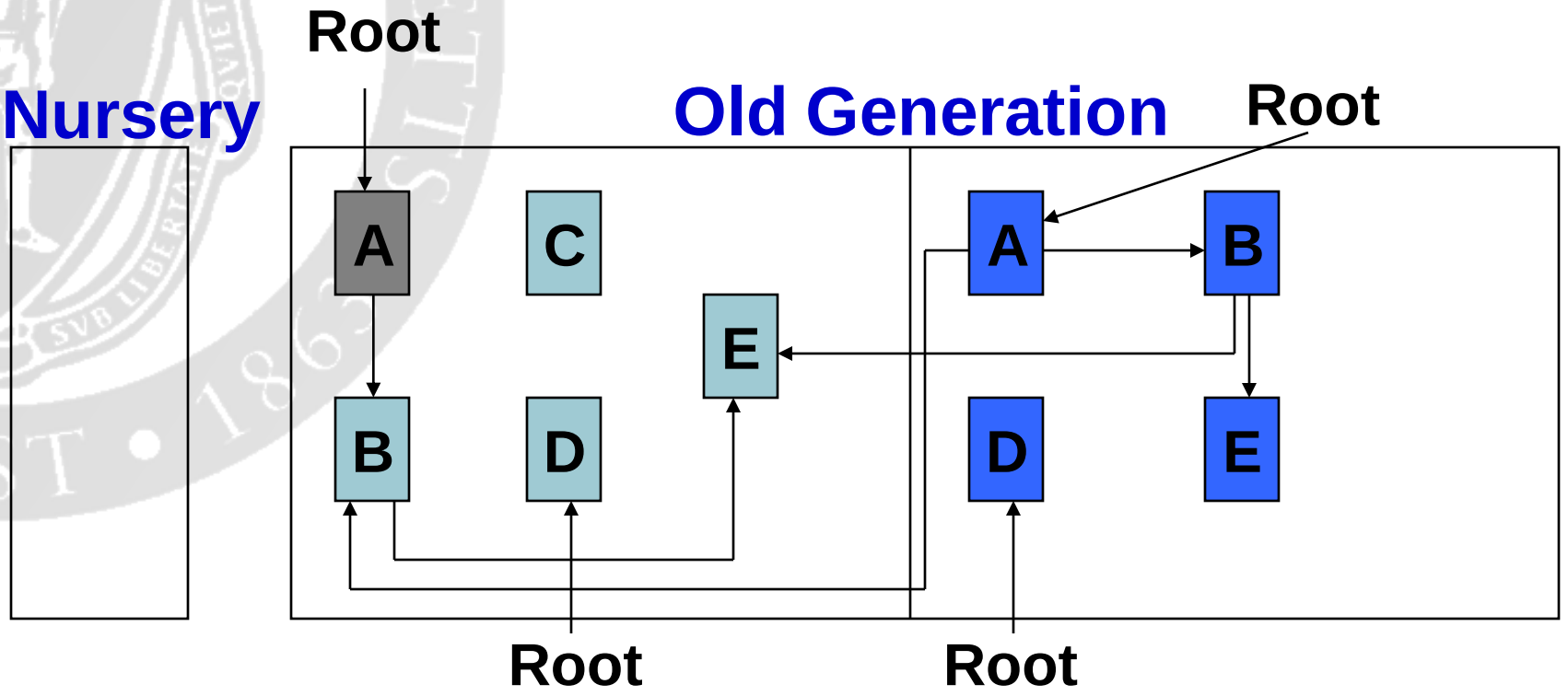
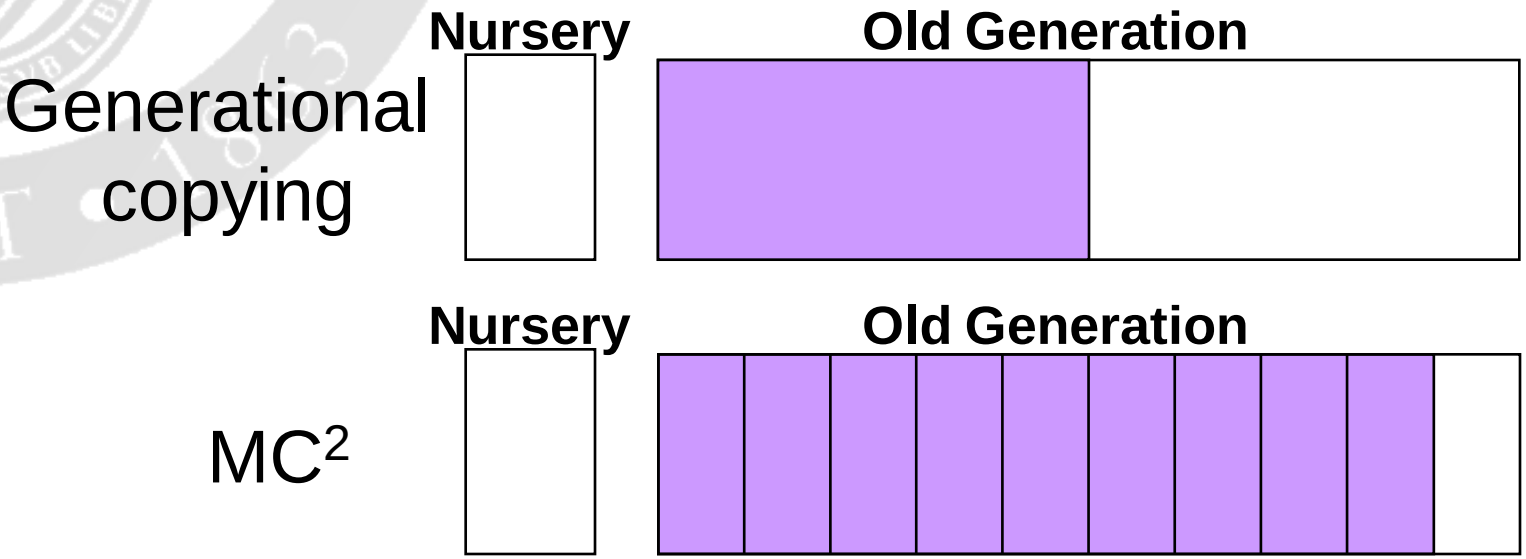


MC² Overview

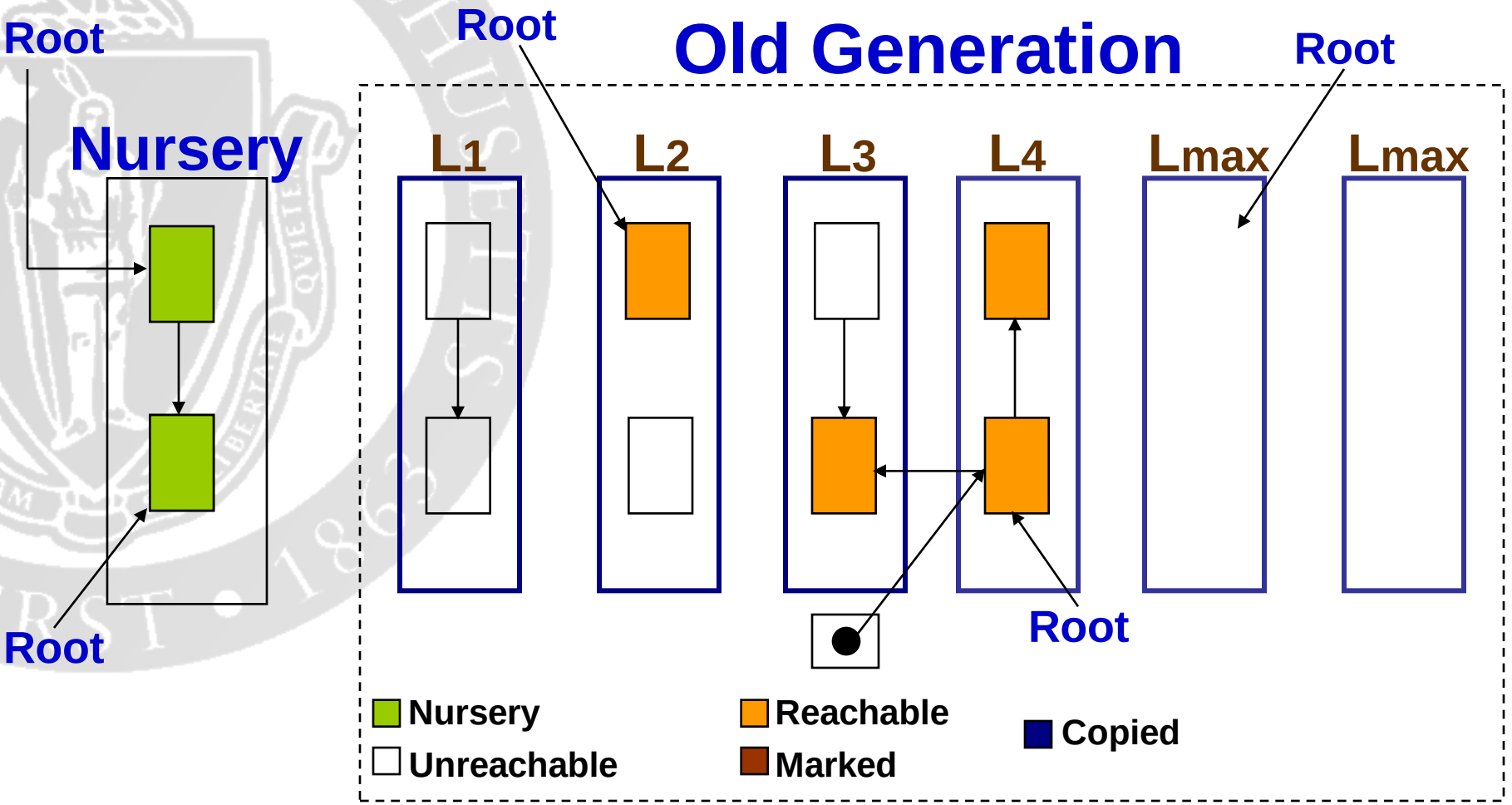
- Extends gen. copying, overcomes 2X overhead
- Divides space into equal size *windows*
- Reserves one or more windows for copying
- Collects in two phases: *mark* and *copy*



MC² – Mark Phase

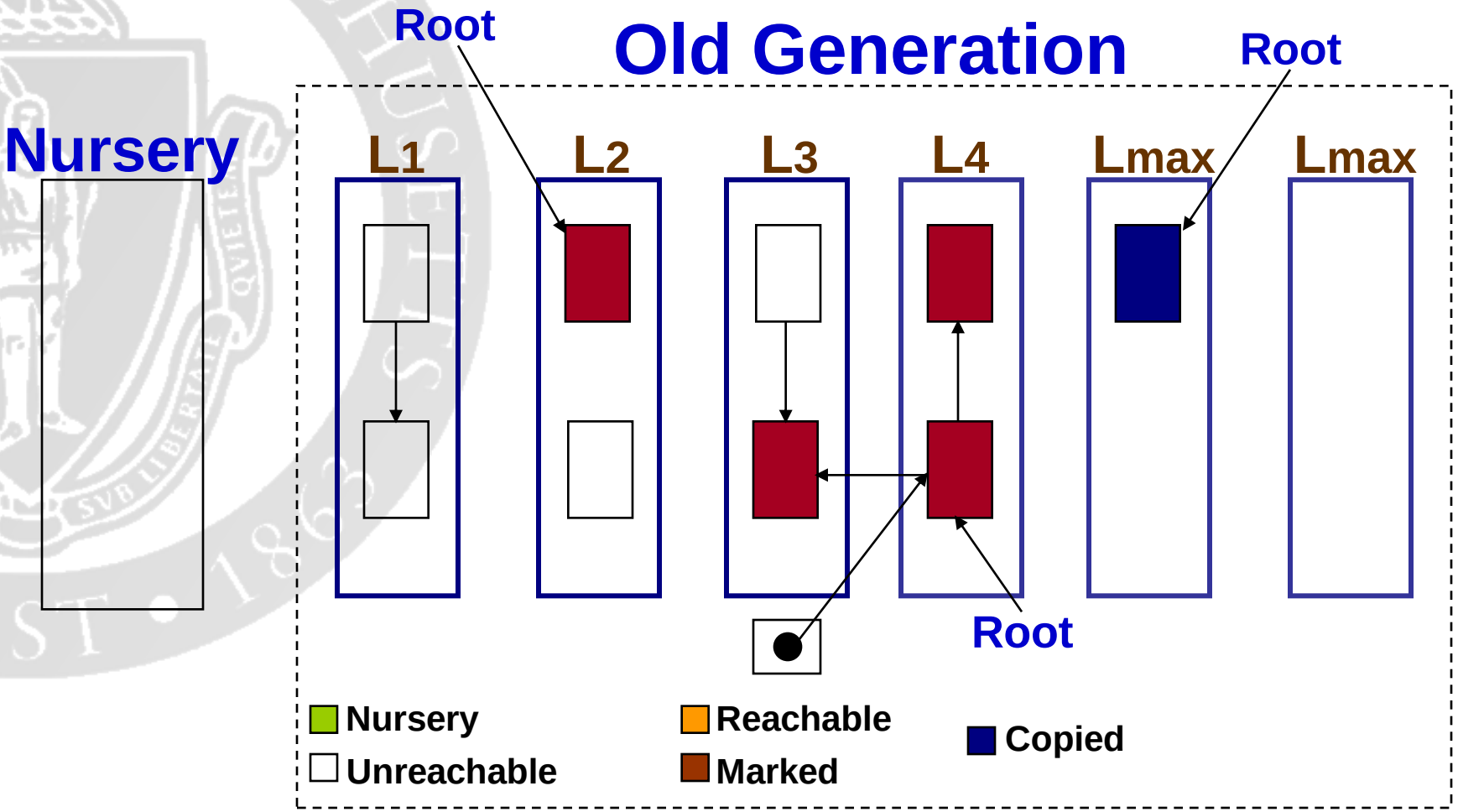
- Logically order old gen. windows
- Three mark phase tasks:
  - Mark reachable objects
  - Calculate live data volume in each window
  - Build per-window remembered sets
- Start when old gen. getting full: 80%, say
- Interleave marking with nursery allocation
  - Do some marking after every *n* bytes allocated
  - Reduces mark phase pauses

