

**Important:** The question marked with a ♠ is this week's quiz. The start time for the quiz is 1PM and the end time for the quiz is 1:12PM. Questions marked with a \* may be a little harder and can be considered optional.

**Problem 1 [Die25, Prob 12, page 30]**

Show that every 2-connected graph contains a cycle.

**Problem 2 [LLM18, Prob. 12.37]**

A graph  $G$  is called *2-removable* if it contains two vertices  $v \neq w$  such that  $G - v$  and  $G - w$  are both connected. Prove that every connected graph with at least 2 vertices is 2-removable.

**Problem 3**

Given a graph  $G = (V, E)$  and a minimal edge separator  $F \subseteq E$ , show that any cycle of  $G$  contains an even number of edges of  $F$  (this number could be 0 as well).

**Problem 4**

Let  $\bar{G}$  be the complement of the graph  $G$ , i.e., all edges of  $G$  are non-edges of  $\bar{G}$  and vice versa. Show that both  $G$  and  $\bar{G}$  cannot be disconnected, i.e., at least one of them must be connected.

**Problem 5 Related to [LLM18, Prob. 12.38]**

The *n-Hamming cube* is a graph with  $V(G) = \{0, 1\}^n$ , i.e., whose vertices are vectors with  $n$  coordinates, each of which can be either 0 or 1. We put an edge between any two vertices whose vectors differ in exactly one coordinate.

**Problem 5.1**

Prove that the *n-Hamming cube* is a connected graph for any  $n > 0$ .

**Problem 5.2**

Find  $\kappa(G)$  for the *n-Hamming cube*. In order to do this first prove that there are for any two vertices  $x, y$  there is a collection of  $n$   $x - y$  paths that are independent of each other (i.e. that don't share any vertex apart from  $x$  and  $y$ ).

**Problem 6 [LLM18, Prob. 12.42]**

An edge is said to *leave* a set of vertices  $A \subset V$  if only one endpoint of the edge is in  $A$ . A graph is called *k-mangled* if there is at least one edge leaving every subset of vertices of size  $k$  or smaller.

**Problem 6.1 ♠**

Prove that every  $\lfloor |V|/2 \rfloor$ -mangled graph is connected.

**Problem 6.2**

Is every  $\lceil |V|/3 \rceil$ -mangled graph connected?

**Problem 7 [LLM18, Prob. 12.42]**

If a graph has  $e$  edges,  $v$  vertices, and  $k$  connected components, then it has at least  $e - v + k$  cycles. Prove this by induction on the number of edges  $e$ .

### Problem 8

Given a graph  $G = (V, E)$  such that  $|V| = n$ , a cut  $F \subset E$  is called a *balanced cut* if  $G \setminus F$  has exactly 2 components and each of these components has size at least  $n/3$ . Construct graphs on  $n$  vertices whose smallest balanced cut has size (a)  $\theta(1)$ , (b)  $\theta(\sqrt{n})$  and (c)  $\theta(n)$ .

### Problem 9 [Die25, Prob 16, page 30] \*

Suppose that we have a  $k$ -edge connected graph with the property that removing any edge of this graph makes it lose its  $k$ -edge connectedness. Prove that such a graph has a vertex of degree  $k$ . Try to prove this first *without* using Menger's Theorem (c.f. Problem 10) and then see if using Menger's theorem makes it easier to prove.

### Problem 10 (Menger's Theorem)

Prove that a graph  $G$  has  $\lambda(G) = k$  for any  $k \geq 1$  iff there are  $k$  edge-disjoint paths between any pair of vertices in  $G$ . Two paths are said to be edge-disjoint if they don't share any edges. Caution: One direction of this theorem is easy and the other is tricky.

## References

[Die25] Reinhard Diestel, *Graph Theory 6ed.*, Springer, 2025.

[LLM18] E. Lehman, F. T. Leighton, and A. R. Meyer. Mathematics for Computer Science, June 2018, MIT Open Courseware.