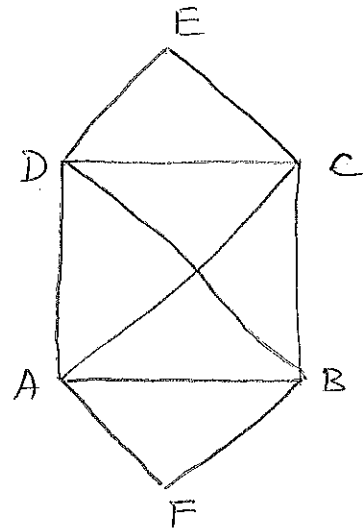
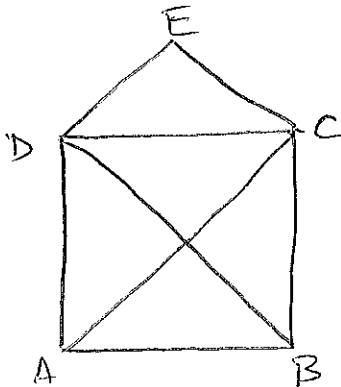


Homework 1

1. Problem Formulation.

Graphs:



The primary data structure that we check here is `EdgeVisited`.

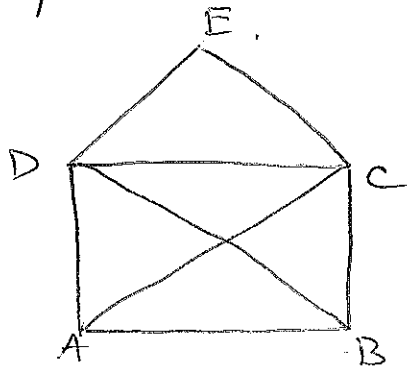
Initial State: `EdgeVisited = []`

Goal State: `EdgeVisited = [(A,B), (B,A), ..., (all edges visited), ...]`

At the goal state, we can verify if we reach back to the source node, then we have the cycle; or else we have the path. This calculation can be pre-empted to know what to expect.

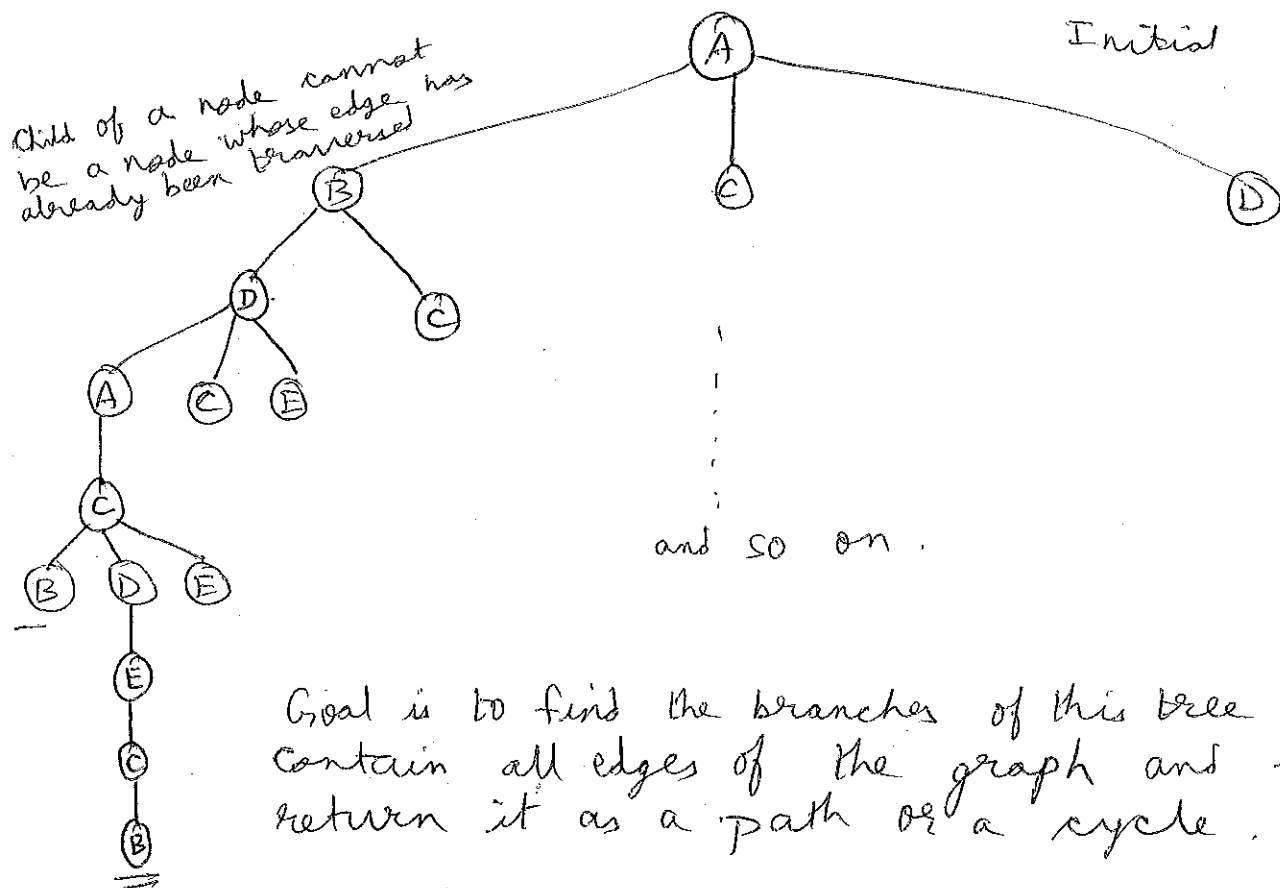
At any arbitrary state: `EdgeVisited = [..., (ai, bi), (bi, ai), ...]`
a_i and b_i are the nodes where b_i ~~was~~ ^{is} the node just entered from a_i.

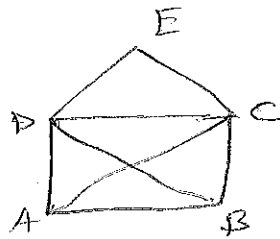
We can ~~use~~ use BFS or DFS policy on the search tree for the graphs.
To show an example; consider the first graph:



$A \rightarrow B, C, D$
 $B \rightarrow A, D, C$
 $C \rightarrow A, B, D, E$
 $D \rightarrow A, B, C, E$
 $E \rightarrow C, D$

The search tree will look as follows:





A sample implementation using DFS would be the following :

Stack	Node_Visited	Edge_Visited	Frequencies	Commands
A	-	-	[3, 3, 4, 4, 2]	Push(A)
-	A	-	[3, 3, 4, 4, 2]	Pop()
B C (D)	A	-	[3, 3, 4, 4, 2]	Push(A.available children)
B C	A, D	AD, DA	[2, 3, 4, 3, 2]	1. Pop() 2. If node-visited has entry, look for last entry(A) and combine with popped node to form edges and add to visited-edges list.
BC BC (E)	A, D	A AD, DA	[2, 3, 4, 3, 2]	Push(D.available children)

This can go on until Edge-Visited has all ~~popped~~ the edges of the graph, and since the given graphs are undirected, we add both front and back edges to the data structure.