MC² Example – Mark Phase



- MC² marks two objects during nursery allocation

MC² Example – Classify Windows



- Locate high-occupancy windows and discard remsets

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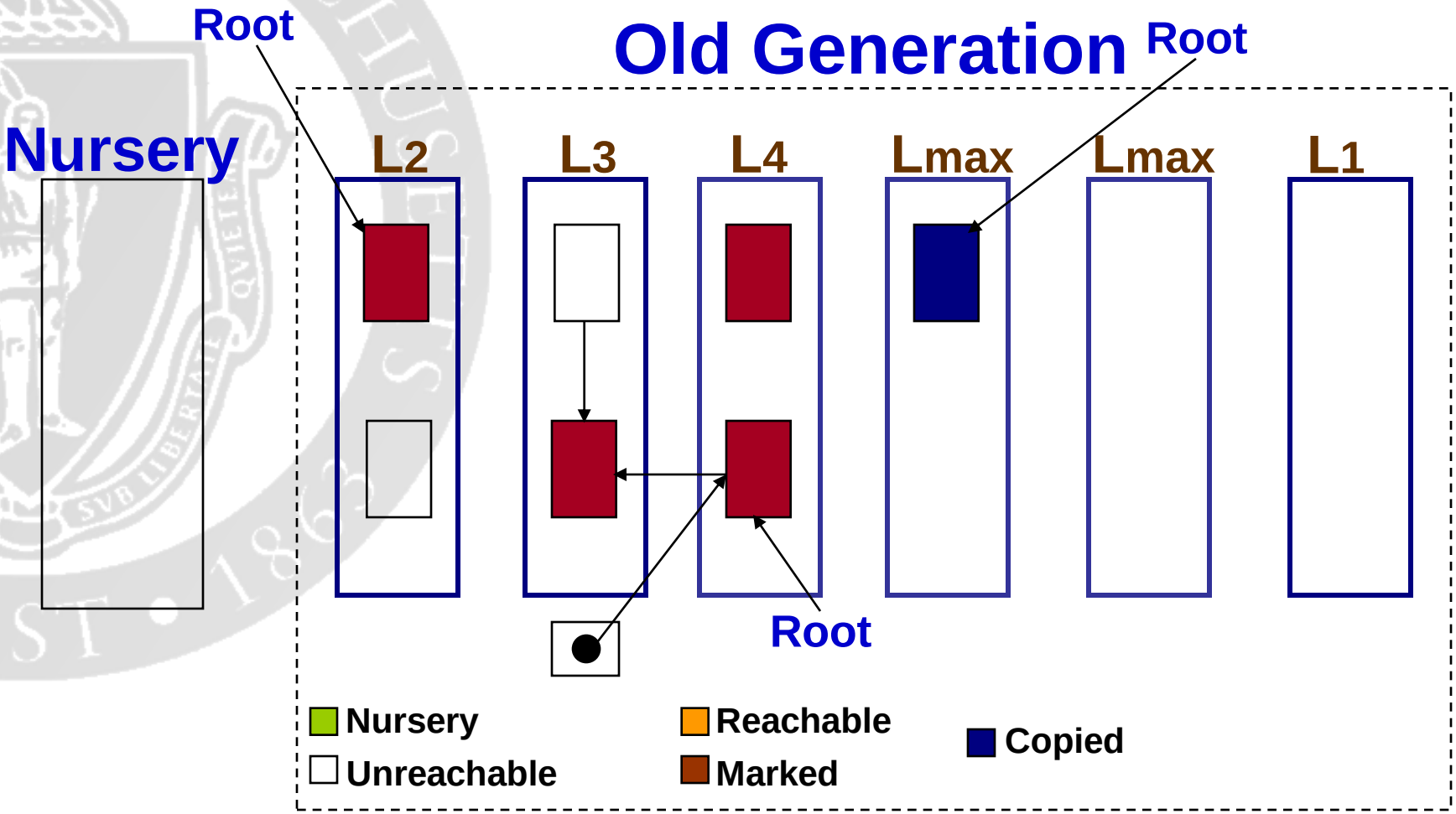

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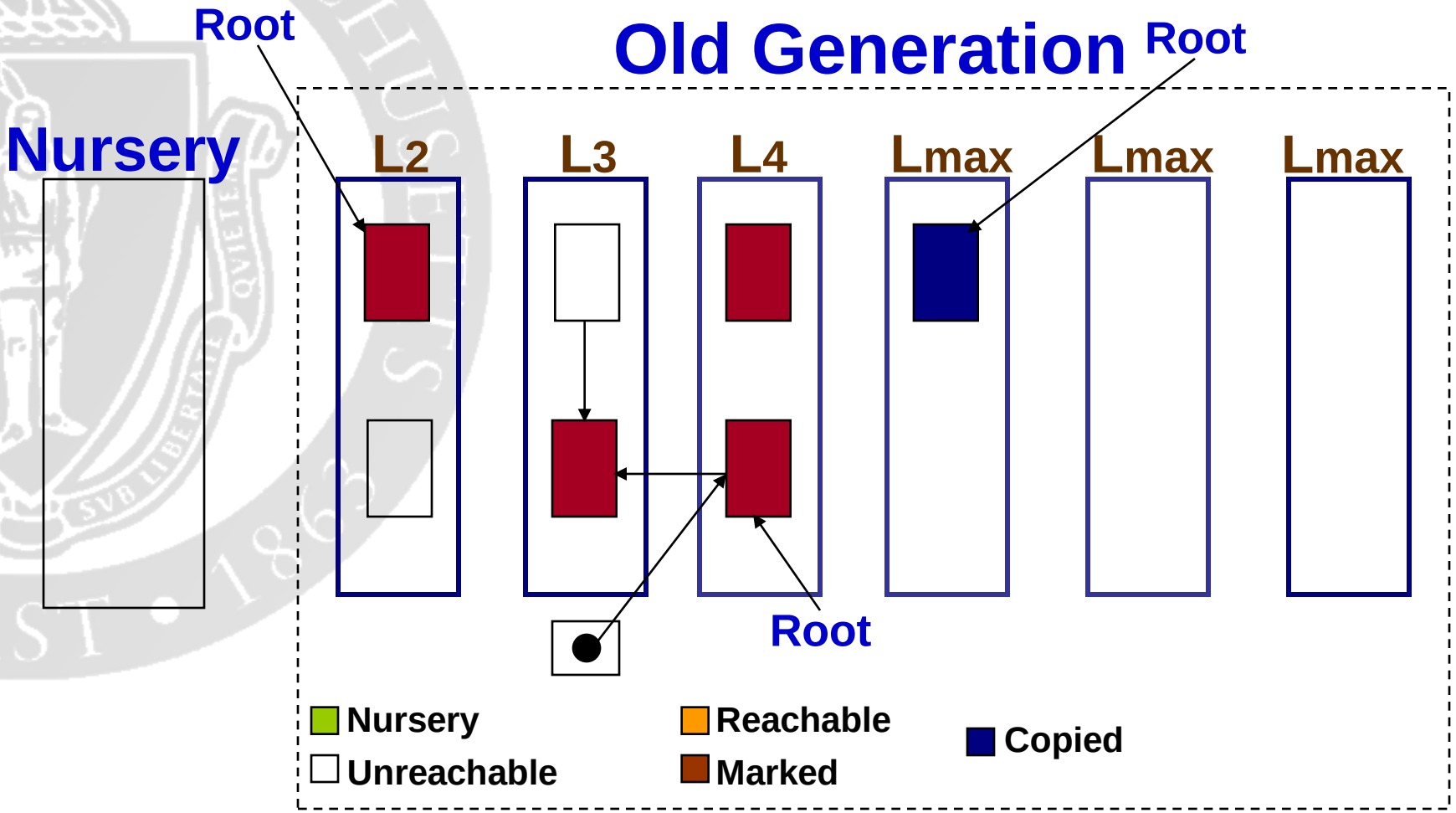
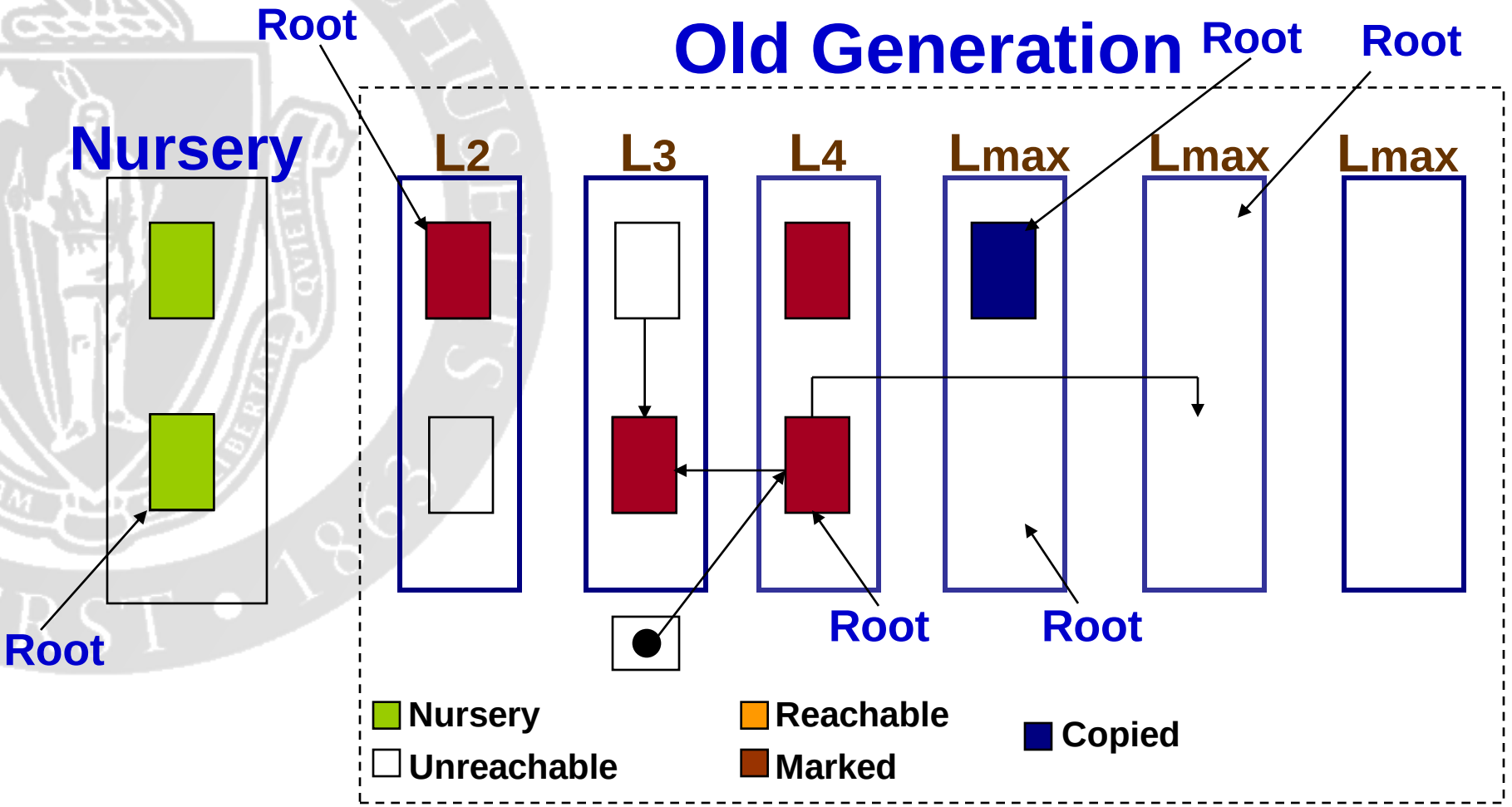


