

COL1000: Introduction to Programming

Algorithmic Thinking (Merge & Merge Sort)

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Other Examples of Greedy

- Classroom Scheduling
- Huffman Encoding
- **Minimum Spanning Tree (on graphs)**
- Load Balancer
- Set Cover
- 0/1 Knapsack
- ...

Divide and Conquer



Divide and Conquer: Merge Sort

- **Divide:** Split the problem into smaller subproblems
- **Conquer:** The subproblems are solved (usually recursively) by applying the same solution
- **Combine:** Fusing the results together of the subproblems.
- **Problem:** sort $A = [8, 3, 1, 7, 0, 10, 2]$ using MergeSort

Divide and Conquer: Merge Sort

- **Problem:** sort A = [8, 3, 1, 7, 0, 10, 2] using MergeSort
- **Algorithm Run (Trace):**
 - **Recursively split the problem to smaller problems**
 - Left: [8,3,1] Right:[7,0,10,2]
 - **Sort Left:** [8,3,1]
 - **Split again:** Left [8] Right[3,1]
 - **Sort Left → sorted → [8] Sort Right: [3,1]**
 - **Split again:** Left [3] Right[1]
 - **Sort left → sorted → [3]; Sort right → sorted → [1]**
 - **Combine (Merge) → [1,3]**
 - **Combine (Merge) → [1,3,8]**

Divide and Conquer: Merge Sort

- **Merge(A, B):**
 - **Base case:**
 - **if not A: return B**
 - **if not B: return A**
 - **Induction step:**
 - **If $A[0] \leq B[0]$: return $[A[0]] + \text{Merge}(A[1:], B)$**
 - **return $[B[0]] + \text{Merge}(A, B[1:])$**
- **Problem:** sort $A = [8, 3, 1, 7, 0, 10, 2]$ using MergeSort
- **Algorithm:**
 - **MergeSort(A):**
 - **Base case: if $\text{len}(A) \leq 1: \text{return } A$**
 - **Induction Step:**
 - **Mid = $\text{len}(A) // 2;$**
 - **LeftLst = MergeSort($A[0:mid]$) //LeftLst is sorted**
 - **RightLst = MergeSort($A[mid:]$) //RightLst is sorted**
 - **return Merge(LeftLst, RightLst)**

Exercise: Write the imperative (non-recursive) version of Merge and MergeSort

Other Examples of Divide-Conquer

- **Binary Search:** Finding a value in a sorted array
- **Maximum Subarray Sum:** Contiguous subarray with max sum
- **Power Function (x^n):** Recursive doubling
- **Closest Pair of Points:** Closest pair in the 2D space
- **Count Inversions in an Array**
- ...

Count Inversions

Brute Force

- Given a list (or array), A , find the number of pairs (i, j) where:
 - $i < j$, and
 - $A[i] > A[j]$
- Brute Force:**
 - Check all pairs

```
def cntInversions(A):  
    n = len(A)  
    for i in range(n):  
        for j in range(i + 1, n):  
            if A[i] > A[j]:  
                count += 1
```

The code runs for nearly n^2 many steps

Count Inversions – Divide & Conquer

- Given a list (or array), A , find the number of pairs (i, j) where:
 - $i < j$, and
 - $A[i] > A[j]$
- **Algorithm: $A = [2, 4, 1, 3, 5]$**
 - Split A in 2 halves; Count recursively the inversions in the left half and right half and sum them up
 - $L = [2, 4], R = [1, 3, 5]$
 - Count the inversions across Left half and Right half (**basically, Merge**)
 - $2 > 1 \rightarrow$ everything after 2 is also greater than 1 \Rightarrow I discovered $\text{len}(L) - i$ many inversions in one shot!
 - **Observation: Similar to Merge Sort** \rightarrow Why not sort the array and count the inversions!

Count Inversions – Divide & Conquer

- Code: