

### 5.1 Show that $EQ_{CFG}$ is undecidable.

Assume  $R$  decides  $EQ_{CFG}$ , and construct TM  $S$  deciding  $ALL_{CFG}$

Let  $S =$  "On input  $\langle G \rangle$ ,

- 1) Construct CFG  $G_{all}$  such that  $L(G_{all}) = \Sigma^*$
- 2) Use  $R$  to determine if  $L(G) = L(G_{all})$
- 3) Accept if yes, otherwise reject."

### 5.2 Show that $EQ_{CFG}$ is co-Turing-recognizable.

Assume  $R$  recognizes  $\overline{EQ_{CFG}}$  and construct TM  $S$  that recognizes it.

Let  $S =$  "On input  $\langle G_1, G_2 \rangle$ ,

- 1) Enumerate all strings  $w \in \Sigma^*$
- 2) For each string  $w$ , simulate a decider for whether  $w \in L(G_1)$  and whether  $w \in L(G_2)$
- 3) If  $w \in L(G_1)$  and  $w \notin L(G_2)$  or  $w \notin L(G_1)$  and  $w \in L(G_2)$
- 4)  $S$  accepts, otherwise reject

### 5.3 Find a match in the following instance of the Post Correspondence Problem.

$$\left\{ \left[ \frac{ab}{abab} \right], \left[ \frac{b}{a} \right], \left[ \frac{aba}{b} \right], \left[ \frac{aa}{a} \right] \right\}$$

The sequence is,  $\frac{ab}{abab}, \frac{ab}{abab}, \frac{aba}{b}, \frac{b}{a}, \frac{b}{a}, \frac{aa}{a}, \frac{aa}{a}$

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$$\Rightarrow [1, 1, 3, 2, 2, 4, 4]$$

5.9: Let  $T = \{ \langle M \rangle \mid M \text{ is a TM that accepts } w^R \text{ whenever it accepts } w \}$ . Show that  $T$  is undecidable. (Some Editions have a different 5.9 Exercise)

Construct  $M'$  such that

is undecidable. (Some Editions have a different 5.9 Exercise)

Construct  $M'$  such that

$M' =$  "on input  $x$ :

- 1) If  $x = w$ , run  $M$  on  $w$
- 2) If  $M$  accepts  $w$ , accepts
- 3) If  $x = w^R$ , rejects
- 4) Otherwise, reject."

$S =$  "On input  $\langle M, w \rangle$ :

- 1) construct  $M'$
- 2) Run  $R$  on  $\langle M' \rangle$
- 3) If  $R$  accepts, reject
- 4) If  $R$  rejects, accept."