- Locate high-occupancy windows and discard remsets

MC² – Copy Phase

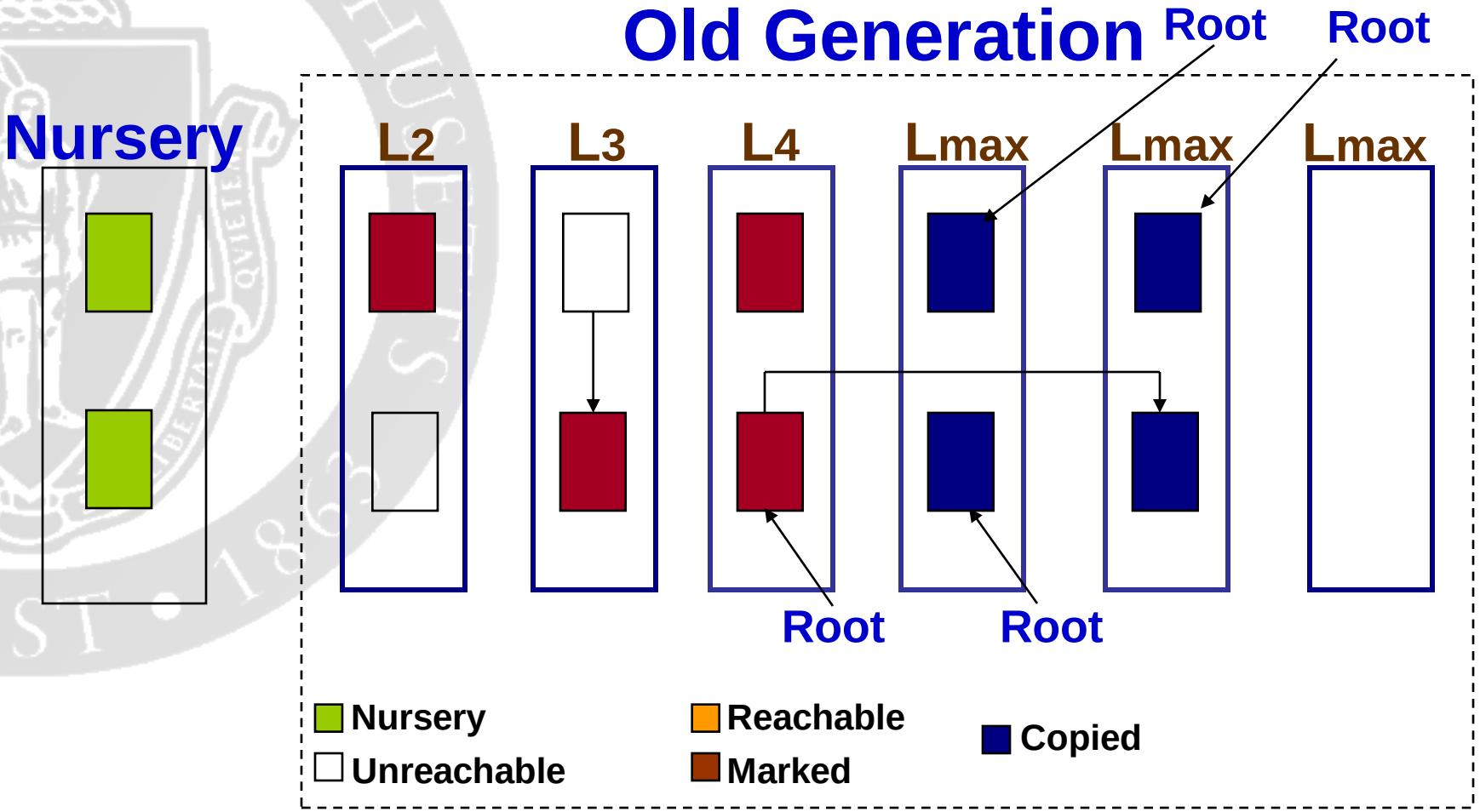
- Copy and compact reachable data
- Performed in small increments
  - One windowful of live data copied per increment
- One increment per nursery collection
- High-occupancy windows copied *logically*

MC² Example – Copy Phase



- Copy a window worth of data during nursery collection

MC² Example – Copy Phase



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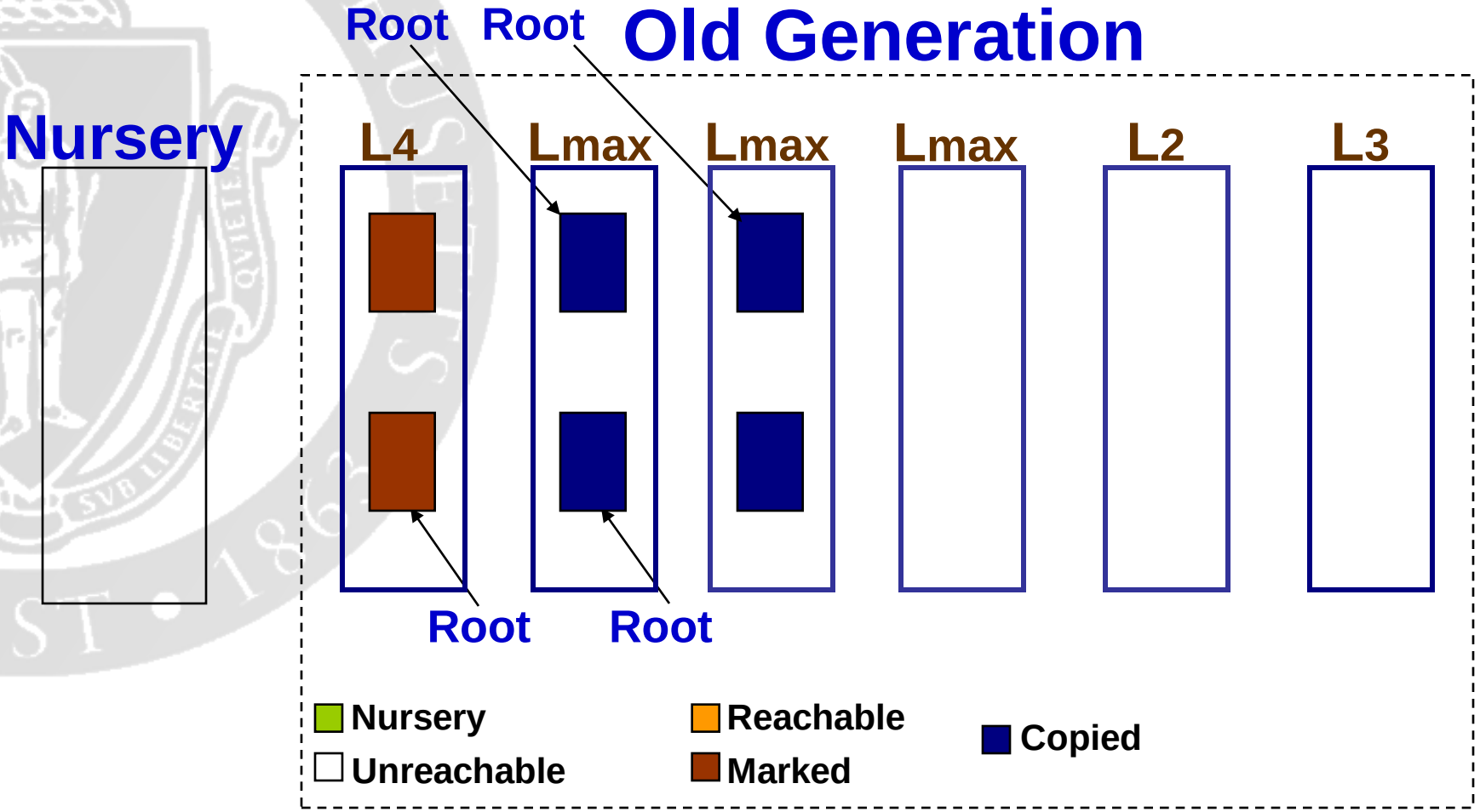

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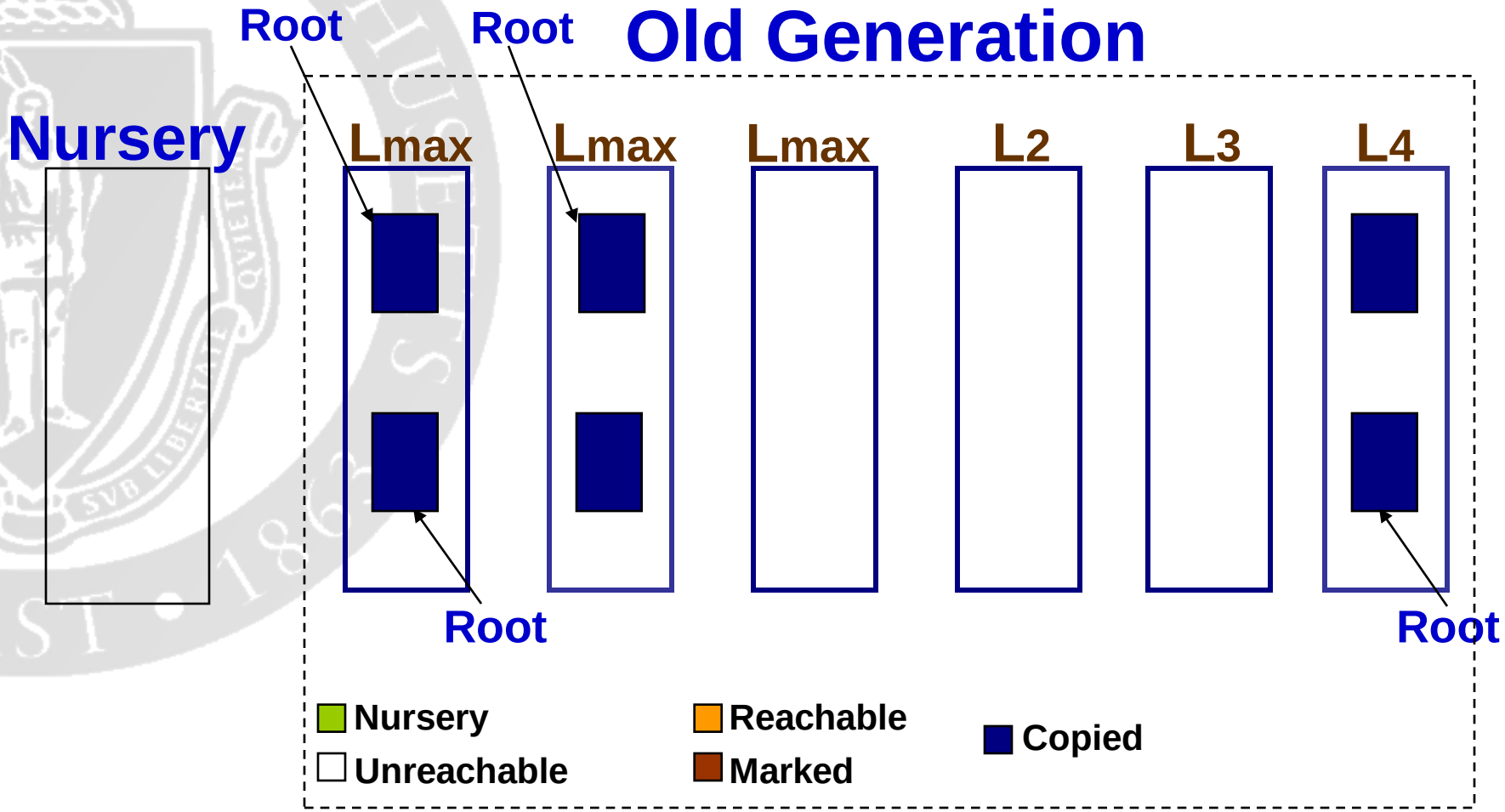

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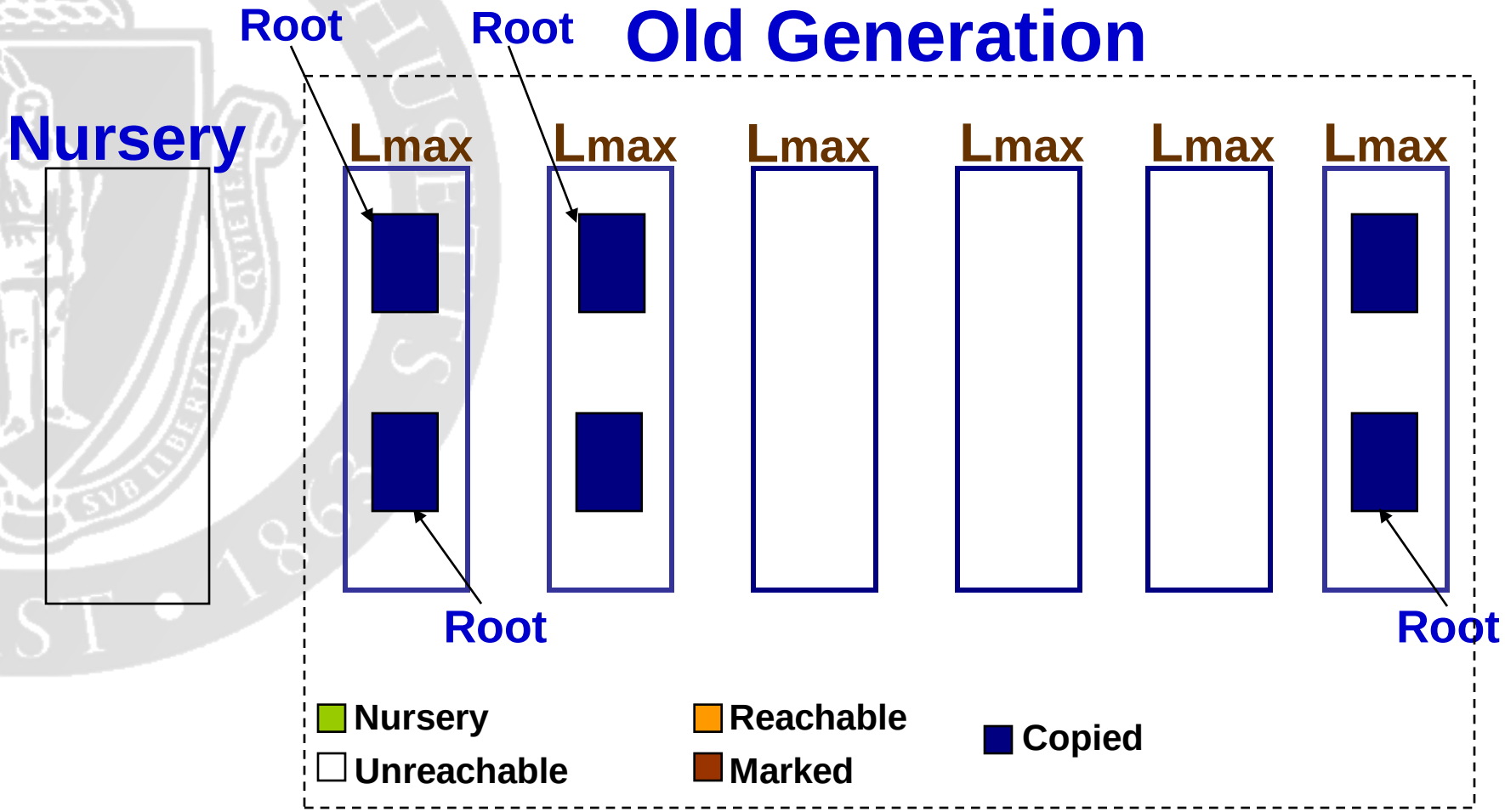

