

**Important:** The question marked with a ♠ is to be written on a sheet of paper and submitted to your TA within the first 10 minutes of the beginning of your tutorial session. Questions marked with a \* are optional challenge problems and are not to be discussed in the tutorial.

### Problem 1 [1, Prob 12, page 30]

Show that every 2-connected graph contains a cycle.

### Problem 2

Show using only the material covered in [1, Ch 1.4] that every connected graph on  $n$  vertices has at least  $n - 1$  edges.

### Problem 3

Generalize the result of Problem 2 to show that every graph on  $n$  vertices and  $m$  edges has at least  $n - m$  components.

### Problem 4

Given a graph  $G = (V, E)$  and a minimal edge separator  $F \subseteq E$ , show that any cycle of  $G$  contains an even number of edges of  $F$  (this number could be 0 as well).

### Problem 5

The  $n$ -Hamming cube is a graph with  $V(G) = \{0, 1\}^n$ , i.e., whose vertices are vectors with  $n$  coordinates, each of which can be either 0 or 1. We put an edge between any two vertices whose vectors differ in exactly one coordinate.

### Problem 5.1 ♠

Prove that the  $n$ -Hamming cube is a connected graph for any  $n > 0$ .

### Problem 5.2 \*

What is the highest value  $k$  such that the  $n$ -Hamming cube is  $k$ -connected?

### Problem 6

Let  $\bar{G}$  be the complement of the graph  $G$ , i.e., all edges of  $G$  are non-edges of  $\bar{G}$  and vice versa. Show that both  $G$  and  $\bar{G}$  cannot be disconnected, i.e., at least one of them must be connected.

### Problem 7

Given a graph  $G = (V, E)$  such that  $|V| = n$ , a cut  $F \subset E$  is called a *balanced cut* if  $G \setminus F$  has exactly 2 components and each of these components has size at least  $n/3$ . Construct graphs on  $n$  vertices whose smallest balanced cut has size (a)  $\theta(1)$ , (b)  $\theta(\sqrt{n})$  and (c)  $\theta(n)$ .

### Problem 8 (Menger's Theorem)

Prove that a graph  $G$  has  $\lambda(G) = k$  for any  $k \geq 1$  iff there are  $k$  edge-disjoint paths between any pair of vertices in  $G$ . Two paths are said to be edge-disjoint if they don't share any edges. Caution: One direction of this theorem is easy and the other is tricky.

## References

- [1] Reinhard Diestel, *Graph Theory 5ed.*, Springer, 2016.