

Tutorial Sheet - 6

2. a) Theorem: Every graph with width at most w is $(w+1)$ -colorable.

Proof: By induction, on number of vertices n in the graph.

$P(n)$: A graph $G = (V, E)$ having n vertices (i.e. $|V|=n$) with width at most w is $(w+1)$ -colorable.

Base case: $n=1$, i.e. G has only 1 vertex, so its width $w=0$.

Hence, it can be colored in $w+1=1$ color. $\therefore P(1)$ is true.

Induction hypothesis: Assume, for some $n \in \mathbb{N}$, $P(n)$ is true.

Induction step: We need to show $P(n+1)$ is true.

Consider graph $G = (V, E)$, with $|V|=n+1$ and width at most w .

That is, the vertices can be arranged in a sequence $v_1, v_2, \dots, v_n, v_{n+1}$ such that $\forall i \in \{1, \dots, n+1\}$, v_i is connected to at most w vertices preceding it.

Remove v_{n+1} and all edges incident to it from G to obtain a new graph G' with n -vertices and width at most w . Since removing v_{n+1} and retaining the same sequence does not affect the number of edges from a vertex v_i ($i \in \{1, 2, \dots, n\}$) to a preceding vertex.

Hence, from induction hypothesis, G' is $(w+1)$ -colorable.

Now, replacing v_{n+1} with all its incident edges. Since G has width at most w and there are no vertices succeeding v_{n+1} , v_{n+1} has at most w neighbours which precede v_{n+1} in the sequence. So, we can colour v_{n+1} differently from all its neighbours and this will require only $(w+1)$ colors.

$\Rightarrow P(n+1)$ is true

Hence Proved.

b) Theorem: The average degree of a graph of width w is atmost $2w$.

Proof: By induction, on number of vertices n in the graph.

$P(n)$: A graph $G = (V, E)$ with $|V| = n$ and width w has an average degree of atmost $2w$.

Base case: $n=1$, width of the graph, $w=0$ and ~~its~~ its average degree is $\frac{0}{1} = 0 = 2w$.

$\therefore P(1)$ is true.

Induction hypothesis: Assume for some $n \in \mathbb{N}$, $P(n)$ is true.

i.e.
$$\frac{\sum_{v \in V} \deg(v)}{|V|} \leq 2w$$

Induction step: We need to show $P(n+1)$ is true.

Consider graph $G = (V, E)$ with $|V| = n+1$ and width w . The vertices can be arranged in a sequence $v_1, v_2, \dots, v_n, v_{n+1}$ such that $\forall i \in \{1, \dots, n+1\}$, v_i is connected to atmost w vertices preceding it.

Remove v_{n+1} and all edges incident to it to obtain a new graph $G' = (V', E')$ with $|V'| = n$ and width w (reason for width $= w$ explained in part a))

From induction hypothesis, $\sum_{v \in V'} \deg(v) \leq 2wn$

Now, replacing v_{n+1} with all its incident edges. Since width of G is w , atmost w edges are incident to v_{n+1} and only the degree of atmost w vertices preceding v_{n+1} increases by 1 on replacing v_{n+1} .

Hence, average degree =
$$\frac{\sum_{v \in V} \deg(v)}{n+1} \leq \frac{\sum_{v \in V'} \deg(v) + 2w}{n+1}$$

$$\leq \frac{2wn + 2w}{n+1} = \frac{2w(n+1)}{n+1} = 2w$$

[sum of degrees of n vertices when v_{n+1} removed + $\deg(v_{n+1})$ + at most w edges to v_{n+1}]

$\therefore P(n+1)$ is true. Hence proved.