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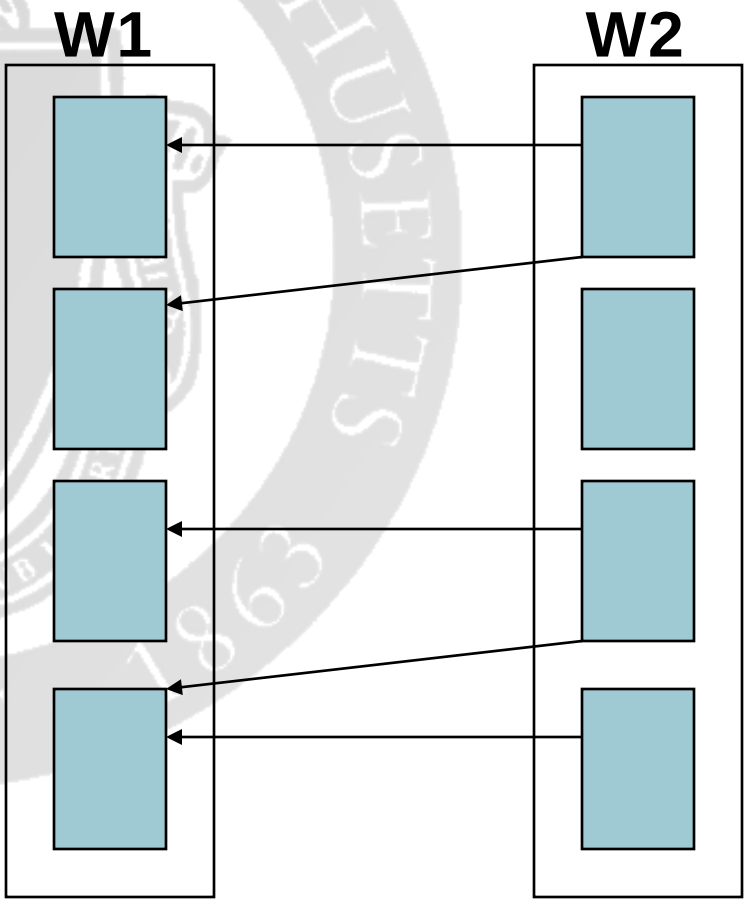
Algorithmic Details

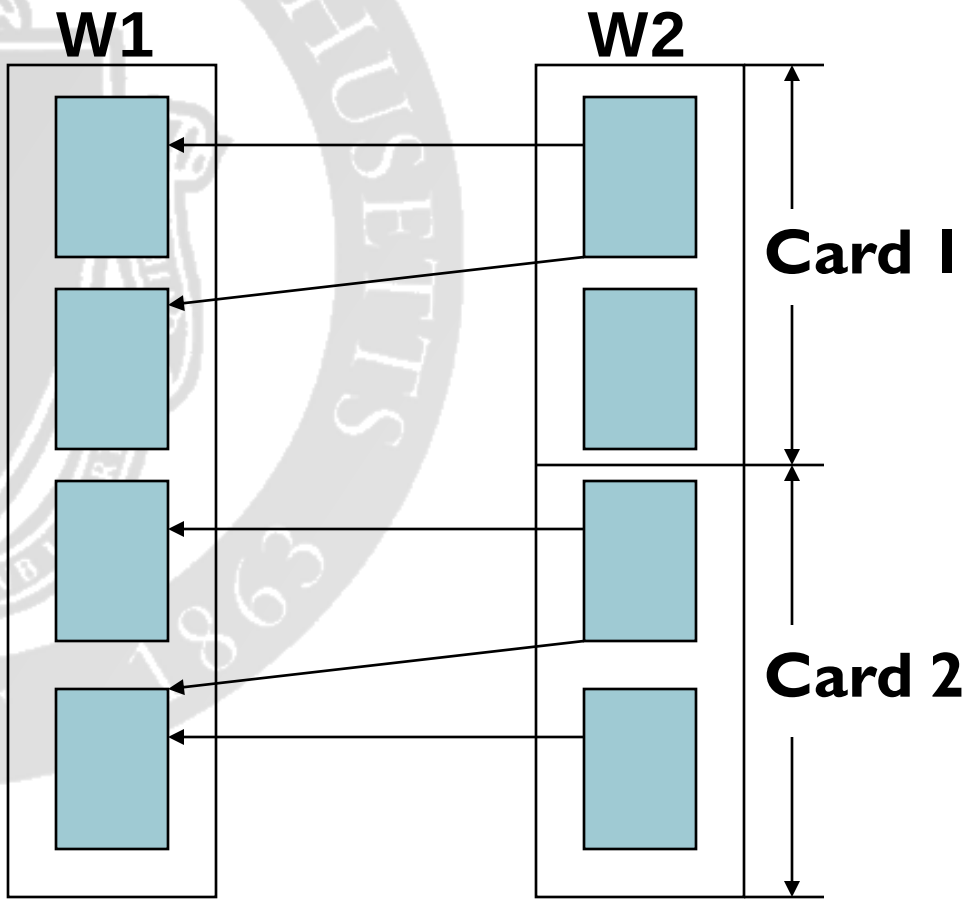
- Large remembered sets
  - Bounding space overhead
- Popular objects
  - Preventing long pauses

Handling large remembered sets



- Normal remembered set for W1 stores 5 pointers (20 bytes on a 32 bit machine)

Handling large remembered sets



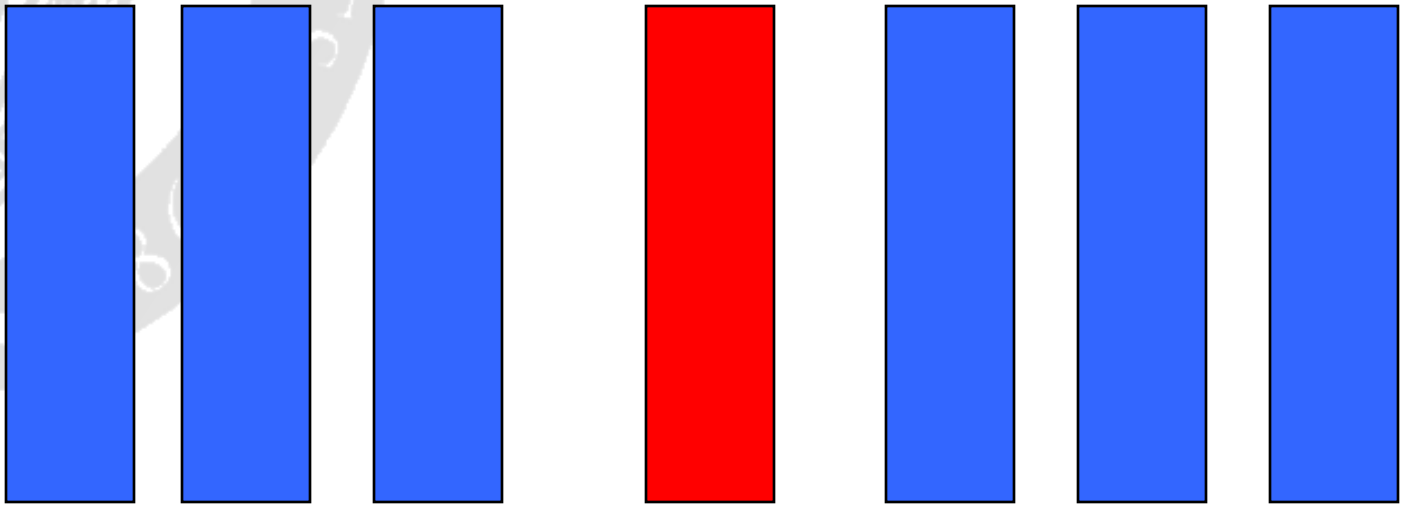
- Card table requires only 2 bytes (one per card)

Handling large remembered sets

- Set a limit on the total remembered set size (e.g., 5% of total heap space).
- Replace large remsets with card table when total size approaches limit
- Good tradeoff between speed and space utilization

Handling popular objects

- Divide heap into small *physical* windows
- Normally copy a group of physical windows
- Popular physical windows not grouped



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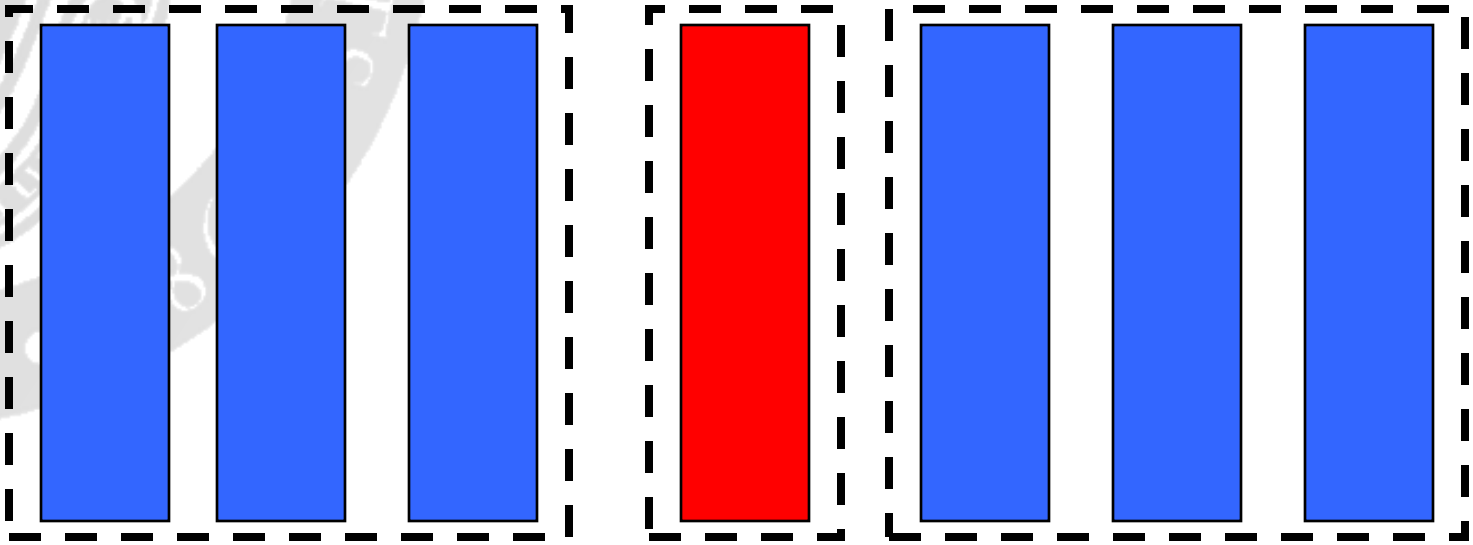


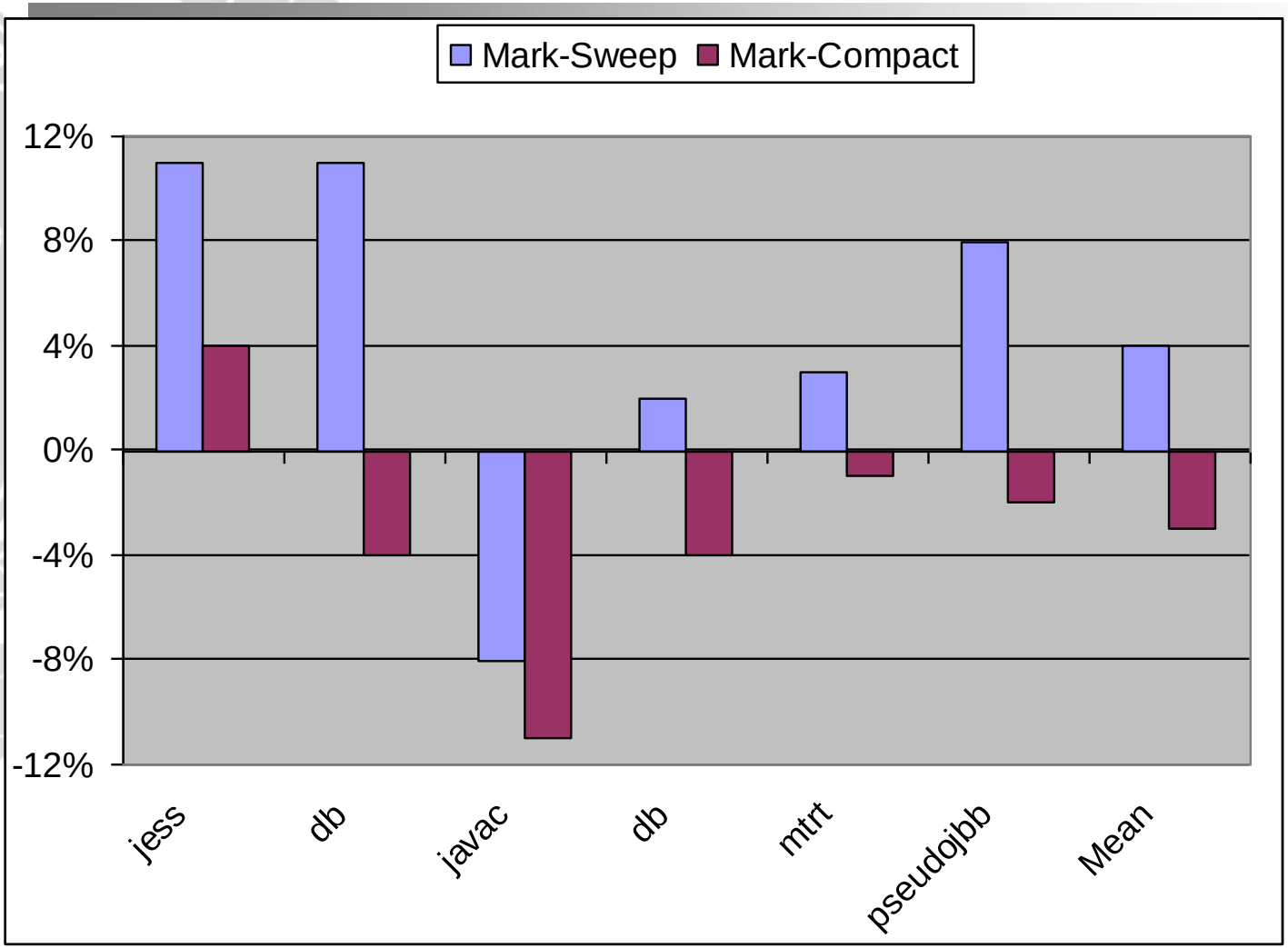
Handling popular objects

- Identify popular object while converting remset to card table
- Isolate popular object at high end of heap
  - Do not need to maintain references to the objects
- Copying a popular object can cause a long pause, but does not recur

