

Part 2 – Searching

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Searches

Beatles

searching for: Beatles

Binary search comparisons: 13

Loop comparisons: 337

Total comparisons: 350

Theoretical complexity ($\log_2 N$): 13.36

Santana

searching for: Santana

Binary search comparisons: 3

Loop comparisons: 114

Total comparisons: 117

Theoretical complexity ($\log_2 N$): 13.36

Arlo

searching for: Arlo

Binary search comparisons: 14

Loop comparisons: 6

Total comparisons: 20

Theoretical complexity ($\log_2 N$): 13.36

A

searching for: A

Binary search comparisons: 13

Loop comparisons: 582

Total comparisons: 595

Theoretical complexity ($\log_2 N$): 13.36

Z
searching for: Z
Binary search comparisons: 14
Loop comparisons: 141
Total comparisons: 155
Theoretical complexity ($\log_2 N$): 13.36

X
searching for: X
Binary search comparisons: 14
Loop comparisons: 2
Total comparisons: 16
Theoretical complexity ($\log_2 N$): 13.36

Task 2 Unit test output from SongCollection

Total songs

10514, first songs:

Aerosmith, "Adam's Apple"

Aerosmith, "Ain't Enough"

Aerosmith, "Ain't that a Bitch"

Aerosmith, "All Your Love"

Aerosmith, "Amazing"
Aerosmith, "Angel"
Aerosmith, "Angel's Eye"
Aerosmith, "Animal Crackers"
Aerosmith, "Attitude Adjustment"
Aerosmith, "Avant Garden"
BUILD SUCCESSFUL (total time: 0 seconds)

3. The $O(K + \log_2 N)$ Goal

This search method does meet that goal since binary search as implemented in the Java Arrays.binarySearch method has a worst case big O of $\log_2 N$ in the worst case scenario. When you combine that with while loops for adjacent elements now that you're in the neighborhood, you find that you will always conduct K number of operations if K is the number of songs that have that prefix. This gives us a big O notation of $O(K + \log_2 N)$

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Incomplete parts: None

Known bugs: None