

Q2.

Theorem: A water level of x L is possible to create in the given game if and only if x is a multiple of $\gcd(a, b, c)$.

Observation: The amount of water in each jug is an integer linear combination of a, b and c .

Corollary: The amount of water in each jug is an integer multiple of $\gcd(a, b, c)$.

Lemma 1: $\gcd(a, b, c) = \gcd(\gcd(a, b), c)$

Proof of Lemma 1:

Let $g_1 = \gcd(a, b, c)$

$$\Rightarrow g_1 | a, g_1 | b, g_1 | c$$

$$\Rightarrow g_1 | \gcd(a, b) \text{ and } g_1 | c$$

$$\Rightarrow g_1 | \gcd(\gcd(a, b), c) \quad \text{--- ①}$$

Also, let $\gcd(\gcd(a, b), c) = g_2$

$$\Rightarrow g_2 | \gcd(a, b) \text{ and } g_2 | c$$

$$\Rightarrow g_2 | a, g_2 | b, g_2 | c$$

$$\Rightarrow g_2 | \gcd(a, b, c) \quad \text{--- ②}$$

From ① and ②: $g_1 | g_2$ and $g_2 | g_1$ and since g_1 and g_2 are positive (both are gcd), $g_1 = g_2$

Hence proved.

Lemma 2: $\gcd(a, b, c) = \text{spr}(a, b, c)$

Proof of Lemma 2: Let $S :=$ set of positive integers of the form $\lambda a + \mu b + \eta c$
 $(\lambda, \mu, \eta \in \mathbb{Z})$
 and we can see that the set S is non-empty (can choose $\lambda = \mu = \eta = 1$)

\therefore By Well Ordering Principle, S has a least element $f = \text{spc}(a, b, c)$

$$\Rightarrow f = \lambda_0 a + \mu_0 b + \eta_0 c$$

To show: $f \mid b$. Assume contradiction.

\therefore By division theorem, $b = qf + r$, $0 < r < f$

$$\begin{aligned}\Rightarrow r &= b - qf \\ &= b(1 - q\mu_0) - q\lambda_0 a - q\eta_0 c\end{aligned}$$

$\Rightarrow r$ is a linear combination of a, b, c and $r < f$ which is a contradiction since $f = \text{spc}(a, b, c)$

$\therefore f \mid b$ and $r = 0$.

Similarly, $f \mid a$ and $f \mid c$

$\therefore f$ is a common divisor of a, b and c

$$\Rightarrow f \leq \text{gcd}(a, b, c) \quad \text{--- (3)}$$

Now, let $g = \text{gcd}(a, b, c)$

$$\Rightarrow g \mid a, g \mid b, g \mid c$$

$$\Rightarrow g \mid (\lambda_0 a + \mu_0 b + \eta_0 c)$$

$$\Rightarrow g \mid f$$

$$\Rightarrow g \leq f \quad \text{--- (4)}$$

From (3) and (4), $\text{gcd}(a, b, c) = \text{spc}(a, b, c)$

Proof of theorem:

(\Rightarrow) Follows from corollary

(\Leftarrow) Suppose x is a multiple of $\gcd(a, b, c)$

Then x is a multiple of $\text{spc}(a, b, c)$ ($\left\{ \begin{array}{l} \text{From Lemma 2,} \\ \gcd = \text{spc} \end{array} \right.$)

w.l.o.g. assume $a \leq b \leq c$

Feasibility: $0 \leq x \leq c$

If $x = c$, it is trivial.

Hence, assume $0 \leq x < c$

Claim: We can replace jugs a and b with a single jug of capacity $\gcd(a, b)$.

We have, $x = s(\alpha a + \beta b) + tc$, where $\gcd(a, b) = \alpha a + \beta b = g$
with $s > 0$ and $\alpha > 0 \rightarrow$ (if not, we can always adjust s, t, α, β to achieve this)
 $\rightarrow \{x \text{ is a linear combination of } a, b, c\}$

Repeat the procedure s times:

* Repeat α times:

• Fill aL jug and pour it into bL jug.

• If bL jug is full, empty it and pour rest of a into b

* Pour bL jug into cL jug.

* If cL jug is full, empty it and pour rest of b into c .

We had already proved for 2 jug case that the above procedure:

* at each ^{iteration} ~~step~~, after α repeats of 1st step, bL jug will have exactly $gL (= \gcd(a, b) = \alpha a + \beta b)$ of water left.

* Hence, 1st and 2nd step combined is equivalent to replacing aL and bL jugs with a single jug of gL .

* Hence, again, after s iterations, we will have exactly $x = [s(\alpha a + \beta b) + tc]L$ of water in jug c {from the two jug case proved in class}

Now, from Lemma 1, $\gcd(\gcd(a, b), c) = \gcd(a, b, c)$

Hence, x is an integer multiple of $\gcd(a, b, c)$ as well as $\gcd(\gcd(a, b), c)$

Hence, if x is an integer multiple of $\gcd(a, b, c)$, then xL is possible to fill.

Hence proved.