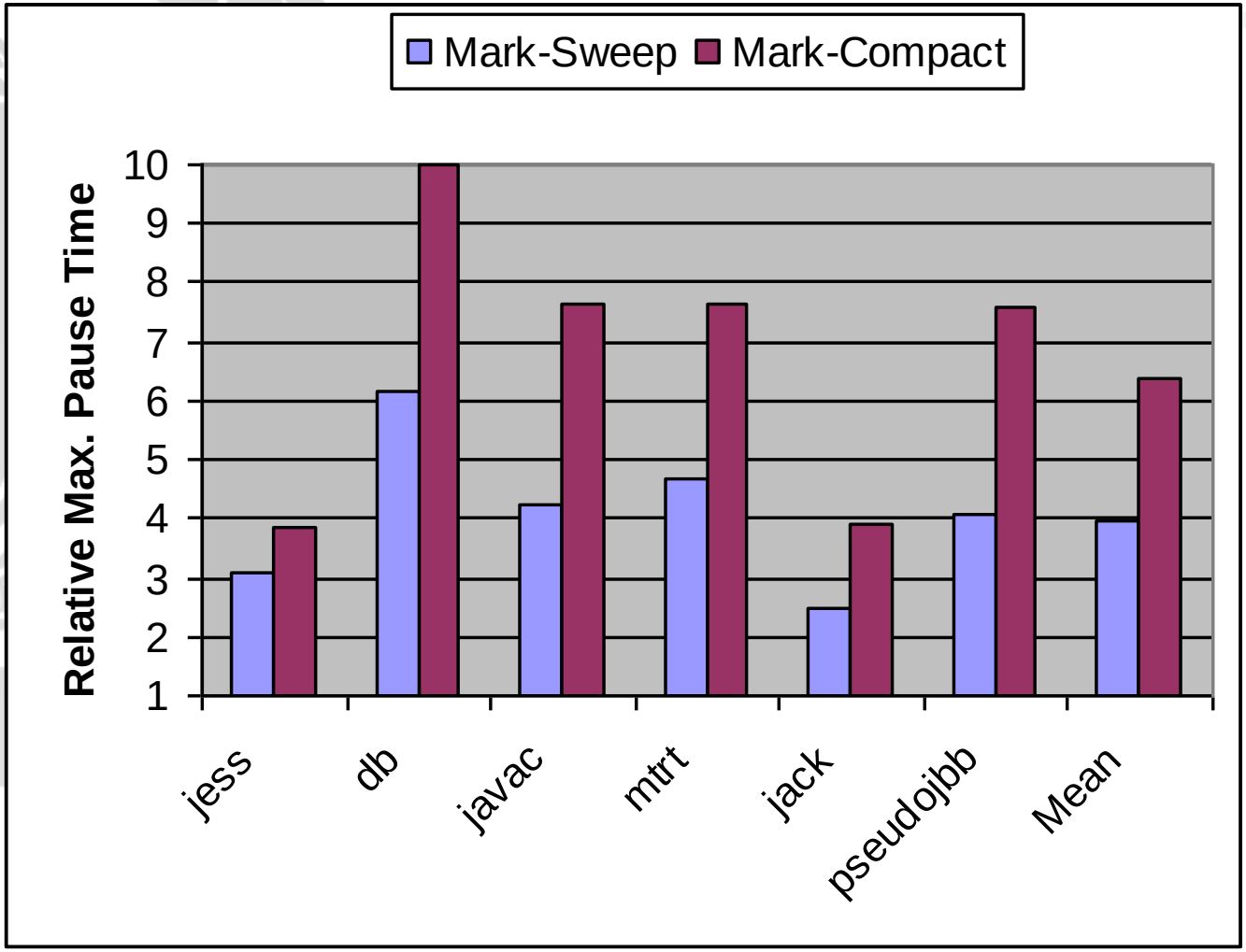
Experimental Results

- Implemented in Jikes RVM (2.2.3)/MMTk
- Pentium 4 1.7 GHz, 512MB memory, RedHat Linux 2.4.7-10
- Benchmarks: SPECjvm98, pseudojbb
- Collectors evaluated
  - Generational Mark-Sweep (MS)
  - Generational Mark-(Sweep)-Compact (MSC)
  - MC^2
- MSC, MC^2 use separate code and data spaces
- Results: Execution time, pause time in a heap 1.8x program live size

Execution Time relative to MC^2 (Heap Size = 1.8x max. live size)

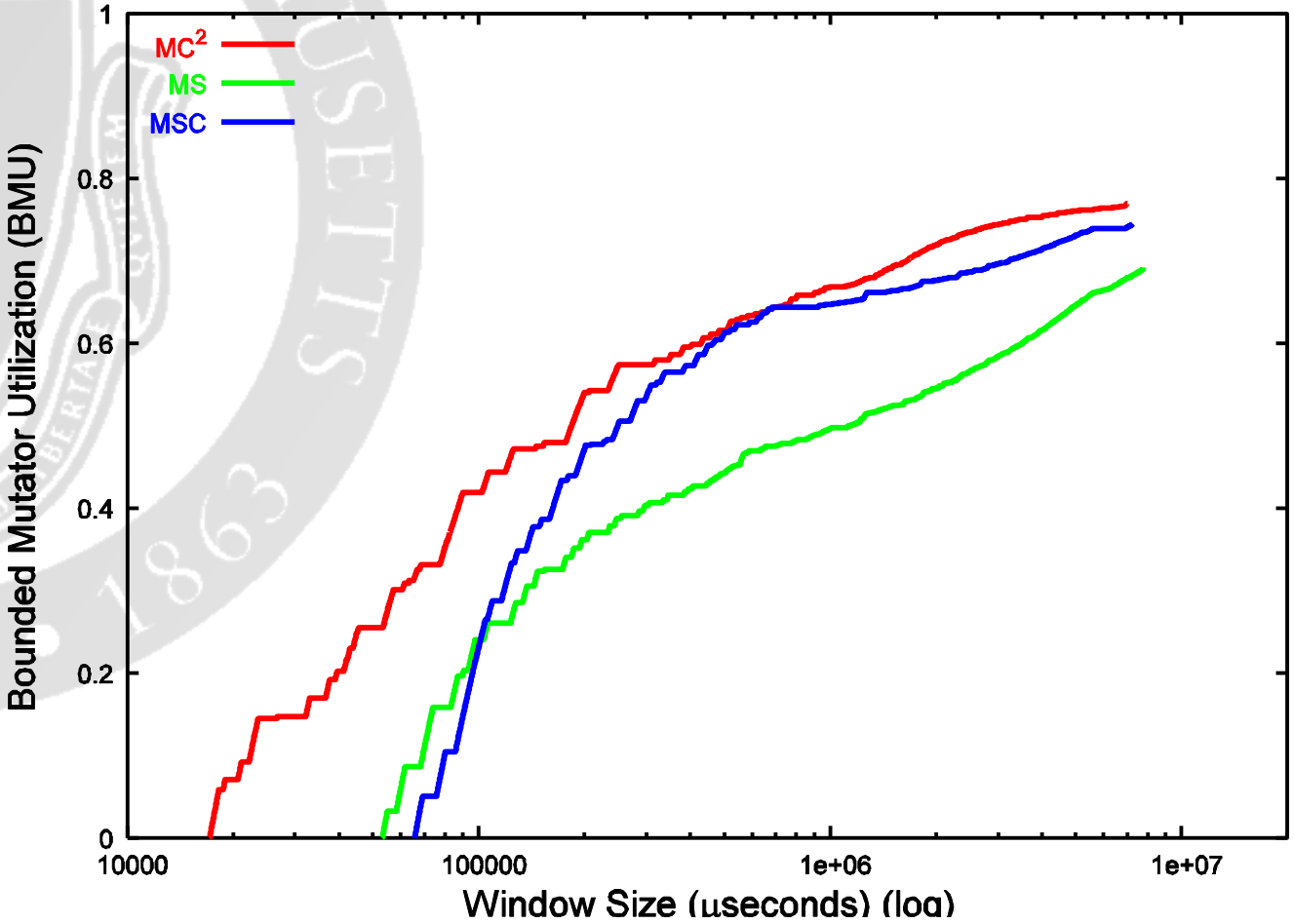


Max. Pause Time relative to MC^2 (Heap Size = 1.8x max. live size)

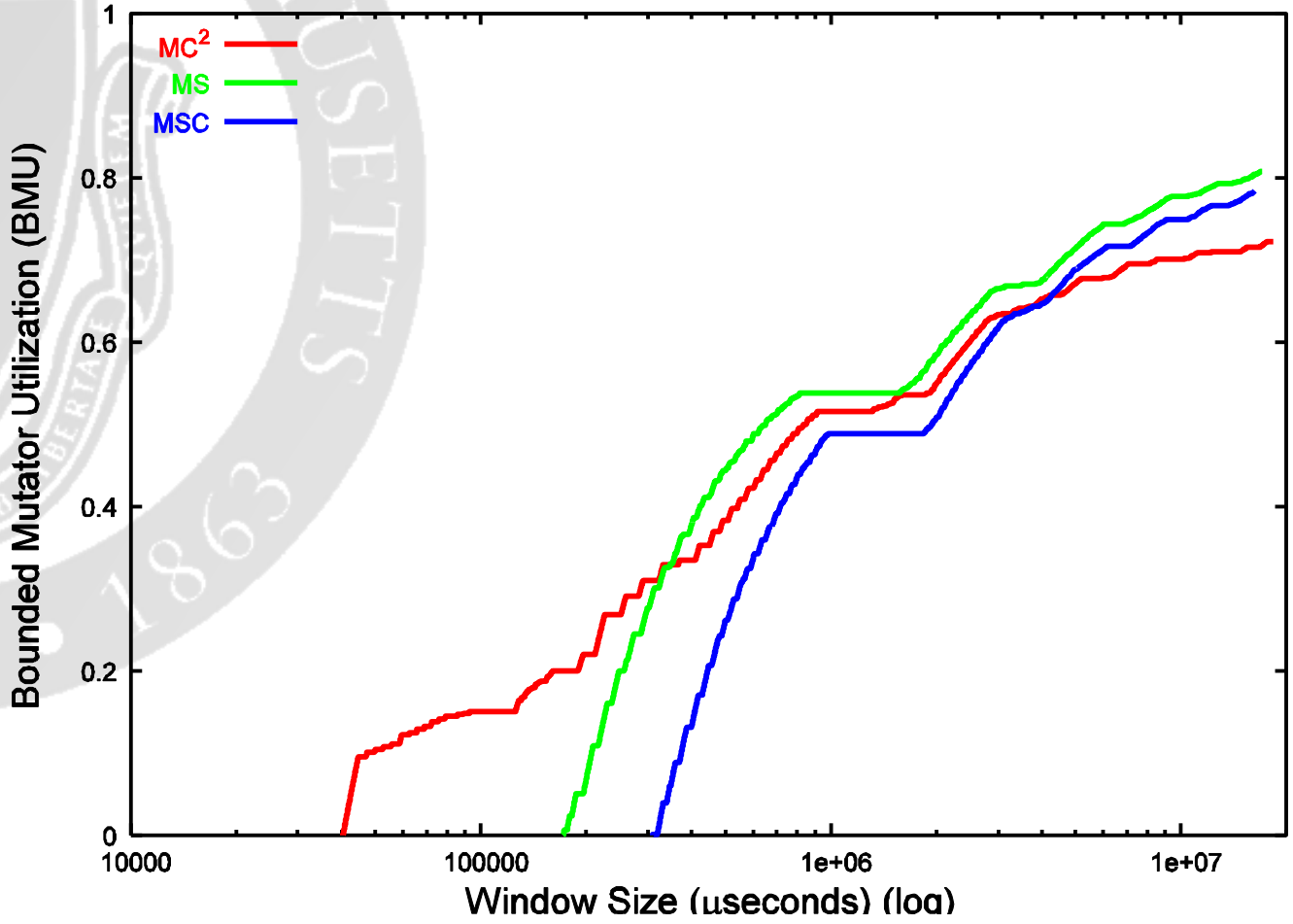


- MC^2 max. pause range: 17-41ms

Bounded Mutator Utilization (jess)



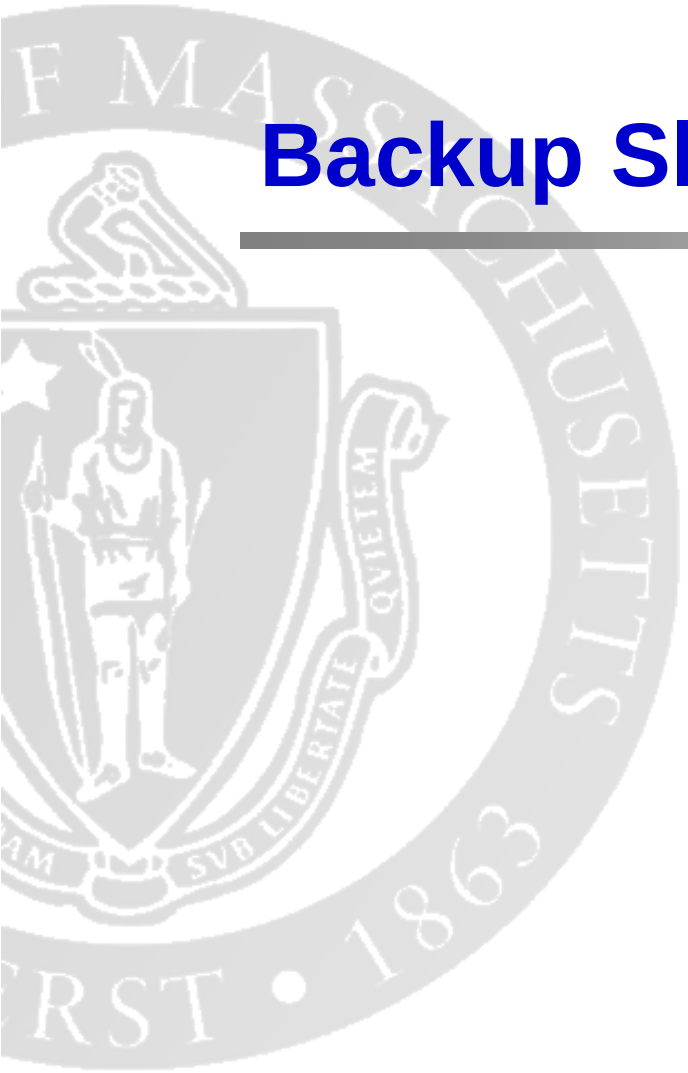
Bounded Mutator Utilization (javac)



Conclusions

- MC^2: suitable for handheld devices with soft real time requirements
  - Low space overhead (50-80%)
  - Good throughput (3-4% slower than non-incremental compacting collector)
  - Short pause times (17-41ms, factor of 6 lower than non-incremental compacting collector)
  - Well distributed pauses
- Also suitable for desktop environments

Backup Slides



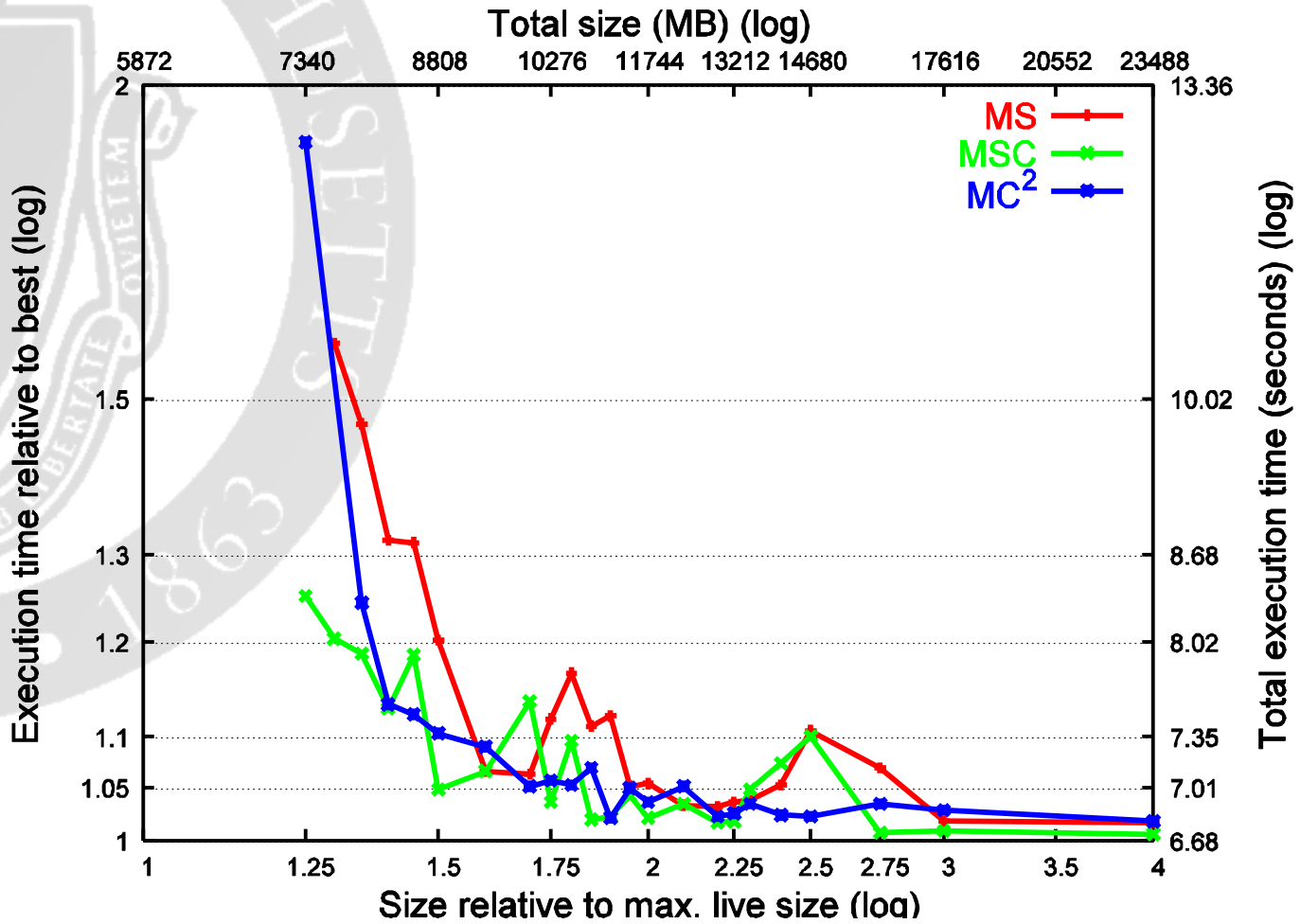
Traditional Tracing Collectors

- Mark-Sweep
  - Fragmentation
  - Locality effects
- Mark-Compact
  - Long pauses
- Copying
  - 2X space overhead

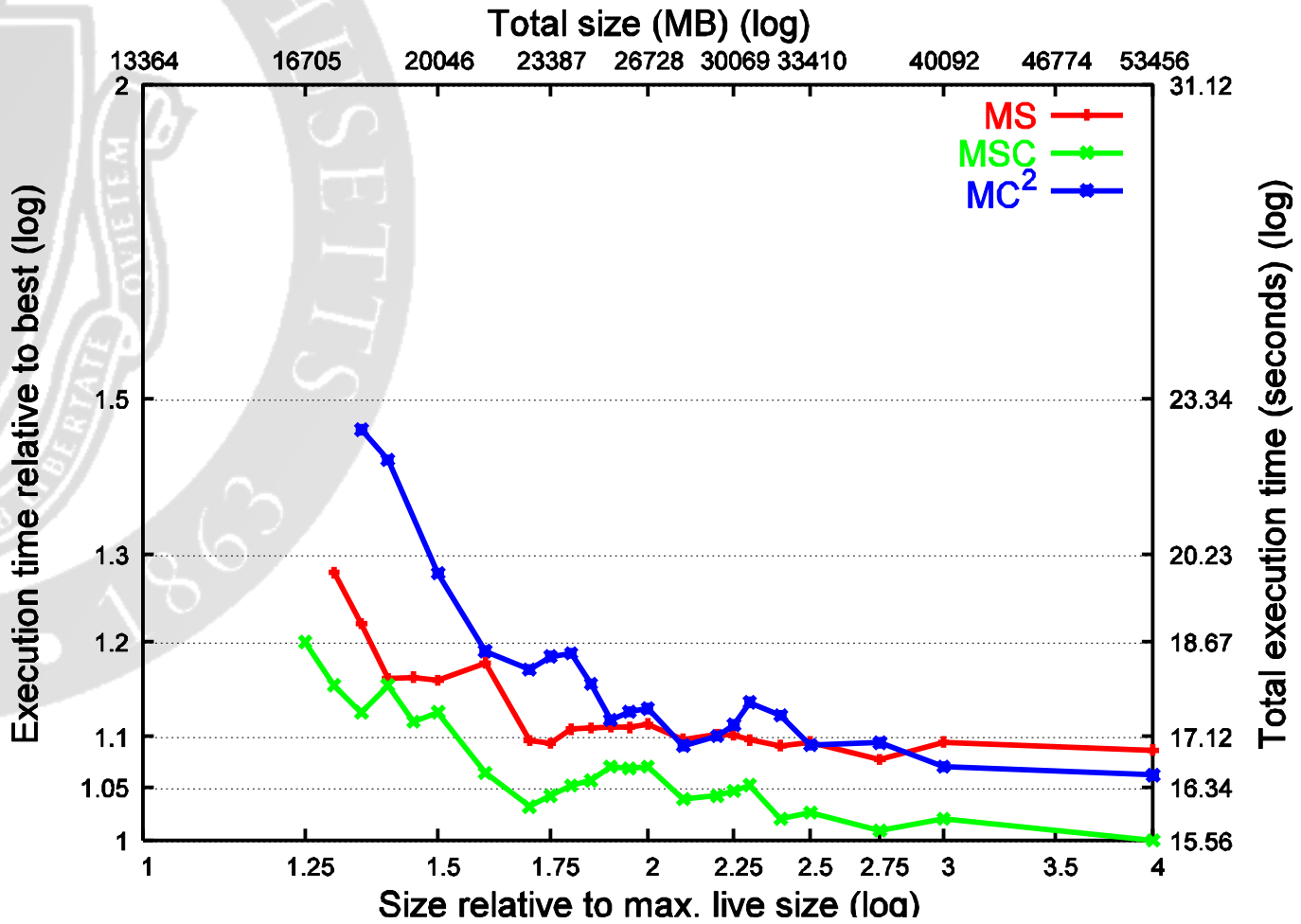
Modern incremental collectors

- Train collector – Hudson and Moss
  - Performs well in moderate size heaps
  - High copying overhead
- Lang and Dupont, Ben Yitzhak et al
  - Good throughput, pause times
  - Do not address metadata overheads
- Bacon, Cheng, Rajan
  - Address memory-constrained device requirements
  - Require advanced compiler optimizations

jess Execution Time



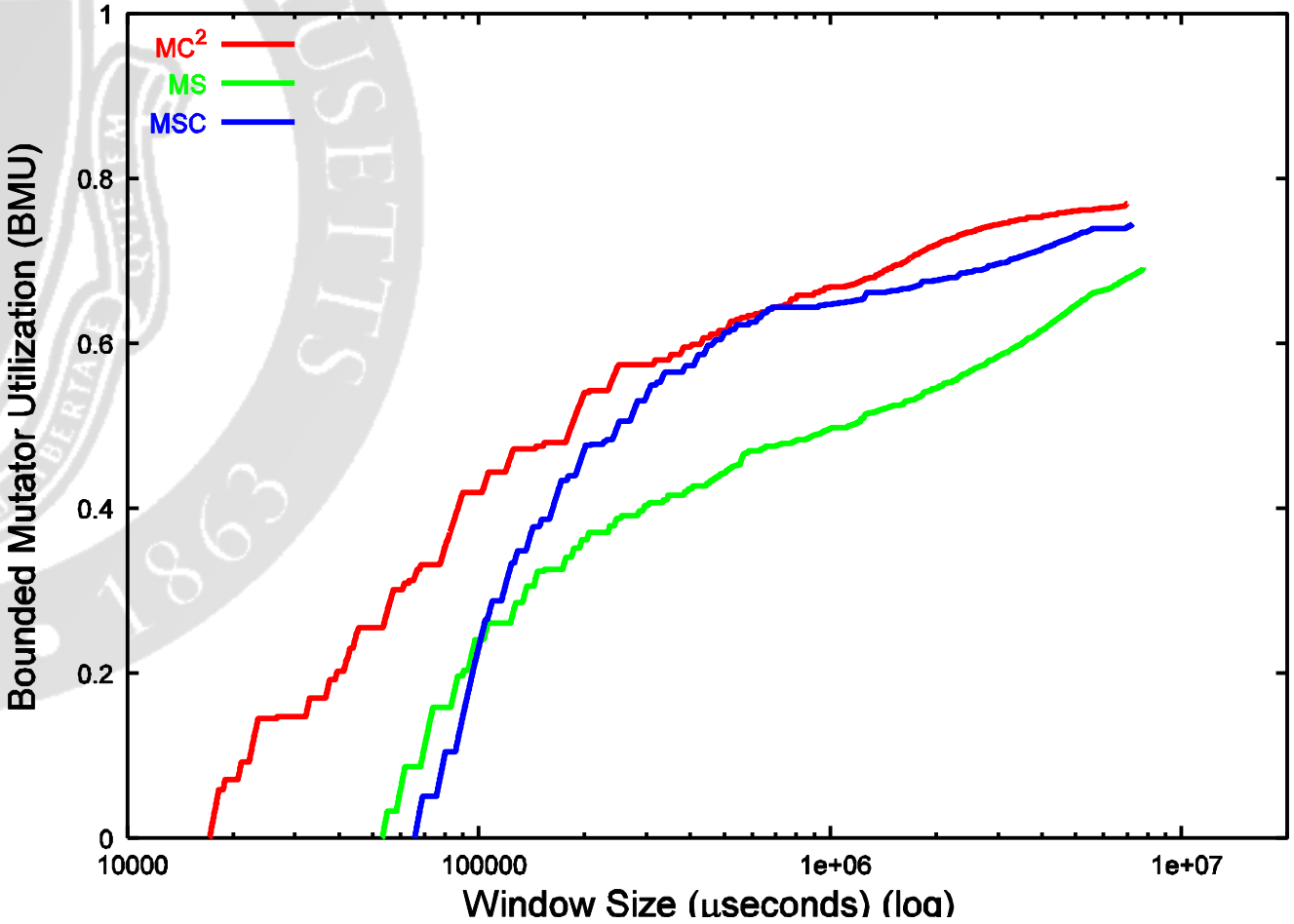
javac Execution Time



Pause Time, Execution Time Comparison (80% space overhead)

Benchmark	MC² MPT (ms)	MS MPT (ms)	MS/MC² ET	MSC MPT (ms)	MSC/MC² ET
_202_jess	17.2	53.2	1.11	65.7	1.04
_209_db	19.9	123.0	1.11	198.5	0.96
_213_javac	40.4	171.9	0.92	308.9	0.89
_227_mtrt	29.6	138.1	1.02	225.2	0.96
_228_jack	23.9	59.7	1.03	92.9	0.99
pseudojbb	41.5	168.2	1.08	314.5	0.98
Geo. Mean	**27.2**	**107.7**	**1.04**	**172.7**	**0.97**

Bounded Mutator Utilization (jess)



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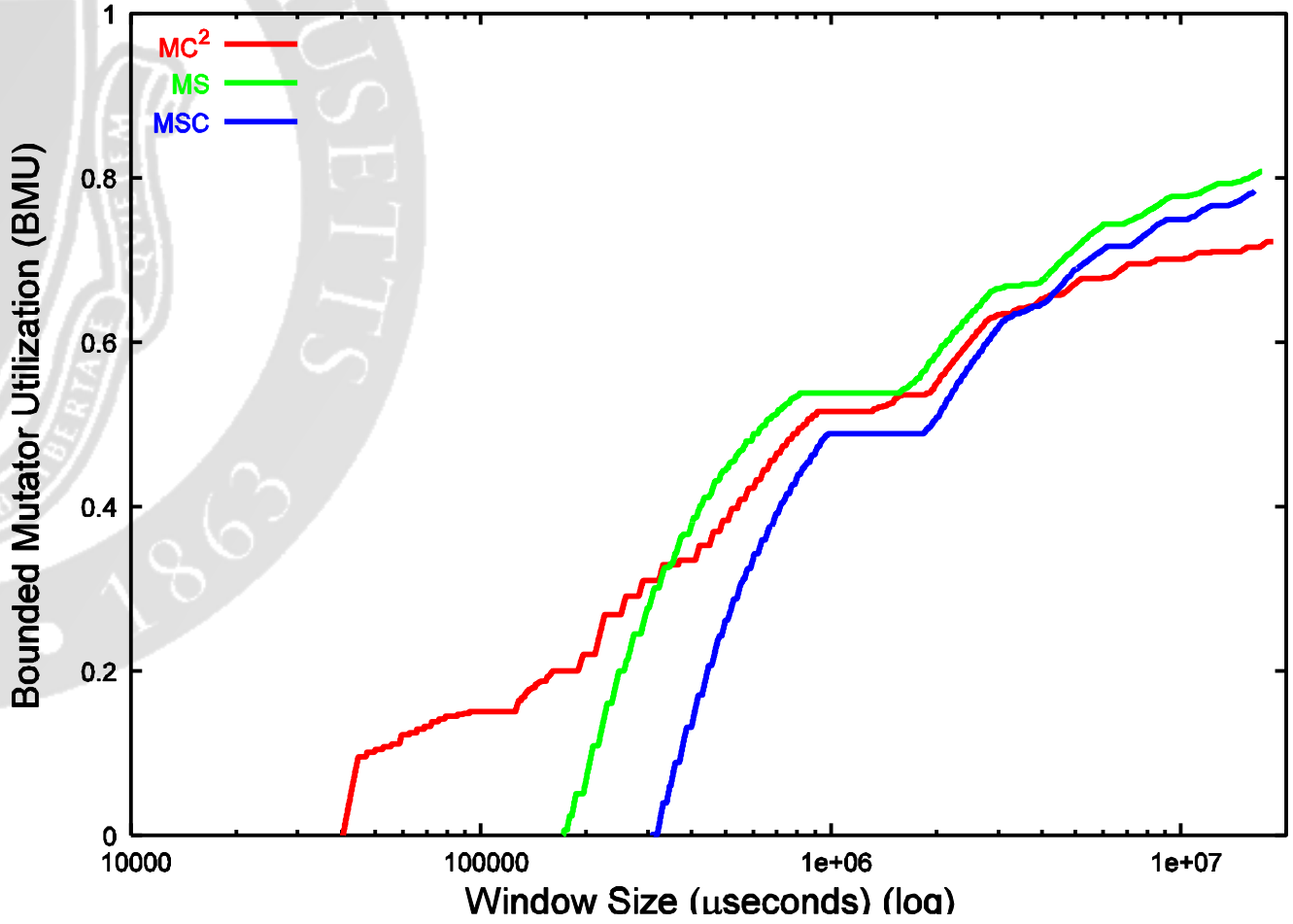

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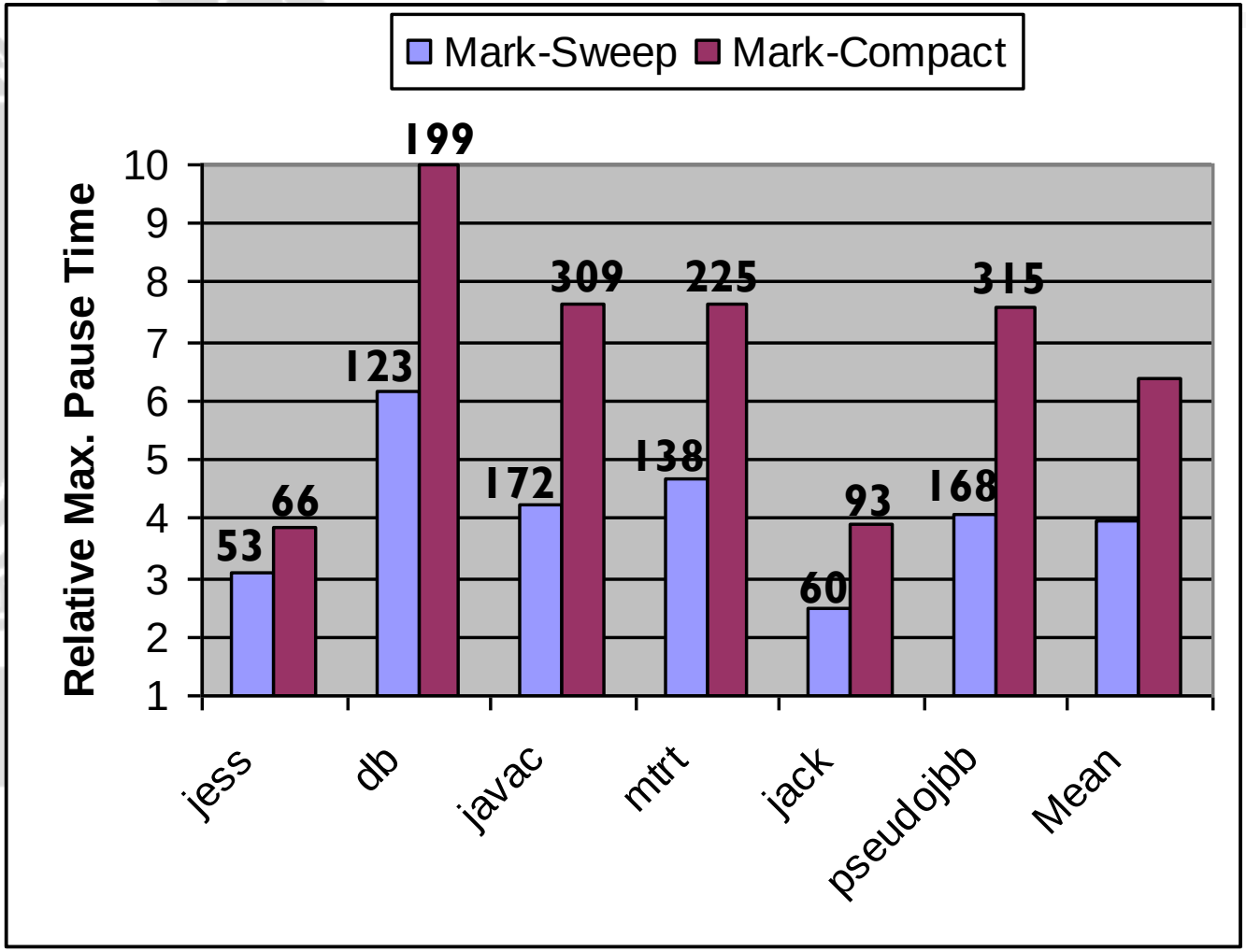

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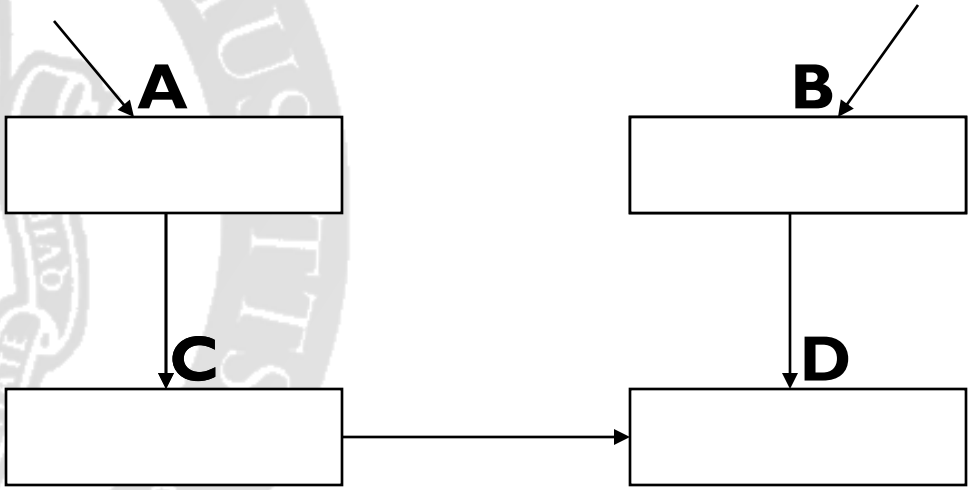
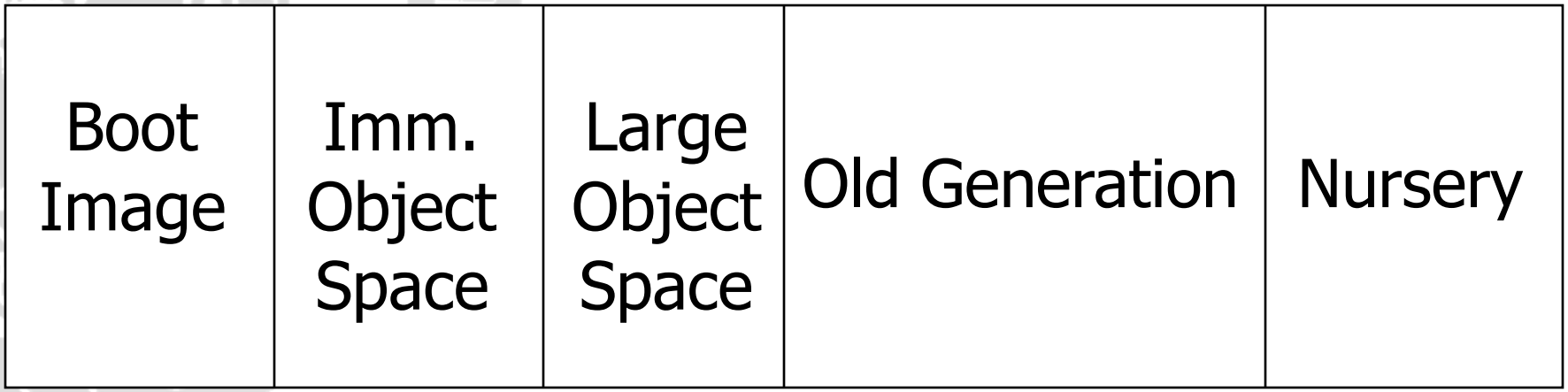
Incremental Marking Error



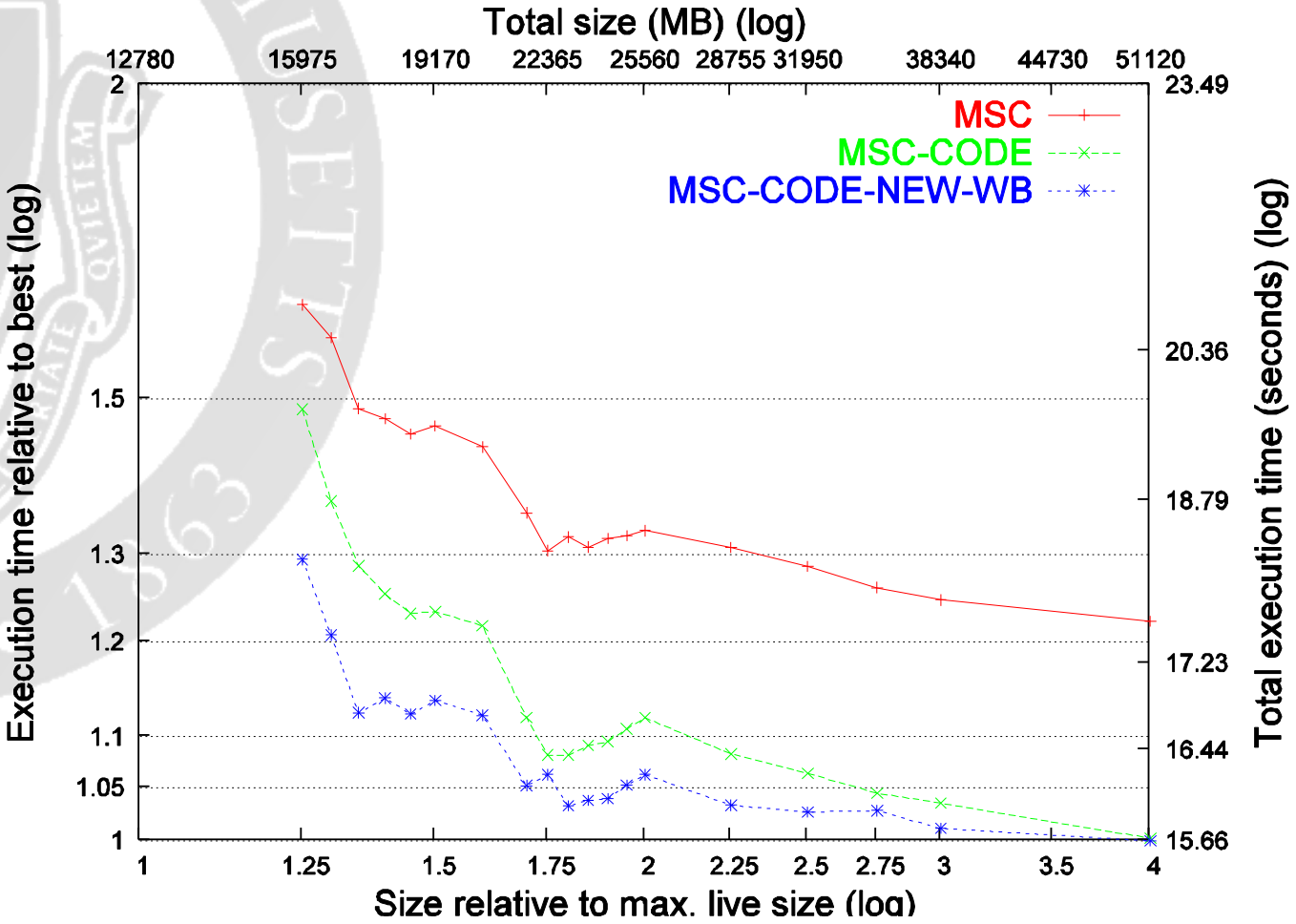
Handling marking error

- Track mutations using write barrier
- Record modified old generation objects
- Scan these modified objects at GC time
- Record interesting slots in remembered sets

JMTk Heap Layout



_213_javac (MSC execution time)



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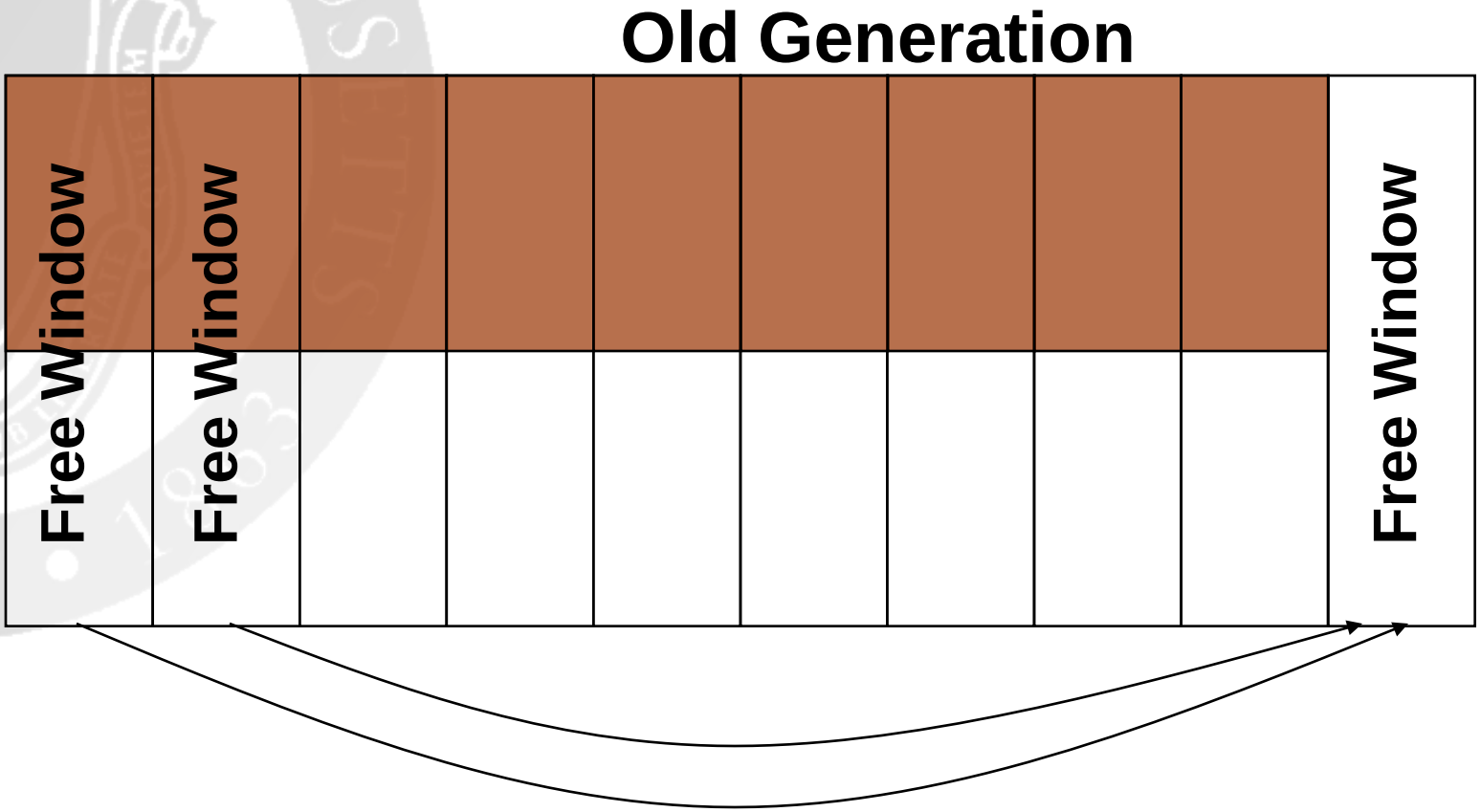

Reachable data Unreachable data Copied data

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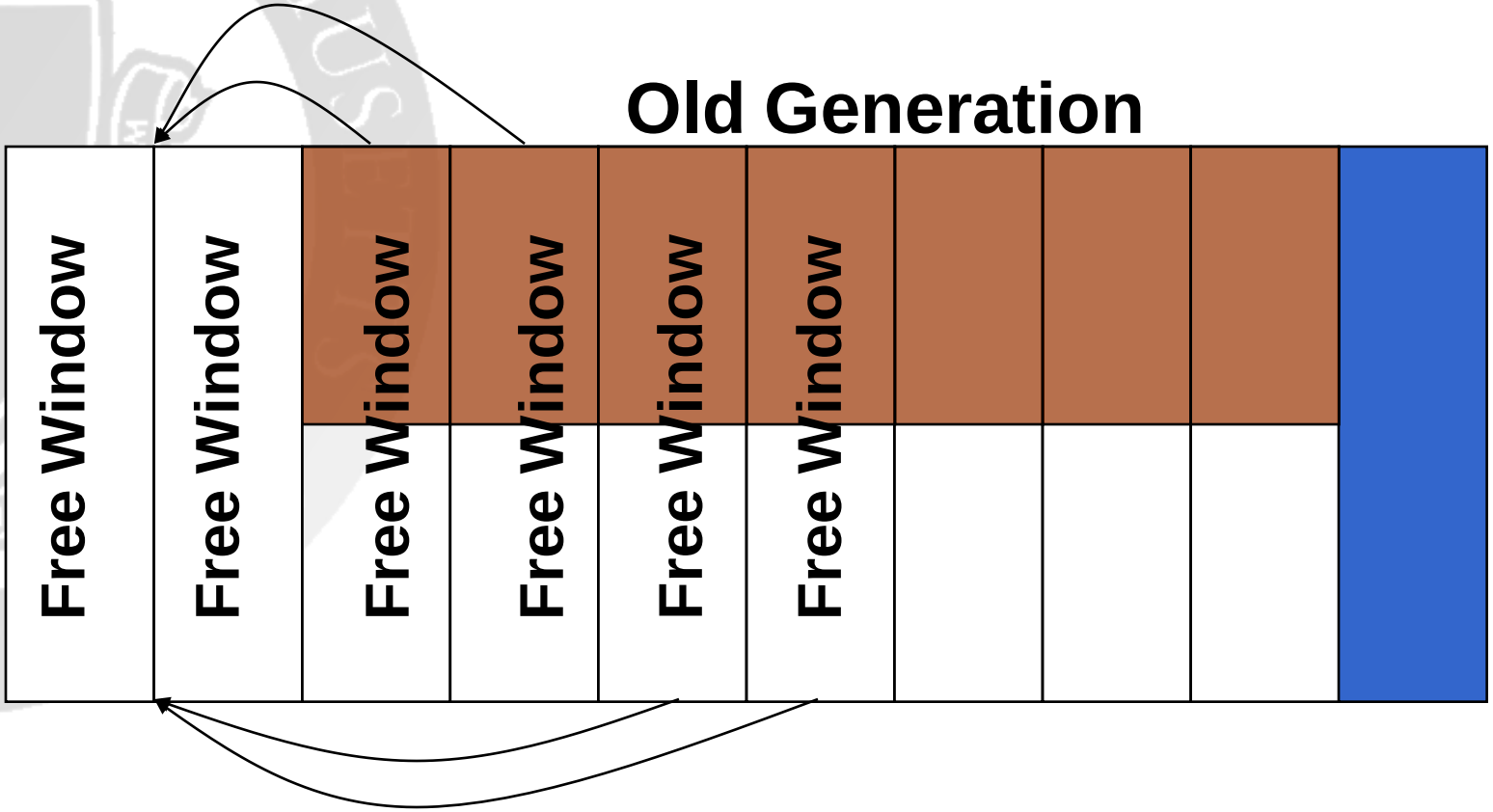

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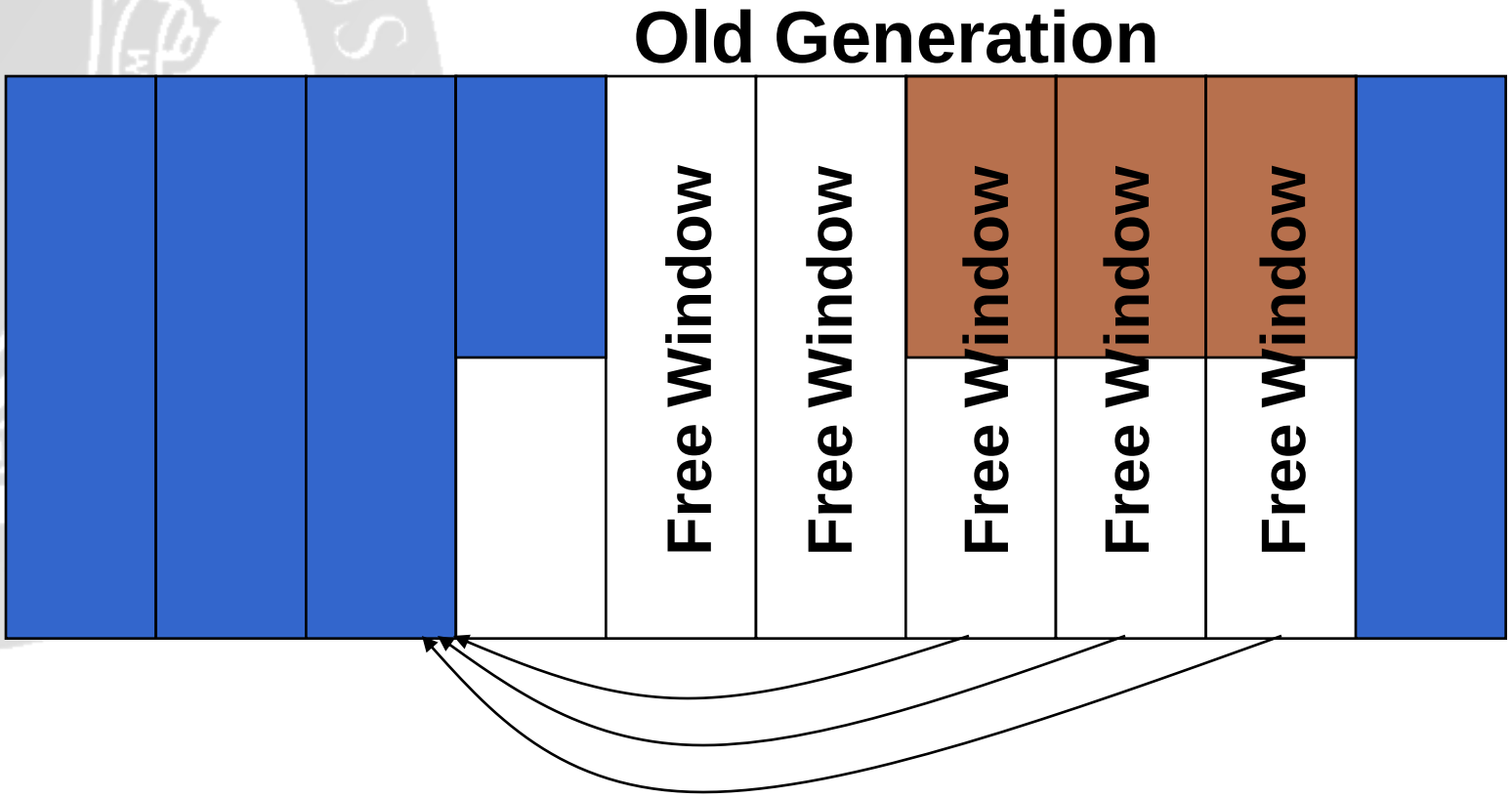

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